

⑦ Suppose within your browser you can click on a link to obtain a webpage. The IP address for the associated URL is not cached in your local host, so a DNS lookup is necessary to obtain the IP address. Suppose that n DNS servers are visited before your host receives the IP address from DNS, the successive visits incur an RTT of RTT_1, \dots, RTT_n . Further suppose that the web page associated with link contains exactly one object, consisting of a small amount of HTML text. Let RTT_0 denote the RTT between the local host & the server containing the object. Assuming zero transmission time of the object, how much time elapses from when the client clicks on a link until the client receives the object?

⇒ consider the IP address of total amount of time:

$$RTT_1 + RTT_2 + \dots + RTT_n$$

Time elapses from when the client clicks on the link until the client receives the object:

$$2RTT_0 + RTT_1 + RTT_2 + \dots + RTT_n$$

⑧ Consider figure 2.12 for which there is an institutional network connected to Internet. Suppose that the avg object size is 850,000 bits and that the average request rate from the institution's browser to the origin server is 16 requests per second. Also suppose that the amount of time it takes from router on the Internet side of the access link forward an HTTP request until it receives the response is 3 sec on average. Use $\Delta / (1 - \Delta/\beta)$, where Δ is the average time required to send an object over the access link & β is the arrival rate of object to access link.

a) Find the total avg. response time.

b) Now suppose a cache installed in the institutional LAN, suppose the miss rate is 0.4. Find total response time.

⇒ a) The time to transmit an object of size L over a link of rate R is L/R . The avg time is the average size of the object divided by R .

$$= \frac{850,000 \text{ bits}}{(15,000,000 \text{ bits/sec})} \Rightarrow \underline{0.0567 \text{ sec}}$$

⊗ The traffic intensity on the link is given by $\rho \Delta^2$
 $(16 \text{ requests/sec}) (0.0567 \text{ sec/request})$
 $= \underline{0.907}$

Thus avg. access delay is $\frac{(0.0567 \text{ sec})}{(1 - 0.907)}$
 $\approx \underline{0.6 \text{ sec}}$

∴ The total average response time is $0.6 \text{ sec} + 3 \text{ sec}$
 $= \underline{3.6 \text{ sec}}$

b) The traffic intensity on the access link is reduced by 60% since the 60% of the requests are satisfied within institutional network. Thus the avg. access delay is $\frac{(0.0567 \text{ sec})}{(4)(0.907)} = \underline{0.0089 \text{ sec}}$.

The response time is approximately zero if the request is satisfied by the cache.
 the average response time is $0.0089 \text{ sec} + 3 \text{ sec} \approx 3.0089 \text{ sec}$
 for cache misses (which happens 40% of the time).

So, the average response time is $(0.67 \text{ sec}) + (4)(3.0089 \text{ sec}) = \underline{1.24 \text{ sec}}$

Thus the avg response time is reduced from ~~3.6~~ 3.65 to 1.24 sec

22) Consider distributing a file of $F = 15$ Gbits to N peers. The server has an upload rate of $u_s = 30$ Mbps, and each peer has a download rate of $d_i = 2$ Mbps & an upload rate of u_i . For $N = 10, 100, 1000$ and $u_i = 300$ kbps, 700 kbps & 2 Mbps, prepare a chart giving in the minimum distribute time for each of the combinations of N and u_i for both client-server distribution and P2P distribution.

⇒ For calculating the minimum distribute time for client-server distribution, we use the following formula:

$$D_{cs} = \max\{NF/u_s, F/d_{\min}\}$$

Similarly for calculating the minimum distribute time for P2P distribution, we use the formula:

$$D_{p2p} = \max\{F/u_s, F/d_{\min}, NF/(u_s + \sum_{i=1}^N u_i)\}$$

Given

$$F = 15 \text{ Gbits} = 15360 \text{ Mbits. } u_s = 30 \text{ Mbps } d_i = 2 \text{ Mbps} = d_{\min}$$

$$\text{Note: } 300 \text{ kbps} = 300/1024 \text{ Mbps}$$

client-server

	10	100	1000
300 kbps	7680	51200	512000
700 kbps	7680	51200	512000
2 Mbps	7680	51200	512000

example to solve

$$\begin{aligned} D_{cs} &= \max\left\{\frac{NF}{u_s}, \frac{F}{d_{\min}}\right\} \\ &= \max\left\{\frac{10 \times 15360}{30}, \frac{15360}{2}\right\} \\ &= \max\{5120, 7680\} \\ &= 7680 \text{ sec} \end{aligned}$$

