

Cyber Physical Systems - Discrete Models

Exercise Sheet 8 Solution

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2.5/4

15.75/22

Exercise 1: Prefixes and Closure I

1/1 Part A

$$P \subseteq \text{cl}(P)$$

Solution

Let $\omega \in P$, then $\text{pref}(\omega) \subseteq \text{pref}(P)$ trivially. Since $\text{pref}(\omega) \subseteq \text{pref}(P)$ is the predicate for closure, then we can conclude $\forall \omega \in P \rightarrow \omega \in \text{cl}(P)$, therefore $P \subseteq \text{cl}(P)$. ✓

1.5/2 Part B

$$\text{pref}(\text{cl}(P)) = \text{pref}(P)$$

Solution

If we prove both $\text{pref}(\text{cl}(P)) \subseteq \text{pref}(P)$ and $\text{pref}(P) \subseteq \text{pref}(\text{cl}(P))$, then we can conclude $\text{pref}(\text{cl}(P)) = \text{pref}(P)$.

Direction 1: $\text{pref}(\text{cl}(P)) \subseteq \text{pref}(P)$

Let $\omega \in \text{pref}(\text{cl}(P))$, then $\exists \sigma \in \text{cl}(P) \rightarrow \omega \in \text{pref}(\sigma)$. By the definition of closure, $\forall \sigma \in \text{cl}(P) \rightarrow \text{pref}(\sigma) \subseteq \text{pref}(P)$. Hence, $\omega \in \text{pref}(P)$. Which concludes that $\forall \omega \in \text{pref}(\text{cl}(P)) \rightarrow \omega \in \text{pref}(P)$. ✓

Direction 2: $\text{pref}(P) \subseteq \text{pref}(\text{cl}(P))$

Let $\omega \in \text{pref}(P)$, then $\exists \sigma \in P \rightarrow \omega \in \text{pref}(\sigma)$. By proof in part a, we can claim $\forall \sigma \in P \rightarrow \sigma \in \text{cl}(P)$. Therefore, $\omega \in \text{pref}(P) \rightarrow \omega \in \text{pref}(\text{cl}(P))$. ✓

Your line of thinking is correct, but the way you wrote it down is not syntactically correct. -0.5

For example: $\forall \sigma \in \text{cl}(P) \rightarrow \text{pref}(\sigma) \subseteq \text{pref}(P)$

is an implication with no precondition (only a quantifier), thus makes no syntactical sense.

To be safe in the exam, express as much in words as possible.

Exercise 2: Prefixes and Closure II

3.75/6

Part A

1.25/2

$$\begin{aligned}
 P_1 &= \{A_0 A_1 \dots \mid \exists S \subseteq \mathbb{N} \cdot (|S| = 1 \wedge \forall i \in S \cdot a \in A_i)\} \quad \sim 0.5 \\
 P_2 &= \{A_0 A_1 \dots \mid \forall i \in \mathbb{N} \cdot (a \in A_i \rightarrow b \in A_{i+1})\} \quad -0.25 \quad (\checkmark) \\
 P_3 &= \{A_0 A_1 \dots \mid \exists i \in \mathbb{N} \cdot (\forall j \in \mathbb{N} \cdot (j \geq i \rightarrow a \notin A_j))\} \quad (\checkmark) \\
 P_4 &= \{A_0 A_1 \dots \mid a \in A_0 \wedge \exists i \in \mathbb{N} \cdot a \in A_i\} \quad (\checkmark)
 \end{aligned}$$

Part B

1.25/2

$$\begin{aligned}
 \text{pref}(P_1) &= \{A_0 A_1 \dots A_k \mid \exists S \subseteq \mathbb{N} \cdot (|S| \leq 1 \wedge \forall i \in S \cdot a \in A_i)\} \quad \sim 0.5 \\
 \text{pref}(P_2) &= \{A_0 A_1 \dots A_k \mid \forall i \cdot (i < k \wedge a \in A_i \rightarrow b \in A_{i+1})\} \quad (\checkmark) \quad -0.25 \\
 \text{pref}(P_3) &= \{A_0 A_1 \dots A_k \mid \text{true}\} \quad (\checkmark) \\
 \text{pref}(P_4) &= \{A_0 A_1 \dots A_k \mid a \in A_0\} \quad (\checkmark)
 \end{aligned}$$

