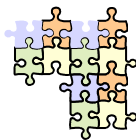


# Game Trees and Heuristics

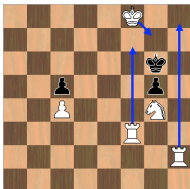
15-211: Fundamental Data Structures and Algorithms

Margaret Reid-Miller  
7 April 2005

2



## Games



## Announcements

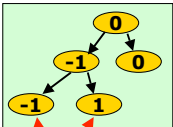
- **HW5** Due Monday 11:59 pm

## Game Properties

- Two players
- Deterministic
  - no randomness, e.g., dice
- Perfect information
  - No hidden cards or hidden Chess pieces...
- Finite
  - Game must end in finite number of moves
- Zero sum
  - Total winnings of all players

4

## MiniMax Algorithm

- Player A (us) **maximize** goodness.
  - Player B (opponent) **minimizes** goodness
- Player A          **maximize** (draw)
- Player B     **minimize** (lose, draw)
- Player A     **maximize** (lose, win)
- At a leaf (a **terminal** position) A wins, loses, or draws.
    - Assign a score: 1 for win; 0 for draw; -1 for lose.
  - At **max** layers, take node score as **maximum** of the child node scores
  - At **min** layers, take nodes score as **minimum** of the child node scores

5

## Heuristics

- A heuristic is an approximation that is typically fast and used to aid in optimization problems.
- In this context, heuristics are used to "rate" board positions based on local information.
- For example, in Chess I can "rate" a position by examining who has more pieces. The difference in black's and white's pieces would be the evaluation of the position.

6

## Heuristics and Minimax

- When dealing with game trees, the heuristic function is generally referred to as the **evaluation function**, or the static evaluator.
- The static evaluation takes in a board position, and gives it a score.
- The higher the score, the better it is for you, the lower, the better for the opponent.

7

## Heuristics and Minimax

- Each layer of the tree is called a **ply**.
- We cut off the game tree at a certain maximum depth,  $d$ . Called a **d-ply** search.
- At the bottom nodes of the tree, we apply the heuristic function to those positions.
- Now instead of just Win, Loss, Tie, we have a score.

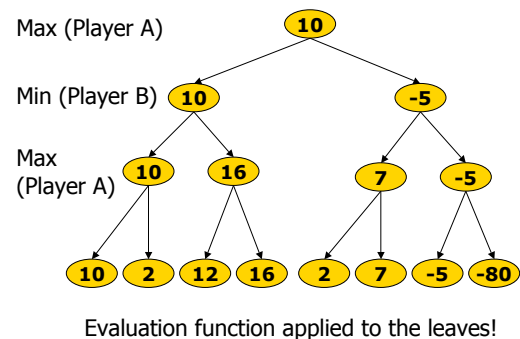
8

## Evaluation Function

- Guesses the outcome of an incomplete search
- $\text{Eval}(\text{Position}) = \sum_i w_i * f_i(\text{Position})$
- Weights  $w_i$  may depend on the phase of the game
- Features for chess  $f_i$ :
  - #of Pawns (material terms)
  - Centrality
  - Square control
  - Mobility
  - Pawn structure

9

## Minimax in action



10

## How fast?

- Minimax is pretty slow even for a modest depth.
- It is basically a brute force search.
- What is the running time?
  - Each level of the tree has some average  $b$  moves per level. We have  $d$  levels. So the running time is  $O(b^d)$ .

11

## Alpha Beta Pruning

- Idea: Track "window" of expectations. Two variables:
  - $\alpha$  – Best score so far at a **max** node: increases.
    - Score can be forced on our opponent.
  - $\beta$  – Best score so far at a **min** node: decreases
    - Opponent can force a situation no worse than this score. Any move with better score is too good for opponent to allow
- Either case: If  $\alpha \geq \beta$ :
  - Stop searching further subtrees of that child. Opponent won't let you get that high a score.
- Start the process with an infinite window ( $\alpha = -\infty, \beta = \infty$ ).

12



## Timed moves

- Not uncommon to have a limited time to make a move. May need to produce a move in say 2 minutes.
- How do we ensure that we have a good move before the timer goes off?

19

## Iterative Deepening

- Evaluate moves to successively deeper and deeper depths:
  - Start with 1-ply search and get best move(s). **Fast**
  - Next do 2-ply search using the previous best moves to order the search.
  - Continue to increased depth of search
- If some depth takes too long, fall back to previous results (timed moves).

20

## Iterative Deepening

- Save best moves of shallower searches to control move ordering.
- Need to search the same part of the tree multiple times but improved move ordering more than makes up for this redundancy
- Difficulty:
  - Time control: each iteration needs about 6x more time
  - What to do when time runs out?

