

CS 374 HW 3 Problem 3

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TOTAL POINTS

100 / 100

QUESTION 1

1 3A 40 / 40

✓ - 0 pts Correct

- 30 pts IDK

- 5 pts empty string not accepted by grammar

- 10 pts recursive rule for $0^i 1^i(A)$ incorrect

- 10 pts recursive rule for $1^i 2^i(B)$ incorrect

- 5 pts A: $0^i 1^i$ does not terminate (epsilon rule not added)

- 5 pts B: $1^i 2^i$ does not terminate (epsilon rule not added)

- 10 pts recursive rule to combine A and B wrong

- 40 pts Ineligible: We can't read the answer/ follow the logic/ its way too long. You might want to consider the IDK option in future

- 40 pts Late

QUESTION 2

2 3B 60 / 60

✓ - 0 pts Correct

- 45 pts IDK

- 20 pts L in L(G) wrong

- 5 pts induction variable to prove L(A) incorrect

- 5 pts base case of L(A) induction incorrect

- 5 pts hypothesis of L(A) induction incorrect

- 10 pts inductive proof of L(A) incorrect

- 5 pts any error in proof of L(B), or did not state L(B) belongs to $1^j 2^j$

- 10 pts error in proving $S \rightarrow AB$, or did not state and prove L(G) is in L, by combining L(A) and L(B)

- 60 pts Unreadable/Illegible

(3)

$$\begin{aligned}(A) \quad S &\rightarrow X \cdot Y && \parallel \{0^i 1^j 2^k \mid j = i + k\} \\ X &\rightarrow 0X1 \mid \epsilon && \parallel \{0^i 1^i \mid i \geq 0\} \\ Y &\rightarrow 1Y2 \mid \epsilon && \parallel \{1^j 2^j \mid j \geq 0\}\end{aligned}$$

• Non-Terminals: $S; X; Y$

• Terminals: $0; 1; 2$

- ~~When~~ In production rule S we move onto to production rule X and concatenate these results with the results of going into production rule Y
- In production rule X we add a 1 for every 0 & we can ~~re~~ repeat production rule X to repeat the number of 0's & 1's (ie 00...11) or we can be finished with this production by ~~going to~~ adding nothing
- In production rule Y we add a 2 for every 1 & we can repeat this production rule Y in order to repeat the number of 1's & 2's or we can be finished with the production by ~~go~~ adding nothing

13A 40 / 40

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- 5 pts B: 1^i2^i does not terminate(epsilon rule not added)

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(Q3)

(A) $L = \{0^i 1^j 2^k \mid j = i + k\}$

~~$S \rightarrow X \cdot Y \quad \parallel \{0^i 1^j 2^k \mid j = i + k\}$~~

~~$X \rightarrow 0X1 \mid \epsilon \quad \parallel \{0^i 1^i \mid i \geq 0\}$~~

~~$Y \rightarrow 1Y2 \mid \epsilon \quad \parallel \{1^j 2^j \mid j \geq 0\}$~~

(B) Prove that $w \in L$ iff $S \xrightarrow{*} w$

(\Rightarrow) Suppose $w \in L$, prove $S \xrightarrow{*} w$:

• We have $w \in L \Rightarrow w = 0^i 1^j 2^k = 0^i 1^i 1^k 2^k$ (because $i+k=j$)

• We have $X \rightarrow 0X1 \mid \epsilon$

$\Rightarrow X \xrightarrow{i} 0^i X 1^i$ (Apply $X \rightarrow 0X1$ i times)

and $0^i X 1^i \rightarrow 0^i 1^i$ (Apply $X \rightarrow \epsilon$)

$\Rightarrow X \xrightarrow{i+1} 0^i 1^i$

• We have $Y \rightarrow 1Y2 \mid \epsilon$

$\Rightarrow Y \xrightarrow{k} 0^k Y 1^k$ (Apply $Y \rightarrow 0Y1$ k times)

and $0^k Y 1^k \rightarrow 0^k 1^k$ (Apply $Y \rightarrow \epsilon$)

$\Rightarrow Y \xrightarrow{k+1} 0^k 1^k$

$\Rightarrow S \rightarrow XY \xrightarrow{i+1} 0^i 1^i Y \xrightarrow{k+1} 0^i 1^i 1^k 2^k = w$

Therefore, $S \xrightarrow{*} w$

(\Leftarrow) Suppose $S \xrightarrow{*} w$, prove $w \in L$:

- Prove that: if $X \xrightarrow{*} x$, then $x \in \{0^i 1^i \mid i \geq 0\}$
 - Induction on number of steps in derivation of x .
 - Base case: Derivation with 1 step
The only 1 step is $X \rightarrow \epsilon \Rightarrow x = \epsilon \Rightarrow x \in \{0^i 1^i \mid i \geq 0\}$
 \Rightarrow Base case is true
 - I.H.: Assume if number of steps is smaller or equal n ($n \geq 1$)
then $x \in \{0^i 1^i \mid i \geq 0\}$
- Let number of steps be $(n+1)$, prove $X \xrightarrow{n+1} x \in \{0^i 1^i \mid i \geq 0\}$
Because number of steps $= (n+1) > 1$
 \Rightarrow First step cannot be $X \rightarrow \epsilon$
 \Rightarrow First step must be: $X \rightarrow 0X1$
- Based on I.H., we have:
if $X \xrightarrow{n} x_1$, then $x_1 \in \{0^i 1^i \mid i \geq 0\}$
- Therefore:
 $X \rightarrow 0X1 \xrightarrow{n} 0x_11 = 00^i 1^i 1 = 0^{i+1} 1^{i+1}$
 $\Rightarrow X \xrightarrow{n+1} x = 0^{i+1} 1^{i+1} \Rightarrow x \in \{0^i 1^i \mid i \geq 0\}$
- Repeat the same argument, we have:
if $Y \xrightarrow{*} y$, then $y \in \{1^k 2^k \mid k \geq 0\}$
- We have: $S \rightarrow XY$ and $S \xrightarrow{*} w$
Therefore: $w = x \cdot y$ with $X \xrightarrow{*} x$; $Y \xrightarrow{*} y$.
And we have: $X \xrightarrow{*} x \Rightarrow x \in \{0^i 1^i \mid i \geq 0\}$ (Proven above)
 $Y \xrightarrow{*} y \Rightarrow y \in \{1^k 2^k \mid k \geq 0\}$
 $\Rightarrow w = 0^i 1^i 1^k 2^k$ with $i \geq 0$ and $k \geq 0$
 $\Rightarrow w = 0^i 1^{i+k} 2^k \in L$

Therefore $w \in L$.

Therefore, our grammar is correct!

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