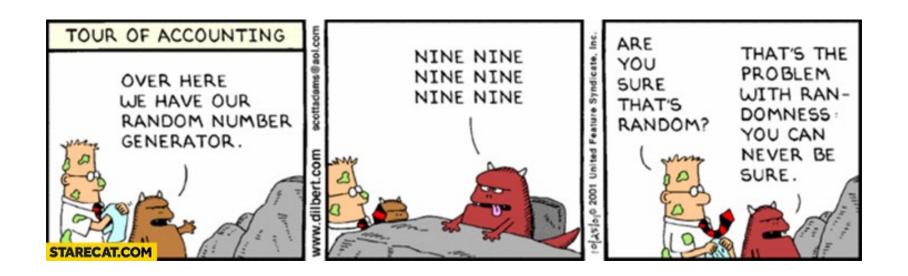
DS8104: Network Science

Class 4: Random Graphs



https://alexandergates.net

https://uvads8104.github.io/

Random graphs

Hypothesis testing: compare the observed value to a randomized null model

Goal: We want to quantify what network properties are surprising and what is expected.

Random graph models allow comparisons between properties of real-world networks and random assumptions

Random graph history

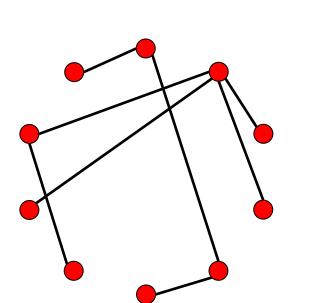
First used in: JL Moreno, HH Jennings, 1938. Statistics of social configurations. Sociometry, 1(3/4):342–374.

They manually generated (7) random graph instances to compare the observed network.

Much more detailed properties derived in:

R Solomonoff, A Rapoport, 1951. Connectivity of random nets. Bull. Math. Biophysics 13:107–117.

Pál Erdös (1913-1996)



p=1/6 N=10

< k > ~ 1.5



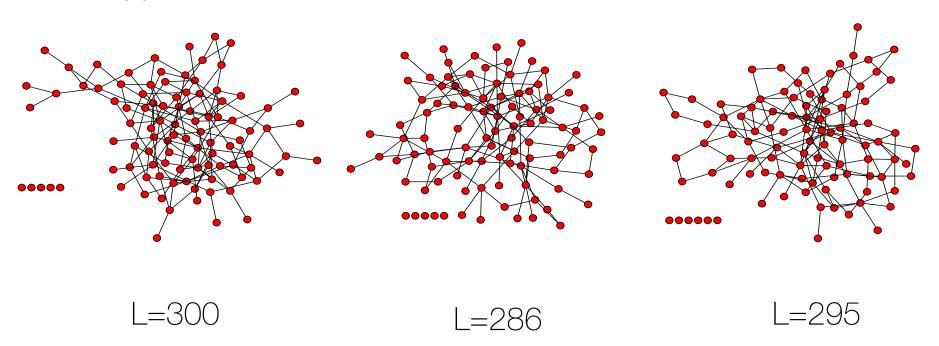


Alfréd Rényi (1921-1970)

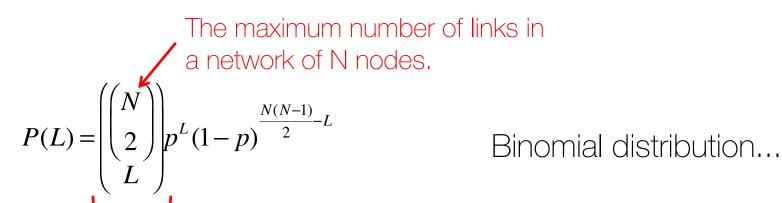
G(N, L): L links are dropped randomly without replacement between N nodes

G(N, p): Each edge has probability p to be created between N nodes

p=0.03 N=100



P(L): the probability to have exactly L links in a network of N nodes and probability p:



Number of different ways we can choose L links among all potential links.

