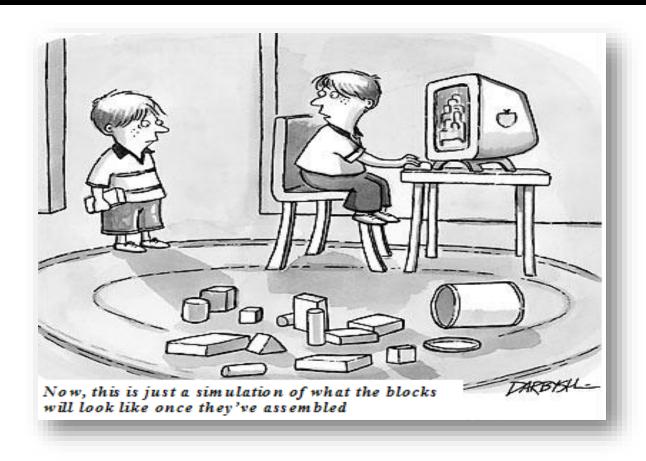
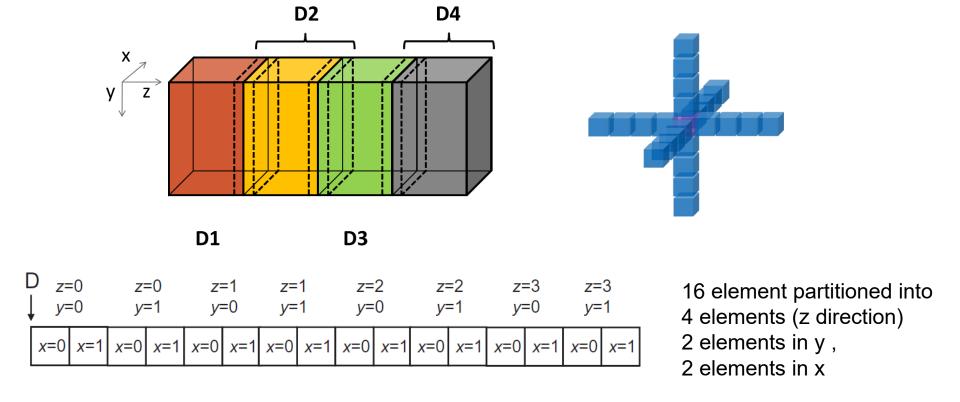
ECE569 Module 50



MPI Example: 3D Heat Transfer

3D Heat Transfer

- Heat transfer based on Jacobi Iterative Method
 - in each iteration or time step, value of a point calculated as a weighted sum of 4 neighbors in each direction



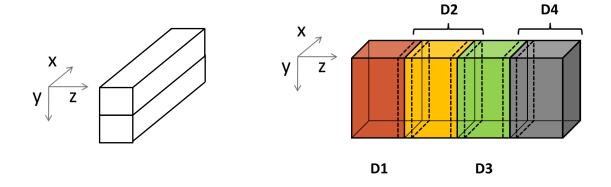
Wave Propagation: Kernel Code

```
/* Coefficients used to calculate the laplacian */
constant float coeff[5];
 global void wave propagation(float *next, float *in,
                 float *prev, float *velocity, dim3 dim)
    unsigned x = threadIdx.x + blockIdx.x * blockDim.x;
    unsigned y = threadIdx.y + blockIdx.y * blockDim.y;
    unsigned z = threadIdx.z + blockIdx.z * blockDim.z;
    /* Point index in the input and output matrixes */
    unsigned n = z * dim.x * dim.y + y * dim.x + x;
    /* Only compute for points within the matrixes */
    if(x < dim.x \&\& y < dim.y \&\& z < dim.z) {
        /* Calculate the contribution of each point to the laplacian */
        float laplacian = coeff[0] + in[n];
                                                             D2
```