Peer to peer

	10	100	1000
300 kbps	7680	25906	47569
700 kbps	7680	15616	21525
2 Mbps	7680	7680	7680

solved one case.

$$\begin{aligned} u_s &= 30 \text{ Mbps} \\ F &= 15360 \text{ Mbits} \\ u_i &= 300 \text{ kbps} \\ &= \frac{300}{1024} = 0.2929 \text{ Mbps} \\ \therefore \sum_{i=1}^N u_i &= \sum_{i=1}^{100} 0.2929 \end{aligned}$$

$$\begin{aligned} D_{p2p} &= \max\left\{\frac{F}{u_s}, \frac{F}{d_{\min}}, \frac{NF}{u_s + \sum_{i=1}^N u_i}\right\} = 100 \times 0.2929 \\ &= \max\left\{\frac{15360}{30}, \frac{15360}{2}, \frac{100 \times 15360}{30 + 29.29}\right\} \\ &= \max\{512, 7680, 25906\} = 25906 \text{ sec} \end{aligned}$$

26) Suppose Bob joins a bit-torrent torrent, but he does not want to upload any data to any other peers.

a) Bob claims that he can receive a complete copy of the file that is shared by the swarm. Is Bob's claim possible? why or why not?

b) Bob further claims that he can further make his "free-riding" more efficient by using a collection of multiple computers (with distinct IP addresses) in the computer lab in his department. How can he do that?

⇒ a) Yes his first claim is possible as long as there are enough peers staying in the swarm for a long enough time Bob can always receive data through optimistic unchoking by other peers.

b) His second claim is also true. he can run a client on each host, let each client "free-ride," and combine chunks from the different hosts into a single file. He can even write a small scheduling program to make the different hosts ask for different chunks of file. This is actually a kind of Sybil attack in P2P network.

28). Install & compile the python prog. TCP client & UDP client on the host & TCP server & UDP server on another host.

a. Suppose you run TCP client before you run TCP server, what happens? why?

b. Suppose you run UDP client before you run UDP server, what happens? why?

c. What happens if you use different port numbers for the client & server sides?

⇒ a) If you run TCP client first, then client will attempt to make a TCP connection with a non-existent server process. A TCP connection will not be made.

- b) UDP client doesnot establish a TCP connection with the server. That everything should work fine if you first run UDP client, then run UDP server, and then type some input into the keyboard.
- c) If you use a different port number, then the client will attempt to establish a TCP connection with the wrong process or a non-existent process. Error will occur.
-

2d) Continue
TCP client

from socket import *

```
def client():
```

```
    HOST = "10.211.55.91"
```

```
    PORT = 10521
```

```
    BUFSIZE = 1024
```

```
    clientsocket = socket(AF_INET, SOCK_STREAM)
```

```
    clientsocket.connect((HOST, PORT))
```

```
    while True:
```

```
        data = input('>')
```

```
        if not data:
```

```
            break
```

```
        clientsocket.send(data.encode())
```

```
        data = clientsocket.recv(BUFSIZE).decode()
```

```
    if not data:
```

```
        break
```

```
    print(data)
```

```
client()
```

TCP server

from socket import *

from time import time

```
def server():
```

```
    HOST = ''
```

```
    PORT = 10521
```

```
    BUFSIZE = 1024
```

```
    ADDR = (HOST, PORT)
```

```
server_socket = socket(AF_INET, SOCK_STREAM)
```

```
server_socket.bind(ADDR)
```

```
server_socket.listen(5)
```

```
while True:
```

```
    print('waiting for connecting...')
```

```
    tcp_client_socket, addr = server_socket.accept()
```

```
    print('connected from', addr)
```

```
    while True:
```

```
        data = tcp_client_socket.recv(BUFSIZE).decode()
```

```
        if not data:
```

```
            break
```

```
        print(data)
```

```
        data = input('>')
```

```
        tcp_client_socket.send(('[ %s] %s' % (time(), data)).encode())
```

```
    tcp_client_socket.close()
```

```
server_socket.close()
```

```
server()
```