Binomial distributions

$$P(x) = \binom{N}{x} p^{x} (1-p)^{N-x}$$

$$\langle x \rangle = Np$$

$$< x^2 > = p(1-p)N + p^2N^2$$

$$\sigma_x = (\langle k^2 \rangle - \langle k \rangle^2)^{1/2} = [p(1-p)N]^{1/2}$$

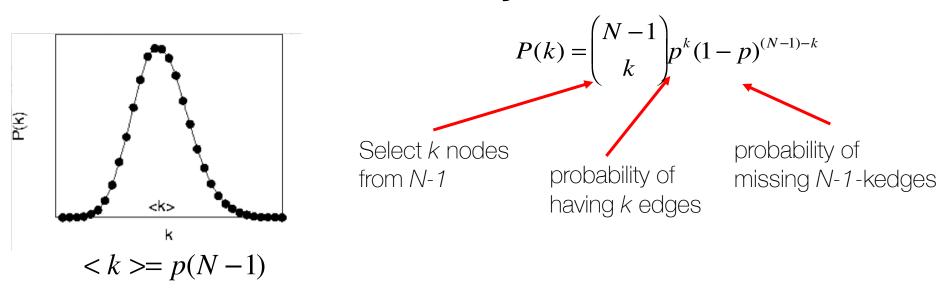
P(L): the probability to have exactly L links in a network of N nodes and probability p: ((N))

 $P(L) = \begin{pmatrix} N \\ 2 \\ L \end{pmatrix} p^{L} (1-p)^{\frac{N(N-1)}{2}-L}$

The average number of links <L> in a random graph:

$$< L> = \sum_{L=0}^{\frac{N(N-1)}{2}} LP(L) = p \frac{N(N-1)}{2}$$
 $< k> = 2L/N = p(N-1)$

$$\sigma^2 = p(1-p)\frac{N(N-1)}{2}$$



As the network size increases, the distribution becomes increasingly narrow—we are increasingly confident that the degree of a node is in the vicinity of <k>.

 $\sigma_{k}^{2} = p(1-p)(N-1)$

 $\frac{\sigma_k}{\langle k \rangle} = \left| \frac{1-p}{p} \frac{1}{(N-1)} \right|^{1/2} \approx \frac{1}{(N-1)^{1/2}}$

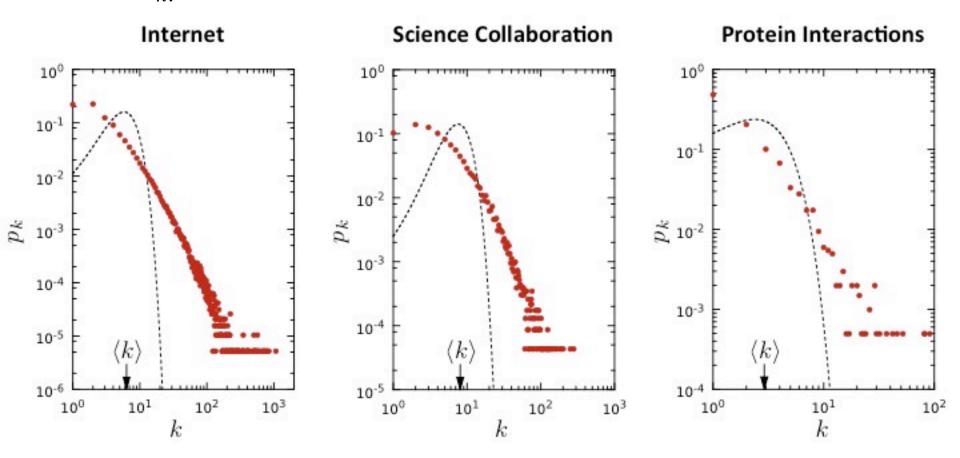
$$P(k) = \binom{N-1}{k} p^k (1-p)^{(N-1)-k}$$
 $< k >= p(N-1)$ $p = \frac{< k >}{(N-1)}$

