

# Orders

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# 1 Preorders

We begin by studying the properties of preorders. Basically, we define a *property* as a class which is closed under isomorphism. We then define the sum operation on preorders. This will be used to create new properties from old ones.

**Definitions 1.1** (Preorder). *A preorder is a set  $M$  together with a binary relation  $\leq$  on  $M$  such that  $\leq$  is reflexive and transitive, possibly endowed with monadic predicates (labels) over some first-order monadic signature.*

**Definition 1.2** (Property of preorders). *A property  $\mathbf{P}$  of preorders is a class of preorders which is closed under isomorphism.*

**Definition 1.3.** *A property  $\mathbf{P}$  of preorders is monotone if for every preorder  $M$ ,  $M \in \mathbf{P}$  implies that every suborder of  $M$  is in  $\mathbf{P}$ .*

**Definition 1.4.** *Let  $M$  be a preorder.*

*Then  $M^*$  is the dual/reverse preorder of  $M$ .*

**Definition 1.5** (Sum of preorders). *Let  $I$  be a preorder, and let  $\{M_i\}_{i \in I}$  be a family of preorders over a disjoint signature (i.e., for every  $i \in I$ ,  $I$  and  $M_i$  have disjoint sets of labels).*

*The sum  $M = \sum_{i \in I} M_i$  is defined as follows:*

*The domain is  $M = \bigsqcup_{i \in I} M_i$  (a disjoint union).*

*Let  $\leq_i$  be the preorder on  $M_i$ .*

*Let  $x \in M_i$  and  $y \in M_j$ .*

*Then we define  $x \leq y$  iff either  $i = j$  and  $x \leq_i y$  or  $i < j$ .*

*The labels are inherited from either  $I$  or the  $M_i$ 's.*

*If  $I = 2$ , we define  $M_1 + M_2 := \sum_{i \in 2} M_i$ .*

**Lemma 1.6.** *Let  $I$  be a preorder, and let  $\{M_i\}_{i \in I}$  be a family of preorders, over a disjoint signature.*

*Then  $M = \sum_{i \in I} M_i$  is a preorder.*

*Proof.* Reflexivity is clear.

For transitivity, suppose  $x \leq y$  and  $y \leq z$ .

Suppose  $x \in M_i$ ,  $y \in M_j$ ,  $z \in M_k$ .

Then  $i \leq j$  and  $j \leq k$ , so  $i \leq k$ . If  $i = k$ , then necessarily  $i = j = k$ , and so  $x \leq_i y$  and  $y \leq_i z$ , so  $x \leq_i z$ , so  $x \leq z$ , as required.

Otherwise,  $i < k$ , and thus  $x \leq z$ , as required.  $\square$

**Definition 1.7.** *Let  $\mathbf{P}_1$  and  $\mathbf{P}_2$  be properties of preorders.*

*Then we define*

$$\mathbf{P}_1 + \mathbf{P}_2 := \{M_1 + M_2 : M_1 \in \mathbf{P}_1 \wedge M_2 \in \mathbf{P}_2\}$$

*The labels are inherited from either  $\mathbf{P}_1$  or  $\mathbf{P}_2$ .*

**Definition 1.8.** *A property  $\mathbf{P}$  of preorders is an additive property if for every preorders  $M_1$  and  $M_2$ ,  $M_1 + M_2 \in \mathbf{P}$  iff  $M_1, M_2 \in \mathbf{P}$ .*

**Definition 1.9** (Kleene plus). *Let  $\mathbf{P}$  be a property of preorders.*

*We define its Kleene plus as the smallest property of preorders  $\mathbf{P}^+$  which contains  $\mathbf{P}$  and is closed under finite sums.*

*That is,  $1^+ = \{1, 2, \dots\}$ , and  $\mathbf{P}^+ = \sum_{1^+} \mathbf{P}$ .*

**Definition 1.10** (Sum of a property over a preorder). *Let  $I$  be a preorder.*

*Let  $\mathbf{Q}$  be a property of preorders.*

*Then we define*

$$\sum_I \mathbf{Q} := \left\{ \sum_{i \in I} M_i : \forall i \in I. M_i \in \mathbf{Q} \right\}$$

**Definition 1.11** (Sum of a family of properties over a preorder). *Let  $I$  be a preorder.*

*Let  $\{\mathbf{Q}_i\}_{i \in I}$  be a family of properties of preorders over a disjoint signature.*

*Then we define*

$$\sum_{i \in I} \mathbf{Q}_i := \left\{ \sum_{i \in I} M_i : \forall i \in I. M_i \in \mathbf{Q}_i \right\}$$

*The labels are inherited from either  $I$  or the  $\mathbf{Q}_i$ 's.*

**Note 1.12.** *By the previous two definitions, if  $I$  has no labels,*

$$\sum_I \mathbf{Q} = \sum_{i \in I} \mathbf{Q}$$

**Definition 1.13** (Sum of a property over a property). *Let  $\mathbf{P}$  be a property of preorders.*

*Let  $\mathbf{Q}$  be a property of preorders over a disjoint signature.*

*Then we define,*

$$\sum_{\mathbf{P}} \mathbf{Q} := \left\{ \sum_I \mathbf{Q} : I \in \mathbf{P} \right\}$$

## 2 Linear Orders

**Definitions 2.1** (Linear order). A linear order is a preorder which is symmetric and total.

**Definition 2.2** (Property of linear orders). A property  $\mathbf{P}$  of linear orders is a class of linear orders which is closed under isomorphism.

**Definition 2.3.** Subintervals Let  $M$  be a linear order, and let  $x, y \in M$ , such that  $x \leq y$ .

Then we define the bounded subintervals  $[x, y]$ ,  $(x, y]$ ,  $[x, y)$  and  $(x, y)$  as usual.

We also define the semi-bounded subintervals  $(-\infty, x]$ ,  $[x, \infty)$ ,  $(-\infty, x)$  and  $(x, \infty)$  as usual.

We also define the unbounded subinterval  $(-\infty, \infty)$  as the whole linear order  $M$ , as usual.

A subinterval is either a bounded subinterval, a semi-bounded subinterval or the unbounded subinterval.

