Spring - Mar 7

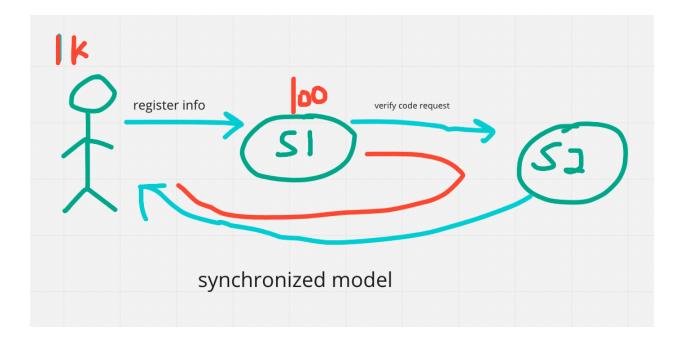
Mar 8 9am summary

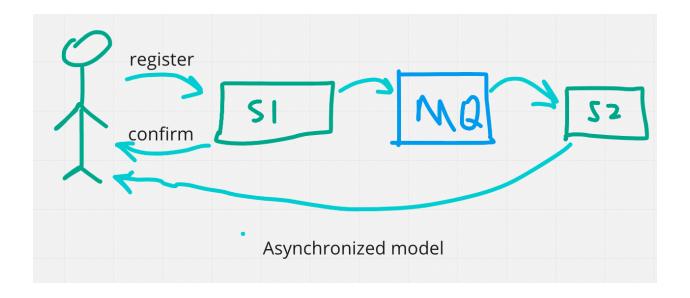
- review message Q
- Kafka architecture

•

37, Kafka

why message queue (message Q)





sever needs to register user's info and store in db

• for register, we need to verify user

user start register \rightarrow S1 request verification \rightarrow S2 send email to user for verification \rightarrow User click email link to verify

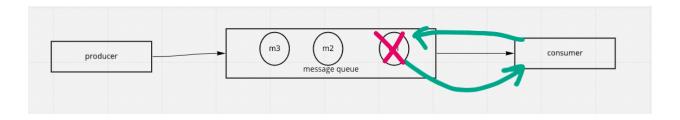
Problem: everyone want to access s1(limited threads)

use asynchronized model, s1 can send the confirm message to user immediately, and not waiting the whole process to finish

- · loosely coupled
- cache
- flexible
- asynchronized communication

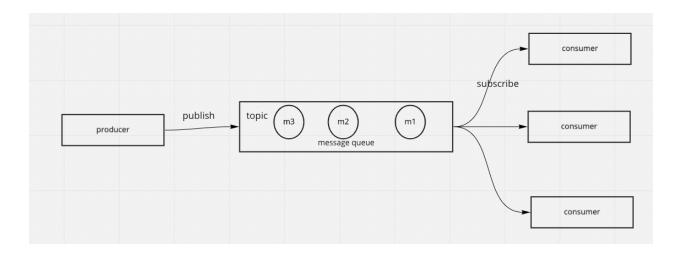
Message queue model

1. point to point (one to one)



- pull model (trigger by consumer)
 - consumer pulls the message from MQ
- push model (trigger by message queue is push model)
 - MQ pushes the message to consumer
 - When the consumer receives the data, it will send acknowledge message back to MQ. Then MQ deletes the m1.
- each message only can be consume once

2, publish and subscribe model (one to many)

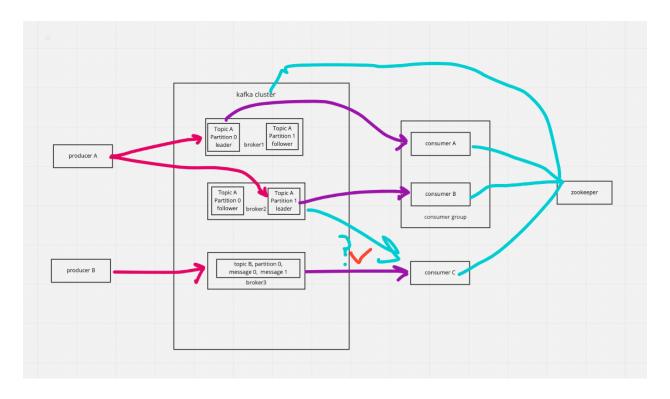


• producer publishes multiple messages to MQ, we called topic

- o consumer will register/subscribe to the topic, and retrieve the message from MQ
- message can be consumed multiple times

