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# **Controllability and Reachability**

Master degree in Automation Engineering

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# Controllability/Reachability

$$\begin{aligned} \dot{x}(t) \\ x(t+1) \end{aligned} \left\{ = Ax(t) + Bu(t) \quad x(0) = x_0 \right.$$

*free evolution*      *convolution sum*

$$x(t) = \phi(t) x(t_0) + \Psi(t) u([0,t))$$

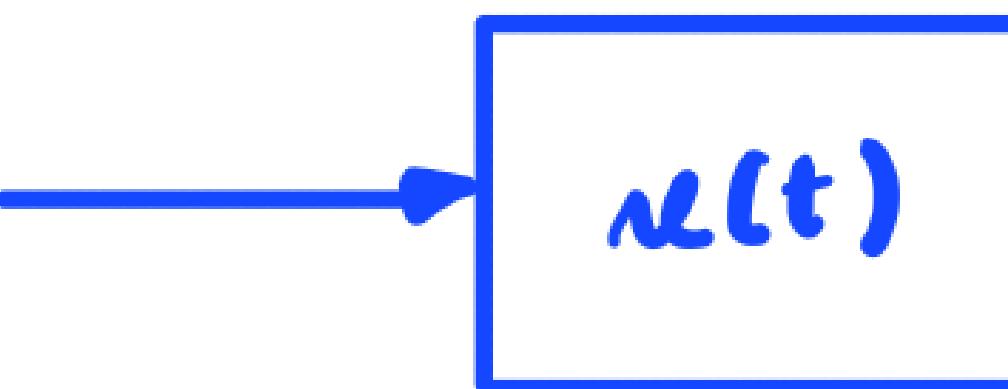
*exp. with power of A*

No outputs involved in the following analysis

**Main question to be answered:** given the system with a given initial condition (maybe the origin) can we reach in an appropriate time interval any “target state” in  $\mathbb{R}^n$ ? Are there some “regions” of  $\mathbb{R}^n$  that cannot be reached or regions where we remain “trapped”?

**Reachability:** Which state can be reached starting from the origin by playing with  $u(\cdot)$ ?  $\rightarrow \mathcal{R}^+$

**Controllability:** Which initial states can be steered to the origin by playing with  $u(\cdot)$ ?  $\rightarrow \mathcal{R}^-$



# Reachability

$$\begin{aligned} \dot{x}(t) \\ x(t+1) \end{aligned} \left\} = Ax(t) + Bu(t) \quad x(0) = x_0 \right.$$

$$x(t) = \phi(t)x(t_0) + \Psi(t)u([0,t))$$

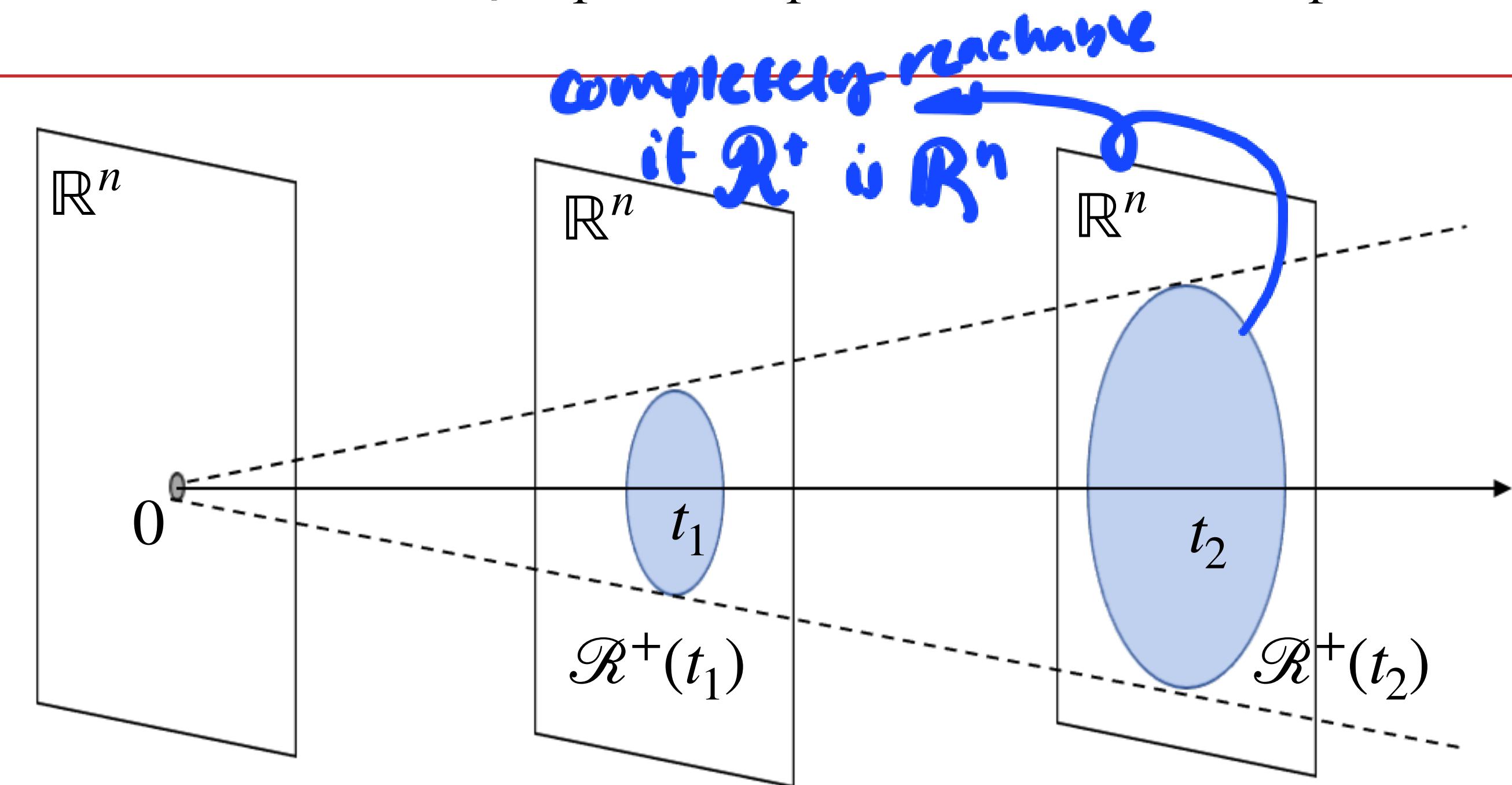
No outputs involved in the following analysis

**Set of reachable states at time  $t_1$ :**  $\mathcal{R}^+(t_1) = \{x \in \mathbb{R}^n : x = \psi(t_1)u([0,t_1]) \text{ for some } u([0,t_1]) \in \mathbb{R}^m\}$

If  $t_2 > t_1$  then  $\mathcal{R}^+(t_1) \subseteq \mathcal{R}^+(t_2)$

**Reachable set:**  $\mathcal{R}^+ := \mathcal{R}^+(\infty)$

**Definition:** The system is said to be completely reachable if  $\mathcal{R}^+ := \mathbb{R}^n$



# Reachability

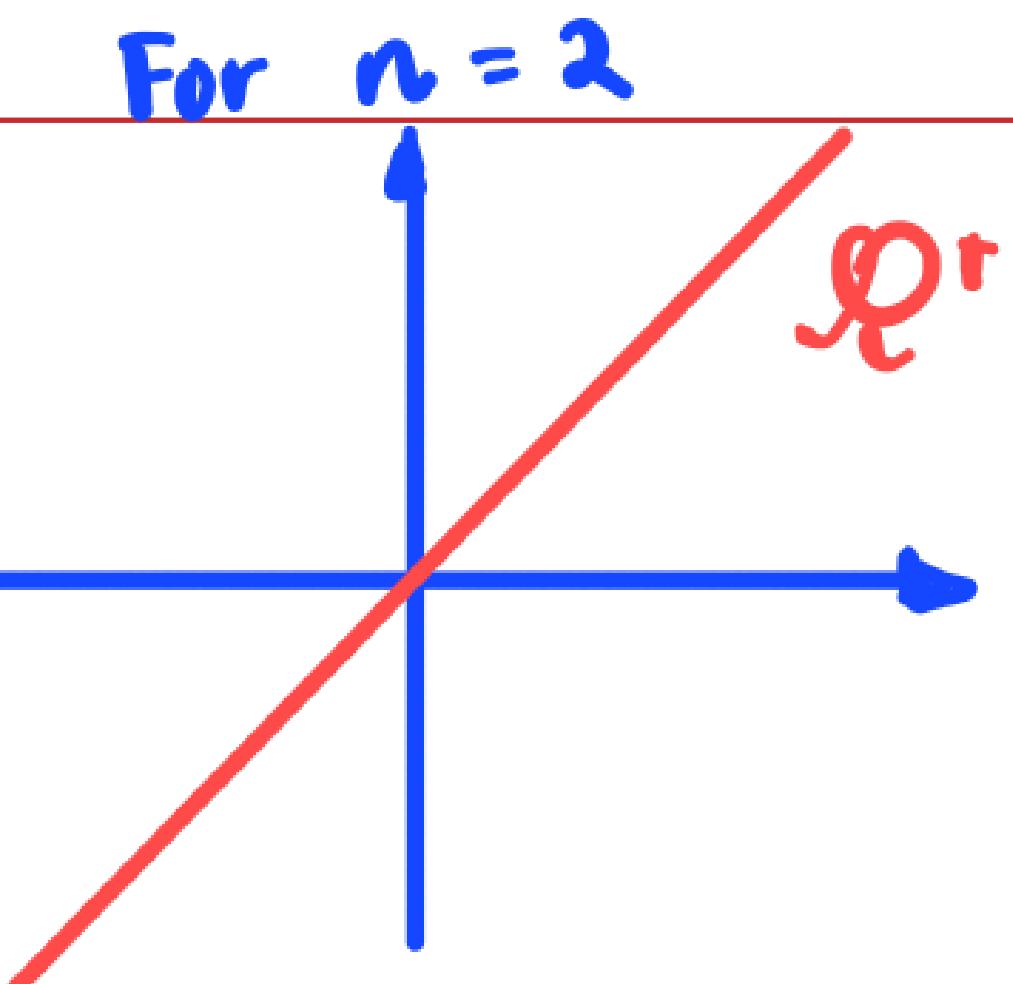
**Result:** the reachable set  $\mathcal{R}^+$  is a subspace of  $\mathbb{R}^n$  (much more than a set!)