If  $x > y$  then we define the intervals as follows:

$$[x, y] := [y, x]$$

$$(x, y] := (y, x]$$

$$[x, y) := [y, x)$$

$$(x, y) := (y, x)$$

**Definition 2.4.** Let  $M$  be a linear order.

A set  $A \subseteq M$  is left cofinal in  $M$  if for every  $x \in M$ , there exists  $y \in A$  such that  $y < x$ .

A set  $A \subseteq M$  is right cofinal in  $M$  if for every  $x \in M$ , there exists  $y \in A$  such that  $x < y$ .

A set  $A \subseteq M$  is bi-directionally cofinal in  $M$  if it is both left and right cofinal.

**Lemma 2.5.** Let  $\mathbf{P}$  be an additive property of linear orders.

Let  $M \in \mathbf{P}$  be a linear order.

Let  $x, y \in M$  be any two points in a linear order  $M$ .

Then,  $[x, y] \in \mathbf{P}$ .

*Proof.* WLOG, suppose  $x \leq y$ .

Note that,

$$M = (-\infty, \infty) = (-\infty, x) + [x, y] + (y, \infty)$$

when  $(-\infty, x)$  and/or  $(y, \infty)$  may be empty.

Since  $\mathbf{P}$  is an additive property, we conclude that  $[x, y] \in \mathbf{P}$ .  $\square$

**Corollary 2.6.** Let  $\mathbf{P}$  be a nontrivial additive property of linear orders.

Then  $1 \in \mathbf{P}$ .

*Proof.* Let  $M \in \mathbf{P}$  be any linear order and let  $x \in M$  be any point in  $M$ .

Apply lemma 2.5 to the linear order  $M$ , and the points  $x$  and  $x$ , to conclude that  $[x, x] \equiv 1 \in \mathbf{P}$ .  $\square$

**Note 2.7.** Note that corollary 2.6 is false if we do not restrict ourselves to linear orders.

For example,  $(1 \oplus 1)^+$  is a property of preorders which is additive, but does not contain 1.

**Corollary 2.8.** Let  $\mathbf{P}$  be an additive property of linear orders.

Let  $M$  be a linear order.

Let  $x, y \in M$  be any two points in a linear order  $M$ . Then the following are equivalent:

1.  $(x, y) \in \mathbf{P}$
2.  $(x, y] \in \mathbf{P}$
3.  $[x, y) \in \mathbf{P}$
4.  $[x, y] \in \mathbf{P}$

*Proof.* This is just applying the definition of an additive property to the orders  $[x, y]$  and 1.  $\square$

**Corollary 2.9.** Let  $\mathbf{P}$  be an additive property of linear orders.

Let  $M$  be a linear order.

Let  $x, y, z \in M$  be any three points in a linear order  $M$ , such that  $[x, y] \in \mathbf{P}$  and  $[y, z] \in \mathbf{P}$ .

Then  $[x, z] \in \mathbf{P}$ .

*Proof.* If  $y \in [x, z]$ , then  $[x, z] = [x, y] + (y, z]$ , and  $(y, z] \in \mathbf{P}$  by corollary 2.8.

Otherwise, either  $x \in [y, z]$  or  $z \in [x, y]$ . WLOG, suppose  $z \in [x, y]$ .

Then  $[x, y] = [x, z] + (z, y]$ , so  $[x, z] \in \mathbf{P}$  by the fact that  $\mathbf{P}$  is additive.  $\square$

**Definitions 2.10.** Let  $\mathbf{P}$  be a property of linear orders.

We define the following properties of linear orders:

- $\mathcal{B}[\mathbf{P}]$  is the class of linear orders  $M$  such that for every  $x, y \in M$ , the bounded subinterval  $[x, y]$  is in  $\mathbf{P}$ .
- $\mathcal{L}[\mathbf{P}]$  is the class of linear orders  $M$  such that for every  $x \in M$ , the left-bounded ray  $[x, \infty)$  is in  $\mathbf{P}$ .
- $\mathcal{R}[\mathbf{P}]$  is the class of linear orders  $M$  such that for every  $x \in M$ , the right-bounded ray  $(-\infty, x]$  is in  $\mathbf{P}$ .

**Definition 2.11.** A property  $\mathbf{P}$  of linear orders is a star property if for every linear orders  $M$ , and every family  $\mathcal{F} \subseteq \mathbf{P}$  of subintervals of  $M$  such that  $J_1 \cap J_2 \neq \emptyset$  for every  $J_1, J_2 \in \mathcal{F}$ , we have that  $\bigcup \mathcal{F} \in \mathbf{P}$ .

**Lemma 2.12.** *Let  $\mathbf{P}$  be a star property.*

*Then for every linear order  $M$ , and every point  $x \in M$ , there exists a largest subinterval  $J \subseteq M$  such that  $J \in \mathbf{P}$ .*

*Equivalently, we can define a convex equivalence relation  $\sim_{\mathbf{P}}$  on  $M$  such that  $x \sim_{\mathbf{P}} y$  iff  $[x, y] \in \mathbf{P}$ .*

*That is,  $x \sim_{\mathbf{P}} y$  iff  $x$  and  $y$  are in the same largest  $\mathbf{P}$ -subinterval.*

*Proof.* Let  $J \subseteq M$  be the union of all  $\mathcal{B}[\mathbf{P}]$ -subintervals containing  $x$ . All such subintervals intersect at  $x$ .

Therefore, by the star lemma,  $J$  is in  $\mathcal{B}[\mathbf{P}]$ , and by definition  $J$  is the largest  $\mathbf{P}$ -subinterval containing  $x$ .

Thus we can define the equivalence relation  $\sim_{\mathbf{P}}$  as above.  $\square$

**Lemma 2.13** (Star Lemma). *Let  $\mathbf{P}$  be an additive property of linear orders.*

*Then the property  $\mathcal{B}[\mathbf{P}]$  is a star property.*

*Proof.* Let  $M$  be a linear order, and let  $\mathcal{F} \subseteq \mathcal{B}[\mathbf{P}]$  be a family of subintervals of  $M$ .

Let  $[x, y] \subseteq \bigcup \mathcal{F}$  be any bounded subinterval. We need to prove it is in  $\mathbf{P}$ .

Suppose  $x \in J_1$  and  $y \in J_2$  for  $J_1, J_2 \in \mathcal{F}$ .

