

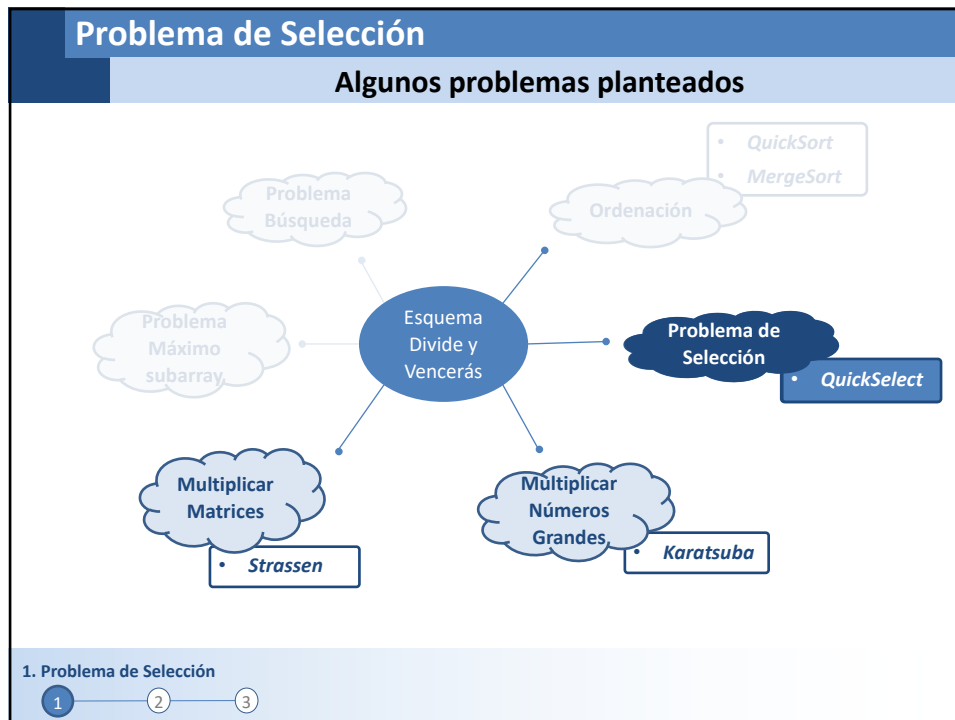


Universidad Politécnica de Madrid
Escuela Técnica Superior de Ingeniería de Sistemas Informáticos

Tema 7. Otros algoritmos basados en Divide y Vencerás

Algorítmica y Complejidad

1



2

Problema de Selección

Enunciado del Problema

Dado un *array*: un número k : 1

| | | | | | | | | | |
|----|-----|---|---|----|---|----|---|----|---|
| -2 | -21 | 3 | 4 | 12 | 8 | 10 | 5 | -6 | 1 |
|----|-----|---|---|----|---|----|---|----|---|

Devolver el k -ésimo número mayor del *array*

12

1. Problema de Selección

1
2
3

3

Problema de Selección

Enunciado del Problema

Dado un *array*: un número k : 3

| | | | | | | | | | |
|----|-----|---|---|----|---|----|---|----|---|
| -2 | -21 | 3 | 4 | 12 | 8 | 10 | 5 | -6 | 1 |
|----|-----|---|---|----|---|----|---|----|---|

Devolver el k -ésimo número mayor del *array*

8

1. Problema de Selección

1
2
3

4

Problema de Selección

Enunciado del Problema

Dado un *array*: un número k : 5

| | | | | | | | | | |
|----|-----|---|---|----|---|----|---|----|---|
| -2 | -21 | 3 | 4 | 12 | 8 | 10 | 5 | -6 | 1 |
|----|-----|---|---|----|---|----|---|----|---|

Devolver el k -ésimo número mayor del *array*

4

1. Problema de Selección

1
2
3

5

Problema de Selección

?

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

Dado un *array*: un número k : 5

| | | | | | | | | | |
|----|-----|---|---|----|---|----|---|----|---|
| -2 | -21 | 3 | 4 | 12 | 8 | 10 | 5 | -6 | 1 |
|----|-----|---|---|----|---|----|---|----|---|

Devolver el k -ésimo número mayor del *array*

4

1. Problema de Selección

1
2
3

6

Problema de Selección

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

max: 10
j: 2
j_max: 8

i: 6

k: 5

| | | | | | | | | | |
|----|-----|---|---|----|---|----|---|----|---|
| -2 | -21 | 3 | 4 | 12 | 8 | 10 | 5 | -6 | 1 |
|----|-----|---|---|----|---|----|---|----|---|

```

int k_esimoValor(int[] vector, int k){
    int max= Integer.MAX_VALUE;
    for (int j=0; j<k; j++) {
        int j_max = Integer.MIN_VALUE;
        for (int i=0; i<vector.length; i++)
            if (vector[i]< max && vector[i]> j_max)
                j_max = vector[i];
        max = j_max;
    }
    return max;
}

```

ERROR!!!
No funciona si hay elementos repetidos

1. Problema de Selección

1 2 3

7

Problema de Selección

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

k: 5

j: 0

| | | | | | | | | | |
|----|-----|---|---|----|---|----|---|----|---|
| -2 | -21 | 3 | 4 | 12 | 8 | 10 | 5 | -6 | 1 |
|----|-----|---|---|----|---|----|---|----|---|

```

int k_esimoValor(int[] vector, int k){
    int max = Integer.MIN_VALUE;
    ➔ for (int j=0; j<k; j++) {
        max = vector[j];
        int j_max = j;
        for (int i=j; i<vector.length; i++)
            if (vector[i] > max) {
                j_max = i; max = vector[i];
            }
        int aux = vector[j]; vector[j] = vector[j_max]; vector[j_max]= aux;
    }
    return max;
}

```

1. Problema de Selección

1 2 3

8

Problema de Selección

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

k: 5

j:0

| | | | | | | | | | |
|----|-----|---|---|----|---|----|---|----|---|
| -2 | -21 | 3 | 4 | 12 | 8 | 10 | 5 | -6 | 1 |
|----|-----|---|---|----|---|----|---|----|---|

```

int k_esimoValor(int[] vector, int k){
    int max = Integer.MIN_VALUE;
    for (int j=0; j<k; j++) {
        max = vector[j];
        int j_max = j;
        ➔ for (int i=j; i<vector.length; i++)
            if (vector[i] > max) {
                j_max = i; max = vector[i];
            }
        int aux = vector[j]; vector[j] = vector[j_max]; vector[j_max]= aux;
    }
    return max;
}

```

Encuentra el elemento más grande del subvector

1. Problema de Selección

1 2 3

9

Problema de Selección

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

k: 5

j:0 j_max:4

| | | | | | | | | | |
|----|-----|---|---|----|---|----|---|----|---|
| -2 | -21 | 3 | 4 | 12 | 8 | 10 | 5 | -6 | 1 |
|----|-----|---|---|----|---|----|---|----|---|

```

int k_esimoValor(int[] vector, int k){
    int max = Integer.MIN_VALUE;
    for (int j=0; j<k; j++) {
        max = vector[j];
        int j_max = j;
        for (int i=j; i<vector.length; i++)
            if (vector[i] > max) {
                j_max = i; max = vector[i];
            }
        ➔ int aux = vector[j]; vector[j] = vector[j_max]; vector[j_max]= aux;
    }
    return max;
}

```

Intercambiamos los elementos del vector

1. Problema de Selección

1 2 3

10

Problema de Selección

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

k: 5

j:0 j_max:4

| | | | | | | | | | |
|----|-----|---|---|----|---|----|---|----|---|
| 12 | -21 | 3 | 4 | -2 | 8 | 10 | 5 | -6 | 1 |
|----|-----|---|---|----|---|----|---|----|---|

```

int k_esimoValor(int[] vector, int k){
    int max = Integer.MIN_VALUE;
    ➔ for (int j=0; j<k; j++) {
        max = vector[j];
        int j_max = j;
        for (int i=j; i<vector.length; i++)
            if (vector[i] > max) {
                j_max = i; max = vector[i];
            }
        int aux = vector[j]; vector[j] = vector[j_max]; vector[j_max]= aux;
    }
    return max;
}

```

1. Problema de Selección

1 2 3

11

Problema de Selección

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

k: 5

j:1

| | | | | | | | | | |
|----|-----|---|---|----|---|----|---|----|---|
| 12 | -21 | 3 | 4 | -2 | 8 | 10 | 5 | -6 | 1 |
|----|-----|---|---|----|---|----|---|----|---|

```

int k_esimoValor(int[] vector, int k){
    int max = Integer.MIN_VALUE;
    for (int j=0; j<k; j++) {
        max = vector[j];
        int j_max = j;
        ➔ for (int i=j; i<vector.length; i++)
            if (vector[i] > max) {
                j_max = i; max = vector[i];
            }
        int aux = vector[j]; vector[j] = vector[j_max]; vector[j_max]= aux;
    }
    return max;
}

```

Encuentra el elemento más grande del subvector

1. Problema de Selección

1 2 3

12

Problema de Selección

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

k: 5

| | | | | | | | | | |
|----|-----|---|---|----|---|----|---|----|---|
| 12 | -21 | 3 | 4 | -2 | 8 | 10 | 5 | -6 | 1 |
|----|-----|---|---|----|---|----|---|----|---|

j:1 j_max:6

```

int k_esimoValor(int[] vector, int k){
    int max = Integer.MIN_VALUE;
    for (int j=0; j<k; j++) {
        max = vector[j];
        int j_max = j;
        for (int i=j; i<vector.length; i++)
            if (vector[i] > max) {
                j_max = i; max = vector[i];
            }
        ➔ int aux = vector[j]; vector[j] = vector[j_max]; vector[j_max]= aux;
    }
    return max;
}

```

Intercambiamos los elementos del vector

1. Problema de Selección

1 2 3

13

Problema de Selección

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

k: 5

| | | | | | | | | | |
|----|----|---|---|----|---|-----|---|----|---|
| 12 | 10 | 3 | 4 | -2 | 8 | -21 | 5 | -6 | 1 |
|----|----|---|---|----|---|-----|---|----|---|

j:1 j_max:6

```

int k_esimoValor(int[] vector, int k){
    int max = Integer.MIN_VALUE;
    ➔ for (int j=0; j<k; j++) {
        max = vector[j];
        int j_max = j;
        for (int i=j; i<vector.length; i++)
            if (vector[i] > max) {
                j_max = i; max = vector[i];
            }
        int aux = vector[j]; vector[j] = vector[j_max]; vector[j_max]= aux;
    }
    return max;
}

```

1. Problema de Selección

1 2 3

14

Problema de Selección

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

k: 5

| | | | | | | | | | |
|----|----|---|---|----|---|-----|---|----|---|
| 12 | 10 | 3 | 4 | -2 | 8 | -21 | 5 | -6 | 1 |
|----|----|---|---|----|---|-----|---|----|---|

j:2

```

int k_esimoValor(int[] vector, int k){
    int max = Integer.MIN_VALUE;
    for (int j=0; j<k; j++) {
        max = vector[j];
        int j_max = j;
        ➔ for (int i=j; i<vector.length; i++)
            if (vector[i] > max) {
                j_max = i; max = vector[i];
            }
        int aux = vector[j]; vector[j] = vector[j_max]; vector[j_max]= aux;
    }
    return max;
}

```

Encuentra el elemento más grande del subvector

1. Problema de Selección

1 2 3

15

Problema de Selección

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

k: 5

| | | | | | | | | | |
|----|----|---|---|----|---|-----|---|----|---|
| 12 | 10 | 3 | 4 | -2 | 8 | -21 | 5 | -6 | 1 |
|----|----|---|---|----|---|-----|---|----|---|

j:2 j_max:5

```

int k_esimoValor(int[] vector, int k){
    int max = Integer.MIN_VALUE;
    for (int j=0; j<k; j++) {
        max = vector[j];
        int j_max = j;
        for (int i=j; i<vector.length; i++)
            if (vector[i] > max) {
                j_max = i; max = vector[i];
            }
        ➔ int aux = vector[j]; vector[j] = vector[j_max]; vector[j_max]= aux;
    }
    return max;
}

```

Intercambiamos los elementos del vector

1. Problema de Selección

1 2 3

16

Problema de Selección

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

k: 5

| | | | | | | | | | |
|----|----|---|---|----|---|-----|---|----|---|
| 12 | 10 | 8 | 4 | -2 | 3 | -21 | 5 | -6 | 1 |
|----|----|---|---|----|---|-----|---|----|---|

j:2 j_max:5

```

int k_esimoValor(int[] vector, int k){
    int max = Integer.MIN_VALUE;
    ➔ for (int j=0; j<k; j++) {
        max = vector[j];
        int j_max = j;
        for (int i=j; i<vector.length; i++)
            if (vector[i] > max) {
                j_max = i; max = vector[i];
            }
        int aux = vector[j]; vector[j] = vector[j_max]; vector[j_max]= aux;
    }
    return max;
}
  
```

1. Problema de Selección

1 — 2 — 3

17

Problema de Selección

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

k: 5

| | | | | | | | | | |
|----|----|---|---|----|---|-----|---|----|---|
| 12 | 10 | 8 | 4 | -2 | 3 | -21 | 5 | -6 | 1 |
|----|----|---|---|----|---|-----|---|----|---|

j:3

```

int k_esimoValor(int[] vector, int k){
    int max = Integer.MIN_VALUE;
    for (int j=0; j<k; j++) {
        max = vector[j];
        int j_max = j;
        ➔ for (int i=j; i<vector.length; i++)
            if (vector[i] > max) {
                j_max = i; max = vector[i];
            }
        int aux = vector[j]; vector[j] = vector[j_max]; vector[j_max]= aux;
    }
    return max;
}
  
```

Encuentra el elemento más grande del subvector

1. Problema de Selección

1 — 2 — 3

18

Problema de Selección

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

k: 5

| | | | | | | | | | |
|----|----|---|---|----|---|-----|---|----|---|
| 12 | 10 | 8 | 4 | -2 | 3 | -21 | 5 | -6 | 1 |
|----|----|---|---|----|---|-----|---|----|---|

j:3 j_max:7

```

int k_esimoValor(int[] vector, int k){
    int max = Integer.MIN_VALUE;
    for (int j=0; j<k; j++) {
        max = vector[j];
        int j_max = j;
        for (int i=j; i<vector.length; i++)
            if (vector[i] > max) {
                j_max = i; max = vector[i];
            }
        ➔ int aux = vector[j]; vector[j] = vector[j_max]; vector[j_max]= aux;
    }
    return max;
}

```

Intercambiamos los elementos del vector

1. Problema de Selección

1 — 2 — 3

19

Problema de Selección

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

k: 5

| | | | | | | | | | |
|----|----|---|---|----|---|-----|---|----|---|
| 12 | 10 | 8 | 5 | -2 | 3 | -21 | 4 | -6 | 1 |
|----|----|---|---|----|---|-----|---|----|---|

j:3 j_max:7

```

int k_esimoValor(int[] vector, int k){
    int max = Integer.MIN_VALUE;
    ➔ for (int j=0; j<k; j++) {
        max = vector[j];
        int j_max = j;
        for (int i=j; i<vector.length; i++)
            if (vector[i] > max) {
                j_max = i; max = vector[i];
            }
        int aux = vector[j]; vector[j] = vector[j_max]; vector[j_max]= aux;
    }
    return max;
}