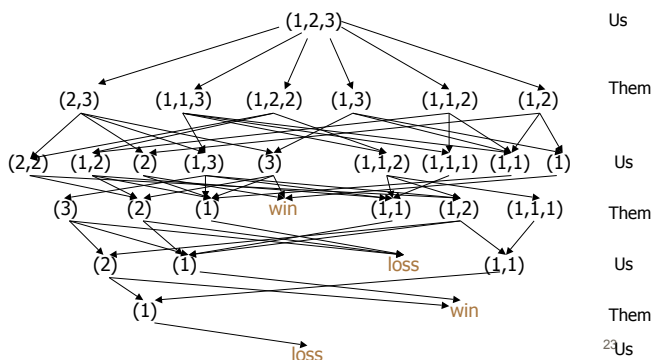
21

## Transposition Tables

- Minimax and Alpha Beta (implicitly) build a tree. But what is the underlying structure of a game?
- Different sequences of moves can lead to the same position.
- Several game positions may be functionally equivalent (e.g. symmetric positions).

22

## Nim Game Graph



Us

Them

Us

Them

Us

Them

Us

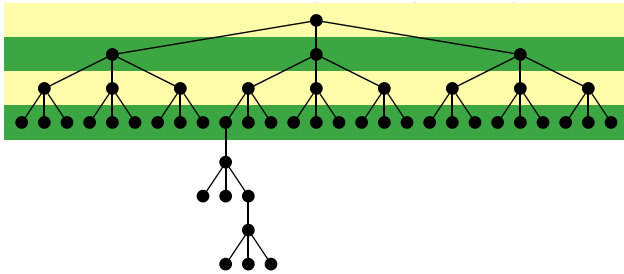
## Transposition Tables

- Memoize:** hash table of board positions to get
  - Value for the node
    - Upper Bound, Lower Bound, or Exact Value
    - Be **extremely careful** with alpha beta as may only know a bound at that position.
  - Best move at the position
    - Useful for move ordering for greater pruning!
- Which positions to save?
  - Sometimes obvious from context
  - Ones in which more search time has been invested
  - Collisions: Simply overwrite

24

## Limited Depth Search Problems

- Horizon effect: push bad news over the search depth



25

## Alpha Beta Pruning

**Theorem:** Let  $v(P)$  be the value of a position  $P$ . Let  $X$  be the value of a call to  $AB(\alpha, \beta)$ . Then one of the following holds:

$$\begin{aligned} v(P) \leq \alpha \text{ and } X \leq \alpha \\ \alpha < v(P) < \beta \text{ and } X = v(P) \\ \beta \leq v(P) \text{ and } X \geq \beta \end{aligned}$$

Suppose we take a chance and reduce the size of the infinite window. What might happen?

26

## Aspiration Window:

- Suppose you had information that value of a position was probably close to 2 (say from the result of shallower search)
- Instead of starting with an infinite window, start with an "aspiration window" around 2 (e.g., (1.5, 2.5)) to get more pruning.
  - If the result is in that range you are done.
  - If outside the range you don't know the exact value, only a bound. Repeat with a different range.
- How might this technique be use for parallel evaluation?

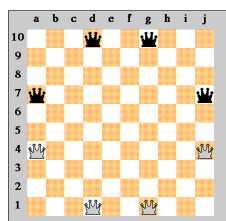
27

## Tricks

- Many tricks and heuristics have been added to chess program, including this tiny subset:
  - Opening Books
  - Avoiding mistakes from earlier games
  - Endgame databases (Ken Thompson)
  - Singular Extensions
  - Think ahead
  - Contempt factor
  - Strategic time control

28

## Game of Amazons



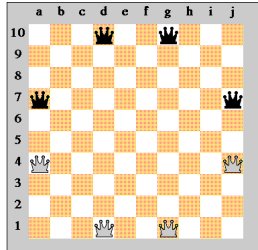
## Game of Amazons

- Invented in 1988 by Argentinian Walter Zamkaskas.
- Several programming competitions and yearly championships.
- Distantly related to Go.
- Active area in combinatorial game theory.

30

## Amazons Board

- Chess board 10 x 10
- 4 White and 4 Black chess Queens (Amazons) and Arrows
- Starting configuration
- White moves first



31

## Amazons Rules

- Each move consists of two steps:
  1. Move an amazon of own color.
  2. This amazon has to throw an arrow to an empty square where it stays.
- Amazons and arrows move as a chess Queen as long as no obstacle blocks the way (amazon or arrow)
- Players alternate moves.
- Player who makes last move wins.