Part C

1.25/2

$$\begin{aligned}
 \text{cl}(P_1) &= \{A_0 A_1 \dots \mid \exists S \subseteq \mathbb{N} \cdot (|S| \leq 1 \wedge \forall i \in S \cdot a \in A_i)\} \quad \sim 0.5 \\
 \text{cl}(P_2) &= \{A_0 A_1 \dots \mid \forall i \cdot (a \in A_i \rightarrow b \in A_{i+1})\} \quad (\checkmark) \quad -0.25 \\
 \text{cl}(P_3) &= \{A_0 A_1 \dots \mid \text{true}\} \quad (\checkmark) \\
 \text{cl}(P_4) &= \{A_0 A_1 \dots \mid a \in A_0\} \quad (\checkmark)
 \end{aligned}$$

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Exercise 3: Safety & Liveness Properties

Part A

P_1

- Is an invariant.
- Invariant condition: $a \notin S$.

The invariant condition is a propositional logic formula ϕ over the set of atomic propositions, here $\{a, b\}$.
 $\phi = \neg a$ - 0.5

P_2

- Not an invariant.
- Example trace: $(w = \{a\}^\omega) \notin P_2$. However, $(\forall \sigma \in w \cdot \sigma = \{a\})$, and we can give trace $\{a\}\{b\}^\omega \in P_2$ as a counter example, since the set $\{a\}$ also occurs in $\{a\}\{b\}^\omega$. - 0.5

P_3

- Is an invariant.
- Invariant condition: $a \in S \rightarrow b \in S$. $\phi = b \rightarrow a$ - 0.5

P_4

- Not an invariant.
- Example trace: $(w = \{b\}^\omega) \notin P_4$. However, $(\forall \sigma \in w \cdot \sigma = \{b\})$, and we can give trace $\{b\}\{a\}^\omega \in P_2$ as a counterexample, since the set $\{b\}$ also occurs in $\{b\}\{a\}^\omega$. - 0.5

Part B

P_1

- Is a safety property.
- Set of bad prefixes: $\text{BadPref} = \{A_0 A_1 \dots A_n \mid \exists i \in \{0, \dots, n\} \cdot a \in A_i\}$. $\in (2^{AP})^+ - 0.25$

P_2

- Not a safety property.
- Example trace: $(\sigma = \{a\}^\omega) \in (2^{AP})^\omega \setminus P_2$. But $\forall w \in \text{pref}(\sigma)$, we can always extend it with $\sigma\{b\}^\omega \in P_2$ so no σ can be a bad prefix. ✓

P_3

- Is a safety property.
- Set of bad prefixes:
 $\text{BadPref} = \{A_0 A_1 \dots A_n \mid \exists i \in \{0, \dots, n\} \cdot b \in A_i \wedge a \notin A_i\}$. ✓

P_4

- Is a safety property.
- Set of bad prefixes:
 $\text{BadPref} = \{A_0 A_1 \dots A_n \mid \exists S \subseteq \{0, \dots, n\} \cdot |S| > 1 \wedge (\forall i \in S \cdot b \in S_i)\}$ ✓

Part C

P_1

- Is not a liveness property.
- Example bad prefix: $\{a\}$. $\forall w \in (2^{\text{AP}})^\omega \cdot (\{a\}w) \notin P_1$. ✓

P_2

- Is a liveness property.
- We can extend any finite trace with $\{b\}^\omega$. $\forall w \in (2^{\text{AP}})^* \cdot w\{b\}^\omega \in P_2$. ✓

P_3

- Is not a liveness property.
- Example bad prefix: $\{b\}$. $\forall w \in (2^{\text{AP}})^\omega \cdot (\{b\}w) \notin P_3$. ✓

P_4

- Is not a liveness property.
- Example bad prefix: $\{b\}\{b\}$. $\forall w \in (2^{\text{AP}})^\omega \cdot (\{b\}\{b\}w) \notin P_4$. ✓