```

1. Problema de Selección

1 — 2 — 3

20

Problema de Selección

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

k: 5

| | | | | | | | | | |
|----|----|---|---|----|---|-----|---|----|---|
| 12 | 10 | 8 | 5 | -2 | 3 | -21 | 4 | -6 | 1 |
|----|----|---|---|----|---|-----|---|----|---|

j:4

```

int k_esimoValor(int[] vector, int k){
    int max = Integer.MIN_VALUE;
    for (int j=0; j<k; j++) {
        max = vector[j];
        int j_max = j;
        ➔ for (int i=j; i<vector.length; i++)
            if (vector[i] > max) {
                j_max = i; max = vector[i];
            }
        int aux = vector[j]; vector[j] = vector[j_max]; vector[j_max]= aux;
    }
    return max;
}

```

Encuentra el elemento más grande del subvector

1. Problema de Selección

1 2 3

21

Problema de Selección

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

k: 5

| | | | | | | | | | |
|----|----|---|---|----|---|-----|---|----|---|
| 12 | 10 | 8 | 5 | -2 | 3 | -21 | 4 | -6 | 1 |
|----|----|---|---|----|---|-----|---|----|---|

j:4 j_max:7

```

int k_esimoValor(int[] vector, int k){
    int max = Integer.MIN_VALUE;
    for (int j=0; j<k; j++) {
        max = vector[j];
        int j_max = j;
        for (int i=j; i<vector.length; i++)
            if (vector[i] > max) {
                j_max = i; max = vector[i];
            }
        ➔ int aux = vector[j]; vector[j] = vector[j_max]; vector[j_max]= aux;
    }
    return max;
}

```

Intercambiamos los elementos del vector

1. Problema de Selección

1 2 3

22

Problema de Selección

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

k: 5

| | | | | | | | | | |
|----|----|---|---|---|---|-----|----|----|---|
| 12 | 10 | 8 | 5 | 4 | 3 | -21 | -2 | -6 | 1 |
|----|----|---|---|---|---|-----|----|----|---|

j:4 j_max:7

```

int k_esimoValor(int[] vector, int k){
    int max = Integer.MIN_VALUE;
    ➔ for (int j=0; j<k; j++) {
        max = vector[j];
        int j_max = j;
        for (int i=j; i<vector.length; i++)
            if (vector[i] > max) {
                j_max = i; max = vector[i];
            }
        int aux = vector[j]; vector[j] = vector[j_max]; vector[j_max]= aux;
    }
    return max;
}

```

1. Problema de Selección

1 — 2 — 3

23

Problema de Selección

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

k: 5

| | | | | | | | | | |
|----|----|---|---|---|---|-----|----|----|---|
| 12 | 10 | 8 | 5 | 4 | 3 | -21 | -2 | -6 | 1 |
|----|----|---|---|---|---|-----|----|----|---|

j:5

```

int k_esimoValor(int[] vector, int k){
    int max = Integer.MIN_VALUE;
    for (int j=0; j<k; j++) {
        max = vector[j];
        int j_max = j;
        for (int i=j; i<vector.length; i++)
            if (vector[i] > max) {
                j_max = i; max = vector[i];
            }
        int aux = vector[j]; vector[j] = vector[j_max]; vector[j_max]= aux;
    }
    ➔ return max;
}

```

1. Problema de Selección

1 — 2 — 3

24

Problema de Selección

Actividad 7.1. Diseña e implementa un algoritmo que resuelva el problema de selección.

k: 5

j:5

| | | | | | | | | | |
|----|----|---|---|---|---|-----|----|----|---|
| 12 | 10 | 8 | 5 | 4 | 3 | -21 | -2 | -6 | 1 |
|----|----|---|---|---|---|-----|----|----|---|

```

int k_esimoValor(int[] vector, int k){
    int max = Integer.MIN_VALUE;
    for (int j=0; j<k; j++) {
        max = vector[j];
        int j_max = j;
        for (int i=j; i<vector.length; i++)
            if (vector[i] > max) {
                j_max = i; max = vector[i];
            }
        int aux = vector[j]; vector[j] = vector[j_max]; vector[j_max]= aux;
    }
    return max;
}

```

Parecido a ordenación por selección.

Solo ordenamos los k primeros elementos del array

1. Problema de Selección

1 — 2 — 3

25

Problema de Selección

Algoritmo Simple

Actividad 7.2. Calcula la complejidad del algoritmo.

j:5

| | | | | | | | | | |
|----|----|---|---|---|---|-----|----|----|---|
| 12 | 10 | 8 | 5 | 4 | 3 | -21 | -2 | -6 | 1 |
|----|----|---|---|---|---|-----|----|----|---|

```

int k_esimoValor(int[] vector, int k){
    int max = Integer.MIN_VALUE;
    for (int j=0; j<k; j++) {
        max = vector[j];
        int j_max = j;
        for (int i=j; i<vector.length; i++)
            if (vector[i] > max) {
                j_max = i; max = vector[i];
            }
        int aux = vector[j]; vector[j] = vector[j_max]; vector[j_max]= aux;
    }
    return max;
}

```

Complejidad: $\Theta(N \cdot k)$

1. Problema de Selección

1 — 2 — 3

26

Problema de Selección

Algoritmo Simple

Actividad 7.3. Diseña e implementa un algoritmo que tenga complejidad $O(N \cdot \log N)$. ?

$j:5$

| | | | | | | | | | |
|----|----|---|---|---|---|-----|----|----|---|
| 12 | 10 | 8 | 5 | 4 | 3 | -21 | -2 | -6 | 1 |
|----|----|---|---|---|---|-----|----|----|---|

```
int k_esimoValor (int[] vector, int k){
    ordenarHeapSortMayorMenor(vector);
    return vector[k-1];
}
```

Complejidad:
 $\Theta(N \cdot \log N)$

1. Problema de Selección

1

2

3

27

Problema de Selección

QuickSelect: Algoritmo basado en Divide y Vencerás

| | | | | | | | | | |
|----|----|---|---|-----|---|----|---|----|---|
| -2 | 12 | 3 | 4 | -21 | 8 | 10 | 5 | -6 | 1 |
|----|----|---|---|-----|---|----|---|----|---|

elementos ≤ 1
↓ Fase Dividir
elementos > 1

$k: 8$

| | | | | | | | | | |
|----|----|-----|---|---|---|----|---|----|---|
| -2 | -6 | -21 | 1 | 3 | 8 | 10 | 5 | 12 | 4 |
|----|----|-----|---|---|---|----|---|----|---|

3 elementos
 $k: 8-6-1=1$
6 elementos

Fase Conquistar ↓

| | | |
|-----|----|----|
| -21 | -6 | -2 |
|-----|----|----|

Fase Combinar: ↓

| |
|----|
| -2 |
|----|

Caso 1 El k-ésimo valor está en el 1^{er} subarray

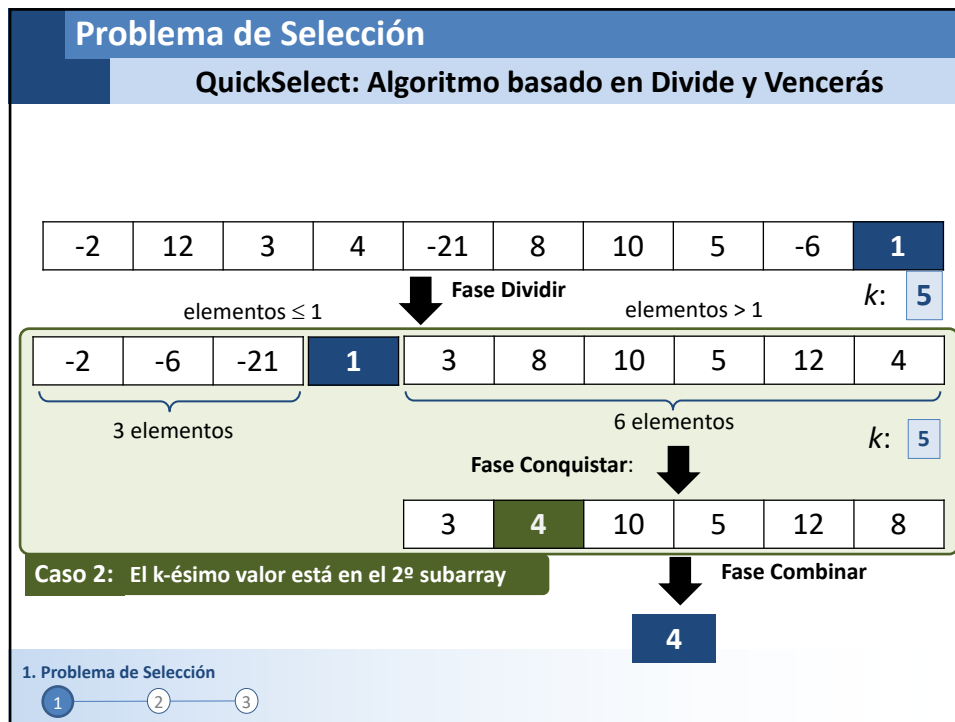
1. Problema de Selección

1

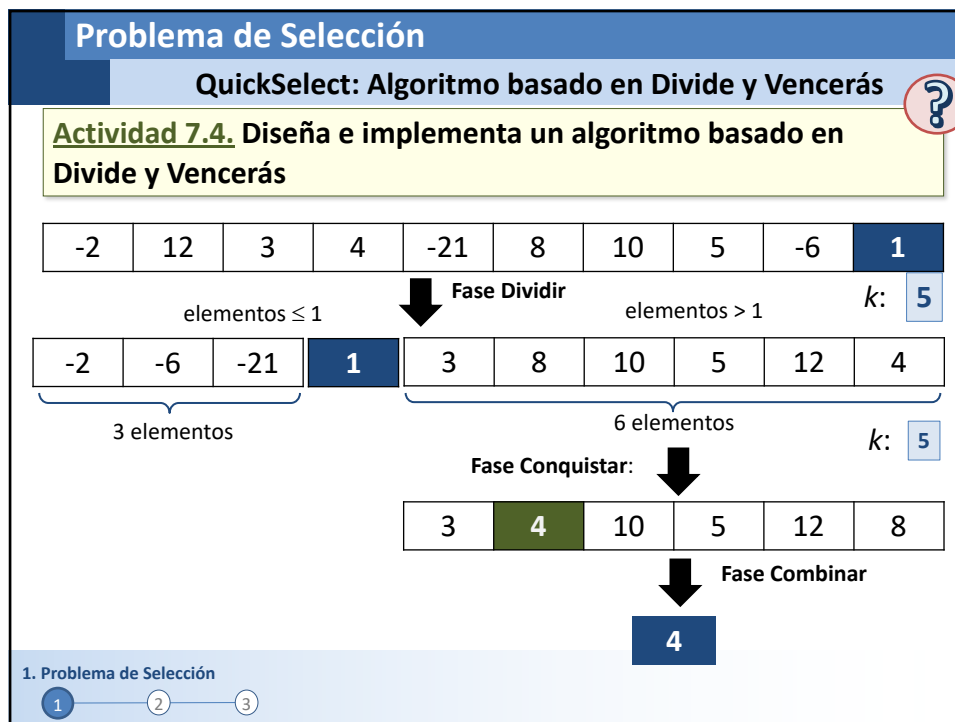
2

3

28



29



30

Problema de Selección

Algoritmo basado en Divide y Vencerás: Implementación

| | | | | | | | | | |
|-------|-----|-------|-------|-------|----------|----------|----------|-----|----------|
| x_0 | ... | x_7 | x_8 | x_9 | x_{10} | x_{11} | x_{12} | ... | x_{50} |
|-------|-----|-------|-------|-------|----------|----------|----------|-----|----------|

```

int k_esimoValorAux (int[] vector, int i0, int iN, int k){
    if (i0 == iN)
        return vector[i0];
    else {
        int m = ordenarPorPivote(vector, i0, iN);
        if (k <= iN-m)
            return k_esimoValorAux(vector, m+1, iN, k);
        else if (k == iN-m+1)
            return vector[m];
        else
            return k_esimoValorAux(vector, i0, m-1, k-(iN-m+1));
    }
}

int k_esimoValor(int[] vector, int k){
    return k_esimoValorAux(vector, 0, vector.length-1, k);
}
  
```

Utilizado en QuickSort
Complejidad: $\Theta(N)$

1. Problema de Selección

1 — 2 — 3

31

Problema de Selección

Algoritmo basado en Divide y Vencerás: Implementación

$i0: 7$

$iN: 12$

| | | | | | | | | | |
|-------|-----|-------|-------|-------|----------|----------|----------|-----|----------|
| x_0 | ... | x_7 | x_8 | x_9 | x_{10} | x_{11} | x_{12} | ... | x_{50} |
|-------|-----|-------|-------|-------|----------|----------|----------|-----|----------|

```

int k_esimoValorAux (int[] vector, int i0, int iN, int k){
    if (i0 == iN)
        return vector[i0];
    else {
        int m = ordenarPorPivote(vector, i0, iN);
        if (k <= iN-m)
            return k_esimoValorAux(vector, m+1, iN, k);
        else if (k == iN-m+1)
            return vector[m];
        else
            return k_esimoValorAux(vector, i0, m-1, k-(iN-m+1));
    }
}

int k_esimoValor(int[] vector, int k){
    return k_esimoValorAux(vector, 0, vector.length-1, k);
}
  
```

1. Problema de Selección

1 — 2 — 3

32

Problema de Selección

Algoritmo basado en Divide y Vencerás: Implementación

$i0:7$ $iN:7$

↓

| | | | | | | | | | |
|-------|-----|-------|-------|-------|----------|----------|----------|-----|----------|
| x_0 | ... | x_7 | x_8 | x_9 | x_{10} | x_{11} | x_{12} | ... | x_{50} |
|-------|-----|-------|-------|-------|----------|----------|----------|-----|----------|

```

int k_esimoValorAux (int[] vector, int i0, int iN, int k){
    if (i0 == iN)
        return vector[i0];
    else {
        int m = ordenarPorPivote(vector, i0, iN);
        if (k<=iN-m)
            return k_esimoValorAux(vector, m+1, iN, k);
        else if (k==iN-m+1)
            return vector[m];
        else
            return k_esimoValorAux(vector, i0, m-1, k-(iN-m+1));
    }
}

int k_esimoValor(int[] vector, int k){
    return k_esimoValorAux(vector,0, vector.length-1, k);
}

```

Caso Base

1. Problema de Selección

① — ② — ③

33

Problema de Selección

Algoritmo basado en Divide y Vencerás: Implementación

$i0: 7$ $iN: 12$

↓ ↓

| | | | | | | | | | |
|-------|-----|-------|-------|-------|----------|----------|----------|-----|----------|
| x_0 | ... | x_7 | x_8 | x_9 | x_{10} | x_{11} | x_{12} | ... | x_{50} |
|-------|-----|-------|-------|-------|----------|----------|----------|-----|----------|

```

int k_esimoValorAux (int[] vector, int i0, int iN, int k){
    if (i0 == iN)
        return vector[i0];
    else {
        int m = ordenarPorPivote(vector, i0, iN);
        if (k<=iN-m)
            return k_esimoValorAux(vector, m+1, iN, k);
        else if (k==iN-m+1)
            return vector[m];
        else
            return k_esimoValorAux(vector, i0, m-1, k-(iN-m+1));
    }
}

int k_esimoValor(int[] vector, int k){
    return k_esimoValorAux(vector,0, vector.length-1, k);
}

