

## Problem Set 2

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### 1 Instructions

- The solutions **should be typed**, using proper mathematical notation. We cannot accept hand-written solutions. Here's a short intro to  $\text{\LaTeX}$ .
- You should submit your work through the **class Canvas page** only. Please submit one PDF file, compiled using this  $\text{\LaTeX}$  template.
- You may not need a full page for your solutions; pagebreaks are there to help Gradescope automatically find where each problem is. Even if you do not attempt every problem, please submit this document with no fewer pages than the blank template (or Gradescope has issues with it).
- You are welcome and encouraged to collaborate with your classmates, as well as consult outside resources. You must **cite your sources in this document**. **Copying from any source is an Honor Code violation**. Furthermore, all submissions must be in your own words and reflect your understanding

**of the material.** If there is any confusion about this policy, it is your responsibility to clarify before the due date.

- Posting to **any** service including, but not limited to Chegg, Reddit, StackExchange, etc., for help on an assignment is a violation of the Honor Code.
- You **must** virtually sign the Honor Code (see Section 2). Failure to do so will result in your assignment not being graded.

## 2 Honor Code (Make Sure to Virtually Sign)

**Problem 1.**     • My submission is in my own words and reflects my understanding of the material.

- Any collaborations and external sources have been clearly cited in this document.
- I have not posted to external services including, but not limited to Chegg, Reddit, StackExchange, etc.
- I have neither copied nor provided others solutions they can copy.

*Agreed (Alex Ojemann).* I agree to the above, Alex Ojemann

□

### 3 Standard 5- BFS and DFS

#### 3.1 Problem 2

**Problem 2.** Consider a Modified Connectivity problem:

- Instance: Let  $G(V, E)$  be a simple, undirected graph. Let  $s, t \in V(G)$ .
- Decision: Given an integer  $k \geq 1$ , is there a shortest path from  $s$  to  $t$  in  $G$  that consists of  $k$  edges? Here the length/weight of a path is defined as the number of edges of the path.

Do the following. [**Note:** There are parts (a) and (b). Part (b) is on the next page.]

- (a) Design an algorithm to solve the Modified Connectivity problem. Your solution should provide enough detail that a CSCI 2270 student could reasonably be expected to implement your solution.

*Answer for Part (a).* Create a queue  $Q$  and an array of booleans  $visited$  with  $V$  elements initially all set to false.

Enqueue  $s$  into  $Q$ .

Create current vertex variable initially set to NULL.

While  $Q$  isn't empty, repeat these steps:

Dequeue a vertex from  $Q$  and set current vertex variable to that vertex.

Set  $visited$  at the current vertex to be true.

If current vertex is  $t$ , break from the while loop. There is a shortest path because you started at  $s$  and found the path to  $t$  the least number of edges away from  $s$ .

If  $visited$  is true for all nodes in the graph, break from the while loop. There is no shortest path because every node has been visited and  $t$  hasn't been found thus  $s$  and  $t$  aren't connected.

Enqueue each vertex that is connected to current vertex that has not been visited.

End while loop.

□

- (b) We say that the graph  $G$  is *connected* if for every pair of vertices  $s, t \in V(G)$ , there exists a path from  $s$  to  $t$ . Design an algorithm to determine whether  $G$  is connected. Your algorithm should only traverse the graph once - this means that you should **not** apply BFS or DFS more than once. Your solution should provide enough detail that a CSCI 2270 student could reasonably be expected to implement your solution.

*Answer for Part (b).* Create a queue  $Q$ , and a boolean array `visited` of size  $V$  and initially all set to false.

Enqueue  $s$  into  $Q$

While  $Q$  is not empty repeat these steps:

Dequeue a vertex from  $Q$  and set current vertex equal to that vertex.

Set `visited` at the current vertex to be true.

Enqueue each vertex that is connected to current vertex that has not been visited.

End while loop.

