

# Cyber Physical Systems - Discrete Models

## Exercise Sheet 8 Solution

Alper Ari  
aa508@uni-freiburg.edu

Onur Sahin  
os141@uni-freiburg.de

December 10, 2023

### Exercise 1: Prefixes and Closure I

#### Part A

$$P \subseteq \text{cl}(P)$$

#### Solution

Let  $\omega \in P$ , then  $\text{pref}(\omega) \subseteq \text{pref}(P)$  trivially. Since  $\text{pref}(\omega) \subseteq \text{pref}(P)$  is the predicate for closure, then we can conclude  $\forall \omega \in P \rightarrow \omega \in \text{cl}(P)$ , therefore  $P \subseteq \text{cl}(P)$ .

#### Part B

$$\text{pref}(\text{cl}(P)) = \text{pref}(P)$$

#### Solution

If we prove both  $\text{pref}(\text{cl}(P)) \subseteq \text{pref}(P)$  and  $\text{pref}(P) \subseteq \text{pref}(\text{cl}(P))$ , then we can conclude  $\text{pref}(\text{cl}(P)) = \text{pref}(P)$ .

#### Direction 1: $\text{pref}(\text{cl}(P)) \subseteq \text{pref}(P)$

Let  $\omega \in \text{pref}(\text{cl}(P))$ , then  $\exists \sigma \in \text{cl}(P) \rightarrow \omega \in \text{pref}(\sigma)$ . By the definition of closure,  $\forall \sigma \in \text{cl}(P) \rightarrow \text{pref}(\sigma) \subseteq \text{pref}(P)$ . Hence,  $\omega \in \text{pref}(P)$ . Which concludes that  $\forall \omega \in \text{pref}(\text{cl}(P)) \rightarrow \omega \in \text{pref}(P)$ .

#### Direction 2: $\text{pref}(P) \subseteq \text{pref}(\text{cl}(P))$

Let  $\omega \in \text{pref}(P)$ , then  $\exists \sigma \in P \rightarrow \omega \in \text{pref}(\sigma)$ . By proof in part a, we can claim  $\forall \sigma \in P \rightarrow \sigma \in \text{cl}(P)$ . Therefore,  $\omega \in \text{pref}(P) \rightarrow \omega \in \text{pref}(\text{cl}(P))$ .

## Exercise 2: Prefixes and Closure II

### Part A

$$P_1 = \{A_0A_1\ldots \mid \exists S \subseteq N \cdot (|S| = 1 \wedge \forall i \in S a \in A_i)\}$$

$$P_2 = \{A_0A_1\ldots \mid \forall i \cdot (a \in A_i \rightarrow b \in A_{i+1})\}$$

$$P_3 = \{A_0A_1\ldots \mid \exists i \cdot (\forall j \cdot (j \geq i \rightarrow a \notin A_j))\}$$

$$P_4 = \{A_0A_1\ldots \mid a \in A_0 \wedge \exists^\infty i \cdot a \in A_i\}$$

### Part B

$$\text{pref}(P_1) = \{A_0A_1\ldots A_k \mid \exists S \subseteq N \cdot (|S| \leq 1 \wedge \forall i \in S a \in A_i)\}$$

$$\text{pref}(P_2) = \{A_0A_1\ldots A_k \mid \forall i \cdot (i < k \wedge a \in A_i \rightarrow b \in A_{i+1})\}$$

$$\text{pref}(P_3) = \{A_0A_1\ldots A_k \mid \text{true}\}$$

$$\text{pref}(P_4) = \{A_0A_1\ldots A_k \mid a \in A_0\}$$

### Part C

$$\text{cl}(P_1) = \{A_0A_1\ldots \mid \exists S \subseteq N \cdot (|S| \leq 1 \wedge \forall i \in S a \in A_i)\}$$

$$\text{cl}(P_2) = \{A_0A_1\ldots \mid \forall i \cdot (a \in A_i \rightarrow b \in A_{i+1})\}$$

$$\text{cl}(P_3) = \{A_0A_1\ldots \mid \text{true}\}$$

$$\text{cl}(P_4) = \{A_0A_1\ldots \mid a \in A_0\}$$

## Exercise 3: Safety & Liveness Properties

### Part A

$P_1$

- Is an invariant.
- Invariation condition:  $a \notin S$ .

$P_2$

- Not an invariant.
- Example trace:  $(w = \{a\}^\omega) \notin P_2$ . However,  $\forall \sigma \in w \cdot \sigma = \{a\}$ , and we can give trace  $\{a\}\{b\}^\omega \in P_2$  as a counter example.

$P_3$

- Is an invariant.
- Invariant condition:  $a \in S \rightarrow b \in S$ .

$P_4$

- Not an invariant.
- Example trace:  $(w = \{b\}^\omega) \notin P_4$ . However,  $\forall \sigma \in w \cdot \sigma = \{b\}$ , and we can give trace  $\{b\}\{a\}^\omega \in P_2$  as a counterexample.

### Part B

$P_1$

- Is a safety property.
- Set of bad prefixes:  $\text{BadPref} = \{A_0A_1\dots A_n \mid \exists i \in \{0, \dots, n\} \cdot a \in A_i\}$ .

$P_2$

- Not a safety property.
- Example trace:  $(\sigma = \{a\}^\omega) \in (2^{\text{AP}})^\omega \setminus P_2$ . But  $\forall w \in \text{pref}(\sigma)$ , we can always extend it with  $\sigma\{b\}^\omega \in P_2$  so no  $\sigma$  can be a bad prefix.

$P_3$

- Is a safety property.
- Set of bad prefixes:  
 $\text{BadPref} = \{A_0A_1\dots A_n \mid \exists i \in \{0, \dots, n\} \cdot b \in A_i \wedge a \notin A_i\}$ .

$P_4$

- Is a safety property.
- Set of bad prefixes:  
 $\text{BadPref} = \{A_0A_1\dots A_n \mid \exists S \subseteq \{0, \dots, n\} \cdot |S| > 1 \wedge (\forall i \in S \cdot b \in S_i)\}$

## Part C

$P_1$

- Is not a liveness property.
- Example bad prefix:  $\{a\}$ .  $\forall w \in (2^{\text{AP}})^\omega \cdot (\{a\}w) \notin P_1$ .

$P_2$

- Is a liveness property.
- We can extend any finite trace with  $\{b\}^\omega$ .  $\forall w \in (2^{\text{AP}})^* \cdot w\{b\}^\omega \in P_2$ .

$P_3$

- Is not a liveness property.
- Example bad prefix:  $\{b\}$ .  $\forall w \in (2^{\text{AP}})^\omega \cdot (\{b\}w) \notin P_3$ .

$P_4$

- Is not a liveness property.
- Example bad prefix:  $\{b\}\{b\}$ .  $\forall w \in (2^{\text{AP}})^\omega \cdot (\{b\}\{b\}w) \notin P_4$ .