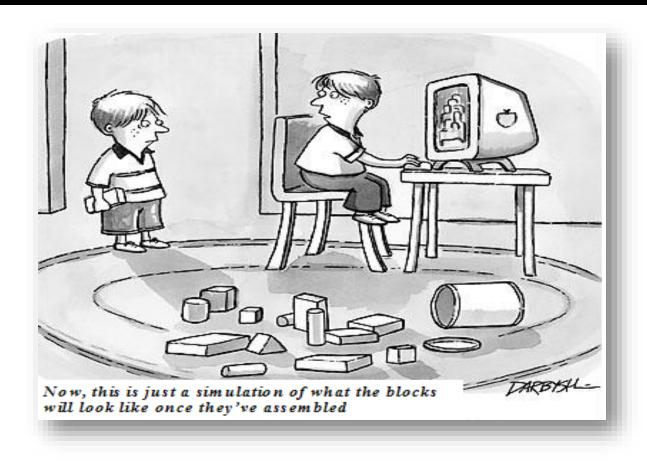
#### ECE569 Module 25



• Matrix Multiplication – Boundary Conditions

1

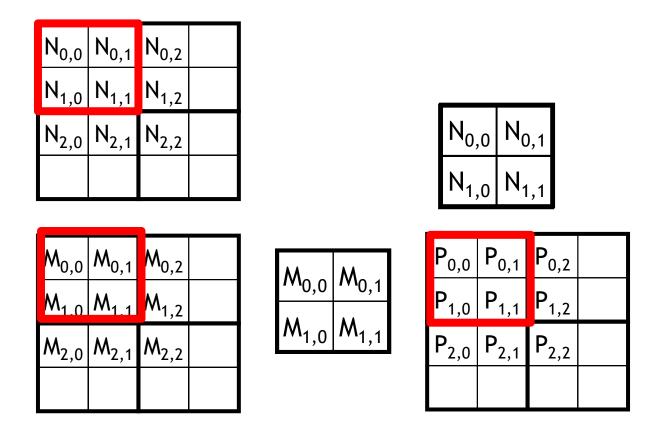
### **Tiled Matrix Multiplication Kernel**

```
global void MatrixMulKernel(float* M, float* N, float* P, int Width)
  shared float ds M[TILE WIDTH][TILE WIDTH];
  shared float ds N[TILE WIDTH][TILE WIDTH];
int bx = blockIdx.x; int by = blockIdx.y;
 int tx = threadIdx.x; int ty = threadIdx.y;
 int Row = by * TILE WIDTH + ty;
 int Col = bx * TILE WIDTH + tx;
 float Pvalue = 0;
// Loop over the M and N tiles required to compute the P element
for (int p = 0; p < Width/TILE WIDTH; ++p) {
   // Collaborative loading of M and N tiles into shared memory
   ds M[ty][tx] = M[Row*Width + p*TILE WIDTH+tx];
   ds N[ty][tx] = N[(p*TILE WIDTH+ty)*Width + Col];
   syncthreads();
   for (int i = 0; i < TILE WIDTH; ++i)
     Pvalue += ds M[ty][i] * ds N[i][tx];
   synchthreads();
P[Row*Width+Col] = Pvalue;
```

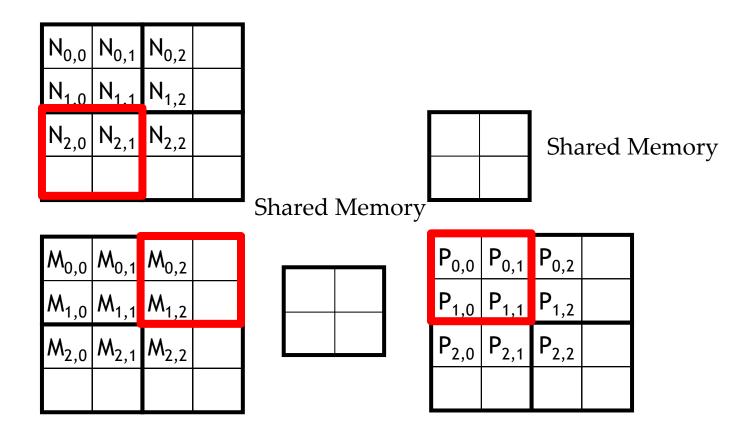
## **Handling Matrix of Arbitrary Size**

