smsa2

February 11, 2020

1 Lecture 2: Introduction to Linear Algebra and MATLAB

Sound and Music Signal Analysis, SMC8, Aalborg University, 2020

By Jesper Kjær Nielsen (jkn@create.aau.dk), Audio Analysis Lab, Aalborg University.

Last edited: 2020-02-11

Table of Contents

1 Notation

2 Basic operations

2.1 Matrix addition and subtraction

2.2 Matrix multiplication

2.3 (Hermitian) transpose of a matrix

2.4 Matrix Inversion

3 Structured matrices

3.1 Symmetric and Hermitian matrix

3.2 Toeplitz matrix

3.3 Circulant matrix

4 Bonus: The eigenvalue decomposition

4.1 The EVD of special matrices

1.1 Notation

An N-dimensional **column** vector \boldsymbol{v} is written as

$$oldsymbol{v} = egin{bmatrix} v_1 \ v_2 \ dots \ v_N \end{bmatrix}$$

where the v_n 's are all scalers.

An $N \times M$ matrix \boldsymbol{A} is written as

$$\mathbf{A} = \begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1M} \\ a_{21} & a_{22} & \cdots & a_{2M} \\ \vdots & \vdots & \ddots & \vdots \\ a_{N1} & a_{N2} & \cdots & a_{NM} \end{bmatrix}$$

where the a_{nm} 's are all scalers.

Note that

- an N-dimensional column vector is an $N \times 1$ matrix
- an $N \times M$ matrix contains M N-dimensional column vectors $\{a_m\}_{m=1}^M$, i.e.,

$$m{A} = egin{bmatrix} m{a}_1 & m{a}_2 & \cdots & m{a}_M \end{bmatrix}$$
 .

- instead of writing $N \times M$ matrix A, people often write $A \in \mathbb{R}^{N \times M}$ instead where \mathbb{R} describes that all entries are **real-valued** scalars (\mathbb{C} is used for complex-valued scalars).
- if
- -N < M, the matrix is said to be **fat**
- -N=M, the matrix is said to be **square**
- -N>M, the matrix is said to be **skinny**
- please avoid using **row** vectors if you are already using column vectors. It will only confuse the reader (and yourself)!

v1 =

1

2

3

v2 =

1 2 3

```
21, 22, 23;
31, 32, 33;
41, 42, 43;
]
```

A =

```
    11
    12
    13

    21
    22
    23

    31
    32
    33

    41
    42
    43
```

```
[4]: % what size is my matrix?
size(v1)
size(v2)
size(A)
```

ans =

3 1

ans =

1 3

ans =

4 3

Example: writing a finite data set as a vector Assume that we observe x(n) for $n=0,1,\ldots,N-1$.

How can we write this as a vector?

We simply stack the observations as

$$\boldsymbol{x} = \begin{bmatrix} x(0) \\ x(1) \\ \vdots \\ x(N-1) \end{bmatrix} .$$

ans =

20480 1

1.2 Basic operations

1.2.1 Matrix addition and subtraction

You add/subtract two matrices $\mathbf{A} \in \mathbb{C}^{N \times M}$ and $\mathbf{B} \in \mathbb{C}^{N \times M}$ by doing adding/subtracting all entries elementwise, i.e.,

$$\mathbf{A} \pm \mathbf{B} = \begin{bmatrix}
a_{11} & a_{12} & \cdots & a_{1M} \\
a_{21} & a_{22} & \cdots & a_{2M} \\
\vdots & \vdots & \ddots & \vdots \\
a_{N1} & a_{N2} & \cdots & a_{NM}
\end{bmatrix} \pm \begin{bmatrix}
b_{11} & b_{12} & \cdots & b_{1M} \\
b_{21} & b_{22} & \cdots & b_{2M} \\
\vdots & \vdots & \ddots & \vdots \\
b_{N1} & b_{N2} & \cdots & b_{NM}
\end{bmatrix} \\
= \begin{bmatrix}
a_{11} \pm b_{11} & a_{12} \pm b_{12} & \cdots & a_{1M} \pm b_{1M} \\
a_{21} \pm b_{21} & a_{22} \pm b_{22} & \cdots & a_{2M} \pm b_{2M} \\
\vdots & \vdots & \ddots & \vdots \\
a_{N1} \pm b_{N1} & a_{N2} \pm b_{N2} & \cdots & a_{NM} \pm b_{NM}
\end{bmatrix}$$
(1)

Z =

Some basic rules of matrix additions and subtractions are

$$A + (B + C) = (A + B) + C \tag{3}$$

$$A + B = B + A \tag{4}$$

$$A - A = 0 \tag{5}$$

where $\mathbf{0}$ is the zero-matrix given by

$$\mathbf{0} = \begin{bmatrix} 0 & \cdots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \cdots & 0 \end{bmatrix} .$$

Therefore, you can treat matrices as normal scalars when you add and subtract them.