**Theorem:**  $\mathcal{R}^+ = \text{Im } R$

$$R = [B \ AB \ A^2B \ \dots \ A^{n-1}B] \underbrace{\}_{mn}}$$

Reachability matrix

Dimension =  $n \times mn$ :  $\begin{cases} \text{"Fat" matrix if } m > 1, \\ \text{"Square" matrix if } m = 1 \end{cases}$



Proof ... (for D-T systems and intuition for C-T systems)

Image  $M = \{ v \in \mathbb{R}^{n \times m}; v = M w, w \in \mathbb{R}^m \}$

Remarks:

- If rank( $R$ ) =  $n$  (full row rank) then  $\mathcal{R}^+ = \mathbb{R}^n$  (all the states in  $\mathbb{R}^n$  can be reached from the origin by applying a certain control input)
- For C-T (D-T) systems if a state can be reached from the origin (namely if it is in the span of  $R$ ), then it can be reached in an arbitrarily small amount of time (in at most  $n$  steps)

$$x(t+1) = Ax(t) + Bu(t), \quad x(0) = 0$$

$$x(t) = \sum_{i=0}^{t-1} A^{t-i-1} Bu(i) \quad t \geq 1$$

$$x(1) = Bu(0) \quad \mathcal{R}^+(1) = \text{Image } B$$

$$x(2) = ABu(0) + Bu(1) \quad \mathcal{R}^+(2) = \text{Image } [B \ AB]$$

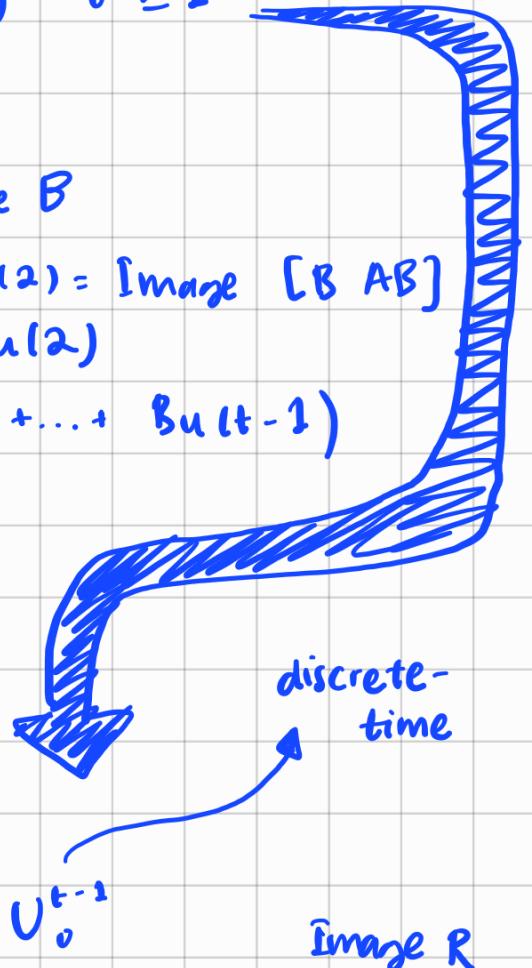
$$x(3) = A^2Bu(0) + ABu(1) + Bu(2)$$

$$x(t) = A^{t-1}Bu(0) + A^{t-2}Bu(1) + \dots + Bu(t-1)$$

$$R_t \triangleq [B : AB : A^{t-1}B]$$

$$U_0^{t-1} \triangleq \begin{bmatrix} u(t-1) \\ \vdots \\ u(0) \end{bmatrix}$$

collecting  
the input  
from  $t-1$   
up to 0



$$\text{Image } R_1 \subseteq \text{Image } R_2 \subseteq \dots \subseteq \text{Image } R_t \subseteq \text{Image } R_n$$

$$\text{Image } B \quad \text{Image } [B \ AB] \quad \text{Image } [B \ A^{t-1}B]$$

$$\mathcal{R}^+ = \text{Image } R$$



Goal to be  
proved set of  
states reachable  
in  $n$  steps.

$A^n B$  ← linearly  
dependent  
on  $[B \ AB \ A^{n-1}B]$

$$\text{Image } R_n \subseteq \text{Image } R_{n+1} = \dots = \text{Im } R_{n+2}$$

Cayley - Hamilton:

$$\varphi_A(\lambda) = \det(\lambda I - A) = A^{n-1} + \dots + \alpha_n I$$

$$= \lambda^n + \alpha_1 \lambda^{n-1} + \dots + \alpha_n$$

$$A^n B = -\alpha_1 A^{n-1} B - \alpha_n B$$

$$A^n B = -\alpha_1 A^{n-1} B - \dots - \alpha_n AB$$

Image  $R_n$  & Image  $R_{n+1}$

$$= \text{Image } R_{n+2} = [R_{n+1} : \boxed{A^{n+2} B}]$$

Homework:

For cont. time system,

$$x^{[k]}(0) = R_k U_0^{k-1} = \begin{bmatrix} U^{(k-1)}(0) \\ \dots \\ U^{(2)(0)} \\ U^{(1)}(0) \end{bmatrix}$$

For discrete time system:

$$x(t) = R_t U_0^{t-1}$$

For continuous-time

$$x^{[k]}(0) = R_k U_0^{k-1} \quad \text{where } U_0^{k-1} = \begin{bmatrix} U^{(k-1)}(0) \\ \vdots \\ U^{(2)(0)} \\ U^{(1)}(0) \end{bmatrix}$$

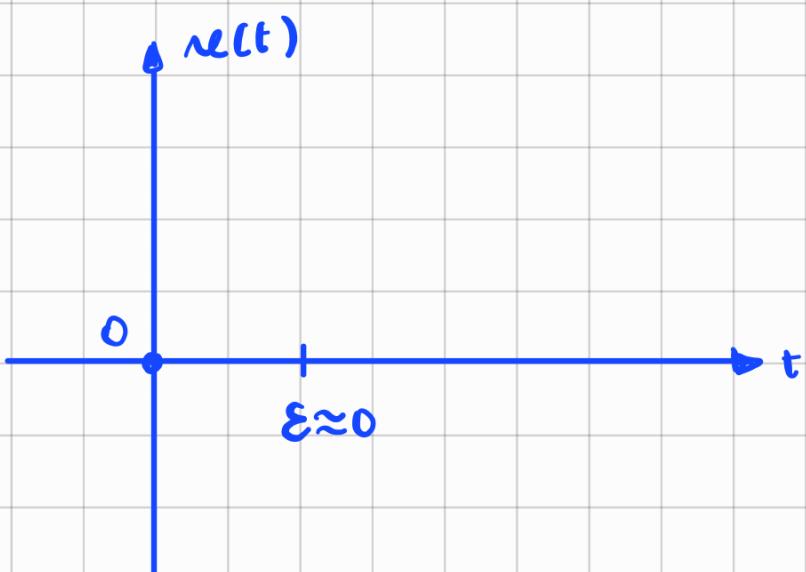
→ reachability matrix

Suppose that  $\mathbb{R}^+ = \mathbb{R}^n$ , the image of  $R = \mathbb{R}^n$ , by Taylor,

$$x(t) = x(t_0) + \frac{1}{1!} x^{(1)}(t_0)(t-t_0) +$$

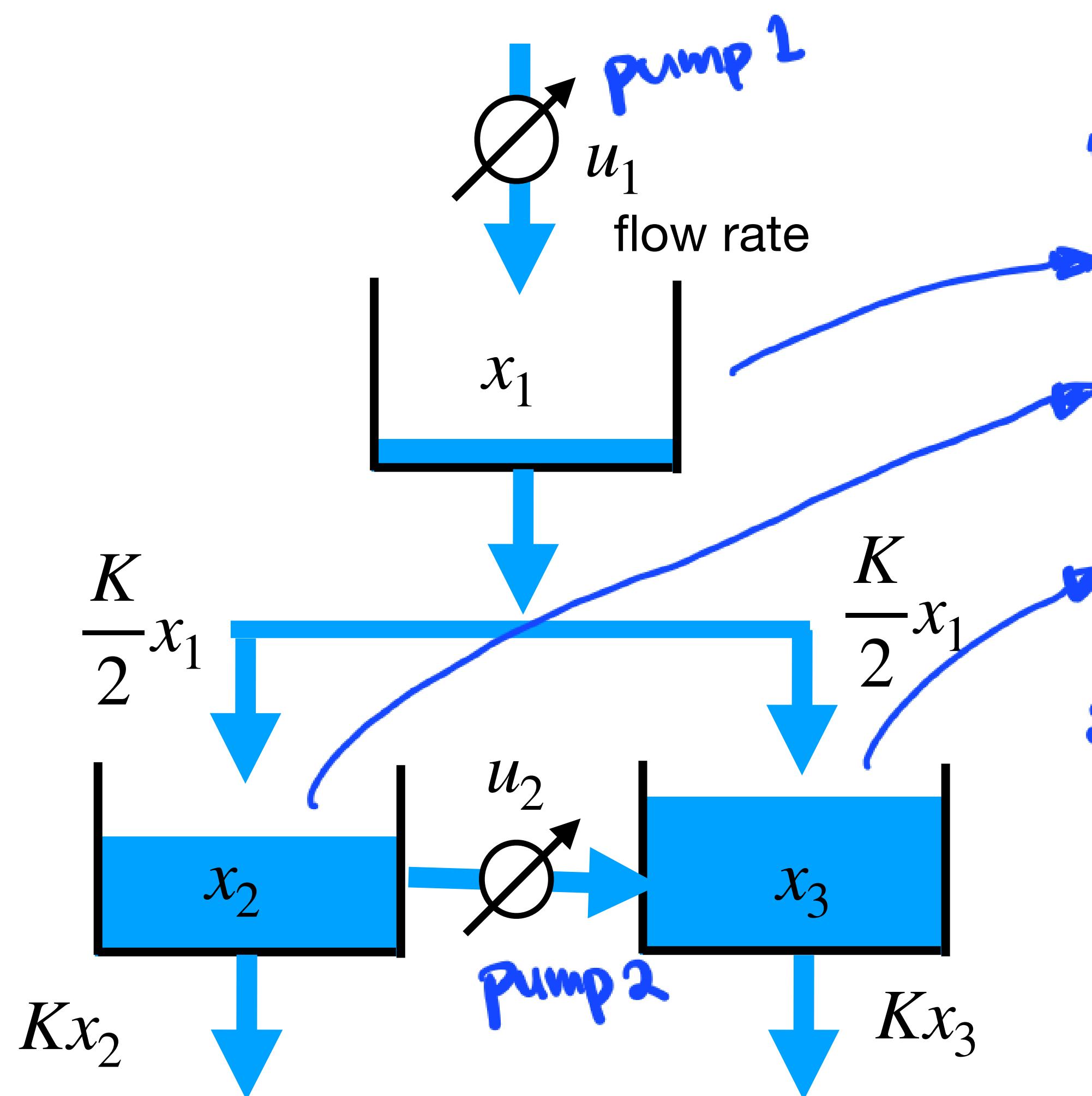
$$\frac{1}{2!} x^{(2)}(t_0)(t-t_0)^2 + \dots + \frac{1}{n!} x^{(n)}(t_0)(t-t_0)^n$$

