## Cyber Physical Systems - Discrete Models Exercise Sheet 3 Solution

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## Exercise 1: Intersection of $\omega$ -regular languages

(a)  $L_1$ : It is certain that a is not infinite. However, b or c can be infinite.  $L_2$ : It is certain that b is infinite. However, a and c can also be infinite.  $L_1 \cap L_2$ : It is certain that a is not infinite and b is infinite. However, c can be infinite.

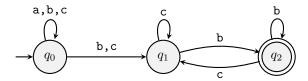


Figure 1: Büchi automaton A that accepts intersection of  $L_1$  and  $L_2$ 

(b)  $L_1$ : It is certain that a is infinite. However, b can also be infinite.  $L_2$ : It is certain that b is infinite. However, a can also be infinite.  $L_1 \cap L_2$ : It is certain that a and b infinite.

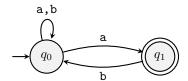


Figure 2: Büchi automaton A that accepts intersection of  $L_1$  and  $L_2$ 

(c)  $L_1$ : It is certain that a is not infinite. However, b can be infinite.  $L_2$ : It is certain that if there is b at any position, regardless of finite or infinite words, there must

be a right next to it. However, a can be infinite. Yet, b can also be infinite as long as it follows a, which means both are infinite only if they are together.

 $L_1 \cap L_2$ :  $\emptyset$  (It is certain that a is not infinite. And b cannot be infinite since a is not infinite)

## Exercise 2: Transition Systems

(a) The model represents a case or lock which can be opened by inserting a ticket. Red light appears if locked. Otherwise, green light appears.

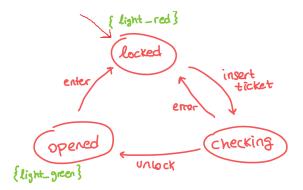


Figure 3: Transition System

(b) Mathematical definition of given transition system is as follows:

$$S = \{1, 2, 3, 4\}$$

$$Act = \{\text{close\_door}, \text{open\_door}, \text{go\_up}, \text{go\_down}\}$$

$$S_0 = \{4, 1\}$$

$$L(4) = \{\text{open}, \text{top\_floor}\}$$

$$L(3) = \{\text{top\_floor}\}$$

$$L(2) = \{\text{ground\_floor}\}$$

$$L(1) = \{\text{open}, \text{ground\_floor}\}$$

Transitions are as follows:

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\longrightarrow = \{(4, \operatorname{close\_door}, 3), (3, \operatorname{open\_door}, 4), (3, \operatorname{go\_down}, 2), \\ (2, \operatorname{go\_up}, 3), (2, \operatorname{open\_door}, 1), (1, \operatorname{close\_door}, 2)\}
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And the door is open at states 2 and 3.

## Exercise 3: Crossroads Traffic Lights

(a) The traffic lights do not synchronize with each other.

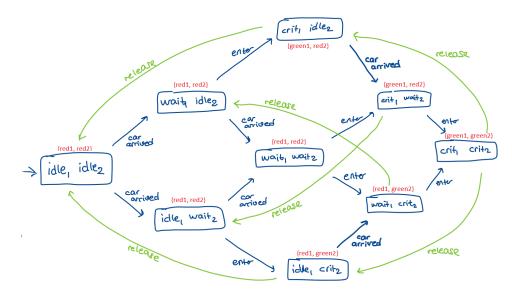


Figure 4: Interleaving transition systems  $TS_1 |||TS_2||$ 

(b) The traffic lights do not synchronize with each other, but they both synchronize with the arbiter.

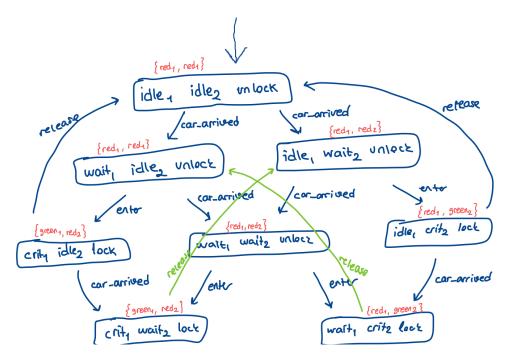


Figure 5: Parallel composition of systems  $(TS_1|||TS_2)$  || Arbiter

(c) Yes, the system is safe when synchronized with an Arbiter. There is no such a state where the atomic proposition is {green<sub>1</sub>, green<sub>2</sub>}.

However, in this design, the lights immediately switch from green to red. We would expect another state(s), e.g. a yellow light state, which would result in a slower and safer switch from

green to red.