


Colocolo



Optimizations on Ocelot - a Relational Logic Solver in Rosette

Ocelot

- An embedding of relational logic in Rosette
 - Used in **MemSynth**, a tool for reasoning about memory consistency.
 - Ocelot works fast and well in its domain, but does not implement certain general optimizations, and lacks benchmarks.
 - Skolemization
 - Optimized CNF-SAT translation
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SKOLEMIZATION

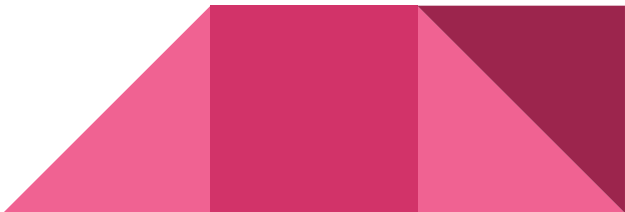
$$\forall x \exists y. P(x,y) \Leftrightarrow \exists f \forall x. P(x,f(x))$$

Skolemization in Relational Logic

- In the case of a bounded universe, we translate $\forall x \exists y . P(x, y)$ to

$$\bigwedge_{i=1}^n \bigvee_{j=1}^n P(A_i, A_j)$$

- After skolemizing,

$$\bigwedge_{i=1}^n P(A_i, f(A_i))$$


Skolemization in Relational Logic

- Skolemize existentially quantified declarations.
- Let \mathbf{S} be the set of \forall or $\neg\exists$ quantified variables in scope.
- For each \exists or $\neg\forall$ quantified declaration $v : \text{mult } R$, where mult is a multiplicity (e.g. one, some) and R is a relation of arity k ,
 - Introduce a new relation R_v of with arity equal to $|\mathbf{S}| + k$.
 - Introduce a new expression E_v which will replace any occurrence of v down the tree.
 - $E_v = a_n.(a_{n-1}.\dots(a_1.R_v)\dots)$, where $a_i \in \mathbf{S}$
 - $\text{upper-bound}(R_v) = \text{upper-bound}(a_1) \rightarrow \dots \rightarrow \text{upper-bound}(a_n) \rightarrow \text{upper-bound}(R)$
 - Domain constraint $:= R_v.U. \underbrace{\dots.U}_{k \text{ times}} \subseteq \{ a_1 : \text{mult}_1 R_1, \dots, a_n : \text{mult}_n R_n \mid \top \}$
 - Range constraint: $(E_v \text{ in } R)$ and $(m E_v)$

Translation to CNF

- $F_1 \wedge F_2 \wedge \dots \wedge F_n \rightsquigarrow (F_1 \vee \neg o) \wedge \dots \wedge (F_n \vee \neg o) \wedge (\neg F_1 \vee \dots \vee \neg F_n \vee o).$
- $F_1 \vee F_2 \vee \dots \vee F_n \rightsquigarrow (\neg F_1 \vee o) \wedge \dots \wedge (\neg F_n \vee o) \wedge (F_1 \vee \dots \vee F_n \vee \neg o).$
- o is a fresh auxiliary variable; true iff left hand formula is true.
- Apply recursively on each F_i and substitute new CNF-SAT variable
- Used internally by Kodkod



Meow: KodKod AST to Colocolo

- Written in Java
- Traverses Kodkod AST, which is actually a DAG due to node-sharing
 - DAG structure is preserved by caching nodes
 - Sort the nodes topologically to ensure correct compiled declaration order.
- Used to compile Kodkod benchmarks into Colocolo code



Demo/Reflection

- An engineering problem
 - Focused on performance improvement
 - Ocelot focused specifically on MemSynth
 - Difficulty with higher arity relations, especially when quantifying
 - No Integers
 - No Compact Boolean Circuits
- Performance Improvements
 - Reduction to SAT - not as good as we hoped
 - Skolemization: 2 - 4x improvement
 - Symmetry Breaking
 - Mixed bag



Average Solver Time

