

On the complexity of unconditional convex bodies

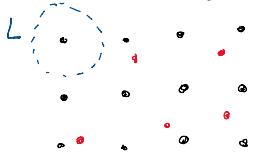
Thursday, April 15, 2021

Borodinov and Veomett question: Can any n -dimensional convex body be approximated by the projection of a section of an n -dimensional simplex with controlled number of facets?

■ Rekha Thomas (on the previous week)

$$\begin{array}{l} \xrightarrow{\text{Lifts}} \\ \xrightarrow{\text{Cross-Polytope}} \end{array} \left\{ \begin{array}{l} C_n = \{x \in \mathbb{R}^n : \sum_{i=1}^n |x_i| \leq 1\} \rightarrow 2^n \text{ facets} \\ Q_n = \{(x, y) \in \mathbb{R}^{2n} : \sum_{i=1}^n y_i = 1, -y_i \leq x_i \leq y_i \forall i\} \Rightarrow C_n = \pi_x(Q_n) \text{ for } \pi_x(x, y) = x. \\ \downarrow 2n \text{ facets.} \end{array} \right.$$

Answer: Not in general



■ Litvak+Rudelson+Tomczak-Jaegermann '14: $\exists L \in \mathcal{K}^n, L = -L$, s.t. $\forall K \in \mathcal{K}^n$ produced as a projection of a section of S in \mathbb{R}^N

$$d(L, K) \geq c \sqrt{\frac{n}{\log \frac{2N \log 2N}{n}}} , \quad c \text{ abs. constant.}$$

To prove this they estimated the complexity of the set of all Banach spaces

Pisier '13: \uparrow Such complexity is doubly exponential in n .

- See also survey by Mankiewicz+Tomczak-Jaegermann (Quotients of finite-dimensional Banach spaces)

■ Rudelson '16:

→ Unconditional

→ Completely symmetric

Def: Let (X, d) be a metric space, $\varepsilon > 0$, then

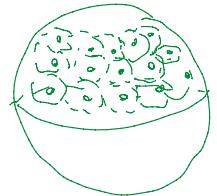
- $N_\varepsilon \subset X$ is an ε -net if $\forall x \in X, \exists y \in N_\varepsilon : d(x, y) \leq \varepsilon$.

• If X is compact then the minimal cardinality, $N(X, \varepsilon)$, of a net is called the covering number

- $\log_2 N(X, \varepsilon)$ is the metric entropy of X . In geometric functional analysis

In some sense it measures the complexity of the metric space.

Nets are a convenient way to discretize compact sets



E.g. (Covering number of S^n):

Let $\varepsilon > 0$, and N_ε a maximal ε -separated set ($\forall x, y \in N_\varepsilon, d(x, y) > \varepsilon$ and no other subset of N_ε has this property)

- N_ε is a net: If $\exists x \in X$ s.t. $\forall y \in N_\varepsilon d(x, y) \leq \varepsilon \Rightarrow N_\varepsilon \cup \{x\}$ is ε -separated.

$$\bullet B_2^n(x, \frac{\varepsilon}{2}) \cap B_2^n(y, \frac{\varepsilon}{2}) = \emptyset \quad \forall x, y \in N_\varepsilon, x \neq y,$$

$$\bullet B_2^n(x, \frac{\varepsilon}{2}) \subset (1 - \frac{\varepsilon}{2}) B_2^n$$

$$\Rightarrow \left(\frac{\varepsilon}{2}\right)^n |B_2^n| \cdot |N_\varepsilon| = \left|\frac{\varepsilon}{2} B_2^n\right| \cdot |N_\varepsilon| \leq \left|(1 + \frac{\varepsilon}{2}) B_2^n\right| = \left(1 + \frac{\varepsilon}{2}\right)^n |B_2^n|$$

$$\Rightarrow |N_\varepsilon| \leq \left(1 + \frac{2}{\varepsilon}\right)^n$$

E.g. (Spectral norm on a net):

Let A be an $N \times n$ matrix, N_ε a net for S^{n-1} , $\varepsilon \in (0, 1)$. Then,

- $\|A\| := \sup_{x \in S^{n-1}} \{ \|Ax\|_2\} \Rightarrow \max_{x \in N_\varepsilon} \{ \|Ax\|_2\}$