```

1. Problema de Selección

① — ② — ③

34

Problema de Selección

Algoritmo basado en Divide y Vencerás: Implementación

$i0: 7$ $m: 10$ $iN: 12$

↓ ↓ ↓

| | | | | | | | | | |
|-------|-----|-------|-------|-------|----------|----------|----------|-----|----------|
| x_0 | ... | x_7 | x_8 | x_9 | x_{10} | x_{11} | x_{12} | ... | x_{50} |
|-------|-----|-------|-------|-------|----------|----------|----------|-----|----------|

```

int k_esimoValorAux (int[] vector, int i0, int iN, int k){
    if (i0 == iN)
        return vector[i0];
    else {
        int m = ordenarPorPivote(vector, i0, iN);
        if (k <= iN-m)
            return k_esimoValorAux(vector, m+1, iN, k);
        else if (k == iN-m+1)
            return vector[m];
        else
            return k_esimoValorAux(vector, i0, m-1, k-(iN-m+1));
    }
}

int k_esimoValor(int[] vector, int k){
    return k_esimoValorAux(vector, 0, vector.length-1, k);
}
  
```

} Dividir
Utilizado en QuickSort

1. Problema de Selección

1 — 2 — 3

35

Problema de Selección

Algoritmo basado en Divide y Vencerás: Implementación

$i0: 7$ $m: 10$ $iN: 12$

↓ ↓ ↓

| | | | | | | | | | |
|-------|-----|-------|-------|-------|----------|----------|----------|-----|----------|
| x_0 | ... | x_7 | x_8 | x_9 | x_{10} | x_{11} | x_{12} | ... | x_{50} |
|-------|-----|-------|-------|-------|----------|----------|----------|-----|----------|

```

int k_esimoValorAux (int[] vector, int i0, int iN, int k){
    if (i0 == iN)
        return vector[i0];
    else {
        int m = ordenarPorPivote(vector, i0, iN);
        if (k <= iN-m)
            return k_esimoValorAux(vector, m+1, iN, k);
        else if (k == iN-m+1)
            return vector[m];
        else
            return k_esimoValorAux(vector, i0, m-1, k-(iN-m+1));
    }
}

int k_esimoValor(int[] vector, int k){
    return k_esimoValorAux(vector, 0, vector.length-1, k);
}
  
```

} Conquistar

1. Problema de Selección

1 — 2 — 3

36

Problema de Selección

Algoritmo basado en Divide y Vencerás: Implementación

Actividad 7.5. Calcula la complejidad en el mejor caso. ?

```

int k_esimoValorAux (int[] vector, int i0, int iN, int k){
    if (i0 == iN)
        return vector[i0];
    else {
        int m = ordenarPorPivote(vector, i0, iN);
        if (k <= iN-m)
            return k_esimoValorAux(vector, m+1, iN, k);
        else if (k == iN-m+1)
            return vector[m];
        else
            return k_esimoValorAux(vector, i0, m-1, k-(iN-m+1));
    }
}

int k_esimoValor(int[] vector, int k){
    return k_esimoValorAux(vector, 0, vector.length-1, k);
}

```

Complejidad:
 $\Theta(N)$

1. Problema de Selección

1 — 2 — 3

37

Problema de Selección

Algoritmo basado en Divide y Vencerás: Implementación

Actividad 7.5. Calcula la complejidad en el peor caso. ?

```

int k_esimoValorAux (int[] vector, int i0, int iN, int k){
    if (i0 == iN)
        return vector[i0];
    else {
        int m = ordenarPorPivote(vector, i0, iN);
        if (k <= iN-m)
            return k_esimoValorAux(vector, m+1, iN, k);
        else if (k == iN-m+1)
            return vector[m];
        else
            return k_esimoValorAux(vector, i0, m-1, k-(iN-m+1));
    }
}

```

Complejidad:
 $\Theta(N^2)$

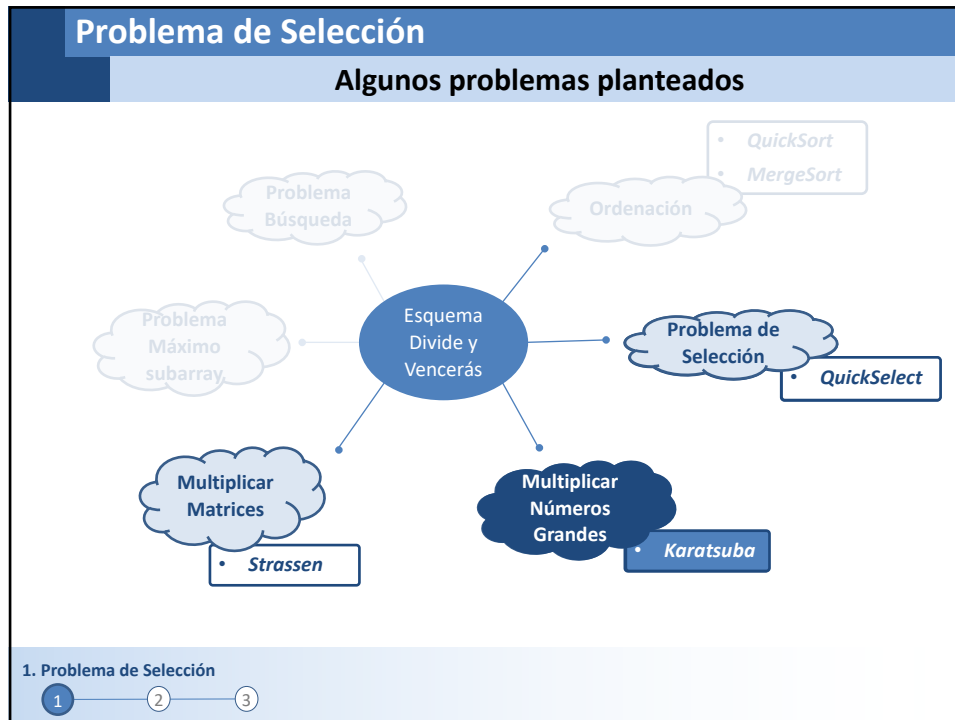
Ecuación de Recurrencia

$$T(N) = \begin{cases} \Theta(1) & N = 1 \\ T(N-1) + \Theta(N) & N > 1 \end{cases}$$

1. Problema de Selección

1 — 2 — 3

38



39

Multiplicar números grandes

Enunciado del Problema

Dados dos números:

$$a = \begin{bmatrix} 1 & 2 & 3 & 4 \end{bmatrix} \qquad b = \begin{bmatrix} 4 & 1 & 5 & 2 \end{bmatrix}$$

Multiplicarlos:

$$c = a \cdot b = \begin{bmatrix} 5 & 1 & 2 & 3 & 5 & 6 & 8 \end{bmatrix}$$

2. Multiplicar Números Grandes

1 — 2 — 3

40 de 98

40

Multiplicar números grandes

Algoritmo Tradicional

| | | | |
|---|---|---|---|
| 1 | 2 | 3 | 4 |
|---|---|---|---|

 = a
 x

| | | | |
|---|---|---|---|
| 4 | 1 | 5 | 2 |
|---|---|---|---|

 = b

| | | | |
|---|---|---|---|
| 2 | 4 | 6 | 8 |
|---|---|---|---|

| | | | |
|---|---|---|---|
| 6 | 1 | 7 | 0 |
|---|---|---|---|

| | | | |
|---|---|---|---|
| 1 | 2 | 3 | 4 |
|---|---|---|---|

 +

| | | | |
|---|---|---|---|
| 4 | 9 | 3 | 6 |
|---|---|---|---|

| | | | | | | |
|---|---|---|---|---|---|---|
| 5 | 1 | 2 | 3 | 5 | 6 | 8 |
|---|---|---|---|---|---|---|

 = c = a · b

2. Multiplicar Números Grandes

① — ② — ③

41 de 98

41

Multiplicar números grandes

Actividad 7.6. Calcula la complejidad del Algoritmo de acuerdo con el número de dígitos, N , de los multiplicandos.

Complejidad:
 $\Theta(N^2)$

| | | | |
|---|---|---|---|
| 1 | 2 | 3 | 4 |
|---|---|---|---|

 = a
 x

| | | | |
|---|---|---|---|
| 4 | 1 | 5 | 2 |
|---|---|---|---|

 = b

| | | | |
|---|---|---|---|
| 2 | 4 | 6 | 8 |
|---|---|---|---|

| | | | |
|---|---|---|---|
| 6 | 1 | 7 | 0 |
|---|---|---|---|

| | | | |
|---|---|---|---|
| 1 | 2 | 3 | 4 |
|---|---|---|---|

 +

| | | | |
|---|---|---|---|
| 4 | 9 | 3 | 6 |
|---|---|---|---|

| | | | | | | |
|---|---|---|---|---|---|---|
| 5 | 1 | 2 | 3 | 5 | 6 | 8 |
|---|---|---|---|---|---|---|

 = c = a · b

2. Multiplicar Números Grandes

① — ② — ③

42 de 98

42

Multiplicar números grandes

Algoritmo inmediato basado en Divide y Vencerás

N dígitos

$$a = \begin{array}{|c|c|c|c|} \hline a_1 & a_2 & & \\ \hline 1 & 2 & 3 & 4 \\ \hline \end{array}$$

$$a = a_1 \cdot 10^{N/2} + a_2$$

N dígitos

$$b = \begin{array}{|c|c|c|c|} \hline b_1 & b_2 & & \\ \hline 4 & 1 & 5 & 2 \\ \hline \end{array}$$

$$b = b_1 \cdot 10^{N/2} + b_2$$

$$\begin{array}{r}
 a_1 \cdot b_1 = \begin{array}{|c|c|c|} \hline 4 & 9 & 2 \\ \hline \end{array} \\
 a_1 \cdot b_2 = \begin{array}{|c|c|c|} \hline 6 & 2 & 4 \\ \hline \end{array} \\
 a_2 \cdot b_1 = \begin{array}{|c|c|c|c|} \hline 1 & 3 & 9 & 4 \\ \hline \end{array} \\
 + \quad a_2 \cdot b_2 = \begin{array}{|c|c|c|c|} \hline 1 & 7 & 6 & 8 \\ \hline \end{array} \\
 \hline
 \begin{array}{|c|c|c|c|c|c|} \hline 5 & 1 & 2 & 3 & 5 & 6 & 8 \\ \hline \end{array}
 \end{array}$$

$$c = a \cdot b = a_1 \cdot b_1 \cdot 10^N + (a_1 \cdot b_2 + a_2 \cdot b_1) \cdot 10^{N/2} + a_2 \cdot b_2$$

2. Multiplicar Números Grandes

① — ② — ③

43 de 98

43

Multiplicar números grandes

Enunciado del Problema

N dígitos

$$a = \begin{array}{|c|c|c|c|} \hline a_1 & a_2 & & \\ \hline 1 & 2 & 3 & 4 \\ \hline \end{array}$$

$$a = a_1 \cdot 10^{N/2} + a_2$$

N dígitos

$$b = \begin{array}{|c|c|c|c|} \hline b_1 & b_2 & & \\ \hline 4 & 1 & 5 & 2 \\ \hline \end{array}$$

$$b = b_1 \cdot 10^{N/2} + b_2$$

$$\begin{array}{r}
 a_1 \cdot b_1 = \begin{array}{|c|c|c|} \hline 4 & 9 & 2 \\ \hline \end{array} \\
 a_1 \cdot b_2 = \begin{array}{|c|c|c|} \hline 6 & 2 & 4 \\ \hline \end{array} \\
 a_2 \cdot b_1 = \begin{array}{|c|c|c|c|} \hline 1 & 3 & 9 & 4 \\ \hline \end{array} \\
 + \quad a_2 \cdot b_2 = \begin{array}{|c|c|c|c|} \hline 1 & 7 & 6 & 8 \\ \hline \end{array} \\
 \hline
 \begin{array}{|c|c|c|c|c|c|} \hline 5 & 1 & 2 & 3 & 5 & 6 & 8 \\ \hline \end{array}
 \end{array}$$

$$c = a \cdot b = a_1 \cdot b_1 \cdot 10^N + (a_1 \cdot b_2 + a_2 \cdot b_1) \cdot 10^{N/2} + a_2 \cdot b_2$$

Fase Dividir

2. Multiplicar Números Grandes

① — ② — ③

44 de 98

44

45

46

Actividad 7.7. Calcula la complejidad de este algoritmo basado en divide y vencerás para multiplicar números

N dígitos

$a = \begin{array}{|c|c|c|c|} \hline a_1 & a_2 & & \\ \hline 1 & 2 & 3 & 4 \\ \hline \end{array}$
 $a = a_1 \cdot 10^{N/2} + a_2$

N dígitos

$b = \begin{array}{|c|c|c|c|} \hline b_1 & b_2 & & \\ \hline 4 & 1 & 5 & 2 \\ \hline \end{array}$
 $b = b_1 \cdot 10^{N/2} + b_2$

$a_1 \cdot b_1 = \begin{array}{|c|c|c|} \hline 4 & 9 & 2 \\ \hline \end{array}$
 $a_1 \cdot b_2 = \begin{array}{|c|c|c|} \hline 6 & 2 & 4 \\ \hline \end{array}$
 $a_2 \cdot b_1 = \begin{array}{|c|c|c|c|} \hline 1 & 3 & 9 & 4 \\ \hline \end{array}$
 $a_2 \cdot b_2 = \begin{array}{|c|c|c|c|} \hline 1 & 7 & 6 & 8 \\ \hline \end{array}$

+

$\begin{array}{|c|c|c|c|c|c|} \hline 5 & 1 & 2 & 3 & 5 & 6 & 8 \\ \hline \end{array}$

$c = a \cdot b = a_1 \cdot b_1 \cdot 10^N + (a_1 \cdot b_2 + a_2 \cdot b_1) \cdot 10^{N/2} + a_2 \cdot b_2$

2. Multiplicar Números Grandes

1 2 3

47 de 98

47

Actividad 7.7. Calcula la complejidad de este algoritmo basado en divide y vencerás para multiplicar números

Fase Dividir

N dígitos

$a = \begin{array}{|c|c|c|c|} \hline a_1 & a_2 & & \\ \hline 1 & 2 & 3 & 4 \\ \hline \end{array}$
 $a = a_1 \cdot 10^{N/2} + a_2$

N dígitos

$b = \begin{array}{|c|c|c|c|} \hline b_1 & b_2 & & \\ \hline 4 & 1 & 5 & 2 \\ \hline \end{array}$
 $b = b_1 \cdot 10^{N/2} + b_2$

Fase Conquistar

$a_1 \cdot b_1 = \begin{array}{|c|c|c|} \hline 4 & 9 & 2 \\ \hline \end{array}$
 $a_1 \cdot b_2 = \begin{array}{|c|c|c|} \hline 6 & 2 & 4 \\ \hline \end{array}$
 $a_2 \cdot b_1 = \begin{array}{|c|c|c|c|} \hline 1 & 3 & 9 & 4 \\ \hline \end{array}$
 $a_2 \cdot b_2 = \begin{array}{|c|c|c|c|} \hline 1 & 7 & 6 & 8 \\ \hline \end{array}$

