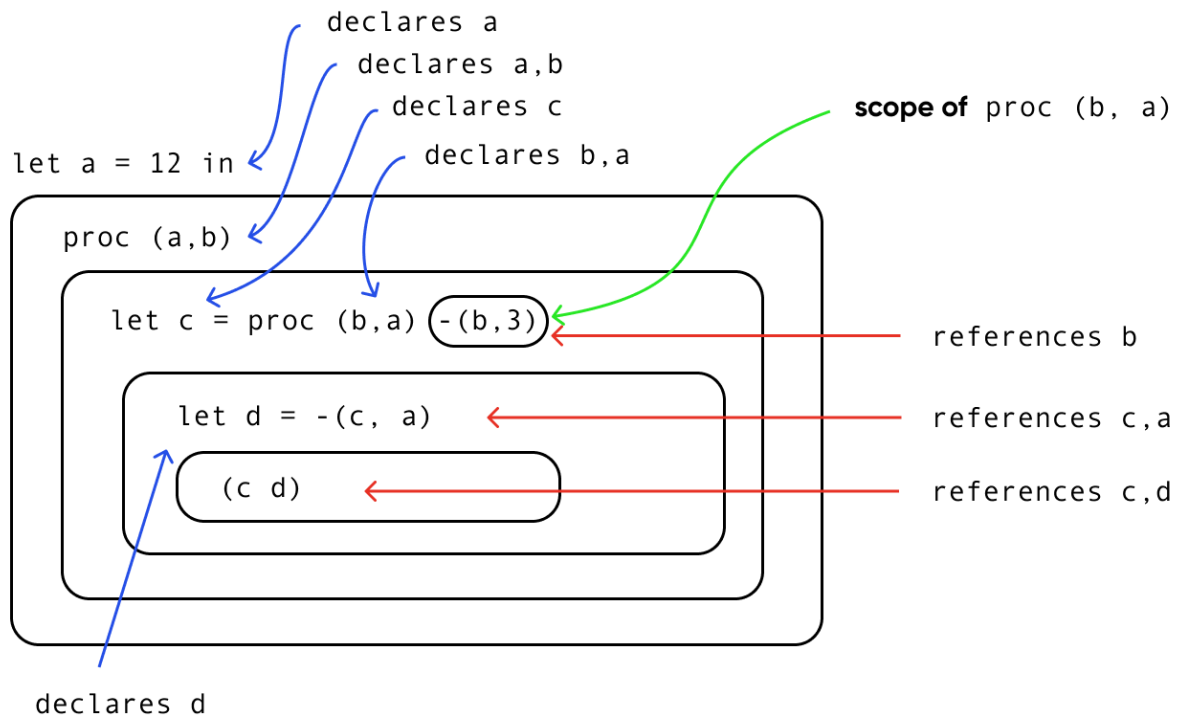


Problem 1A.