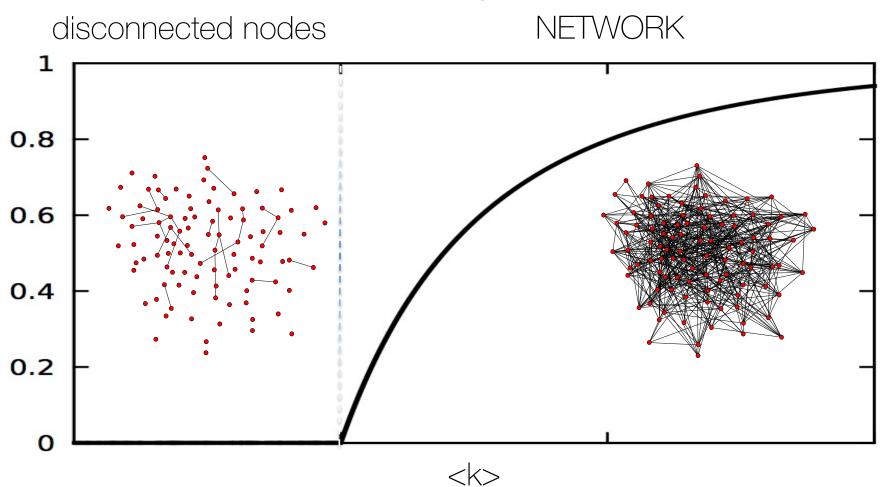
For large N and small k(p), we arrive to the Poisson distribution:

$$P(k) = e^{-\langle k \rangle} \frac{\langle k \rangle^k}{k!}$$

Real networks are not Poisson

$$P(k) = e^{-\langle k \rangle} \frac{\langle k \rangle^k}{k!}$$



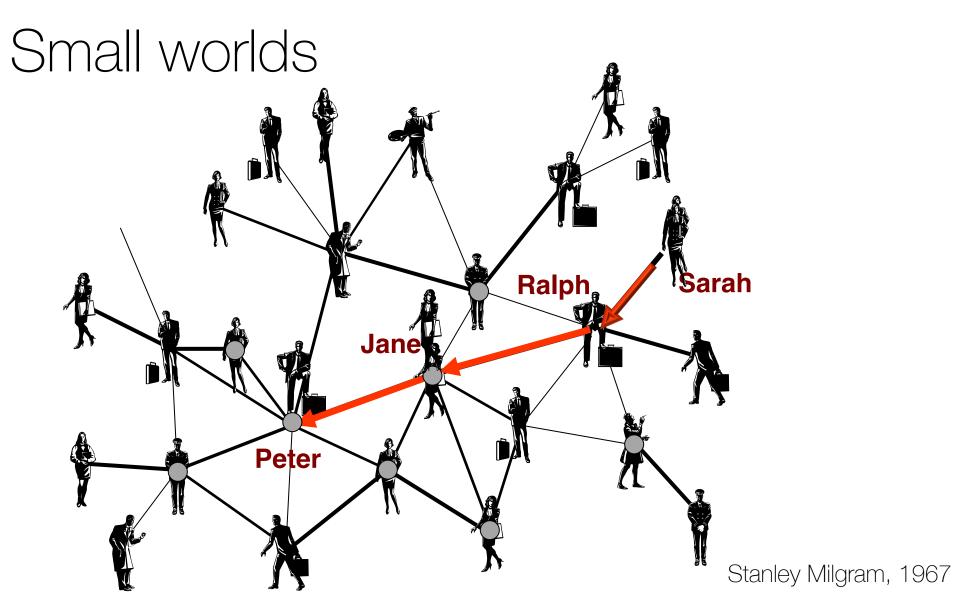


<k_c>=1 (Erdos and Renyi, 1959)

The fact that at least one link per node is necessary to have a giant component is not unexpected.

It is somewhat unexpected, however that one link is sufficient for the emergence of a giant component.

It is equally interesting that the emergence of the giant cluster is not gradual, but follows what physicists call a second order phase transition at <k>=1.