Indeed,  $\nu(t) \equiv$  linear combination of  $\nu^{[k]}(0)$ ,  
where  $k = 1, 2, \dots, n, n+1, \dots, \infty$



## Example

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Is this system completely reachable?

1. Modeling the system

$$\dot{x}_1 = u_1 - \frac{K}{2}x_1 - \frac{K}{2}x_1 = u_1 - x_1$$

$$\dot{x}_2 = \frac{K}{2}x_1 - Kx_2 - u_2$$

$$\dot{x}_3 = \frac{K}{2}x_2 - Kx_3 + u_2$$

2. A, B matrix

$$A = \begin{bmatrix} -K & 0 & 0 \\ K/2 & -K & 0 \\ K/2 & 0 & -K \end{bmatrix}$$

$$B = \begin{bmatrix} 1 & 0 \\ 0 & -1 \\ 0 & 1 \end{bmatrix}$$

### 3. Reachability matrix

$$R = \begin{bmatrix} 1 & 0 & -K & * \\ 0 & -1 & K/2 & * \\ 0 & 1 & K/2 & * \end{bmatrix}$$

$\brace{B}$ 
 $\brace{AB}$ 
 $\brace{A^2B}$

→ this col. is not needed  
because we already  
know the rank of  $R = 3$

According to the theory, the system is reachable.

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Now, how if pump 2 is faulty?

1. New B matrix:

$$B = \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix}$$

NOTE :  
this analysis  
only applies  
on linear Syst.

2. Reachability matrix:

$$R = \begin{bmatrix} 1 & -K & K^2 \\ 0 & K/2 & -K^2 \\ 0 & K/2 & -K^2 \end{bmatrix}$$

$\brace{B}$ 
 $\brace{AB}$ 
 $\brace{A^2B}$

3rd col
-2K
 $\begin{bmatrix} -K \\ K/2 \\ K/2 \end{bmatrix}$

$\begin{bmatrix} K^2 \\ -K^2 \\ -K^2 \end{bmatrix}$ 
 $= -2K^2$ 
 $\begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix}$

rank  $R = 2$ 
 $\mathcal{R}^+ = \text{image } R$

**Popov - Belevitch- Hautus** test: it presents a sufficient and necessary condition for complete ~~observability~~ **reachability**, not requiring the computation of the reachability matrix

The pair  $(A, B)$  is completely reachable ( $\text{rank}R = n$ ;  $\mathcal{R}^+ = \mathbb{R}^n$ ) iff

$$\text{rank } [\lambda I - A \ B] = n \quad \forall \lambda (\in \sigma(A))$$

*pick a generic  $A$ , eigenvalue  $\lambda$ , and multiply with  $I$  matrix.*

full row-rank (if  $\lambda$  is not in  $\sigma(A)$  the condition is always fulfilled)

**Proof** ( $\text{rank}R = n \Rightarrow \text{rank}[\lambda I - A \ B] = n$ )

The proof:

$$\text{rank} \left( \begin{bmatrix} (\lambda I - A) & | & B \\ \downarrow & | & \swarrow \\ \text{the determinant} & | & m \\ \text{should be } \neq 0 & & \end{bmatrix}_{n \times n} \right) = n, \forall \lambda$$

If  $\lambda \notin \sigma(A)$ ,  $\det(\lambda I - A) \neq 0 \Rightarrow (\lambda I - A)$  is not singular. As long as the previous condition holds, the test is always fulfilled.

Suppose that the system is reachable,  $R^+ = R^n$ ,  $\text{rank } R = n$ , and there exists  $\exists \lambda^* : \text{rank } (\lambda^* I - A : B) < n$ , by contradiction.

$$\Rightarrow \exists W^* \in \mathbb{R}^n, W^{*T}(\lambda^* I - A : B) = 0$$

$$\begin{aligned} &\text{non-zero} \\ &W^{*T}(\lambda^* I - A)^B = 0^B \\ &W^{*T} \underbrace{\lambda^* B}_{=0} = \boxed{W^{*T} AB = 0} \\ &\boxed{W^{*T} A^2 B = 0} \end{aligned}$$

reachability matrix  $R$

By these relations,

$$\text{we can conclude } W^{*T} [B : AB : \dots : A^{n-1} B] = 0$$

## Computing the input - The D-T case

Suppose  $AB$  is reachable

How to practically compute the input steering the state of a system from the origin to a target state in  $\mathcal{R}^+$ ?  
 $= \mathbb{R}^n$

We know that  $x(t) = R_t u_0^{t-1}$  with  $R_t = [B \ AB \ \dots \ A^{t-1}B]$  and  $u_0^{t-1} = [u(t-1) \ u(t-2) \ \dots \ u(0)]^T$

$n \times t \cdot m$

- Case  $\mathcal{R}^+ = \mathbb{R}^n$  (all states of  $\mathbb{R}^n$  are reachable in at most  $n$  steps). Let  $\bar{x}$  a target state.

$$t \geq n, m \geq 1$$

generic

$$\bar{x} = R_t u_0^{t-1}$$

**“Fat” matrix** (full row rank)

$$u_0^{t-1} = R_t^T (R_t R_t^T)^{-1} \bar{x}$$

Right inverse of  $R_t$

**Result:** Given any full row rank matrix  $M$ , the square matrix  $MM^T$  is nonsingular

Special case  $t = n, m = 1$

generic

$$\bar{x} = R_n u_0^{n-1} = R u_0^{n-1}$$

**square** (not singular)

$$u_0^{n-1} = R^{-1} \bar{x}$$

$$\bar{R} = R(t) = R_t \quad u^{t-1}$$

fixed      fixed      fixed

*to be computed*

Discrete-time :

$$t \geq n$$



$$\bar{R} = \left( \underbrace{\begin{matrix} R_t \\ \vdots \\ R_t \end{matrix}}_{n-mt} \right)_n \quad \left( \underbrace{\begin{matrix} R_t^T \\ \vdots \\ R_t^T \end{matrix}}_{n-mt} \right)^n \quad \left| \begin{matrix} u^{t-1} \end{matrix} \right|$$

how to  
invert  $R_t$ ?

NB :  $[R_t R_t^T] [R_t R_t^T]^{-1} = I$

$$[R_t R_t^T]^{-1}$$

In case of  $t \geq n$

$$R_t = \left[ \underbrace{\begin{matrix} B & AB & \dots & A^{n-1}B & A^nB & A^{t-1}B \end{matrix}}_{n \times nm} \right]$$

## Computing the input - The D-T case

**no assumption on t**

- In all the other cases (namely  $t < n$  or  $\mathcal{R}^+ \subset \mathbb{R}^n$ ) the target state  $\bar{x}$  is not generic but, to be reached, must fulfil  $\bar{x} \in \text{Im } R_t$ . In these cases the equation  $\bar{x} = R_t u_0^{t-1}$  can be solved for  $u_0^{t-1}$ :

→ must be  
the image  
of  $R_t$

$$u_0^{t-1} = R_t^\dagger \bar{x}$$

Generalized (Moore Penrose) inverse of  $R_t$

$\bar{u} \in \mathcal{R}_2$   
 $\parallel$   
 $[B]$  → single  
col.  
tall  
matrix

The Moore-Penrose pseudoinverse is defined for any matrix and is unique. If the matrix is (right) invertible it boils down to the canonical (right) inverse. Otherwise, it provides the unique solution to  $\bar{x} = R_t u_0^{t-1}$  if the latter has a solution (namely  $\bar{x} \in \text{Im } R_t$ ). If the equation in question does not admit a solution (namely  $\bar{x} \notin \text{Im } R_t$ ) the Moore-Penrose inverse provides the “closest” solution to  $\bar{x}$  in the Euclidean sense