D1

D3

Wave Propagation: Time Integration



Wave Propagation: Main Process

```
int main(int argc, char *argv[]) {
    int pad = 0, dimx = 480+pad, dimy = 480, dimz = 400, nreps = 100;
    int pid=-1, np=-1;
    MPI Init(&argc, &argv);
    MPI Comm rank(MPI COMM WORLD, &pid);
    MPI Comm size(MPI COMM WORLD, &np);
    if(np < 3) {
        if(0 == pid) printf("Needed 3 or more processes.\n");
        MPI Abort( MPI_COMM_WORLD, 1 ); return 1;
    if(pid < np - 1)
        compute_node(dimx, dimy, dimz / (np - 1), nreps);
    else
        data_server( dimx,dimy,dimz, nreps );
                                                              number of
    MPI Finalize();
                                                              iterations that need
    return 0;
                                                              to be done for
                                                              all the data points
                                                              in the grid.
```

Stencil Code: Server Process (I)

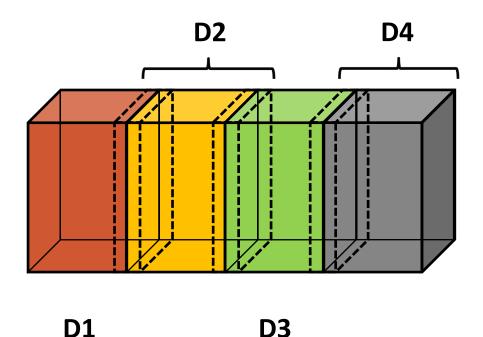
```
void data server(int dimx, int dimy, int dimz, int nreps) {
1
   int np, num comp nodes = np - 1, first node = 0, last node = np - 2;
2
   unsigned int num points = dimx * dimy * dimz;
   unsigned int num bytes = num points * sizeof(float);
   float *input=0, *output = NULL, *velocity = NULL;
5
    /* Set MPI Communication Size */
6
   MPI Comm size (MPI COMM WORLD, &np);
   /* Allocate input data */
8
   input = (float *)malloc(num bytes);
   output = (float *)malloc(num bytes);
10
   velocity = (float *)malloc(num bytes);
      if(input == NULL || output == NULL || velocity == NULL) {
11
12
       printf("Server couldn't allocate memory\n");
13
       MPI Abort( MPI COMM WORLD, 1 );
14
15
   /* Initialize input data and velocity */
16
   random data(input, dimx, dimy, dimz, 1, 10);
17
   random data(velocity, dimx, dimy, dimz, 1, 10);
18 float *send address = input;
```

Stencil Code: Server Process (II)

```
/* Calculate number of shared points */
   int edge_num_points = dimx * dimy * (dimz / num_comp_nodes + 4);
19
   int int num points = dimx * dimy * (dimz / num comp nodes + 8);
20
                               D2
```

Internal partitions receive data from left and right

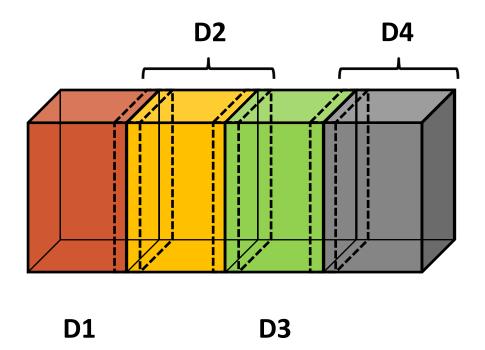
Stencil Code: Server Process (III)



Stencil Code: Server Process (IV)

```
/* adjust send address for internal nodes */
22
      send address =
    /* Send input data to "internal" compute nodes */
   for(int process = 1; process < last node; process++) {</pre>
23
24
       MPI Send(send address, int num points, MPI FLOAT, process,
               DATA DISTRIBUTE, MPI COMM WORLD);
       send address += dimx * dimy * (dimz / num comp nodes);
25
26
                     D1
                                  D3
```

Stencil Code: Server Process (V)



Stencil Code: Server Process (VI)

```
28 float *velocity send address = velocity;
    /* Send velocity data to compute nodes */
29
   for(int process = 0; process < last node + 1; process++) {</pre>
30
            MPI Send(velocity send address, edge num points,
              MPI FLOAT, process, DATA DISTRIBUTE, MPI COMM WORLD);
           velocity send address += dimx * dimy *
31
                                             (dimz / num comp nodes);
32 }
   /* Wait for nodes to compute */
33 MPI Barrier(MPI COMM WORLD);
   /* Collect output data */
34 MPI Status status;
35 for(int process = 0; process < num comp nodes; process++)</pre>
36
       MPI Recv(output + process * num points / num comp nodes,
           num points / num comp nodes, MPI FLOAT, process,
           DATA COLLECT, MPI COMM WORLD, &status );
37 }
```

