CS 374 HW 0 Problem 1

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TOTAL POINTS

100 / 100

QUESTION 1

1 A 40 / 40

√ - 0 pts Correct

- 20 pts Pigeonhole Principle but not proving n/m
- 10 pts Correct floor/ceiling
- 15 pts Proof by contradiction
- 5 pts M = β α + 1 possible values
- 10 pts Correct conclusion & proper formal

language

- 10 pts IDK

QUESTION 2

2 B 40 / 40

√ - 0 pts Correct - uses part (a)

- 15 pts Differentiating between 0-degree and no 0-

degree cases

- 20 pts There are at most n-1 degree values
- 5 pts Using part (a.)
- 10 pts Bad arithmetic with ceiling
- 10 pts Correct, but not using part a.
- 5 pts Argument by contradiction
- 10 pts All degree values appear once
- 15 pts Contradiction in deg = 0 and deg = n-1
- 10 pts Not using part a.
- 10 pts IDK

QUESTION 3

3 C 20 / 20

√ - 0 pts Correct - multiple components

- 5 pts Noticing every component has at least 2

vertices

- 10 pts The size of a component bounds the

degree of the vertices (define the max)

- 5 pts applying the same argument as in part (b)
- 0 pts Correct can just restrict to a connected

component

- 5 pts Every component has at least 2 vertices
- **5 pts** Noting that the degrees in the component equals the degrees in G.
- 10 pts Arguing correctly that (b) applies
- 5 pts IDK

1)(H)

• Prove that: $\lceil \frac{a+1}{b} \rceil = \lceil \frac{a}{b} \rceil + 1$, if a mod b = 0Q1)(A) la, bare positive => a = b.q (with q is an interger) · Care 1: a mod b = 0 =) $\frac{a+1}{b} = q + \frac{1}{b} = 3$ $q < \frac{a+1}{b} < q+1$ (because b > 0) . We also have that: a mod b = 0 $\frac{a}{b} = q$: an integer Therefore, $\lceil \frac{a+1}{b} \rceil = \lceil \frac{a}{b} \rceil + 1$ • Care 2: a mod $b \neq 0$:=) $a = b \cdot q + r$ (with q and r be integers)

=) a + 1 - r + r + 1 (when $\Rightarrow q \left\langle \frac{a+1}{b} \right\rangle \left\langle q+1 \right\rangle \left(\frac{because 0 \left\langle r \right\rangle \left\langle b \right\rangle}{and} \frac{r+1 \leq b}{b} \right) \leq 1$ • We also have that: $a = b.9 + r = \frac{a}{b} = 9 + \frac{c}{b}$ = 9< $\frac{a}{b}$ <9+1 (because 0<r
f<1) Therefore $\lceil \frac{a+1}{b} \rceil = \lceil \frac{a}{b} \rceil$ (because both = (q+1)) Baselon those 2 cases, we can conclude that:

Induction on n

Base case: n = 0:

$$\left\lceil \frac{0}{\beta - \alpha + 1} \right\rceil = 0$$

In the sequence of 0 number, there's always 0 numbers that are all equal. Therefore, base case is correct

<u>Hypothesis</u>: Assume for any non-negative integer $n \le k$, in the sequence $x_0, x_1, ... x_k$, such that $\alpha \le x_i \le \beta$, there are at least $\lceil n/(\beta - \alpha + 1) \rceil$ numbers that are all equals.

For n = k+1: Need to prove that the sequence $x_0, x_1, ... x_k x_{k+1}$, such that $\alpha \le x_i \le \beta$, has at least $\lceil (k+1)/(\beta - \alpha + 1) \rceil$ numbers that are all equals

Based on the proof above, we have that for any non-negative integers a, b, we have the fact that:

$$\left\lceil \frac{a+1}{b} \right\rceil = \begin{cases} \left\lceil \frac{a}{b} \right\rceil & a \mod b \neq 0 \\ \left\lceil \frac{a}{b} \right\rceil + 1 & a \mod b = 0 \end{cases}$$

Case 1: k mod $(\beta - \alpha + 1) \neq 0$

Because k mod $(\beta - \alpha + 1) \neq 0$, we have $\lceil k/(\beta - \alpha + 1) \rceil = \lceil (k+1)/(\beta - \alpha + 1) \rceil$. Therefore, we only need to prove that the sequence $x_0, x_1, ... x_k, x_{k+1}$, such that $\alpha \leq x_i \leq \beta$, has at least $\lceil k/(\beta - \alpha + 1) \rceil$ numbers that are all equals.

Based on the hypothesis, the sequence $x_0,x_1,...x_k$, such that $\alpha \le x_i \le \beta$, there are at least $\lceil k/(\beta-\alpha+1) \rceil$ numbers that are all equals. Therefore, adding any integer to the sequence will give us the sequence $x_0,x_1,...x_k,x_{k+1}$ that has at least $\lceil k/(\beta-\alpha+1) \rceil$ numbers that are equal. Therefore, the statement is true for this case.