Kafka architecture(popular)

- another: rabbitMQ (message tool)



kafka use publish/subscribe model → different topic

• every cluster represent a sever

producer produce a message will send to leader partition the follower will "sync up" topic B just has one partition

zookeeper: coordinator will coordinate all the stuff

- manage kafka cluster
 - When leader node is down, zookeeper will handle it
 - the follow node for backup
- manage consumers

producer sends messages to leader node(partition) directly

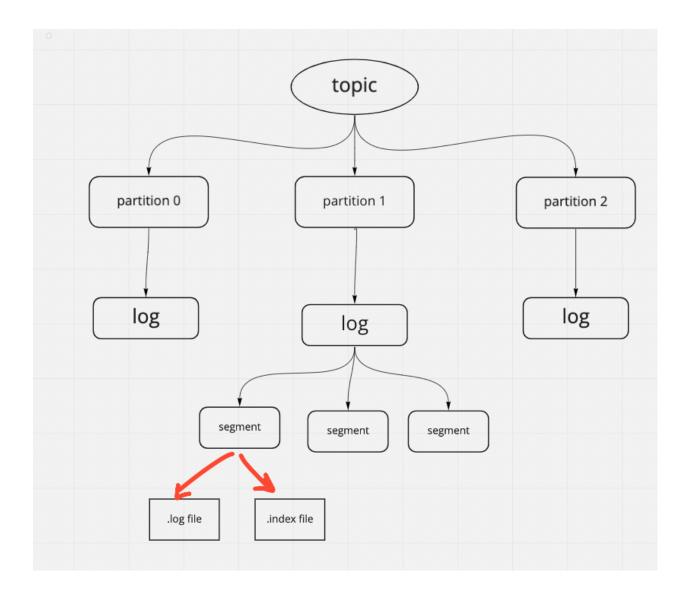
- Producer
 - produce some messages
- consumer
 - betry information from
- · consumer group
 - each consumer in the group can not consume same partition in the same time
- broker(physical sever)
 - kafka sever
 - each broker can carry different topic, and have different topics
- topic
 - theoretical idea
- partition
 - 0
- replica
 - backup
- leader
 - interacted leader and consumer
- follower

• "thing up" data following the leader node

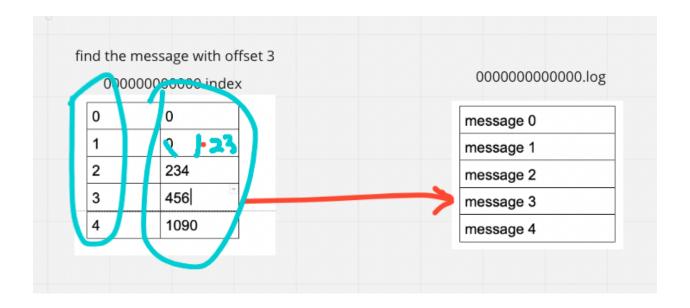
Dive deep into Kafka

file storage mechanism

file the data



we call all messages are "log"