Stencil Code: Server Process (VII)

```
/* Store output data */
38 store_output(output, dimx, dimy, dimz);

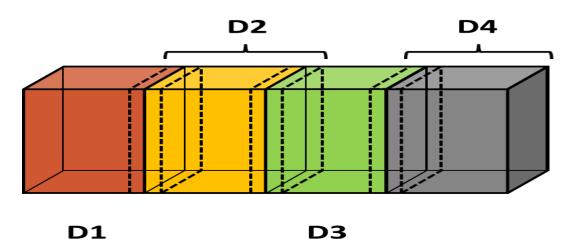
/* Release resources */
39 free(input);
40 free(velocity);
41 free(output);
42}
```

Stencil Code: Compute Process (I)

```
void compute node stencil(int dimx, int dimy, int dimz, int nreps ) {
   int np, pid;
   MPI Comm rank (MPI COMM WORLD, &pid);
   MPI Comm size(MPI COMM WORLD, &np);
                                  = dimx * dimy * (dimz + 8);
   unsigned int num points
   unsigned int num bytes
                                  = num points\* sizeof(float);
   unsigned int num ghost points = 4 * dimx * dimy;
6
                                  = num_ghost_points *
   unsigned int num ghost bytes
                                                         For simplicity both
                                     sizeof(float);
                                                         edge and internal
8
   int server process = np-1;
                                                         nodes allocate same
    /* Alloc host memory */
                                                         amount, +8 cells
    float *h input = (float *)malloc(num bytes);
    /* Alloc device memory for input and output data *>
10 float *rcv address = h input + num_ghost_points * (0 == pid);
11
    MPI Recv(rcv address, num points, MPI FLOAT, server process,
             MPI ANY TAG, MPI COMM WORLD, &status );
```

Allocated +8 for node 0 receive buffer, left padding not used!

Stencil Code: Compute Process (II)



MPI Sending and Receiving Data

- int MPI_Sendrecv(void *sendbuf, int sendcount, MPI_Datatype sendtype, int dest, int sendtag, void *recvbuf, int recvcount, MPI_Datatype recvtype, int source, int recvtag, MPI_Comm comm, MPI_Status *status)
 - Sendbuf: Initial address of send buffer (choice)
 - Sendcount: Number of elements in send buffer (integer)
 - Sendtype: Type of elements in send buffer (handle)
 - Dest: Rank of destination (integer)
 - Sendtag: Send tag (integer)
 - Recvcount: Number of elements in receive buffer (integer)
 - Recvtype: Type of elements in receive buffer (handle)
 - Source: Rank of source (integer)
 - Recvtag: Receive tag (integer)
 - Comm: Communicator (handle)
 - Recvbuf: Initial address of receive buffer (choice)
 - Status: Status object (Status). This refers to the receive operation.

Stencil Code: Compute Process (III)

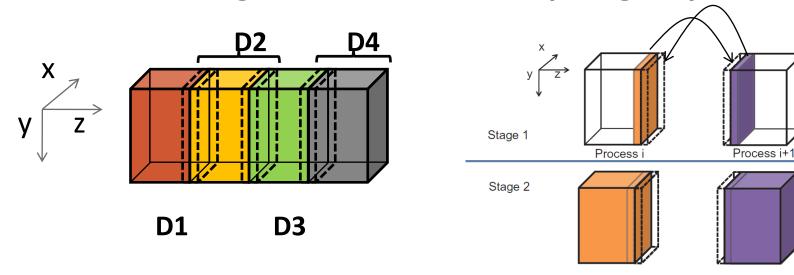
```
/* Send data to left, get data from right */
MPI Sendrecv(h left boundary, num halo points, MPI FLOAT,
           left neighbor, i, h right halo,
           num halo points, MPI FLOAT, right neighbor, i,
           MPI COMM WORLD, &status);
/* Send data to right, get data from left */
MPI Sendrecv(h right boundary, num halo points, MPI FLOAT,
           right neighbor, i, h left halo,
           num halo points, MPI FLOAT, left neighbor, i,
           MPI COMM WORLD, &status );
cudaMemcpy(d output+left halo offset, h left halo,
           num halo bytes, cudaMemcpyHostToDevice);
cudaMemcpy(d output+right ghost offset, h right ghost,
           num halo bytes, cudaMemcpyHostToDevice);
cudaDeviceSynchronize();
float *temp = d output;
d output = d input; d input = temp;
```

