# Competitive Programming Notes

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## 1 Algorithms

Algorithm	Time	Space
Articulations and Bridges	$\mathcal{O}(V+E)$	$\mathcal{O}(V+E)$
Bellman-Ford	$\mathcal{O}(VE)$	$\mathcal{O}(V+E)$
Dijksta	$\mathcal{O}((V+E)\log V)$	$\mathcal{O}(V^2)$
Edmond Karp	$\mathcal{O}(VE^2)$	$\mathcal{O}(V+E)$
Euler Tour	$\mathcal{O}(E^2)$	
Floyd Warshall	$\mathcal{O}(V^3 + E)$	$\mathcal{O}(V^2+E)$
Graph Check	$\mathcal{O}(V+E)$	$\mathcal{O}(V+E)$
Kahn	$\mathcal{O}(VE)$	$\mathcal{O}(V+E)$
Kruskal	$\mathcal{O}(E \log V)$	$\mathcal{O}(V+E)$
LCA	$\mathcal{O}(N \log N)$	$\mathcal{O}(N \log N)$
MCBM	$\mathcal{O}(VE)$	
Prim	$\mathcal{O}(E \log V)$	$\mathcal{O}(V+E)$
Tarjan	$\mathcal{O}(V+E)$	$\mathcal{O}(V+E)$
Extended Euclid	$\mathcal{O}(\log \min(a, b))$	$\mathcal{O}(1)$
Floyd (cycle)	$\mathcal{O}(V)$	$\mathcal{O}(1)$
Prime Fac. w/ Opt. Trial Div.	$\mathcal{O}(\pi(\sqrt{n}))$	$\mathcal{O}(n)$
Sieve of Eratosthenes	$\mathcal{O}(n \log \log n)$	$\mathcal{O}(n)$
Binary Search	$\mathcal{O}(\log N)$	
Coordinate Compression	$\mathcal{O}(N \log N)$	
KMP	$\mathcal{O}(N)$	
MUF	$\mathcal{O}(AM)$	$\mathcal{O}(N)$
Bottom-Up SegTree	$\mathcal{O}(\log N)$	$\mathcal{O}(N)$

A: Ackermann function

## 1.1 Graph

#### 1.1.1 Articulations and Bridges

If vertex v is an **articulation point** and you remove it, the connected component to which it belongs becomes disconnected

If edge u, v is a  $\mathbf{bridge}$  and you remove it, you can't reach v from u

### 1.1.2 Edmond Karp MaxFlow

Ford-Fulkerson's method with BFS  $\rightarrow \mathcal{O}(VE)$  BFS calls,  $\mathcal{O}(E)$  per BFS

**Vertex weights:** if vertex V has a weight, create  $V_{in}$  (receives all in-edges of V and has an edge to  $V_{out}$ ) and  $V_{out}$  (receives an edge from  $V_{in}$  and has all out-edges of V); edge  $\{V_{in}, V_{out}\}$  has the weight from V

**MinCut:** run EdmondKarp; S-T sets are: all V that you can reach from the source with edges of positive residual capacity and all other V

MultiSource/MultiSink: create a super source with infinite capacity pointing to all sources, analogous for sinks

Max Cardinality Bipartite Matching: use capacity 1 on all edges and apply the multi-source and multi-sink strategies

#### 1.1.3 Euler Tour

Find the closest neighbor that has a path back to the current vertex to build an euler tour

Euler path: visits each edge once

 $\mathbf{Tour}/\mathbf{cycle}/\mathbf{circuit}$  euler path that starts and ends at same node

Undirected and has path: every vertex has even degree or two have odd degree

Undirected and has circuit: every vertex has even degree

**Directed and has path:**  $\delta^+(v) - \delta^-(v) = 1$  for at most one v, = -1 for at most one v, = 0 for all other v

Directed and has circuit:  $\delta^+(v) = \delta^-(v) \forall v \in V$ 

### 1.1.4 Floyd Warshall

Also works for SSSP ( $V \le 400$ )

**Printing path:** p[i][j] set to i (last node that appears before j on the path), then p[i][j] = p[k][j] on update.

**Transitive Closure:** weight is boolean (init as 1 if there's an edge), update with bitwise OR

Minimax/Maximin: w[i][j] will be min(w[i][j],
max(w[i][k], w[k][j]))

Finding negative/cheapest cycle: init w[i][i] =
inf; run(); any w[i][i] != inf is a cycle and the smallest
is the cheapest; any w[i][i] < 0 is negative cycle</pre>

This can also be used for finding SCCs (check with transitive closure)  $\,$ 

#### 1.1.5 Kahn's Topological Sort

Particular order (alphabetical)

## 1.1.6 Kruskal

min)

Order edges by increasing weight, then use a MUF to know if each edge is useful (if it connects two previously disconnected vertices)

Min Span Subgraph: previously process fixed edges Min Span Forest: count number of sets on the MUF 2nd Best MST: run kruskal; for each chosen edge, flag it as unavailable and run it without using that edge (O(VE)) Minimax: max edge weight on the MST (maximin:

## 1.1.7 Lowest Common Ancestor

Binary lift to binary search the LCA or Euler Path

## 1.1.8 Max Cardinality Bipartite Matching

Jump from free to matched edges until you've used them all

#### 1.1.9 Prim's Algorithm

Take smallest edge that leads to vertex v

#### 1.1.10 Tarjan

A node can reach any other node in its own SCC (DFS + stack)

#### 1.2 Math

#### 1.2.1 Floyd

Slow and fast (tortoise and hare)

## 1.3 Paradigm

## 1.3.1 Coordinate Compression

Normalize vector access; can also be done with map/set but high constant

#### 1.3.2 128 Bit Integers

GCC extension;  $2^{127}$   $10^{38}$ 

### 1.4 String

#### 1.4.1 Prefix Function (KMP)

To find ocurrences of s in t, use the string s+%+t, then look for pi[i] = s.length() on the "t side"