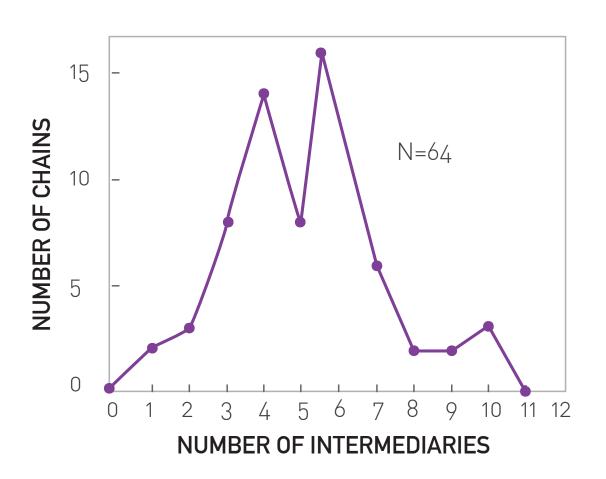
The Small-World Problem By Stanley Milgram

HOW TO TAKE PART IN THIS STUDY

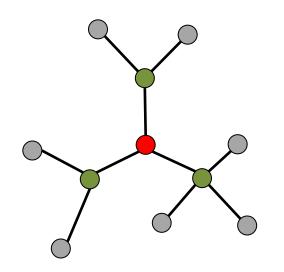
- 1. ADD YOUR NAME TO THE ROSTER AT THE BOTTOM OF THIS SHEET, so that the next person who receives this letter will know who it came from.
- 2. DETACH ONE POSTCARD. FILL IT AND RETURN IT TO HARVARD UNIVERSITY. No stamp is needed. The postcard is very important. It allows us to keep track of the progress of the folder as it moves toward the target person.
- 3. IF YOU KNOW THE TARGET PERSON ON A PERSONAL BASIS, MAIL THIS FOLDER DIRECTLY TO HIM (HER). Do this only if you have previously met the target person and know each other on a first name basis.
- 4. IF YOU DO NOT KNOW THE TARGET PERSON ON A PERSONAL BASIS, DO NOT TRY TO CONTACT HIM DIRECTLY. INSTEAD, MAIL THIS FOLDER (POST CARDS AND ALL) TO A PERSONAL ACQUAINTANCE WHO IS MORE LIKELY THAN YOU TO KNOW THE TARGET PERSON. You may send the folder to a friend, relative or acquaintance, but it must be someone you know on a first name basis.

The Small-World Problem

By Stanley Milgram



Random graphs tend to have a tree-like topology with almost constant node degrees.



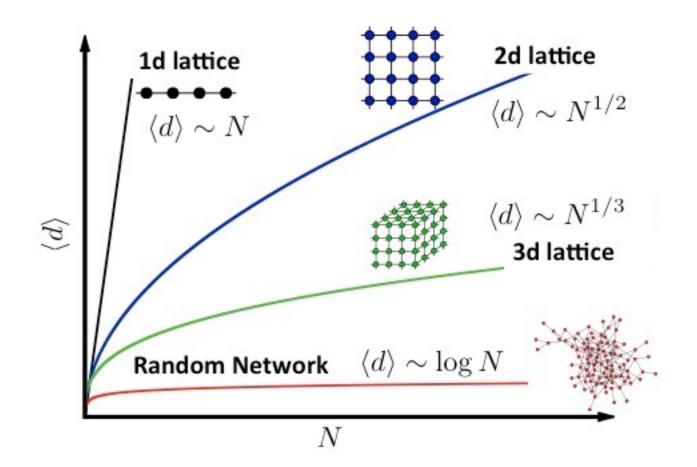
$$<$$
k $>$ nodes at distance one (d=1)

$$<$$
k $>$ ² nodes at distance two (d=2)

<k> dmax nodes at distance d_{max}

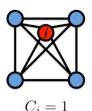
$$N = 1 + \langle k \rangle + \langle k \rangle^2 + \dots + \langle k \rangle^{d_{\text{max}}} = \frac{\langle k \rangle^{d_{\text{max}} + 1} - 1}{\langle k \rangle - 1} \approx \langle k \rangle^{d_{\text{max}}} \qquad \Longrightarrow \qquad d_{\text{max}} = \frac{\log N}{\log \langle k \rangle}$$

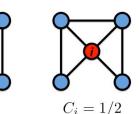
Why are small worlds surprising?