1.2.2 Matrix multiplication

If we multiply a matrix $\mathbf{A} \in \mathbb{R}^{N \times M}$ by a scalar c, we obtain

$$c\mathbf{A} = c \begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1M} \\ a_{21} & a_{22} & \cdots & a_{2M} \\ \vdots & \vdots & \ddots & \vdots \\ a_{N1} & a_{N2} & \cdots & a_{NM} \end{bmatrix} = \begin{bmatrix} ca_{11} & ca_{12} & \cdots & ca_{1M} \\ ca_{21} & ca_{22} & \cdots & ca_{2M} \\ \vdots & \vdots & \ddots & \vdots \\ ca_{N1} & ca_{N2} & \cdots & ca_{NM} \end{bmatrix}.$$

Thus, we simply multiply every entry in A with c.

Some basic rules for multiplying a scalar with a matrix are

$$c\mathbf{A} = \mathbf{A}c\tag{6}$$

$$(c_1c_2)\mathbf{A} = c_1(c_2\mathbf{A}) = c_2(c_1\mathbf{A})$$
 (7)

$$1A = A \tag{8}$$

$$(c_1 \pm c_2)\mathbf{A} = c_1\mathbf{A} \pm c_2\mathbf{A} \tag{9}$$

$$c(\mathbf{A} \pm \mathbf{B}) = c\mathbf{A} \pm c\mathbf{B} . \tag{10}$$

Active break: writing a phasor of finite duration as a vector We would like to write

$$x(n) = Ae^{j(\omega_0 n + \phi)}$$

for n = 0, 1, ..., N - 1 in matrix form.

First, we write x(n) as

$$x(n) = Ae^{j\phi}e^{j\omega_0 n} .$$

Let us now define $c = Ae^{j\phi}$ and

$$\boldsymbol{x} = \begin{bmatrix} x(0) \\ x(1) \\ \vdots \\ x(N-1) \end{bmatrix} \qquad \boldsymbol{z}(\omega_0) = \begin{bmatrix} 1 \\ e^{j\omega_0} \\ \vdots \\ e^{j\omega_0(N-1)} \end{bmatrix} . \tag{11}$$

Then,

$$\boldsymbol{x} = c\boldsymbol{z}(\omega_0)$$
.

You multiply two matrices $\boldsymbol{A} \in \mathbb{R}^{N \times K}$ and $\boldsymbol{B} \in \mathbb{R}^{K \times M}$ by

$$\mathbf{AB} = \begin{bmatrix} a_{11} & \cdots & a_{1K} \\ \vdots & \ddots & \vdots \\ a_{N1} & \cdots & a_{NK} \end{bmatrix} \begin{bmatrix} b_{11} & \cdots & b_{1M} \\ \vdots & \ddots & \vdots \\ b_{K1} & \cdots & b_{KM} \end{bmatrix}
= \begin{bmatrix} \sum_{k=1}^{K} a_{1k} b_{k1} & \cdots & \sum_{k=1}^{K} a_{1k} b_{kM} \\ \vdots & \ddots & \vdots \\ \sum_{k=1}^{K} a_{Nk} b_{k1} & \cdots & \sum_{k=1}^{K} a_{Nk} b_{kM} \end{bmatrix} .$$
(12)

$$= \begin{bmatrix} \sum_{k=1}^{K} a_{1k} b_{k1} & \cdots & \sum_{k=1}^{K} a_{1k} b_{kM} \\ \vdots & \ddots & \vdots \\ \sum_{k=1}^{K} a_{Nk} b_{k1} & \cdots & \sum_{k=1}^{K} a_{Nk} b_{kM} \end{bmatrix} .$$
 (13)

Note that

- matrix multiplication is **not** elementwise multiplication
- dimensions have to fit, i.e., the number of columns in A and the number of rows in B must be the same.

