

Week 7



STEVENS
INSTITUTE *of* TECHNOLOGY
THE INNOVATION UNIVERSITY®



An Introduction to Cyber Security – CS 573

Instructor: Dr. Edward G. Amoroso
eamoroso@tag-cyber.com

Required Week Seven Readings

**1. “THE POSSIBILITY OF SECURE NON-SECRET DIGITAL ENCRYPTION
by J. H. Ellis, January 1970**

<https://cryptocellar.org/cesg/possnse.pdf>

**2. Finish Reading “*From CIA to APT: An Introduction*”
to *Cyber Security*, E. Amoroso & M. Amoroso**

Twitter: @hashtag_cyber

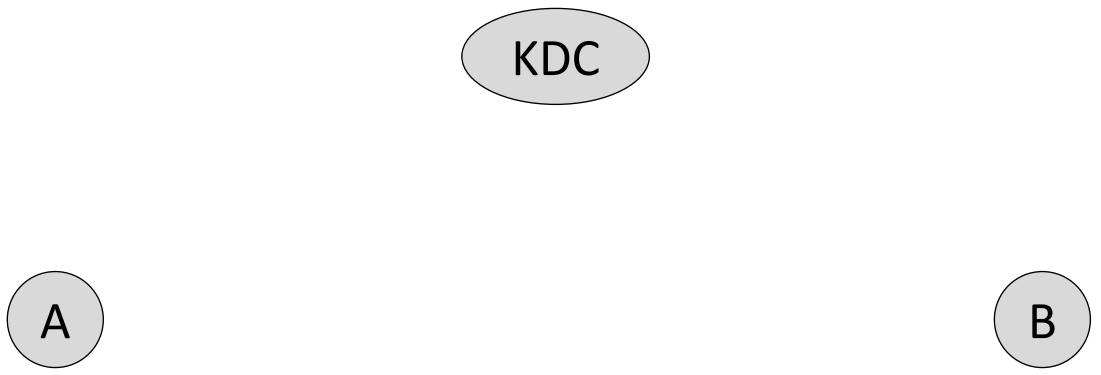
LinkedIn: Edward Amoroso



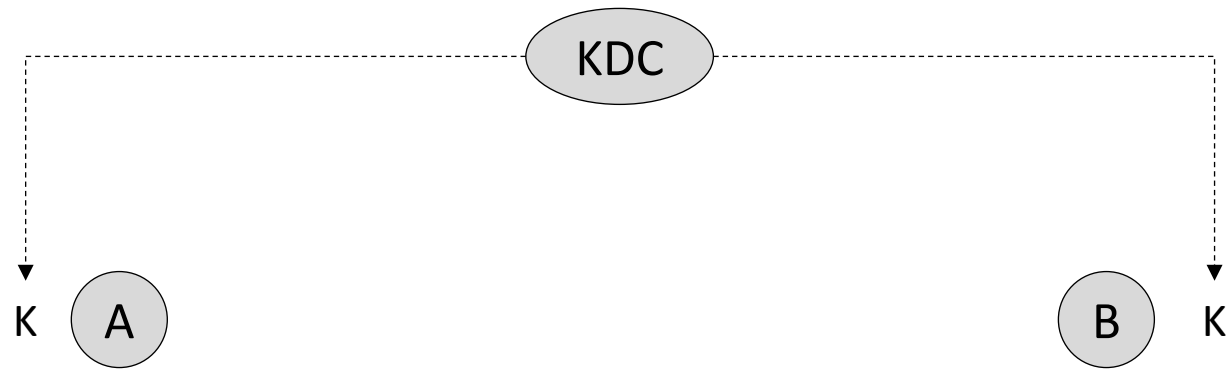
Week 7: Public Key Cryptography

What Properties of Conventional Cryptography
Must Be Maintained?

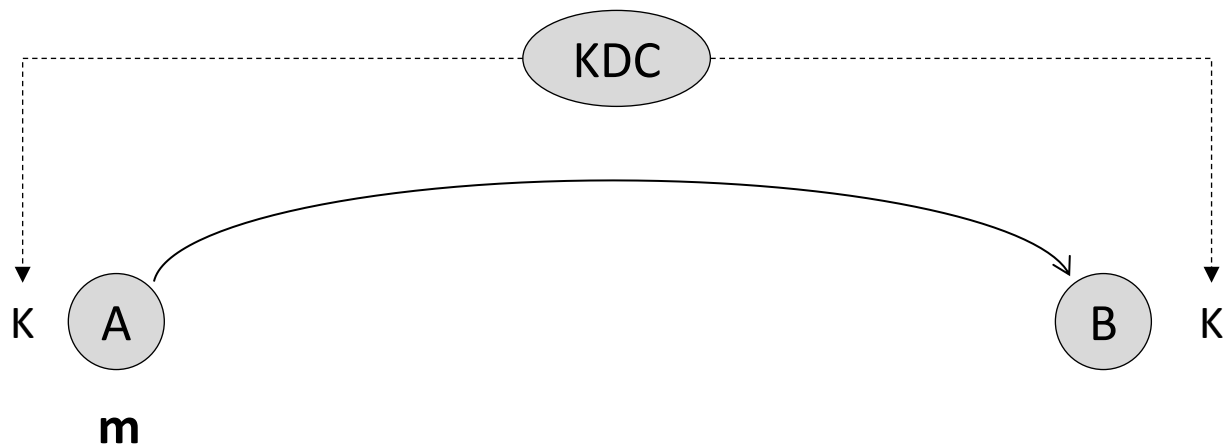
Conventional Cryptography



Conventional Cryptography

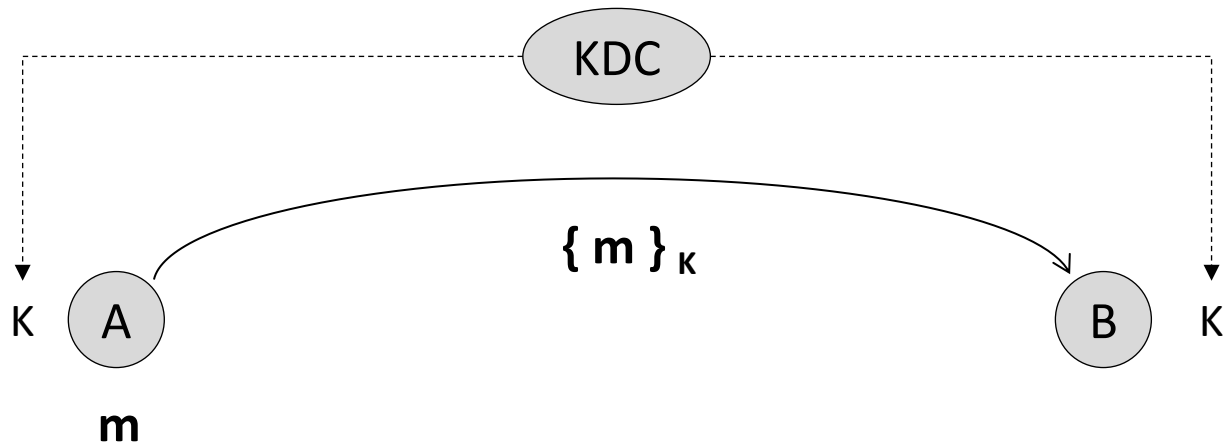


Conventional Cryptography



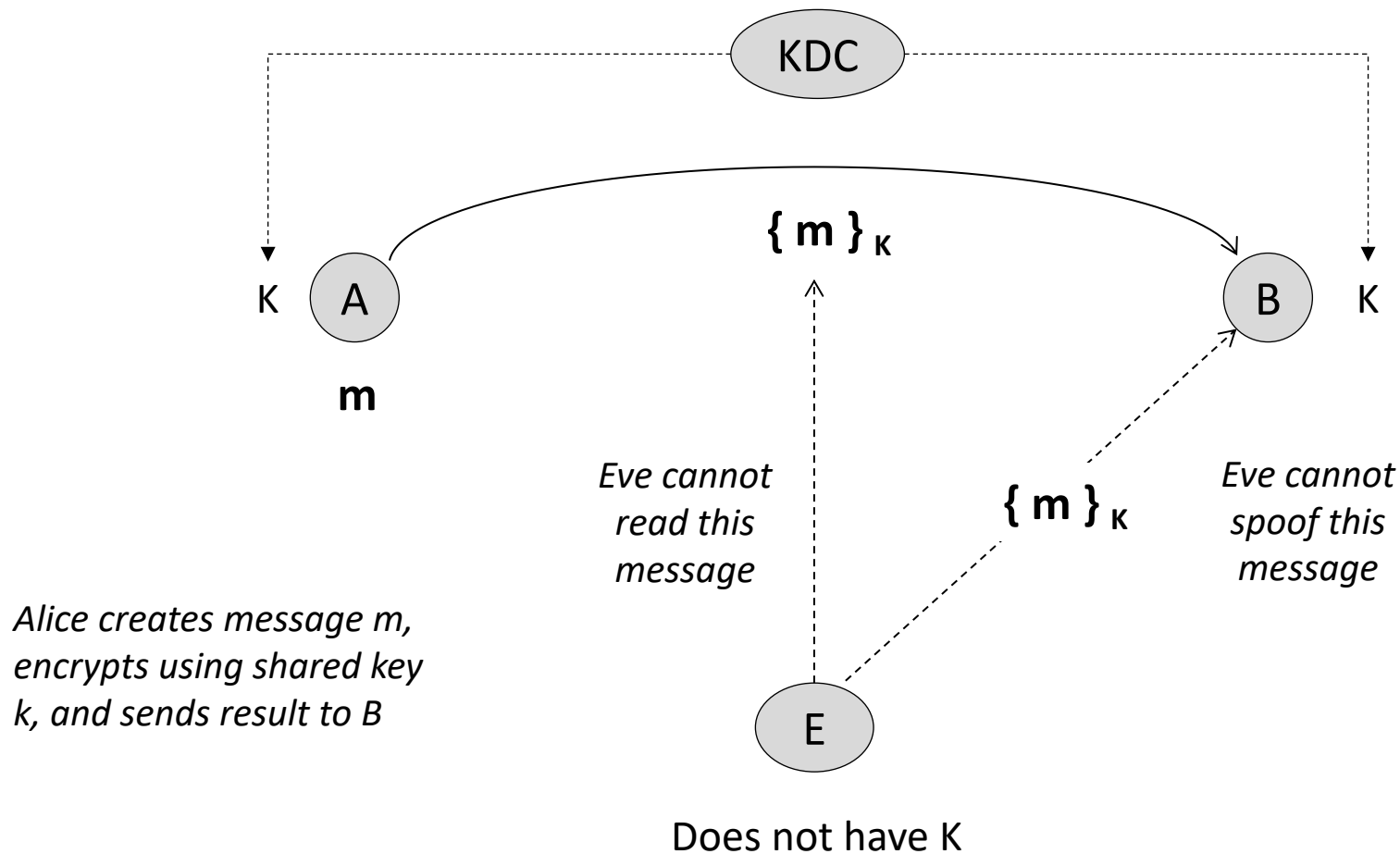
Alice creates message m . . .

