

COMSC-210 Lecture Topic 14

$n \log(n)$ Sorting Algorithms

Reference

Childs Ch. 14

sorting-algorithms.com

No Holes

sort up to a "fill-line"

The "Heapsort" Algorithm

uses a binary "tree" structure

(stored in array, but thought of as a binary tree)

first: arrange array into a heap (highest value at top)

then: repeatedly remove top of heap

to its position in the array

starting from the end,

until the heap is empty.

best suited to arrays

Viewing An Array As A "Heap"

heaps are *arrays*, but all arrays are not heaps

"heap" is another way of looking at an array, as is "table"

each position in array corresponds to a

position in a binary tree

$a[0]$ is the top; $a[1]$ and $a[2]$ are its children

$a[3]$ and $a[4]$ are children of $a[1]$

$a[5]$ and $a[6]$ are children of $a[2]$

...and so it continues, doubling the #of children

in each successive generation

the children of $a[i]$ are $a[2i+1]$ and $a[2i+2]$

Building A Heap

in any array, all elements from $n/2$ to $n-1$

are valid heaps: no children

but including $n/2-1$, it may no longer be a heap

because it has children

so working from $n/2-1$ to 0, one-at-a-time,

consider the array from there to $n-1$ as a heap,

remove the top, promote remaining elements,

and reinsert the removed top

after removing $a[0]$ and reinserting, we have heap!

click [here](#) and [here](#) for array code sample

Sorting The Heap

one-by-one, fill from the back of the array

top of heap is highest value: belongs at end of array

so: (1) save value at end of array,

(2) copy top to end of array

(3) reinsert saved value into heap

repeat until < 3 are left in heap, array is sorted

each cycle, heap size in left half of array is *decreased by 1*

and the sorted array size in the right half is *increased by 1*

click [here](#) for array code sample

click [here](#) for tree-based heapsort visualization

The "Quicksort" Algorithm

arrayed or linked

public function: `void quickSort(int);` // #of elements to sort

divide values into a pivot and 2 separate collections

swap values so that:

...all values in 1st collection are *less than* the pivot

...all values in 2nd collection are *greater*

concatenate 1st collection, pivot, and 2nd collection

recurse 1st and 2nd collections until

...each has *size=1*

suitable for arrays *and* linked-lists

click [here](#) for list code sample

click [here](#) for array code sample

The "Mergesort" Algorithm

divide values into separate collections

subdivide each collection

repeat subdividing until each collection has *size=1*

merge collections into one big collection

suitable for arrays and linked-lists

for lists with *more than one* value:

click [here](#) for code sample

$O(n)$ Sorting!

radix (or "bogo") sort ([ref](#))

click [here](#) and [here](#) for radix sort visualizations