- Let $y \in S^{n-1}$: $\|Ay\|_2 = \|A^T y\|_2 \Rightarrow \exists x \in N_\varepsilon : \|x - y\|_2 \leq \varepsilon$

$$\Rightarrow \|Ax - Ay\|_2 \leq \|A\| \cdot \|x - y\|_2 \leq \varepsilon \cdot \|A\|.$$

Thus,

$$\|A\| \geq \|A\| \cdot \|A^T y\|_2 \geq \|A\| - \varepsilon \cdot \|A\| = (1 - \varepsilon) \|A\|$$

Thus,

$$\|Ax\|_2 \geq \|A\bar{x}\|_2 - \|A(x - \bar{x})\|_2 \geq \|A\bar{x}\| - \varepsilon \|A\bar{x}\| = (1-\varepsilon) \|A\bar{x}\|.$$

Then,

$$\max_{x \in N_\varepsilon} \{\|Ax\|_2\} \leq \|A\bar{x}\| \leq (1-\varepsilon) \max_{x \in N_\varepsilon} \{\|Ax\|_2\}$$

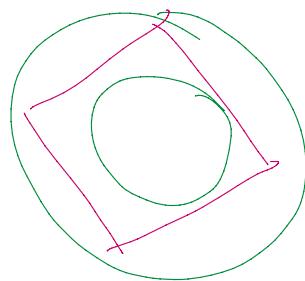
Thus, instead of taking supremum over the whole sphere to compute the spectral norm, we can estimate it on a net.

Def: Let X, Y be two n -dim normed spaces,

Geometric $d_{BM}(X, Y) = \inf \left\{ \frac{R}{r} : rB_Y \subseteq TB_X \subseteq RB_Y, \text{ for some } T \in GL(n) \right\}$

unit balls of
Y and X

Analytic $d_{BM}(X, Y) = \inf \left\{ \|T\| \cdot \|T^{-1}\| : T: X \rightarrow Y \text{ isomorphism} \right\}$



The Banach-Mazur distance defines a metric on a set of equivalent classes of normed spaces. It shows how far the spaces are from being isometric.

Remarks:

- ① $d(X, Y) \geq 1$
- ② $d(X, Y) = d(Y, X)$
- ③ $d(X, Z) \leq d(X, Y) d(Y, Z)$

Resembles a metric

- ④ Log d_{BM} defines a metric on $B_n \leftarrow$ set of all n -dim Banach spaces.
- ⑤ (B_n, d_{BM}) is a compact metric space \leftarrow Banach Mazur Compactum

John '48: $d(X, \ell_2^n) \geq \sqrt{n}, \forall x \in B_n$. In particular,

$$d(p_1^n, p_2^n) \quad \sqrt{n} ? \leq d(x, y) \leq n \quad \forall x, y \in B_n$$

↳ By multiplicative triangle inequality.

Gluskin '81: $\exists c > 0$, abs. constant, $\forall n \in \mathbb{N}$, $\exists x, y \in B_n$ with $d(x, y) \geq cn$.

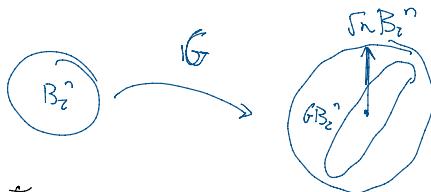
Remark: diameter of $B_n = n$.



Gluskin's Idea:

Let $G = [g_1, \dots, g_n]$ Gaussian matrix, $g_i \sim N(0, I_n)$ $\forall i \in [n]$, $g_i = (Y_{i1}, \dots, Y_{in})$ for $Y_{ij} \sim N(0, 1)$.
smallest scaling needed of the ball

- $\mathbb{E} \|G: \ell_1^n \rightarrow \ell_2^n\| = \mathbb{E} \sup_{x \in \ell_1^n} \|Gx\|_2 \approx \sqrt{n}$