Stencil Code: Compute Process (IV)

```
/* Wait for previous communications */
MPI Barrier(MPI COMM WORLD);
float *temp = d output;
d output = d input;
d input = temp;
/* Send the output, skipping halo points */
cudaMemcpy(h output, d output, num bytes,
          cudaMemcpyDeviceToHost);
float *send address = h output + num ghost points;
MPI Send(send address, dimx * dimy * dimz, MPI REAL,
       server process, DATA COLLECT, MPI COMM WORLD);
MPI Barrier(MPI COMM WORLD);
/* Release resources */
free(h input); free(h output);
cudaFreeHost(h left ghost own);
cudaFreeHost(h right ghost own);
cudaFreeHost(h_left_ghost); cudaFreeHost(h_right_ghost);
cudaFree( d input ); cudaFree( d output );
```

Computation Efficiency

- Rather than having each node process all data, let each node first compute for its edge cells (Stage 1)
- Stream edge cell values to neighbors while processing rest of the data (Stage 2)



```
float *h output = NULL, *d output = NULL, *d vsq = NULL;
   float *h output = (float *)malloc(num bytes);
   cudaMalloc((void **)&d output, num bytes );
   float *h left boundary = NULL, *h right boundary = NULL;
   float *h left halo = NULL, *h right halo = NULL;
   /* Alloc host memory for ghost data */
   cudaHostAlloc((void **)&h left boundary, num ghost bytes,
cudaHostAllocDefault);
   cudaHostAlloc((void **)&h right boundary,num ghost bytes,
cudaHostAllocDefault);
   cudaHostAllocDefault);
   cudaHostAllocDefault);
                                             Left halo offset right halo offset
   /* Create streams used for stencil computation */
   cudaStream t stream0, stream1;
   cudaStreamCreate(&stream0);
   cudaStreamCreate(&stream1);
                                          left stage 1 offset
                                                    right stage 1 offset
```

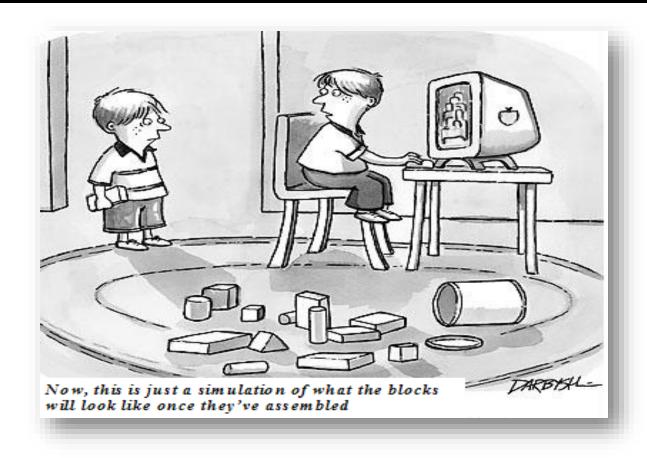
```
MPI Status status;
int left neighbor = (pid > 0) ? (pid - 1) : MPI PROC NULL;
int right neighbor = (pid < np - 2) ? (pid + 1) : MPI PROC NULL;
/* Upload stencil cofficients */
upload coefficients (coeff, 5);
int left halo offset = 0;
int right halo offset = dimx * dimy * (4 + dimz);
int left stage1 offset = 0;
int right stage1 offset = dimx * dimy * (dimz - 4);
int stage2 offset = num halo points;
MPI Barrier( MPI COMM WORLD );
                                     Left halo offset
                                                   right halo offset
/* Compute boundary*/
/* Stage 1 */
/* Compute remaining points */
/* Stage 2 */
                                 left stage 1 offset
                                                 right stage 1 offset
```

```
/* Send data to left, get data from right */
MPI Sendrecv(h left boundary, num halo points, MPI FLOAT,
           left neighbor, i, h right halo,
           num halo points, MPI FLOAT, right neighbor, i,
           MPI COMM WORLD, &status);
/* Send data to right, get data from left */
MPI Sendrecv(h right boundary, num halo points, MPI FLOAT,
           right neighbor, i, h left halo,
           num halo points, MPI FLOAT, left neighbor, i,
           MPI COMM WORLD, &status );
cudaMemcpyAsync(d output+left halo offset, h left halo,
           num halo bytes, cudaMemcpyHostToDevice, stream0);
cudaMemcpyAsync(d output+right ghost offset, h right ghost,
           num halo bytes, cudaMemcpyHostToDevice, stream0 );
cudaDeviceSynchronize();
float *temp = d output;
d output = d input; d input = temp;
```

