# **Real Time Programming** with Ada

Part 2: Real time facilities

## System Time

A timer circuit programmed to interrupt the processor at fixed rate.

- To approximate the universial time
- For distributed systems, we need clock synchronization

Each time interrupt is called a system tick (time resolution):

- Normally, the tick can vary 1-50ms, even microseconds in RTOS
- Linux 2.4, 10ms (100HZ), Linux 2.6, 1ms (1000HZ)
- The tick may be selected by the user
- All time parameters for tasks should be the multiple of the tick
- System time = 32 bits

  - One tick = 1ms: your system can run 50 days
    One tick = 20ms: your system can run 1000 days = 2.5 years
    One tick = 50ms: your system can run 2500 days = 7 years
- In Ada95, it is required that the system time should last at least 50 years

#### Real Time Programming: we need support for

- · Concurrency (Ada tasking)
- · Communication & synchronization (Ada Rendezvous)
- Consistency in data sharing (Ada protected data type)
- Real time facilities (Ada real time packages and delay statements) - accessing system time so that the passage of time can be measured
  - delaying processes until some future time
  - Timeouts: waiting for or running some action for a given time period

## Real-Time Support in Ada

- · Two pre-defined packages to access the system clock
  - Ada.Calendar and Ada.Real\_Rime
  - Both based on the same hardware clock
- · There are two delay-statements
  - Delay time\_expression (in seconds)
  - Delay until time\_expression
- · The delay statements can be used together with select to program timeouts, timed entry etc.

#### Package calendar in Ada: specification

```
package Ada.Calendar is
       ackage Ada.Calendar is

type Time is private;

--- time is pre-defined based on the system clock
subtype Year_Number is Integer range 1901 .. 2099;
subtype Month_Number is Integer range 1 .. 12;
subtype Day_Number is Integer range 1 .. 31;
subtype Day_Duration is Duration range 0.0 .. 86_400.0;

--- Duration is pre-defined type (length of interval,

--- expressed in sec's) declared in the package: Standard
function Clock return Time;
function Year (Date: Time) return Year Number;
        function Clock return !:me; function Year (Date : Time) return Year Number; function Month (Date : Time) return Month Number; function Day (Date : Time) return Day_Number; function Seconds (Date : Time) return Day_Duration;
        function Seconds(Date : Time) return Day_Dural
procedure Split (Date : in Time;
Year : out Year_Number;
Month : out Month Number;
Day : out Day_Number;
Seconds : out Day_Duration);
```

## Package calendar in Ada: specification (ctn.)

```
function Time Of (Year
                                                 : Year Number;
                                    Month : Month_Number;
                                   Day : Day_Number;
Seconds : Day_Duration := 0.0)
    return Time:
    function "+" (Left : Time; Right : Duration) return Time;
    function "+" (Left: Duration; Right: Time) return Time;
function "-" (Left: Time; Right: Duration) return Time;
function "-" (Left: Time; Right: Duration) return Duration;
   function "-" (Left : Time; Right : Time) return
function "<" (Left, Right : Time) return Boolean;
function "<=" (Left, Right : Time) return Boolean;
function ">=" (Left, Right : Time) return Boolean;
function ">=" (Left, Right : Time) return Boolean;
Time_Error: exception;
private
      -- not specified by the language
      -- implementation dependent
end Ada.Calendar:
```

#### Package Real\_Time in Ada: specification

```
package Ada.Real_Time is

type Time is private;

Time First : constant Time;

Time Jast : constant Time;

Time Unit : constant Time;

Time Unit : constant Time;

Time Span is private;

--- as Duration, a Time_Span value M representing the length of an interval, corresponding to the length of an interval, corresponding to the real time duration M*Time_Unit.

Time_Span First : constant Time_Span;

Time_Span_Zero : constant Time_Span;

Time_Span_Zero : constant Time_Span;

Time_Span_Zero : constant Time_Span;

Time_Span_Zero : constant Time_Span;

Time_Span_Constant Time_Span;

Time_Span_Unit : constant Time_Span;

Timetion Clock return Time;

function "+" (Left : Time; Right : Time_Span) return Time;

function "+" (Left : Time; Right : Time_Span) return Time;

function "-" (Left : Time; Right : Time) return Time_Span;

function "<="(Left, Right : Time) return Boolean;

function ">" (Left, Right : Time) return Boolean;
```