Conventional Cryptography

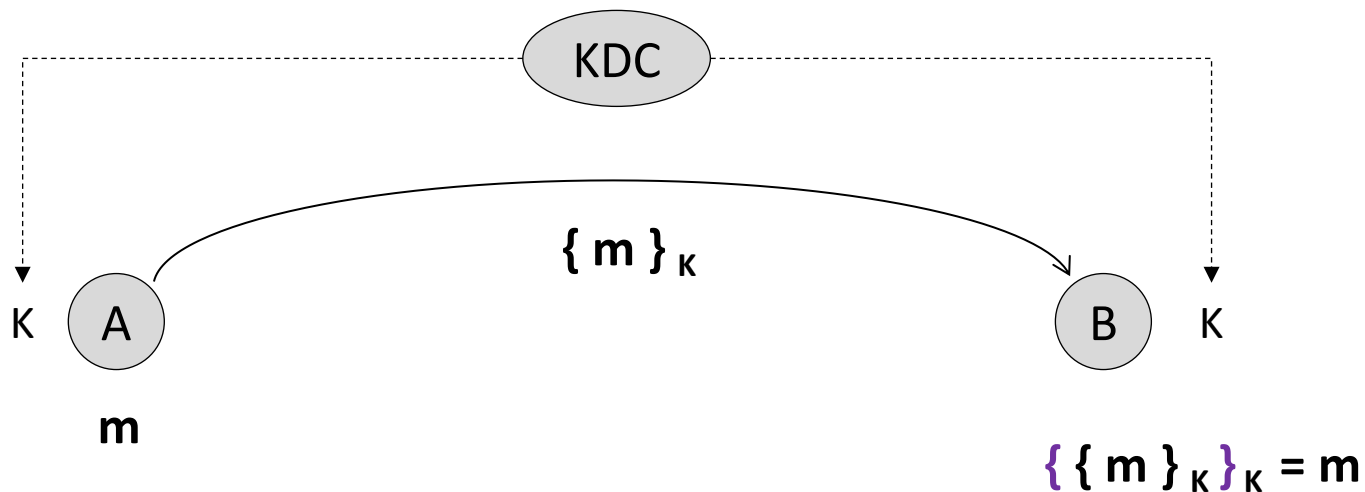


*Alice creates message m ,
encrypts using shared key
 k , and sends result to B*

Conventional Cryptography

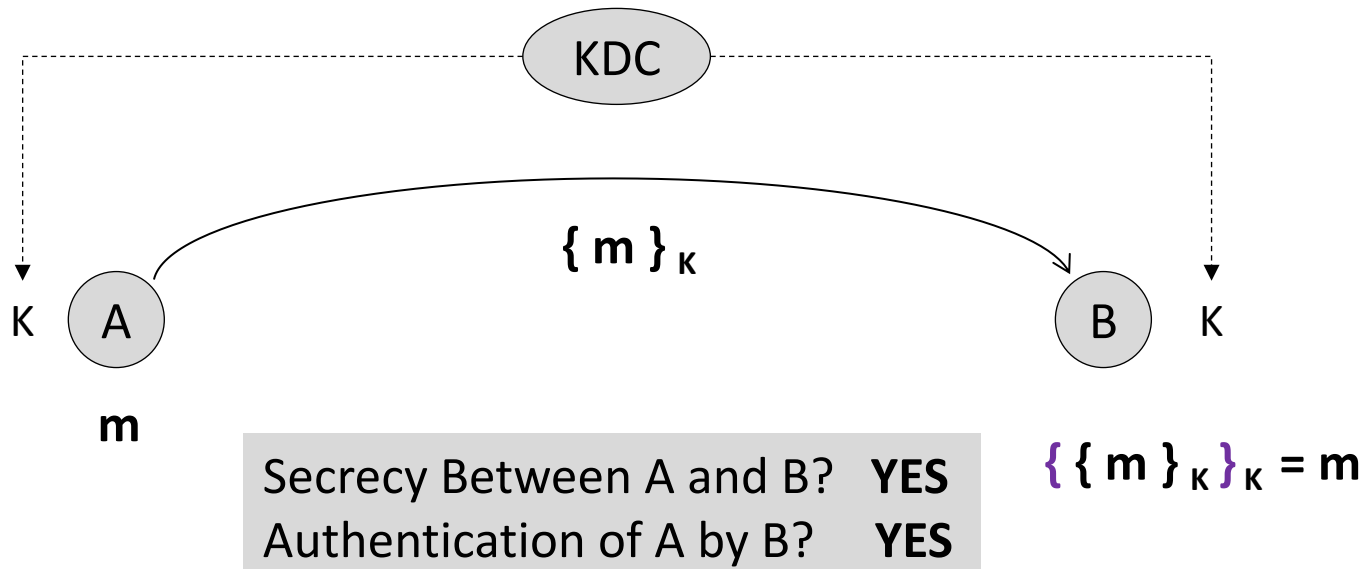


Conventional Cryptography

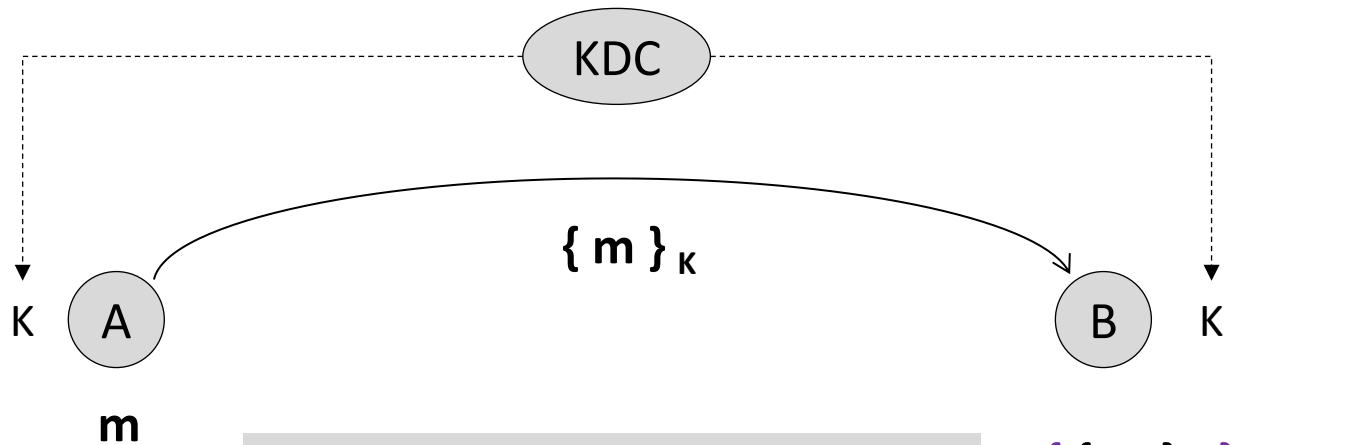


Bob receives encrypted message, and decrypts using shared key k , and obtains message m

Conventional Cryptography



Conventional Cryptography



Secrecy Between A and B? **YES**
Authentication of A by B? **YES**

$$\{\{m\}_K\}_K = m$$

Does this approach scale? **NO**

What are the Basic Properties of
Public Key Cryptography?

Public Key Cryptography Basics

Two Communicants: A and B

1. A generates pair of keys P_A and S_A
2. B generates pair of keys P_B and S_B

Public Key Cryptography Basics

Two Communicants: A and B

1. A generates pair of keys PA and SA
2. B generates pair of keys PB and SB
3. Properties:

$$\{ \{ m \}_{PA} \}_{SA} = m$$

$$\{ \{ m \}_{SA} \}_{PA} = m$$

$$\{ \{ m \}_{PA} \}_X = m \Rightarrow (X = SA)$$

$$\{ \{ m \}_{SA} \}_X = m \Rightarrow (X = PA)$$

*Concept proposed by Whit Diffie
and Marty Hellman, Stanford and
Ralph Merkle, UC Berkeley – circa 1976*



Public Key Cryptography Basics

Two Communicants: A and B

1. A generates pair of keys P_A and S_A
2. B generates pair of keys P_B and S_B
3. Properties:

$$\{ \{ m \}_{P_A} \}_{S_A} = m$$

$$\{ \{ m \}_{S_A} \}_{P_A} = m$$

$$\{ \{ m \}_{P_A} \}_X = m \Rightarrow (X = S_A)$$

$$\{ \{ m \}_{S_A} \}_X = m \Rightarrow (X = P_A)$$

*Concept proposed by Whit Diffie
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Requirements:

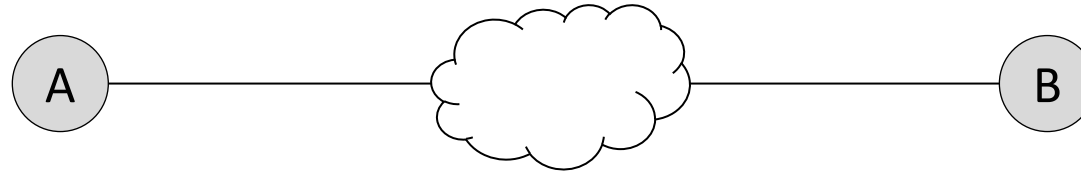
- (i) Keep S_A , S_B secret to A, B
- (ii) Make P_A , P_B public to all
- (iii) No KDC required to generate keys

“Address Scaling Issue”