+

Fase Combinar

$\begin{array}{|c|c|c|c|c|c|} \hline 5 & 1 & 2 & 3 & 5 & 6 & 8 \\ \hline \end{array}$

$c = a \cdot b = a_1 \cdot b_1 \cdot 10^N + (a_1 \cdot b_2 + a_2 \cdot b_1) \cdot 10^{N/2} + a_2 \cdot b_2$

$$T(N) = \begin{cases} \Theta(1) & N = 1 \\ 4 \cdot T\left(\frac{N}{2}\right) + \Theta(N) & N > 1 \end{cases}$$

48 de 98

48

Actividad 7.7. Calcula la complejidad de este algoritmo basado en divide y vencerás para multiplicar números

Fase Dividir

N dígitos

$a = \begin{matrix} a_1 & a_2 \\ 1 & 2 & 3 & 4 \end{matrix}$
 $a = a_1 \cdot 10^{N/2} + a_2$

N dígitos

$b = \begin{matrix} b_1 & b_2 \\ 4 & 1 & 5 & 2 \end{matrix}$
 $b = b_1 \cdot 10^{N/2} + b_2$

Fase Conquistar

$a_1 \cdot b_1 = \begin{matrix} 4 & 9 & 2 \end{matrix}$
 $a_1 \cdot b_2 = \begin{matrix} 6 & 2 & 4 \end{matrix}$
 $a_2 \cdot b_1 = \begin{matrix} 1 & 3 & 9 & 4 \end{matrix}$
 $a_2 \cdot b_2 = \begin{matrix} 1 & 7 & 6 & 8 \end{matrix}$

Fase Combinar

$c = a \cdot b = a_1 \cdot b_1 \cdot 10^N + (a_1 \cdot b_2 + a_2 \cdot b_1) \cdot 10^{N/2} + a_2 \cdot b_2$

$T(N) = \begin{cases} \theta(1) & N = 1 \\ 4 \cdot T\left(\frac{N}{2}\right) + \theta(N^1) & N > 1 \end{cases}$

Teorema Maestro
 (1^{er} Caso)
 $\log_2(4) = 2 > 1$

49 de 98

49

Actividad 7.7. Calcula la complejidad de este algoritmo basado en divide y vencerás para multiplicar números

Fase Dividir

N dígitos

$a = \begin{matrix} a_1 & a_2 \\ 1 & 2 & 3 & 4 \end{matrix}$
 $a = a_1 \cdot 10^{N/2} + a_2$

N dígitos

$b = \begin{matrix} b_1 & b_2 \\ 4 & 1 & 5 & 2 \end{matrix}$
 $b = b_1 \cdot 10^{N/2} + b_2$

Fase Conquistar

$a_1 \cdot b_1 = \begin{matrix} 4 & 9 & 2 \end{matrix}$
 $a_1 \cdot b_2 = \begin{matrix} 6 & 2 & 4 \end{matrix}$
 $a_2 \cdot b_1 = \begin{matrix} 1 & 3 & 9 & 4 \end{matrix}$
 $a_2 \cdot b_2 = \begin{matrix} 1 & 7 & 6 & 8 \end{matrix}$

Fase Combinar

$c = a \cdot b = a_1 \cdot b_1 \cdot 10^N + (a_1 \cdot b_2 + a_2 \cdot b_1) \cdot 10^{N/2} + a_2 \cdot b_2$

$T(N) = \begin{cases} \theta(1) & N = 1 \\ 4 \cdot T\left(\frac{N}{2}\right) + \theta(N^1) & N > 1 \end{cases}$

Teorema Maestro
 (1^{er} Caso)
 $\log_2(4) = 2 > 1$

NO MEJORA EL ALGORITMO NATURAL

Complejidad:
 $\Theta(N^2)$

50

Multiplicar números grandes

Algoritmo de Karatsuba

N dígitos

$a = \begin{array}{|c|c|c|c|} \hline a_1 & a_2 & & \\ \hline 1 & 2 & 3 & 4 \\ \hline \end{array}$

$a = a_1 \cdot 10^{N/2} + a_2$

N dígitos

$b = \begin{array}{|c|c|c|c|} \hline b_1 & b_2 & & \\ \hline 4 & 1 & 5 & 2 \\ \hline \end{array}$

$b = b_1 \cdot 10^{N/2} + b_2$

$a_1 \cdot b_1 = \begin{array}{|c|c|c|} \hline 4 & 9 & 2 \\ \hline \end{array}$

$a_1 \cdot b_2 = \begin{array}{|c|c|c|} \hline 6 & 2 & 4 \\ \hline \end{array}$

$a_2 \cdot b_1 = \begin{array}{|c|c|c|c|} \hline 1 & 3 & 9 & 4 \\ \hline \end{array}$

$a_2 \cdot b_2 = \begin{array}{|c|c|c|c|} \hline 1 & 7 & 6 & 8 \\ \hline \end{array}$

4 multiplicaciones

$$c = a \cdot b = a_1 \cdot b_1 \cdot 10^N + (a_1 \cdot b_2 + a_2 \cdot b_1) \cdot 10^{N/2} + a_2 \cdot b_2$$

2. Multiplicar Números Grandes

① — ② — ③

51 de 98

51

Multiplicar números grandes

Algoritmo de Karatsuba

N dígitos

$a = \begin{array}{|c|c|c|c|} \hline a_1 & a_2 & & \\ \hline 1 & 2 & 3 & 4 \\ \hline \end{array}$

$a = a_1 \cdot 10^{N/2} + a_2$

N dígitos

$b = \begin{array}{|c|c|c|c|} \hline b_1 & b_2 & & \\ \hline 4 & 1 & 5 & 2 \\ \hline \end{array}$

$b = b_1 \cdot 10^{N/2} + b_2$

$a_1 \cdot b_1 = \begin{array}{|c|c|c|} \hline 4 & 9 & 2 \\ \hline \end{array}$

$p_3 = (a_1 + a_2) \cdot (b_1 + b_2) = \begin{array}{|c|c|c|c|} \hline 4 & 2 & 7 & 8 \\ \hline \end{array}$

$a_2 \cdot b_2 = \begin{array}{|c|c|c|c|} \hline 1 & 7 & 6 & 8 \\ \hline \end{array}$

3 multiplicaciones

$$c = a \cdot b = a_1 \cdot b_1 \cdot 10^N + (a_1 \cdot b_2 + a_2 \cdot b_1) \cdot 10^{N/2} + a_2 \cdot b_2$$

2. Multiplicar Números Grandes

① — ② — ③

52 de 98

52

Multiplicar números grandes

Algoritmo de Karatsuba

N dígitos

$a = \begin{array}{|c|c|c|c|} \hline a_1 & a_2 & & \\ \hline 1 & 2 & 3 & 4 \\ \hline \end{array}$

$a = a_1 \cdot 10^{N/2} + a_2$

N dígitos

$b = \begin{array}{|c|c|c|c|} \hline b_1 & b_2 & & \\ \hline 4 & 1 & 5 & 2 \\ \hline \end{array}$

$b = b_1 \cdot 10^{N/2} + b_2$

$a_1 \cdot b_1 = \begin{array}{|c|c|c|} \hline 4 & 9 & 2 \\ \hline \end{array}$

$p_3 = (a_1 + a_2) \cdot (b_1 + b_2) = \begin{array}{|c|c|c|c|} \hline 4 & 2 & 7 & 8 \\ \hline \end{array}$

$a_2 \cdot b_2 = \begin{array}{|c|c|c|c|} \hline 1 & 7 & 6 & 8 \\ \hline \end{array}$

3 multiplicaciones

$c = a \cdot b = a_1 \cdot b_1 \cdot 10^N + (p_3 - a_1 \cdot b_1 - a_2 \cdot b_2) \cdot 10^{N/2} + a_2 \cdot b_2$

2. Multiplicar Números Grandes

① — ② — ③

53 de 98

53

Multiplicar números grandes

Algoritmo de Karatsuba

N dígitos

$a = \begin{array}{|c|c|c|c|} \hline a_1 & a_2 & & \\ \hline 1 & 2 & 3 & 4 \\ \hline \end{array}$

$a = a_1 \cdot 10^{N/2} + a_2$

N dígitos

$b = \begin{array}{|c|c|c|c|} \hline b_1 & b_2 & & \\ \hline 4 & 1 & 5 & 2 \\ \hline \end{array}$

$b = b_1 \cdot 10^{N/2} + b_2$

$p_1 = a_1 \cdot b_1 = \begin{array}{|c|c|c|} \hline 4 & 9 & 2 \\ \hline \end{array}$

$p_2 = a_2 \cdot b_2 = \begin{array}{|c|c|c|c|} \hline 1 & 7 & 6 & 8 \\ \hline \end{array}$

$p_3 = (a_1 + a_2) \cdot (b_1 + b_2) = \begin{array}{|c|c|c|c|} \hline 4 & 2 & 7 & 8 \\ \hline \end{array}$

$c = a \cdot b = a_1 \cdot b_1 \cdot 10^N + (p_3 - a_1 \cdot b_1 - a_2 \cdot b_2) \cdot 10^{N/2} + a_2 \cdot b_2$

$\begin{array}{|c|c|c|c|c|c|c|} \hline 5 & 1 & 2 & 3 & 5 & 6 & 8 \\ \hline \end{array}$

2. Multiplicar Números Grandes

① — ② — ③

54 de 98

54

Multiplicar números grandes

Algoritmo de Karatsuba

N dígitos

$a = \begin{array}{|c|c|c|c|} \hline a_1 & a_2 & & \\ \hline 1 & 2 & 3 & 4 \\ \hline \end{array}$

$a = a_1 \cdot 10^{N/2} + a_2$

N dígitos

$b = \begin{array}{|c|c|c|c|} \hline b_1 & b_2 & & \\ \hline 4 & 1 & 5 & 2 \\ \hline \end{array}$

$b = b_1 \cdot 10^{N/2} + b_2$

Fase Dividir

$p_1 = a_1 \cdot b_1 = \begin{array}{|c|c|c|} \hline 4 & 9 & 2 \\ \hline \end{array}$

$p_2 = a_2 \cdot b_2 = \begin{array}{|c|c|c|c|} \hline 1 & 7 & 6 & 8 \\ \hline \end{array}$

$p_3 = (a_1 + a_2) \cdot (b_1 + b_2) = \begin{array}{|c|c|c|c|} \hline 4 & 2 & 7 & 8 \\ \hline \end{array}$

↓

$c = a \cdot b = \begin{array}{|c|c|c|c|c|c|} \hline 5 & 1 & 2 & 3 & 5 & 6 & 8 \\ \hline \end{array}$

$c = a \cdot b = a_1 \cdot b_1 \cdot 10^N + (p_3 - a_1 \cdot b_1 - a_2 \cdot b_2) \cdot 10^{N/2} + a_2 \cdot b_2$

2. Multiplicar Números Grandes

① — ② — ③

55 de 98

55

Multiplicar números grandes

Algoritmo de Karatsuba

N dígitos

$a = \begin{array}{|c|c|c|c|} \hline a_1 & a_2 & & \\ \hline 1 & 2 & 3 & 4 \\ \hline \end{array}$

$a = a_1 \cdot 10^{N/2} + a_2$

N dígitos

$b = \begin{array}{|c|c|c|c|} \hline b_1 & b_2 & & \\ \hline 4 & 1 & 5 & 2 \\ \hline \end{array}$

$b = b_1 \cdot 10^{N/2} + b_2$

↓

$p_1 = a_1 \cdot b_1 = \begin{array}{|c|c|c|} \hline 4 & 9 & 2 \\ \hline \end{array}$

$p_2 = a_2 \cdot b_2 = \begin{array}{|c|c|c|c|} \hline 1 & 7 & 6 & 8 \\ \hline \end{array}$

$p_3 = (a_1 + a_2) \cdot (b_1 + b_2) = \begin{array}{|c|c|c|c|} \hline 4 & 2 & 7 & 8 \\ \hline \end{array}$

Fase Conquistar

$c = a \cdot b = \begin{array}{|c|c|c|c|c|c|} \hline 5 & 1 & 2 & 3 & 5 & 6 & 8 \\ \hline \end{array}$

$c = a \cdot b = a_1 \cdot b_1 \cdot 10^N + (p_3 - a_1 \cdot b_1 - a_2 \cdot b_2) \cdot 10^{N/2} + a_2 \cdot b_2$

2. Multiplicar Números Grandes

① — ② — ③

56 de 98

56

Multiplicar números grandes

Algoritmo de Karatsuba

N dígitos

$a = \begin{bmatrix} a_1 & a_2 \end{bmatrix}$

$a = \begin{bmatrix} 1 & 2 & 3 & 4 \end{bmatrix}$

$a = a_1 \cdot 10^{N/2} + a_2$

N dígitos

$b = \begin{bmatrix} b_1 & b_2 \end{bmatrix}$

$b = \begin{bmatrix} 4 & 1 & 5 & 2 \end{bmatrix}$

$b = b_1 \cdot 10^{N/2} + b_2$

$p_1 = a_1 \cdot b_1 = \begin{bmatrix} 4 & 9 & 2 \end{bmatrix}$

$p_2 = a_2 \cdot b_2 = \begin{bmatrix} 1 & 7 & 6 & 8 \end{bmatrix}$

$p_3 = (a_1 + a_2) \cdot (b_1 + b_2) = \begin{bmatrix} 4 & 2 & 7 & 8 \end{bmatrix}$

$c = a \cdot b = \begin{bmatrix} 5 & 1 & 2 & 3 & 5 & 6 & 8 \end{bmatrix}$

$c = a \cdot b = a_1 \cdot b_1 \cdot 10^N + (p_3 - a_1 \cdot b_1 - a_2 \cdot b_2) \cdot 10^{N/2} + a_2 \cdot b_2$

Fase Combinar

os Grandes

① — ② — ③

57 de 98

57

Multiplicar números grandes

Actividad 7.8. Calcula la complejidad del algoritmo de Karatsuba ?

N dígitos

$a = \begin{bmatrix} a_1 & a_2 \end{bmatrix}$

$a = \begin{bmatrix} 1 & 2 & 3 & 4 \end{bmatrix}$

$a = a_1 \cdot 10^{N/2} + a_2$

N dígitos

$b = \begin{bmatrix} b_1 & b_2 \end{bmatrix}$

$b = \begin{bmatrix} 4 & 1 & 5 & 2 \end{bmatrix}$

$b = b_1 \cdot 10^{N/2} + b_2$

$p_1 = a_1 \cdot b_1 = \begin{bmatrix} 4 & 9 & 2 \end{bmatrix}$

$p_2 = a_2 \cdot b_2 = \begin{bmatrix} 1 & 7 & 6 & 8 \end{bmatrix}$

$p_3 = (a_1 + a_2) \cdot (b_1 + b_2) = \begin{bmatrix} 4 & 2 & 7 & 8 \end{bmatrix}$

$c = a \cdot b = \begin{bmatrix} 5 & 1 & 2 & 3 & 5 & 6 & 8 \end{bmatrix}$

$c = a \cdot b = a_1 \cdot b_1 \cdot 10^N + (p_3 - a_1 \cdot b_1 - a_2 \cdot b_2) \cdot 10^{N/2} + a_2 \cdot b_2$

Fase Combinar

os Grandes

① — ② — ③

58 de 98

58

Multiplicar números grandes

Actividad 7.8. Calcula la complejidad del algoritmo de Karatsuba ?