## **Programming Delays**

#### Semantics of Delay(20); Action Granularity difference Time specified by between clock and program delay Executing the 20 sec Ready to run here Action schedulled Interrupts disabled Time

#### Package Real\_Time in Ada: specification (cnt.)

```
function "+" (Left, Right : Time Span) return Time Span; function "-" (Left, Right : Time Span) return Time Span; function "-" (Right : Time Span) return Time Span; function "-" (Left : Time Span) return Time Span; function "*" (Left : Time Span) Right : Integer) return Time Span; function "*" (Left : Time Span) Right : Integer; Right : Time Span) return Time Span; function "/" (Left : Right : Time Span) return Time Span; function "/" (Left : Time Span; Right : Integer) return Time Span; function "C" (Left : Right : Time Span) return Boolean; function "C" (Left : Right : Time Span) return Boolean; function "C" (Left : Right : Time Span) return Boolean; function "C" (Left : Right : Time Span) return Boolean; function "C" (Left : Right : Time Span) return Boolean; function To Time Span (D : Duration) return Time Span; function To Time Span (D : Duration) return Time Span; function Microseconds (US : Integer) return Time Span; function Time Span (SC) cout Seconds Count; TS : Time Span) return Time Span; function Microseconds Count Span; function Time Span; function Time Span; function Microseconds Count Span; function Time Span; function Microseconds Count Span; fu
```

## Relative Delays

- · Delay the execution of a task for a given period
- · Relative delays (using clock access)

```
Start := Clock;
loop
  exit when (Clock - Start) > 10.0; -- bust waiting
end loop;
ACTION;
```

- To avoid busy-waiting, most languages and OS provide some form of delay primitive
  - In Ada, this is a delay statement delay 10.0In UNIX, sleep(10)

#### Absolute Delays

 To delay the execution of a task to a specified time point (using clock access):

```
Start := Clock;
FIRST ACTION;
loop
  exit when Clock > Start+10.0; -- busy waiting
end loop;
SECOND_ACTION;
```

To avoid busy-wait (access "clock" all time every tick!):

```
START := Clock;
FIRST ACTION;
delay until START + 10.0; (this is by interrupt)
SECOND ACTION;
```

• As with delay, delay until is accurate only in its lower bound

#### Absolute Delays: Example

```
task Ticket_Agent is
entry Registration(...);
end Ticket_Agent;
task body Ticket_Agent is
  -- declarations
  Shop_Open : Boolean := True;
begin
  while Shop_Open loop
    select
      accept Registration(...) do
            log details
      end Registration;
      delay until Closing_Time;
            Open := False;
    end select;
      - process registrations
end loop;
end Ticket Agent;
```

## Periodic Task

```
task body Periodic_T is
Next_Release : Time;
ReleaseInterval : Duration := 10
begin
Next_Release := Clock + ReleaseInterval;
loop
--Action
delay until Next_Release;
Next_Release := Next_Release + ReleaseInterval;
end loop;
end Periodic T;

Will run on average every 10 seconds

If Action takes 11 seconds, the delay
statement will have no effect

Will run on average local drift only
```

#### Control Example I

```
with Ada.Real_Time; use Ada.Real_Time;
with Data_Types; use Data_Types;
with IO; use IO;
with Control_Procedures;
use Control_Procedures;

procedure Controller is
   task Temp_Controller;
   task Pressure_Controller;
```

#### Control Example II

```
task body Temp_Controller is
  TR : Temp_Reading; HS : Heater_Setting;
  Next : Time;
  Interval : Time_Span := Milliseconds(30);
begin
  Next := Clock; -- start time
  loop
    Read(TR);
    Temp_Convert(TR,HS);
    Write(HS);
    Write(TR);
    Next := Next + Interval;
    delay until Next;
  end loop;
end Temp_Controller;
```