Understanding Public Key Technology

*“Assume A
is a client”*



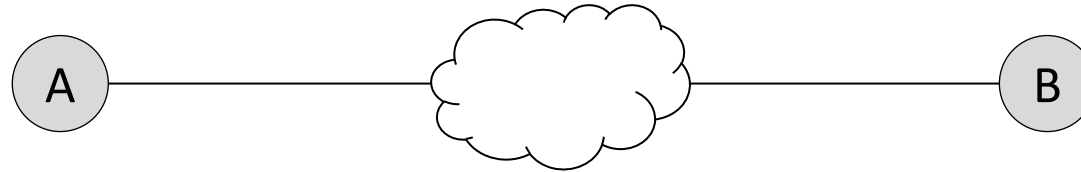
*“Assume B
is a server”*

Understanding Public Key Technology

*"Assume A
is a client"*

*No Key Distribution
Center (KDC) Required*

*"Assume B
is a server"*



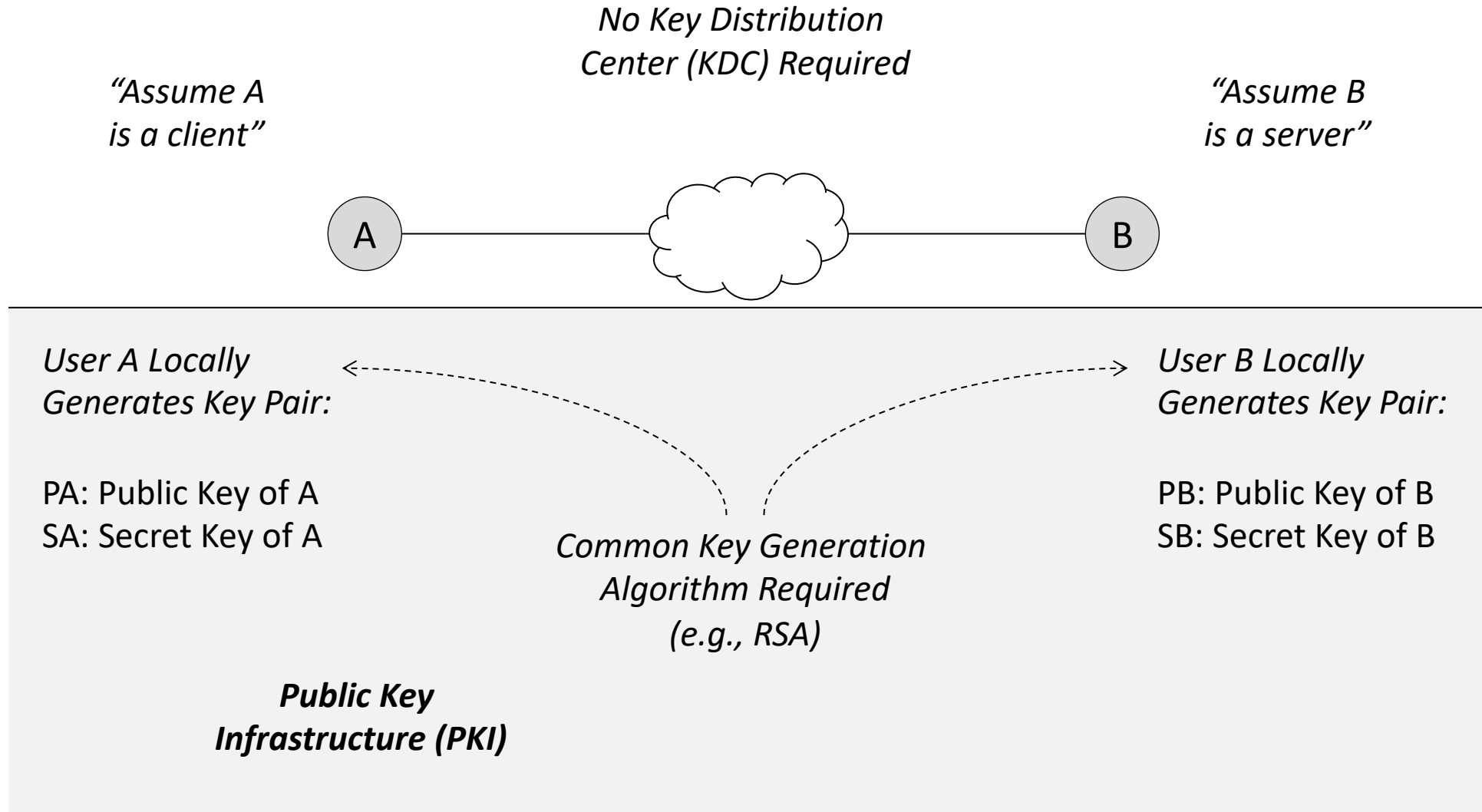
*User A Locally
Generates Key Pair:*

PA: Public Key of A
SA: Secret Key of A

*User B Locally
Generates Key Pair:*

PB: Public Key of B
SB: Secret Key of B

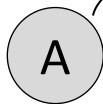
Understanding Public Key Technology



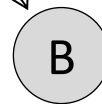
Sending a Secret Message

Alice creates message m . . .

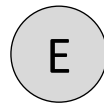
PA, SA, PB



m



PB, SB, PA

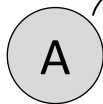


PA, PB, PE, SE

Sending a Secret Message

*Alice creates message m ,
encrypts using Bob's public key
 PB , and sends result to B*