Fase Dividir

N dígitos

$a = \begin{bmatrix} 1 & 2 & 3 & 4 \end{bmatrix}$
 $a = a_1 \cdot 10^{N/2} + a_2$

N dígitos

$b = \begin{bmatrix} 4 & 1 & 5 & 2 \end{bmatrix}$
 $b = b_1 \cdot 10^{N/2} + b_2$

Fase Conquistar

$p_1 = a_1 \cdot b_1 = \begin{bmatrix} 4 & 9 & 2 \end{bmatrix}$

$p_2 = a_2 \cdot b_2 = \begin{bmatrix} 1 & 7 & 6 & 8 \end{bmatrix}$

$p_3 = (a_1 + a_2) \cdot (b_1 + b_2) = \begin{bmatrix} 4 & 2 & 7 & 8 \end{bmatrix}$

Fase Combinar

$c = a \cdot b = \begin{bmatrix} 5 & 1 & 2 & 3 & 5 & 6 & 8 \end{bmatrix}$
 $c = a \cdot b = a_1 \cdot b_1 \cdot 10^N + (p_3 - a_1 \cdot b_1 - a_2 \cdot b_2) \cdot 10^{N/2} + a_2 \cdot b_2$

$$T(N) = \begin{cases} \Theta(1) & N = 1 \\ 3 \cdot T\left(\frac{N}{2}\right) + \Theta(N) & N > 1 \end{cases}$$

59 de 98

59

Multiplicar números grandes

Actividad 7.8. Calcula la complejidad del algoritmo de Karatsuba ?

Fase Dividir

N dígitos

$a = \begin{bmatrix} 1 & 2 & 3 & 4 \end{bmatrix}$
 $a = a_1 \cdot 10^{N/2} + a_2$

N dígitos

$b = \begin{bmatrix} 4 & 1 & 5 & 2 \end{bmatrix}$
 $b = b_1 \cdot 10^{N/2} + b_2$

Fase Conquistar

$p_1 = a_1 \cdot b_1 = \begin{bmatrix} 4 & 9 & 2 \end{bmatrix}$

$p_2 = a_2 \cdot b_2 = \begin{bmatrix} 1 & 7 & 6 & 8 \end{bmatrix}$

$p_3 = (a_1 + a_2) \cdot (b_1 + b_2) = \begin{bmatrix} 4 & 2 & 7 & 8 \end{bmatrix}$

Fase Combinar

$c = a \cdot b = \begin{bmatrix} 5 & 1 & 2 & 3 & 5 & 6 & 8 \end{bmatrix}$
 $c = a \cdot b = a_1 \cdot b_1 \cdot 10^N + (p_3 - a_1 \cdot b_1 - a_2 \cdot b_2) \cdot 10^{N/2} + a_2 \cdot b_2$

$$T(N) = \begin{cases} \Theta(1) & N = 1 \\ 3 \cdot T\left(\frac{N}{2}\right) + \Theta(N^1) & N > 1 \end{cases}$$

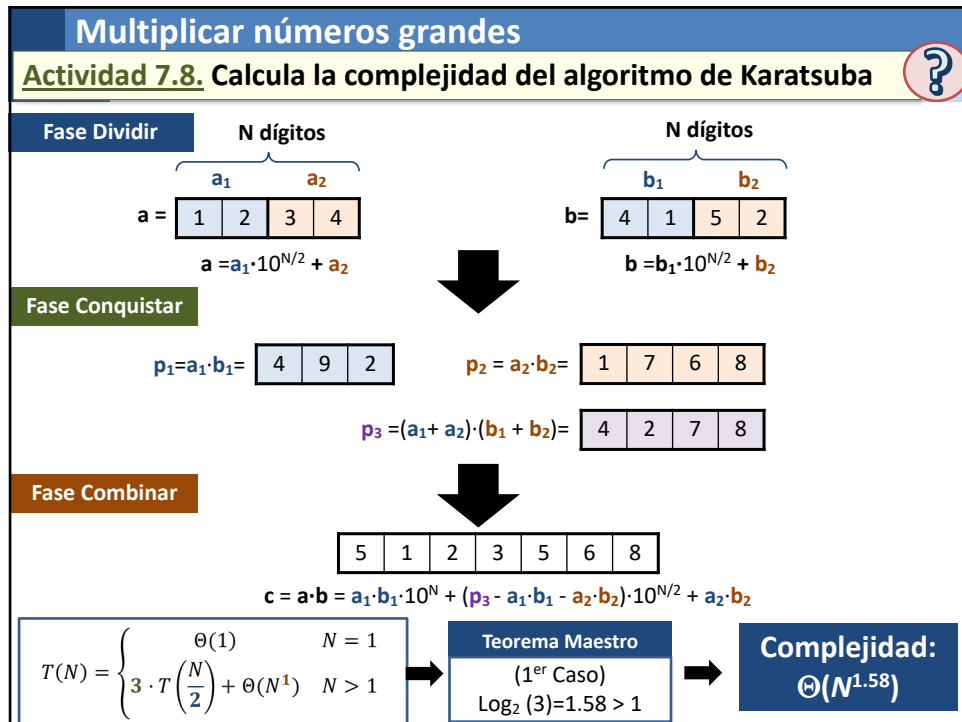
Teorema Maestro

(1^{er} Caso)

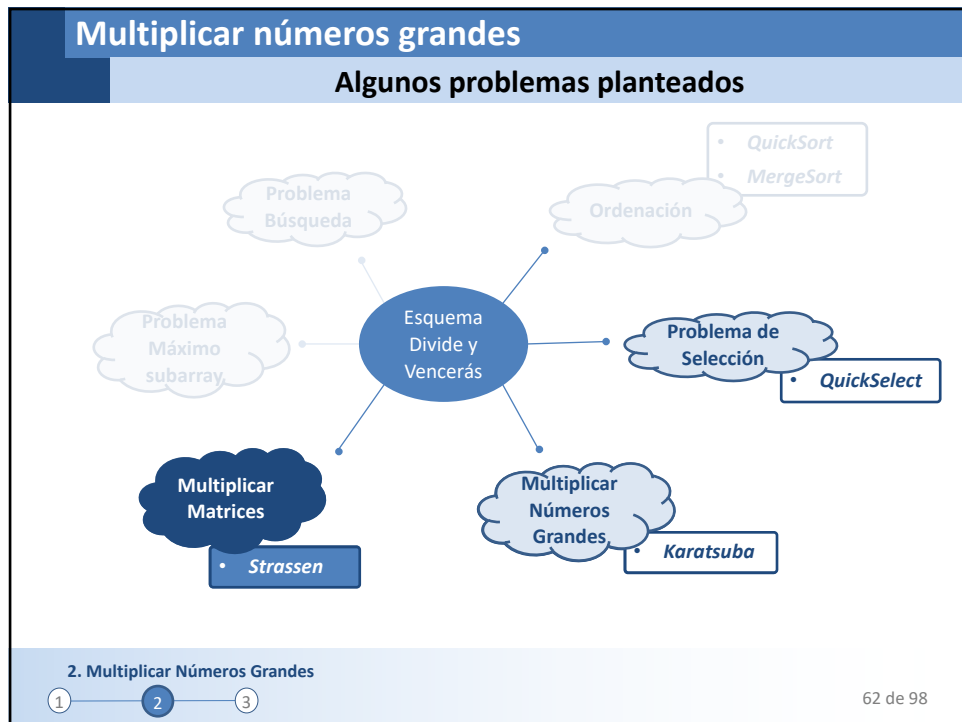
$\log_2(3) = 1.58 > 1$

60 de 98

60



61



62

Multiplicar Matrices

Enunciado del Problema

Dadas dos matrices cuadradas:

A =

| | | | | |
|---|---|---|---|---|
| 1 | 2 | 3 | 4 | 5 |
| 3 | 7 | 5 | 2 | 7 |
| 2 | 3 | 3 | 4 | 5 |
| 4 | 1 | 6 | 8 | 0 |
| 1 | 0 | 3 | 4 | 3 |

B =

| | | | | |
|---|---|---|---|---|
| 1 | 1 | 0 | 3 | 4 |
| 5 | 2 | 8 | 1 | 2 |
| 3 | 3 | 3 | 8 | 1 |
| 4 | 0 | 5 | 4 | 5 |
| 2 | 6 | 2 | 2 | 3 |

↓

Multiplicarlas:

A·B =

| | | | | |
|----|----|----|----|----|
| 46 | 44 | 55 | 55 | 46 |
| 75 | 74 | 95 | 78 | 62 |
| 52 | 47 | 63 | 59 | 52 |
| 59 | 24 | 66 | 93 | 64 |
| 32 | 28 | 35 | 49 | 36 |

1
2
3

3. Multiplicar Matrices

63

Multiplicar Matrices

Algoritmo Natural

Actividad 7.9. Diseña e implementa un algoritmo que multiplique dos matrices

```
int[][] multiplicar(int[][] A, int[][] B){
    int[][] resul= new int[A.length][A[0].length];
    for (int f=0; f<A.length; f++)
        for (int c=0; c<A[0].length; c++) {
            resul[f][c]=0;
            for (int i=0; i<A[0].length; i++)
                resul[f][c]+=A[f][i]*B[i][c];
        }
    return resul;
}
```

1
2
3

3. Multiplicar Matrices

64

Multiplicar Matrices

Algoritmo Natural

Actividad 7.10. Calcula la complejidad del Algoritmo

```

int[][] multiplicar(int[][] A, int[][] B){
    int[][] resul= new int[A.length][A[0].length];
    for (int f=0; f<A.length; f++)
        for (int c=0; c<A[0].length; c++) {
            resul[f][c]=0;
            for (int i=0; i<A[0].length; i++)
                resul[f][c]+=A[f][i]*B[i][c];
        }
    return resul;
}
    
```

Complejidad:
 $\Theta(N^3)$

3. Multiplicar Matrices

1
2
3

65

Multiplicar Matrices

Algoritmo inmediato basado en Divide y Vencerás

| | | | |
|---|---|---|---|
| 1 | 2 | 3 | 5 |
| 3 | 7 | 5 | 7 |
| 2 | 3 | 3 | 5 |
| 1 | 0 | 3 | 3 |

A

| | | | |
|---|---|---|---|
| 1 | 1 | 0 | 4 |
| 5 | 2 | 8 | 2 |
| 3 | 3 | 3 | 1 |
| 2 | 6 | 2 | 3 |

B

Vamos a considerar una matriz cuadrada de dimensión 2^k

| | | | |
|----|----|----|----|
| 30 | 44 | 35 | 26 |
| 67 | 74 | 85 | 52 |
| 36 | 47 | 43 | 32 |
| 16 | 28 | 15 | 16 |

C = A · B

3. Multiplicar Matrices

1
2
3

66

Multiplicar Matrices

Algoritmo inmediato basado en Divide y Vencerás

| | | | |
|---|---|---|---|
| 1 | 2 | 3 | 5 |
| 3 | 7 | 5 | 7 |
| 2 | 3 | 3 | 5 |
| 1 | 0 | 3 | 3 |

| | | | |
|---|---|---|---|
| 1 | 1 | 0 | 4 |
| 5 | 2 | 8 | 2 |
| 3 | 3 | 3 | 1 |
| 2 | 6 | 2 | 3 |

A

B

| | | | |
|----|----|----|----|
| 30 | 44 | 35 | 26 |
| 67 | 74 | 85 | 52 |
| 36 | 47 | 43 | 32 |
| 16 | 28 | 15 | 16 |

C = A · B

3. Multiplicar Matrices

① — ② — ③

67

Multiplicar Matrices

Algoritmo inmediato basado en Divide y Vencerás

A_{11}

| | | | |
|---|---|---|---|
| 1 | 2 | 3 | 5 |
| 3 | 7 | 5 | 7 |

 A_{12}

B_{11}

| | | | |
|---|---|---|---|
| 1 | 1 | 0 | 4 |
| 5 | 2 | 8 | 2 |

 B_{12}

A_{21}

| | | | |
|---|---|---|---|
| 2 | 3 | 3 | 5 |
| 1 | 0 | 3 | 3 |

 A_{22}

B_{21}

| | | | |
|---|---|---|---|
| 3 | 3 | 3 | 1 |
| 2 | 6 | 2 | 3 |

 B_{22}

A

B

| | | | |
|----|----|----|----|
| 30 | 44 | 35 | 26 |
| 67 | 74 | 85 | 52 |
| 36 | 47 | 43 | 32 |
| 16 | 28 | 15 | 16 |

C = A · B

$C_{11} = A_{11} \cdot B_{11} + A_{12} \cdot B_{21}$

 $C_{21} = A_{21} \cdot B_{11} + A_{22} \cdot B_{21}$

$C_{12} = A_{11} \cdot B_{12} + A_{12} \cdot B_{22}$

 $C_{22} = A_{21} \cdot B_{12} + A_{22} \cdot B_{22}$

3. Multiplicar Matrices

① — ② — ③

68

Multiplicar Matrices

Algoritmo inmediato basado en Divide y Vencerás

A_{11}

| | | | |
|---|---|---|---|
| 1 | 2 | 3 | 5 |
| 3 | 7 | 5 | 7 |
| 2 | 3 | 3 | 5 |
| 1 | 0 | 3 | 3 |

 A_{12}
 A_{21}
 A_{22}
A

B_{11}

| | | | |
|---|---|---|---|
| 1 | 1 | 0 | 4 |
| 5 | 2 | 8 | 2 |
| 3 | 3 | 3 | 1 |
| 2 | 6 | 2 | 3 |

 B_{12}
 B_{21}
 B_{22}
B

$C_{11} = A_{11} \cdot B_{11} + A_{12} \cdot B_{21}$
 $C_{21} = A_{21} \cdot B_{11} + A_{22} \cdot B_{21}$

| | | | |
|----|----|----|----|
| 30 | 44 | 35 | 26 |
| 67 | 74 | 85 | 52 |
| 36 | 47 | 43 | 32 |
| 16 | 28 | 15 | 16 |

 $C = A \cdot B$

$C_{12} = A_{11} \cdot B_{12} + A_{12} \cdot B_{22}$
 $C_{22} = A_{21} \cdot B_{12} + A_{22} \cdot B_{22}$

3. Multiplicar Matrices

1
2
3

69

Multiplicar Matrices

Algoritmo inmediato basado en Divide y Vencerás

A_{11}

| | | | |
|---|---|---|---|
| 1 | 2 | 3 | 5 |
| 3 | 7 | 5 | 7 |
| 2 | 3 | 3 | 5 |
| 1 | 0 | 3 | 3 |

 A_{12}
 A_{21}
 A_{22}
A

B_{11}

| | | | |
|---|---|---|---|
| 1 | 1 | 0 | 4 |
| 5 | 2 | 8 | 2 |
| 3 | 3 | 3 | 1 |
| 2 | 6 | 2 | 3 |

