

Fundamentals of Information & Network Security

ECE 471/571



Lecture #18: Hash and Message Authentication Code

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Message Authentication Code

- $MD(m)$?
- $MD(K_{AB} || m)$: only the one who knows the secret can compute/verify
- Problem?

The Problem with keyed hash $h(\text{Key} \parallel m)$

- A feature of message digest algorithms
 - In order to compute the message digest through chunk n , all that you need to know is the message digest through chunk $n-1$, plus the chunk n of the padded message.
- An Attack
 - Someone gets m , and $\text{digest}(\text{Key} \parallel m)$
 - He first pads m according the used hash function, and then adds another message M at the end. The result is $m \parallel \text{pad} \parallel M$.
 - $\text{digest}(\text{Key} \parallel m \parallel \text{pad} \parallel M)$ can be calculated from $\text{digest}(\text{Key} \parallel m)$, which is the intermediate digest.

Solutions

- Use $h(m \parallel \text{Key})$
- HMAC

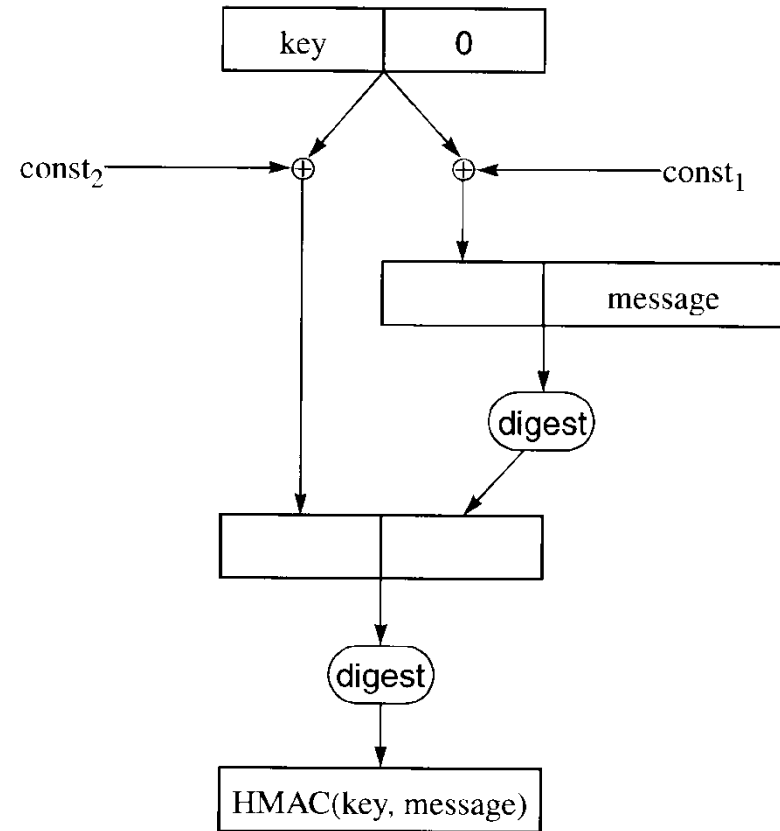


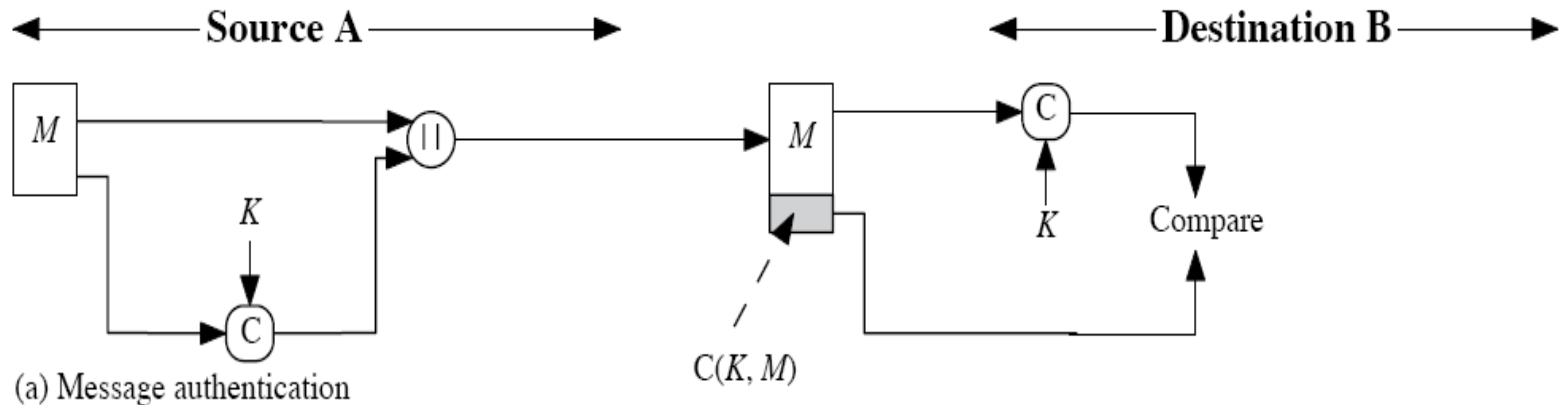
Figure 5-10. HMAC

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$$\text{HMAC}_K(x) = \text{SHA-1}((K \oplus \text{opad}) \parallel \text{SHA-1}((K \oplus \text{ipad}) \parallel x)).$$