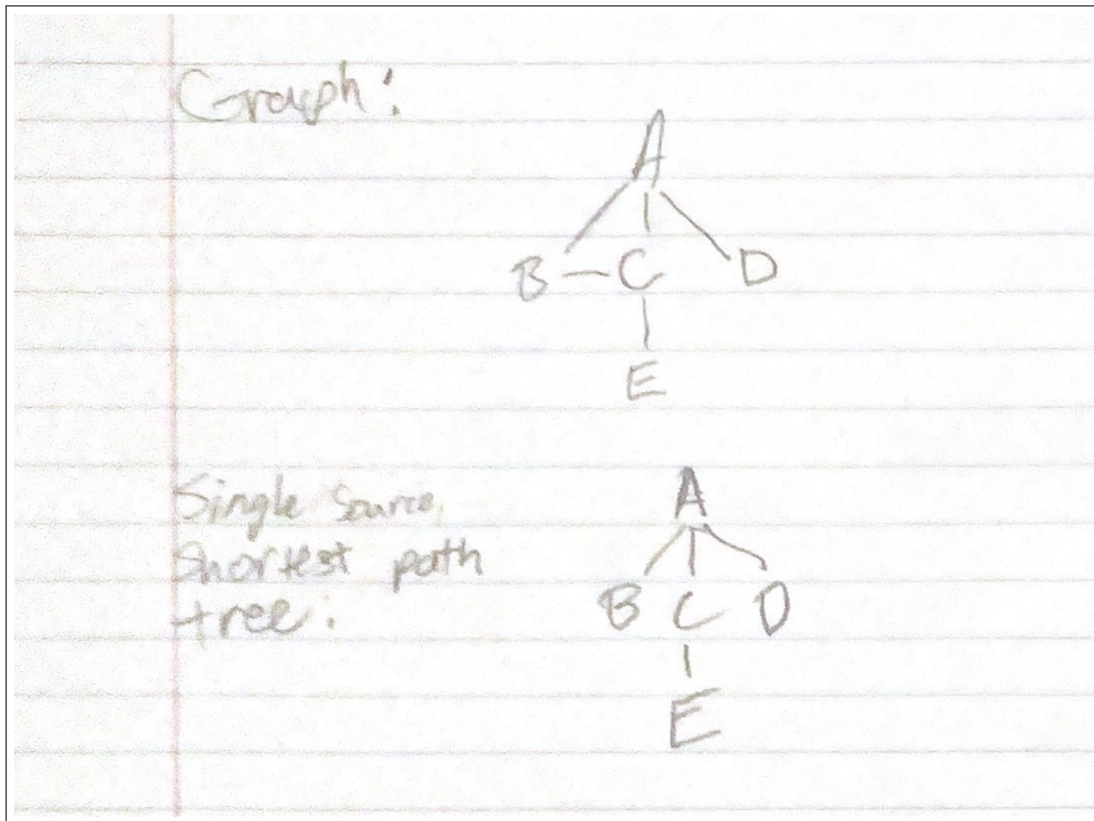
If any false value still exists in `visited`, it means  $G$  is not connected, else it is

□

### 3.2 Problem 3

**Problem 3.** Give an example of a simple, undirected, and unweighted graph  $G(V, E)$  that has a single source shortest path tree which a **depth-first traversal** will not return for any ordering of its vertices. Your answer must