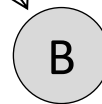
PA, SA, PB



$\{ m \}_{PB}$

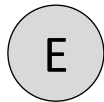
B

PB, SB, PA

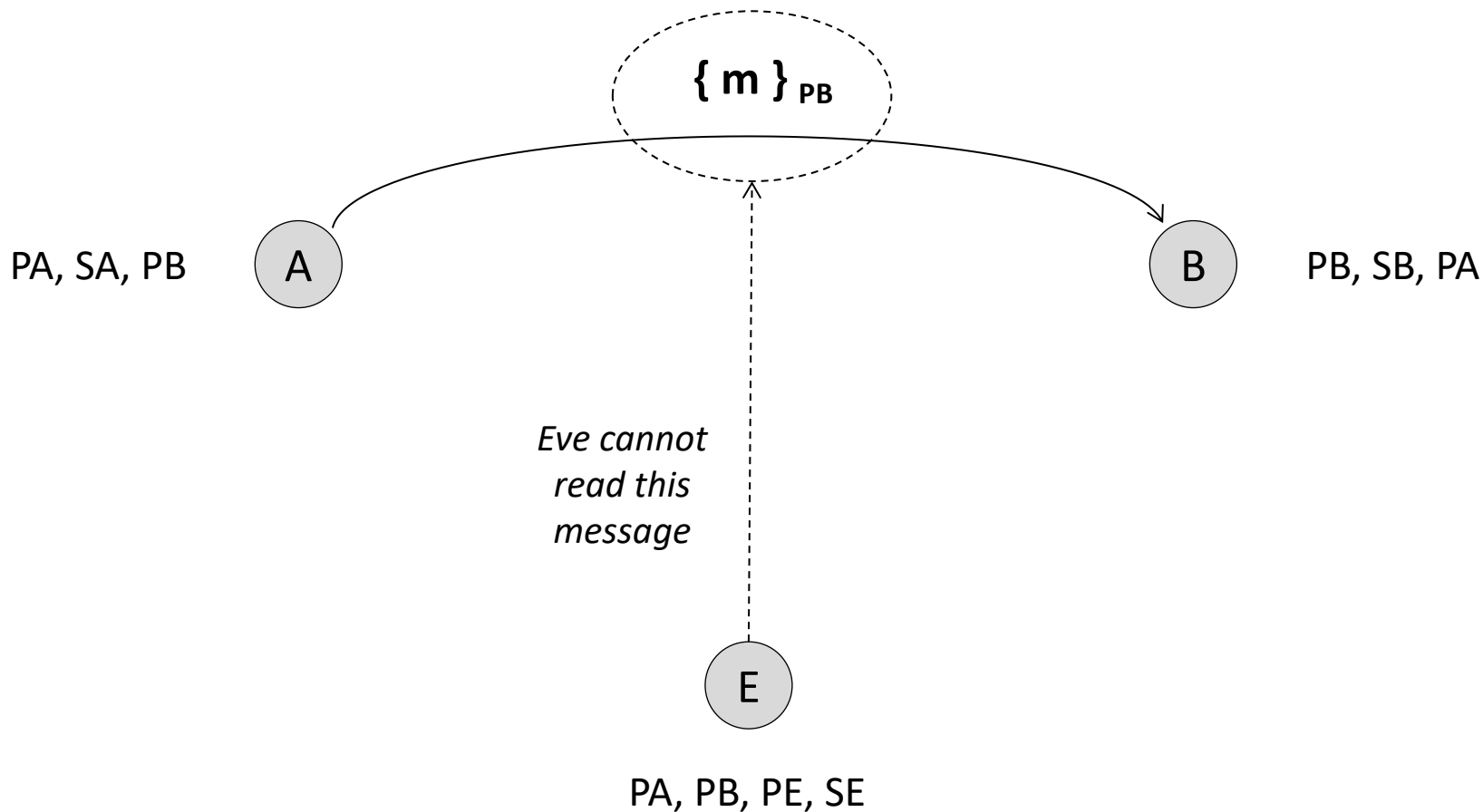


E

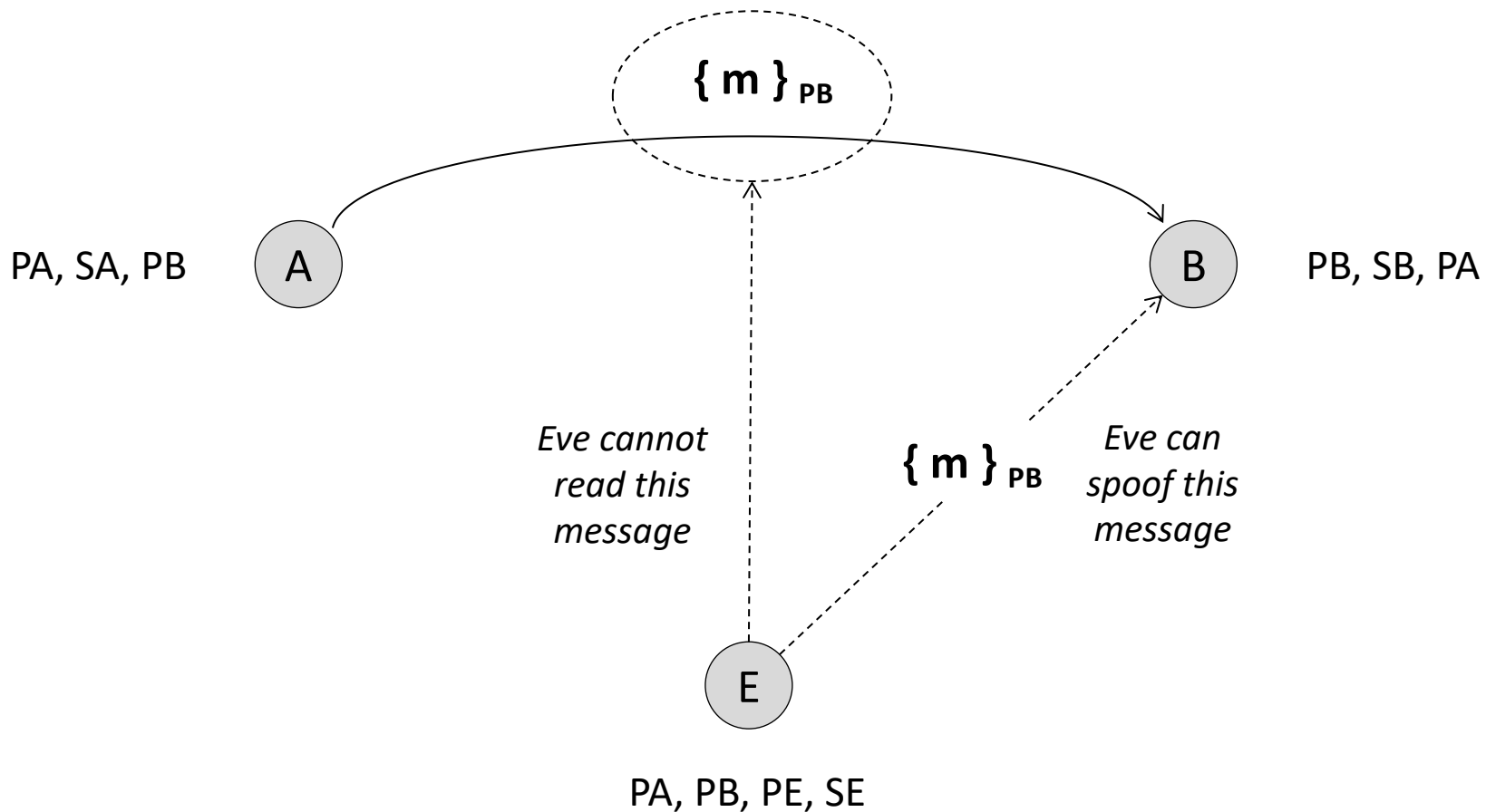
PA, PB, PE, SE



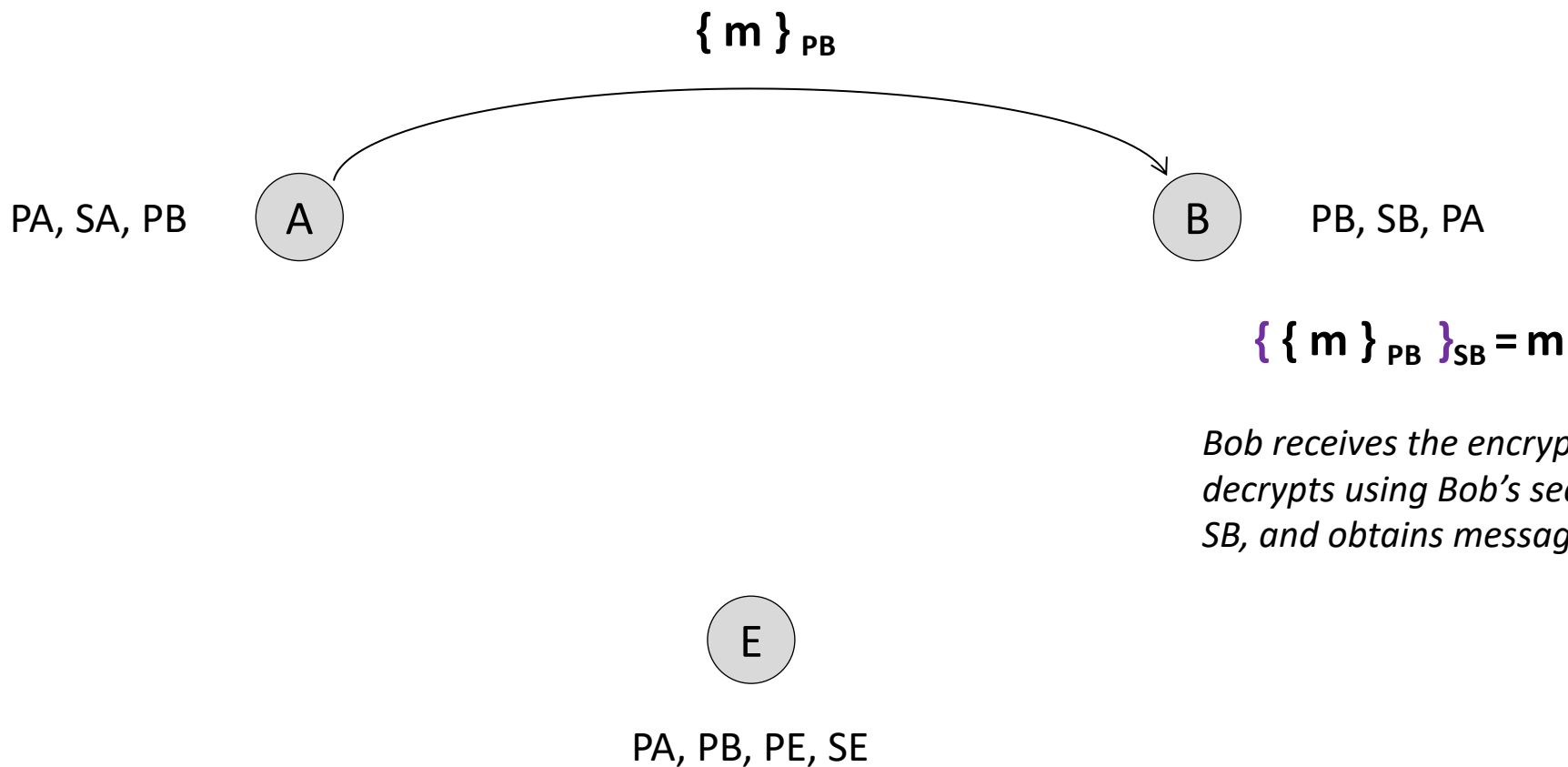
Sending a Secret Message



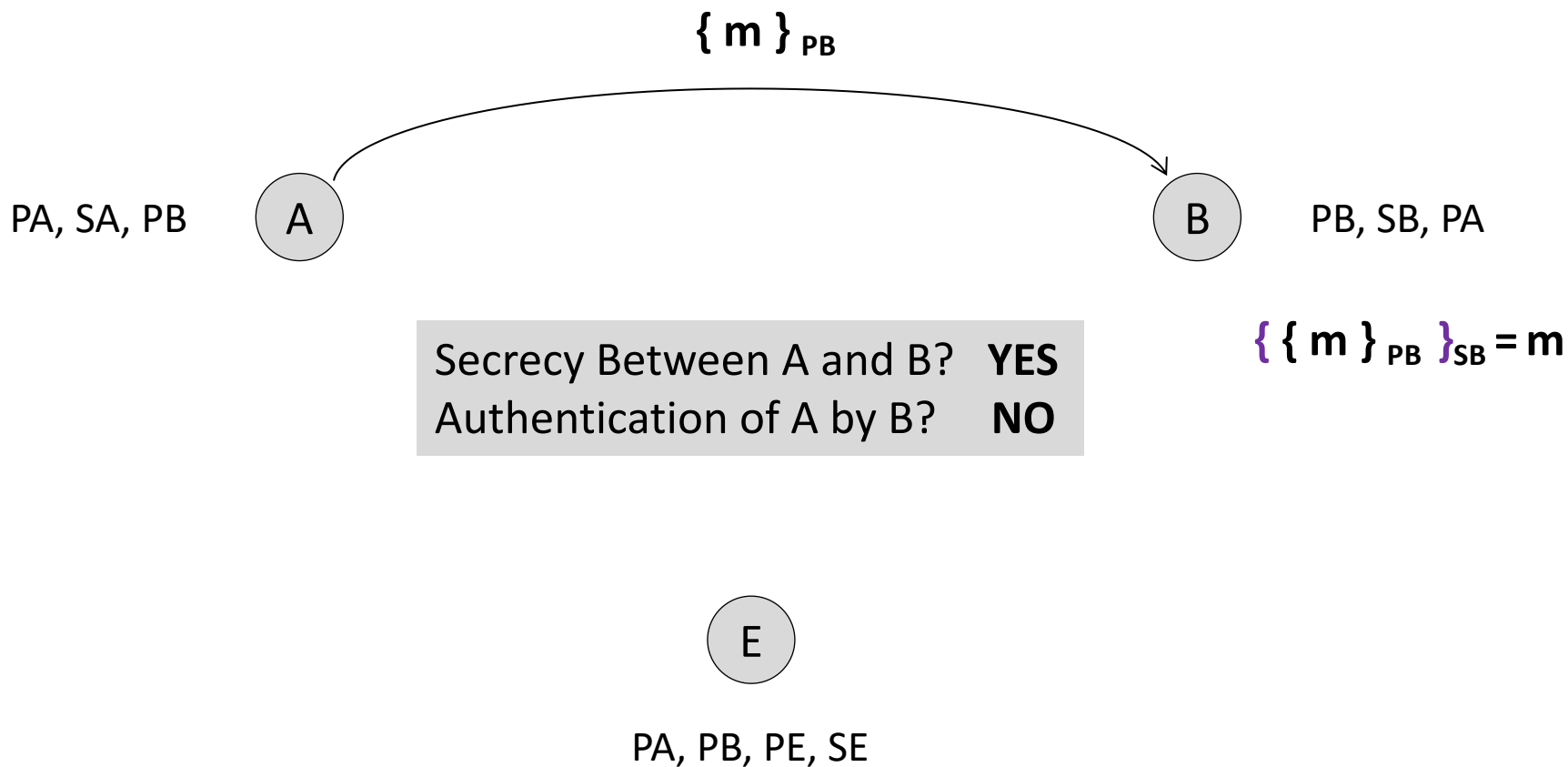
Sending a Secret Message



Sending a Secret Message

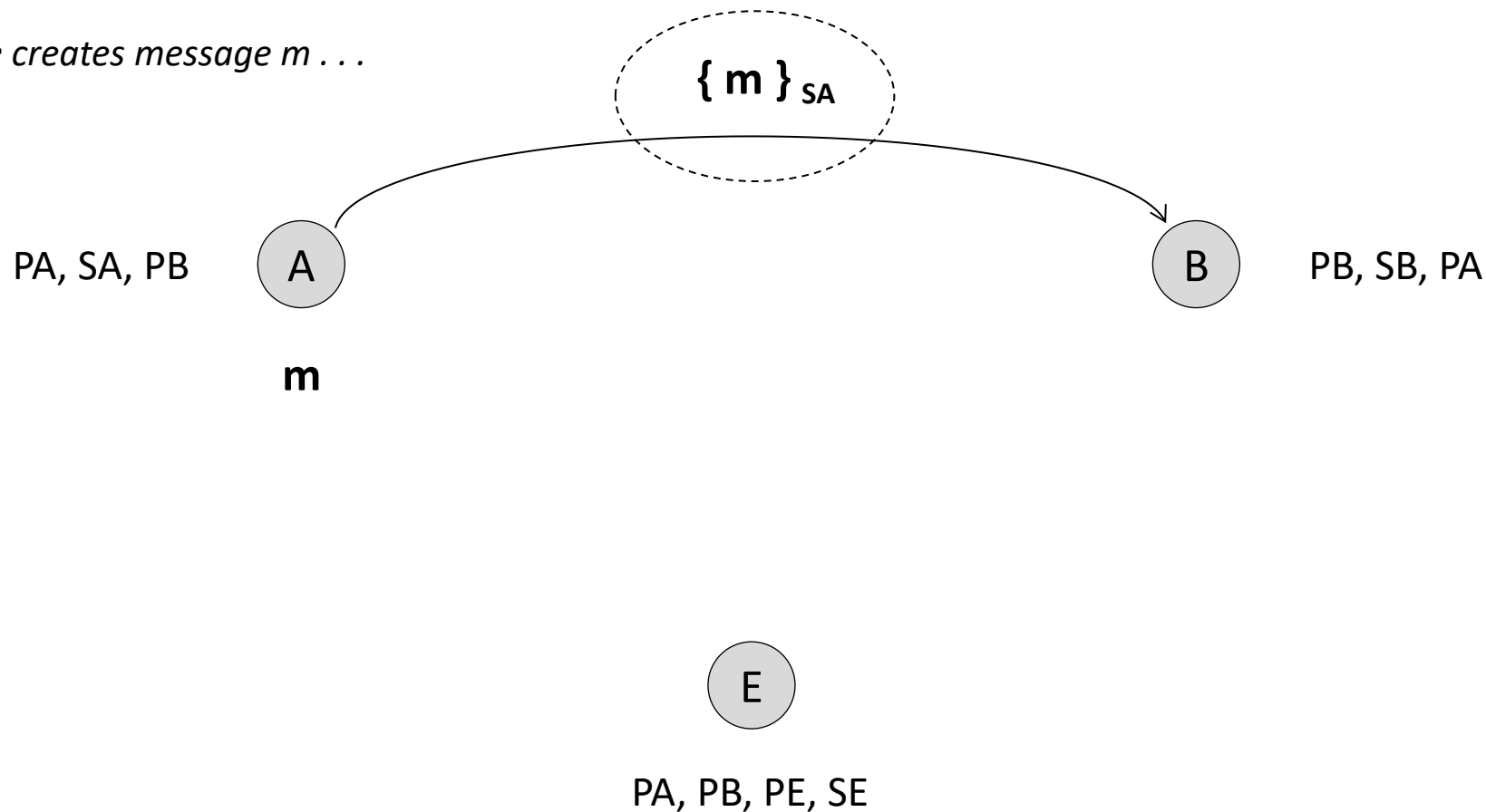


Sending a Secret Message



Sending a Signed Message

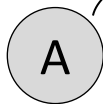
Alice creates message m . . .



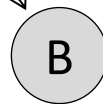
Sending a Signed Message

*Alice creates message m ,
encrypts using Alice's secret key
 SA , and sends result to B*