A very important matrix is the **identity matrix** which is a **square** $N \times N$ matrix given by

$$m{I}_N = egin{bmatrix} 1 & 0 & \cdots & 0 \ 0 & 1 & \cdots & 0 \ dots & dots & \ddots & dots \ 0 & 0 & \cdots & 1 \end{bmatrix} \;.$$

The identity matrix corresponds to a '1' in the scalar case since

$$I_N A = A \tag{14}$$

$$\mathbf{A}\mathbf{I}_{M} = \mathbf{A} \tag{15}$$

when $\mathbf{A} \in \mathbb{R}^{N \times M}$.

ans =

8 5 -1 0

ans =

ans =

Some basic rules for multiplying matrices (assuming that the dimensions fit) are

$$(c\mathbf{A})\mathbf{B} = \mathbf{A}(c\mathbf{B}) = c(\mathbf{A}\mathbf{B}) \tag{16}$$

$$A(B+C) = AB + AC \tag{17}$$

$$(A+B)C = AC + BC \tag{18}$$

$$(AB)C = A(BC) \tag{19}$$

$$AB \neq BA$$
. (20)

1.2.3 (Hermitian) transpose of a matrix

The transpose $(\cdot)^T$ of a matrix $\mathbf{A} \in \mathbb{C}^{N \times M}$ flips the row and column indices, i.e.,

$$\mathbf{A} = \begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1M} \\ a_{21} & a_{22} & \cdots & a_{2M} \\ \vdots & \vdots & \ddots & \vdots \\ a_{N1} & a_{N2} & \cdots & a_{NM} \end{bmatrix}$$
 (21)

$$\mathbf{A} = \begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1M} \\ a_{21} & a_{22} & \cdots & a_{2M} \\ \vdots & \vdots & \ddots & \vdots \\ a_{N1} & a_{N2} & \cdots & a_{NM} \end{bmatrix}$$

$$\mathbf{A}^{T} = \begin{bmatrix} a_{11} & a_{21} & \cdots & a_{N1} \\ a_{12} & a_{22} & \cdots & a_{N2} \\ \vdots & \vdots & \ddots & \vdots \\ a_{1M} & a_{2M} & \cdots & a_{NM} \end{bmatrix} .$$

$$(21)$$

The **Hermitian** transpose $(\cdot)^H$ of a matrix $\mathbf{A} \in \mathbb{C}^{N \times M}$ takes the transpose and the complex conjugate of all elements, i.e.,

$$\mathbf{A}^{H} = (\mathbf{A}^{T})^{*} = \begin{bmatrix} a_{11}^{*} & a_{21}^{*} & \cdots & a_{N1}^{*} \\ a_{12}^{*} & a_{22}^{*} & \cdots & a_{N2}^{*} \\ \vdots & \vdots & \ddots & \vdots \\ a_{1M}^{*} & a_{2M}^{*} & \cdots & a_{NM}^{*} \end{bmatrix} . \tag{23}$$

Note that

ullet taking the transpose and Hermitian transpose are equivalent if all elements in $oldsymbol{A}$ are real-valued

Z =

```
11.0000 +11.0000i 12.0000 +12.0000i
21.0000 +21.0000i 22.0000 +22.0000i
31.0000 +31.0000i 32.0000 +32.0000i
```

ans =

```
11.0000 +11.0000i 21.0000 +21.0000i 31.0000 +31.0000i 12.0000 +12.0000i 22.0000 +22.0000i 32.0000 +32.0000i
```

ans =

```
11.0000 -11.0000i 21.0000 -21.0000i 31.0000 -31.0000i 12.0000 -12.0000i 22.0000 -22.0000i 32.0000 -32.0000i
```

Some basic rules for the Hermitian/transpose of a matrix are

$$(\boldsymbol{A}^T)^* = \boldsymbol{A}^H \tag{24}$$

$$(\mathbf{A}^H)^H = \mathbf{A} \tag{25}$$

$$(\mathbf{A} + \mathbf{B})^H = \mathbf{A}^H + \mathbf{B}^H \tag{26}$$

$$(c\mathbf{A})^H = c^* \mathbf{A}^H \tag{27}$$

$$(\mathbf{A}\mathbf{B})^H = \mathbf{B}^H \mathbf{A}^H \tag{28}$$

$$(\mathbf{A}^H)^T = (\mathbf{A}^T)^H = \mathbf{A}^* .$$
 (29)

1.2.4 Matrix Inversion

An $N \times N$ matrix **A** is invertible, if another $N \times N$ matrix **B** exists such that

$$AB = BA = I_N$$
.

