### LoRa Quake

2023 Alphan Şahin

#### Acknowledgement

- This research is supported by "TUBITAK 2221 Call for Earthquake Research Opens within the Scope of Supporting Scientist with Guest or Academic Leave Program" between June and July 2023
- The research is conducted at CoSiNC 6G research lab directed by Prof. Hüseyin Arslan, Istanbul Medipol University
  - Many thanks to the students and staff at CoSiNC family!







### **Agenda**

- Part 1: Motivation & Challenges
- Part 2: Proposed solutions
- Part 3: Demonstration

### **Motivation & Challenges (1/2)**

- February 6th Kahramanmaraş Earthquakes (2023 Turkey–Syria Earthquake) showed that it is very difficult to establish timely contact and locate those trapped under the rubble quickly
- Challenge: Even if a trapped person has his/her cell phone nearby, establishing a connection through cellular networks is quite challenging:
  - Significantly attenuated radio signals due to the debris (30-60 dB [1])
  - Damaged, inaccessible, and/or congested cellular networks due to the earthquake

[1] Holloway, C., Koepke, G., Camell, D., Remley, C., Schima, S., McKinley, M. and Johnk, R. (2006), Propagation and Detection of Radio Signals Before, During, and After the Implosion of a Large Convention Center, Technical Note (NIST TN), National Institute of Standards and Technology, Gaithersburg, MD

### **Motivation & Challenges (2/2)**

- Establishing contact largely depends on rescue teams and their equipment
  - Methods using heat, sound, or CO2 levels [1]
  - Systems based on electromagnetic or acoustic waves using radar principles
    - Ground penetration radar (e.g., NASA FINDER working at 3 GHz band) [2], [3]
- Challenge: Detection via systems deployed from the surface into the buildings' interiors is limited due to the weakened signals
  - We need innovative communication and sensing technologies to significantly reduce our losses in future disaster situations

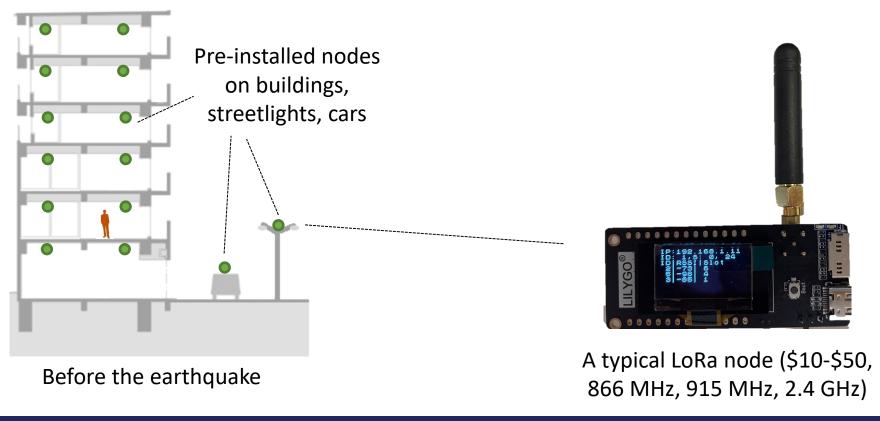
<sup>[1]</sup> Zhang, Di, Salvatore Sessa, Ritaro Kasai, Sarah Cosentino, Cimarelli Giacomo, Yasuaki Mochida, Hiroya Yamada, Michele Guarnieri, and Atsuo Takanishi. 2018. "Evaluation of a Sensor System for Detecting Humans Trapped under Rubble: A Pilot Study" Sensors 18, no. 3: 852.

<sup>[2]</sup> NASA, https://spinoff.nasa.gov/FINDER-Finds-Its-Way-into-Rescuers-Toolkits, Last accessed: 01/31/2023

<sup>[3]</sup> V. Cable, J. Lux and S. Haque, "Target & propagation models for the FINDER radar," 2013 IEEE Antennas and Propagation Society International Symposium (APSURSI), Orlando, FL, USA, 2013, pp. 1614-1615

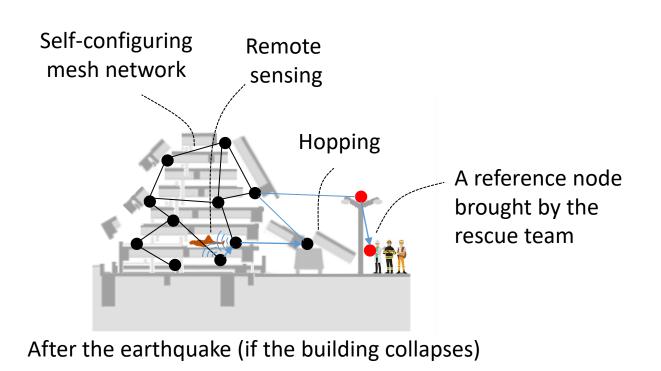
#### **Proposed Approach**

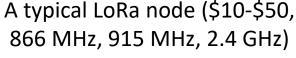
- Deploy low-cost Long-Range (LoRa) nodes with sensing capabilities to the buildings and establish a mesh network along with the reference nodes brought by the rescue teams if the building collapses
  - Nodes sense their environments and convey sensing results to the reference nodes