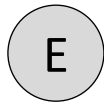
PA, SA, PB



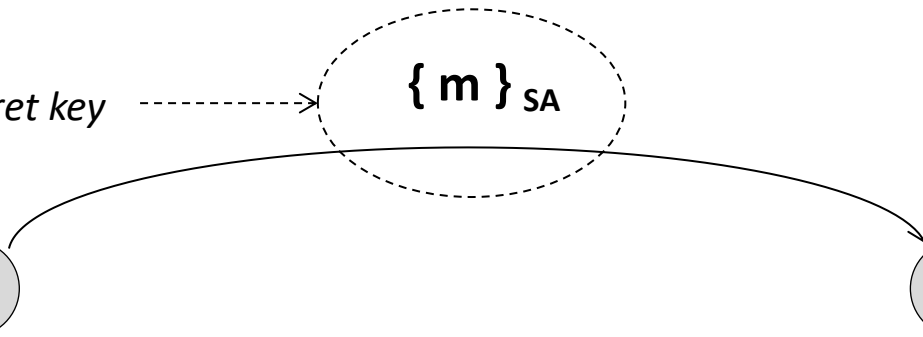
$\{ m \}_{SA}$



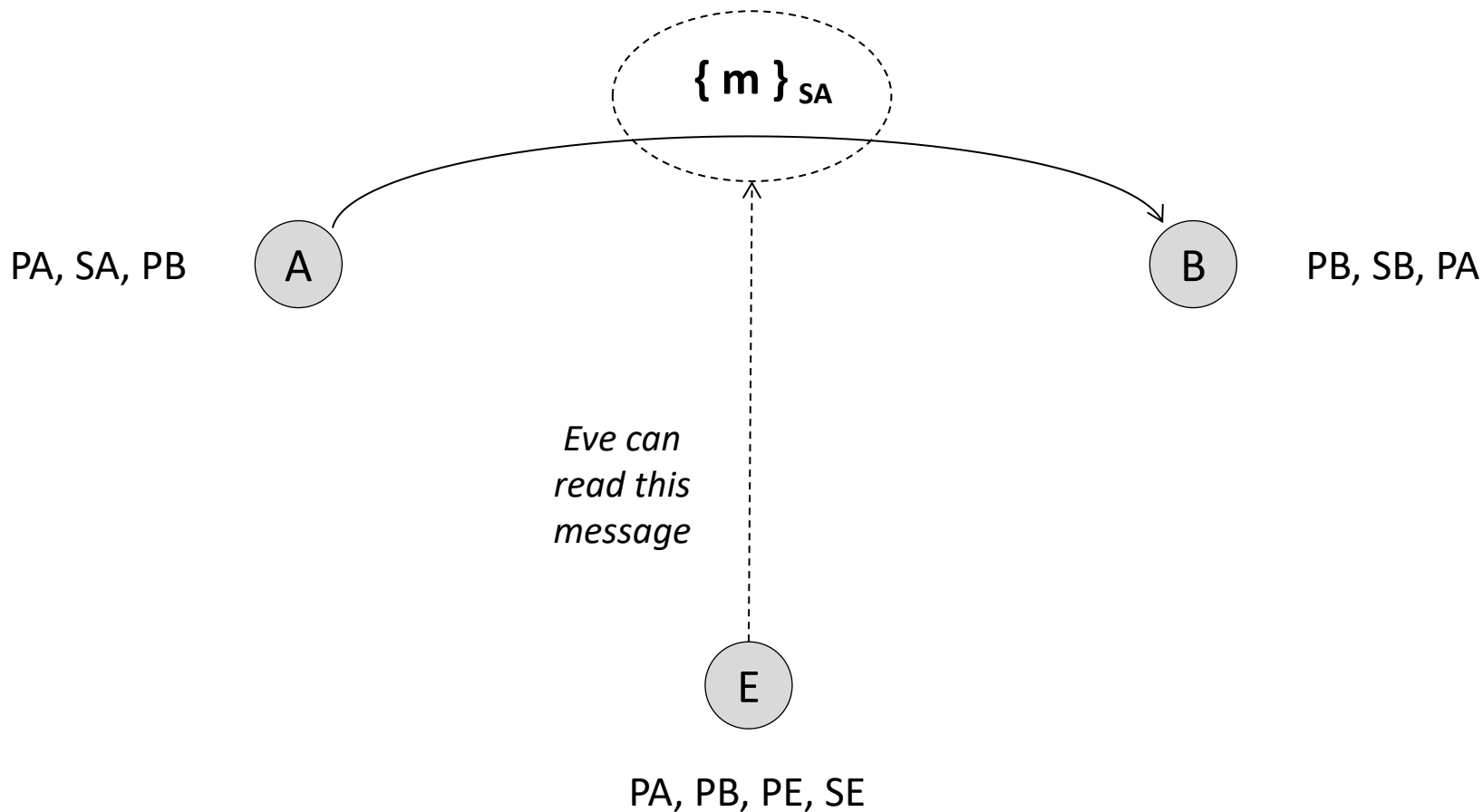
PB, SB, PA



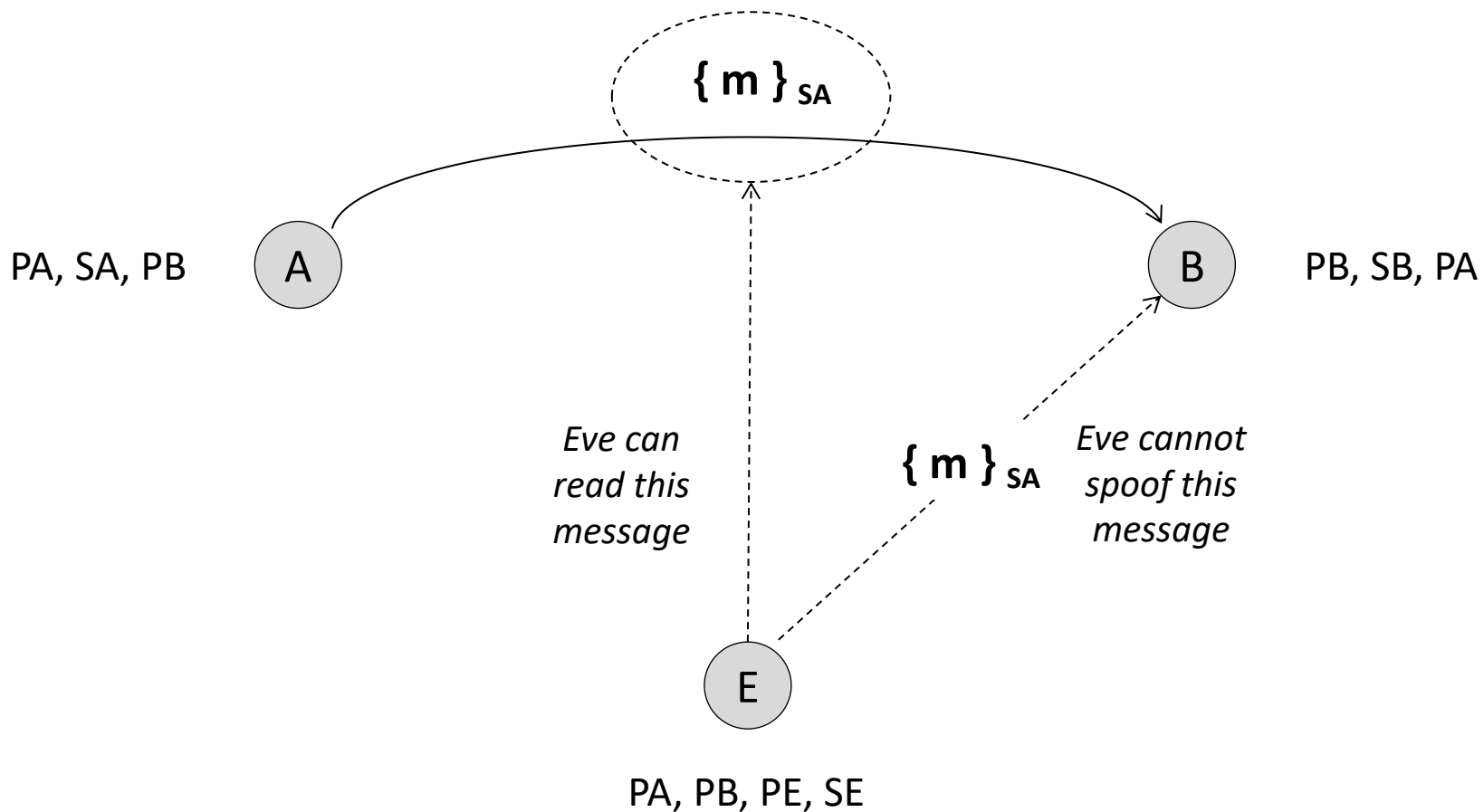
PA, PB, PE, SE



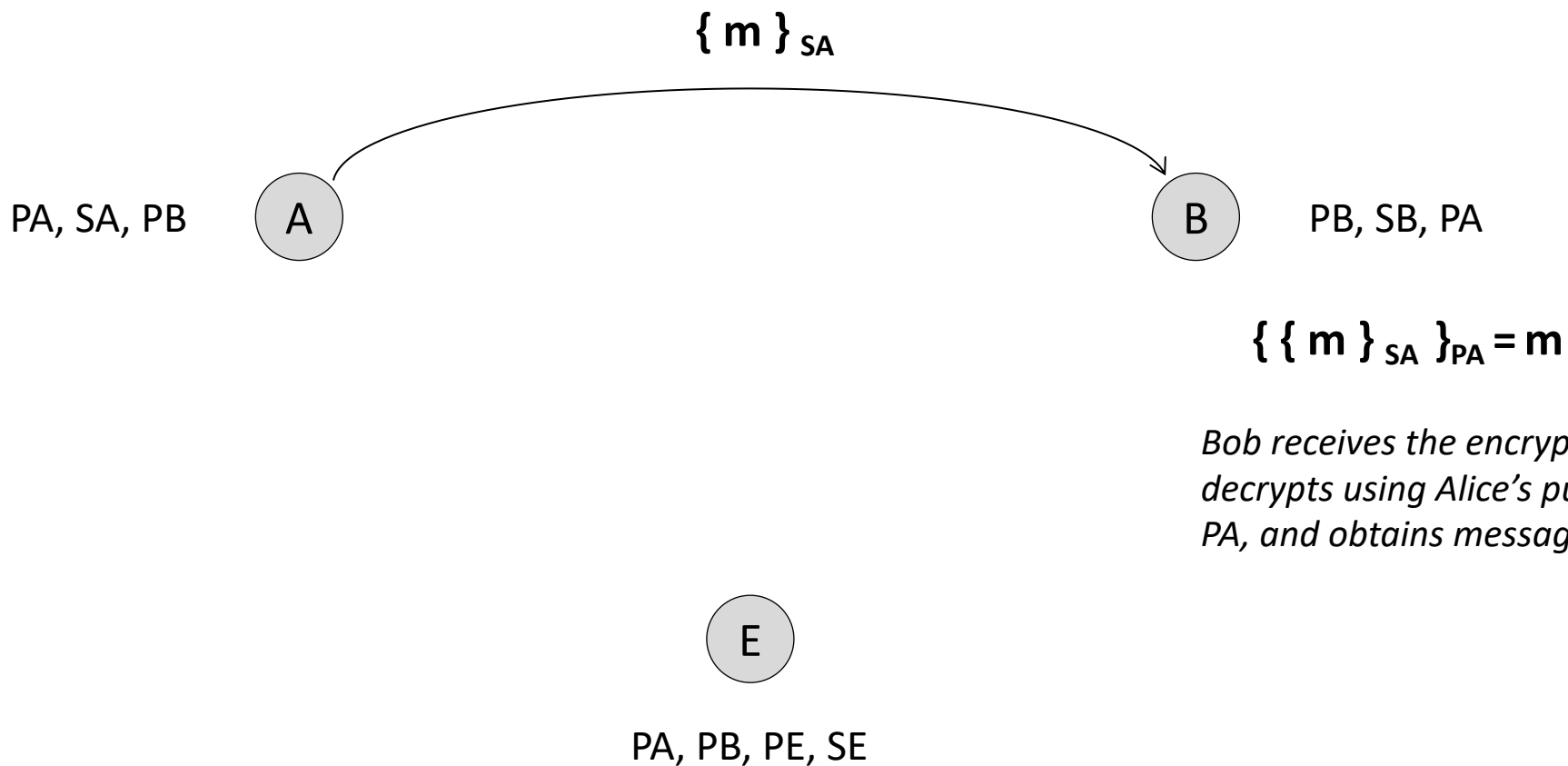
Sending a Signed Message



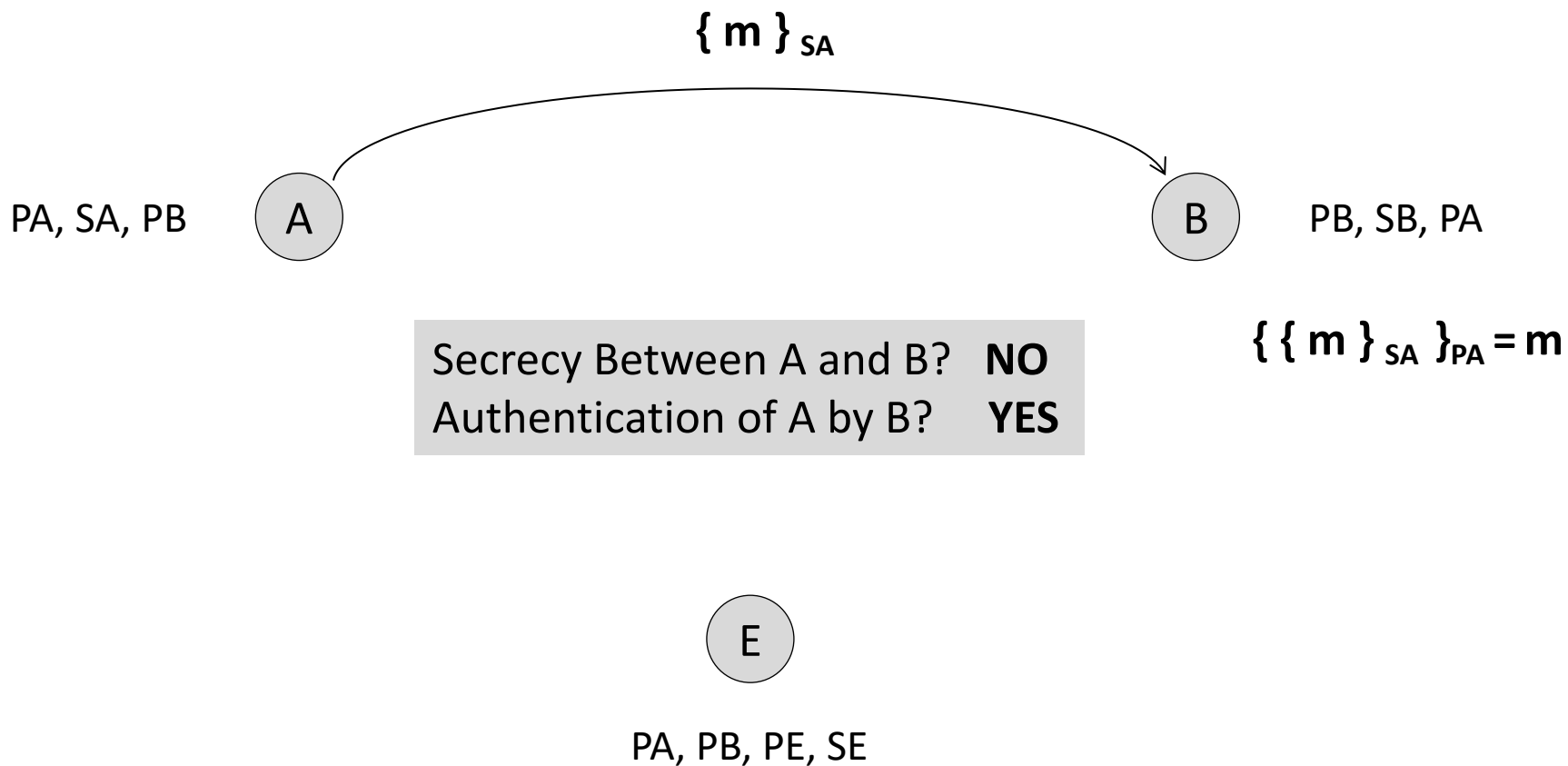
Sending a Signed Message