Problem 2B.

```
(value-of <<let z = 4 in
      letrec f(x) = if zero?(x) then 1 else (f -(x, 1)) in
      (f z)>> p0)
```

```
= (value-of <<letrec f(x) = if zero?(x) then 1 else (f -(x, 1)) in
      (f z)>>
  (extend-env-rec z (value-of <<4>>) p0))
```

```
Let p1 = (extend-env-rec z (value-of <<4>>) p0)
        = [z=[4]]p0
```

```
= (value-of <<(f z)>>
  (extend-env-rec f << if zero?(x) then 1 else (f -(x, 1))>> p1))
```

```
Let p2 = (extend-env-rec f << if zero?(x) then 1 else (f -(x, 1))>> p1)
```

```
= (apply-procedure (value-of <<f>>) (value-of <<z>>) p2)
```

```
= (apply-procedure (procedure x <<if zero?(x) ...>> p2) [4] p2)
```

```
= (value-of <<if zero?(x) ...>> [x=[4]]p2)
```

```
... = (value-of <<(f -(x, 1))>> [x=[4]]p2)
```

```

... = (apply-procedure (procedure x <<if zero?(x) ...>> [x=[4]]p2)
      [3] [x=[4]]p2)

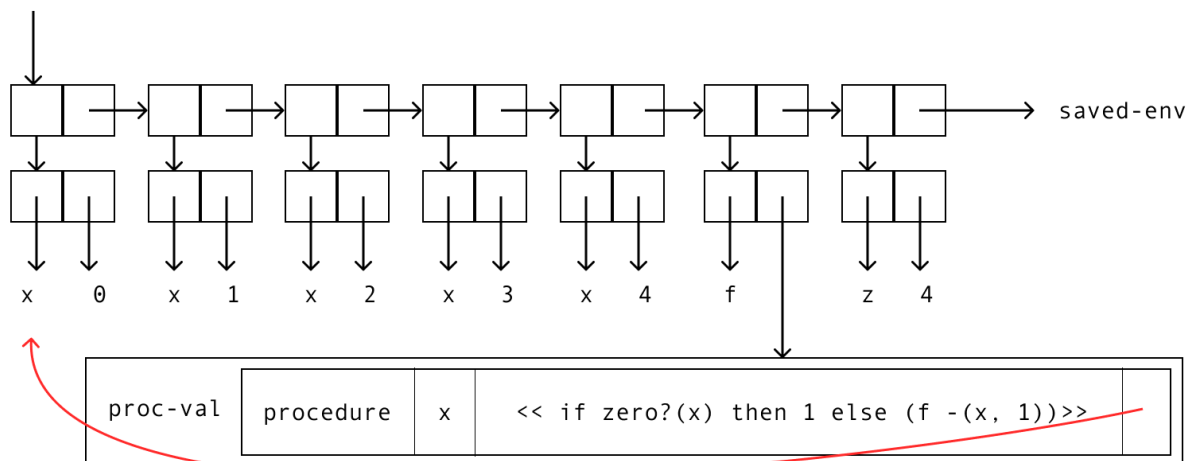
... = (value-of <<(f -(x, 1))>> [x=[3]][x=[4]]p2)

... = (value-of <<(f -(x, 1))>> [x=[2]][x=[3]][x=[4]]p2)

... = (value-of <<(f -(x, 1))>> [x=[1]][x=[2]][x=[3]][x=[4]]p2)

... = (value-of <<1>> [x=[0]][x=[1]][x=[2]][x=[3]][x=[4]]p2)

```



Problem 2A.

```

%let %nameless-var 6 in
  %let %nameless-var 4 in
    %let %nameless-var -(#2,#1) in
      %let %nameless-var -(#2,#1) in
        %nameless-var %lexproc %lexproc %lexproc -(#2,-(#1,#0))
          in (((#0 #1) #3) #2)

```

Problem 2B.

```

let a = 13 in
  let b = 7 in
    let c = 12 in
      let f = proc (x) -(b,x) in
        let d = 5 in
          proc (x) -(x,-(b,c))

```

Problem 3A.

Main function for letrec-m

```
(define apply-extend-env-rec*  
  (lambda (env p-names b-vars bodies saved-env search-var)  
    (let ((procs (member search-var  
                          (map list p-names b-vars bodies)  
                          (lambda (key proc) (eqv? key (car proc))))))  
      (if procs  
          (proc-val (procedure (caddr procs) (caddar procs) env))  
          (apply-env saved-env search-var))))))
```

Problem 3B.

Yes, certainly! But, it depends on the requirements. “letrec” is essentially, “letrec-m” with only one procedure declaration, which means “letrec-m” can easily replace “letrec”. In its current implementation, “letrec-m” extends the environment by declaring all the procedures at once within the same scope, which makes it possible to have mutually recursive procedures.

Depending on the requirement, letrec-m can easily be modified to bind procvals either in parallel, much like scheme’s “let”, or bind sequentially like “let*”.

This is how a parallel binding would look:

```
(let-exp (vars exps body)  
  (let ((vals (map (lambda (exp) (value-of exp env)) exps)))  
    (value-of body (append-env vars vals env))))
```

append-env => binds pre-evaluated values to the environment.

This is how a sequential binding would look:

```
(let*-exp (vars exps body)  
  (value-of body (fold-env vars exps env)))
```

fold-env => iteratively evaluates and binds values to the environment.

Depending on the requirements, a binding mechanism can be implemented.