

What is interpolation?

Given a set of pts $\{x_j\}_{j=0}^n$ & function evaluations $\{f(x_j)\}_{j=0}^n$ construct an approximation of the function that goes through the pts

In APPM 4600, we do interpolation w/ polynomials & trig polynomials.

Naive approach: Given $\{(x_j, f(x_j))\}_{j=0}^n$
We can construct a degree n polynomial.

$$P_n(x) = a_0 + a_1 x + \dots + a_n x^n$$

Plug in each of that a pts.

$$P_n(x_j) = a_0 + a_1 x_j + \dots + a_n x_j^n = f(x_j)$$

This is one row egn. but we have n rows.

$$\begin{bmatrix} 1 & x_0 & \cdots & x_0^n \\ 1 & x_1 & \cdots & x_1^n \\ \vdots & & & \\ 1 & x_n & & x_n^n \end{bmatrix} \begin{bmatrix} a_0 \\ \vdots \\ a_n \end{bmatrix} = \begin{bmatrix} f(x_0) \\ \vdots \\ f(x_n) \end{bmatrix}$$

(A blue bracket groups the columns from 1 to x_n^n , and an orange wavy line groups the rows from 1 to $f(x_n)$.)

This is called the Vandermonde matrix.

This is ill-conditioned for large n .

Def: We call the pts $\{x_j\}_{j=0}^n$ the interpolation nodes.

How do we know that this idea of interpolation is sane?

Thm (Weierstrass approximation Thm)

Suppose f is continuous on $[a, b]$. For $\epsilon > 0$,

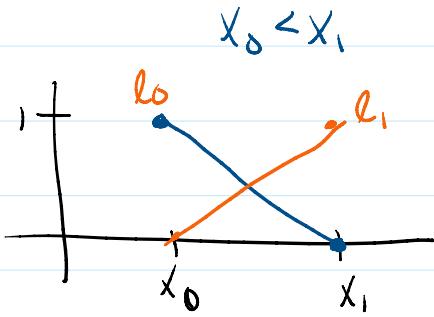
There exists a polynomial $p(x)$ st

$$|p(x) - f(x)| < \epsilon \quad \forall x \in [a, b]$$

How we build this polynomial?

let's do something other than monomials.

let's start w/ a linear approximation.



Data: $(x_0, f(x_0)), (x_1, f(x_1))$

our polynomial is going to be the sum of
2 lines: $l_0(x) \leq l_1(x)$
that satisfy $\begin{aligned} l_0(x_0) &= 1 \\ l_0(x_1) &= 0 \end{aligned}$ $\begin{aligned} l_1(x_0) &= 0 \\ l_1(x_1) &= 1 \end{aligned}$

$$l_0(x) = \frac{x - x_1}{x_0 - x_1} \quad l_1(x) = \frac{x - x_0}{x_1 - x_0}$$

we can easily write our polynomial

$$P_1(x) = f(x_0) l_0(x) + f(x_1) l_1(x)$$

How do we extend this to build higher order polynomials?

Data: $\{(x_j, f(x_j))\}_{j=0}^n$ How do build my basis of $l_j(x)$ functions?

Start w/ l_0 : We want

$l_0(x_0) = 1 \quad l_0(x_j) = 0 \quad j \neq 0$
 need $l_0(x)$ must be degree n.

$$l_0(x) = \frac{(x-x_1)(x-x_2)\cdots(x-x_n)}{(x_0-x_1)(x_0-x_2)\cdots(x_0-x_n)}$$

↑ makes roots
← scales for 1 at x_0

$$l_j(x) = \frac{(x-x_0)(x-x_1)\cdots(x-x_{j-1})(x-x_{j+1})\cdots(x-x_n)}{(x_j-x_0)(x_j-x_1)\cdots(x_j-x_{j-1})(x_j-x_{j+1})\cdots(x_j-x_n)}$$

for $j = 0, \dots, n$

$$l_j(x) = \prod_{\substack{i=0 \\ i \neq j}}^n \frac{(x-x_i)}{(x_j-x_i)}$$

Def: $l_j(x)$ is called a Lagrange Polynomial of degree n.

our interpolating polynomial

$$P_n(x) = f(x_0)l_0(x) + f(x_1)l_1(x) + \cdots + f(x_n)l_n(x)$$

$$= \sum_{j=0}^n f(x_j)l_j(x)$$

Ex. Consider the function $f(x) = e^x$

(a) construct $p_1(x)$ w/ $x_0=0$, $x_1=1$

Soln: $p_1(x) = mx + b$

$$= m(x - x_0) + f(x_0)$$

solve for m .

$$p_1(x) = f(x_0) \frac{(x - x_0)}{x_0 - x_1} + f(x_1) \frac{(x - x_0)}{(x_1 - x_0)}$$

$$= \frac{x-1}{-1} + e^{\frac{x}{1}}$$

$$\begin{matrix} l_0 & & l_1 \\ \diagdown & & \diagup \end{matrix}$$

(b) construct $p_2(x)$ w/ $x_0=0$, $x_1=\frac{1}{2}$, $x_2=1$

Soln: $p_2(x) = f(x_0) l_0(x) + f(x_1) l_1(x) + f(x_2) l_2(x)$

$$= 1 \frac{(x - \frac{1}{2})(x-1)}{(0 - \frac{1}{2})(0-1)} + e^{\frac{1}{2}} \frac{(x-0)(x-1)}{(\frac{1}{2}-0)(\frac{1}{2}-1)}$$

$$+ e^{\frac{(x-0)(x-\frac{1}{2})}{2}}$$

$$+ e \frac{(x-\delta)(x-1/2)}{(1-\delta)(1-1/2)}$$

How accurate is this?

Thm 3.3 Suppose x_0, \dots, x_n are distinct numbers in the interval $[a, b]$ & $f \in C^{(n+1)}[a, b]$

Then, for each $x \in [a, b]$, $\exists \eta$ (generally not known) between x_0, x_1, \dots, x_n in $[a, b]$ st

$$f(x) = P_n(x) + \frac{f^{(n+1)}(\eta)}{(n+1)!} (x-x_0) \dots (x-x_n)$$

where $P_n(x)$ is the Lagrange approximation of degree n through the pts $\{(x_j, f(x_j))\}_{j=0}^n$

Compare w/ Taylor: for $a \in [a-\varepsilon, a+\varepsilon]$

$$f(x) = T_n(x) + \frac{f^{(n+1)}(a)}{(n+1)!} (x-a)^{n+1}$$