# COMSC-210 Lecture Topic 15 Binary Trees

## ☐ Reference Childs Ch. 15 YouTube DFS

## ■ Binary Tree As Linked Structure

```
data structure: the "root" pointer
bool empty() const {return root == 0;}
the node:

template <class T>
struct Node
{
   T data;
   Node* left;
   Node* right;
}
```

#### ■ Binary Tree Iteration

```
#1 level order (array-based tree only)
all recursive (array-based or linked):
    #2 "Preorder": vertex, left child, right-child (VLR)
    #3 "Inorder": left child, vertex, right-child (LVR)
    #4 "Postorder": left child, right-child, vertex (LRV)

template <class T>
void inorderOutput(const Node* node) const
{
    if (!node) return;
    inorderOutput(node->left);
    cout << node->data;
    inorderOutput(node->right);
```

#### **■** Binary Tree Dynamic Memory Management

generic recursive copy: from the root, using preorder sequence recursive clear: from the root, using postorder sequence copy constructor needs to copy destructor needs to clear assignment operator needs to clear and copy

## ■ Binary Tree Basic Getter Functions

bool empty() const; // true if root is zero
int size() const; // track #of entries in "operator[]" and "deleteKey"
int height() const; // a new one!
 counts number of "levels" PDF

### ☐ Binary Search Trees (BST) PDF

balance issues *view* full trees: 1 3 7 15... complete trees balanced trees *PPT* "horizontal ordering" "inorder" iteration

#### ■ **BST** operator[ ]

the easy way (see sample code)
simple iteration from root, looking to match a key
max #of operations = height of tree, O(log n)

the hard way: maintain balance, and maintain O(log n)
AVL height balancing PDF
balance factor = right-child height - left-child height
...maintain in {-1, 0, +1} PDF
red-black height balancing

## ■ **BST** deleteKey

maintain "horizontal ordering" "seek & swap" algorithm PDF

the hard way: maintain balance, and maintain O(log n) AVL height balancing red-black height balancing