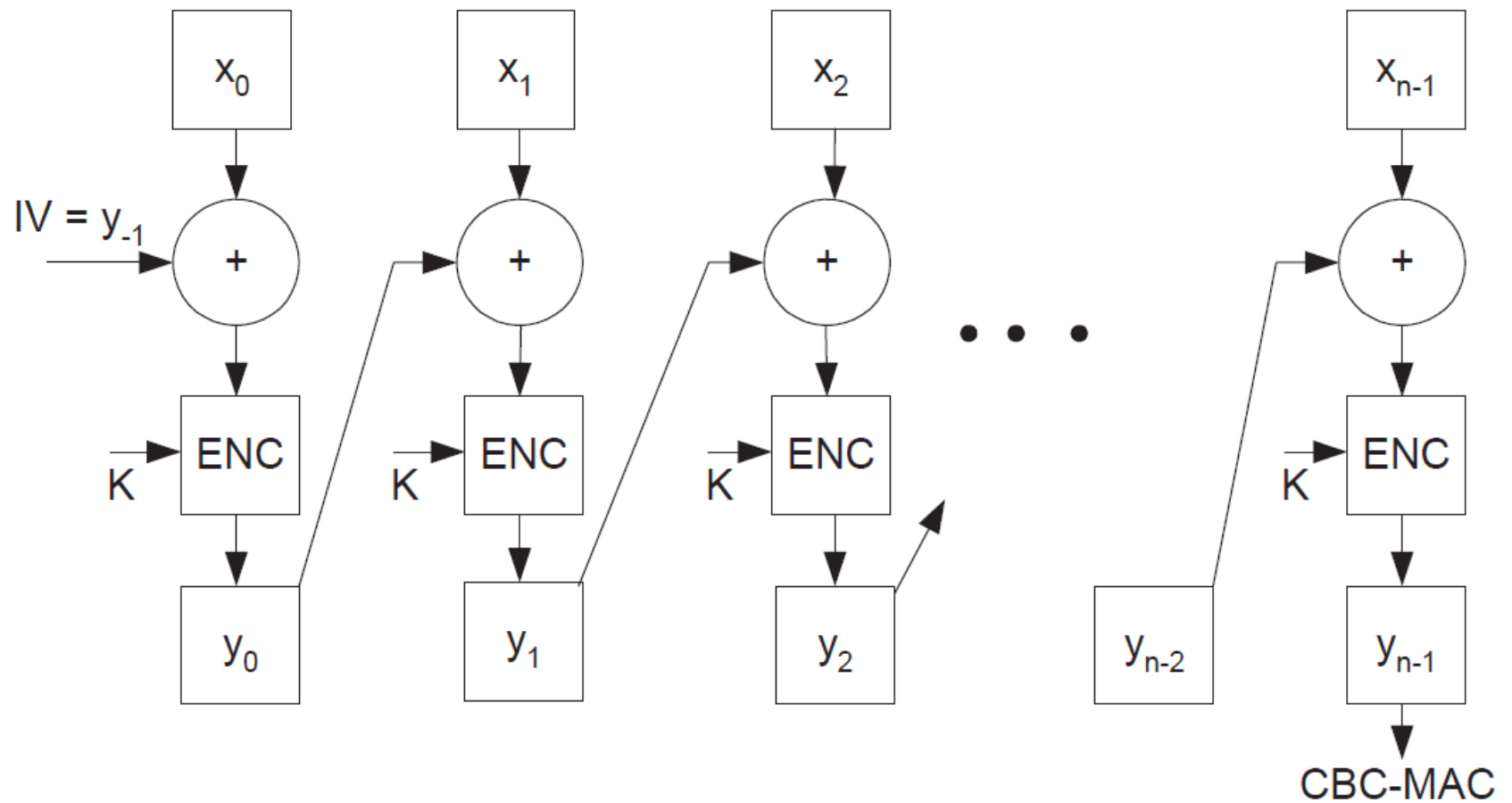
Properties of MAC

Two properties a MAC should have

- Key non-recovery: it is hard to compute the secret key from observed (message, MAC value) pairs.
- Computation resistance: even if many (message, MAC value) pairs are observed, it is hard to compute an as yet unobserved (message, MAC value) pair that verifies correctly.