Since  $J_1 \cap J_2 \neq \emptyset$ , we can take  $z \in J_1 \cap J_2$ .

Then  $[x, z] \subseteq J_1$  and  $[z, y] \subseteq J_2$ , and thus by the definition of  $\mathcal{B}[\mathbf{P}]$ ,  $[x, z], [z, y] \in \mathbf{P}$ . Since  $\mathbf{P}$  is additive, by corollary 2.9, we have  $[x, y] \in \mathbf{P}$ .  $\square$

**Lemma 2.14.** *Let  $\mathbf{P}$  be an additive property of linear orders.*

*Then,*

1.  $\mathcal{L}[\mathbf{P}] = \{M : M + 1 \in \mathcal{B}[\mathbf{P}]\}$
2.  $\mathcal{R}[\mathbf{P}] = \{M : 1 + M \in \mathcal{B}[\mathbf{P}]\}$
3.  $\mathbf{P} = \mathcal{L}[\mathbf{P}] \cap \mathcal{R}[\mathbf{P}] = \{M : 1 + M + 1 \in \mathcal{B}[\mathbf{P}]\}$

*Proof.* Let  $M$  be a linear order.

1. Suppose  $M + \{\infty\} \in \mathcal{B}[\mathbf{P}]$ . Then for every  $x \in M$ , we have  $[x, \infty] \in \mathbf{P}$ , and thus  $[x, \infty) \in \mathbf{P}$ . Therefore,  $M \in \mathcal{L}[\mathbf{P}]$ .

Conversely, if  $M \in \mathcal{L}[\mathbf{P}]$ , let  $x, y \in M$  be any two points in  $M + 1$ .

If  $y < \infty$ , then  $[x, y] \subseteq [x, \infty)$ . Since  $[x, \infty) \in \mathbf{P}$ , we conclude that  $[x, y] \in \mathbf{P}$ . Otherwise, if  $y = \infty$ , then  $[x, y] = [x, \infty] = [x, \infty) + \{\infty\}$ , and thus  $[x, y] \in \mathbf{P}$ .

2. The second case is dual to the first case.

3. We will show a triple inclusion.

If  $M \in \mathbf{P}$ , then by additivity,  $1 + M \in \mathbf{P}$  and  $M + 1 \in \mathbf{P}$ , and thus  $M \in \mathcal{L}[\mathbf{P}] \cap \mathcal{R}[\mathbf{P}]$ .

If  $M \in \mathcal{L}[\mathbf{P}] \cap \mathcal{R}[\mathbf{P}]$ , then by lemma 2.13,  $1 + M + 1 \in \mathcal{B}[\mathbf{P}]$ .

If  $1 + M + 1 \in \mathcal{B}[\mathbf{P}]$ , then  $M$  is a bounded subinterval of  $1 + M + 1$ , so  $M \in \mathcal{B}[\mathbf{P}]$ .

□

**Lemma 2.15.** *Let  $\mathbf{P}$  be an additive property of linear orders.*

*Then,*

$$\begin{aligned}\mathcal{B}[\mathbf{P}] &= \mathbf{P} \\ &\uplus (\mathcal{L}[\mathbf{P}] \setminus \mathcal{R}[\mathbf{P}]) \\ &\uplus (\mathcal{R}[\mathbf{P}] \setminus \mathcal{L}[\mathbf{P}]) \\ &\uplus (\mathcal{B}[\mathbf{P}] \setminus (\mathcal{L}[\mathbf{P}] \cup \mathcal{R}[\mathbf{P}]))\end{aligned}$$

*Proof.* By lemma 2.14, we conclude that  $\mathcal{L}[\mathbf{P}], \mathcal{R}[\mathbf{P}] \subseteq \mathcal{B}[\mathbf{P}]$ , since  $M + 1 \in \mathbf{P}$  and  $1 + M \in \mathbf{P}$  both imply  $1 + M + 1 \in \mathbf{P}$ .

Thus,

$$\begin{aligned}\mathcal{B}[\mathbf{P}] &= (\mathcal{L}[\mathbf{P}] \cap \mathcal{R}[\mathbf{P}]) \\ &\uplus (\mathcal{L}[\mathbf{P}] \setminus \mathcal{R}[\mathbf{P}]) \\ &\uplus (\mathcal{R}[\mathbf{P}] \setminus \mathcal{L}[\mathbf{P}]) \\ &\uplus (\mathcal{B}[\mathbf{P}] \setminus (\mathcal{L}[\mathbf{P}] \cup \mathcal{R}[\mathbf{P}]))\end{aligned}$$

Since by lemma 2.14  $\mathbf{P} = \mathcal{L}[\mathbf{P}] \cap \mathcal{R}[\mathbf{P}]$ , we conclude what we wanted to prove. □

**Lemma 2.16.** *Let  $\mathbf{P}$  be a property of linear orders.*

*Let  $M \in \mathcal{B}[\mathbf{P}]$  be a linear order.*

*Let  $x \in M$  be a non-extreme point in  $M$ .*

*Then  $[x, \infty) \in \mathcal{L}[\mathbf{P}]$  and  $(-\infty, x) \in \mathcal{R}[\mathbf{P}]$ .*

*Furthermore,  $[x, \infty) \in \mathcal{R}[\mathbf{P}]$  iff  $M \in \mathcal{R}[\mathbf{P}]$ , and  $(-\infty, x) \in \mathcal{L}[\mathbf{P}]$  iff  $M \in \mathcal{L}[\mathbf{P}]$ .*

*Proof.* This follows immediately from definitions 2.10. □

**Corollary 2.17.**

$$\mathcal{B}[\mathbf{P}] \setminus (\mathcal{L}[\mathbf{P}] \cup \mathcal{R}[\mathbf{P}]) = (\mathcal{L}[\mathbf{P}] \setminus \mathcal{R}[\mathbf{P}]) + (\mathcal{R}[\mathbf{P}] \setminus \mathcal{L}[\mathbf{P}])$$

**Lemma 2.18** (Associativity of sum). *Let  $\mathbf{P}_1, \mathbf{P}_2$  and  $\mathbf{P}_3$  be properties.*

*Then  $\sum_{\mathbf{P}_1} \sum_{\mathbf{P}_2} \mathbf{P}_3 = \sum_{\sum_{\mathbf{P}_1} \mathbf{P}_2} \mathbf{P}_3$ .*