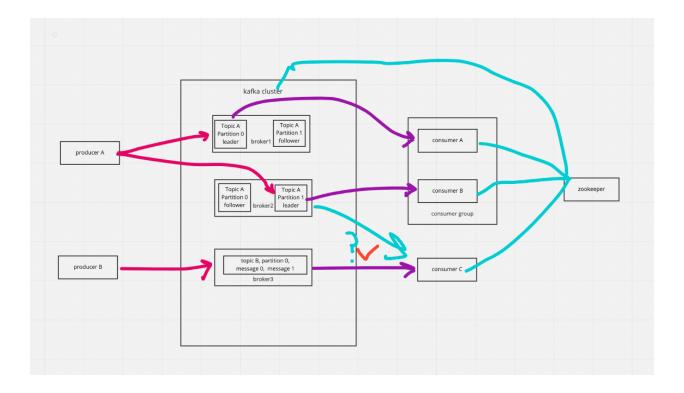


key-value

find the index first and jump the message by address

Producer

- why we need partition
 - divide big data to small partition
- ProderRecord(topic, partition, key)
 - → producer put the message into ProduerRecord, and then send the ProduerRecord to the sever.
 - the partition principle
 - if tell producer → partition number
 - else if use key -> hashcode -> %
 - else round robin algorithm
 - $P0 \rightarrow P1 \rightarrow P0 \rightarrow P1...$

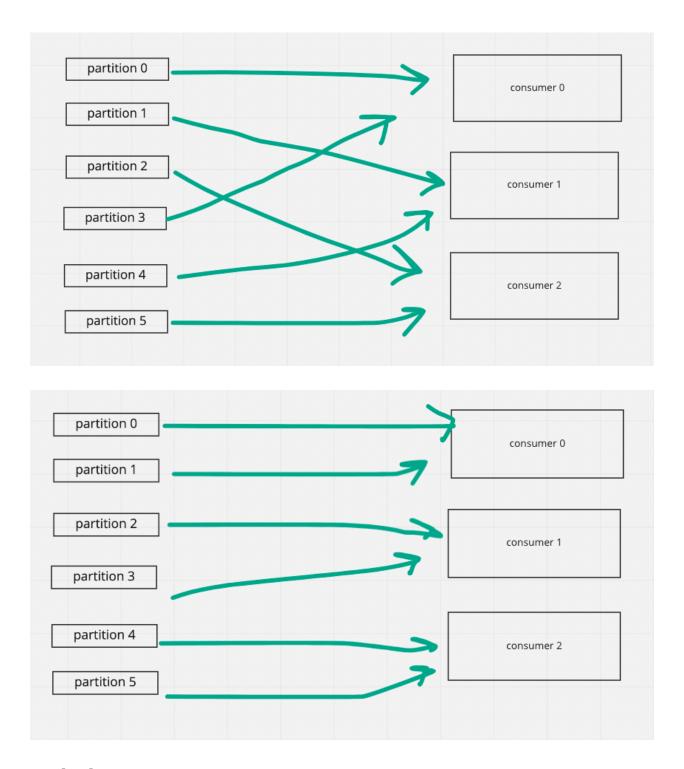


- how to guarantee the reliability
 - ack mechanism (producer)
 - When broker receives message, it will send acknowledge back to producer
 - three levels acknowledgement to handle a long waiting time of producer(wait ack):
 - 0, : at most once (producer does care if broker receives message) → non important data
 - may lose data
 - 1, P just care whether leader received
 - followers may lose data, cuz when leader send ack back, it will shut down. At that time, followers may not finish syncing data
 - -1: at least once → wait all the follows to finish syncing data
 - potentially duplicated data
 - ISR
 - In sync replica set

- all the follow nodes in the ISR will sync up data from leaders
- after sync, it will send ack back to producer
- exactly once (de-duplicated data)
 - at least once + incompetence
 - enable.idempotence = true;
 - broker do: <PID, Partition, SeqNumber>
 each time producer sends message to broker. Broker does the de-duplication,
 based on primary key (<PID, Partition, SeqNumber>). check they are three
 whether exist
 - set each producer a PID
 - set Partition number
 - set sequence Number for each message

Consumer

- partition assignment strategy
 - round robin
 - range



Rebalance

- ConsumerRebalanceListener
 - OnPartitionRevoked (before rebalance, commit the offset)
 - o OnParitionAssigned (after rebalance, seek offset)