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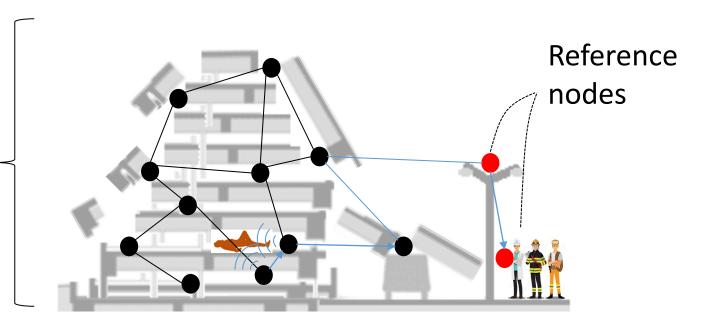




#### Goals

- Goal 1: Establishing a self-configuring mesh network
  - How to deal with interference problems without a coordinator?
- Goal 2: Transferring messages from nodes to the rescue teams
  - How to decide where to send the message in the mesh?
- Goal 3: Finding the positions of the nodes to find the trapped people
  - How to locate the nodes?

- 1) Interference?
- 2) Reaching reference nodes? —
- 3) Localization?

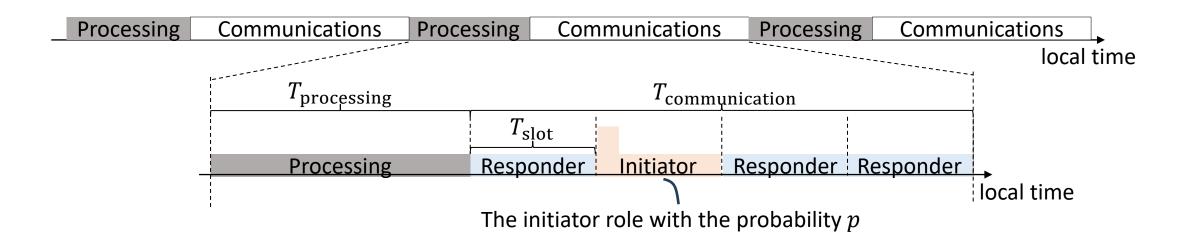


### Solutions

Goal 1: Establishing a self-configuring mesh network

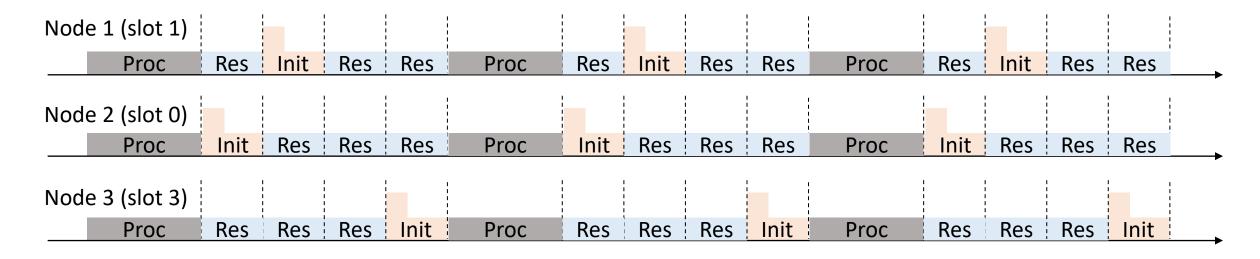
#### **Periodically Alternating States**

- A typical LoRa device equipped with a low-cost microprocessor
  - Cannot handle communication (TX and RX) and processing tasks simultaneously
- We propose to alternate processing and communication periodically, further divide communications into  $N_{\rm slot}$  slots, and define three states:
  - 1) Processing, 2) Initiator, and 3) Responder
  - A node is granted to be an initiator only with one of the slots with the probability of p