For the lower bound we plug a vector, and $\mathbb{E} \|gy\|_2 \approx \sqrt{n}$

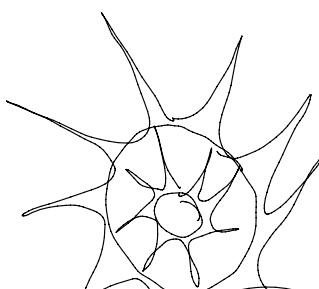
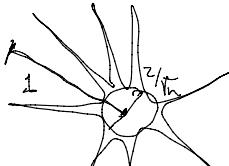
For the upper bound, we discretize on the sphere as on the previous example and we use Gaussians decay very fast (concentration),

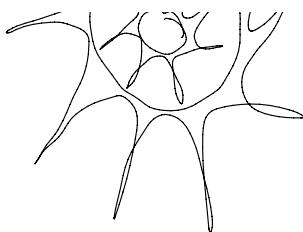
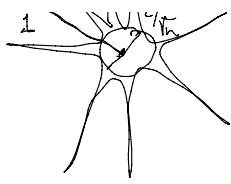
- $\mathbb{E} \|G: \ell_1^n \rightarrow \ell_2^n\| = \mathbb{E} \sup_{\substack{i \in [n] \\ \text{by boundary is} \\ \text{on a vertex}}} \|Ge_i\|_1 \geq n \approx n$

\uparrow
 $n\sqrt{\frac{2}{\pi}}$

Idea! Put the randomness in the space!

what if we throw some extra points to ℓ_1^n ?





$$\cdot |\mathcal{B}_1^n| \approx \frac{c}{n}$$

$$\cdot |\mathcal{B}_{\frac{1}{\sqrt{n}}}^n| = \left[\left(\frac{1}{\sqrt{n}} \right)^n |\mathcal{B}_1^n| \right]^{\frac{1}{n}} \approx \frac{1}{\sqrt{n}} \cdot \frac{c}{\sqrt{n}} \approx \frac{c}{n}$$

The feature of these bodies is that they have a bulk that makes up most of the volume, with little diameter, and spikes with little volume but big diameter.

Gaussian Shows: $X_m, Y_m \in \mathbb{R}^n$ s.t. $\mathcal{H}\mathcal{T}$ isomorphism

$$\begin{cases} \|T: X_m \rightarrow Y_m\| \geq c\sqrt{n} \\ \|T: Y_m \rightarrow X_m\| \geq c\sqrt{n} \end{cases} \Rightarrow \mathcal{C}_{BM}(X_m, Y_m) \geq c \cdot n$$

Outline of the construction:

$$\mathcal{B}_{X_m} = \text{Conv} \left\{ \underbrace{\frac{g_1}{\sqrt{n}}, \dots, \frac{g_n}{\sqrt{n}}}_{\text{random points}}, \underbrace{e_1, \dots, e_n}_{\text{deterministic}} \right\} \text{ for } g_1, \dots, g_n \sim N(0, I_n)$$

$$\mathcal{B}_{Y_m} = \text{Conv} \left\{ \underbrace{\frac{g'_1}{\sqrt{n}}, \dots, \frac{g'_n}{\sqrt{n}}}_{?}, e_1, \dots, e_n \right\}$$

- We divide by \sqrt{n} so we're roughly on the sphere, just normalization.
- If we didn't add $\pm e_i$'s we'd need the measure of the ball of all Gaussian vectors. Including the ℓ_1 ball automatically gives us a ball of radius $\frac{1}{\sqrt{n}}$.
- In particular, as we need to deal with all operators from one space to the other, we want to be able to consider one not far off them.

The identity:

by independence

$$= \mathbb{E} \left[\left(\frac{g_1}{\sqrt{n}}, \dots, \frac{g_n}{\sqrt{n}}, e_1, \dots, e_n \right)^\top \right] = \mathbb{E} \left[\left(\frac{g_1}{\sqrt{n}}, \dots, \frac{g_n}{\sqrt{n}} \right)^\top \right]$$

$$\begin{aligned} P(B_{X_n} \subseteq RB_{Y_n}) &\leq P\left(\bigcup_{i=1}^{2n} RB_{Y_n} \cap \{x_i\}\right) \stackrel{\text{by independence}}{\leq} \prod_{i=1}^{2n} P(g_i \in RB_{Y_n}) \\ &\stackrel{\text{cover set so include endpoints}}{=} \left[P(g_i \in RB_{Y_n})\right]^{2n} \\ &\stackrel{\text{identically dist.}}{=} p^{2n} \end{aligned}$$