*Proof.* It follows directly from the associativity of the sum operation on linear orders. Actually, it generalizes to any algebraic equation which holds on linear orders. □

**Lemma 2.19** (Sum and union commute). *Let  $\mathcal{P}$  be a family of properties.*

*Let  $\mathbf{Q}$  be a property.*

*Then  $\sum_{\bigcup \mathcal{P}} \mathbf{Q} = \bigcup_{\mathbf{P} \in \mathcal{P}} \sum_{\mathbf{P}} \mathbf{Q}$ .*

*Proof.* This is obvious from the definition of the sum operation. □

**Definition 2.20.** *We define  $\mathbf{CNT}$  as the class of all countable linear orders.*

**Definition 2.21.** *Let  $\gamma \geq \omega$  be a limit ordinal.*

*We define  $\Gamma_\gamma := \{\beta : \beta \subseteq \gamma^* + \gamma\}^+$ .*

*We define  $\Omega := \Gamma_\omega$ .*

**Example 2.22.**

$$\Omega = \{1, \omega, \omega^*\}^+$$

**Observation 2.23.** *Let  $\gamma \geq \omega$  be a limit ordinal.*

*Then  $\Gamma_\gamma$  is a monotone, additive property of linear orders.*



### 3 General Hausdorff Rank

**Definition 3.1.** Let  $\mathbf{Q}$  be a property of linear orders.

We define a property  $\mathbf{Q}^{<\alpha}$  for every ordinal  $\alpha$  as follows:

- For  $\alpha = 0$ ,  $\mathbf{Q}^0 = \{1\}$ .
- For  $\alpha = \gamma + 1$ ,

$$\mathbf{Q}^{<\alpha} = \sum_{\mathbf{Q}} \mathbf{Q}^{<\gamma}$$

- For  $\alpha$  a limit ordinal,

$$\mathbf{Q}^{<\alpha} = \bigcup_{\beta < \alpha} \mathbf{Q}^{<\beta}$$

**Example 3.2.** Let  $\mathbf{Q}$  be a property of linear orders.

Then  $\mathbf{Q}^1 = \mathbf{Q}$ .

**Lemma 3.3.** Let  $\mathbf{Q}$  be a property of linear orders.

Let  $\alpha, \delta$  be ordinals.

Then,

$$\mathbf{Q}^{<\alpha+\delta} = \sum_{\mathbf{Q}^{<\delta}} \mathbf{Q}^{<\alpha}$$

*Proof.* We shall prove this by induction on  $\delta \geq 0$ .

For  $\delta = 0$  we need to prove

$$\mathbf{Q}^{<\alpha} = \sum_{\mathbf{Q}^0} \mathbf{Q}^{<\alpha}.$$

Which is true by definition, since  $\mathbf{Q}^0 = \{1\}$ .

For  $\delta = \gamma + 1$ , using the induction hypothesis,

$$\begin{aligned} \mathbf{Q}^{<\alpha+\delta} &= \mathbf{Q}^{<\alpha+\gamma+1} \\ &= \sum_{\mathbf{Q}} \mathbf{Q}^{<\alpha+\gamma} \\ &= \sum_{\mathbf{Q}} \sum_{\mathbf{Q}^{<\gamma}} \mathbf{Q}^{<\alpha} \\ &= \sum_{\sum_{\mathbf{Q}} \mathbf{Q}^{<\gamma}} \mathbf{Q}^{<\alpha} \\ &= \sum_{\mathbf{Q}^{<\gamma+1}} \mathbf{Q}^{<\alpha} \\ &= \sum_{\mathbf{Q}^{<\delta}} \mathbf{Q}^{<\alpha} \end{aligned}$$

For  $\delta$  a limit ordinal, using the induction hypothesis,

$$\begin{aligned}
\mathbf{Q}^{<\alpha+\delta} &= \bigcup_{\gamma < \delta} \mathbf{Q}^{<\alpha+\gamma} \\
&= \bigcup_{\gamma < \delta} \sum_{\mathbf{Q}^{<\gamma}} \mathbf{Q}^{<\alpha} \\
&= \sum_{\bigcup_{\gamma < \delta} \mathbf{Q}^{<\gamma}} \mathbf{Q}^{<\alpha} \\
&= \sum_{\mathbf{Q}^{<\delta}} \mathbf{Q}^{<\alpha}
\end{aligned}$$

□

**Definition 3.4.** Let  $\mathbf{Q}$  be a property of linear orders.

Let  $\alpha$  be an ordinal.

We define  $\mathbf{Q}^{=\alpha} := \mathbf{Q}^{<\alpha+1} \setminus \mathbf{Q}^{<\alpha}$ .

**Definition 3.5.** Let  $\mathbf{Q}$  be a property of linear orders.

Let  $M$  be a linear order, such that  $M \in (\mathbf{Q}^{<\alpha})^+$  for some ordinal  $\alpha$ .

We define the  $\mathbf{Q}$ -Hausdorff rank of  $M$  as

$$\mathbf{hrank}_{\mathbf{Q}}(M) = \sup \left\{ \beta : M \notin (\mathbf{Q}^{<\beta})^+ \right\}$$

where the supremum is taken over all ordinals  $\beta$ . (Recall that the supremum of the empty set is defined to be 0.)

**Example 3.6.** Let  $\mathbf{Q}$  be a property of linear orders.

Let  $M$  be a linear order.

Then  $\mathbf{hrank}_{\mathbf{Q}}(M) = 0$  if and only if  $M$  is finite.

## 4 $\omega$ -Hausdorff rank

**Definitions 4.1.** Let  $\alpha > 0$  be an ordinal.

We define:

1.  $\mathcal{S}_\alpha^1 := \Omega^\alpha$
2.  $\mathcal{S}_\alpha^\omega := \mathcal{L}[\Omega^\alpha] \setminus \mathcal{R}[\Omega^\alpha]$
3.  $\mathcal{S}_\alpha^{\omega^*} := \mathcal{R}[\Omega^\alpha] \setminus \mathcal{L}[\Omega^\alpha]$
4.  $\mathcal{S}_\alpha^{\omega^* + \omega} := \mathcal{B}[\Omega^\alpha] \setminus (\mathcal{L}[\Omega^\alpha] \cup \mathcal{R}[\Omega^\alpha])$

The names will soon be justified.