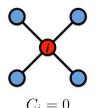


Clustering coefficient

$$C_i \equiv \frac{2 < L_i >}{k_i (k_i - 1)}$$





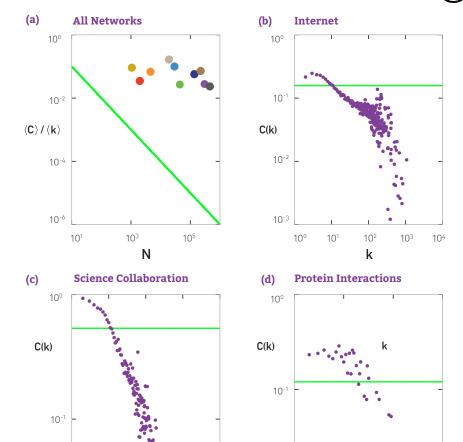


Edges are independent and have the same probability p

$$\langle L_i \rangle \cong p \frac{k_i(k_i - 1)}{2}$$
 \Longrightarrow $C_i = \frac{2\langle L_i \rangle}{k_i(k_i - 1)} = p = \frac{\langle k \rangle}{N}.$

The clustering coefficient of a random graph is small (we don't expect many triangles!)

Clustering coefficient



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10¹

k

10°

 10^{1}

 10^{2}

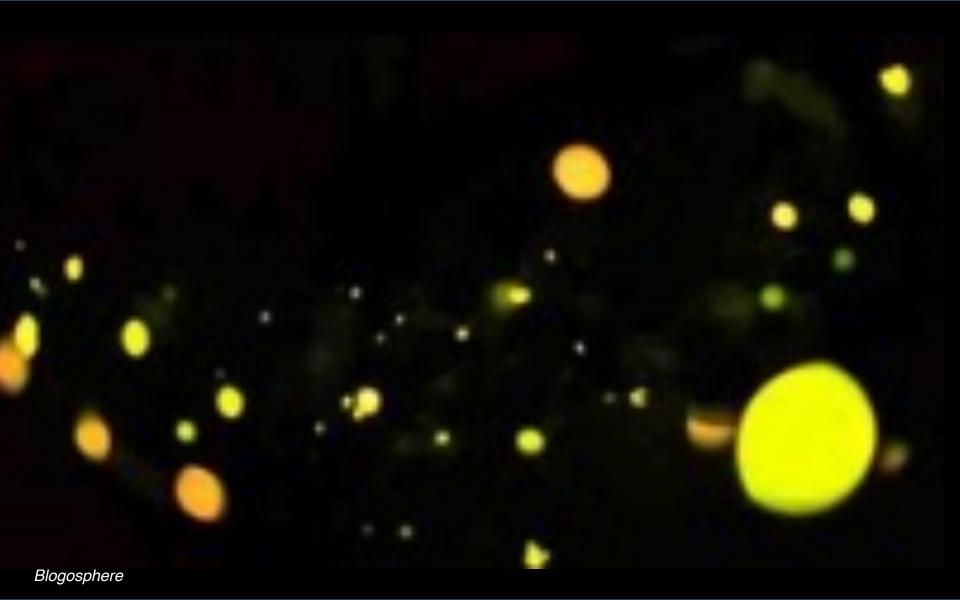
k

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$$C_i = \frac{2\langle L_i \rangle}{k_i(k_i - 1)} = p = \frac{\langle k \rangle}{N}.$$

C decreases with the system size N.

C is independent of a node's degree k.



Collective dynamics of 'small-world' networks

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