IMO Shortlist Writeups

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1 Introduction

Here's a compiled list of my typed up solutions to various IMO shortlist problems during my preparation for the 66th IMO.

2 Problems

2.1 ISL 2022

Problem 2.1–2022 A1 Let $(a_n)_{n\geq 1}$ be a sequence of positive real numbers with the property that

$$(a_{n+1})^2 + a_n a_{n+2} \le a_n + a_{n+2}$$

for all positive integers n. Show that $a_{2022} \leq 1$.

Solution. Define a sequence $b_i = a_i - 1 \ \forall i$. The given condition is equivalent to

$$b_n b_{n+2} + b_{n+1} (b_{n+1} + 2) \le 0$$

Where $b_i > -1 \ \forall i$. Now FTSOC $b_{2022} > 0$. Notice we also have

$$b_{n-1}b_{n+1} + b_n(b_n + 2) \le 0$$

Summing the two gives

$$b_n(b_{n-1} + b_n + 2) + b_{n+1}(b_{n+1} + b_{n-1} + 2) \le 0$$

Substituting n=2022 and n=2021 into the above equation, we get that $b_{2021}<0$ and also $b_{2023}<0$ respectively. As a result, considering n=2021 in the first equation gives us a contradiction, and we're done.

2.2 ISL 2020

Problem 2.2–2020 A3 Suppose that a, b, c, d are positive real numbers satisfying (a +

c)(b+d) = ac + bd. Find the smallest possible value of

$$\frac{a}{b} + \frac{b}{c} + \frac{c}{d} + \frac{d}{a}.$$

Proof.

$$\frac{a}{b} + \frac{d}{c} = \frac{ac + bd}{bc}$$

$$\frac{b}{c} + \frac{a}{d} = \frac{ac + bd}{cd}$$

$$\frac{c}{d} + \frac{b}{a} = \frac{ac + bd}{ad}$$

$$\frac{d}{a} + \frac{c}{b} = \frac{ac + bd}{ab}$$

Summing the 4 expressions gives:

$$\begin{split} \left(\frac{a}{b} + \frac{b}{c} + \frac{c}{d} + \frac{d}{a}\right) + \left(\frac{b}{a} + \frac{c}{b} + \frac{d}{c} + \frac{a}{d}\right) &= (ac + bd) \left(\frac{1}{bc} + \frac{1}{cd} + \frac{1}{ad} + \frac{1}{ab}\right) \\ &= (a + c)(b + d) \left(\frac{1}{bc} + \frac{1}{cd} + \frac{1}{ad} + \frac{1}{ab}\right) \\ &= (ab + ad + bc + cd) \left(\frac{1}{bc} + \frac{1}{cd} + \frac{1}{ad} + \frac{1}{ab}\right) \\ &> 16 \end{split}$$

Where the last inequality is by Cauchy-Schwarz. Let $x=\frac{a}{b},y=\frac{b}{c},z=\frac{c}{d},w=\frac{d}{a}.$ So we have xyzw=1 and

$$(x+y+z+w) + \left(\frac{1}{x} + \frac{1}{y} + \frac{1}{z} + \frac{1}{w}\right) \ge 16$$

$$\sum_{cuc} \left(x + \frac{1}{x} - 2\right) = \sum_{cuc} \left(\frac{(x-1)^2}{x}\right) = \sum_{cuc} x \left(1 - \frac{1}{x}\right)^2 \ge 8$$

Finally, notice that

$$x+y+z+w \geq \sum_{cw} x \left(1-\frac{1}{x}\right)^2 \geq 8$$

Finally, we can take $(a, b, c, d) = (1, 2 + \sqrt{3}, 1, 2 + \sqrt{3})$ to achieve equality.

Problem 2.3–2020 C2 In a regular 100-gon, 41 vertices are colored black and the remaining 59 vertices are colored white. Prove that there exist 24 convex quadrilaterals Q_1, \ldots, Q_{24} whose corners are vertices of the 100-gon, so that

ullet the quadrilaterals Q_1,\ldots,Q_{24} are pairwise disjoint, and

• every quadrilateral Q_i has three corners of one color and one corner of the other color.

Proof. Here's a really clever trick... We will completely ignore one of the white vertices. Now we have 41 black vertices and 58 white vertices. Notice now that no matter how we choose the quadraliterals, due to parity reasons, the number of whites unchosen is never equal to the number of blacks unchosen.

Lemma: □

2.3 ISL 2019

Problem 2.4–2019 G4 Let P be a point inside triangle ABC. Let AP meet BC at A_1 , let BP meet CA at B_1 , and let CP meet AB at C_1 . Let A_2 be the point such that A_1 is the midpoint of PA_2 , let B_2 be the point such that B_1 is the midpoint of PB_2 , and let C_2 be the point such that C_1 is the midpoint of PC_2 . Prove that points A_2, B_2 , and C_2 cannot all lie strictly inside the circumcircle of triangle ABC.