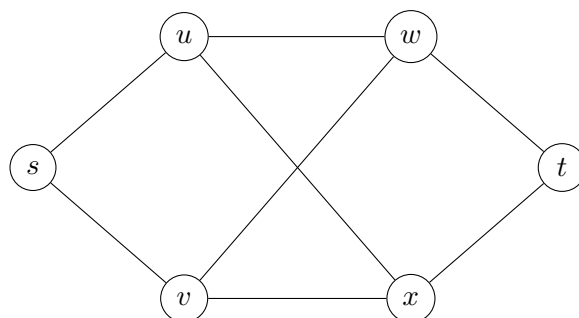
- Provide a drawing of the graph  $G$ . [Note: We have provided TikZ code below if you wish to use L<sup>A</sup>T<sub>E</sub>X to draw the graph. Alternatively, you may hand-draw  $G$  and embed it as an image below, provided that (i) your drawing is legible and (ii) we do not have to rotate our screens to grade your work.]
- Specify the single source shortest path tree  $T = (V, E_T)$  by specifying  $E_T$  and also specifying the root  $s \in V$ . [Note: You may again hand-draw this tree. If you wish, you may clearly mark the edges of  $T$  on your drawing of  $G$ . Please make it easy on the graders to identify the edges of  $T$ .]
- Include a clear explanation of why the depth-first search algorithm we discussed in class will never produce  $T$  for any orderings of the vertices.



Answer.

Depth first search will never find the single source shortest path tree for this graph because nodes  $B$  and  $C$  are connected. Thus, if you start at  $A$  and go to  $B$  you will go to  $C$  from  $B$  next which isn't the shortest path to  $C$  because  $A$  goes directly to  $C$ , and if you start at  $A$  and go to  $C$  You will go to  $B$  from  $C$  which isn't the shortest path to  $B$  because  $B$  is directly connected to  $A$ .  $\square$

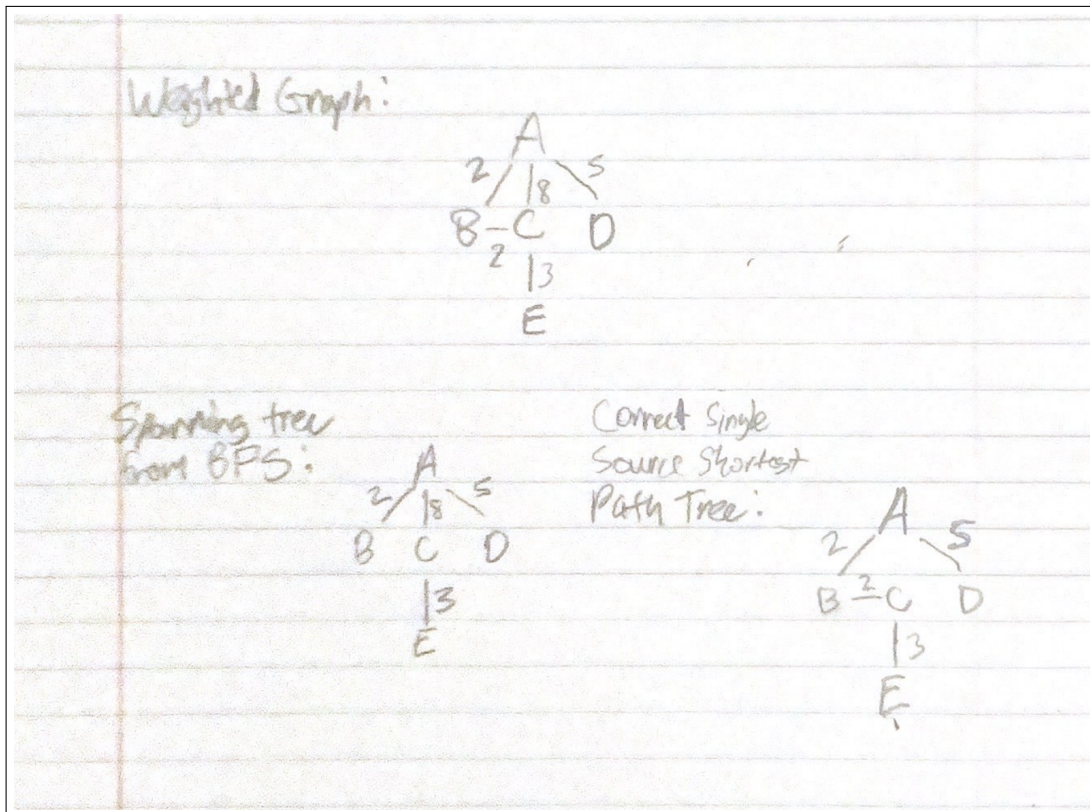
The following provides a sample of how to draw graphs with L<sup>A</sup>T<sub>E</sub>X. You may use the graph below, or you may come up with your own example.



