CS 202, Spring 2022 Homework 1

Algorithm Efficiency

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Question 1

a) Show that $f(n) = 8n^4 + 5n^3 + 7$ is $O(n^5)$ by specifying appropriate c and n_0 values in Big-O definition.

We need to show that $f(n) = O(n^5)$ if \exists positive constants c, n_0 such that

(Eq 1)
$$0 \le f(n) \le cn^5, \forall n \ge n_0$$

If we choose c and n_0 such that c = 20 and $n_0 = 1$, then our equation becomes

(Eq 2)
$$0 \le f(n) = 8n^4 + 5n^3 + 7 \le 20n^5, \forall n \ge 1$$

(Eq 3)
$$0 \le f(n) = 8n^4 \le 8n^5, \forall n \ge 1$$

(Eq 4)
$$0 \le f(n) = 5n^3 \le 5n^5, \forall n \ge 1$$

(Eq 5)
$$0 \le f(n) = 7 \le 7n^5, \forall n \ge 1$$

If we add Eq (3-5), we get our equation Eq 2. Therefore, Eq 2 is correct and $f(n) = 8n^4 + 5n^3 + 7$ is $O(n^5)$.

- b) Trace the following sorting algorithms to sort the array [22, 8, 49, 25, 18, 30, 20, 15, 35, 27] in ascending order. Use the array implementation of the algorithms as described in the textbook and show all major steps.
- Selection sort
- Bubble sort

1) Selection Sort

• Initial: All Unsorted

22 8 49 25	18 30	20 15	35 27	
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• Largest -> 49: (49 - 27 swap)

Sorted

22	8	27	25	18	30	20	15	35	49

• Largest -> 35

| Sorted

22	8	27	25	18	30	20	15	35	49
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• Largest -> 30: (30 - 15 swap)

Sorted

П										
	99	Q	97	25	1Ω	1.5	20	20	25	4.0
П	44	0	47	4.5	10	13	40	30	33	43
- 1										4

• Largest -> 27: (27 - 20 swap)

Sorted

ī										
ı										
1	99	Ω	20	25	1.2	15	97	30	35	40
1	44	0	40	4.5	10	13	47	30	33	13
-1										

• Largest -> 25: (25 - 15 swap)

Sorted

|--|

• Largest -> 22: (22 - 18 swap)

	0 1
ı	Sorted

18	8	20	15	22	25	27	30	35	49

• Largest -> 20: (20 - 15 swap)

Sorted

10 0 13 20 22 23 27 30 33 19

• Largest -> 20: (18 - 15 swap)

Sorted

15	8	18	20	22	25	27	30	35	49

• Largest -> 15: (15 - 8 swap)

Sorted

8	15	18	20	22	25	27	30	35	49

2) Bubble Sort

• Initial: All Unsorted

22	8	49	25	18	30	20	15	35	27	

• Pass 1: 22 > 8: swap

• Pass 1: 49 > 25: swap

8	22	25	49	18	30	20	15	35	27
	44	4.5	13	10	30	40	10		

•	Pass	1:	49	>	18:	swap
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8 22 25 18	49 30	20 15	35 27	1
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• Pass 1: 49 > 30: swap

8 22 25 18	30 49	20 15	35 27	
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• Pass 1: 49 > 20: swap

• Pass 1: 49 > 15: swap

8 22 25 18	30	20	15	49	35	27

• Pass 1: 49 > 35: swap

8	22	25	18	30	20	15	35	49	27
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• Pass 1: 49 > 27: swap

8 22 25 18 30 20 15 35 27 49

• Pass 2: 25 > 18: swap

8	22	18	25	30	20	15	35	27	49
0	44	10	4.5	30	40	13	33	47	43

• Pass 2: 30 > 20: swap

8 22 18 25 2	30 15 35 27 49
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• Pass 2: 30 > 15: swap

	8	22	18	25	20	15	30	35	27	49
•	Pass 2	2: 35 > 27	: swap						Contad	
									Sorted	
L	8	22	18	25	20	15	30	27	35	49
•	Pass 3	3 : 22 > 18	: swap							
Г	8	18	22	25	20	15	30	27	35	49
•	Pass 3	3 : 25 > 20	: swap							
	8	18	22	20	25	15	30	27	35	49
										-
•	Pass 3	3 : 25 > 15	: swap							
	8	18	22	20	15	25	30	27	35	49
•	Pass 3	s: 30 > 27	: swap							
								Sorted		
	8	18	22	20	15	25	27	30	35	49
										_
•	Pass 4	: 22 > 20	: swap							
	8	18	20	22	15	25	27	30	35	49
•	Pass 4	e: 22 > 15	: swap							
							Sorted			
	8	18	20	15	22	25	27	30	35	49

