Problem Spaces

Given an intersection:

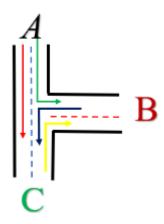


Figure 1: Intersection

Design a control pattern for traffic lights so that the intersection as safe as possible.

The first thing we need to find out is which turns are compatible.

Identify safe and unsafe turns.

	AB	AC	BC	СВ
AB		0	×	×
AC	0		×	0
BC	×	×		0
СВ	×	0	0	

× Unsafe

O Safe

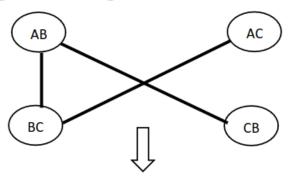
	AB	AC	BC	СВ
AB	-	Ο	X	X
AC	O	-	X	O
BC	X	X	-	Ο
СВ	X	O	Ο	-

Convert the table into a graph.

$$G = \langle V, E \rangle$$

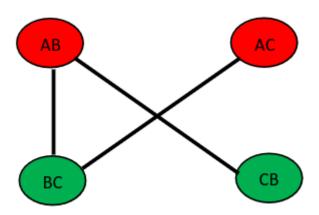
Vertices → trans

Edges → let the edge connect trans that are not compatible (unsafe)



What does each edge signify?

Each edge displays a pair of incompatible turns.



Color the Graph

Graph Coloring

Coloring parts of a graph so that no to edges/nodes are the same color.

Picking a set of parameters that you want to prioritize and use them efficiently.

• Greedy optimization

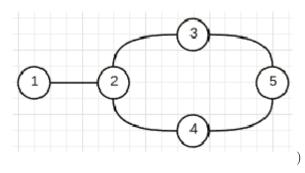
- Opportunistic Approach

How do we color the graph?

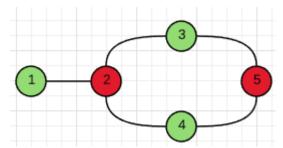
- Pick a node and color it with a new color.
- Scan the list of nodes not connected to this node and paint it with the same color.
 - Repeat this step until there are no more nodes that can be painted with this color.
- Pick up a new color and repeat 1 3 until you are done.

What are some pitfalls of this greedy algorithm?

Given the following graph labeled {1, 2, 3, 4, 5}

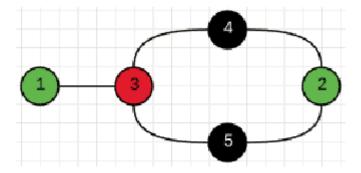


When completing the algorithm in the sequential order:



Everything turns out fine.

However, when we change the order that the notes are colored in, the minimum number of colors cannot be achieved with the greedy approach.



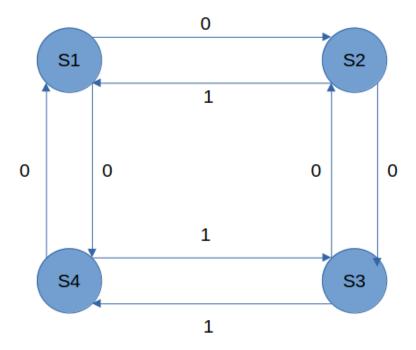
Creating a formal system to analyze this process: State Spaces

State Spaces

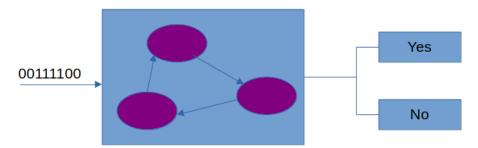
Capturing the condition/situation of a system at a point in time. State paces are a representation method.

Given the graph representation of the finite state acceptor. In a device that accepts or rejects strings.

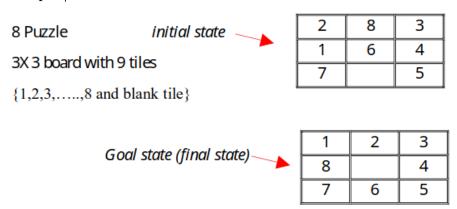
$$Q = \langle S_1, S_2, S_3, S_4 \rangle$$



Any string that has an even number of zeroes and even number of ones, then it is accepted. If the string contains an odd number of zeroes or odd number of ones, then the string is not accepted.



Example | 8-Puzzle



To solve it; need to make a series of moves

Abstraction for moves

Instead of moving tiles next to the black, you move the black tile.

MOVE BLACK UP, DOWN, LEFT, RIGHT (U, D, L, R)

A solution is the sequence of moves that gives me the goal state from the initial state.

Back to the example:

up, up, left, down, right.

How To Do This Systematically

1. Define the state space (DB)

• A 3×3 board

$$DB = \{T_{i,j} | T_{i,j} \in \{1, 2, ..., 8, \text{blank} | i, j = 1, 2, 3, 4, 5, 6, 7, 8\}$$

And not true that there is: $T_{i,j} = T_{m,n} \in \{1,2,3\}$

Is it possible that given a Goal and Initial state: Is it possible that there does not a exist a path to the goal state from the initial state?

Yes, it is possible.

Operators: "Moves" that allow you to transform/transition between states.

AI Production System (AIPS) <- Problem Solving Paradigm

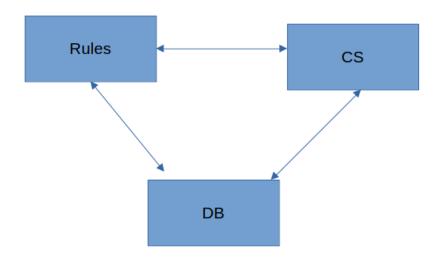


Figure 2: AIPS Block Diagram

DB: Data base/state space Rules: operators that allow you to change states CS: Control strategy

Goal state will sometimes be referred to as goal state/states/termination condition.

So, what are the operators of a given system?

Rules are typically given as an if <condition> then <action>, in terms of state machine representation, the action that is completed when a set of conditions is met changes the state of the system.

Control Strategies are a method that select and apply the "best" possible rule so that a solution can be found in the most efficient way.

The possible states of a given system can be arranged into a tree. You can take care to avoid cycles in the tree.

Does the tree give us a solution?

Yes

Is the solution from the tree efficient?

It is not the most efficient, but it is good enough, we can change what we prioritize if we're smart about it.

Procedure Reduction

Let's take D to be a state from the database.

R is a rule from the set of all possible rules that apply/match to D.

while R is not empty and D does not satisfy the termination condition:

- Select a rule to apply to the database
- Apply the rule to the database
- Update the set of rules that are available

The rule that you select to perform is dependent on how "smart" you want to be to receive the most payoff from your transformation.

Water Jug Problem

Given two jugs of water of capacities $\langle C_1, C_2 \rangle$ and no markings on them, how can you get precisely $\langle x, y \rangle$ amounts of water in them starting with state $\langle i_1, i_2 \rangle$

For us, the following instance:

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initial state: <0,0>
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goal state: <2, y>, where y is any amount of water $0 \le y \le 3$

Let's use the AIPS method to approach this problem.

Such that:

$$DB = \{ \langle x, y \rangle | 0 \le x \le 4 \text{ and } 0 \le y \le 3 \}$$

initial state: <0,0>

goal state: $\langle 2, y \rangle$, where y is any amount of water $0 \leq y \leq 3$

How to go about the solution?