Next

Optimizing GPU Program

ECE569 Module 51



Optimizing GPU Programs

Fundamental Issues

Parallel computing requires that

- The problem can be decomposed into sub-problems that can be safely solved at the same time
- The programmer structures the code and data to solve these sub-problems concurrently
- The problems must be large enough to justify parallel computing and to exhibit exploitable concurrency.

Optimizing GPU Programs

- ☐ Decrease arithmetic intensity
- ☐ Decrease time spent on memory operations
- ☐ Coalesce global memory accesses
- ☐ Do fewer memory operations per thread
- **□** Avoid thread divergence
- ☐ Move all data to shared memory

Levels of Optimization

- Algorithm selection
- Basic efficiency principles
- Architecture specific optimizations
- Instruction level optimizations

Establishing a Strategy

- Analyze, Parallelize, Optimize, Deploy
 - Profile guided optimization
 - Deploy early
- Amdahl's Law

Establishing an Upper Limit

- Theoretical Peak Bandwidth
 - Memory clock: 2GHz
 - Memory bus: 128 bits

- 40-60% okay
- 60-75% not bad
- >75% very good

Establishing an Upper Limit

- Theoretical Peak Bandwidth
 - Memory clock: 1.6GHz
 - Memory bus: 128 bits
 - Kernel
 - Data size 1024x1024 elements (4 byte each)
 - 2 memory transactions per element
 - Execution time 0.5ms

DRAM Utilization (Global Memory)

Coalescing

Utilization of the bytes delivered by memory

What can we do to improve bandwidth?

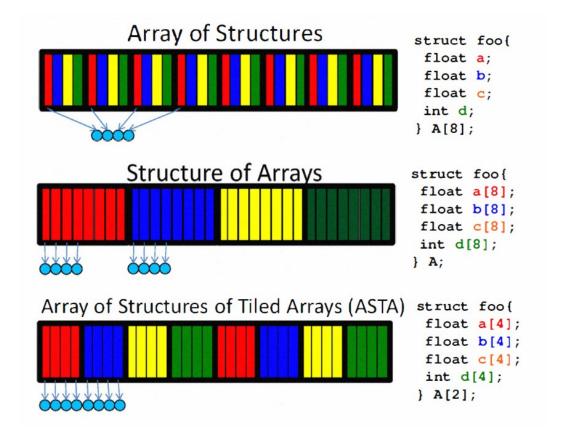
- a. Increase the number of bytes delivered
- b. Increase the latency (time between transactions)
- c. Decrease the number of bytes delivered
- d. Decrease the latency (time between transactions)

Data Transformation

```
struct my Data {
                                      struct my Data {
 float a;
                                       float a[128];
 float b;
                                       float b[128];
 float c;
                                       float c[128];
 float d;
                                       float d[128];
};
                                      };
struct my Data data[128];
                                      struct my Data data;
Array of Structures
                                      Structure of arrays
int i = threadIdx.x;
                                      int i = threadIdx.x;
data[i].a++;
                                      data.a[i]++;
data[i].b += data[i].*
                                      data.b[i] += data.c[i]*
data[i].d;
                                      data.d[i];
```

In case your shared memory is performing worse

Data layout



. A. Stratton *et al.*, "Algorithm and Data Optimization Techniques for Scaling to Massively Threaded Systems," in *Computer*, vol. 45, no. 8, pp. 26-32, August 2012.