## How to Synchronize the Nodes and Choose Slots (1/2)

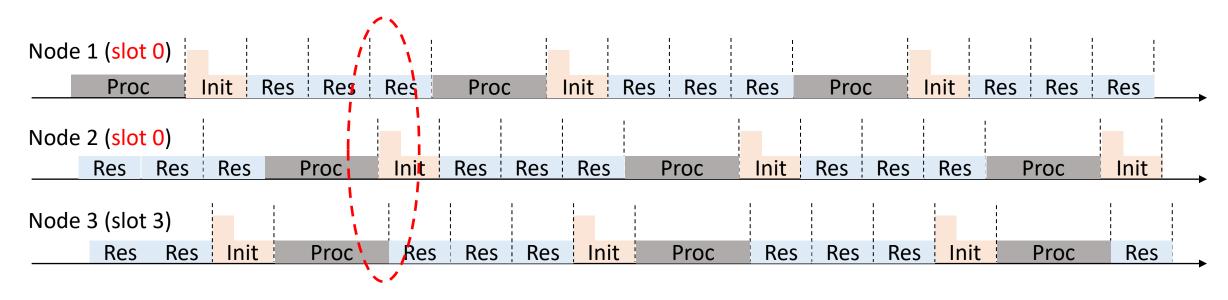
- The nodes wake up at different times
  - There is no global clock that they can use to align their slots
- There is no coordinator to assign slots to the nodes



An ideal slot boundary alignment and slot assignment

## How to Synchronize the Nodes and Choose Slots (1/2)

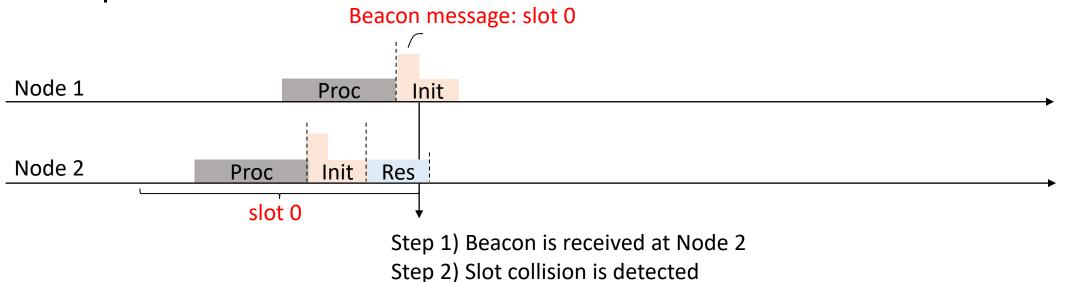
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Time synchronization and slot assignment problems

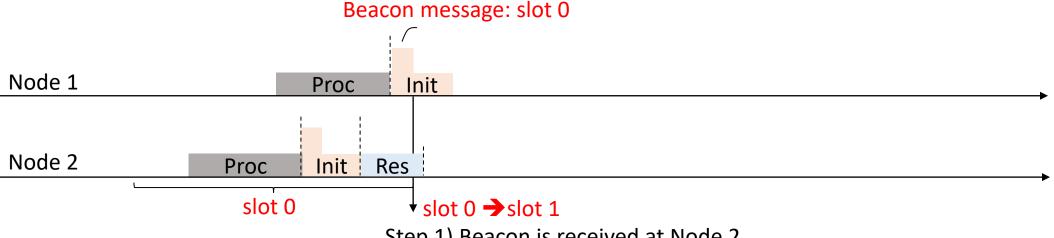
- We propose a joint synchronization & slot adjustment method:
  - 1) All the nodes broadcast a beacon signal which contains its own slot and its neighbors' slots (to address hidden node problems)
  - 2) If a node receives a beacon signal, the node selects a slot not listed in the beacon signal and adjusts its timing

#### Example:



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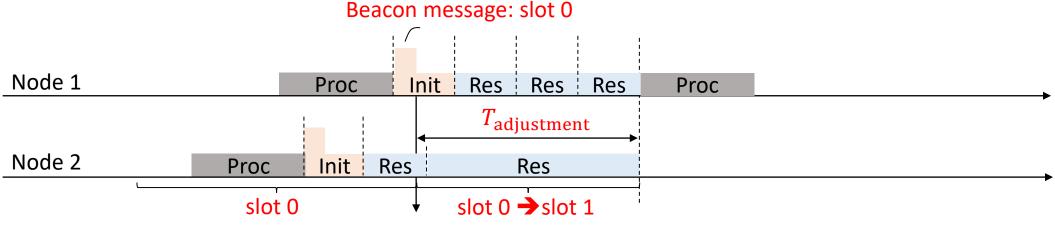