• Pass 5: 20 > 15: swap

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1										
1	8	18	15	20	22	25	27	30	35	49
-										

• Pass 6: 18 > 15: swap

Sorted

8	3	15	18	20	22	25	27	30	35	49
		10					٦,	0 0	00	10

• Pass 7: No swap: Finish

Sorted

0	1.5	1.0	20	00	05	07	20	25	40
Ö	13	10	20	44	23	47	30	33	49

Question 2

Screenshots for the requested outputs of the program are listed below.

----Insertion Sort----Comparison Count: 69 Move Count: 88 [1,2,4,5,6,7,8,9,11,12,13,16,16,17,18,20] ----Bubble Sort----Comparison Count: 110 Move Count: 174 [1,2,4,5,6,7,8,9,11,12,13,16,16,17,18,20] ----Merge Sort-----Comparison Count: 47 Move Count: 128 [1,2,4,5,6,7,8,9,11,12,13,16,16,17,18,20] ----Quick Sort----Comparison Count: 50 Move Count: 116 [1,2,4,5,6,7,8,9,11,12,13,16,16,17,18,20]

Output 1. Array Sorting Results

	Insertion Sort Elapsed time		compCount		moveCount
Array Size 5000	15263 ms		compCount 6198993	6203997	llloveCount
10000	60343 ms		24984764	0203997	24994770
15000	132682 ms		56059223		56074233
20000	236596 ms		99610742		99630752
25000	368020 ms		156813750		156838753
30000	528064 ms		224861836		224891844
35000	724337 ms		307993705		308028712
40000 	941167 ms		399617471		399657477
Analysis of	Bubble Sort				
Array Size	Elapsed time		compCount		moveCount
5000	44602 ms		12495154		18581997
10000	207984 ms		49984269		74924316
15000	529845 ms		112468629		168132705
20000	996093 ms		199982374		298772262
25000	1613127 ms		312472794		470366265
30000	2368750 ms		449946219		674495538
35000	3297113 ms		612465664		923876142
40000	4335733 ms		799933640		1198732437
 Analysis of	Merge Sort				
Array Size	Elapsed time		compCount		moveCount
5000	732 ms	55250	123616		morecourre
10000	1476 ms	120511	267232		
15000	2341 ms	189259	417232		
20000	3141 ms	260865	574464		
25000	3821 ms	334163	734464		
30000	4747 ms	408755	894464		
35000	5528 ms	484664	1058928		
40000	6453 ms	561905	1228928		
 Analysis of	Ouick Sort				
Array Size	Elapsed time		compCount		moveCount
5000	449 ms	66435	113918		o t c c c c c c i i c
10000	971 ms	171886	277808		
15000	1469 ms	239528	386254		
20000	2036 ms	346591	568652		
25000 25000	2612 ms	441529	666306		
30000	3101 ms	514365	858243		
35000 35000	3726 ms	623024	1031365		
33000 40000	4347 ms	727982	1239837		
+V1V1V1V1	434/ IIIS	12/902	123983/		

Output 2. Performance Analysis for Random Array Scenario

Analysis of	Incortion Cort				
Analysis of Array Size 5000 10000 15000 20000 25000 35000 35000	Insertion Sort Elapsed time 1964 ms 7593 ms 16214 ms 28590 ms 44673 ms 67480 ms 91439 ms 115543 ms	805951 3142005	compCount 810950 3152004 6882767 12077141 18932441 28582649 38713069 49134849	6897766	12097140 18957440 28612648 38748068 49174848
					49174040
Analysis of Array Size 5000 10000 15000 20000 25000 30000 35000 40000	Bubble Sort Elapsed time 33199 ms 130333 ms 290717 ms 522758 ms 824692 ms 1192574 ms 1627187 ms 2118706 ms		compCount 12475345 49928570 112335980 199730440 312294369 449743835 611942220 798912009		moveCount 2402856 9396018 20603304 36171426 56722326 85657950 116034210 147284550
Analysis of I					
Array Size 5000 10000 15000 20000 25000 30000 35000 40000	Elapsed time 494 ms 1030 ms 1554 ms 2142 ms 2757 ms 3260 ms 3904 ms 4462 ms	50728 111862 174810 244943 310343 379756 451400 526799	compCount 123616 267232 417232 574464 734464 894464 1058928 1228928		moveCount
Analysis of Array Size			compCount		moveCount
5000 10000 15000 20000 25000 30000 35000 40000	Elapsed time 720 ms 1345 ms 2229 ms 3481 ms 3436 ms 7155 ms 7673 ms 8943 ms	219085 428898 681569 1122294 1027109 2515687 2716333 3161653	1294360 1846758		movecount