### 3.3 Problem 4

**Problem 4.** Give an example of a simple, undirected, weighted graph such that a breadth-first traversal outputs a search-tree that is not a single source shortest path tree. (That is, BFS is not sufficiently powerful to solve the shortest-path problem on weighted graphs. This motivates Dijkstra's algorithm.) Your answer must

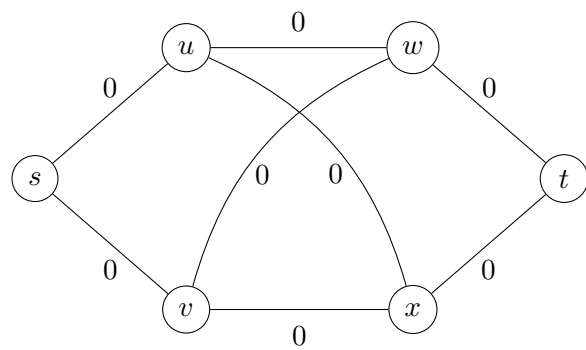
- Draw the graph  $G = (V, E, w)$  by specifying  $V$  and  $E$ , clearly labeling the edge weights. [Note: We have provided TikZ code below if you wish to use L<sup>A</sup>T<sub>E</sub>X to draw the graph. Alternatively, you may hand-draw  $G$  and embed it as an image below, provided that (i) your drawing is legible and (ii) we do not have to rotate our screens to grade your work.]
- Specify a spanning tree  $T(V, E_T)$  that is returned by BFS, but is not a single-source shortest path tree. [Note: You may again hand-draw this tree. If you wish, you may clearly mark the edges of  $T$  on your drawing of  $G$ . Please make it easy on the graders to identify the edges of  $T$ .]
- Specify a valid single-source shortest path tree  $T' = (V, E_{T'})$ . [Note: You may again hand-draw this tree. If you wish, you may clearly mark the edges of  $T$  on your drawing of  $G$ . Please make it easy on the graders to identify the edges of  $T$ .]
- Include a clear explanation of why the search-tree output by breadth-first search is not a valid single-source shortest path tree of  $G$ .



Answer.

Breadth First search doesn't return the correct single source shortest path tree for this graph because the tree it generates has a path from A to E of  $A \rightarrow C \rightarrow E$  which has a total weight of 11; however, the path from A to E with the least total weight is  $A \rightarrow B \rightarrow C \rightarrow E$  with a total weight of 7. This happens because BFS gives you the path to each vertex with the fewest number of edges but in this case a path with more edges has a smaller weight.  $\square$

The following provides a sample of how to draw graphs with edge weights using L<sup>A</sup>T<sub>E</sub>X. You may use the graph below, or you may come up with your own example.