- Real applications need to handle arbitrary sized matrices.
  - One could pad (add elements to) the rows and columns into multiples of the tile size, but would have significant space and data transfer time overhead.
- We will take a different approach.
  - Boundary condition checking
  - Regularizing tile contents

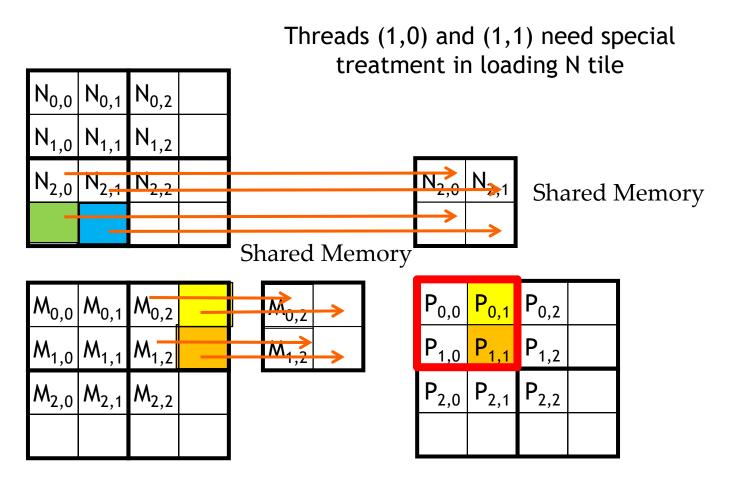
## Phase 0 Loads for Block (0,0) for a 3x3 Example



## Phase 1 Loads for Block (0,0) for a 3x3 Example



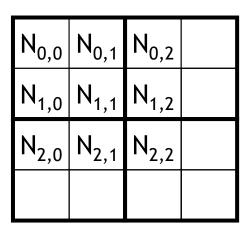
#### Phase 1 Loads for Block (0,0) for a 3x3 Example



Threads (0,1) and (1,1) need special treatment in loading M tile

If we let thread(0,1) read from M(0,3), what will actually be read into SM(0,1)?

## Phase 0 Use for Block (0,0) (iterations 0 and 1)



## Which threads participate in computations?

N <sub>0,0</sub>	N <sub>0,1</sub>
N <sub>1,0</sub>	N <sub>1,1</sub>

Shared Memory

Shared Memory

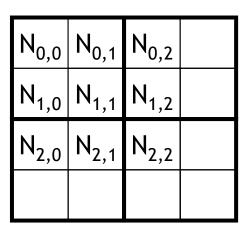
Man Man

$M_{0,0}$	$M_{0,1}$	$M_{0,2}$	
$M_{1,0}$		$M_{1,2}$	
$M_{2,0}$	M <sub>2,1</sub>	M <sub>2,2</sub>	

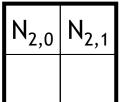
$M_{0,0}$	$M_{0,1}$
$M_{1,0}$	$M_{1,1}$

P <sub>0,0</sub>	P <sub>0,1</sub>	P <sub>0,2</sub>	
P <sub>1,0</sub>	P <sub>1,1</sub>	P <sub>1,2</sub>	
P <sub>2,0</sub>	P <sub>2,1</sub>	P <sub>2,2</sub>	

## Phase 1 Use for Block (0,0) (iteration 0)



## Which threads participate in computations?



Shared Memory

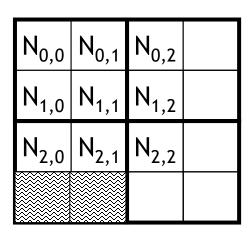
Shared Memory

$M_{0,0}$	$M_{0,1}$	$M_{0,2}$	
$M_{1,0}$	$M_{1,1}$	$M_{1,2}$	
$M_{2,0}$	$M_{2,1}$	M <sub>2,2</sub>	

$M_{0,2}$	
$M_{1,2}$	

P <sub>0,0</sub>	P <sub>0,1</sub>	P <sub>0,2</sub>	
P <sub>1,0</sub>	P <sub>1,1</sub>	P <sub>1,2</sub>	
P <sub>2,0</sub>	P <sub>2,1</sub>	P <sub>2,2</sub>	