 B_{12}
 B_{21}
 B_{22}
B

$C_{11} = A_{11} \cdot B_{11} + A_{12} \cdot B_{21}$
 $C_{21} = A_{21} \cdot B_{11} + A_{22} \cdot B_{21}$

| | | | |
|----|----|----|----|
| 30 | 44 | 35 | 26 |
| 67 | 74 | 85 | 52 |
| 36 | 47 | 43 | 32 |
| 16 | 28 | 15 | 16 |

 $C = A \cdot B$

$C_{12} = A_{11} \cdot B_{12} + A_{12} \cdot B_{22}$
 $C_{22} = A_{21} \cdot B_{12} + A_{22} \cdot B_{22}$

3. Multiplicar Matrices

1
2
3

70

Multiplicar Matrices

Algoritmo inmediato basado en Divide y Vencerás

A_{11}

| | | | |
|---|---|---|---|
| 1 | 2 | 3 | 5 |
| 3 | 7 | 5 | 7 |

 A_{12}

B_{11}

| | | | |
|---|---|---|---|
| 1 | 1 | 0 | 4 |
| 5 | 2 | 8 | 2 |

 B_{12}

A_{21}

| | | | |
|---|---|---|---|
| 2 | 3 | 3 | 5 |
| 1 | 0 | 3 | 3 |

 A_{22}

B_{21}

| | | | |
|---|---|---|---|
| 3 | 3 | 3 | 1 |
| 2 | 6 | 2 | 3 |

 B_{22}

A **B**

$C_{11} = A_{11} \cdot B_{11} + A_{12} \cdot B_{21}$

| | | | |
|----|----|----|----|
| 30 | 44 | 35 | 26 |
| 67 | 74 | 85 | 52 |

$C_{12} = A_{11} \cdot B_{12} + A_{12} \cdot B_{22}$

$C_{21} = A_{21} \cdot B_{11} + A_{22} \cdot B_{21}$

| | | | |
|----|----|-----------|-----------|
| 36 | 47 | 43 | 32 |
| 16 | 28 | 15 | 16 |

$C_{22} = A_{21} \cdot B_{12} + A_{22} \cdot B_{22}$

C = A · B

3. Multiplicar Matrices
 ① — ② — ③

71

Multiplicar Matrices

Algoritmo inmediato basado en Divide y Vencerás

A_{11}

| | | | |
|---|---|---|---|
| 1 | 2 | 3 | 5 |
| 3 | 7 | 5 | 7 |

 A_{12}

B_{11}

| | | | |
|---|---|---|---|
| 1 | 1 | 0 | 4 |
| 5 | 2 | 8 | 2 |

 B_{12}

A_{21}

| | | | |
|---|---|---|---|
| 2 | 3 | 3 | 5 |
| 1 | 0 | 3 | 3 |

 A_{22}

B_{21}

| | | | |
|---|---|---|---|
| 3 | 3 | 3 | 1 |
| 2 | 6 | 2 | 3 |

 B_{22}

A **B**

$C_{11} = A_{11} \cdot B_{11} + A_{12} \cdot B_{21}$

| | | | |
|----|----|----|----|
| 30 | 44 | 35 | 26 |
| 67 | 74 | 85 | 52 |

$C_{12} = A_{11} \cdot B_{12} + A_{12} \cdot B_{22}$

$C_{21} = A_{21} \cdot B_{11} + A_{22} \cdot B_{21}$

| | | | |
|-----------|-----------|----|----|
| 36 | 47 | 43 | 32 |
| 16 | 28 | 15 | 16 |

$C_{22} = A_{21} \cdot B_{12} + A_{22} \cdot B_{22}$

C = A · B

3. Multiplicar Matrices
 ① — ② — ③

72

Multiplicar Matrices

Algoritmo inmediato basado en Divide y Vencerás

A_{11}

| | | | |
|---|---|---|---|
| 1 | 2 | 3 | 5 |
| 3 | 7 | 5 | 7 |
| 2 | 3 | 3 | 5 |
| 1 | 0 | 3 | 3 |

 A_{12}

B_{11}

| | | | |
|---|---|---|---|
| 1 | 1 | 0 | 4 |
| 5 | 2 | 8 | 2 |
| 3 | 3 | 3 | 1 |
| 2 | 6 | 2 | 3 |

 B_{12}

A_{21}

| | | | |
|---|---|---|---|
| 2 | 3 | 3 | 5 |
| 1 | 0 | 3 | 3 |

 A_{22}

B_{21}

| | | | |
|---|---|---|---|
| 3 | 3 | 3 | 1 |
| 2 | 6 | 2 | 3 |

 B_{22}

| | | | |
|----|----|----|----|
| 30 | 44 | 35 | 26 |
| 67 | 74 | 85 | 52 |
| 36 | 47 | 43 | 32 |
| 16 | 28 | 15 | 16 |

$C = A \cdot B$

$C_{11} = A_{11} \cdot B_{11} + A_{12} \cdot B_{21}$

$C_{12} = A_{11} \cdot B_{12} + A_{12} \cdot B_{22}$

$C_{21} = A_{21} \cdot B_{11} + A_{22} \cdot B_{21}$

$C_{22} = A_{21} \cdot B_{12} + A_{22} \cdot B_{22}$

3. Multiplicar Matrices

① — ② — ③

73

Multiplicar Matrices

Algoritmo inmediato basado en Divide y Vencerás

| | | | |
|---|---|---|---|
| 1 | 2 | 3 | 5 |
| 3 | 7 | 5 | 7 |
| 2 | 3 | 3 | 5 |
| 1 | 0 | 3 | 3 |

A

| | | | |
|---|---|---|---|
| 1 | 1 | 0 | 4 |
| 5 | 2 | 8 | 2 |
| 3 | 3 | 3 | 1 |
| 2 | 6 | 2 | 3 |

B

3. Multiplicar Matrices

① — ② — ③

74

Multiplicar Matrices

Algoritmo inmediato basado en Divide y Vencerás

A_{11}

| | |
|---|---|
| 1 | 2 |
| 3 | 7 |

 A_{12}

| | |
|---|---|
| 3 | 5 |
| 5 | 7 |

B_{11}

| | |
|---|---|
| 1 | 1 |
| 5 | 2 |

 B_{12}

| | |
|---|---|
| 0 | 4 |
| 8 | 2 |

A_{21}

| | |
|---|---|
| 2 | 3 |
| 1 | 0 |

 A_{22}

| | |
|---|---|
| 3 | 5 |
| 3 | 3 |

B_{21}

| | |
|---|---|
| 3 | 3 |
| 2 | 6 |

 B_{22}

| | |
|---|---|
| 3 | 1 |
| 2 | 3 |

Fase Dividir

3. Multiplicar Matrices

① — ② — ③

75

Multiplicar Matrices

Algoritmo inmediato basado en Divide y Vencerás

A_{11}

| | |
|---|---|
| 1 | 2 |
| 3 | 7 |

 A_{12}

| | |
|---|---|
| 3 | 5 |
| 5 | 7 |

B_{11}

| | |
|---|---|
| 1 | 1 |
| 5 | 2 |

 B_{12}

| | |
|---|---|
| 0 | 4 |
| 8 | 2 |

A_{21}

| | |
|---|---|
| 2 | 3 |
| 1 | 0 |

 A_{22}

| | |
|---|---|
| 3 | 5 |
| 3 | 3 |

B_{21}

| | |
|---|---|
| 3 | 3 |
| 2 | 6 |

 B_{22}

| | |
|---|---|
| 3 | 1 |
| 2 | 3 |

$A_{11} \cdot B_{11} =$

| | |
|----|----|
| 11 | 5 |
| 38 | 17 |

$A_{12} \cdot B_{21} =$

| | |
|----|----|
| 19 | 39 |
| 29 | 57 |

$A_{11} \cdot B_{12} =$

| | |
|----|----|
| 16 | 8 |
| 56 | 26 |

$A_{12} \cdot B_{22} =$

| | |
|----|----|
| 19 | 18 |
| 29 | 26 |

$A_{21} \cdot B_{11} =$

| | |
|----|---|
| 17 | 8 |
| 1 | 1 |

$A_{22} \cdot B_{21} =$

| | |
|----|----|
| 19 | 39 |
| 15 | 27 |

$A_{21} \cdot B_{12} =$

| | |
|----|----|
| 24 | 14 |
| 0 | 4 |

$A_{22} \cdot B_{22} =$

| | |
|----|----|
| 19 | 18 |
| 15 | 12 |

Fase Conquistar

3. Multiplicar Matrices

① — ② — ③

76

Multiplicar Matrices

Algoritmo inmediato basado en Divide y Vencerás

$A_{11} = \begin{bmatrix} 1 & 2 \\ 3 & 7 \end{bmatrix}$
 $A_{21} = \begin{bmatrix} 2 & 3 \\ 1 & 0 \end{bmatrix}$

$A_{12} = \begin{bmatrix} 3 & 5 \\ 5 & 7 \end{bmatrix}$
 $A_{22} = \begin{bmatrix} 3 & 5 \\ 3 & 3 \end{bmatrix}$

$B_{11} = \begin{bmatrix} 1 & 1 \\ 5 & 2 \end{bmatrix}$
 $B_{21} = \begin{bmatrix} 3 & 3 \\ 2 & 6 \end{bmatrix}$

$B_{12} = \begin{bmatrix} 0 & 4 \\ 8 & 2 \end{bmatrix}$
 $B_{22} = \begin{bmatrix} 3 & 1 \\ 2 & 3 \end{bmatrix}$

$A_{11} \cdot B_{11} = \begin{bmatrix} 11 & 5 \\ 38 & 17 \end{bmatrix}$
 $A_{21} \cdot B_{11} = \begin{bmatrix} 17 & 8 \\ 1 & 1 \end{bmatrix}$

$A_{12} \cdot B_{21} = \begin{bmatrix} 19 & 39 \\ 29 & 57 \end{bmatrix}$
 $A_{22} \cdot B_{21} = \begin{bmatrix} 19 & 39 \\ 15 & 27 \end{bmatrix}$

$A_{11} \cdot B_{12} = \begin{bmatrix} 16 & 8 \\ 56 & 26 \end{bmatrix}$
 $A_{21} \cdot B_{12} = \begin{bmatrix} 24 & 14 \\ 0 & 4 \end{bmatrix}$

$A_{12} \cdot B_{22} = \begin{bmatrix} 19 & 18 \\ 29 & 26 \end{bmatrix}$
 $A_{22} \cdot B_{22} = \begin{bmatrix} 19 & 18 \\ 15 & 12 \end{bmatrix}$

$C_{11} = A_{11} \cdot B_{11} + A_{12} \cdot B_{21}$
 $C_{21} = A_{21} \cdot B_{11} + A_{22} \cdot B_{21}$

$\begin{bmatrix} 30 & 44 \\ 67 & 74 \end{bmatrix}$
 $\begin{bmatrix} 36 & 47 \\ 16 & 28 \end{bmatrix}$

$C_{12} = A_{11} \cdot B_{12} + A_{12} \cdot B_{22}$
 $C_{22} = A_{21} \cdot B_{12} + A_{22} \cdot B_{22}$

$\begin{bmatrix} 35 & 26 \\ 85 & 52 \end{bmatrix}$
 $\begin{bmatrix} 43 & 32 \\ 15 & 16 \end{bmatrix}$

① Fase Combinar

77

Multiplicar Matrices

Algoritmo inmediato basado en Divide y Vencerás

$A_{11} = \begin{bmatrix} 1 & 2 \\ 3 & 7 \end{bmatrix}$
 $A_{21} = \begin{bmatrix} 2 & 3 \\ 1 & 0 \end{bmatrix}$

$A_{12} = \begin{bmatrix} 3 & 5 \\ 5 & 7 \end{bmatrix}$
 $A_{22} = \begin{bmatrix} 3 & 5 \\ 3 & 3 \end{bmatrix}$

$B_{11} = \begin{bmatrix} 1 & 1 \\ 5 & 2 \end{bmatrix}$
 $B_{21} = \begin{bmatrix} 3 & 3 \\ 2 & 6 \end{bmatrix}$

$B_{12} = \begin{bmatrix} 0 & 4 \\ 8 & 2 \end{bmatrix}$
 $B_{22} = \begin{bmatrix} 3 & 1 \\ 2 & 3 \end{bmatrix}$

$A_{11} \cdot B_{11} = \begin{bmatrix} 11 & 5 \\ 38 & 17 \end{bmatrix}$
 $A_{21} \cdot B_{11} = \begin{bmatrix} 17 & 8 \\ 1 & 1 \end{bmatrix}$

$A_{12} \cdot B_{21} = \begin{bmatrix} 19 & 39 \\ 29 & 57 \end{bmatrix}$
 $A_{22} \cdot B_{21} = \begin{bmatrix} 19 & 39 \\ 15 & 27 \end{bmatrix}$

$A_{11} \cdot B_{12} = \begin{bmatrix} 16 & 8 \\ 56 & 26 \end{bmatrix}$
 $A_{21} \cdot B_{12} = \begin{bmatrix} 24 & 14 \\ 0 & 4 \end{bmatrix}$

$A_{12} \cdot B_{22} = \begin{bmatrix} 19 & 18 \\ 29 & 26 \end{bmatrix}$
 $A_{22} \cdot B_{22} = \begin{bmatrix} 19 & 18 \\ 15 & 12 \end{bmatrix}$

$C_{11} = A_{11} \cdot B_{11} + A_{12} \cdot B_{21}$
 $C_{21} = A_{21} \cdot B_{11} + A_{22} \cdot B_{21}$

$\begin{bmatrix} 30 & 44 \\ 67 & 74 \end{bmatrix}$
 $\begin{bmatrix} 36 & 47 \\ 16 & 28 \end{bmatrix}$

$C_{12} = A_{11} \cdot B_{12} + A_{12} \cdot B_{22}$
 $C_{22} = A_{21} \cdot B_{12} + A_{22} \cdot B_{22}$

$\begin{bmatrix} 35 & 26 \\ 85 & 52 \end{bmatrix}$
 $\begin{bmatrix} 43 & 32 \\ 15 & 16 \end{bmatrix}$

① Fase Combinar

78

Multiplicar Matrices

Algoritmo inmediato basado en Divide y Vencerás

$A_{11} = \begin{bmatrix} 1 & 2 \\ 3 & 7 \end{bmatrix}$
 $A_{21} = \begin{bmatrix} 2 & 3 \\ 1 & 0 \end{bmatrix}$

$A_{12} = \begin{bmatrix} 3 & 5 \\ 5 & 7 \end{bmatrix}$
 $A_{22} = \begin{bmatrix} 3 & 5 \\ 3 & 3 \end{bmatrix}$

$B_{11} = \begin{bmatrix} 1 & 1 \\ 5 & 2 \end{bmatrix}$
 $B_{21} = \begin{bmatrix} 3 & 3 \\ 2 & 6 \end{bmatrix}$

$B_{12} = \begin{bmatrix} 0 & 4 \\ 8 & 2 \end{bmatrix}$
 $B_{22} = \begin{bmatrix} 3 & 1 \\ 2 & 3 \end{bmatrix}$

