# IMO Shortlist Writeups

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## 1 Introduction

Here's a compiled list of my typed up solutions to various IMO shortlist problems during my preparation for the 66th IMO.

## 2 Problems

### 2.1 ISL 2022

**Problem 2.1–2022 A1** Let  $(a_n)_{n\geq 1}$  be a sequence of positive real numbers with the property that

$$(a_{n+1})^2 + a_n a_{n+2} \le a_n + a_{n+2}$$

for all positive integers n. Show that  $a_{2022} \leq 1$ .

Solution. Define a sequence  $b_i = a_i - 1 \ \forall i$ . The given condition is equivalent to

$$b_n b_{n+2} + b_{n+1} (b_{n+1} + 2) \le 0$$

Where  $b_i > -1 \ \forall i$ . Now FTSOC  $b_{2022} > 0$ . Notice we also have

$$b_{n-1}b_{n+1} + b_n(b_n+2) \le 0$$

Summing the two gives

$$b_n(b_{n-1} + b_n + 2) + b_{n+1}(b_{n+1} + b_{n-1} + 2) \le 0$$

Substituting n=2022 and n=2021 into the above equation, we get that  $b_{2021}<0$  and also  $b_{2023}<0$  respectively. As a result, considering n=2021 in the first equation gives us a contradiction, and we're done.

## 2.2 ISL 2019

**Problem 2.2–2019 G4** Let P be a point inside triangle ABC. Let AP meet BC at  $A_1$ , let BP meet CA at  $B_1$ , and let CP meet AB at  $C_1$ . Let  $A_2$  be the point such that  $A_1$  is the midpoint of  $PA_2$ , let  $B_2$  be the point such that  $B_1$  is the midpoint of  $PB_2$ , and let  $C_2$ 

be the point such that  $C_1$  is the midpoint of  $PC_2$ . Prove that points  $A_2, B_2$ , and  $C_2$  cannot all lie strictly inside the circumcircle of triangle ABC.

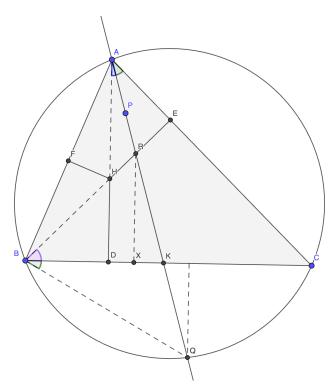


Fig 1: Diagram for 2019 G4

*Proof.* We will first prove the case where ABC is acute, and then deal with the obtuse case.  $\Box$ 

**Remark:** As I was attempting this problem, I told myself that barycentric coordinates would be an easy way to solve this, but I didn't know how they worked :( It turns out, this problem is much easier with a bary bash.

**Problem 2.3–2019 N3** We say that a set S of integers is *rootiful* if, for any positive integer n and any  $a_0, a_1, \dots, a_n \in S$ , all integer roots of the polynomial  $a_0 + a_1x + \dots + a_nx^n$  are also in S. Find all rootiful sets of integers that contain all numbers of the form  $2^a - 2^b$  for positive integers a and b.

*Proof.* I claim that the answer is  $S \in \mathbb{Z}$ . Clearly this works. Now we prove that if we start with the set  $S = \{2^a - 2^b \mid a, b \in \mathbb{Z}^+\}$ , we can then get every integer with the appropriate choices of coefficients.

First we see 1 is in S by taking P(x) = 2x - 2.

Now, notice that if k is in S, -k must also be in S by taking P(x) = x + k. Therefore, we restrict our search to only positive integers, as the negative integers will follow.

We will take the minimal positive integer m such that  $m \notin S$  and aim to find a contradiction. Let  $m = 2^e p$ , where p is odd. Consider the number  $M = 2^{e+\varphi(p)+1} - 2^{e+1} = 2^{e+1}(2^{\varphi(p)} - 1)$ . Clearly  $M \in S$  and  $m \mid M$  (by Euler's Theorem).