Step 1) Beacon is received at Node 2

Step 2) Slot collision is detected, hence choose a different slot, e.g., slot 1

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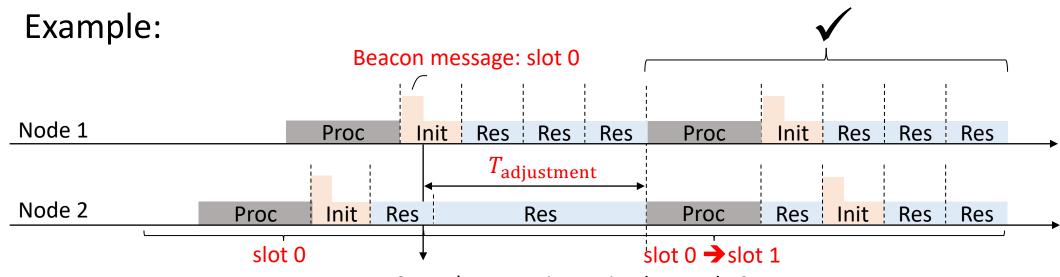


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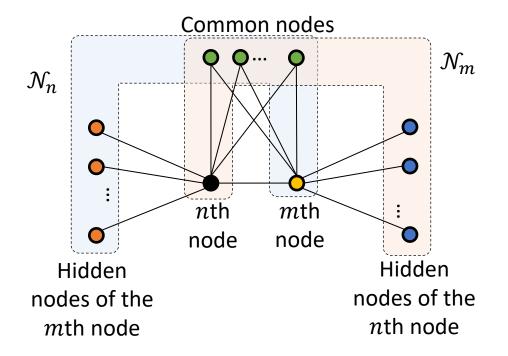
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### Addressing Hidden Node Problem (1/2)

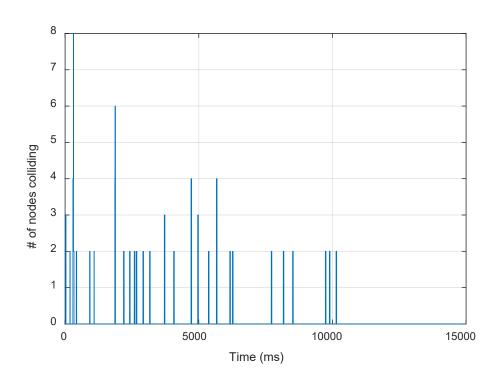
 The proposed approach reduces the collisions over time and addresses the hidden node problem

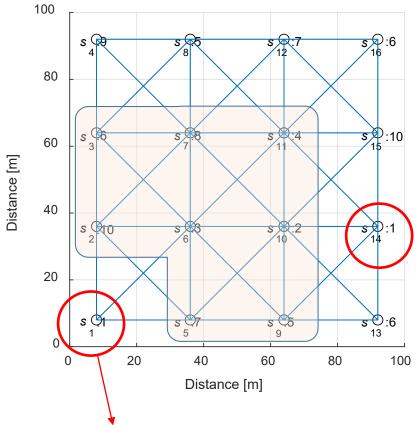


All the nodes in  $\mathcal{N}_n \cup \mathcal{N}_m$  should use a different slot to avoid interference between the mth and the nth node

## Addressing Hidden Node Problem (2/2)

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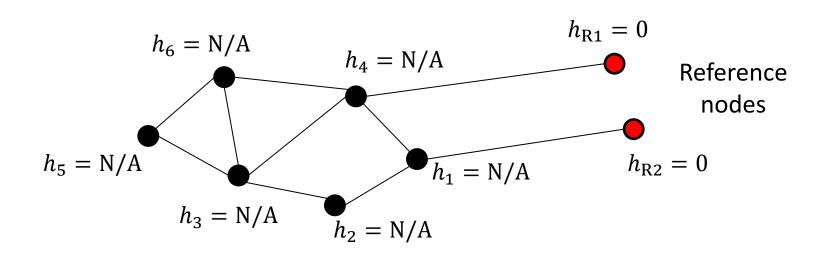
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An example of slots assignment after the proposed method runs for 15 seconds (10 slots, 16 nodes): Slot 1 is not used at the neighbors and neighbors' neighbors