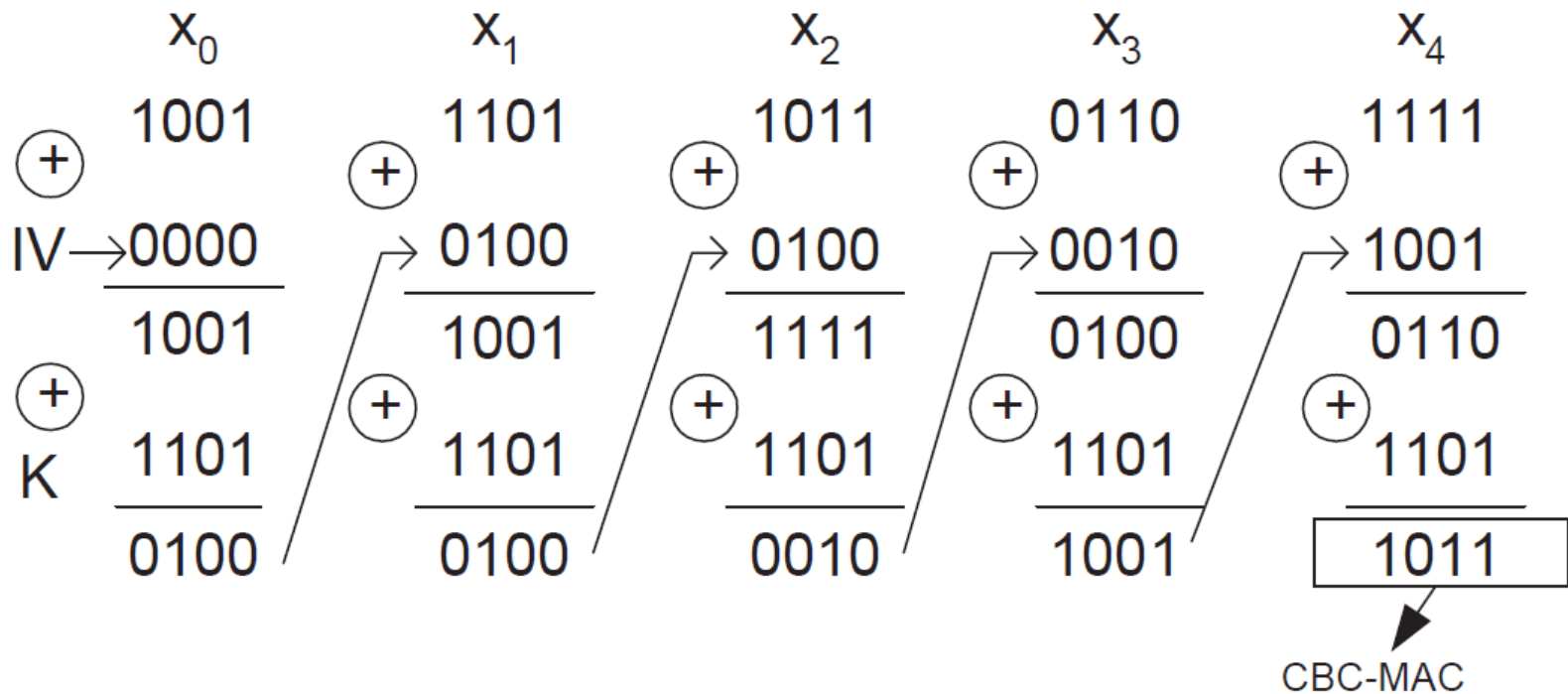


CBC-MAC



Example

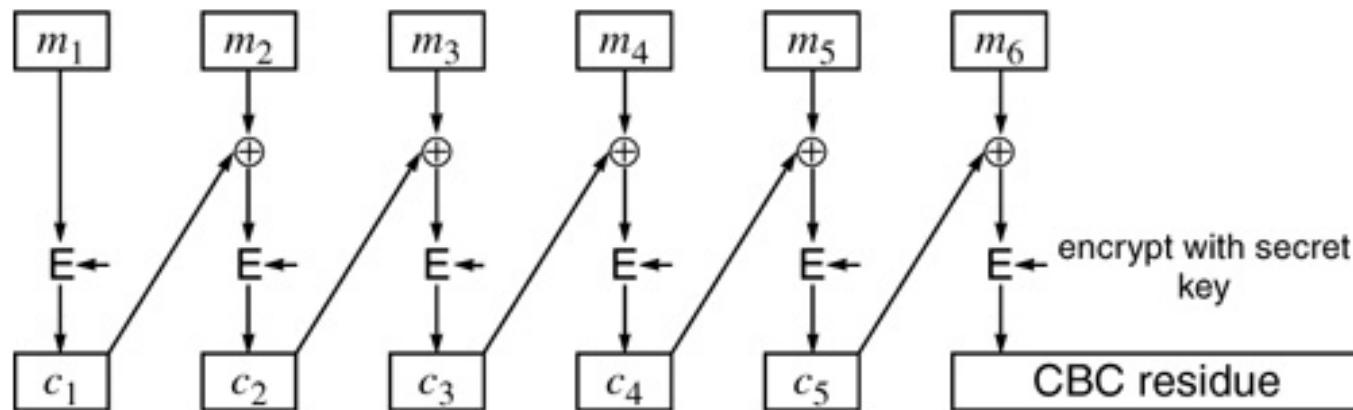
- How does Bob know the message length?
- Is this a good CBC-MAC?



