

# Fundamentals of Information & Network Security

## ECE 471/571



Lecture #6: Early Ciphers and Cryptanalysis

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# Vigenere Cipher

- Best known polyalphabetic cipher. Key is m-length vector. Use different monoalphabetic substitutions as one proceeds through the plaintext message

$$y = e_K(x_1, x_2, \dots, x_m) = (x_1 + K_1, x_2 + K_2, \dots, x_m + K_m) \bmod 26,$$
$$d_K(y_1, y_2, \dots, y_m) = (y_1 - K_1, y_2 - K_2, \dots, y_m - K_m) \bmod 26.$$

**Example:** key: “deceptive”; plaintext: “we are discovered save yourself”, obtain the ciphertext:

|             |                                    |
|-------------|------------------------------------|
| key:        | <i>deceptivedeceptivedeceptive</i> |
| plaintext:  | <i>wearediscoveredsaveyourself</i> |
| ciphertext: | <i>ZICVTWQNGRZGVTWAVZHCQYGLMGJ</i> |

# Hill Cipher

- Takes  $m$  successive plaintext letters and substitutes for them  $m$  ciphertext letters
- The substitution is determined by  $m$  linear equations in which each character is assigned a numerical value

$$\mathcal{P} = \mathcal{C} = (\mathbb{Z}_{26})^m$$

$K$  must be invertible

$$e_K(x) = \underline{x}K, \\ d_K(y) = \underline{y}K^{-1}.$$

- For example:

plaintext: test

$$K = \begin{pmatrix} 2 & 3 \\ 5 & 7 \end{pmatrix}$$

- Question: Is the Hill cipher encryption an injective function?

# Stream Ciphers

- Generate a keystream  $z = z_1 z_2 \cdot \cdot \cdot$  and encrypt each character  $x_i$  of the plaintext with a different key  $z_i$ .

$$y = y_1 y_2 \cdots = e_{z_1}(x_1) e_{z_2}(x_2) \cdots$$

- Shift and Vigenere ciphers can be viewed as special cases of stream cipher

- Vernam Cipher

- One-Time Pad
  - Key is as long as the message

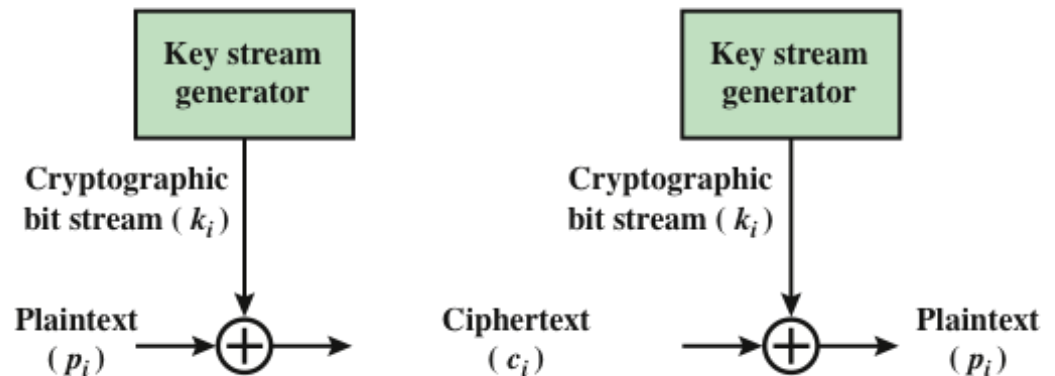


Figure 3.7 Vernam Cipher

# Permutation Ciphers

- Permutation (transposition)
  - Rearranging of the symbols, a.k.a Permutation
  - **Diffusion**: widely spreading the information from the message or the key across the ciphertext
  - Break the establish pattern
- Formal definition:  $\mathcal{P} = \mathcal{C} = (\mathbb{Z}_{26})^m$ 

$$y = e_{\pi}(x_1, \dots, x_m) = (x_{\pi(1)}, \dots, x_{\pi(m)})$$

$$x = d_{\pi}(y_1, \dots, y_m) = (y_{\pi^{-1}(1)}, \dots, y_{\pi^{-1}(m)}).$$

|          |   |   |   |   |   |   |
|----------|---|---|---|---|---|---|
| $j$      | 1 | 2 | 3 | 4 | 5 | 6 |
| $\pi(j)$ | 3 | 5 | 1 | 6 | 4 | 2 |

|               |   |   |   |   |   |   |
|---------------|---|---|---|---|---|---|
| $j$           | 1 | 2 | 3 | 4 | 5 | 6 |
| $\pi^{-1}(j)$ | 3 | 6 | 1 | 5 | 2 | 4 |

Example: plaintext is: followashore

# Permutation (cont'd)

- Rail Fence Cipher

- Plaintext is written down as a sequence of diagonals and then read off as a sequence of rows

Plaintext:    m e m a t r h t g p r y  
                 e t e f e t e o a a t

Ciphertext:  M E M A T R H T G P R Y E T E F E T E O A A T

- Columnar Transpositions

- A rearrangement of the characters of the plain text into columns

- Example

Plaintext: We are discovered. Flee at once.

Ciphertext: WIREEESEAACDTROFOEVLNDEEC

|   |   |   |   |   |   |
|---|---|---|---|---|---|
| W | E | A | R | E | D |
| I | S | C | O | V | E |
| R | E | D | F | L | E |
| E | A | T | O | N | C |
| E |   |   |   |   |   |

# Cryptanalysis

- The goal of cryptanalysis?



The Enigma machine

Question: should the crypto algorithms be published or not?

# Kerckhoffs' principle



Auguste Kerckhoffs (1883):

*The enemy knows the system*

The cipher should remain secure even if **the adversary knows the specification of the cipher.**

The only thing that is **secret** is a

short key **k**

that is **usually chosen uniformly at random**



# Kerckhoffs' Principle – the Motivation

1. In commercial products it is unrealistic to assume that the design details remain secret (**reverse-engineering!**)
2. Short keys are easier to **protect**, **generate** and **replaced**.
3. The design details can be discussed and **analyzed in public**.