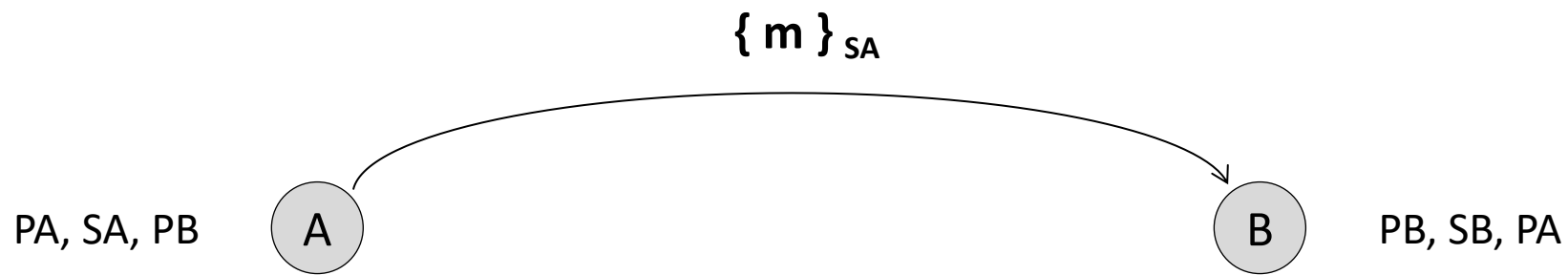
Sending a Signed Message



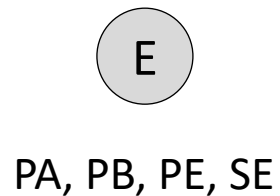
Sending a Signed Message



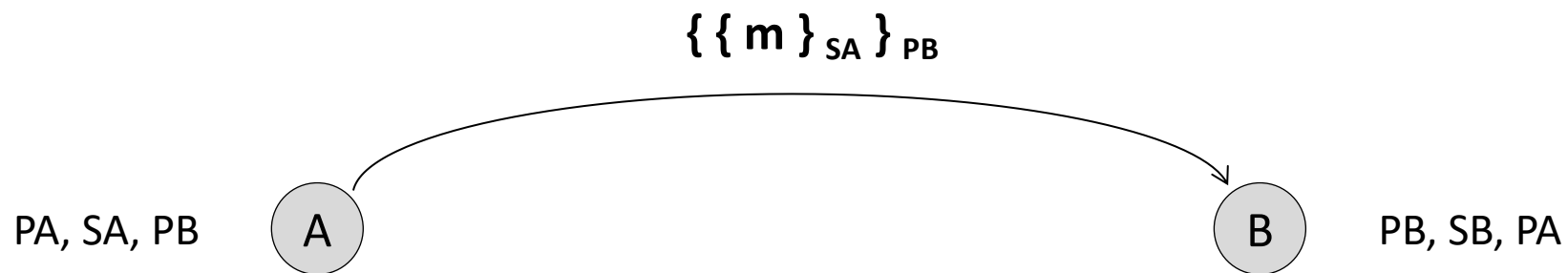
Secure Message Exchange



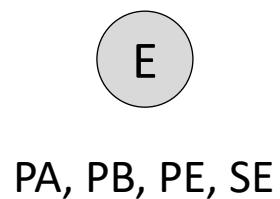
*Alice creates a message m ,
encrypts it with a public key algorithm
using her secret key SA . . .*



