# Provably Efficient Q-Learning with Low Switching Cost

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Draft poster, work in progress.

## **Motivation**

In many RL domains, executing a new policy is expensive. **Off-policy RL**: Find  $\pi_\star$  given only off-line data from  $\mu$ 

Challenging... Relax



**Limited Adaptivity RL:** Find  $\pi_{\star}$  with online data from  $\{\mu_1,\ldots,\mu_n\}$  for some small n

# **Local Policy Switch**

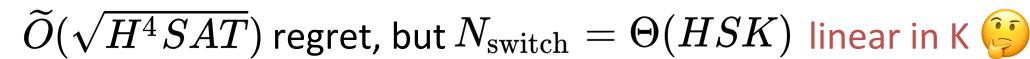
Setup: episodic MDP with horizon H, play K episodes Def: The **number of local policy switches** for an RL algorithm is

$$N_{ ext{switch}} = \sum_{k=1}^{K-1} \left| \{(h,s): \pi_h^k(s) 
eq \pi_h^{k+1}(s) \} 
ight|$$

where  $\pi^k$ is the (deterministic) policy it plays at episode K.

Smaller  $N_{\rm switch} \Rightarrow$  closer to off-policy

Prior work: Q-Learning with UCB-Hoeffding exploration<sup>[1]</sup>:



Any sublinear regret algo such that  $N_{
m switch}$  sublinear in K?

## **UCB2** for Bandits

**Algorithm** (UCB2):

- Select arm j that maximizes the UCB
- If this is the r-th time it's selected, play the arm exactly au(r) au(r-1)times, where  $\tau(r) = (1+\alpha)^r$

**Theorem**<sup>[2]</sup>: UCB2 achieves same regret as UCB & only log(K) policy

switches:  $N_{\mathrm{switch}} = O(A \log(K/A))$ 

Idea: Integrate UCB2 into Q-Learning!

# Algorithm: Q-Learning with UCB2 Scheduling

Idea: update the policy according to Q only when Q has been updated  $\tau(r) = (1 + \alpha)^r$  times.

```
Algorithm 2 Q-learning with UCB2 scheduling
input Parameter \alpha \in (0,1) and c > 0.
                                                                                                                                // Two sets of Q:
   Initialize: \widetilde{Q}_h(x,a) \leftarrow H, Q_h \leftarrow \widetilde{Q}_h, N_h(x,a) \leftarrow 0 for all (x,a,h) \in \mathcal{S} \times \mathcal{A} \times [H].
                                                                                                                                 Running estimate \hat{Q}
                                                                                                                                 Policy network Q
   for episode k = 1, \dots, K do
      Receive x_1.
      for step h = 1, \dots, H do
          Take action a_h \leftarrow \arg \max_{a'} Q_h(x_h, a'), and observe x_{h+1}. // Take action according to Q
          t = N_h(x_h, a_h) \leftarrow N_h(x_h, a_h) + 1; b_t = c\sqrt{H^3\ell/t}.
          \widetilde{Q}_h(x_h,a_h) \leftarrow (1-\alpha_t)\widetilde{Q}_h(x_h,a_h) + \alpha_t[r_h(x_h,a_h) + \widetilde{V}_{h+1}(x_{h+1}) + b_t]. // Update \widetilde{Q} via Q-Learning
          \widetilde{V}_h(x_h) \leftarrow \min \left\{ H, \max_{a' \in \mathcal{A}} \widetilde{Q}_h(x_h, a') \right\}.
          if t = \tau(r) for some r then
                                                                          // Set Q to be \widetilde{Q} occasionally according to UCB2 scheduling
             (Update policy) Q(x_h, \cdot) \leftarrow \widetilde{Q}(x_h, \cdot).
          end if
      end for
   end for
```

## Theoretical Result

**Theorem 1**: Q-Learning with UCB2 scheduling achieves regret  $O(\sqrt{H^4SAT})$  and policy switch bound  $N_{
m switch} \leq O(H^3 SA \log(K/A))$ 

which is logarithmic in K.



**Proof highlight**: improved *propagation of error argument* under delayed Q updates.

"Theorem" 2 (lower bound): Any sublinear regret or PAC algorithm must have

 $N_{
m switch} \geq \Omega(HSA)$ 

Mild gap:  $O(H^2\log(K/A))$ , conjecture that log is also necessary & gap is at most  $O(H^2)$ 

#### **Discussion & Future Work**

- Algorithms with tighter regret bounds (e.g. tighten the  $H^4$ )?
- Close the gap between lower and upper bound.
- Model-based algorithms with limited adaptivity? Better bounds?

#### References

[1] Jin, C., Allen-Zhu, Z., Bubeck, S. and Jordan, M.I., 2018. Is q-learning provably efficient?. In Advances in Neural Information Processing Systems (pp. 4868-4878).

[2] Auer, P., Cesa-Bianchi, N. and Fischer, P., 2002. Finite-time analysis of the multiarmed bandit problem. *Machine learning*, 47(2-3), pp.235-256.