Dynamic Programming Project

Recursive Implementation

Below is the code I used for my recursive implementation of the algorithm. When run with the parameters $p=3,\ t=16$ the function is called 753,665 times.

```
26
         # recursive algorithm
27
        def recursive(self, p: int, t: int) -> int:
28
            self.recursive_counter += 1 # increment counter for report
29
            # base cases
            if (t == 0 or t == 1):
30
31
                return t
             if (p == 1):
32
33
                return t
34
35
            # general cases for all values of x from 1 to t, inclusive
            results = [] # list to hold calculated costs, used to choose minimum
36
37
            for x in range(1, t + 1):
38
                breaksCase = self.recursive(p=p - 1, t=x - 1)
39
                 intactCase = self.recursive(p=p, t=t - x)
40
                maxThrows = max(breaksCase, intactCase) # maximum of the two cases is chosen
41
                results.append(maxThrows)
42
            results.sort() # sorts in place in ascending order
            return 1 + results[0] # gets the minimum value from all values of x + 1
```

Recursive Runtime Analysis (measured in ms)

$\frac{t}{p}$	10	12	15	18	20
2	1	2	20	170	716
4	7	39	499	5,942	29,772
8	13	117	2,958	65,222	
12	13	134	3,398		

As can be seen in the table above, the runtime increases greatly with slight changes in t. Additionally, higher values of p begin to affect the runtime more as t increases. For example, the runtime of p=8, t=10 is fairly close to that of p=4, t=10, but p=8, t=15 is far greater than p=4, t=15.

Dynamic Programming Implementation

Below is the code I used for my dynamic programming implementation of the algorithm.

```
45
         # dynamic programming algorithm
46
        def dynamic(self, p: int, t: int):
47
            self.dynamic_counter += 1 # implement counter for report
48
            # base cases
49
            if (t == 0 or t == 1):
50
                self.table[p-1][t-1] = t
51
                return
52
            if (p == 1):
53
                self.table[p-1][t-1] = t
54
                return
55
            # general case for all values of x from 1 to t
56
57
            results = []
58
            for x in range(1, t + 1):
59
                # if the value is none, it is calculated by calling the dynamic function on that spot
60
                if self.table[p - 2][x - 2] == None:
61
                    self.dynamic(p=p-1, t=x-1)
                # after that, the value is read for comparison
62
63
                breaks\_case = self.table[p - 2][x - 2]
64
                # same as above but for the case where the pumpkin stays intact
65
                if self.table[p - 1][t - \times - 1] == None:
66
67
                    self.dynamic(p=p, t=t - x)
                intact_case = self.table[p - 1][t - x - 1]
68
69
                max_throws = max(breaks_case, intact_case)
70
71
                results.append(max_throws)
72
            results.sort()
73
74
            # assigns minimum value to the table slot corresponding to p and t
75
            self.table[p-1][t-1] = 1 + results[0]
```

Dynamic Runtime Analysis (measured in ms)

$\frac{t}{p}$	20	40	60	80	100	120	140	160	180	200
20	1	6	14	26	41	64	86	114	141	176
40	1	7	20	39	70	107	148	202	256	328
80	1	7	20	47	89	143	217	308	411	538
120	1	7	21	47	91	153	239	353	481	634
160	1	7	21	48	91	153	239	359	517	695

As can be seen in the above table, growth in t leads to an increased runtime with p held constant. Similarly, an increase p can lead to an increased runtime, however this value seems to max out as $p \to t$. If p and t are increasing simultaneously, then the runtime will increase dramatically.