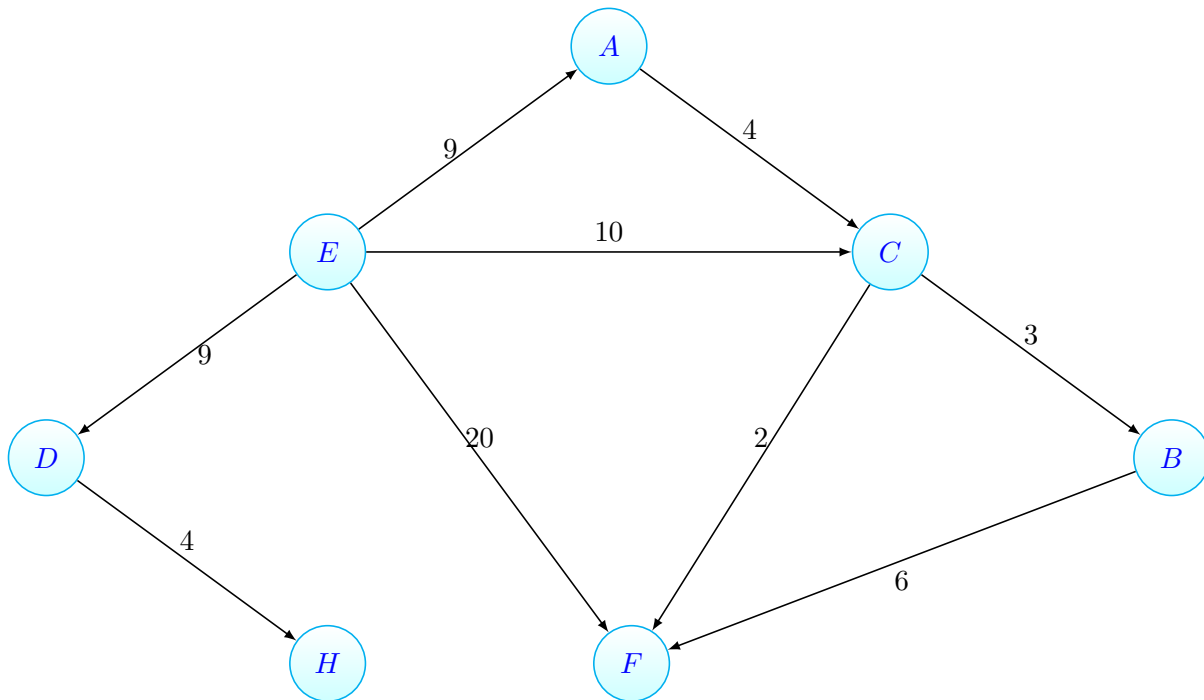


## 4 Standard 6- Dijkstra's Algorithm

### 4.1 Problem 5

**Problem 5.** Consider the weighted graph  $G(V, E, w)$  pictured below. Work through Dijkstra's algorithm on the following graph, using the source vertex  $E$ .

- Clearly include the contents of the priority queue, as well as the distance from  $E$  to each vertex at each iteration.
- If you use a table to store the distances, clearly label the keys according to the vertex names rather than numeric indices (i.e., `dist['B']` is more descriptive than `dist['1']`).
- You do **not** need to draw the graph at each iteration, though you are welcome to do so. [This may be helpful scratch work, which you do not need to include.]



*Answer.* Start:  $Q = [(E, 0)]$

Dequeue  $E$  and add  $E$ 's neighbors:  $Q = [(A, 9), (D, 9), (C, 10), (F, 20)]$

Dequeue  $A$ .  $A$ 's only neighbor is  $C$  and  $C$  already has a shorter path than the path through  $A$ . So  $Q = [(D, 9), (C, 10), (F, 20)]$

Dequeue  $D$  and add  $D$ 's neighbors:  $Q = [(C, 10), (H, 13), (F, 20)]$

Dequeue  $C$ . Update  $F$ 's shortest path because  $C$  the shortest path to  $F$  through  $C$  is shorter than  $F$ 's current shortest path. Then add  $C$ 's unexplored neighbors. So  $Q = [(F, 12), (H, 13), (B, 13)]$

Dequeue  $F$ .  $F$  has no neighbors so  $Q = [(H, 13), (B, 13)]$

Dequeue  $H$ .  $H$  has no neighbors so  $Q = [(B, 13)]$

Dequeue  $B$ .  $B$ 's only neighbor is  $F$  and  $F$  already has a shorter path than the path through  $B$ . So  $Q = []$

$Q$  is now empty. The final weights of the shortest paths from  $E$  to each node are:  $[(E, 0), (A, 9), (D, 9), (C, 10), (F, 12), (H, 13), (B, 13)]$   $\square$



## 4.2 Problem 6

**Problem 6.** You have three batteries, with capacities of 40, 25, and 16 Ah (Amp-hours), respectively. The 25 and 16-Ah batteries are fully charged (containing 25 Ah and 16 Ah, respectively), while the 40-Ah battery is empty, with 0 Ah. You have a battery transfer device which has a “source” battery position and a “target” battery position. When you place two batteries in the device, it instantaneously transfers as many Ah from the source battery to the target battery as possible. Thus, this device stops the transfer either when the source battery has no Ah remaining or when the destination battery is fully charged (whichever comes first).