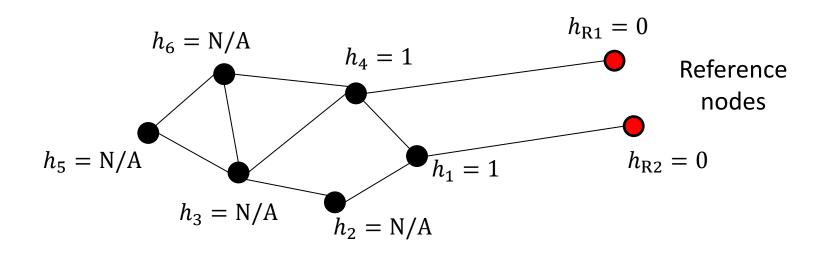
### **Solutions**

Goal 2: Transferring messages from nodes to the reference nodes

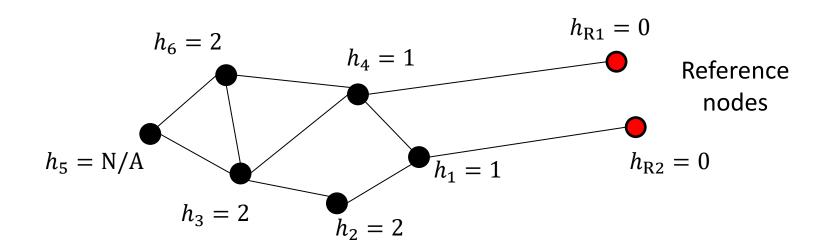
- Our approach is simple:
  - We define "hopping number" as # of hops to reach a reference node
    - The hoping number of a reference node is 0
  - The nth node determines its hopping number  $h_n$  based on its neighbors' hopping numbers as  $h_n \triangleq 1 + \min_{m \in \mathcal{N}_n} h_m$
  - Example:



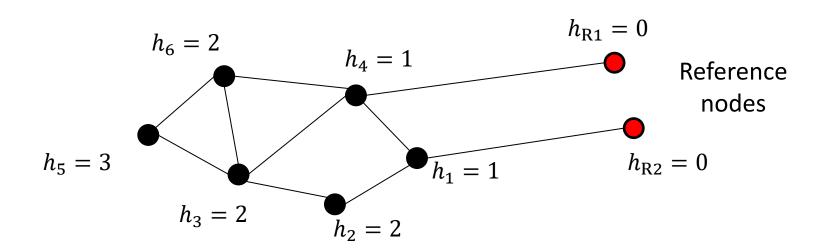
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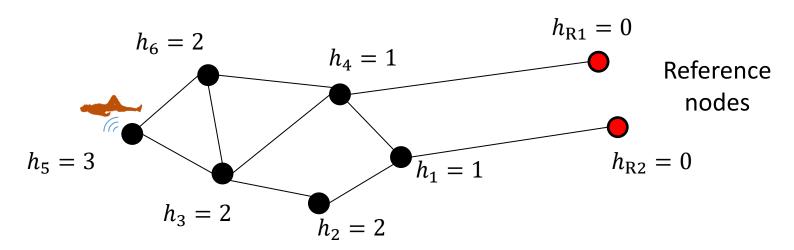
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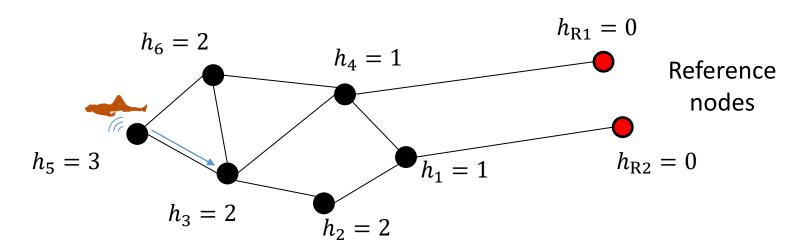
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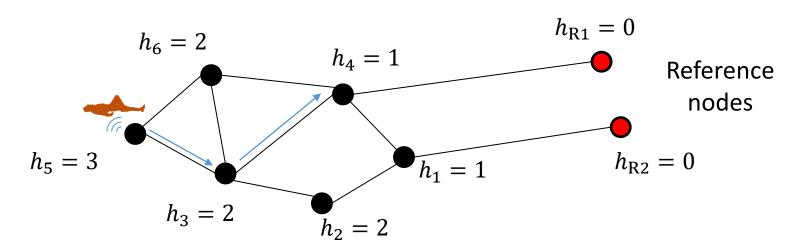
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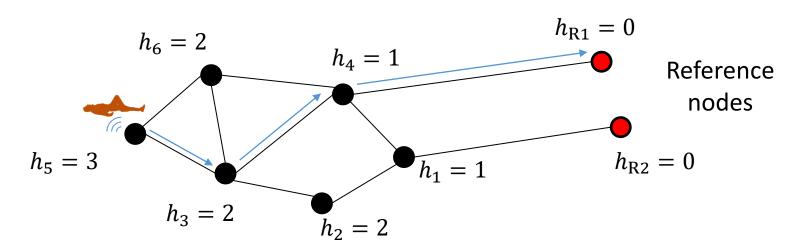
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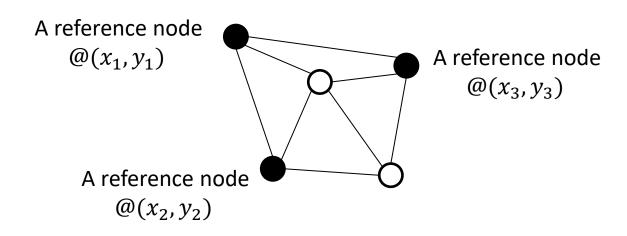
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### Solutions