**Lemma 4.2.** Let  $\alpha > 0$  be an ordinal.

Then,

1.  $\mathcal{L}[\Omega^\alpha] = \sum_\omega \Omega^\alpha$ .
2.  $\mathcal{R}[\Omega^\alpha] = \sum_{\omega^*} \Omega^\alpha$ .
3.  $\mathcal{B}[\Omega^\alpha] = \sum_{\omega^* + \omega} \Omega^\alpha$ .

*Proof.* Let us prove the first part. ( $\supseteq$ ) Let  $M \in \sum_\omega \Omega^\alpha$  be a linear order.

Let  $M = \sum_{i \in \omega} M_i$  be the decomposition of  $M$ , where  $M_i \in \Omega^\alpha$ .

Let  $x, y \in M$  be any two points in  $M$ . WLOG  $x \leq y$ .

Suppose  $x \in M_i$  and  $y \in M_j$  for  $i, j \in \omega$ .

Since  $i$  and  $j$  have a finite distance in  $\omega$ , we conclude  $[x, y] \subseteq M_i + \dots + M_j$ ,

and thus  $[x, y] \subseteq (\Omega^\alpha)^+ = \Omega^\alpha$ .

( $\subseteq$ ) Let  $M \in \mathcal{B}[\Omega^\alpha]$  be a linear order.

Since  $M$  is countable, let  $\{x_i\}_{i \in \omega} M$  be a bidirectionally cofinal  $\omega$ -sequence in  $M$ .

Then  $M = \sum_{i \in \omega} M_i$  where  $M_i = [x_i, x_{i'}]$  for  $i'$  the successor of  $i$  in  $I$ .

But  $M_i$  is a bounded interval and thus  $M_i \in \Omega^\alpha$ , so  $M \in \sum_\omega \Omega^\alpha$ .

The second part is symmetric.

The third part follows from corollary 2.17:

$$\begin{aligned} \mathcal{B}[\Omega^\alpha] &= \mathcal{R}[\Omega^\alpha] + \mathcal{L}[\Omega^\alpha] \\ &= \sum_{\omega^*} \Omega^\alpha + \sum_{\omega} \Omega^\alpha \\ &= \sum_{\omega^* + \omega} \Omega^\alpha \end{aligned}$$

□

**Lemma 4.3.** Let  $\alpha > 0$  be an ordinal.

Let  $s \in \{\omega, \omega^*, \omega^* + \omega\}$ .

Suppose that  $\alpha = \sup_{i \in s} (\alpha_i + 1)$  for ordinals  $\{\alpha_i\}_{i \in s}$ .

Then, we have the following:

$$\mathcal{S}_\alpha^s = \sum_{i \in s} \Omega^{\alpha_i}$$

**Note 4.4.** For the proof of this lemma, we actually use the fact that we work over  $\Omega$ . This proof would not have worked over  $\Gamma_\beta$  for  $\beta > \omega$ .

*Proof.* TBC. □

**Corollary 4.5.** Let  $\alpha, \delta > 0$  be ordinals.

Let  $s \in \{1, \omega, \omega^*, \omega^* + \omega\}$

Then,

$$\mathcal{S}_{\alpha+\delta}^s = \sum_{\mathcal{S}_\delta^s} \Omega^\alpha$$

*Proof.* For  $s = 1$ , it follows from lemma 3.3.

Otherwise, suppose that  $\delta = \sup_{i \in s} (\delta_i + 1)$ .

Then  $\alpha + \delta = \sup_{i \in s} (\alpha_i + \delta_i + 1)$ .

$$\mathcal{S}_{\alpha+\delta}^s = \sum_{i \in s} \mathcal{S}_{\alpha+\delta_i+1}^s = \sum_{i \in s} \sum_{\Omega^{\delta_i+1}} \Omega^\alpha = \sum_{\sum_{i \in s} \Omega^{\delta_i+1}} \Omega^\alpha = \sum_{\mathcal{S}_\delta^s} \Omega^\alpha$$

□

## 5 Type Theory

**Definition 5.1.** Let  $\mathbf{P}$  be a property of preorders.

Let  $n \in \mathbb{N}$ .

We define  $\mathbf{type}_n[\mathbf{P}]$  as the set of all  $n$ -types satisfiable in  $\mathbf{P}$ .

**Definition 5.2.** A property  $\mathbf{P}$  of preorders is computable if  $n \mapsto \mathbf{type}_n[\mathbf{P}]$  is a computable function.

**Lemma 5.3.** Let  $\mathbf{Q}$  be a property of preorders.

There exists a computable function  $f_{\mathbf{Q}} = f : \mathbb{N} \rightarrow \mathbb{N}$  such that for every  $n \in \mathbb{N}$  and every ordinal  $\alpha \geq f(n)$ ,  $\mathbf{type}_n[\mathbf{Q}^{<\alpha}] = \mathbf{type}_n[\mathbf{Q}^{f(n)}]$ .

*Proof.* Since there are only finitely many  $n$ -types, and the ordinal sequence

$$\{\mathbf{type}_n[\mathbf{Q}^{<\kappa}]\}_{\kappa}$$

is monotone, there must be some minimal  $\kappa_0 \in \omega$  where the sequence stabilizes.

This  $\kappa_0$  is computable as a function of  $n$ , because  $\mathbf{type}_n[\mathbf{Q}^{<\kappa}]$  is computable for every finite  $\kappa$ .  $\square$

**Lemma 5.4.** There exist global computable functions  $a, b : \mathbb{N} \rightarrow \mathbb{N}$  such that for all  $n, c_1, c_2 \in \mathbb{N}$  such that  $c_1, c_2 \geq a(n)$  and  $c_1 \equiv c_2 \pmod{b(n)}$ ,

$$\mathbf{type}_n[\mathbf{Q}^{=c_1}] = \mathbf{type}_n[\mathbf{Q}^{=c_2}]$$

*Proof.* Let  $n \in \mathbb{N}$ .