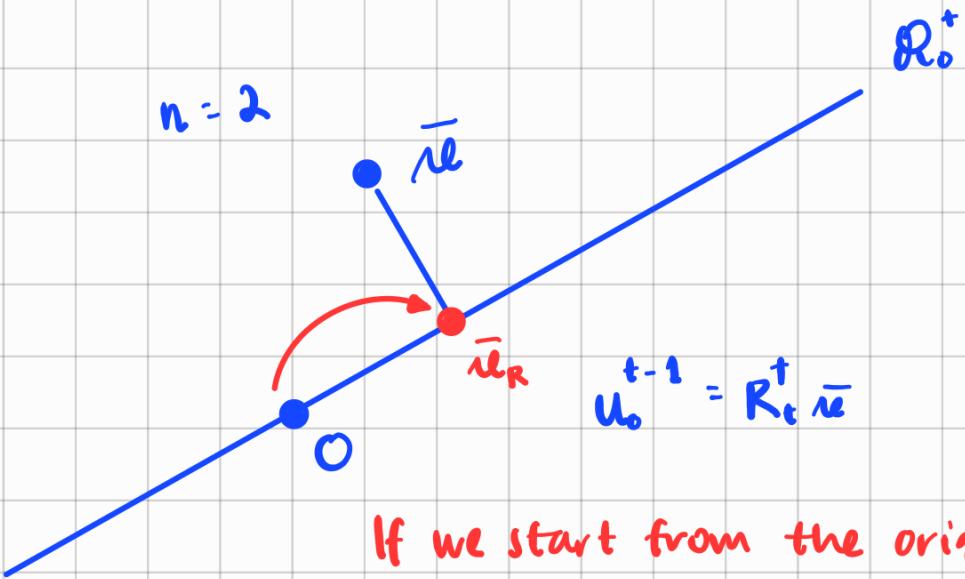
$$\|\bar{x} - R_t R_t^\dagger \bar{x}\| \leq \|\bar{x} - R_t u_0^t\|$$

generic input sequence

$$\bar{u}_t = R_t u_0^{t-1}$$

$$u_0^{t-1} = R_t^\dagger \bar{u}$$

Suppose that the system is not completely reachable



If we start from the origin,  
we will reach the state  
close to the  $\bar{x}$ , which is  
projected by Moore-Penrose  
matrix  $R_o^t$

## Example: Taxi company

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$u(t)$ : # new cars bought at year  $t$

$x_1(t)$ : # of 1-year old cars at year  $t$

$x_2(t)$ : # of cars older than 1 year at year  $t$

$P$ : Probability that in a generic year a car has not an irreparable accident



**Problem:** Compute the number of car to be bought so that after two years the park has  $x_1^*$  1-Y old cars, and  $x_2^*$  older cars (suppose the park is empty at  $t = 0$ ). What's the input profile if the park at  $t = 0$  is not empty?

$$x_1(t+1) = Pu(t)$$

$$x_2(t+1) = Px_1(t) + Px_2(t)$$

1. Modelling the eq.  $\rightarrow$  new cars bought

$$n_1(t+1) = P u(t)$$

$$n_2(t+1) = P n_1(t) + P n_2(t)$$

2. The matrix

$$A = \begin{bmatrix} 0 & 0 \\ P & P \end{bmatrix} \quad B = \begin{bmatrix} P \\ 0 \end{bmatrix}$$

3. Reachability matrix

$$R = \begin{bmatrix} P & 1 & 0 \\ 0 & 1 & P \\ 0 & 0 & P^2 \end{bmatrix} \quad \rightarrow \text{rank} = 2$$

$\underbrace{\phantom{0}}_{B} \quad \underbrace{\phantom{0}}_{AB}$

$$R^{-1} = \begin{bmatrix} 1/P & 0 \\ 0 & 1/P^2 \end{bmatrix}$$

The system is reachable,  
so it is possible to  
reach any number of  
cars.

4. To reach the  $n_1^*$  and  $n_2^*$ :

$$\begin{bmatrix} u(1) \\ u(0) \end{bmatrix} = \begin{bmatrix} 1/P & 0 \\ 0 & 1/P^2 \end{bmatrix} \begin{bmatrix} n_1^* \\ n_2^* \end{bmatrix}$$

$\underbrace{\phantom{0}}_{U_0}$

↑  
any

5. In case of three years

$$\begin{bmatrix} u(3) \\ u(2) \\ u(1) \\ u(0) \end{bmatrix} = R_3^T [R_3 R_3^T] \begin{bmatrix} u_1^* \\ u_2^* \end{bmatrix}$$

## Computing the input - The C-T case

 **initial state**

$$\dot{x} = Ax + Bu, \quad x(0) = 0 \quad x(t) = \int_0^t e^{A(t-s)} B u(s) ds$$

How to compute the input  $u(s)$ ,  $s \in [0,t]$  steering the state of a system from the origin to a target state in  $x(t) = \bar{x} \in \mathcal{R}^+$ ?

Reachability/Controllability Gramian at time  $t$

$$W(t) = \int_0^t e^{A(t-s)} B B^T e^{A^T(t-s)} ds$$

**Suppose the system is completely reachable**

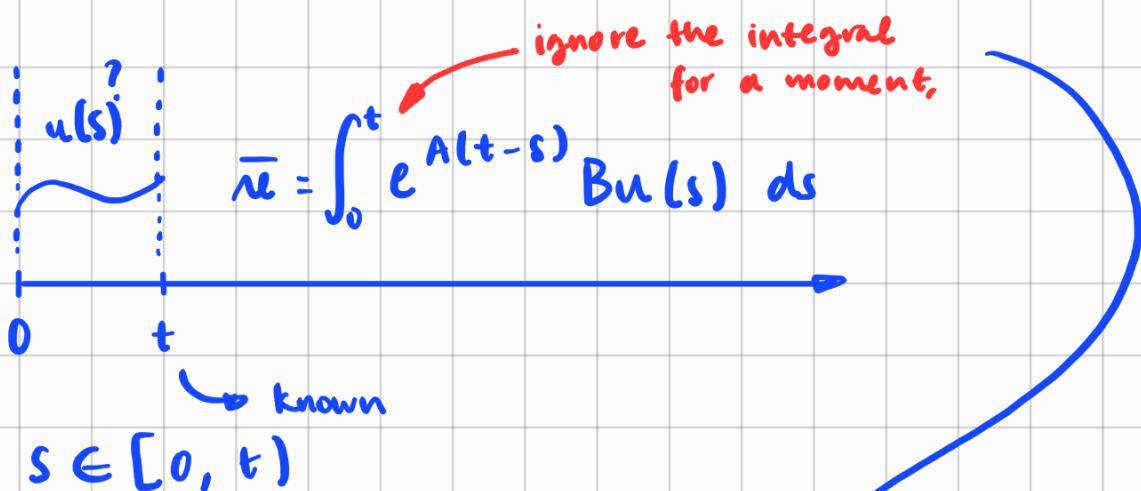
**Theorem:** The Gramian is not singular for each  $t > 0$  iff  $\mathcal{R}^+ = \mathbb{R}^n$  

Proof...

**Result:** Assume that the system is completely reachable. Let  $\bar{x} \in \mathbb{R}^n$  a generic target state. Let  $t > 0$  be an arbitrary time. then the control input able to steer the state of  $(A, B)$  from the origin to  $\bar{x}$  in the interval  $[0,t)$  is

 **target state**

$$u(s) = B^T e^{A^T(t-s)} W^{-1}(t) \bar{x} \quad s \in [0,t)$$



$$m = 1$$

$$\left[ e^{A(t-s)} B \right] \left[ B^T e^{A^T(t-s)} \right] \left[ e^{A(t-s)} B B^T e^{A^T(t-s)} \right] \bar{x}$$

$n \times n$

rank  $[v \ v^T] =$   
rank  $v$   
is not  
an invertible  
matrix

$$W(t) = \int_0^t e^{A(t-s)} B B^T e^{A^T(t-s)} ds$$

$n \times n$  singular

Controllability/  
reachability Gramian

$$\bar{x} = \int_0^t e^{A(t-s)} B \left[ B^T e^{A^T(t-s)} \right] ds W^{-1}(t) \bar{x}$$

$w(t)$

$$u(s) = B^T e^{A^T(t-s)} W^{-1}(t) \bar{w}$$

for  $s \in [0, t]$

$$w(t) = \int_0^t e^{A(t-s)} B u(s) ds$$

A	B

$$\xrightarrow{?} w(t)$$

## Computing the input - The C-T case

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**General case:** How to compute the control input steering the state of a linear system from a generic initial condition (not necessarily the origin) to a generic final state in a predetermined time interval?



$$x(t) - e^{At} x_0 = \int_0^t e^{A(t-s)} B u(s) ds$$

Because of linearity the problem can be cast as the problem of starting the state of the system from the origin to the final target  $\bar{x}' = \bar{x} - e^{At} x_0$  in a predetermined interval.

$$u(s) = B^T e^{A^T(t-s)} W^{-1}(t) (\bar{x} - e^{At} x_0) \quad s \in [0, t]$$

For the Discrete-Time case:

Lagrange:  $x(t) = A^t x_0 + \sum_{i=0}^{t-1} A^{t-i} B u(i)$  ✓  
formula

$\underbrace{R_t u_0^{t-1}}$

$x_0$

$x(t) - A^t x_0 = R_t u_0^{t-1}$   $R_t u_0^{t-1}$

$u_0^{t-1} = R_t^+ (\bar{x} - A^t x_0)$

$t \geq n$   $\mathcal{R}^+ = \mathbb{R}^n$

For the Continuous-Time case:

$$x(t) = e^{At} x_0 + \int_0^t e^{A(t-s)} B u(s) ds$$
 ✓

$$x(t) - e^{At} x_0 = \int_0^t e^{A(t-s)} B u(s) ds$$

# Controllability

$$\begin{aligned} \dot{x}(t) \\ x(t+1) \end{aligned} \Big\} = Ax(t) + Bu(t) \quad x(0) = x_0$$