Case 2: k mod $(\beta - \alpha + 1) = 0$, we have $\lceil (k+1)/(\beta - \alpha + 1) \rceil = 1 + \lceil k/(\beta - \alpha + 1) \rceil$. Let m be the maximum number of numbers that are all equal in the sequence $x_0, x_1, ... x_k$. Based on the hypothesis, we have that $m \ge \lceil k/(\beta - \alpha + 1) \rceil$

<u>Case a:</u> $m > \lceil k/(\beta-\alpha+1) \rceil \Leftrightarrow m \ge \lceil k/(\beta-\alpha+1) \rceil + 1$. Therefore, the sequence $x_0,x_1,...x_k$ has at least $\lceil k/(\beta-\alpha+1) \rceil + 1$ numbers that are all equal. Therefore, adding any integer to the sequence will give us the sequence $x_0,x_1,...x_k,x_{k+1}$ that has at least $\lceil k/(\beta-\alpha+1) \rceil$ numbers that are equal. Therefore, the statement is true for this case.

Case b: $m = [k/(\beta - \alpha + 1)] = k/(\beta - \alpha + 1)$ because $k \mod (\beta - \alpha + 1) = 0$. The sequence $x_0, x_1, ... x_k$ has k elements in the sequence with each element in the range $[\beta, \alpha]$, and $m = k/(\beta - \alpha + 1)$. Therefore, each number in the sequence must be repeated exactly m times. By adding any number in range $[\beta, \alpha]$ to that sequence, we have the sequence $x_0, x_1, ... x_k, x_{k+1}$ that has (m+1) numbers that are all equal. Therefore, the sequence $x_0, x_1, ... x_k, x_{k+1}$ has at least $(\lceil k/(\beta - \alpha + 1) \rceil + 1) = \lceil (k+1)/(\beta - \alpha + 1) \rceil$ numbers that are equal.

In all cases, the claim is true.

Therefore, the statement is true for all value of n.

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(21) (B) Because IVI>, 2, there are 3 cares: Carel: There are at least 2 distinct nodes in G that have degree O => There are 2 distinct nodes u and v with deg(u) = deg(v) = 0 => There are 2 distinct nodes u and v that degree of u is equal to degree of v Case 2: There exeisexactly 1 node that has degree O. · For this case, |V| cannot be 2 because if there are only 2 nodes in the graph, and 4 of them has degree 0, the other node must also have degree O · Because there is exactly I node that has degree of, there are (IVI-1) nodes are connected with the degree of each node is at least 1 · Because there's no loop, each vertex in the (171-1) nodes can be adjacent to at most (IVI-2) vertices =) The degree of those connected (IVI-1) is in range [1, IVI-2] · Let n=|V|-1; $\alpha=1$; $\beta=|V|-2$ (we have $\beta \geqslant \alpha$ because |V|>2· Let x1, x1, x1 be the sequence of integers numbers which represent the degree of (IVI-1) connected nodes in G with 2 Exi EB · Based on the result from part (A), we have that there are at least $\lceil \frac{n}{\beta-\alpha+1} \rceil$ numbers in the sequence that reall equ $\left(\frac{n}{\beta-\alpha+1}\right) = \left(\frac{|V|-1}{(|V|-2)-1+1}\right) = \left(\frac{|V|-1}{(|V|-2)}\right) = 2$ Therefore, there're at least 2 numbers in the sequence =) There 're at least 2 distinct nodes that have the same dease. Coex 3: G is a connected sock graph. This means that all vertices · Because there is no loop, and there are IVI nodes (NI>2) in G, in G have degree that 's greater or equal 1.

each vertex can be adjacent to at most (IVI-1) nodes

=> All nodes in G have degree in range [1, 1VI-1]

Repeat the same argument for case 2 above, we have that there are at least $\lceil \frac{|V|}{(|V|-1)-1+1} \rceil = \lceil \frac{|V|}{|V|-1} \rceil = 2$ nodes that have the same degree. For all cases, the claim is tous.

(1 C): Because all vertices in G are of degree at least one, there are at least 2 vertices in G. => 1V1>, 2.

Also, there exist a subgraph G' of G, such that all vertices in G' are connected

Let |V'| be the number of vertices in G'. Because all nodes in G' are connect and has degree of at least 1 =) |V'| >> 2.

Based on part (B), the graph G' has IV'l > 2, therefore there are at least 2 distinct nodes u an v in G' such that degree of u is equal to degree of v. And we have that G' is a sex connected graph. Therefore, there is a simple path between u an v

Therefore, there is a simple path between 2 distinct nodes that have same degree in G'.

Because G'is a sub-graph of G, there's also a simple path between 2 distinct nodes that have the Same degree in G.

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