DRAM Utilization (Global Memory)

Coalescing

Utilization of the bytes delivered by memory

Tiling

Thread synchronization

What can we do to reduce the average time per thread?

- a. Eliminate Syncthreads call
- b. Reduce the number of threads per block
- c. Increase the number of threads per block
- d. Increase the number of blocks per SM

Tiling

- If threads access overlapping parts of a data
 - Buffering data into fast on-chip storage for repeated access
 - In CPU: resize data so that it fits into cache
 - Not so good for GPU with thousands of threads
 - Instead explicitly copy the data into the shared memory

Which one will benefit from tiling?

```
__global__ void avg(float* out, float* a, float* b, float* c, float*d, float*e) {
    int i = threIdx.x;
    out[i] = (a[i]+b[i]+c[i]+d[i]+e[i])/5.0f;
}

__global__ void avg(float* out, float* a) {
    int i = threIdx.x;
    out[i] = (a[i-2]+a[i-1]+a[i]+a[i+1]+a[i+2])/5.0f;
}
```

DRAM Utilization (Global Memory)

Coalescing

Utilization of the bytes delivered by memory

Tiling

Thread synchronization

Occupancy

- Thread blocks (16 /32)
- Threads per block (1024)
- Threads per SM (2048)
- Registers per SM (64K)
- Register per thread (255)
- Shared memory per thread block(48KB/96KB/64KB)

Occupancy

- Given thread blocks per SM (16), threads per block (1024), Threads per SM (2048), Registers per SM (65,536 registers), Register per thread (255), Shared memory per thread block(48KB), calculate the maximum number of thread blocks per SM for a kernel assuming grid size is (32,32,1), block size of (32,32,1), 7 registers per thread, and 4KB shared memory per block. Which resource prevents us from running more?
 - Maximum number of threads/SM
 - Maximum number of registers/SM
 - Maximum shared memory/SM
 - Maximum thread blocks/SM

Thread divergence

- Maximizing useful computations/second
 - Minimize thread divergence
- What is the maximum branch divergence penalty (slow down factor) for a CUDA thread block with 1024 threads?

```
Switch(expression) {
  case1: ...break;
  case2: ...break;
  :
  case32:...break;
}
```

Thread divergence

 What will be the slowdown factor for the following expression in switch statement?

```
switch (threadIdx.x%32) {
   case 0:
   case 1:
   :
   case 31:
}
```

Assume kernel is launched as 1 block, 1024 threads per block.

Thread divergence

 What will be the slowdown factor for the following expression in switch statement?

```
switch (threadIdx.x%64) {
   case 0:
   case 1:
   :
   :
   case 63:
```

Assume kernel is launched as 1 block, 1024 threads per block.

Scatter-to-Gather transformation

Scatter:

 Threads are assigned to the inputs and each one is deciding where it needs to write

Gather:

 Threads are assigned to the output elements and each one is deciding where it needs to read from

Which one will run more efficiently? (i = threadIdx.x)

Privatization

- Tiling
- Problem: what if threads write into the same address?
 - Histogram
 - Local (per thread and thread blocks) and then global

Common optimization techniques

Suggested Reading:

- J. A. Stratton *et al.*, "Algorithm and Data Optimization Techniques for Scaling to Massively Threaded Systems," in *Computer*, vol. 45, no. 8, pp. 26-32, August 2012.
- J. A. Stratton *et al.*, "Optimization and architecture effects on GPU computing workload performance," 2012 Innovative Parallel Computing (InPar), San Jose, CA, 2012, pp. 1-10.

Summary

Measure and improve memory bandwidth

- Assure sufficient occupancy
- Coalesce global memory access
- Minimize latency between accesses
 - Synchronization: too many threads waiting
- Minimize thread divergence in a warp
- Avoid code with if, switch
 - It is ok to have if statement for edge cases
- Avoid workload imbalance across threads
- Consider using built-in math functions ..sin();
- Single precision vs. double precision
 - 3.14 vs 3.14f (single precision!)
- Streams
 - Overlap computation and memory transfers

Easy to Learn

Takes time to master