$$x(t) = \phi(t)x(t_0) + \Psi(t)u([0,t))$$

No outputs involved in the following analysis

**Set of controllable states at time  $t_1$ :**

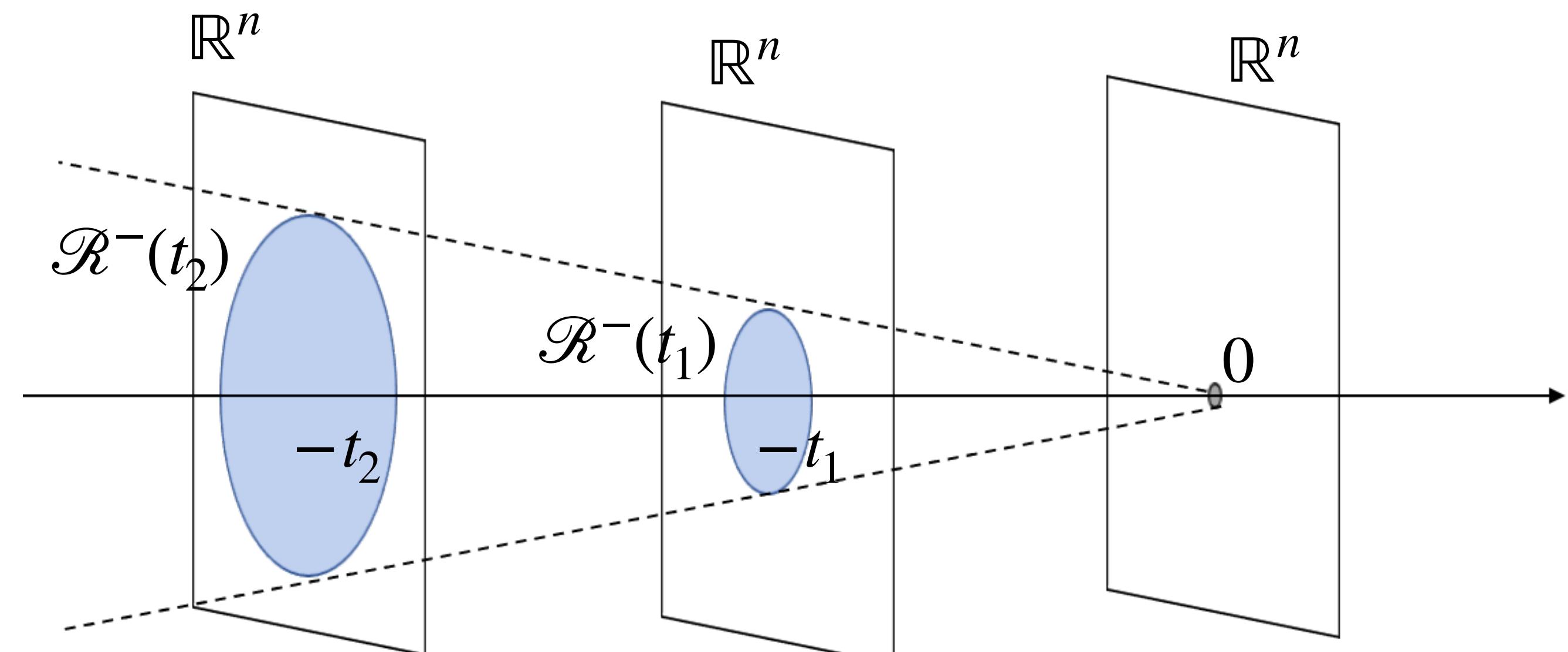
$$\mathcal{R}^-(t_1) = \{x \in \mathbb{R}^n : 0 = \phi(t_1)x + \psi(t_1)u([0,t_1]) \text{ for some } u([0,t_1]) \in \mathbb{R}^m\}$$

If  $t_2 > t_1$  then  $\mathcal{R}^-(t_1) \subseteq \mathcal{R}^-(t_2)$

**Controllable set:**

$$\mathcal{R}^- := \mathcal{R}^-(\infty) \quad (\text{reverse from } \mathcal{R}^+)$$

**Definition:** The system is said to be completely controllable if  $\mathcal{R}^- := \mathbb{R}^n$



## Controllability

not only a set,  
but a subspace

**Result:** the reachable set  $\mathcal{R}^-$  is a subspace of  $\mathbb{R}^n$  (much more than a set!)

**Theorem :**  $\mathcal{R}^+ \subseteq \mathcal{R}^-$  in general.  $\mathcal{R}^+ = \mathcal{R}^-$  if the system is reversible

**Remark:** There could be systems with state that are controllable to the origin but that cannot be reached from the origin. **Complete reachability  $\Rightarrow$  complete controllability**

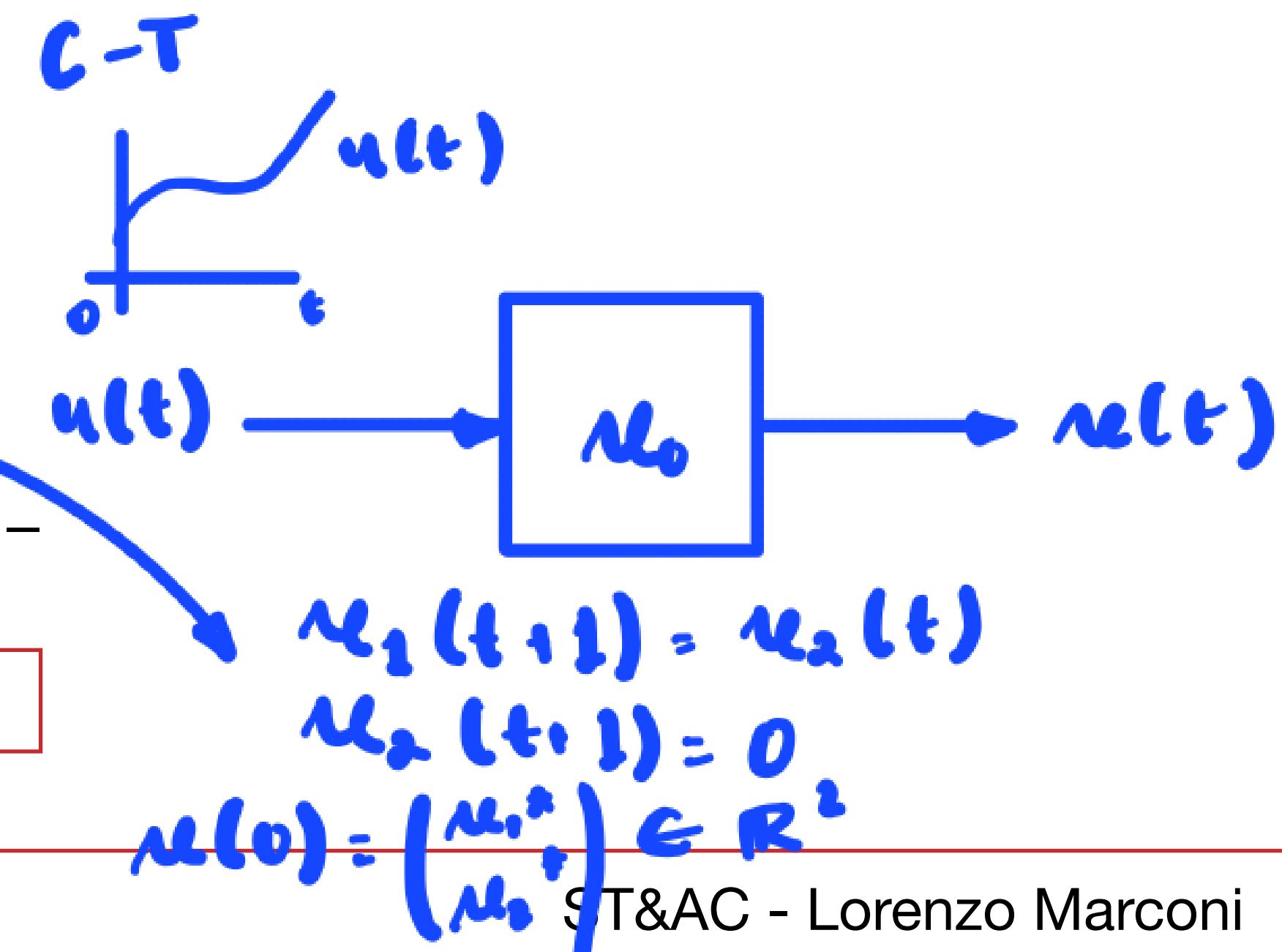
**Remark:** All continuous-time systems are reversible, and thus  $\mathcal{R}^+ = \mathcal{R}^-$ . However there could be discrete time systems for which the inclusion holds true

**Example:** trivially  $\mathbb{R}^1$  is the origin for this case

$$x(t+1) = Ax(t) + Bu(t) \quad A = \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix} \quad B = \begin{pmatrix} 0 \\ 0 \end{pmatrix}$$

**Remark:** D-T sampled-data systems are reversible and thus  $\mathcal{R}^+ = \mathcal{R}^-$

**Remark:** From now on controllability and reachability will be confused



For the Discrete-Time case:

$$n_1(t+1) = n_2(t)$$

$$n_2(t+1) = 0 \quad \longrightarrow \quad n_2(1) = 0$$

$$n_2(0) = \begin{bmatrix} n_1^* \\ n_2^* \end{bmatrix} \in \mathbb{R}^2$$

$$n_1(1) = n_2^*$$

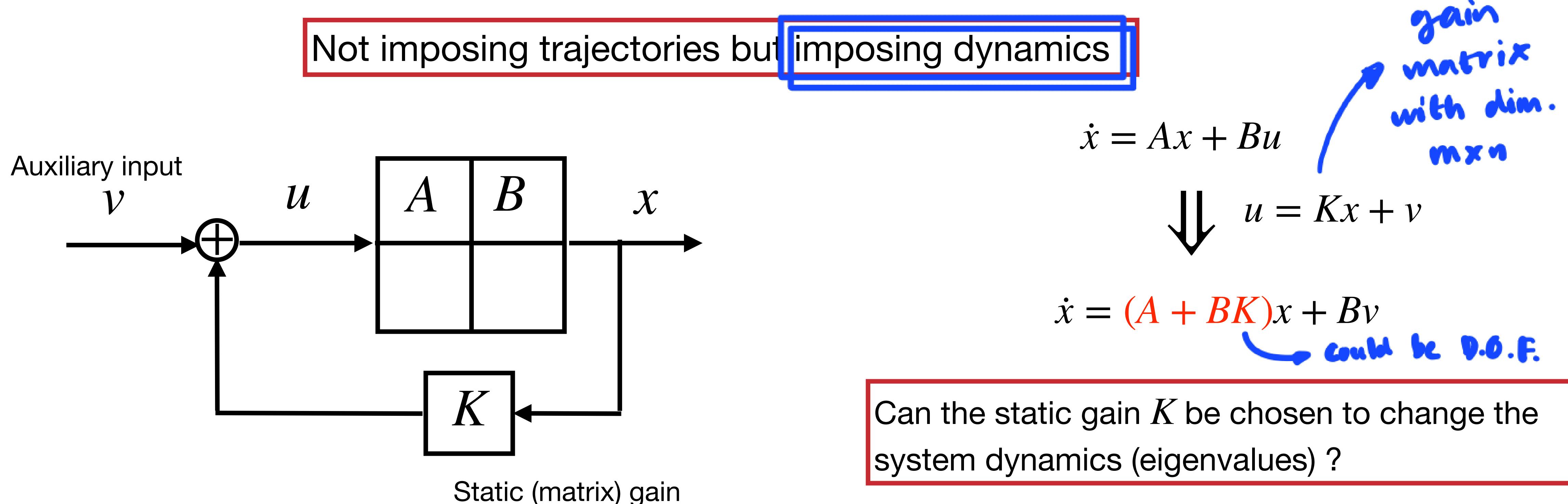
$$n_1(2) = 0$$

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## Controllability and state feedback