Alice sends to Bob: 10011101101101101111**1011**

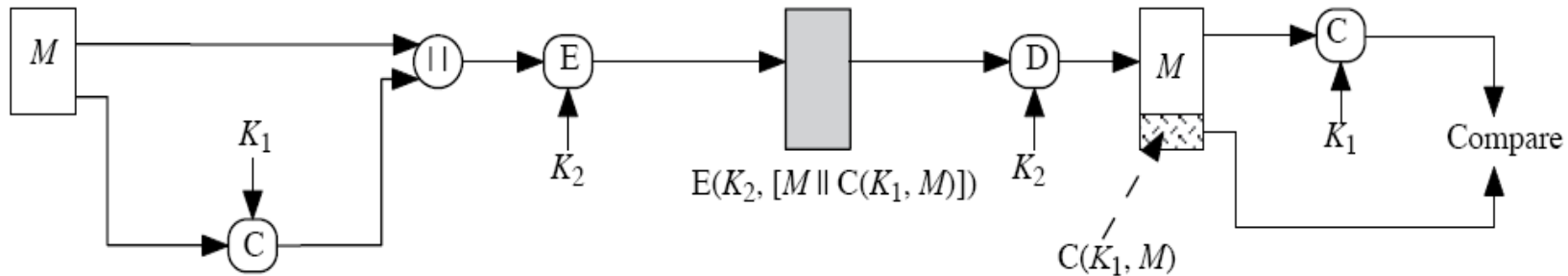
MACs based on Block Ciphers

- Protect against undetected modifications
- **Plaintext** + CBC residue (when message not secret)

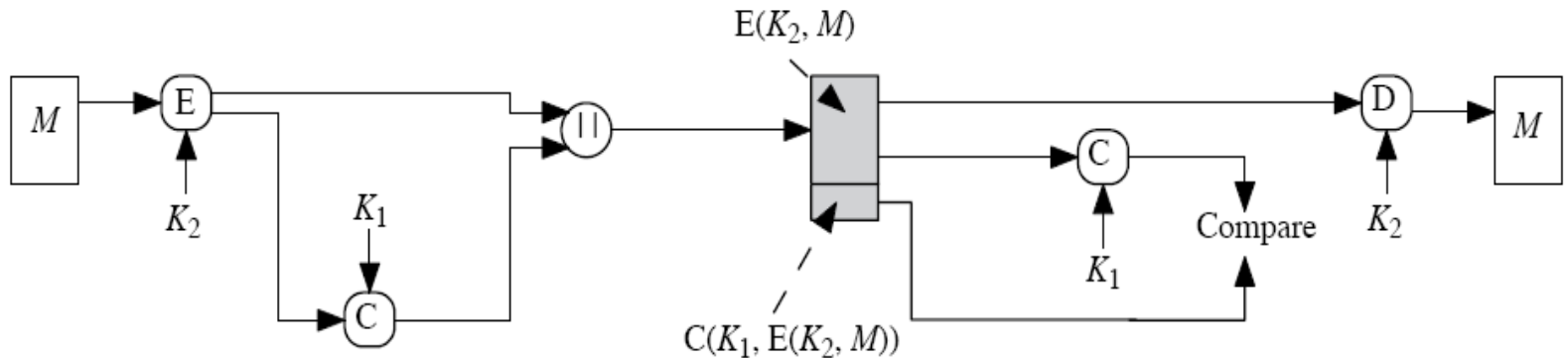


Authenticated Encryption

(a) Hash then Encrypt (not very secure)



(b) Message authentication and confidentiality; authentication tied to plaintext



(c) Message authentication and confidentiality; authentication tied to ciphertext

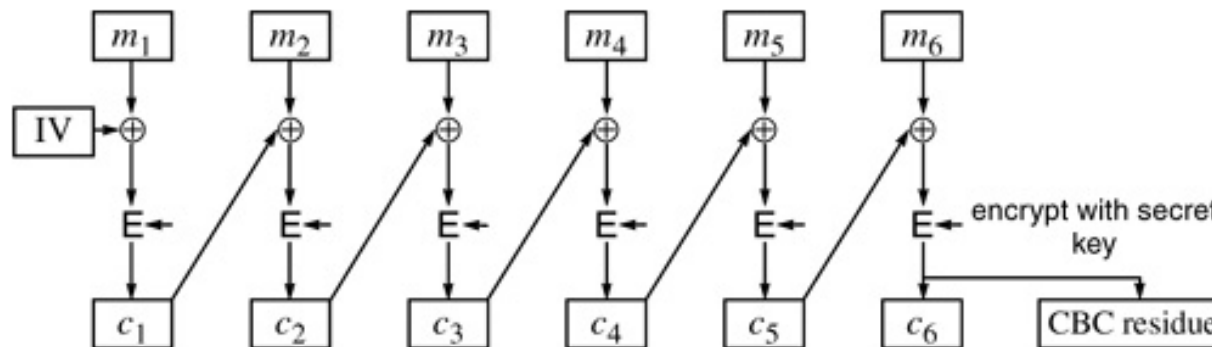
(d) Authentication (MAC) + Encryption (separately)

Authenticated Encryption based on Block Cipher Modes of Encryption

- Example:
 - Privacy: CBC encryption
 - Integrity: CBC residue
- Ciphertext + CBC residue ?
- Encrypt {plaintext + CBC residue} ?
- Encrypt {plaintext + CRC} ?