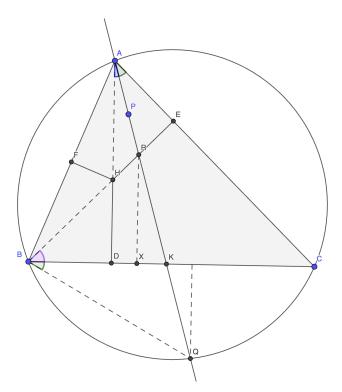


Fig 1: Diagram for 2019 G4

Proof. We will first prove the case where ABC is acute, and then deal with the obtuse case.

Let H denote the orthocenter of the triangle ABC. Drop the perpendiculars from H to each of the sides BC, AC, AB to be D, E, F respectively. Now, WLOG point P lies inside the quadrilateral AEFH. We will prove that the reflection of point P across K will be outside of (ABC) (notice I have renamed a few points for convenience). In fact, we just need to prove KR > KQ for this to work.

Now, WLOG FH < EH. P can lie on either the same side of AH as F, or the opposite. If P lies on the same side of AH as F, we have KR > HD > KQ so we're done. Henceforth we assume P lies ont he same side of AH as E.

In the above diagram, let

$$\alpha = \angle QBC = \angle QAC \text{ (green)}$$

$$\beta = \angle CBE = \angle DAC \text{ (purple)}$$

$$\gamma = \angle DAK \text{ (blue)}$$

We start with the following manipulation:

$$\begin{split} \sin(2\beta)\cos(2\gamma) - \cos(2\beta)\sin(2\gamma) &< \sin(2\beta) \\ \sin(2(\beta - \gamma)) &< \sin(2\beta) \\ \sin(2\alpha) &< \sin(2\beta) \\ \cos(\alpha)\sin(\alpha) &< \cos(\beta)\sin(\beta) \\ \frac{\cos(\alpha)}{\cos(\beta)} &< \frac{\sin(\beta)}{\sin(\alpha)} \end{split}$$

Now, we also have by sine rule in $\triangle ARB$ and $\triangle AQB$:

$$\begin{split} \frac{BR}{\sin(\angle BAR)} &= \frac{AB}{\sin(\angle ARB)} = \frac{AB}{\sin(\angle ARE)} = \frac{AB}{\sin(90 - \alpha)} = \frac{AB}{\cos(\alpha)} \\ \frac{BQ}{\sin(\angle BAR)} &= \frac{AB}{\sin(\angle C)} = \frac{AB}{\sin(90 - \beta)} = \frac{AB}{\cos(\beta)} \\ \frac{BQ}{BR} &= \frac{\cos(\alpha)}{\cos(\beta)} < \frac{\sin(\beta)}{\sin(\alpha)} \end{split}$$

Hence

and we have

$$BQ\sin(\alpha) < BR\sin(\beta) \iff YQ < RX$$

But this is enough to deduce KQ < KR as $\triangle RKX \sim \triangle QKY$. So we're done in the ABC acute case. Now we handle obtuse:

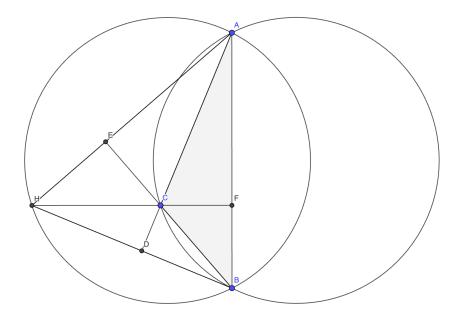


Fig 2: Obtuse case of 2019 G4.

We still construct the orthocenter H. Notice that P must be within the shaded region, and by our work on acute triangles before, the reflection of P across BD or AC will always lie outside of (AHB), as $\triangle AHB$ is acute. Finally, notice that this is enough to finish the problem, as (AHB) completely encloses the arc ACB on the circumcircle of $\triangle ABC$.

Remark: As I was attempting this problem, I told myself that barycentric coordinates would be an easy way to solve this, but I didn't know how they worked :(It turns out, this problem is much easier with a bary bash.

Problem 2.5–2019 N3 We say that a set S of integers is *rootiful* if, for any positive integer n and any $a_0, a_1, \dots, a_n \in S$, all integer roots of the polynomial $a_0 + a_1x + \dots + a_nx^n$ are also in S. Find all rootiful sets of integers that contain all numbers of the form $2^a - 2^b$ for positive integers a and b.