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Up to now we focused on computing a control input steering the state between two desired states in an “open loop” way. If the system is completely controllable/reachable we have full authority to steer the state of the systems between two arbitrary states (also in an arbitrary time for C-T systems). Now we are interested to link the controllability property of a system to the ability of designing a state feedback “improving” in some way the resulting closed-loop



## Controllability and state feedback

**Question:** Are controllability properties affected by change of coordinates? For instance, if  $(A, B)$  is completely controllable, is a “similar pair”  $(\tilde{A}, \tilde{B}) = (TAT^{-1}, TB)$  also completely controllable?

$$\begin{array}{ccc} \tilde{R} & = & T R \\ \nearrow & & \swarrow \\ \tilde{R} = (\tilde{B} \ \tilde{A}\tilde{B} \ \dots \ \tilde{A}^{n-1}\tilde{B}) & & R = (B \ AB \ \dots \ A^{n-1}B) \end{array}$$

$$\Rightarrow \text{rank } R = \text{rank } \tilde{R}$$

Controllability/reachability is a “structural property”, not affected by the coordinate framework used to describe the system

**Controllability/reachability of the system**

**Question:** Are controllability properties affected by static state feedback?

**Result**

$$(A, B) \text{ completely controllable} \Rightarrow (A + BK, B) \text{ completely controllable } \forall K$$

**Homework: prove it**

Proving with any  $K$ , a fully controllable system is controllable.

$$(A, B) \quad R = [B \ AB \ A^2B \ A^{n-1}B]$$

$$\begin{array}{c} \downarrow \\ T \end{array}$$

$$R' = \text{image } R$$

$$(\tilde{A}, \tilde{B}) = (TAT^{-1}, TB)$$

$$\tilde{R} = [TB \underbrace{TAT^{-1}}_{\tilde{A}} \underbrace{TB}_{\tilde{B}} \underbrace{TAT^{-2}TAT^{-1}TB}_{\tilde{A}^2} \underbrace{]}_{\tilde{B}}$$

$$\text{Proven } \tilde{R} = TR \quad \text{rank } \tilde{R} = \text{rank } R$$

$$R = T^{-1}\tilde{R}$$

So, with any  $K$ , it does not change the dynamic of the system.

# Controllability and state feedback

Can we then identify a coordinate framework where the design of  $K$  is easier ?!

**Controllability canonical form  $m = 1$**

$$A_c = \begin{pmatrix} 0 & 1 & 0 & \cdots & 0 \\ 0 & 0 & 1 & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & 1 \\ -\alpha_n & -\alpha_{n-1} & -\alpha_{n-2} & \cdots & -\alpha_1 \end{pmatrix} \quad B_c = \begin{pmatrix} 0 \\ 0 \\ \vdots \\ 0 \\ 1 \end{pmatrix}$$

$\xrightarrow{\text{non-zero}} \text{single column} \quad \xrightarrow{\text{on the last row}}$

$\varphi_A(\lambda) = \det(\lambda I - A)$   
 $= \lambda^n + \alpha_1 \lambda^{n-1} + \dots + \alpha_{n-1} \lambda + \alpha_n$

always 1  
on the last row

**Theorem** ( $m = 1$ )

A pair  $(A, B)$  is similar to  $(A_c, B_c)$  iff the system is completely controllable

$\xrightarrow{\text{change of coord. in controllability canonical form}}$

$$T = T_c := R_c R^{-1}$$

$$R_c = T_c R$$

$$B_c = T_c B$$

$$R_c = (B_c \ A_c B_c \ \dots \ A_c^{n-1} B_c)$$

$$R = (B \ AB \ \dots \ A^{n-1} B)$$

$$A_c = T_c A T_c^{-1}$$

# Controllability and state feedback

**Theorem** ( $m \geq 1$ )

A pair  $(A, B)$  is completely controllable iff for all  $\{\lambda_1^*, \dots, \lambda_n^*\}$  (set of desired eigenvalues) there exists a  $K$  such that  $\sigma(A + BK) = \{\lambda_1^*, \dots, \lambda_n^*\}$

eigenvalues assignment theorem

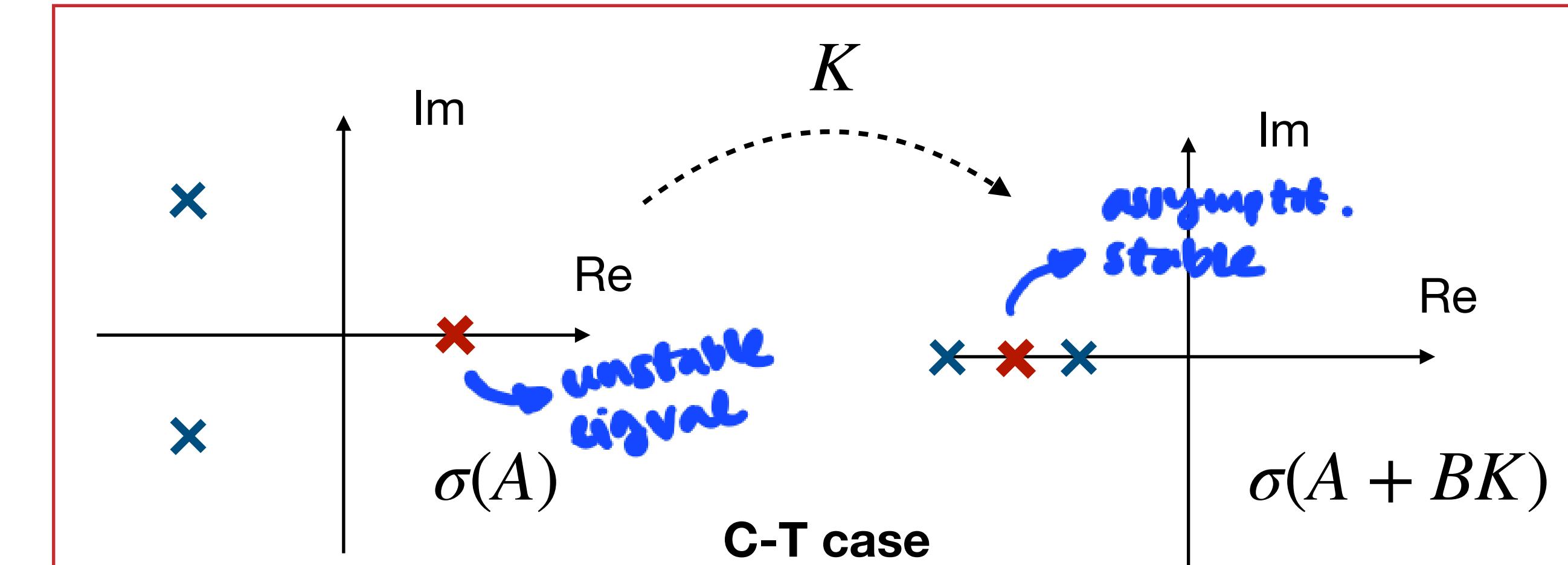
Constructive proof of the if part ( $m = 1$ ):

- Let  $(\alpha_1^*, \dots, \alpha_n^*)$  be such that  $\{\lambda_1^*, \dots, \lambda_n^*\}$  are roots of  $\lambda^n + \alpha_1^* \lambda^{n-1} + \dots + \alpha_{n-1}^* \lambda + \alpha_n^* = 0$
- Let  $K_c = (\alpha_n - \alpha_n^* \ \dots \ \alpha_1 - \alpha_1^*)$
- Pick  $K = K_c T_c$

**Corollary**

A completely controllable system can be always stabilised by static state feedback (...and much more!)

Full authority in eigenvalues assignment



Suppose a system that  $(A, B)$  has a set of eigenvals.  $\{\lambda_1^* \dots \lambda_n^*\} = \sigma(A + BK)$

Recall:

$$K_C = K T_C^{-1}$$

$$K = K_C T_C$$

$$T_C(A + BK) T_C^{-1} = A_C + B_C K T_C^{-1}$$

$$\begin{bmatrix} 0 & 1 & 0 & \dots & 0 \\ 0 & 0 & 1 & \dots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \dots & 1 \\ -a_n & -d_{n-1} & \dots & \dots & -d_1 \end{bmatrix} + \begin{bmatrix} 0 \\ 0 \\ \vdots \\ 0 \\ 1 \end{bmatrix} [K_{c1} \dots K_{cn}]$$

$$= \begin{bmatrix} 0 & 1 & 0 & \dots & 0 \\ 0 & 0 & 1 & \dots & 0 \\ 0 & 0 & 0 & \dots & 1 \\ -a_n + K_{c1} & \dots & \dots & \dots & -d_1 + K_{cn} \\ -a_n^* & & & & -a_1^* \end{bmatrix}$$

The eigenvalues are:

$$\lambda^n + d_2^* \lambda^{n-1} + \dots + d_n^* = 0$$

have roots:

$$\{\lambda_1^*, \dots, \lambda_n^*\}$$

# Kalman decomposition

Suppose now the  $(A, B)$  is not completely controllable, namely  $\text{rank}R = n_r < n$  ( $\dim \mathcal{R}^+ = n_r < n$ )