$A_{11} \cdot B_{11} = \begin{bmatrix} 11 & 5 \\ 38 & 17 \end{bmatrix}$
 $A_{21} \cdot B_{11} = \begin{bmatrix} 17 & 8 \\ 1 & 1 \end{bmatrix}$

$A_{12} \cdot B_{21} = \begin{bmatrix} 19 & 39 \\ 29 & 57 \end{bmatrix}$
 $A_{22} \cdot B_{21} = \begin{bmatrix} 19 & 39 \\ 15 & 27 \end{bmatrix}$

$A_{11} \cdot B_{12} = \begin{bmatrix} 16 & 8 \\ 56 & 26 \end{bmatrix}$
 $A_{21} \cdot B_{12} = \begin{bmatrix} 24 & 14 \\ 0 & 4 \end{bmatrix}$

$A_{12} \cdot B_{22} = \begin{bmatrix} 19 & 18 \\ 29 & 26 \end{bmatrix}$
 $A_{22} \cdot B_{22} = \begin{bmatrix} 19 & 18 \\ 15 & 12 \end{bmatrix}$

$C_{11} = A_{11} \cdot B_{11} + A_{12} \cdot B_{21}$
 $C_{21} = A_{21} \cdot B_{11} + A_{22} \cdot B_{21}$

$C_{12} = A_{11} \cdot B_{12} + A_{12} \cdot B_{22}$
 $C_{22} = A_{21} \cdot B_{12} + A_{22} \cdot B_{22}$

| | |
|----|----|
| 30 | 44 |
| 67 | 74 |

| | |
|----|----|
| 35 | 26 |
| 85 | 52 |

| | |
|----|----|
| 36 | 47 |
| 16 | 28 |

1 **Fase Combinar**

79

Multiplicar Matrices

Algoritmo inmediato basado en Divide y Vencerás

$A_{11} = \begin{bmatrix} 1 & 2 \\ 3 & 7 \end{bmatrix}$
 $A_{21} = \begin{bmatrix} 2 & 3 \\ 1 & 0 \end{bmatrix}$

$A_{12} = \begin{bmatrix} 3 & 5 \\ 5 & 7 \end{bmatrix}$
 $A_{22} = \begin{bmatrix} 3 & 5 \\ 3 & 3 \end{bmatrix}$

$B_{11} = \begin{bmatrix} 1 & 1 \\ 5 & 2 \end{bmatrix}$
 $B_{21} = \begin{bmatrix} 3 & 3 \\ 2 & 6 \end{bmatrix}$

$B_{12} = \begin{bmatrix} 0 & 4 \\ 8 & 2 \end{bmatrix}$
 $B_{22} = \begin{bmatrix} 3 & 1 \\ 2 & 3 \end{bmatrix}$

$A_{11} \cdot B_{11} = \begin{bmatrix} 11 & 5 \\ 38 & 17 \end{bmatrix}$
 $A_{21} \cdot B_{11} = \begin{bmatrix} 17 & 8 \\ 1 & 1 \end{bmatrix}$

$A_{12} \cdot B_{21} = \begin{bmatrix} 19 & 39 \\ 29 & 57 \end{bmatrix}$
 $A_{22} \cdot B_{21} = \begin{bmatrix} 19 & 39 \\ 15 & 27 \end{bmatrix}$

$A_{11} \cdot B_{12} = \begin{bmatrix} 16 & 8 \\ 56 & 26 \end{bmatrix}$
 $A_{21} \cdot B_{12} = \begin{bmatrix} 24 & 14 \\ 0 & 4 \end{bmatrix}$

$A_{12} \cdot B_{22} = \begin{bmatrix} 19 & 18 \\ 29 & 26 \end{bmatrix}$
 $A_{22} \cdot B_{22} = \begin{bmatrix} 19 & 18 \\ 15 & 12 \end{bmatrix}$

$C_{11} = A_{11} \cdot B_{11} + A_{12} \cdot B_{21}$
 $C_{21} = A_{21} \cdot B_{11} + A_{22} \cdot B_{21}$

$C_{12} = A_{11} \cdot B_{12} + A_{12} \cdot B_{22}$
 $C_{22} = A_{21} \cdot B_{12} + A_{22} \cdot B_{22}$

| | |
|----|----|
| 30 | 44 |
| 67 | 74 |

| | |
|----|----|
| 35 | 26 |
| 85 | 52 |

| | |
|----|----|
| 36 | 47 |
| 16 | 28 |

1 **Fase Combinar**

80

Multiplicar Matrices

Algoritmo inmediato basado en Divide y Vencerás

$A_{11} = \begin{bmatrix} 1 & 2 \\ 3 & 7 \end{bmatrix}$
 $A_{21} = \begin{bmatrix} 2 & 3 \\ 1 & 0 \end{bmatrix}$

$A_{12} = \begin{bmatrix} 3 & 5 \\ 5 & 7 \end{bmatrix}$
 $A_{22} = \begin{bmatrix} 3 & 5 \\ 3 & 3 \end{bmatrix}$

$B_{11} = \begin{bmatrix} 1 & 1 \\ 5 & 2 \end{bmatrix}$
 $B_{21} = \begin{bmatrix} 3 & 3 \\ 2 & 6 \end{bmatrix}$

$B_{12} = \begin{bmatrix} 0 & 4 \\ 8 & 2 \end{bmatrix}$
 $B_{22} = \begin{bmatrix} 3 & 1 \\ 2 & 3 \end{bmatrix}$

$A_{11} \cdot B_{11} = \begin{bmatrix} 11 & 5 \\ 38 & 17 \end{bmatrix}$
 $A_{21} \cdot B_{11} = \begin{bmatrix} 17 & 8 \\ 1 & 1 \end{bmatrix}$

$A_{12} \cdot B_{21} = \begin{bmatrix} 19 & 39 \\ 29 & 57 \end{bmatrix}$
 $A_{22} \cdot B_{21} = \begin{bmatrix} 19 & 39 \\ 15 & 27 \end{bmatrix}$

$A_{11} \cdot B_{12} = \begin{bmatrix} 16 & 8 \\ 56 & 26 \end{bmatrix}$
 $A_{21} \cdot B_{12} = \begin{bmatrix} 24 & 14 \\ 0 & 4 \end{bmatrix}$

$A_{12} \cdot B_{22} = \begin{bmatrix} 19 & 18 \\ 29 & 26 \end{bmatrix}$
 $A_{22} \cdot B_{22} = \begin{bmatrix} 19 & 18 \\ 15 & 12 \end{bmatrix}$

$C_{11} = A_{11} \cdot B_{11} + A_{12} \cdot B_{21}$
 $C_{21} = A_{21} \cdot B_{11} + A_{22} \cdot B_{21}$

$C_{12} = A_{11} \cdot B_{12} + A_{12} \cdot B_{22}$
 $C_{22} = A_{21} \cdot B_{12} + A_{22} \cdot B_{22}$

| | |
|----|----|
| 30 | 44 |
| 67 | 74 |

| | |
|----|----|
| 35 | 26 |
| 85 | 52 |

| | |
|----|----|
| 36 | 47 |
| 16 | 28 |

| | |
|----|----|
| 43 | 32 |
| 15 | 16 |

1 **Fase Combinar**

81

Actividad 7.11. Calcula la complejidad de este algoritmo basado en divide y vencerás para multiplicar matrices cuadradas

Fase Dividir

$A_{11} = \begin{bmatrix} 1 & 2 \\ 3 & 7 \end{bmatrix}$
 $A_{21} = \begin{bmatrix} 2 & 3 \\ 1 & 0 \end{bmatrix}$

$A_{12} = \begin{bmatrix} 3 & 5 \\ 5 & 7 \end{bmatrix}$
 $A_{22} = \begin{bmatrix} 3 & 5 \\ 3 & 3 \end{bmatrix}$

$B_{11} = \begin{bmatrix} 1 & 1 \\ 5 & 2 \end{bmatrix}$
 $B_{21} = \begin{bmatrix} 3 & 3 \\ 2 & 6 \end{bmatrix}$

$B_{12} = \begin{bmatrix} 0 & 4 \\ 8 & 2 \end{bmatrix}$
 $B_{22} = \begin{bmatrix} 3 & 1 \\ 2 & 3 \end{bmatrix}$

Fase Conquistar

$A_{11} \cdot B_{11} = \begin{bmatrix} 11 & 5 \\ 38 & 17 \end{bmatrix}$
 $A_{21} \cdot B_{11} = \begin{bmatrix} 17 & 8 \\ 1 & 1 \end{bmatrix}$

$A_{12} \cdot B_{21} = \begin{bmatrix} 19 & 39 \\ 29 & 57 \end{bmatrix}$
 $A_{22} \cdot B_{21} = \begin{bmatrix} 19 & 39 \\ 15 & 27 \end{bmatrix}$

$A_{11} \cdot B_{12} = \begin{bmatrix} 16 & 8 \\ 56 & 26 \end{bmatrix}$
 $A_{21} \cdot B_{12} = \begin{bmatrix} 24 & 14 \\ 0 & 4 \end{bmatrix}$

$A_{12} \cdot B_{22} = \begin{bmatrix} 19 & 18 \\ 29 & 26 \end{bmatrix}$
 $A_{22} \cdot B_{22} = \begin{bmatrix} 19 & 18 \\ 15 & 12 \end{bmatrix}$

Fase Combinar

$C_{11} = A_{11} \cdot B_{11} + A_{12} \cdot B_{21}$
 $C_{21} = A_{21} \cdot B_{11} + A_{22} \cdot B_{21}$

$C_{12} = A_{11} \cdot B_{12} + A_{12} \cdot B_{22}$
 $C_{22} = A_{21} \cdot B_{12} + A_{22} \cdot B_{22}$

| | |
|----|----|
| 30 | 44 |
| 67 | 74 |

| | |
|----|----|
| 35 | 26 |
| 85 | 52 |

| | |
|----|----|
| 36 | 47 |
| 16 | 28 |

| | |
|----|----|
| 43 | 32 |
| 15 | 16 |

3. Multiplicar Matrices

1
2
3

82

$$T(N) = \begin{cases} \Theta(1) & N = 1 \\ 8 \cdot T\left(\frac{N}{2}\right) + \Theta(N^2) & N > 1 \end{cases}$$

Fase Dividir

$A_{11} = \begin{bmatrix} 1 & 2 \\ 3 & 7 \end{bmatrix}$
 $A_{12} = \begin{bmatrix} 3 & 5 \\ 5 & 7 \end{bmatrix}$

$B_{11} = \begin{bmatrix} 1 & 1 \\ 5 & 2 \end{bmatrix}$
 $B_{12} = \begin{bmatrix} 0 & 4 \\ 8 & 2 \end{bmatrix}$

$A_{21} = \begin{bmatrix} 2 & 3 \\ 1 & 0 \end{bmatrix}$
 $A_{22} = \begin{bmatrix} 3 & 5 \\ 3 & 3 \end{bmatrix}$
 $B_{21} = \begin{bmatrix} 3 & 3 \\ 2 & 6 \end{bmatrix}$
 $B_{22} = \begin{bmatrix} 3 & 1 \\ 2 & 3 \end{bmatrix}$

Fase Conquistar

$A_{11} \cdot B_{11} = \begin{bmatrix} 11 & 5 \\ 38 & 17 \end{bmatrix}$
 $A_{21} \cdot B_{11} = \begin{bmatrix} 17 & 8 \\ 1 & 1 \end{bmatrix}$

$A_{12} \cdot B_{21} = \begin{bmatrix} 19 & 39 \\ 29 & 57 \end{bmatrix}$
 $A_{22} \cdot B_{21} = \begin{bmatrix} 19 & 39 \\ 15 & 27 \end{bmatrix}$

$A_{11} \cdot B_{12} = \begin{bmatrix} 16 & 8 \\ 56 & 26 \end{bmatrix}$
 $A_{21} \cdot B_{12} = \begin{bmatrix} 24 & 14 \\ 0 & 4 \end{bmatrix}$

$A_{12} \cdot B_{22} = \begin{bmatrix} 19 & 18 \\ 29 & 26 \end{bmatrix}$
 $A_{22} \cdot B_{22} = \begin{bmatrix} 19 & 18 \\ 15 & 12 \end{bmatrix}$

Fase Combinar

$C_{11} = A_{11} \cdot B_{11} + A_{12} \cdot B_{21} = \begin{bmatrix} 30 & 44 \\ 67 & 74 \end{bmatrix}$
 $C_{21} = A_{21} \cdot B_{11} + A_{22} \cdot B_{21} = \begin{bmatrix} 36 & 47 \\ 16 & 28 \end{bmatrix}$

$C_{12} = A_{11} \cdot B_{12} + A_{12} \cdot B_{22} = \begin{bmatrix} 35 & 26 \\ 85 & 52 \end{bmatrix}$
 $C_{22} = A_{21} \cdot B_{12} + A_{22} \cdot B_{22} = \begin{bmatrix} 43 & 32 \\ 15 & 16 \end{bmatrix}$

3. Multiplicar Matrices

① — ② — ③

83

$$T(N) = \begin{cases} \Theta(1) & N = 1 \\ 8 \cdot T\left(\frac{N}{2}\right) + \Theta(N^2) & N > 1 \end{cases}$$

Teorema Maestro
 (1^{er} Caso)
 $\log_2(8) = 3 > 2$

Fase Dividir

$A_{11} = \begin{bmatrix} 1 & 2 \\ 3 & 7 \end{bmatrix}$
 $A_{12} = \begin{bmatrix} 3 & 5 \\ 5 & 7 \end{bmatrix}$

$B_{11} = \begin{bmatrix} 1 & 1 \\ 5 & 2 \end{bmatrix}$
 $B_{12} = \begin{bmatrix} 0 & 4 \\ 8 & 2 \end{bmatrix}$

$A_{21} = \begin{bmatrix} 2 & 3 \\ 1 & 0 \end{bmatrix}$
 $A_{22} = \begin{bmatrix} 3 & 5 \\ 3 & 3 \end{bmatrix}$
 $B_{21} = \begin{bmatrix} 3 & 3 \\ 2 & 6 \end{bmatrix}$
 $B_{22} = \begin{bmatrix} 3 & 1 \\ 2 & 3 \end{bmatrix}$

Fase Conquistar

$A_{11} \cdot B_{11} = \begin{bmatrix} 11 & 5 \\ 38 & 17 \end{bmatrix}$
 $A_{21} \cdot B_{11} = \begin{bmatrix} 17 & 8 \\ 1 & 1 \end{bmatrix}$

$A_{12} \cdot B_{21} = \begin{bmatrix} 19 & 39 \\ 29 & 57 \end{bmatrix}$
 $A_{22} \cdot B_{21} = \begin{bmatrix} 19 & 39 \\ 15 & 27 \end{bmatrix}$

$A_{11} \cdot B_{12} = \begin{bmatrix} 16 & 8 \\ 56 & 26 \end{bmatrix}$
 $A_{21} \cdot B_{12} = \begin{bmatrix} 24 & 14 \\ 0 & 4 \end{bmatrix}$

$A_{12} \cdot B_{22} = \begin{bmatrix} 19 & 18 \\ 29 & 26 \end{bmatrix}$
 $A_{22} \cdot B_{22} = \begin{bmatrix} 19 & 18 \\ 15 & 12 \end{bmatrix}$