But battery transfers aren’t free! The battery device is also hooked up to your phone by bluetooth, and automatically charges you a number of dollars equal to however many Ah it just transferred.

The goal in this problem is to determine whether there exists a sequence of transfers that leaves exactly 10 Ah either in the 25-Ah battery or the 16-Ah battery, and if so, how little money you can spend to get this result.

Do the following.

### 4.2.1 Problem 6(a)

- (a) Rephrase this as a graph problem. Give a precise definition of how to model this problem as a graph, and state the specific question about this graph that must be answered. [**Note:** While you are welcome to draw the graph, it is enough to provide 1-2 sentences clearly describing what the vertices are and when two vertices are adjacent. If the graph is weighted, clearly specify what the edge weights are.]

*Answer.* Set up a directed graph where the initial node is the state where the 40-Ah battery is empty while the 25-Ah and 16-Ah batteries are fully charged. Neighbors of that node are the states that are attainable via one battery transfer from the starting state and these new nodes have their attainable states adjacent to them and so on. Each edge between states has a weight that represents the cost of making that battery transfer. For the sake of this problem, if a node has 10 Ah in either the 25-Ah or 16-Ah battery it will have no neighbors because once we get to that node there is no purpose of exploring its neighbors.  $\square$

#### 4.2.2 Problem 6(b)

- (b) Clearly describe an algorithm to solve this problem. If you use an algorithm covered in class, it is enough to state that. If you modify an algorithm from class, clearly outline any modifications. Make sure to explicitly specify any parameters that need to be passed to the initial function call.

*Answer.* We can use Dijkstra's algorithm. We would provide the graph I described in 6a and the starting node , (0,25,16), as function parameters. Note that that node and all of the nodes in the graph will be represented as (Number of Ah in 40-Ah battery, Number of Ah in 25-Ah battery, Number of Ah in 16-Ah battery).  $\square$

### 4.2.3 Problem 6(c)

- (c) Apply that algorithm to the question. Report and justify your answer. Here, justification includes the sequences of vertices visited and the total cost.

*Answer.* Start:  $Q = [(0,25,16),0]$

$Q = [(16,25,0),16],[(25,0,16),25]$

$Q = [(25,0,16),25],[(16,9,16),32],[(40,1,0),40]$

$Q = [(16,9,16),32],[(40,1,0),40],[(40,0,1),40],[(25,16,0),41]$

$Q = [(40,1,0),40],[(40,0,1),40],[(25,16,0),41],[(32,9,0),48]$

$Q = [(40,0,1),40],[(25,16,0),41],[(32,9,0),48],[(24,1,16),56]$

$Q = [(25,16,0),41],[(32,9,0),48],[(24,1,16),56],[(15,25,1),65]$

$Q = [(32,9,0),48],[(24,1,16),56],[(15,25,1),65]$

$Q = [(24,1,16),56],[(32,0,9),57],[(15,25,1),65]$

$Q = [(32,0,9),57],[(15,25,1),65],[(24,17,0),72]$

$Q = [(15,25,1),65],[(24,17,0),72],[(7,25,9),82]$

$Q = [(24,17,0),72],[(15,10,16),80],[(7,25,9),82]$

$Q = [(15,10,16),80],[(7,25,9),82],[(8,17,16),88]$

$(15,10,16)$  with a path cost of 80 dollars is the lowest costing state that has 10 Ah in either the 25-Ah or the 16-Ah battery.  $\square$