Note that

- **B** is the **inverse** of **A** and is, therefore, often denoted as A^{-1}
- only square matrices can have an inverse (general matrices might have a pseudo-inverse)
- if A has full rank, then it has an inverse. A matrix has full rank if, e.g.,
 - all its columns are linearly independent
 - all eigenvalues are non-zero

```
[9]: % create a matrix with random numbers (generated from a normal distribution)
A = randn(4,4);
% compute the inverse
iA = inv(A)
% check if it is the inverse (AB must be the indentity matrix)
A*iA
```

iA =

ans =

Some basic rules for the matrix inverse are

$$(A^{-1})^{-1} = A (30)$$

$$(\mathbf{A}\mathbf{B})^{-1} = \mathbf{B}^{-1}\mathbf{A}^{-1} \tag{31}$$

$$(\mathbf{A}^{H})^{-1} = (\mathbf{A}^{-1})^{H} \tag{32}$$

$$(c\mathbf{A})^{-1} = c^{-1}\mathbf{A}^{-1} \tag{33}$$

$$\begin{bmatrix} a & b \\ c & d \end{bmatrix}^{-1} = \frac{1}{ad - cb} \begin{bmatrix} d & -b \\ -c & a \end{bmatrix} . \tag{34}$$

To be true, the above rules require that the inverses exist. For the last rule, this means that $ad \neq cb$.

Example: the **DFT** as a matrix operation The N-point discrete Fourier transform (DFT) of a signal x(n) can be written as

$$X(k) = \sum_{n=0}^{N-1} x(n) e^{-j2\pi kn/N}$$

for k = 0, 1, ..., N - 1. We would like to

- 1. write X(k) as multiplication of a data vector \boldsymbol{x} and a phasor vector \boldsymbol{f}_k
- 2. write X(k) for k = 0, 1, ..., N 1 as a matrix-vector product
- 1. If we define

$$\mathbf{x} = \begin{bmatrix} x(0) & x(1) & \cdots & x(N-1) \end{bmatrix}^T$$
 (35)
 $\mathbf{f}_k = \begin{bmatrix} 1 & e^{-j2\pi k/N} & \cdots & e^{-j2\pi k(N-1)/N} \end{bmatrix}^T$, (36)

$$\mathbf{f}_k = \begin{bmatrix} 1 & e^{-j2\pi k/N} & \cdots & e^{-j2\pi k(N-1)/N} \end{bmatrix}^T$$
, (36)

we can write X(k) as the inner product

$$X(k) = \boldsymbol{f}_k^T \boldsymbol{x}$$
.

2. If we define the matrix

$$m{F} = egin{bmatrix} m{f}_0^T \ m{f}_1^T \ dots \ m{f}_{N-1}^T \end{bmatrix} \; ,$$

we get

$$X = Fx$$

where

$$\boldsymbol{X} = \begin{bmatrix} X(0) & X(1) & \cdots & X(N-1) \end{bmatrix}^T$$
.

X1 =

4.1918 + 0.0000i

-1.5332 - 0.0725i

-1.6220 + 0.0000i

-1.5332 + 0.0725i

X2 =

4.1918 + 0.0000i

-1.5332 - 0.0725i

-1.6220 + 0.0000i

-1.5332 + 0.0725i

1.3 Structured matrices

Matrices with structure are extremely important since such as structure - can say something about the properties of the matrix - can be exploited in fast algorithm

1.3.1 Symmetric and Hermitian matrix

A symmetric matrix $\mathbf{A} \in \mathbb{R}^{N \times N}$ satisfies that

$$\boldsymbol{A}^T = \boldsymbol{A}$$
.

That is, the entries in \boldsymbol{A} all satisfy

$$a_{nm}=a_{mn}$$
.

A Hermitian matrix $\boldsymbol{Z} \in \mathbb{C}^{N \times N}$ satisfies that

$$\mathbf{Z}^H = \mathbf{Z}$$
.

That is, the entries in Z satisfy

$$z_{nm} = z_{mn}^*$$

Note that a symmetric/Hermitian matrix is always a square matrix.

