

Problem Spaces

Given an intersection:

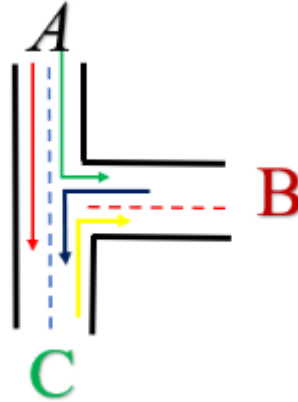


Figure 1: Intersection

Design a control pattern for traffic lights so that the intersection as safe as possible.

The first thing we need to find out is which turns are compatible.

Identify safe and unsafe turns.

| | AB | AC | BC | CB |
|----|----|----|----|----|
| AB | ■ | O | X | X |
| AC | O | ■ | X | O |
| BC | X | X | ■ | O |
| CB | X | O | O | ■ |

× Unsafe

O Safe

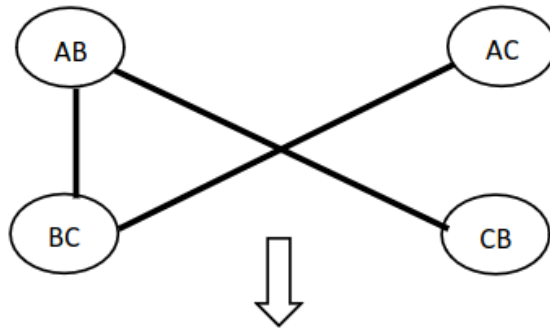
| | AB | AC | BC | CB |
|----|----|----|----|----|
| AB | - | O | X | X |
| AC | O | - | X | O |
| BC | X | X | - | O |
| CB | X | O | O | - |

Convert the table into a graph.

$G = \langle V, E \rangle$

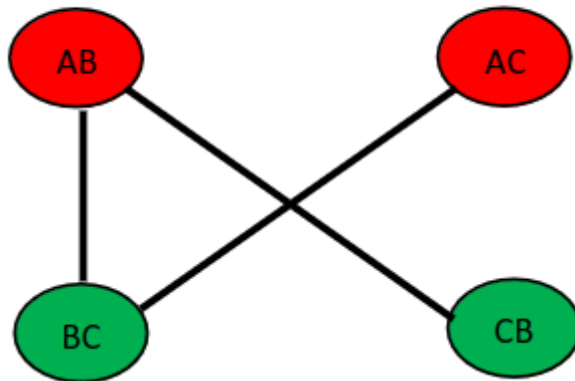
Vertices \rightarrow trans

Edges \rightarrow let the edge connect trans that are not compatible (unsafe)



What does each edge signify?

Each edge displays a pair of incompatible turns.



Color the Graph

Graph Coloring

Coloring parts of a graph so that no two edges/nodes are the same color.

Picking a set of parameters that you want to prioritize and use them efficiently.

- Greedy optimization

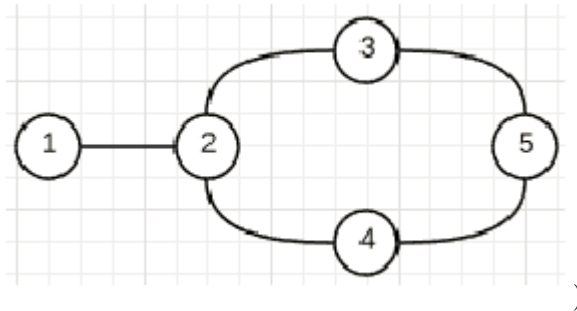
– Opportunistic Approach

How do we color the graph?

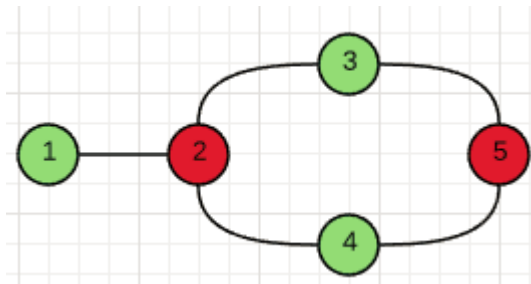
- Pick a node and color it with a new color.
- Scan the list of nodes not connected to this node and paint it with the same color.
 - Repeat this step until there are no more nodes that can be painted with this color.
- Pick up a new color and repeat 1 - 3 until you are done.

What are some pitfalls of this greedy algorithm?