Ciphertext + CBC Residue

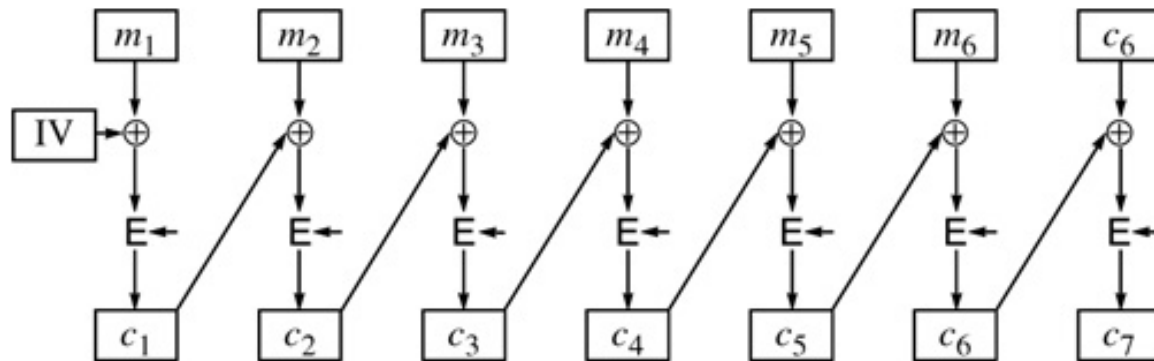
Figure 4-12. Cipher Block Chaining Encryption plus CBC Residue



- Problem?

Encrypt {plaintext + CBC residue}

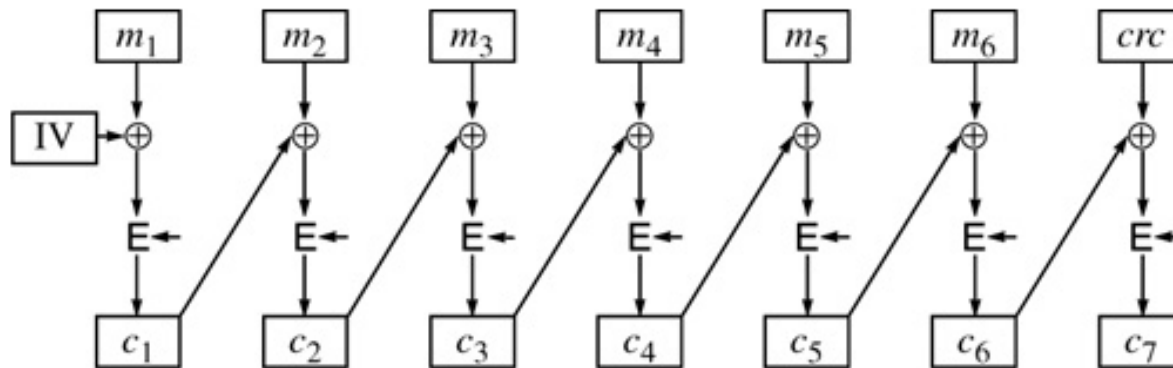
Figure 4-13. Cipher Block Chaining Encryption of Message with CBC Residue



- Problem?

Encrypt {plaintext + CRC}

Figure 4-14. Cipher Block Chaining Encryption of Message with CRC



- Longer CRC maybe Okay

Authenticated Encryption / Confidentiality and Integrity: The Do's

- Confidentiality: CBC encryption + Integrity: CBC residue, but with different keys
- CBC + weak cryptographic checksum
- CBC + cryptographic hash: keyed hash preferred
- CCM: Counter Mode + CMAC (A more secure version of CBC-MAC)
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