1.3.2 Toeplitz matrix

A **Toeplitz** matrix is characterised by that all diagonals of the matrix. As an example, a square Toeplitz matrix $T \in \mathbb{R}^{N \times N}$ can be written as

$$T = \begin{bmatrix} t_0 & t_{-1} & t_{-2} & \cdots & \cdots & t_{-(N-1)} \\ t_1 & t_0 & t_{-1} & \ddots & & \vdots \\ t_2 & t_1 & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & t_{-1} & t_{-2} \\ \vdots & & \ddots & t_1 & t_0 & t_{-1} \\ t_{N-1} & \cdots & \cdots & t_2 & t_1 & t_0 \end{bmatrix}$$

Note that

- a Toeplitz matrix does not have to be square
- is completely specified from the elements in the first column and row

```
[11]: % we specify a Toeplitz matrix via the first column and row of the matrix

T1 = toeplitz([0;1;2],[0,-1,-2,-3])

% if only the first argument is given, a Hermitian Toeplitz matrix is given

T2 = toeplitz([0,1+1i*2])
```

T1 =

0 -1 -2 -3

1 0 -1 -2

T2 =

0.0000 + 0.0000i 1.0000 + 2.0000i
1.0000 - 2.0000i 0.0000 + 0.0000i

1.3.3 Circulant matrix

A **circulant** matrix $C \in \mathbb{R}^{N \times N}$ is a special Toeplitz matrix in which the entries in a column is down-shifted by one compared to the previous column, i.e.,

$$C = \begin{bmatrix} c_0 & c_{N-1} & \cdots & c_2 & c_1 \\ c_1 & c_0 & c_{N-1} & & c_2 \\ \vdots & c_1 & c_0 & \ddots & \vdots \\ c_{N-2} & & \ddots & \ddots & c_{N-1} \\ c_{N-1} & c_{N-2} & \cdots & c_1 & c_0 \end{bmatrix}$$

Note that

- C^T is also a **circulant** matrix
- problems involving circulant matrices can usually be solved extremely efficiently using **an FFT** algorithm (more on this later)

[12]: % create a circulant matrix. Note that you specify the first row in MATLAB - not⊔
→column.

% Therefore, the output is transposed
C = gallery('circul', [0;1;2;3]).'

C =

0 3 2 1
1 0 3 2
2 1 0 3

1.4 Bonus: The eigenvalue decomposition

A square matrix $\mathbf{A} \in \mathbb{C}^{N \times N}$ matrix has an **eigenvalue decomposition** (EVD) if it can be written as

$$AU = U\Lambda$$

or, equivalently,

$$A = U\Lambda U^{-1}$$

where

- $U \in \mathbb{C}^{N \times N}$ is a full rank matrix containing the **eigenvectors** as column vectors
- $\Lambda \in \mathbb{C}^{N \times N}$ is a diagonal matrix containing the **eigenvalues**.

Note that not all square matrices have an EVD. Matrices without an EVD are said to be **defective**.

The EVD is useful for many things such as computing matrix powers since

$$\mathbf{A} = \mathbf{U}\Lambda\mathbf{U}^{-1} \tag{37}$$

$$A^{2} = U\Lambda U^{-1}U\Lambda U^{-1} = U\Lambda^{2}U^{-1}$$
(38)

$$A^n = U\Lambda^n U^{-1} \tag{39}$$

```
[13]: % compute a random matrix
A = randn(3,3);
% compute the EVD
[U, L] = eig(A)
```

U =

L =

1.4.1 The EVD of special matrices

• Symmetric matrices: All symmetric matrices have an eigenvalue decomposition. Moreover, all eigenvalues are real-valued and U is an orthogonal matrix, i.e.,

$$\boldsymbol{U}^{-1} = \boldsymbol{U}^T$$

so that

$$\boldsymbol{A} = \boldsymbol{U}\boldsymbol{\Lambda}\boldsymbol{U}^{-1} = \boldsymbol{U}\boldsymbol{\Lambda}\boldsymbol{U}^T .$$

• Circulant matrices: All circulant matrices have an eigenvalue decomposition. Moreover,

$$\boldsymbol{U} = \boldsymbol{F}^{-1} = N^{-1} \boldsymbol{F}^{H} \tag{40}$$

$$\mathbf{\Lambda} = \operatorname{diag}(\mathbf{F}\mathbf{c}) \tag{41}$$

where

- F is the **DFT matrix** whose (k+1, n+1)'th element is given by

$$[\mathbf{F}]_{k+1,n+1} = e^{-j2\pi nk/N}$$

for $n, k = 0, 2, \dots, N - 1$.

-c is the first column of C.

Consequently, the EVD of a circulant matrix $\boldsymbol{C} \in \mathbb{C}^{N \times N}$ is

$$C = U\Lambda U^{-1} = N^{-1}F^H \operatorname{diag}(Fc)F$$
.

We say that the circulant matrix is **diagonalised** by the DFT. Since the DFT can be computed efficiently using an FFT, this enables the use of fast algorithms in many situations.