Given the following graph labeled $\{1, 2, 3, 4, 5\}$

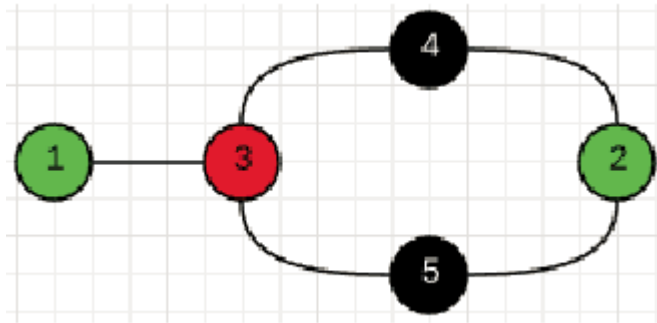


When completing the algorithm in the sequential order:



Everything turns out fine.

However, when we change the order that the nodes are colored in, the minimum number of colors cannot be achieved with the greedy approach.



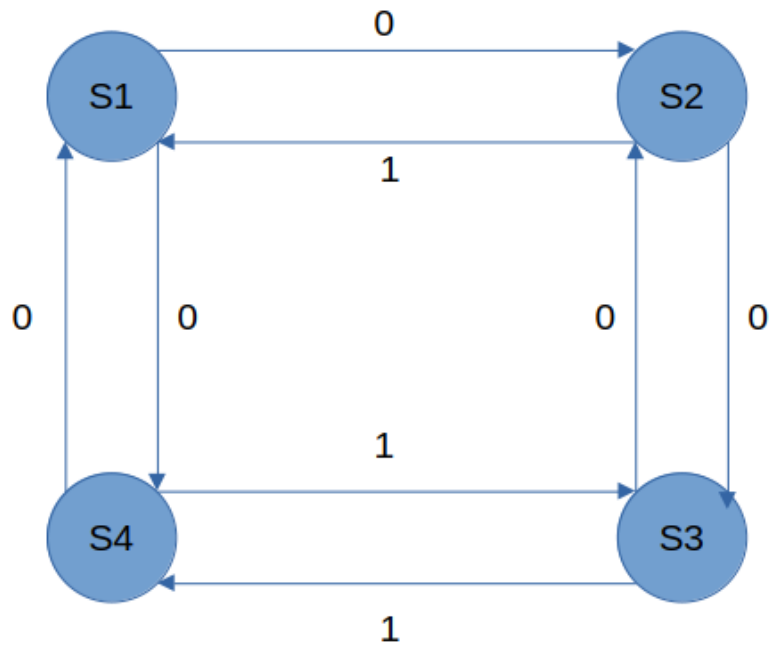
Creating a formal system to analyze this process: State Spaces

State Spaces

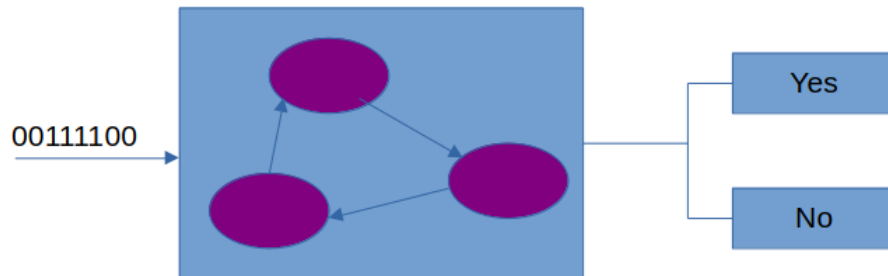
Capturing the condition/situation of a system at a point in time. State spaces are a representation method.

Given the graph representation of the finite state acceptor. In a device that accepts or rejects strings.

$Q = \langle S_1, S_2, S_3, S_4 \rangle$



Any string that has an even number of zeroes and even number of ones, then it is accepted. If the string contains an odd number of zeroes or odd number of ones, then the string is not accepted.



Example | 8-Puzzle

8 Puzzle

initial state

3X 3 board with 9 tiles

{1,2,3,.....,8 and blank tile}

| | | |
|---|---|---|
| 2 | 8 | 3 |
| 1 | 6 | 4 |
| 7 | | 5 |

Goal state (final state)

| | | |
|---|---|---|
| 1 | 2 | 3 |
| 8 | | 4 |
| 7 | 6 | 5 |

- To solve it; need to make a series of moves

Abstraction for moves

Instead of moving tiles next to the blank, you move the blank tile.

MOVE BLANK UP, DOWN, LEFT, RIGHT (U, D, L, R)

A solution is the sequence of moves that gives me the goal state from the initial state.

Back to the example:

up, up, left, down, right.

How To Do This Systematically

1. Define the state space (DB)

- A 3×3 board

$$DB = \{T_{i,j} | T_{i,j} \in \{1, 2, \dots, 8, \text{blank} | i, j = 1, 2, 3, 4, 5, 6, 7, 8\}$$

And not true that there is: $T_{i,j} = T_{m,n} \in \{1, 2, 3\}$

Is it possible that given a Goal and Initial state: Is it possible that there does not exist a path to the goal state from the initial state?

Yes, it is possible.

Operators: “Moves” that allow you to transform/transition between states.