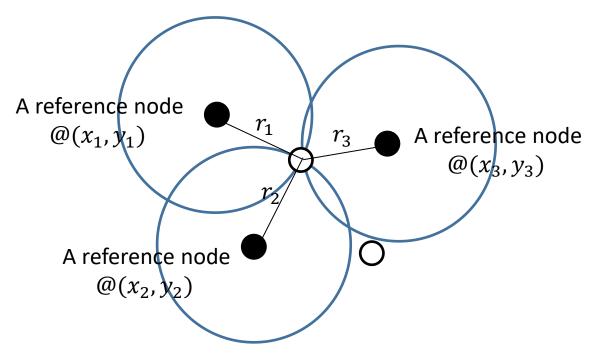
Goal 3: Finding the positions of the nodes to find the trapped people

- If a node has at least four nodes knowing their positions, it can estimate its position in 3D with the state-of-the-art methods (e.g., two-way ranging)
- Our approach relies on discovering the position of the nodes iteratively
  - We assume that all the reference node knows their position information
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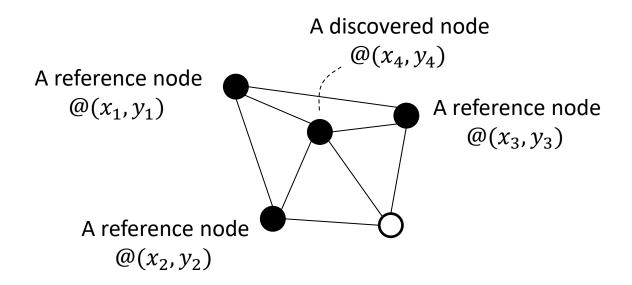
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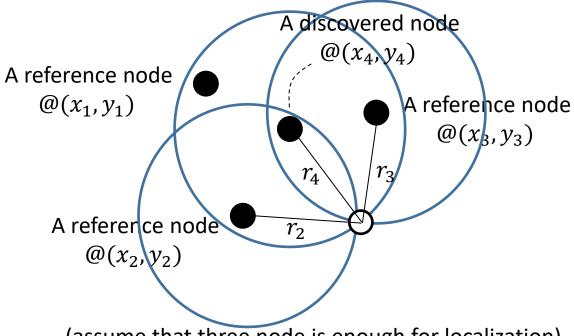
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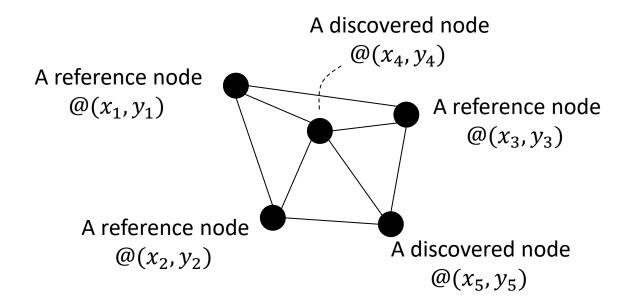


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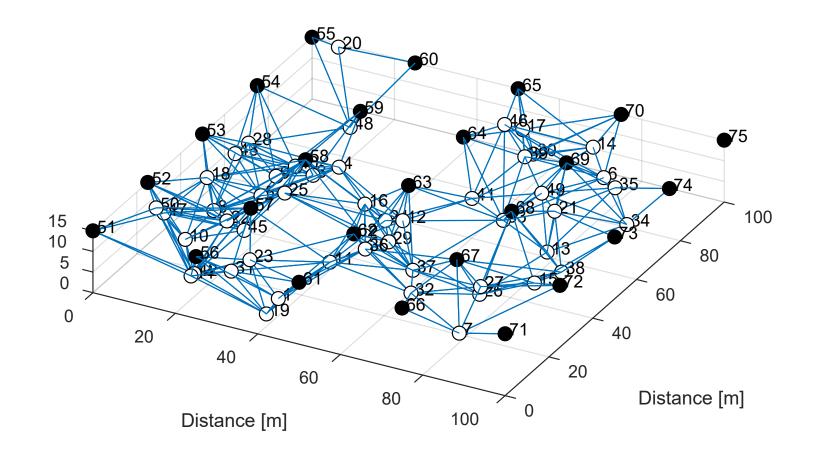