Fase Combinar

$C_{11} = A_{11} \cdot B_{11} + A_{12} \cdot B_{21} = \begin{bmatrix} 30 & 44 \\ 67 & 74 \end{bmatrix}$
 $C_{21} = A_{21} \cdot B_{11} + A_{22} \cdot B_{21} = \begin{bmatrix} 36 & 47 \\ 16 & 28 \end{bmatrix}$

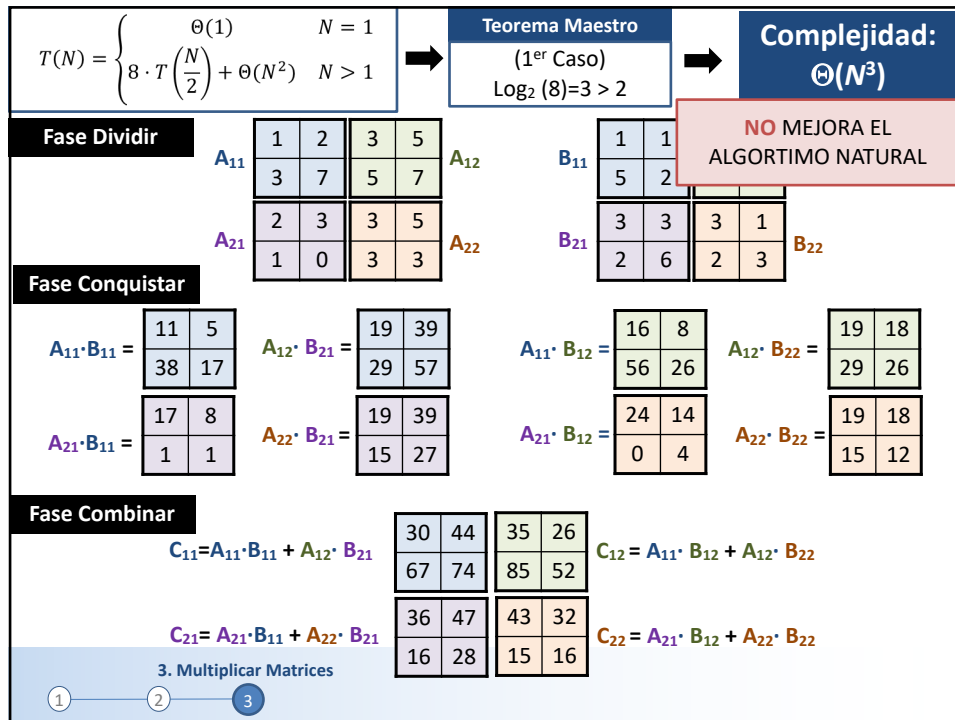
$C_{12} = A_{11} \cdot B_{12} + A_{12} \cdot B_{22} = \begin{bmatrix} 35 & 26 \\ 85 & 52 \end{bmatrix}$
 $C_{22} = A_{21} \cdot B_{12} + A_{22} \cdot B_{22} = \begin{bmatrix} 43 & 32 \\ 15 & 16 \end{bmatrix}$

3. Multiplicar Matrices

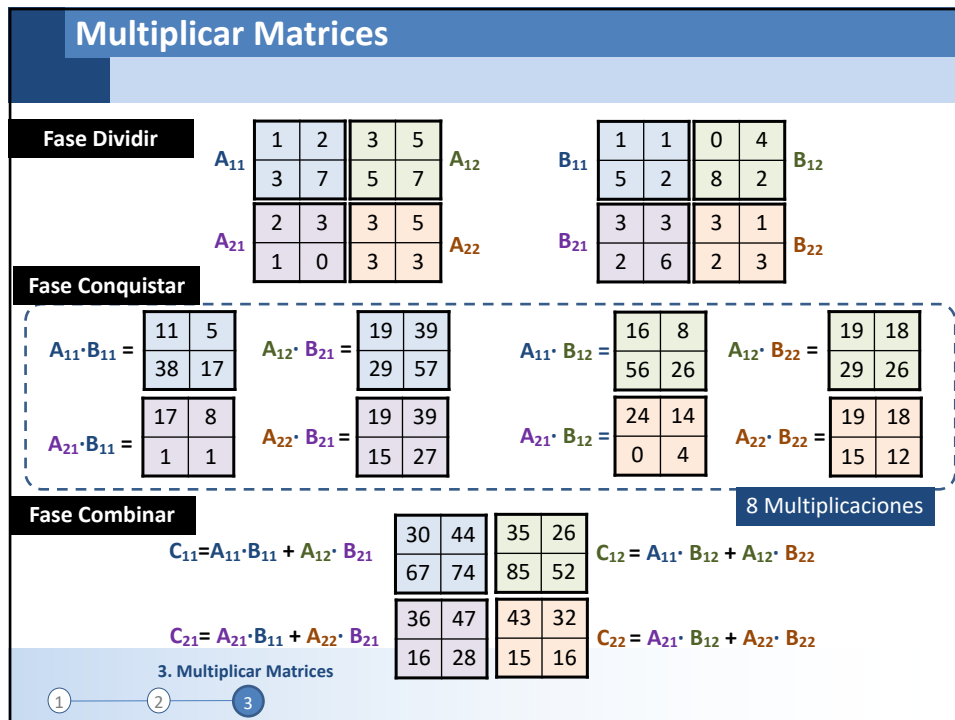
① — ② — ③

84

42



85



86

Multiplicar Matrices

Algoritmo Strassen

Fase Dividir

| | | | | |
|----------|---|---|---|---|
| A_{11} | 1 | 2 | 3 | 5 |
| | 3 | 7 | 5 | 7 |
| A_{21} | 2 | 3 | 3 | 5 |
| | 1 | 0 | 3 | 3 |

A

Fase Dividir

| | | | | |
|----------|---|---|---|---|
| B_{11} | 1 | 1 | 0 | 4 |
| | 5 | 2 | 8 | 2 |
| B_{21} | 3 | 3 | 3 | 1 |
| | 2 | 6 | 2 | 3 |

B

Fase Conquistar

$M_1 = (A_{11} + A_{22})(B_{11} + B_{22})$
 $M_4 = A_{22} \cdot (B_{21} - B_{11})$
 $M_6 = (A_{21} - A_{11})(B_{11} + B_{12})$

$M_2 = (A_{21} + A_{22}) \cdot B_{11}$
 $M_5 = (A_{11} + A_{12}) \cdot B_{22}$
 $M_7 = (A_{12} - A_{22})(B_{21} + B_{22})$

$M_3 = A_{11} \cdot (B_{12} - B_{22})$

7 Multiplicaciones

Fase Combinar

$C_{11} = M_1 + M_4 - M_5 + M_7$
 $C_{21} = M_2 + M_4$

| | | | |
|----|----|----|----|
| 30 | 44 | 35 | 26 |
| 67 | 74 | 85 | 52 |
| 36 | 47 | 43 | 32 |
| 16 | 28 | 15 | 16 |

$C_{12} = M_3 + M_5$
 $C_{22} = M_1 - M_2 + M_3 + M_6$

3. Multiplicar Matrices

1
2
3

87

Multiplicar Matrices

Actividad 7.12. Calcula la complejidad del algoritmo de Strassen

Fase Dividir

| | | | | |
|----------|---|---|---|---|
| A_{11} | 1 | 2 | 3 | 5 |
| | 3 | 7 | 5 | 7 |
| A_{21} | 2 | 3 | 3 | 5 |
| | 1 | 0 | 3 | 3 |

A

Fase Dividir

| | | | | |
|----------|---|---|---|---|
| B_{11} | 1 | 1 | 0 | 4 |
| | 5 | 2 | 8 | 2 |
| B_{21} | 3 | 3 | 3 | 1 |
| | 2 | 6 | 2 | 3 |

B

Fase Conquistar

$M_1 = (A_{11} + A_{22})(B_{11} + B_{22})$
 $M_4 = A_{22} \cdot (B_{21} - B_{11})$
 $M_6 = (A_{21} - A_{11})(B_{11} + B_{12})$

$M_2 = (A_{21} + A_{22}) \cdot B_{11}$
 $M_5 = (A_{11} + A_{12}) \cdot B_{22}$
 $M_7 = (A_{12} - A_{22})(B_{21} + B_{22})$

$M_3 = A_{11} \cdot (B_{12} - B_{22})$

Fase Combinar

$C_{11} = M_1 + M_4 - M_5 + M_7$
 $C_{21} = M_2 + M_4$

| | | | |
|----|----|----|----|
| 30 | 44 | 35 | 26 |
| 67 | 74 | 85 | 52 |
| 36 | 47 | 43 | 32 |
| 16 | 28 | 15 | 16 |

$C_{12} = M_3 + M_5$
 $C_{22} = M_1 - M_2 + M_3 + M_6$

3. Multiplicar Matrices

1
2
3

88

$$T(N) = \begin{cases} \Theta(1) & N = 1 \\ 7 \cdot T\left(\frac{N}{2}\right) + \Theta(N^2) & N > 1 \end{cases}$$

Fase Dividir

A_{11}

| | | | |
|---|---|---|---|
| 1 | 2 | 3 | 5 |
| 3 | 7 | 5 | 7 |
| 2 | 3 | 3 | 5 |
| 1 | 0 | 3 | 3 |

 A_{21}

A_{12}

| | | | |
|---|---|---|---|
| 1 | 1 | 0 | 4 |
| 5 | 2 | 8 | 2 |
| 3 | 3 | 3 | 1 |
| 2 | 6 | 2 | 3 |

 B_{11}

A_{22}

| | | | |
|---|---|---|---|
| 1 | 1 | 0 | 4 |
| 5 | 2 | 8 | 2 |
| 3 | 3 | 3 | 1 |
| 2 | 6 | 2 | 3 |

 B_{21}

B_{12}

| | | | |
|---|---|---|---|
| 1 | 1 | 0 | 4 |
| 5 | 2 | 8 | 2 |
| 3 | 3 | 3 | 1 |
| 2 | 6 | 2 | 3 |

 B_{22}

A **B**

Fase Conquistar

$M_1 = (A_{11} + A_{22})(B_{11} + B_{22})$
 $M_2 = (A_{21} + A_{22}) \cdot B_{11}$
 $M_3 = A_{11} \cdot (B_{12} - B_{22})$

$M_4 = A_{22} \cdot (B_{21} - B_{11})$
 $M_5 = (A_{11} + A_{12}) \cdot B_{22}$

$M_6 = (A_{21} - A_{11})(B_{11} + B_{12})$
 $M_7 = (A_{12} - A_{22})(B_{21} + B_{22})$

Fase Combinar

$C_{11} = M_1 + M_4 - M_5 + M_7$
 $C_{21} = M_2 + M_4$

| | | | |
|----|----|----|----|
| 30 | 44 | 35 | 26 |
| 67 | 74 | 85 | 52 |
| 36 | 47 | 43 | 32 |
| 16 | 28 | 15 | 16 |

 $C_{12} = M_3 + M_5$
 $C_{22} = M_1 - M_2 + M_3 + M_6$

3. Multiplicar Matrices

1

→

2

→

3

89

$$T(N) = \begin{cases} \Theta(1) & N = 1 \\ 7 \cdot T\left(\frac{N}{2}\right) + \Theta(N^2) & N > 1 \end{cases}$$

Teorema del Maestro

(1^{er} Caso)
 $\log_2(7) = 2.80 > 2$

Fase Dividir

A_{11}

| | | | |
|---|---|---|---|
| 1 | 2 | 3 | 5 |
| 3 | 7 | 5 | 7 |
| 2 | 3 | 3 | 5 |
| 1 | 0 | 3 | 3 |

 A_{21}

A_{12}

| | | | |
|---|---|---|---|
| 1 | 1 | 0 | 4 |
| 5 | 2 | 8 | 2 |
| 3 | 3 | 3 | 1 |
| 2 | 6 | 2 | 3 |

 B_{11}

A_{22}

| | | | |
|---|---|---|---|
| 1 | 1 | 0 | 4 |
| 5 | 2 | 8 | 2 |
| 3 | 3 | 3 | 1 |
| 2 | 6 | 2 | 3 |

 B_{21}

B_{12}

| | | | |
|---|---|---|---|
| 1 | 1 | 0 | 4 |
| 5 | 2 | 8 | 2 |
| 3 | 3 | 3 | 1 |
| 2 | 6 | 2 | 3 |

 B_{22}

A **B**

Fase Conquistar

$M_1 = (A_{11} + A_{22})(B_{11} + B_{22})$
 $M_2 = (A_{21} + A_{22}) \cdot B_{11}$
 $M_3 = A_{11} \cdot (B_{12} - B_{22})$

$M_4 = A_{22} \cdot (B_{21} - B_{11})$
 $M_5 = (A_{11} + A_{12}) \cdot B_{22}$

$M_6 = (A_{21} - A_{11})(B_{11} + B_{12})$
 $M_7 = (A_{12} - A_{22})(B_{21} + B_{22})$

Fase Combinar

$C_{11} = M_1 + M_4 - M_5 + M_7$
 $C_{21} = M_2 + M_4$

| | | | |
|----|----|----|----|
| 30 | 44 | 35 | 26 |
| 67 | 74 | 85 | 52 |
| 36 | 47 | 43 | 32 |
| 16 | 28 | 15 | 16 |

 $C_{12} = M_3 + M_5$
 $C_{22} = M_1 - M_2 + M_3 + M_6$

3. Multiplicar Matrices

1

→

2

→

3

90

45

$$T(N) = \begin{cases} \Theta(1) & N = 1 \\ 7 \cdot T\left(\frac{N}{2}\right) + \Theta(N^2) & N > 1 \end{cases}$$

Teorema del Maestro
 (1^{er} Caso)
 $\log_2(7) = 2.80 > 2$

Complejidad:
 $\Theta(N^{2.80})$

Fase Dividir

| | | | |
|---|---|---|---|
| 1 | 2 | 3 | 5 |
| 3 | 7 | 5 | 7 |
| 2 | 3 | 3 | 5 |
| 1 | 0 | 3 | 3 |

| | | | |
|---|---|---|---|
| 1 | 1 | 0 | 4 |
| 5 | 2 | 8 | 2 |
| 3 | 3 | 3 | 1 |
| 2 | 6 | 2 | 3 |

A **B**

Fase Conquistar

$M_1 = (A_{11} + A_{22})(B_{11} + B_{22})$
 $M_2 = (A_{21} + A_{22}) \cdot B_{11}$
 $M_3 = A_{11} \cdot (B_{12} - B_{22})$

$M_4 = A_{22} \cdot (B_{21} - B_{11})$
 $M_5 = (A_{11} + A_{12}) \cdot B_{22}$

$M_6 = (A_{21} - A_{11})(B_{11} + B_{12})$
 $M_7 = (A_{12} - A_{22})(B_{21} + B_{22})$

Fase Combinar

$C_{11} = M_1 + M_4 - M_5 + M_7$

 $C_{21} = M_2 + M_4$

| | | | |
|----|----|----|----|
| 30 | 44 | 35 | 26 |
| 67 | 74 | 85 | 52 |
| 36 | 47 | 43 | 32 |
| 16 | 28 | 15 | 16 |

$C_{12} = M_3 + M_5$

 $C_{22} = M_1 - M_2 + M_3 + M_6$

3. Multiplicar Matrices

1
→
2
→
3