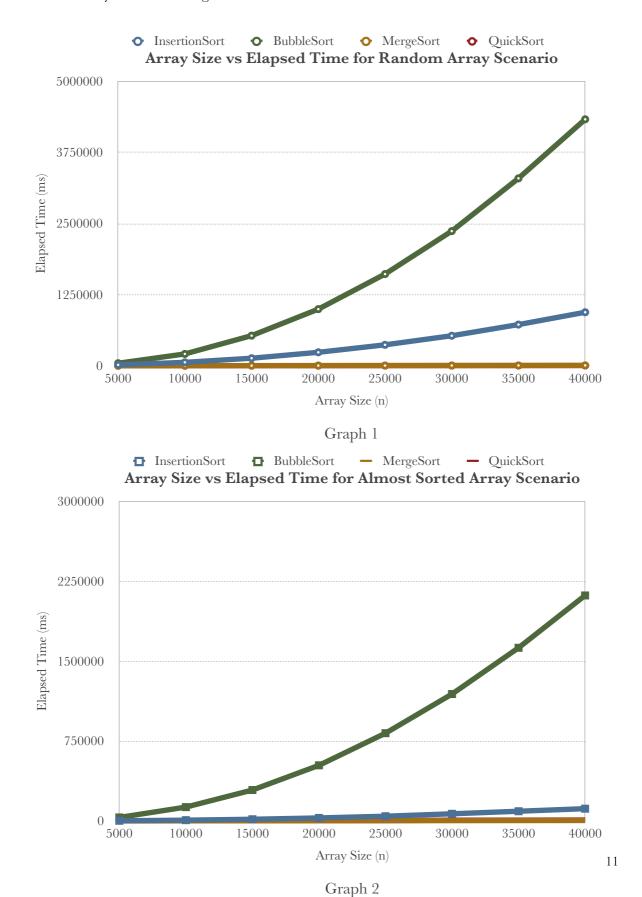
Output 3. Performance Analysis for Almost Sorted Array Scenario

	nsertion Sort			
Array Size	Elapsed time		compCount	moveCount
5000	28305 ms		11749486	11754526
10000	111034 ms		46965684	46975720
15000	251008 ms		105823860	105838912
20000	447707 ms		187548589	187568608
25000	684111 ms		292448938	292473952
30000	991281 ms 1368525 ms		421061959	421091966 574836776
35000			574801767	
40000 	1791252 ms		750541408	750581434
Analysis of B	ubble Sort			
Array Size	Elapsed time		compCount	moveCount
5000	57362 ms		12497500	35233584
10000	229604 ms		49995000	140867166
15000	516003 ms		112492500	317426742
20000	930673 ms		199990000	562585830
25000	1452823 ms		312487500	877271862
30000	2077711 ms		449985000	1263095904
35000	2842616 ms		612482500	1724300334
40000	3748185 ms		799979999	-2043462988
Analysis of M	 erge Sort			
Array Size	Elapsed time		compCount	moveCount
5000	490 ms	49075	123616	
10000	1071 ms	107353	267232	
15000	1569 ms	172032	417232	
20000	2179 ms	236660	574464	
25000	2694 ms	303203	734464	
30000	3273 ms	372575	894464	
35000	4110 ms	444875	1058928	
40000	4563 ms	510406	1228928	
 Analysis of Q	uick Sort			
Array Size	Elapsed time		compCount	moveCount
5000	963 ms	285230	447106	ovccourre
10000	1823 ms	547183	846278	
15000	3672 ms	1129081	1732508	
20000	2325 ms	640509	1024558	
25000	2807 ms	776297	1243358	
30000	3299 ms	918893	1411658	
35000	3772 ms	1036778	1639937	
40000	4946 ms	1390056	2179591	