We write  $M = b_1 m + b_2 m^2 + b_3 m^3 + \cdots + b_n m^n$ . Then we can consider the polynomial

$$P(x) = -M + b_1 x + b_2 x^2 + b_3 x^3 + \dots + b_n x^n$$

This works as  $b_i \in S \ \forall i$  since  $b_i < m$ . Finally, notice that m must be a root to the above polynomial, and so  $m \in S$ , and we're done.

#### 2.3 ISL 2017

**Problem 2.4–2017 A4** A sequence of real numbers  $a_1, a_2, \ldots$  satisfies the relation

$$a_n = -\max_{i+j=n} (a_i + a_j)$$
 for all  $n > 2017$ .

Prove that the sequence is bounded, i.e., there is a constant M such that  $|a_n| \leq M$  for all positive integers n.

*Proof.* Let's denote  $a_x$  to be the element with the maximum absolute value in the set  $\{a_1, a_2, \dots a_{2017}\}$ . We split the problem into cases:

Case 1:  $a_x = 0$ . This case is trivial as all values in the sequence is equal to 0.

Case 2:  $a_x > 0$ . Let  $M = a_x$ , I will prove that  $-2M \le a_i \le M \ \forall i$ . Proof: Notice that

$$\max_{i \neq i-2018} (a_i + a_j) \ge a_x + a_{2018-x} \ge 0$$

So  $a_{2018} = -\max_{i+j=2018}(a_i + a_j) \le 0$ , i.e. It's bounded above by 0. We also know that

$$\max_{i+j=2018} (a_i + a_j) \le M + M = 2M$$

so  $a_{2018} = -\max_{i+j=2018}(a_i + a_j) \ge -2M$ . So we have

$$-2M \le a_{2018} \le 0$$

Now, if  $-M \le a_{2018} \le 0$ , we can carry on this process iteratively to get that the next element also has the bound stated above. Otherwise, assume that  $-2M \le a_{2018} < -M$ . We see that

$$\max_{i+j=2019} (a_i + a_j) \ge a_x + a_{2019-x} \ge M + (-2M) = -M$$

So that means  $a_{2019} = -\max_{i+j=2019} (a_i + a_j) \le M$ . But also,

$$\max_{i+j=2019} (a_i + a_j) \le M + M = 2M$$

So we have

$$-2M \le a_{2019} \le M$$

Thus we may continue this process iteratively to get that  $-2M \le a_i \le M \ \forall i$ .

Case 3:  $a_x < 0$ . Let  $-M = a_x$ , M > 0.

Notice that

$$\max_{i+j=2018} (a_i + a_j) \le 2M$$

$$\max_{i+j=2018} (a_i + a_j) \ge -2M$$

So we achieve the bound that  $-2M \le a_{2018} \le 2M$ .

Case 3.1: If  $M < a_{2018} \le 2M$ , we can refer to Case 2 above to see that the sequence is bounded.

Case 3.2: If  $-M \le a_{2018} \le M$ , We can iterate this process again, as  $a_x = -M$  is still the  $a_i$  with the largest absolute value.

Case 3.3:  $-2M \le a_{2018} \le -M$ .

Let  $a_{2018} = -k$ . Therefore there must exist p, q such that p + q = 2018 and  $a_p + a_q = k$ . WLOG let  $a_p \ge \frac{k}{2}$ .

We see

$$\max_{i+j=2019} (a_i + a_j) \ge a_p + a_{2019-p} \ge \frac{k}{2} + (-k) = \frac{-k}{2} \ge \frac{-2M}{2} = -M$$

and also

$$\max_{i+j=2019} (a_i + a_j) \le M + M = 2M$$

So we actually see that

$$-2M \le a_{2019} \le M$$

But we're done here, by considering the most negative element  $a_n = -N$ . There must be an  $a_i$  with  $a_i > \frac{N}{2}$ , so the lower bound for  $\max_{i+j=n}(a_i+a_j)$  is  $\frac{N}{2}+(-N)=\frac{-N}{2}\geq -M$ . The upper bound of 2M is obvious to see.

So when when calculate the next values of the sequence, the upper and lower bounds for  $\max_{i+j=n}(a_i+a_j)$  are fixed at 2M and -M respectively, so we're done.