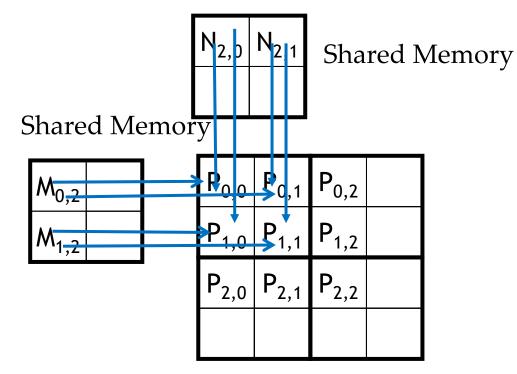
## Phase 1 Use for Block (0,0) (iteration 0)



$M_{0,0}$	$M_{0,1}$	$M_{0,2}$	
$M_{1,0}$	$M_{1,1}$	$M_{1,2}$	
$M_{2,0}$	M <sub>2,1</sub>	M <sub>2,2</sub>	

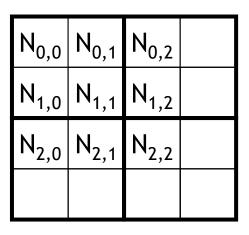
#### **Participation in Computations!**

All threads operate on legitimate data

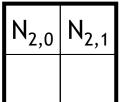


Thread(0,1) participates in [M(0,2)\*N(2,1)] but should not read M(0,3)! Thread(1,1) participates in (M(1,2)\*N(2,1)] but should not read M(1,3)! Similar for Thread(1,0) and Thread(1,1) on N(3,0) ad N(3,1)

## Phase 1 Use for Block (0,0) (iteration 1)



## Which threads participate in computations?



Shared Memory

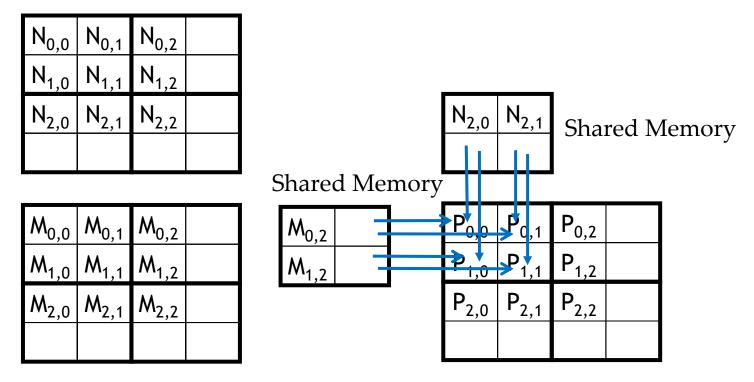
Shared Memory

$M_{0,0}$	$M_{0,1}$	$M_{0,2}$	
$M_{1,0}$		$M_{1,2}$	
$M_{2,0}$			

M <sub>0,2</sub>	
$M_{1,2}$	

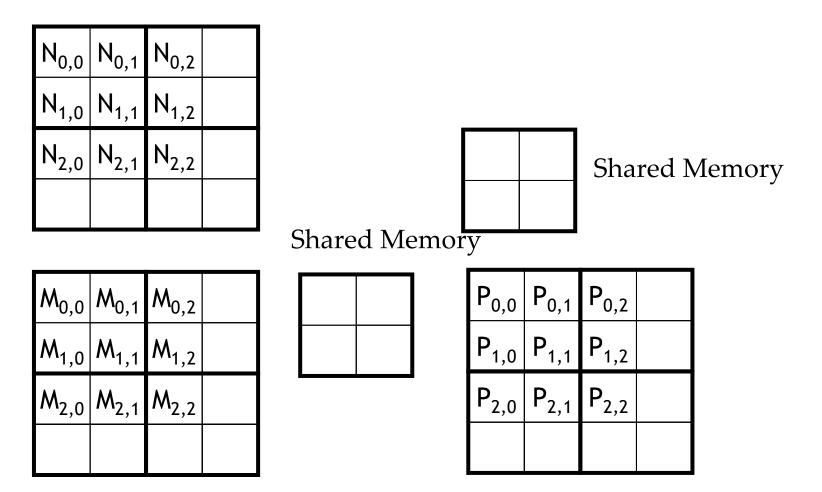
P <sub>0,0</sub>	P <sub>0,1</sub>	P <sub>0,2</sub>	
P <sub>1,0</sub>	P <sub>1,1</sub>	P <sub>1,2</sub>	
P <sub>2,0</sub>	P <sub>2,1</sub>	P <sub>2,2</sub>	

## Phase 1 Use for Block (0,0) (iteration 1)



All Threads need special treatment. None of them should introduce invalidate contributions to their P elements.