Let  $\mathcal{R}_\perp^+$  be the orthogonal complement of  $\mathcal{R}^+$ . It turns out that  $\dim \mathcal{R}_\perp^+ = n - n_r$

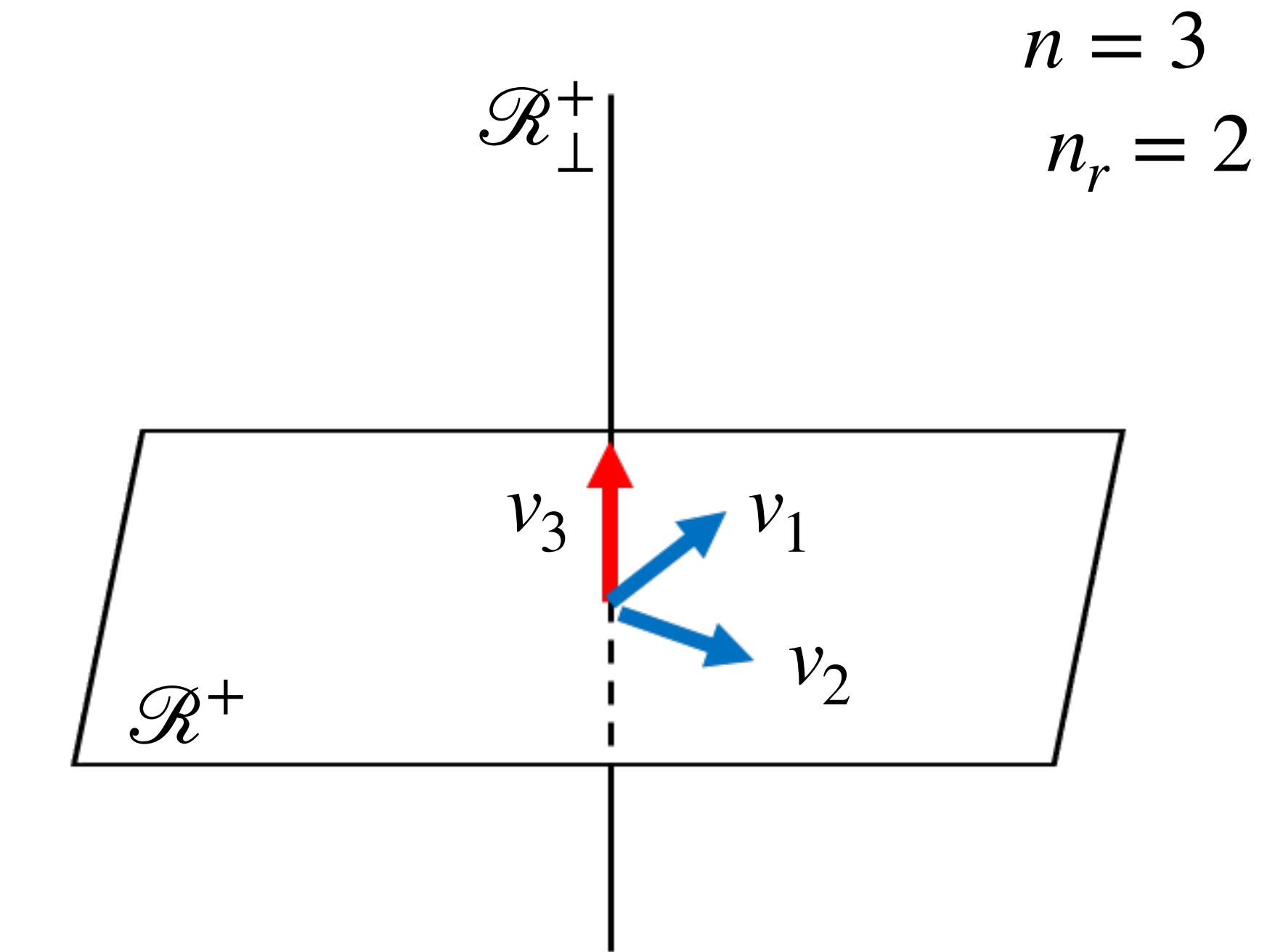
Let  $\{v_1, \dots, v_{n_r}\}$  be a base of  $\mathcal{R}^+$  and let  $\{v_{n_r+1}, \dots, v_n\}$  be a base of  $\mathcal{R}_\perp^+$ . The two sets of vectors are all linearly independent

*base  
of  $\mathcal{R}^+$*

$$\text{Consider the change of variables } T_K^{-1} = \begin{bmatrix} v_1 & \dots & v_{n_r} & v_{n_r+1} & \dots & v_n \end{bmatrix}$$

$$z = T_K x = \begin{pmatrix} z_r \\ z_{nr} \end{pmatrix} \quad x = T_K^{-1} z$$

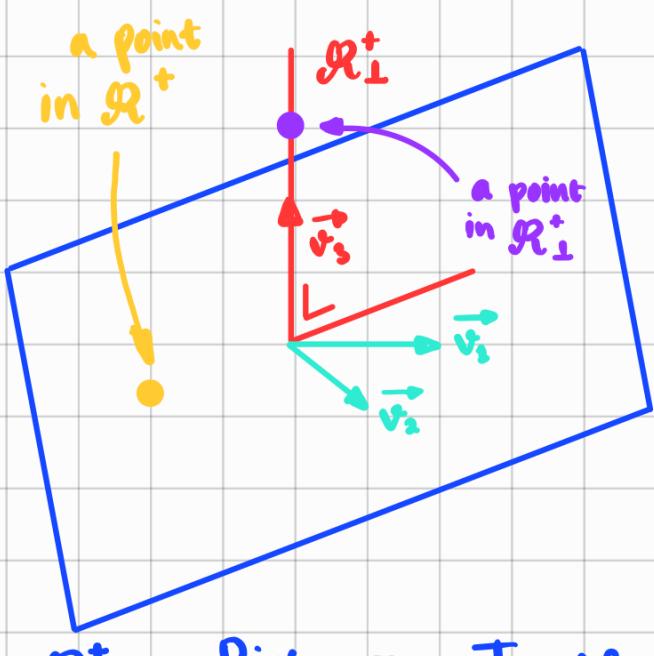
$$\text{It turns out that } x \in \mathcal{R}^+ \Rightarrow z = \begin{pmatrix} \star \\ 0 \end{pmatrix} \quad x \in \mathcal{R}_\perp^+ \Rightarrow z = \begin{pmatrix} 0 \\ \star \end{pmatrix}$$



$$\mathcal{R}^+ = \mathcal{R}^- \subset \mathbb{R}^n$$

rank  $R = n_R < n$   
 $\dim \mathcal{R}^+ = n_R$

Suppose  $n = 3, n_R = 2$



$\{\vec{v}_1, \vec{v}_2\}$  base for  $\mathcal{R}^+$   
 $\{\vec{v}_3\}$  base for  $\mathcal{R}_{\perp}^+$   
 $\{\vec{v}_1, \vec{v}_2, \vec{v}_3\}$  base for  $\mathbb{R}^3$   
 ↪ not necessarily orthogonal but linearly indep.  
 $T^{-1} = [\vec{v}_1, \vec{v}_2, \vec{v}_3]$

Pick  $z = T_K \mathbf{n}, \mathbf{n} = T_K^{-1} z$ , for

an orange point,  $z = \begin{bmatrix} * \\ * \\ 0 \end{bmatrix}$

a violet point,  $z = \begin{bmatrix} 0 \\ 0 \\ * \end{bmatrix}, z = \begin{bmatrix} z_R \\ \vdots \\ z_{n_R} \end{bmatrix} \}^{n_R}$

$$z \in \mathcal{R}^+ \left\{ \begin{array}{l} z_{n_R} = 0 \\ z_R \neq 0 \end{array} \right.$$

$$z \in \mathcal{R}_{\perp}^+ \left\{ \begin{array}{l} z_{n_R} \neq 0 \\ z_R = 0 \end{array} \right.$$

# Kalman decomposition

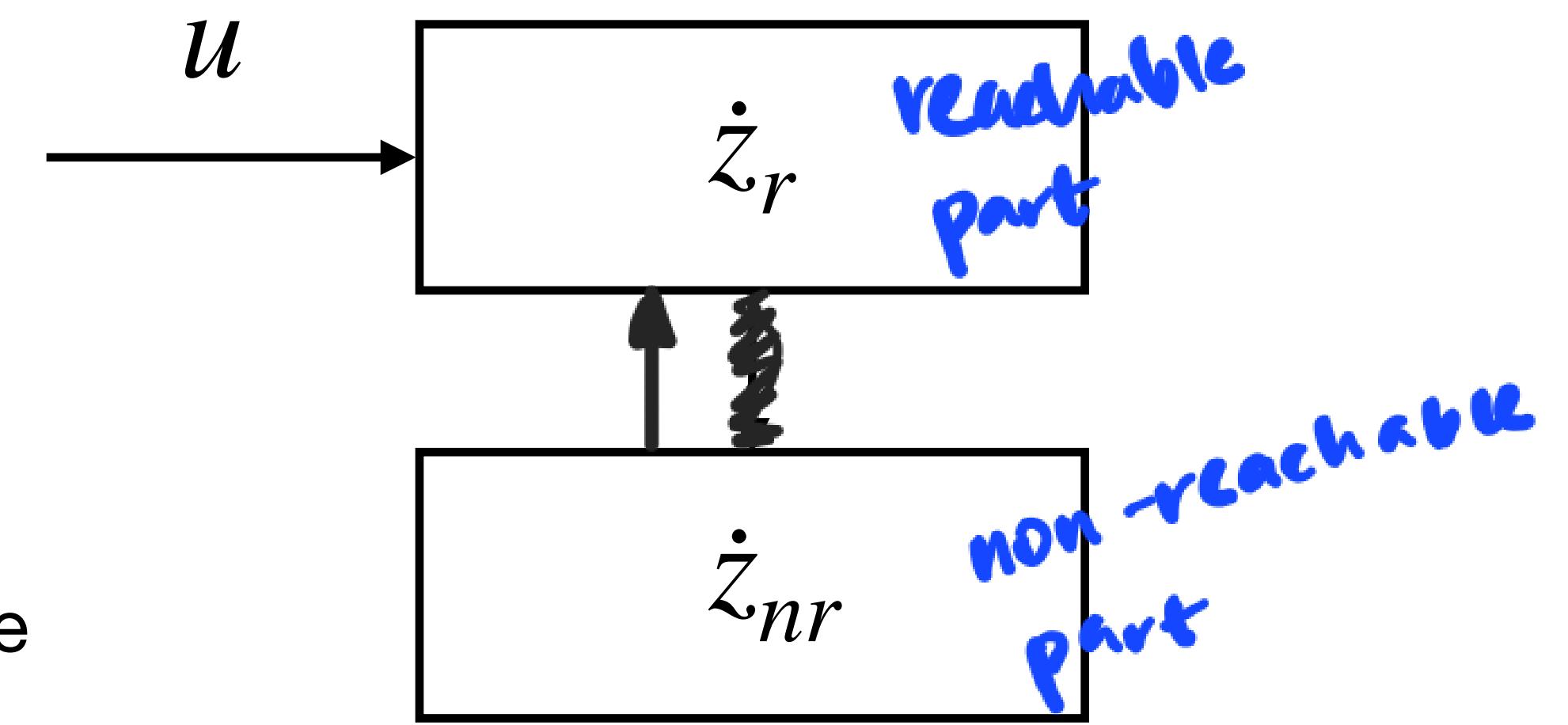
## Result (by Kalman)