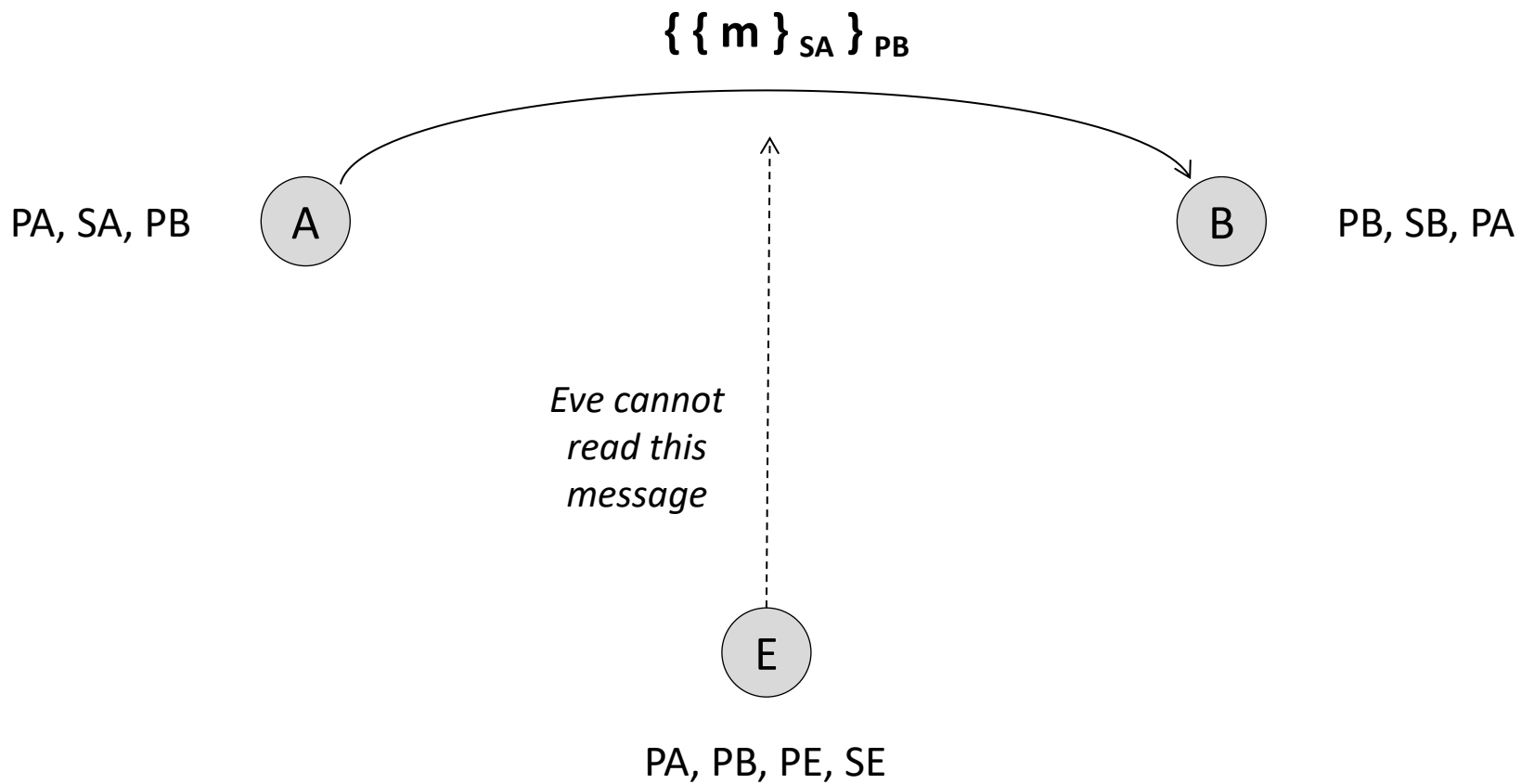
Secure Message Exchange



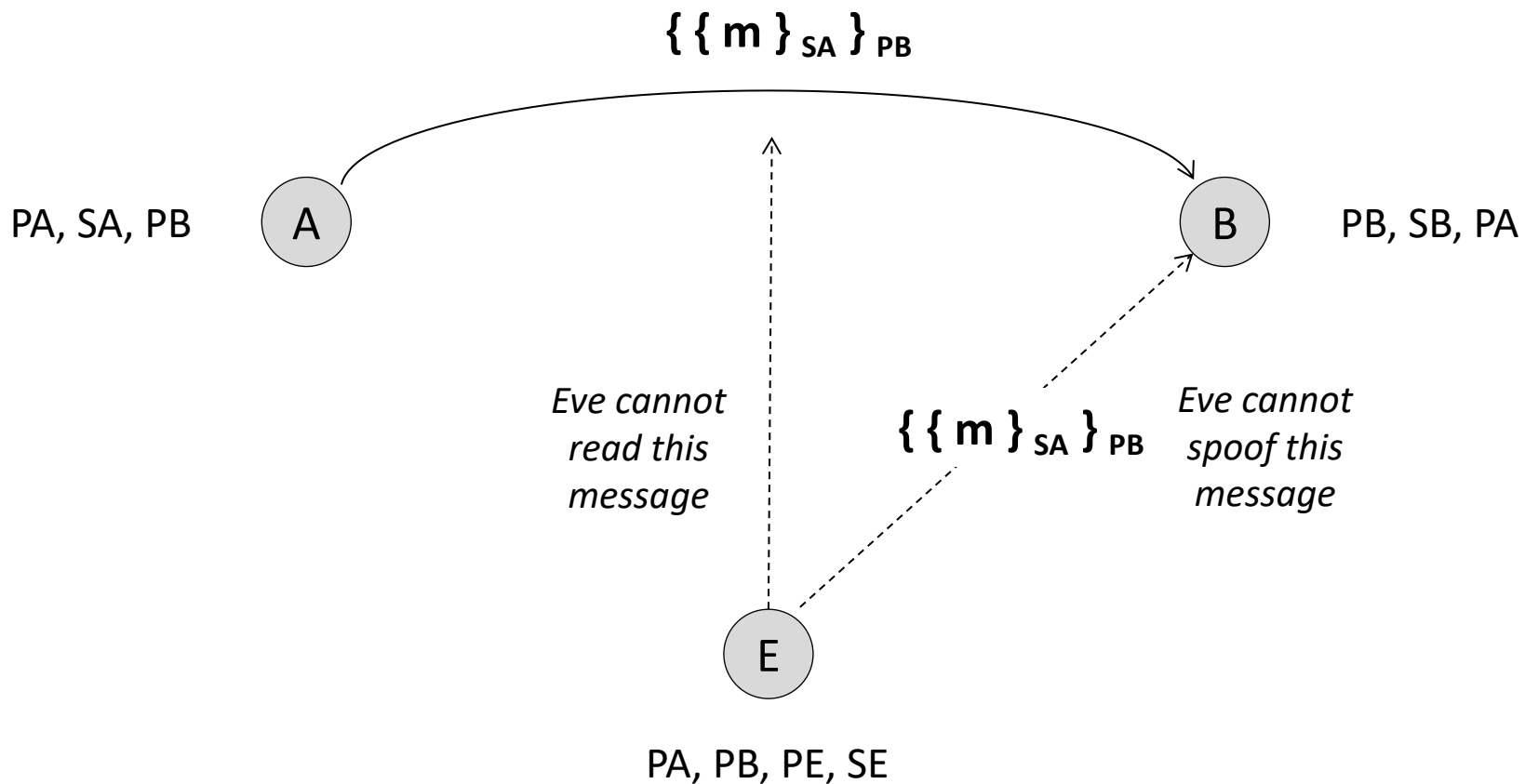
*Alice creates a message m ,
encrypts it with a public key algorithm
using her secret key SA , encrypts it again
using a public key algorithm with Bob's
public key PB , and sends the result to Bob*