## Load problems can occur even during phase 0!

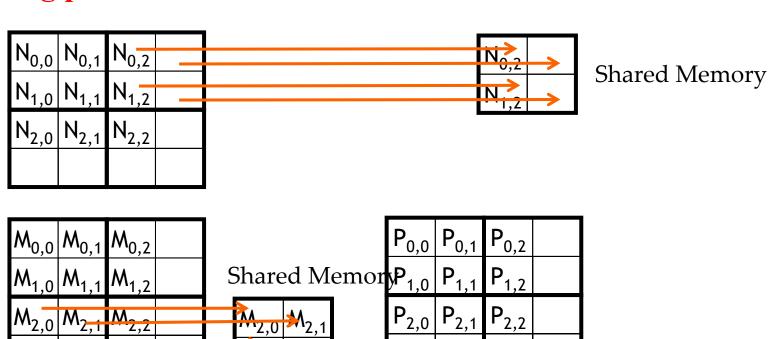


Which thread block will have threads not participating in loads during Phase 0?

### Phase 0 Loads for Block (1,1) for a 3x3 Example

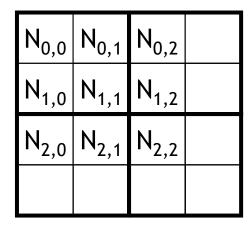
Problems can occur even during phase 0!

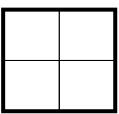
Threads (0,1) and (1,1) need special treatment in loading N tile



Threads (1,0) and (1,1) need special treatment in loading M tile

# Threads not participating in computations in phase 0 may participate in loads. Really???

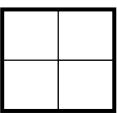




Shared Memory

Shared Memory

$M_{0,0}$	$M_{0,1}$	$M_{0,2}$	
$M_{1,0}$		$M_{1,2}$	
$M_{2,0}$			



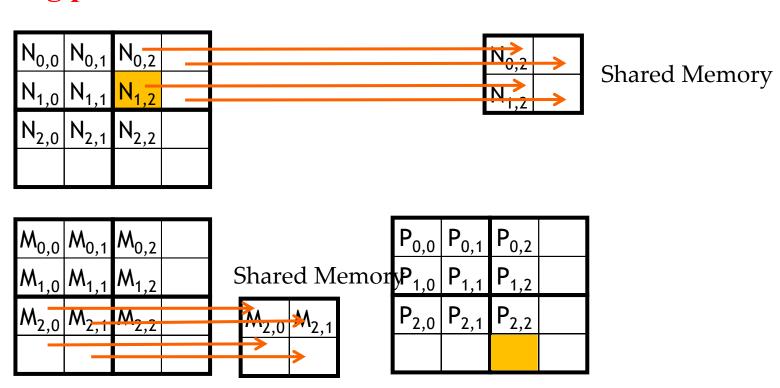
P <sub>0,0</sub>	P <sub>0,1</sub>	P <sub>0,2</sub>	
P <sub>1,0</sub>	P <sub>1,1</sub>	P <sub>1,2</sub>	
P <sub>2,0</sub>	P <sub>2,1</sub>	P <sub>2,2</sub>	

Which thread block?

### Phase 0 Loads for Block (1,1) for a 3x3 Example

Problems can occur even during phase 0!

Threads (0,1) and (1,1) need special treatment in loading N tile



Threads (1,0) and (1,1) need special treatment in loading M tile

Phase 0 of Block(1,1), Thread(1,0), assigned to calculate non-existent P[3,2] but need to participate in loading tile element N[1,2]

## **Major Cases in Toy Example**

- We can not just simply exclude the threads that do not calculate valid P elements
  - they still need to participate in loading the input tiles
  - Phase 0 of Block(1,1), Thread(1,0), assigned to calculate non-existent P[3,2] but need to participate in loading tile element N[1,2]
- Threads that calculate valid P elements may attempt to load non-existing input elements when loading input tiles
  - Phase 0 of Block(0,0), Thread(1,0), assigned to calculate valid P[1,0] but attempts to load nonexisting N[3,0]

#### **Next**

Writing boundary conditions for matrix multiplication