Since there are only finitely many sets of  $n$ -types, there exist (and can be computed) some  $a(n) \geq f(n)$ ,  $a(n) + b(n)$  such that

$$\mathbf{type}_n[\mathbf{Q}^{=a(n)}] = \mathbf{type}_n[\mathbf{Q}^{=a(n)+b(n)}]$$

By induction it follows that for all  $c \geq a(n)$ ,

$$\mathbf{type}_n[\mathbf{Q}^{=c}] = \mathbf{type}_n[\mathbf{Q}^{=c+b(n)}]$$

since  $\mathbf{Q}^{=c+1} = \sum_{\mathbf{Q}} \mathbf{Q}^{=c}$ .  $\square$

**Corollary 5.5.** Let  $n \in \mathbb{N}$ , and let  $\alpha \geq \omega$  be an ordinal.

Let  $s \in \{1, \omega, \omega^*, \omega^* + \omega\}$  be a shape.

Then there exists a computable function  $b(n)$  such that for all  $c_1, c_2 \in \mathbb{N}$  such that  $c_1, c_2 \geq a(n)$  and  $c_1 \equiv c_2 \pmod{b(n)}$ , we have

$$\mathbf{type}_n[\mathcal{S}_{c_1}^s] = \mathbf{type}_n[\mathcal{S}_{c_2}^s]$$

*Proof.* For  $s = 1$ , it follows from lemma 5.3, since  $\mathcal{S}_c^1 = \mathbf{Q}^{<c}$  and  $c \geq a(n) \geq f(n)$  for  $c \in \{c_1, c_2\}$ .

For  $s \in \{\omega, \omega^*, \omega^* + \omega\}$ , it follows easily from lemma 4.3 and lemma 5.4.  $\square$

**Lemma 5.6.** *For every  $n \in \mathbb{N}$  and for every ordinal  $\alpha \geq \omega$ ,*

$$\mathbf{type}_n[\mathbf{Q}^{=\alpha}] = \mathbf{type}_n \left[ \bigcup_{c < b(n)} \mathbf{Q}^{=a(n)+c} \right]$$

*In particular,  $\mathbf{type}_n[\mathbf{Q}^{=\alpha}]$  can be computed, and is independent of the choice of  $\alpha \geq \omega$ .*

*Proof.* By induction on  $\alpha \geq \omega$ .

Let  $\{\alpha_i\}_{i \in \omega}$  be an  $\omega$ -sequence of ordinals such that  $a(n) \leq \alpha_i$  for all  $i \in \omega$ , and  $\sup_{i \in \omega} (\alpha_i + 1) = \alpha$ .

Then  $\mathbf{Q}^{=\alpha} = \sum_{\mathbf{Q}} \bigcup_{i \in \omega} \mathbf{Q}^{=\alpha_i}$  and thus,

$$\begin{aligned} \mathbf{type}_n[\mathbf{Q}^{=\alpha}] &= \mathbf{type}_n \left[ \sum_{\mathbf{Q}} \bigcup_{i \in \omega} \mathbf{Q}^{=\alpha_i} \right] \\ &= \mathbf{type}_n \left[ \sum_{\mathbf{Q}} \bigcup_{i \in \omega} \bigcup_{c < b(n)} \mathbf{Q}^{=a(n)+c} \right] \\ &= \mathbf{type}_n \left[ \sum_{\mathbf{Q}} \bigcup_{c < b(n)} \mathbf{Q}^{=a(n)+c} \right] \\ &= \mathbf{type}_n \left[ \bigcup_{c < b(n)} \sum_{\mathbf{Q}} \mathbf{Q}^{=a(n)+c} \right] \\ &= \mathbf{type}_n \left[ \bigcup_{c < b(n)} \mathbf{Q}^{=a(n)+c+1} \right] \\ &= \mathbf{type}_n \left[ \bigcup_{c < b(n)} \mathbf{Q}^{=a(n)+c} \right] \end{aligned}$$

where the last transition is because  $\mathbf{type}_n[\mathbf{Q}^{=a(n)}] = \mathbf{type}_n[\mathbf{Q}^{=a(n)+b(n)}]$ .  $\square$

**Corollary 5.7.** *Let  $n \in \mathbb{N}$ , and let  $\alpha \geq \omega$  be an ordinal.*

*Let  $s \in \{\omega, \omega^*, \omega^* + \omega\}$  be a shape.*

$$\mathbf{type}_n[\mathcal{S}_\alpha^s] = \mathbf{type}_n \left[ \sum_s \bigcup_{c < b(n)} \Omega^{=a(n)+c} \right]$$

*In particular,  $\mathbf{type}_n[\mathcal{S}_\alpha^s]$  can be computed, and is independent of the choice of  $\alpha \geq \omega$ .*

*Proof.* There exists a sequence  $\{\alpha_i\}_{i \in s}$  such that  $a(n) \leq \alpha_i$  for all  $i \in s$ , and  $\sup_{i \in s} (\alpha_i + 1) = \alpha$ .

Then  $\mathcal{S}_\alpha^s = \sum_{i \in s} \Omega^{\alpha_i}$ , and thus,

$$\begin{aligned}
\mathbf{type}_n [\mathcal{S}_\alpha^s] &= \mathbf{type}_n \left[ \sum_{i \in s} \Omega^{\alpha_i} \right] \\
&= \mathbf{type}_n \left[ \sum_s \bigcup_{c < b(n)} \Omega^{a(n)+c} \right] \\
&= \mathbf{type}_n \left[ \bigcup_{c < b(n)} \sum_s \Omega^{a(n)+c} \right] \\
&= \mathbf{type}_n \left[ \bigcup_{c < b(n)} \mathcal{S}_{a(n)+c+1}^s \right] \\
&= \mathbf{type}_n \left[ \bigcup_{c < b(n)} \mathcal{S}_{a(n)+c}^s \right]
\end{aligned}$$

where the last transition is by corollary 5.5.

□

## 6 Decidability of the rank

**Definition 6.1.** Let  $\mathbf{Q}$  be a property of linear orders.

Let  $M$  be a linear order.

We define the predicate  $\mathbf{Int}_{\mathbf{Q}}(J)$  as true in  $M$  iff  $J$  is a  $\mathbf{Q}$ -subinterval of  $M$ .

**Lemma 6.2.** Let  $\alpha > 0$  be an ordinal.

Then predicates  $\mathbf{Int}_{\Omega \leq \alpha}$ ,  $\mathbf{Int}_{\Omega = \alpha}$  are expressible in  $\mathbf{MSO}[\mathbf{Int}_{\Omega < \alpha}]$ .