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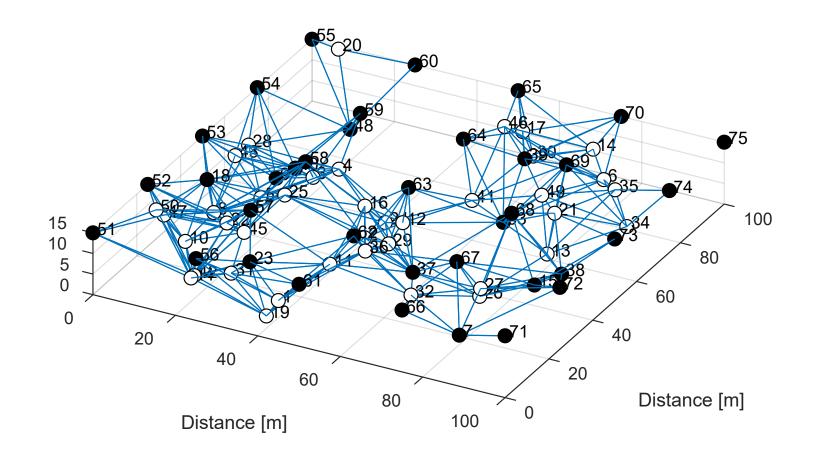


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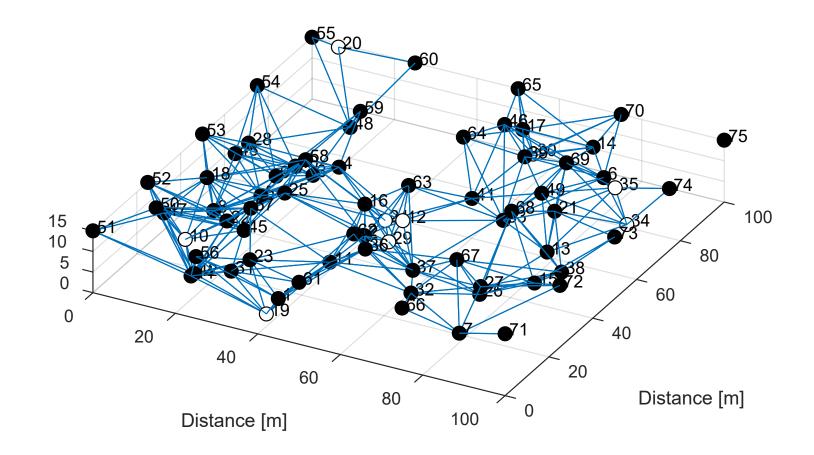
- 25 reference nodes (black ones), 50 nodes under the rubble (white ones)
  - At t = 0, none of the nodes know its position



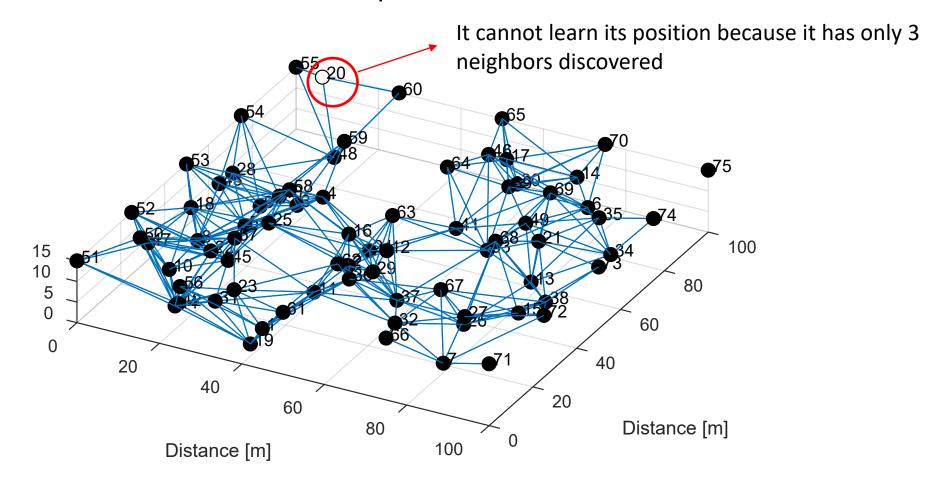
- 25 reference nodes (black ones), 50 nodes under the rubble (white ones)
  - At t = 1, 22% of the nodes learn their positions



- 25 reference nodes (black ones), 50 nodes under the rubble (white ones)
  - At t = 2, 82% of the nodes learn their positions