Proof. I claim that the answer is $S \in \mathbb{Z}$. Clearly this works. Now we prove that if we start with the set $S = \{2^a - 2^b \mid a, b \in \mathbb{Z}^+\}$, we can then get every integer with the appropriate choices of coefficients.

First we see 1 is in S by taking P(x) = 2x - 2.

Now, notice that if k is in S, -k must also be in S by taking P(x) = x + k. Therefore, we restrict our search to only positive integers, as the negative integers will follow.

We will take the minimal positive integer m such that $m \notin S$ and aim to find a contradiction. Let $m = 2^e p$, where p is odd. Consider the number $M = 2^{e+\varphi(p)+1} - 2^{e+1} = 2^{e+1}(2^{\varphi(p)} - 1)$. Clearly $M \in S$ and $m \mid M$ (by Euler's Theorem).

We write $M = b_1 m + b_2 m^2 + b_3 m^3 + \cdots + b_n m^n$. Then we can consider the polynomial

$$P(x) = -M + b_1 x + b_2 x^2 + b_3 x^3 + \dots + b_n x^n$$

This works as $b_i \in S \ \forall i$ since $b_i < m$. Finally, notice that m must be a root to the above polynomial, and so $m \in S$, and we're done.

2.4 ISL 2017

Problem 2.6–2017 A4 A sequence of real numbers a_1, a_2, \ldots satisfies the relation

$$a_n = -\max_{i+j=n} (a_i + a_j)$$
 for all $n > 2017$.

Prove that the sequence is bounded, i.e., there is a constant M such that $|a_n| \leq M$ for all positive integers n.

Proof. Let's denote a_x to be the element with the maximum absolute value in the set $\{a_1, a_2, \dots a_{2017}\}$. We split the problem into cases:

Case 1: $a_x = 0$. This case is trivial as all values in the sequence is equal to 0.

Case 2: $a_x > 0$. Let $M = a_x$, I will prove that $-2M \le a_i \le M \ \forall i$. Proof: Notice that

$$\max_{i+i=2018} (a_i + a_j) \ge a_x + a_{2018-x} \ge 0$$

So $a_{2018} = -\max_{i+j=2018}(a_i + a_j) \le 0$, i.e. It's bounded above by 0. We also know that

$$\max_{i+i=2018} (a_i + a_j) \le M + M = 2M$$

so $a_{2018} = -\max_{i+j=2018}(a_i + a_j) \ge -2M$. So we have

$$-2M < a_{2018} < 0$$

Now, if $-M \le a_{2018} \le 0$, we can carry on this process iteratively to get that the next element also has the bound stated above. Otherwise, assume that $-2M \le a_{2018} < -M$. We see that

$$\max_{i+j=2019} (a_i + a_j) \ge a_x + a_{2019-x} \ge M + (-2M) = -M$$

So that means $a_{2019} = -\max_{i+j=2019} (a_i + a_j) \le M$. But also,

$$\max_{i+j=2019} (a_i + a_j) \le M + M = 2M$$

So we have

$$-2M \le a_{2019} \le M$$

Thus we may continue this process iteratively to get that $-2M \le a_i \le M \ \forall i$.

Case 3: $a_x < 0$. Let $-M = a_x$, M > 0.

Notice that

$$\max_{i+j=2018} (a_i + a_j) \le 2M$$

$$\max_{i+j=2018} (a_i + a_j) \ge -2M$$

So we achieve the bound that $-2M \le a_{2018} \le 2M$.

Case 3.1: If $M < a_{2018} \le 2M$, we can refer to Case 2 above to see that the sequence is bounded.

Case 3.2: If $-M \le a_{2018} \le M$, We can iterate this process again, as $a_x = -M$ is still the a_i with the largest absolute value.

Case 3.3: $-2M \le a_{2018} < -M$.

Let $a_{2018} = -k$. Therefore there must exist p, q such that p + q = 2018 and $a_p + a_q = k$. WLOG let $a_p \ge \frac{k}{2}$.

We see

$$\max_{i+j=2019} (a_i + a_j) \ge a_p + a_{2019-p} \ge \frac{k}{2} + (-k) = \frac{-k}{2} \ge \frac{-2M}{2} = -M$$

and also

$$\max_{i+j=2019} (a_i + a_j) \le M + M = 2M$$

So we actually see that

$$-2M \le a_{2019} \le M$$

But we're done here, by considering the most negative element $a_n = -N$. There must be an a_i with $a_i > \frac{N}{2}$, so the lower bound for $\max_{i+j=n}(a_i+a_j)$ is $\frac{N}{2}+(-N)=\frac{-N}{2}\geq -M$. The upper bound of 2M is obvious to see.

So when when calculate the next values of the sequence, the upper and lower bounds for $\max_{i+j=n}(a_i+a_j)$ are fixed at 2M and -M respectively, so we're done.