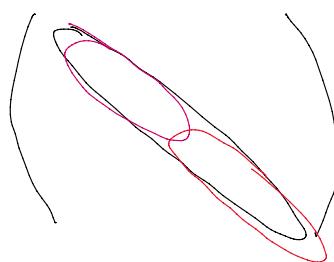
$$P(g_i \in RB_{Y_n}) \leq \frac{1}{(2\pi)^{n/2}} \int_{RB_{Y_n}} e^{-\frac{\|x\|^2}{2}} dx \stackrel{\text{bound density by } l}{\leq} C_n |RB_{Y_n}| = C_n \frac{(lR)^n}{n^n}$$

even when R is of the order of \sqrt{n}
the probability is small.

Initially $|B_{Y_n}| \approx \frac{C}{n}$, the set without the Gaussians is the ℓ_1^n -ball which has small volume, if we add the Gaussians, and they are on the sphere, then adding lots of them is going to have small volume.

$$\text{Say now } T \neq I \Rightarrow \text{SVD } T = UDV^T$$

We're going to assume half of the diagonal values are big and half are small, since when considering T^{-1} we take the reciprocals and the same follows for T^{-1}



In general if T is very degenerate it'd be hard to have a good small ball, but on the inverse we'd have a very good one. So then we just order them and scale the the middle ones are roughly equal.

$$\begin{aligned} T B_{X_n} \subseteq RB_{Y_n} &\Rightarrow P_T B_{X_n} \subseteq P_E B_{Y_n} \Rightarrow P_E T g_i \in \cap P_E RB_{Y_n} \Rightarrow T P_F g_i \in \cap P_E B_{Y_n} \\ &\stackrel{\text{inside the other}}{\stackrel{P}{\Rightarrow}} \stackrel{\text{holds for proj}}{\stackrel{P}{\Rightarrow}} \stackrel{\text{Suppose } T \text{ is}}{\stackrel{P}{\Rightarrow}} \stackrel{\text{diagonal}}{\stackrel{P}{\Rightarrow}} \end{aligned}$$

$$\Rightarrow P_F g_i \in \cap R T^{-1} P_E B_{Y_n} \Rightarrow P_F g_i \in \cap R P_E B_{Y_n}$$

Projection of a gaussian
is gaussian

$\hat{T} = T|_F$

Thus,

Thus,

- $\mathbb{P}(TB_{X_n} \subseteq RB_{Y_n}) \leq \prod_{i=1}^{2n} \mathbb{P}(P_F q_i \in \cap P_E B_{Y_n}) = |\cap P_E B_{Y_n}|^{2n} \leq \left(\frac{n}{c}\right)^{2n} = c^{-n^2}$
 - $|N_1| \leq C^n$

$P(3T : TB_{kn} \subset RB_{kn}) \leq c_1^{n^2} \cdot C^{n^2} \approx c_2^{n^2}$ since we cook up the constants to make some decay

Main Result: Let $1 \leq \ell \leq \tilde{C} n^{1/2} \log^{-5/2} n$, $\mathcal{X}_n^{\text{unc}}$ contains a ℓ -separated set of cardinality at least

$$\exp(\exp(\frac{c}{\epsilon^2 \log(\epsilon^{2+})} n)))$$

Corollary: Let $n \in \mathbb{N}$, $\exists L \in K_n^{\text{loc}}$ s.t. $\forall K \in \mathbb{K}^n$ projection of section of simplex then C_{R^n}

$$d(L, k) \geq c \left(\frac{n}{\log n} \right)^{1/4} \cdot \log^{-1} \left(\frac{n}{\log n} \right)$$

\uparrow
 abs. const.