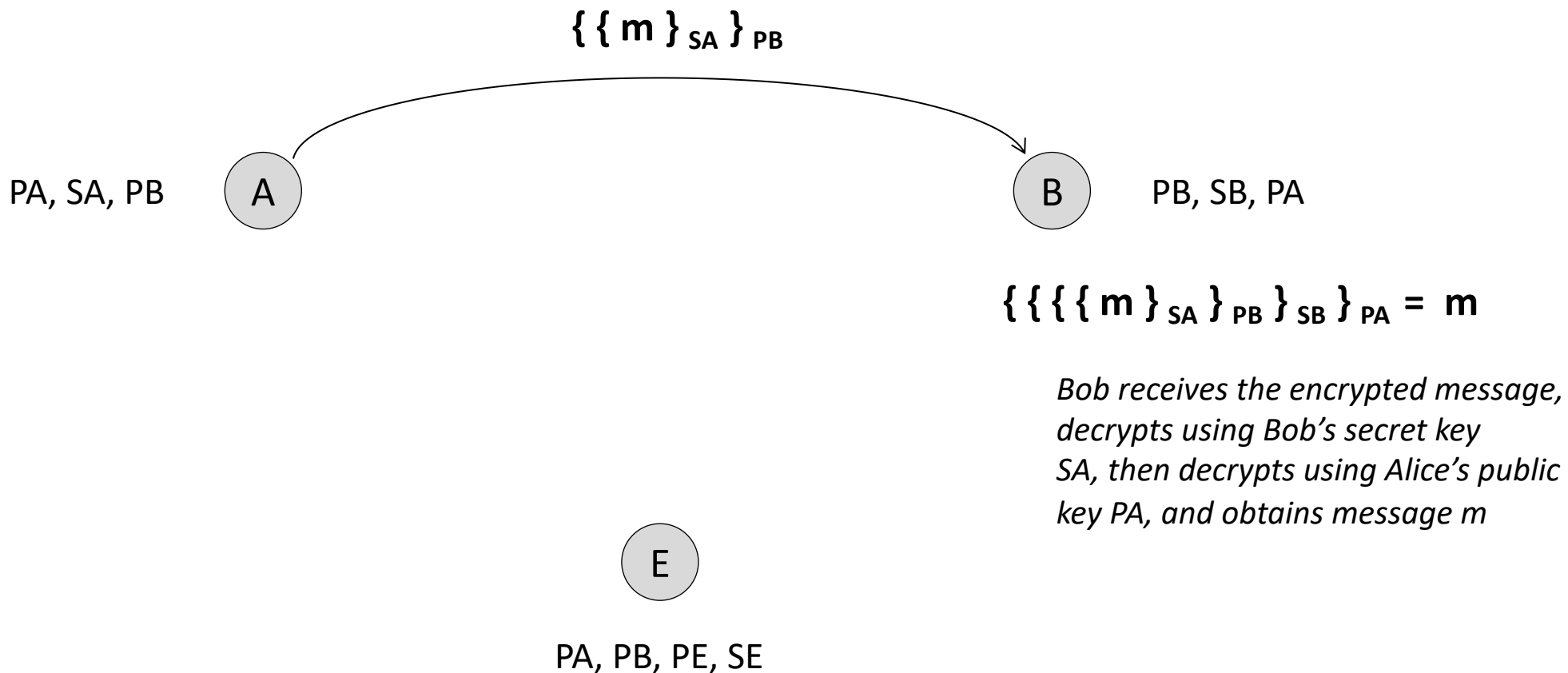
Secure Message Exchange



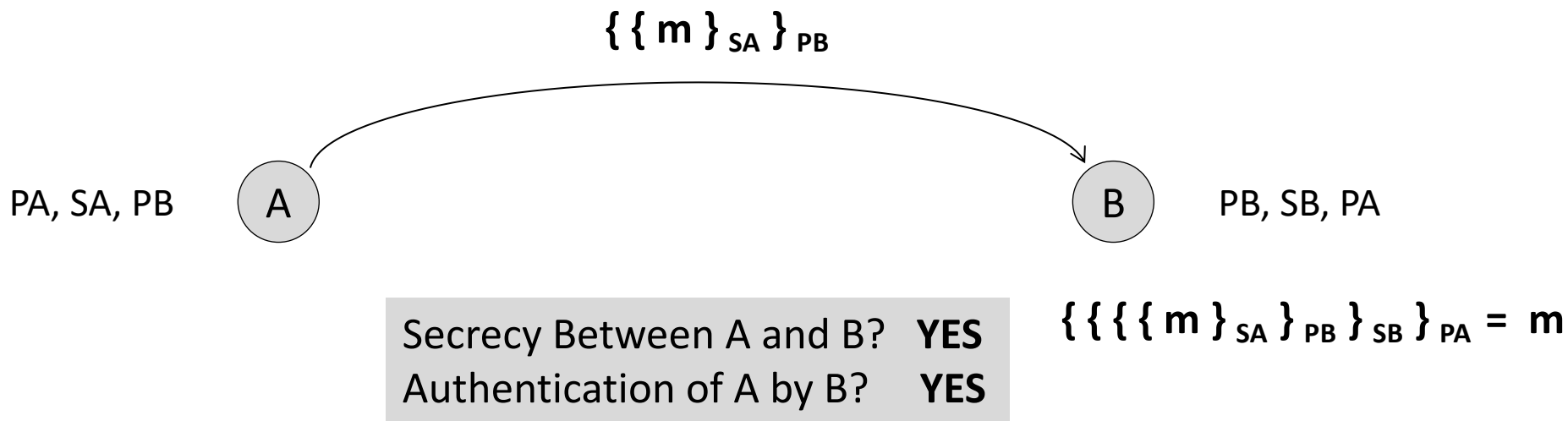
Secure Message Exchange



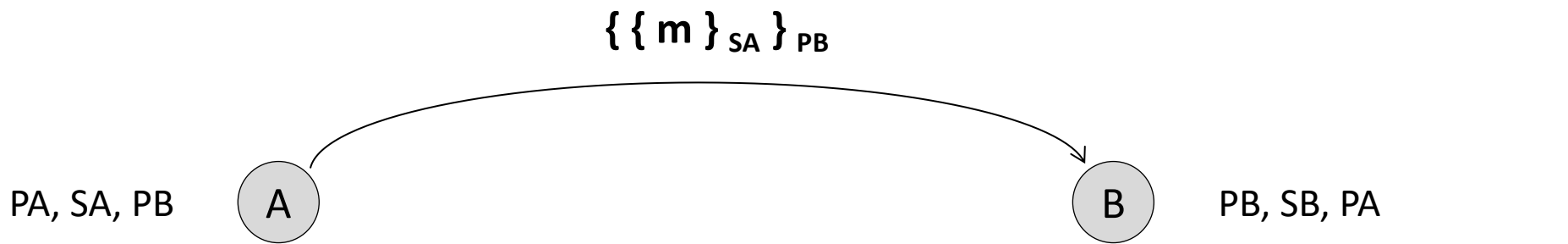
Secure Message Exchange



Secure Message Exchange



Secure Message Exchange

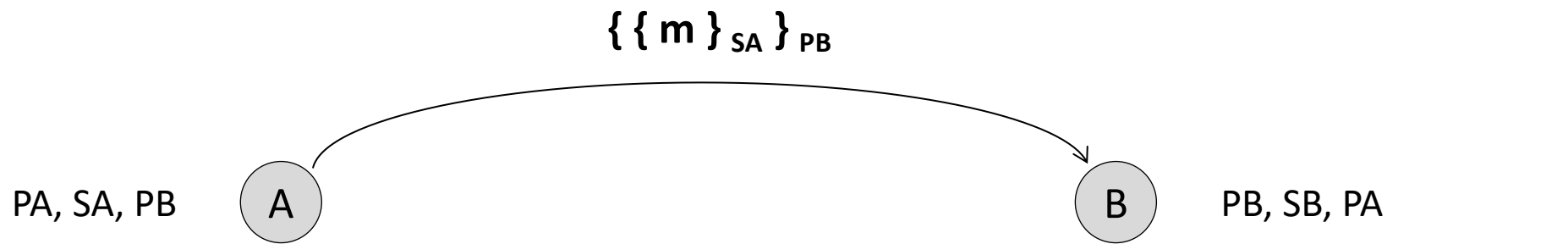


Secrecy Between A and B? **YES**
 Authentication of A by B? **YES**

$$\{\{\{\{m\}_{SA}\}_{PB}\}_{SB}\}_{PA} = m$$

Does this approach scale? **YES**

Secure Message Exchange



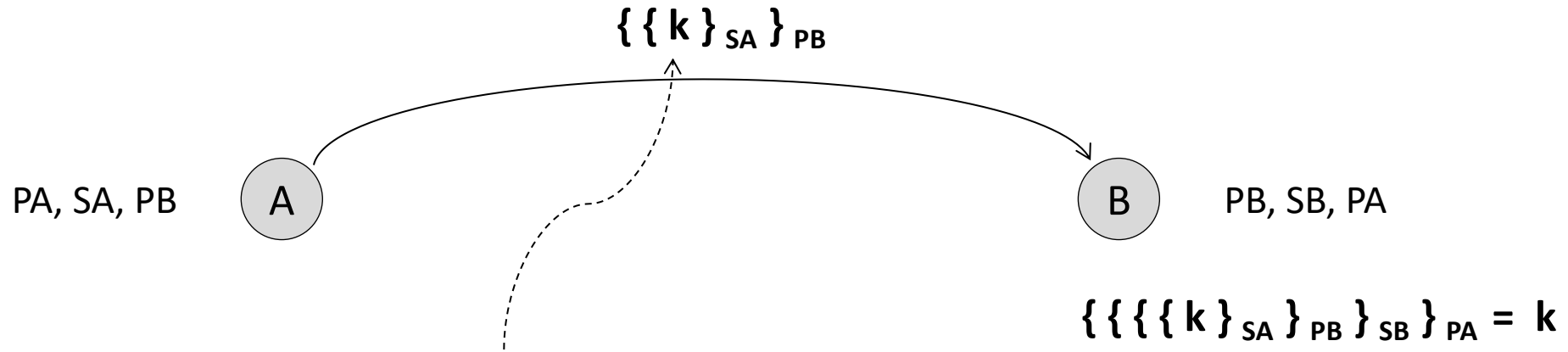
Secrecy Between A and B? **YES**
 Authentication of A by B? **YES**

$$\{ \{ \{ \{ m \}_{SA} \}_{PB} \}_{SB} \}_{PA} = m$$

Does this approach scale? **YES**

Is this approach efficient (cryptographically)? **NO**

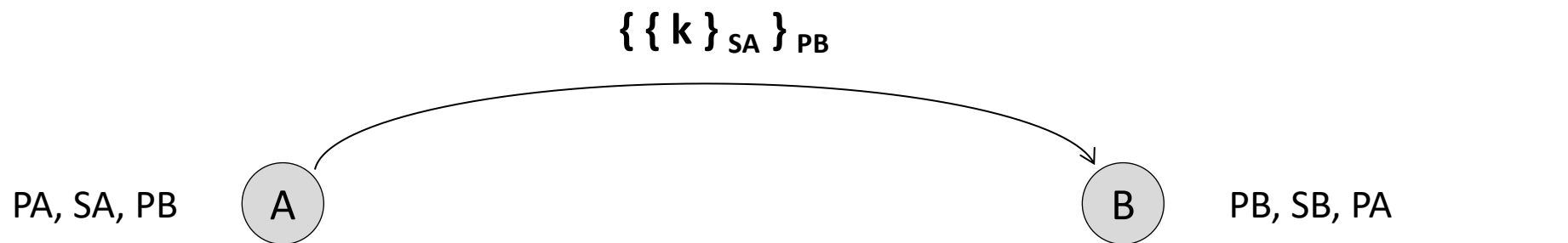
Secure Key Exchange



Alice generates a key k for some bulk encryption algorithm (like 3-DES) and provides this key to B using secure key exchange

- Scalable
- Secret
- Authenticated

Secure Key Exchange



Secrecy Between A and B? **YES**
 Authentication of A by B? **YES**

$$\{\{\{\{k\}_{SA}\}_{PB}\}_{SB}\}_{PA} = k$$

Does this approach scale? **YES**

Is this approach efficient (cryptographically)? **YES**