32

## Amazons challenges

- Absence of opening theory
- Branching factor of more than 1000
- Often 20 reasonable moves
- Need for deep variations
- Opening book >30,000 moves

33

## AMAZONG

- World's best computer player by Jens Lieberum
- Distinguishes three phases of the game:
  - Opening at the beginning
  - Middle game
  - Filling phase at the end

See

<http://jenslieberum.de/amazong/amazong.html>

34

## Amazons Opening

- Main goals in the opening
  - Obtain a good distribution of amazons
  - Build large regions of potential territory
  - Trap opponent's amazons

35

## Amazons Filling Phase

- Filling phase starts when empty squares can be reached by only one player.
- Happens usually after 50 moves.
- Goal is to have access to more empty squares than the other player
- Outcome of game determined by counting number of moves left for each player.
- But...

36

## Defective Territories

- K-defective territory provides k fewer moves than empty squares

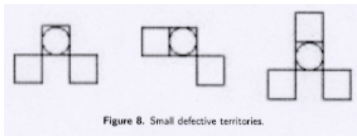
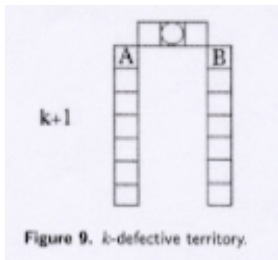


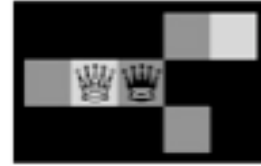
Figure 8. Small defective territories.



**Figure 9.** *k*-defective territory

## Zugzwang

- Seems that black has access to 3 empty squares
- But if black moves first then can only use two.



## More Complex Zugzwang

Player who moves first must either

- take their own region and give region C to the opponent, or
- take region C and block off their own region

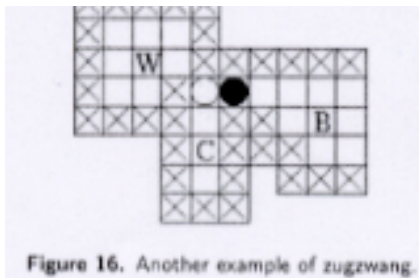


Figure 16. Another example of zugzwang.

## Chiptest, Deep Thought Timeline

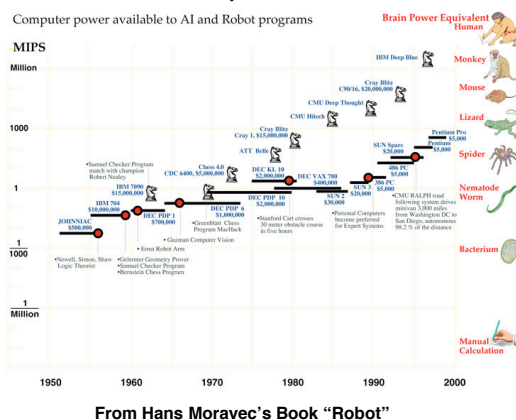


**Story continues  
@ IBM**

- A VLSI design class project evolves into F.H.Hsu move-generator chip
- A productive CS-TG results in ChipTest, about 6 weeks before the ACM computer chess championship
- Chiptest participates and plays interesting (illegal) moves
- Chiptest-M wins ACM CCC
- Redesign becomes DT
- DT participates in human chess championships (in addition to CCC)
- DT wins second Fredkin Prize (\$10K)
- DT is wiped out by Kasparov

## Opinions: Is Computer Chess AI?

- Computer power available to AI and Robot programs



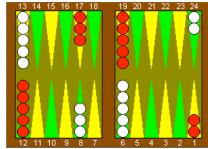
From Hans Moravec's Book "Robot"

So how does Kasparov win?

- Even the best Chess grandmasters say they only look 4 or 5 moves ahead each turn. Deep Junior looks up about 18-25 moves ahead. How does it lose!?
- Kasparov has an unbelievable evaluation function. He is able to assess strategic advantages much better than programs can (although this is getting less true).
- The moral, the evaluation function plays a large role in how well your program can play.

## State-of-the-art: Backgammon

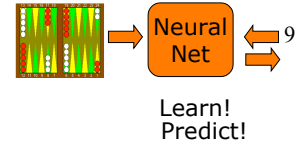
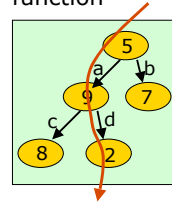
- Gerald Tesauro (IBM)
- Wrote a program which became “overnight” the best player in the world
- Not easy!



43

## State-of-the-art: Backgammon

- Learned the evaluation function by playing 1,500,000 games against itself
  - Temporal credit assignment using reinforcement learning
  - Used Neural Network to learn the evaluation function



44

## State-of-the-art: Go

- Average branching factor 360
- Regular search methods go bust !
- People use higher level strategies
- Systems use vast knowledge bases of rules... some hope, but still play poorly
- \$2,000,000 for first program to defeat a top-level player

45