Def: Let $x_1, \dots, x_m \in \mathbb{R}^n$ with $x_j = (x_{j1}, \dots, x_{jn})$ then

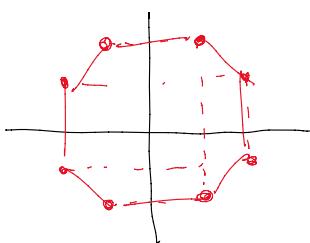
$$\text{Var. Cov} \{x_1, \dots, x_m\} = \text{Cov} \{(\varepsilon_1, x_{11}, \dots, \varepsilon_1 x_{1m}), \dots, (\varepsilon_m, x_{m1}, \dots, \varepsilon_m x_{mn})\}$$

where $E_{ij} = \pm 1$,

Rudelssohn's

random sets

$$K = K(I_1, \dots, I_K)$$



$$= \text{abs. Conv} \{ \cup_{n \in \mathbb{N}} \text{Conv} \{ x_1, \dots, x_N \}, \sqrt{\varepsilon_n} B_1^n, \sqrt{\varepsilon_n} B_2^n \}$$

$$= \text{abs}(\text{Conv}\{ \text{vec}(B_1), \dots, \text{vec}(B_N) \})$$

deterministic

T_p 's ind. random sets in $[n]$

uni. chosen among sets with cardinality s_n , $s > 0$.

$$x_p = \left[e_j \leftarrow \begin{array}{l} \text{coordinates} \\ j \in T_p \end{array} \right. \begin{array}{l} \text{are not independent} \end{array} \right]$$

$$\Pr\left(\bigcap_{j \in S \subseteq [n]} e_j \in R_{S, k'}\right) \leq \exp(-c s_n^2) \quad \leftarrow 4.1$$

$$\rightarrow \Pr(T_k \subseteq R_{S, k'}) \leq \exp(-\exp(-c s_n^2)) \quad \leftarrow 4.2.$$

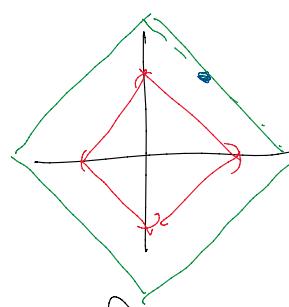
ε 's symmetric ± 1 Bernoulli random variables $\sum \xrightarrow{\text{ns}}$ Decays as Gaussian

J random subset with cardinality s_n .
 $\subseteq [n]$

Intuition: $\Pr(Y \in \text{Conv}\{B_1, \dots, B_N\})$
 Y is o's B_i 's o's ^{Collision spikes} coordinates misalign so we need many

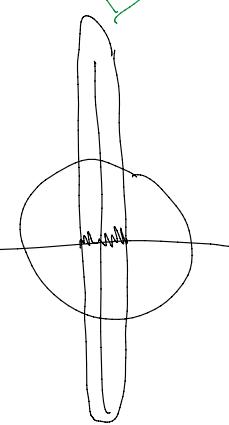
$$\Pr(g \in E) = \Pr(g \in TB_2^n) = \Pr(T^{-1}g \in B_2^n)$$

but $T^{-1}g = \langle T^{-1}g, T^{-1}g \rangle$ T diagonal ^{Hanson-Wright} ineq.



Key idea: we must consider the geometry of tangent space.

\Rightarrow we only treated the $\text{Conv}\{B_1, \dots, B_N\}$ +



$$\Pr(T^{-1}g \in B_2^n) \leq \Pr(T^{-1}g \in B_2^n) \approx \Pr(T^{-1}g \in B_2^n) \approx \Pr(T^{-1}g \in B_2^n)$$

$$\begin{aligned}
 \Pr(TK \subseteq R_s K') &\leq \prod_{j=1}^N \Pr(T \cap e_j \in R_s K') = [\exp(-c's^n)]^N \\
 &= \exp(-N c's^n) = \exp(-c's^n \exp(c's^n)) \\
 &\leq \exp(-\exp(c's^n))
 \end{aligned}$$

$$|N| \leq C^n$$

$$\begin{aligned}
 \Pr(\exists T : TK \subseteq R_s K') &\leq |N| \exp(-\exp(c's^n)) \\
 &\leq \exp(-\exp(c's^n))
 \end{aligned}$$

Thus,

$$\exp(\exp(\frac{cn}{(2d)^4 \log(2d+1)})) \leq \# \left\{ \begin{array}{l} \text{set of uncond. as } c \\ \text{projection of section } d \end{array} \right\} \leq 3 \log n$$

↑ set of uncond. as c
 Rudelson's theorem
 ↑ projection of section d
 ↑ LRTJ's result,