*Proof.* Obviously,

$$\mathbf{Int}_{\Omega = \alpha} \iff \mathbf{Int}_{\Omega \leq \alpha} \wedge \neg \mathbf{Int}_{\Omega < \alpha}$$

So it is enough to express  $\mathbf{Int}_{\Omega \leq \alpha}$ .

Now,  $J$  is a  $\Omega^{\leq \alpha}$ -subinterval of  $M$  iff  $J \in \sum_{\Omega} \Omega^{< \alpha}$ .

But this can be expressed in  $\mathbf{MSO}$  since it is expressible to check whether an arbitrary subset is in  $\Omega$ .  $\square$

**Definition 6.3.** Let  $\alpha > 0$  be an ordinal.

Let  $M$  be a linear order and  $x \in M$ .

We define the convex equivalence relation:

$$\sim_{\alpha} := \sim_{\mathcal{B}[\Omega^{\alpha}]}$$

and  $[x]_{\alpha} := [x]_{\mathcal{B}[\Omega^{\alpha}]}$ .

That is,  $[x]_{\alpha}$  is the largest  $\mathcal{B}[\Omega^{\alpha}]$ -subinterval containing  $x$  in  $M$ .

We define  $\sigma_{\alpha}(x)$  as the  $\alpha$ -shape of  $[x]_{\alpha}$ .

We define  $L_{\alpha}(x) = \mathbf{1}_{[x]_{\alpha} \in \mathcal{L}[\Omega^{< \alpha}]}$  and  $R_{\alpha}(x) = \mathbf{1}_{[x]_{\alpha} \in \mathcal{R}[\Omega^{< \alpha}]}$ .

**Lemma 6.4.** Let  $M$  be a linear order and  $\alpha > 0$  an ordinal.

Let  $J \subseteq M$  be an interval.

Then  $J \in \Omega^{< \alpha}$  iff it is contained in a single  $\sim_{\alpha}$ -equivalence class  $K$ , such that:

- Either  $K \in \mathcal{L}[\Omega^{< \alpha}]$  or there exists some  $x \in K$  such that  $x < J$ .
- Either  $K \in \mathcal{R}[\Omega^{< \alpha}]$  or there exists some  $x \in K$  such that  $x > J$ .

*Proof.* TBC.  $\square$

**Corollary 6.5.** Let  $\alpha > 0$  be an ordinal.

The predicate  $\mathbf{Int}_{\Omega^{< \alpha}}$  is  $\mathbf{MSO}$ -expressible over  $\mathbf{MSO}[[\cdot]_{\alpha}, L_{\alpha}, R_{\alpha}]$ .

**Theorem 6.6.** Let  $\mathbf{P}$  be a computable property of linear orders of some finite signature.

Let  $\{\mathbf{Q}_i\}_{i \in I}$  be a finite family of computable properties of linear orders over some finite signature which is disjoint from the signature of  $\mathbf{P}$ .

Then  $\bigcup_{I \in \mathbf{P}} \sum_{i \in I} \mathbf{Q}_i$  is a computable property of linear orders.



*Proof.* We will use the decomposition theorem. Let  $\tau(X_1, \dots, X_m)$  be an  $n$ -type.

Then we can compute a formula  $\psi(\xi)$  (where  $\xi$  has the type of a coloring whose range is the set of  $n$ -types) such that for any linear order  $M = \sum_{i \in I} M_i$ , and any given  $A_1, \dots, A_m \subseteq M$ ,

$$M \models \tau(A_1, \dots, A_m) \iff I \models \psi(\Xi)$$

where  $\Xi$  is the coloring assigning  $i \in I$  the  $n$ -type of  $M_i$ .

TBC. □

**Lemma 6.7.** *Let  $\alpha$  be an ordinal.*

*Let  $P$ ,  $L$  and  $R$  be first-order unary predicates.*

*Let  $C$  be the class of all countable linear orders labeled with  $P$ ,  $L$  and  $R$ , such that  $P$  represents  $\sim_\alpha$ ,  $L_\alpha(x) \iff [x]_\alpha \in \mathcal{L}[\Omega^{<\alpha}]$  and  $R_\alpha(x) \iff [x]_\alpha \in \mathcal{R}[\Omega^{<\alpha}]$ .*

*Let  $\mathbf{G}$  be the class of all countable linear orders  $I$ , labeled with a  $P$ ,  $L$  and  $R$ , such that for every pair  $i, i' \in I$  such that  $i'$  is the successor of  $i$ ,  $P(i) \neq P(i')$ , and either  $R(i) = 0$  or  $L(i') = 0$ .*

*Let  $\sigma(i) \in \{1, \omega, \omega^*, \omega^* + \omega\}$  be such that  $L(i) = 1$  iff  $\sigma(i) \in \{1, \omega\}$  and  $R(i) = 1$  iff  $\sigma(i) \in \{1, \omega^*\}$ .*

*Then,  $C = \bigcup_{I \in \mathbf{G}} \sum_{i \in I} \mathcal{S}_\alpha^{\sigma(i)}$ .*

*Proof.* ( $\subseteq$ ) Let  $M$  be a countable linear order labeled with  $P$ ,  $L$  and  $R$  as above.

Let  $I = M / \sim_\alpha$  be the quotient of  $M$  by the equivalence relation  $\sim_\alpha$ .

Then  $M = \sum_{i \in I} M_i$ , where  $\{M_i\}_{i \in I}$  are the  $\sim_\alpha$ -equivalence class of  $I$ .

Then for each  $i \in I$ ,  $M_i \in \mathcal{B}[\Omega^{<\alpha}]$ , and by definition  $\sigma(i) = \sigma_\alpha(M_i)$ .

Let  $i'$  be the successor of  $i$  in  $I$ .

Then  $P(i) \neq P(i')$  since  $P$  represents  $\sim_\alpha$ .

Furthermore, suppose  $R(i) = L(i') = 1$  holds. Then  $M_i \in \mathcal{R}[\Omega^{<\alpha}]$  and  $M_{i'} \in \mathcal{L}[\Omega^{<\alpha}]$ . so  $M_i$  and  $M_{i'}$  are the same  $\sim_\alpha$ -equivalence class of  $M$ , which is a contradiction.