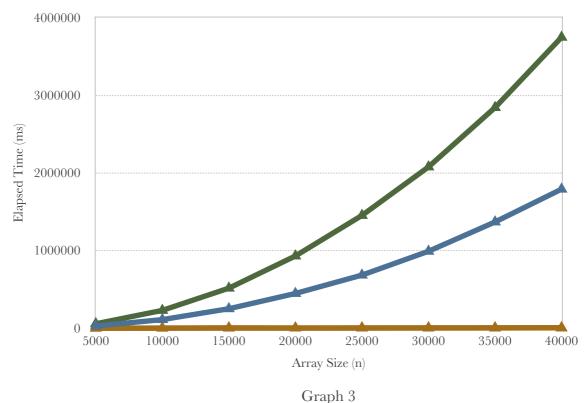
Output 4. Performance Analysis for Almost Unsorted Array Scenario

Question 3

Graphs requested for this part of the homework for random, almost sorted and almost unsorted array scenarios are given below.



▲ InsertionSort ▲ BubbleSort ▲ MergeSort ▲ QuickSort Array Size vs Elapsed Time for Almost Unsorted Array Scenario



While comparing these graphs and sorting algorithms, we need to keep in mind the following time complexities for these algorithms:

- Insertion Sort:
 - Best Case: O(n)
 - Worst Case: O(n²)
 - Average Case: O(n²)
- Bubble Sort:
 - Best Case: O(n)
 - Worst Case: O(n²)
 - Average Case: O(n²)
- Merge Sort:
 - Best Case: O(nlogn)
 - Worst Case: O(nlogn)
 - Average Case: O(nlogn)

Quick Sort:

• Best Case: O(nlogn)

• Worst Case: O(n²)

• Average Case: O(nlogn)

Comments:

- ✓ For the results of insertion sort algorithm, we expect approximately O(n) time complexity for almost sorted array, O(n²) time complexity for other scenarios (worst and average cases). If we observe the graphs, we can see the dramatic increase with array size for almost unsorted and random array scenarios. It supports theoretical expectations for worst and average cases. If we look at the graph for random array, we can see that the increase with array size in the elapsed time is more passive compared to the other graphs. However, its elapsed time is still larger than the merge sort and quick sort algorithms. This can be due to the fact that even though the array is almost sorted, it is not completely sorted. Therefore, this condition is not sufficient for the best case of insertion sort algorithm, and when we look at the comparison and move counts for this scenario, we can see that its counts are larger than merge sort and quick sort algorithm's.
- ✓ For the results of the bubble sort algorithm, we expect O(n) time complexity for the best case, O(n²) time complexity for the worst and average case. When we look at the results bubble sort algorithm is the worst algorithm in terms of runtime. Its comparison and move counts are very large compared to other sorting algorithms, and its time increases dramatically by array size. Even almost sorted array scenario has the same results. Therefore, it can be observed that even if array is slightly unsorted, bubble sort algorithm is badly affected by this and its time complexity increases. Therefore, we cannot observe O(n) time complexity.
- ✓ For the results of merge sort algorithm, we expect similar results, because theoretical expectation is O(nlogn) for the best, worst, and average cases. When we look at the results, move counts of merge sort algorithm remains same for all algorithms. It is normal because merge sort basically moves all array elements to other small arrays and moves them back into original array in order. Comparison counts change slightly, but it does not cause a dramatic change. Therefore, merge sort has similar appearance on all three graphs. Also, its cost is smaller than insertion and bubble sort, which are O(n²) sorting algorithms and similar to the quick sort (O(nlogn)) algorithm.
- ✓ For the results of quick sort algorithm, actually, we expect worst case O(n²) for almost sorted and almost unsorted array scenarios. Because in these scenarios pivot divides the array into two arrays sized 0 and n-1. However, its behavior is not what we expected. It is still very performant compared to bubble sort and insertion sort algorithms and its

performance does not dramatically change with array size. This may show that even if the array is slightly sorted or unsorted, quick sort algorithm still has a good performance. If we compare quick sort algorithm results with merge sort algorithm, we can see that its performance is worse than, even if they have close results, merge sort when array is almost sorted or unsorted. However, its performance is better than merge sort when a random array is used. This behavior is what we expected, because quick sort runs with O(nlogn) in average case like merge sort algorithm and it does not require as much moves or memory allocation as merge sort algorithm.