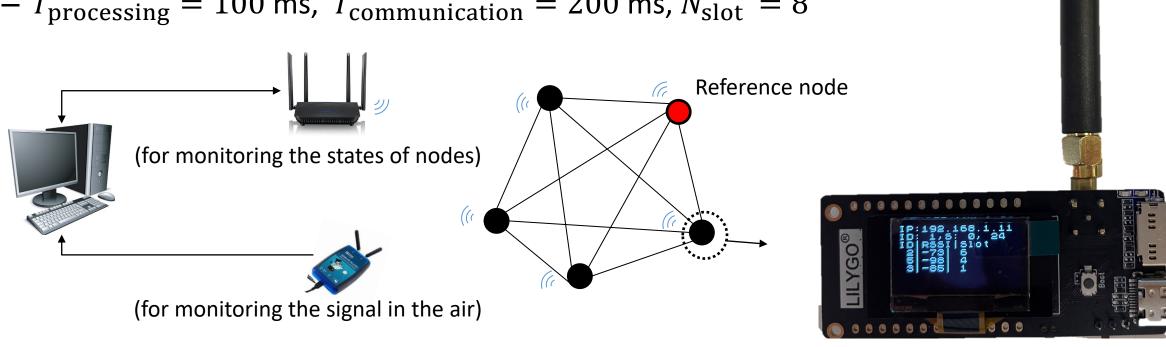
- 25 reference nodes (black ones), 50 nodes under the rubble (white ones)
  - At t = 3, 98% of the nodes learn their positions



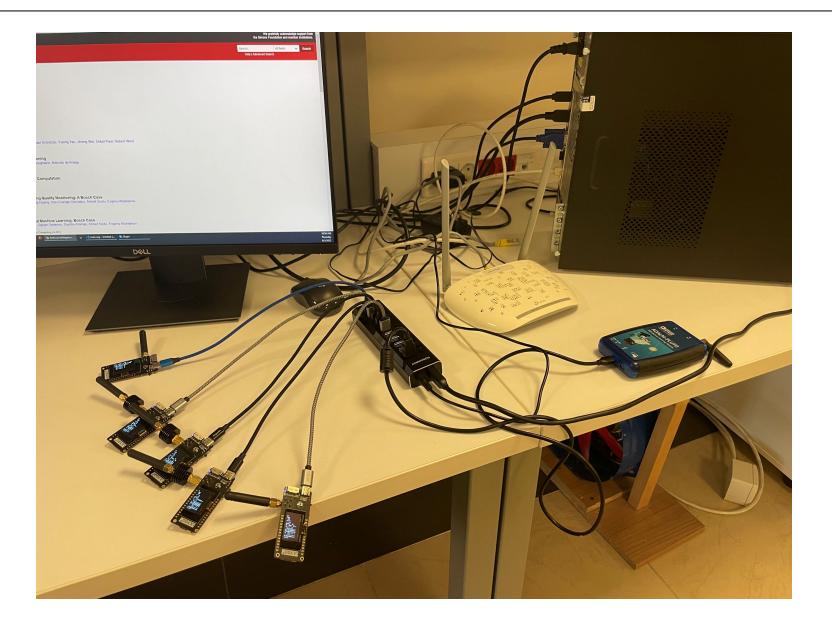
#### Demonstration

#### Setup

- A software defined radio is used for monitoring the signal in the air
- All nodes are connected to a router over Wi-Fi to monitor their states
- 5 nodes (SX1280 LoRa @ 2.5 GHz)
  - One of them is set as a reference node
  - $-T_{\text{processing}} = 100 \text{ ms}, T_{\text{communication}} = 200 \text{ ms}, N_{\text{slot}} = 8$

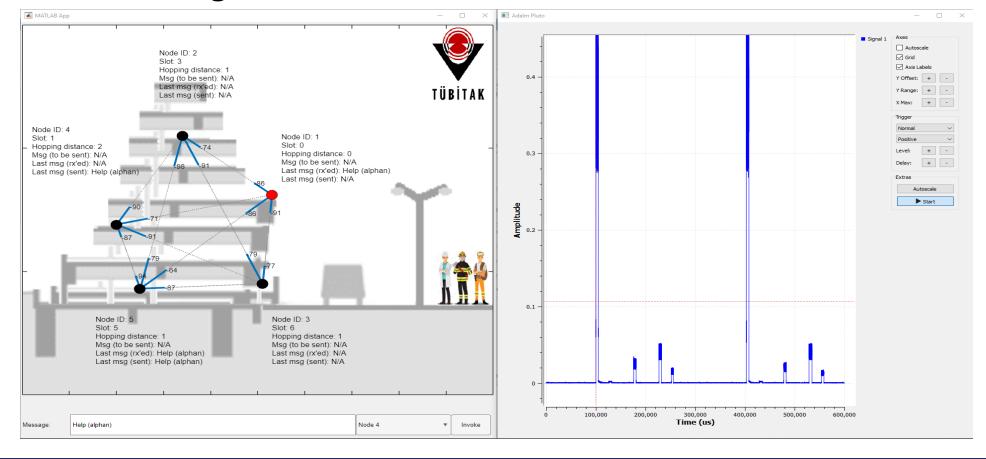


# Setup



### **GUI and Signals Over the Air**

- Result 1: All the nodes find their slots autonomously and adjust their timings
- Result 2: All nodes learn their hopping distance autonomously
- Result 3: The message from Node 4 reaches to the Node 1 over Node 5



# Questions?