## Control Example III

```
task body Pressure_Controller is
   PR : Pressure_Reading; PS : Pressure_Setting;
   Next : Time;
   Interval : Time_Span := Milliseconds(70);
begin
   Next := Clock; -- start time
   loop
        Read(PR);
        Pressure_Convert(PR,PS);
        Write(PS);
        Write(PR);
        Next := Next + Interval;
        delay until Next;
        end loop;
   end Pressure_Controller;
begin
   null;
end Controller;
```

#### Control Example IIII

```
task body Pressure_Controller is
   PR : Pressure_Reading; PS : Pressure_Setting;
    Next : Time;
     Interval : Time Span := Milliseconds(70);
     Next := Clock; -- start time
    loop
       Read(PR);
       Pressure_Convert(PR,PS);
       Write(PS);
Write(PR);
       Next := Next + Interval;
    delay until Next;
end loop;
                                                  Here Temp_Controller & Pressure_Controller
  end Pressure_Controller;
begin
                                                  start concurrently
  null;
end Controller;
```

## Timeout and message passing

```
loop
select
accept Call(T: temperature) do
New_temp:=T;
end Call;
or
delay 10.0;
--action for timeout
end select;
--other actions
end loop;
```

## **Programming Timeouts**

## Timeout (by server)

```
task Server is
  entry Call(T: in Temperature);
  -- other entries
end Server;

task body Server is
  -- declarations
begin
  loop
    select
    accept Call(T: in Temperature) do
        New_Temp := T;
    end Call;
    or
        delay 10.0;
        -- action for timeout
    end select;
    -- other actions
    end loop;
end Server;
```

#### Timeout (by client)

```
loop
   -- get new temperature T
   Server.Call(T);
end loop;
loop
   -- get new temperature T
   select
        Server.Call(T);
   or
        delay 0.5;
        -- other actions
   end select;
end loop;
```

#### Timeouts on Entries

 The above examples have used timeouts on inter-task communication; it is also possible, within Ada, to do timed (and conditional) entry call on protected objects

```
select
  P.E; -- E is an entry in protected object P
or
  delay 0.5;
end select;
```

#### Timeouts on Actions

```
delay 0.1;
then abort
   -- action
end select;
```

- If the action takes too long, the triggering event will be taken and the action will be aborted
- This is clearly an effective way of catching runaway code --- Watchdag

#### SUMMARY: Language support for RT Programming

- · Concurrency: multi-tasking
- Communication & synchronization
- Consistency in data sharing /protected data types
- Real time facilities
  - Access to system clock/time
  - Delay constructs: Delay(10) and Delay until next-time
  - Timely execution of tasks (run-time system)

## The "core" of RT Programming Languages

- · Primitive Types
  - Basic Types: e.g. Integers, reals, lists, ...
  - Abstract data type: Semaphore
- Assignment: X:= E
- Control Statements: If, While, ..., goto
- Sequential composition: P;P
- Concurrent composition: P|| P
- Communication: ale, a?x
- Choice: P or P
- Clock reading: Time
- Delays: Delay(n), Delay until n
- . Exception: Loop P until B

## RT Programming Languages

- "Classic" high-level languages with RT extensions e.g.

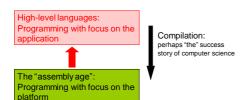
  - Real-Time Java, C + RTOS
  - SDL, Soft RT language for telecommunication systems
- Synchronous Programming (from 1980's)
  - Esterel (Gerard Berry)Lustre (Caspi and Halbwachs)
  - Signal (le Guernic and Benveniste)
- Design, Modeling, Validation, and Code Generation (from models to code)
  - Giotto (Henzinger et al, not quite synchrnous)UPPAAL/TIMES (Uppsala)

  - Real-Time UML
  - SimuLink

#### **RT Programming Languages**

- · "Classic" high-level languages with RT extensions e.g.
  - Ada, Real-Time Java, C + RTOS
  - SDL, Soft RT language for telecommunication systems
- · Synchronous Programming
  - Esterel (Gerard Berry)
  - Lustre ( Caspi and Halbwachs)
  - Signal (le Guernic and Benveniste)
- Towards Real-Time Programming (mostly in research):
  - Giotto (Henzinger et al, not quite synchrnous)
  - TIMES (Uppsala)

The History of Computer Science: Lifting the Level of Abstraction



The History of Computer Science: Lifting the Level of Abstraction