$$\tilde{A}_K = T_K A T_K^{-1} = \begin{pmatrix} \tilde{A}_R & \tilde{A}_J \\ 0 & \tilde{A}_{NR} \end{pmatrix}_{n_r \times n_r} \quad \tilde{B}_K = T_K B = \begin{pmatrix} \tilde{B}_R \\ 0 \end{pmatrix}$$

$(\tilde{A}_R, \tilde{B}_R)$  Completely reachable

$$\dot{x} = Ax + Bu \Rightarrow \begin{cases} \dot{z}_r = \tilde{A}_R z_r + \tilde{A}_J z_{nr} + \tilde{B}_R u \\ \dot{z}_{nr} = \tilde{A}_{nr} z_{nr} \end{cases}$$

$\xrightarrow{\text{reachable}}$   $\xrightarrow{\text{non-reachable}}$



- The set  $\mathcal{R}^+$  is forward invariant for the system dynamics. The “internal dynamics” are completely reachable/controllable
- If  $\tilde{A}_{nr}$  is Hurwitz then trajectories starting outside  $\mathcal{R}^+$  asymptotically converge to it

$$\dot{\mathbf{z}} = \mathbf{A}\mathbf{z} + \mathbf{B}\mathbf{u} \quad \mathbf{z} \in \mathbb{R}^n$$

$\downarrow T_k$

$$\dot{\mathbf{z}} = \tilde{\mathbf{A}}_k \mathbf{z} + \tilde{\mathbf{B}}_k \mathbf{u} \quad \mathbf{z} \in \mathbb{R}^n$$

$$\left\{ \begin{array}{l} \dot{\mathbf{z}}_R = \tilde{\mathbf{A}}_R \mathbf{z}_R + \tilde{\mathbf{B}}_R \mathbf{u} + \tilde{\mathbf{A}}_J \mathbf{z}_{NR} \quad \mathbf{z}_R \in \mathbb{R}^{n_R} \\ \dot{\mathbf{z}}_{NR} = \tilde{\mathbf{A}}_{NR} \mathbf{z}_{NR} \end{array} \right.$$

$$\mathbf{z}_{NR}(0) = 0$$

$$\mathbf{z}_{NR}(t) = 0$$

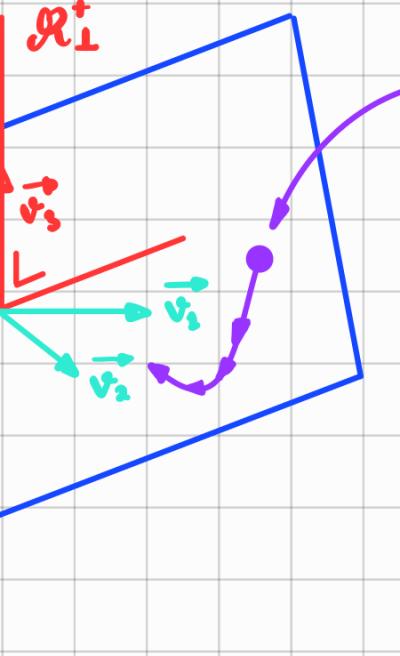
completely  
controllable/reachable

$$\text{For } R_R = [\tilde{\mathbf{B}}_R \tilde{\mathbf{A}}_R \tilde{\mathbf{B}}_R \cdots \tilde{\mathbf{A}}_R^{n_R-1} \tilde{\mathbf{B}}_R],$$

$$\text{rank } R_R = n_R$$

if  $\tilde{\mathbf{A}}_{NR}$   
Hurwitz/Schur,

converge  
to  $\mathcal{P}^+$



For an initial  
condition in  $\mathcal{P}^+$ ,

$$\begin{bmatrix} * \\ 0 \end{bmatrix} \in \mathbb{R}^{n_R} \quad \begin{bmatrix} * \\ 0 \end{bmatrix} \in \mathbb{R}^{n-n_R}$$

$$\dot{\mathbf{z}}_R = \tilde{\mathbf{A}}_R \mathbf{z}_R + \tilde{\mathbf{B}}_R \mathbf{u}$$

Pick  $\alpha(0) \in \mathbb{R}^+$

$\alpha(t) \in \mathbb{R}^+ \forall t, \forall n$

Pick  $\alpha(0) \in \mathbb{R}_L^+$

$z_R(0) = 0, z_{RN}(0) \neq 0,$

$z_R(t) \neq 0, z_{RN}(t) \neq 0, \alpha \notin \mathbb{R}_L^+$

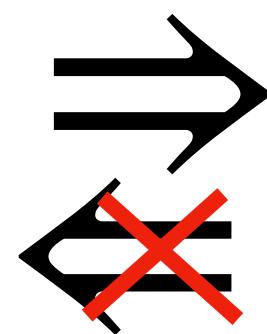
# Stabilisability

## Definition

A system  $(A, B)$  is said to be **stabilizable** if there exists a  $K$  such that  $A + BK$  is Hurwitz (Schur)

## Remark

$(A, B)$  completely reachable



$(A, B)$  stabilizable

## Theorem

$(A, B)$  is stabilizable iff  $\tilde{A}_{nr}$  is Hurwitz (Schur)

**Constructive proof** ( $\tilde{A}_{nr}$  Hurwitz/Schur  $\Rightarrow \exists K : (A + BK)$  is Hurwitz/Schur)

The selection of  $K$  is:  $K = K_K T_K$        $K_K = [K_{Kr} \star]$  with  $K_{Kr}$  so that  $\tilde{A}_r + \tilde{B}_r K_{Kr}$  is Hurwitz

*doesn't play any role*

Suppose that  $(A, B)$  is not controllable but stabilizable. The goal is design  $K$  that makes  $A + BK$  if Hurwitz/Schur.

coord. change on Kalman

$$T_k (A + BK) T_k^{-1} = \tilde{A}_k + \tilde{B}_k K_k \quad \rightarrow \quad K_k = K T_k^{-1}$$

$\underbrace{K T_k^{-1}}$

$$K = K_k T_k$$

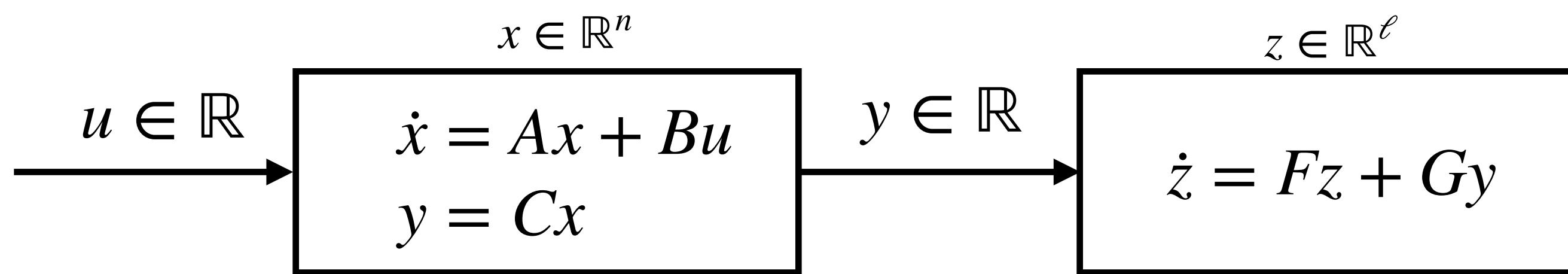
$$\left[ \begin{array}{c|c} \tilde{A}_R & \tilde{A}_J \\ \hline 0 & \tilde{A}_{NR} \end{array} \right] + \left[ \begin{array}{c} \tilde{B}_R \\ 0 \end{array} \right] \left[ \begin{array}{c|c} K_{RR} & K_{RN} \\ \hline 0 & \end{array} \right]$$

assignable

$$\sigma \left[ \begin{array}{c|c} \tilde{A}_R + \tilde{B}_R K_{RR} & \tilde{A}_J + \tilde{B}_R K_{RN} \\ \hline 0 & \tilde{A}_{NR} \end{array} \right] = \sigma(\tilde{A}_R + \tilde{B}_R K_{RR}) \cup \sigma(\tilde{A}_{NR})$$

completely  
controllable

# Controllability of a cascade



$(A, B), (F, G)$  controllable  
↓ ?? **No?**  
 $\begin{pmatrix} A & 0 \\ GC & F \end{pmatrix}$  controllable

## Result

The cascade is controllable iff the pairs  $(A, B), (F, G)$  are controllable and

$$\text{rank} \begin{pmatrix} A - \lambda I & B \\ C & 0 \end{pmatrix} = n + 1 \quad \forall \lambda \in \sigma(F)$$

Non resonance condition