Program Structures and Algorithms Spring 2023(SEC –03)

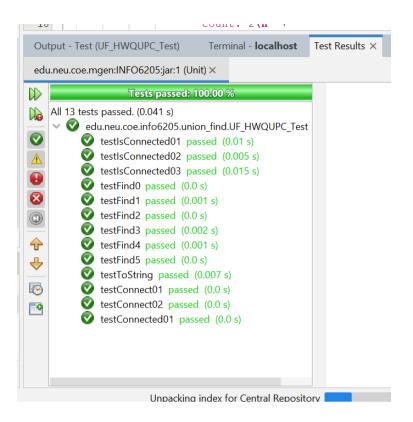
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Task:

- Implement height-weighted Quick Union with Path Compression
- Develop a UF ("union-find") client that takes an integer value n from the command line to determine the number of "sites."
- Determine the relationship between the number of objects (n) and the number of pairs (m) generated

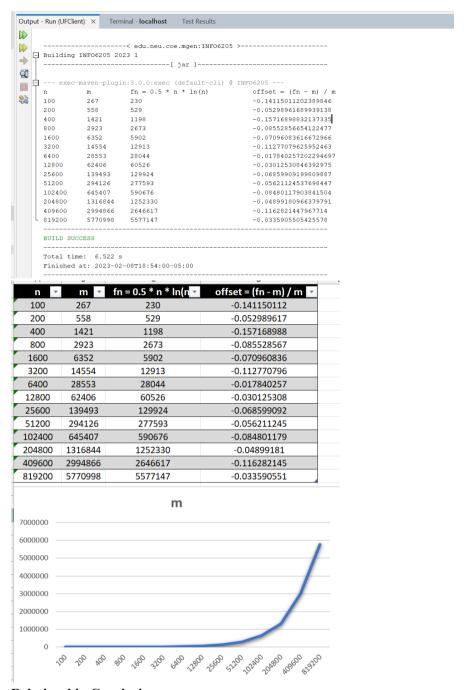
Step 1: UF HWQUPC implementation snippets

```
/** Returns the component identifier for the component containing site ...9 lines */
   public int find(int p) {
       validate(p);
       int root = p;
       while (parent[root] != root) {
            doPathCompression(i:root);
           root = parent[root];
       // END
       return root;
179
L80 📮
         private void mergeComponents(int i, int j) {
181
    // FIXME make shorter root point to taller one
L82
              if (height[j] > height[i]) {
183
L84
                parent[i] = j; height[j] += height[i]; height[i] = 0;
185
                 parent[j] = i; height[i] += height[j]; height[j] = 0;
186
L87
188
              // END
L89
   F
L90
           * This implements the single-pass path-halving mechanism of path
191
           * compression
192
L93
   口
         private void doPathCompression(int i) {
L94
L95
             // FIXME update parent to value of grandparent
196
L97
              if (this.pathCompression) {
L98
                 parent[i] = parent[parent[i]];
199
200
201
              // END
202
203
204
```



Step 2: UF Client implementation

```
UF.java × ☑ TypedUF.java × ☑ TypedUF_HWQUPC.java × ☑ HWQUPC_Solution.java × ☑ UFClient.java ×
Source History | [4] 🕟 • 🔊 • | 🔾 😎 🗗 🗔 | 🔗 👆 🖒 | 💇 💇 | 💿 🖂 | 💯 📑
     * @author altaf
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        public class UFClient {
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             public static int count(int n) {
    UF_HWQUPC uf = new UF_HWQUPC(n);
                   Random random = new Random():
                    while (uf.components() > 1) {
22
                       int i = random.nextInt(bound:n);
23
24
                        int j = random.nextInt(bound:n);
                        if (!uf.connected(p:i, q:j)) {
    uf.union(p:i, q:j);
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                        num++;
                   return num;
34
35
             public static void main(String[] args) {
    System.out.printf(format: "%-10s %-30s %-20s\n", args: "n", args: "n", args: "m", args: "fn = 0.5 * n * ln(n)", args: "offset = (fn - m) / m");
     早
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                   for (int i = 100; i < 1000000; i *= 2) {
                        int sum = 0;
for (int j = 0; j < 10; j++) {
    sum += count(n:i);</pre>
43
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                        double m = sum / 10.0;
int fn = (int) (0.5 * i * Math.log(a:i));
                        System.out.printf(format: "%-10s %-10s %-30s %-20s\n", args: i, (int) m, args: fn, (fn - m) / m);
```



Relationship Conclusion:

The correlation between "n" and "m" can be estimated using the equation "fn = 0.5 * n * ln(n)". "m" represents the number of connections needed to connect all the elements in a union find data structure, while "n" represents the number of elements in the data structure.

The relationship between "n" and "m" can be viewed as <u>logarithmic</u>, meaning that as the number of elements "n" increases, the number of connections "m" required also increases, but at a slower pace. The offset value shows how close the value of "m" is to the calculated value "fn", with a positive offset indicating that "m" is greater than "fn" and a negative offset indicating that "m" is less than "fn".