Thus either  $R(i) = 0$  or  $L(i') = 0$ .

( $\supseteq$ ) Let  $M = \sum_{i \in I} M_i$  be a linear order such that  $I \in \mathbf{G}$  and  $M_i \in \mathcal{S}_\alpha^{\sigma(i)}$  for each  $i \in I$ .

In particular  $M_i \in \mathcal{B}[\Omega^{<\alpha}]$  for each  $i \in I$ , so it is contained in a single  $\sim_\alpha$ -equivalence class of  $M$ .

Suppose that there exist distinct  $j, k \in I$  such that  $j < k$ , and  $M_j, M_k$  are in the same  $\sim_\alpha$ -equivalence class.

Let  $x \in M_j$  and  $y \in M_k$ . Then  $[x, y] \in \Omega^{<\alpha}$ , and thus  $[j, k] \in \Omega^{<\alpha}$ , and in particular it is sparse.

Then there exist some  $j', k' \in I$  such that  $j < j' < k' < k$ , and  $k'$  is the successor of  $j'$  in  $I$ .

Then  $M_{j'}$  and  $M_{k'}$  are in the same  $\sim_\alpha$ -equivalence class. Thus it must be the case that  $M_{j'} \in \mathcal{R}[\Omega^{<\alpha}]$  and  $M_{k'} \in \mathcal{L}[\Omega^{<\alpha}]$ , which implies  $R(j') = L(k') = 1$ , which is a contradiction.

Thus  $\{M_i\}_{i \in I}$  are pairwise distinct  $\sim_\alpha$ -equivalence classes, and obviously the conditions holds, so  $M \in C$  and we are done.  $\square$

**Corollary 6.8.** *Let  $C$  be defined as in lemma 6.7.  
Then  $C$  is a computable property.*

**Theorem 6.9.** *Let  $\alpha > 0$  be an ordinal.*

*Satisfiability of  $\mathbf{MSO}[\mathbf{Int}_{\Omega < \alpha}]$  over all countable linear orders is decidable.*

*Proof.* First, by corollary 6.5, we can convert any formula in  $\mathbf{MSO}[\mathbf{int}_{\Omega < \alpha}]$  to an equivalent formula  $\varphi$  in  $\mathbf{MSO}[[\cdot]_\alpha, L_\alpha, R_\alpha]$ .

Now, we shall replace every occurrence of  $[\cdot]_\alpha$  in  $\varphi$  with  $P$ , every occurrence of  $L_\alpha$  with  $L$ , and every occurrence of  $R_\alpha$  with  $R$ , getting a new formula  $\varphi'$ .

Then, satisfiability of  $\varphi$  over all countable linear orders, amounts to satisfiability of  $\varphi'$  over  $C$ , which is computable by corollary 6.8.

Thus we can compute  $\mathbf{type}_n[C]$  and  $\mathbf{type}_n[\varphi']$ , and thus we can compute whether  $\varphi$  is satisfiable over all countable linear orders, by seeing if these sets intersect.  $\square$

**Theorem 6.10.** *Let  $\alpha, \delta_1, \dots, \delta_k$  be ordinals.*

*Let  $\alpha_i = \alpha + \delta_i$  for  $i = 1, \dots, k$ .*

*Let  $C$  be the class of all countable linear orders labeled with  $\pi_\alpha$  and  $\sigma_\alpha$ , and  $\pi_{\alpha_i}$  and  $\sigma_{\alpha_i}$  for  $i = 1, \dots, k$ .*

*Let  $\mathbf{G}$  be the class of all countable linear orders  $I$ , labeled with a coloring function  $\gamma$  whose range is  $\{1, \omega, \omega^*, \omega^* + \omega\}$ , such that for pair  $i, j \in I$  such that  $j$  is the successor of  $i$ , either  $\gamma(i) \in \{\omega, \omega^* + \omega\}$  or  $\gamma(j) \in \{\omega^*, \omega^* + \omega\}$ .*

*Proof.* TBC.  $\square$

## 7 Everything Better

**Theorem 7.1.** *Let  $\mathcal{C}$  be a computable property of linear orders, such that  $\mathcal{C}$  is closed under taking subintervals, projections and inverse-projections (i.e., of one of the colors), and all finite-sums and  $\mathcal{C}$ -sums.*

*Let  $\mathbf{P}_1, \dots, \mathbf{P}_k \subseteq \mathcal{C}$  be computable properties of linear orders.*

*Let  $\mathbf{MSO}[P_1, \dots, P_k]$  be monadic second order logic of order over  $\mathcal{C}$ , with  $P_1, \dots, P_k$  as monadic predicates whose semantics are:  $P_i(X)$  holds iff  $X$  is a subinterval which satisfies  $\mathbf{P}_i$ .*

*Given  $\phi$  a formula of  $\mathbf{MSO}[P_1, \dots, P_k]$  (possibly with free variables) we define*

$$\mathcal{C}_\phi = \{M \in \mathcal{C} : M \models \phi\}$$

*Then  $\mathcal{C}_\phi$  is a computable property of linear orders.*

*Proof.* By structural induction on  $\phi$ .

Suppose  $\phi$  is an atomic formula. If  $\phi$  is of the form  $X \subseteq Y$  or  $X \leq Y$ ,

$$\mathcal{C}_\phi = \{M \in \mathcal{C} : M \models \phi\}$$

and thus,

$$\mathbf{type}_n[\mathcal{C}_\phi] = \{\tau \in \mathbf{type}_n[\mathcal{C}] : \tau \models \phi\}$$

which is computable since  $\mathbf{type}_n[\mathcal{C}]$  is computable, and we can then compute whether  $\tau \models \phi$  for each  $\tau \in \mathbf{type}_n[\mathcal{C}]$ .

If  $\phi$  is of the form  $P_i(X)$ , then

$$\mathcal{C}_\phi = \{M \in \mathcal{C} : M \models P_i(X)\}$$

and thus,

$$\mathbf{type}_n[\mathcal{C}_\phi] = \mathbf{type}_n[\mathbf{P}_i]$$

which is computable since  $\mathbf{P}_i$  is computable.

If  $\phi = \neg\phi_1$ , then

$$\mathcal{C}_\phi = \mathcal{C} \setminus \mathcal{C}_{\phi_1}$$

□