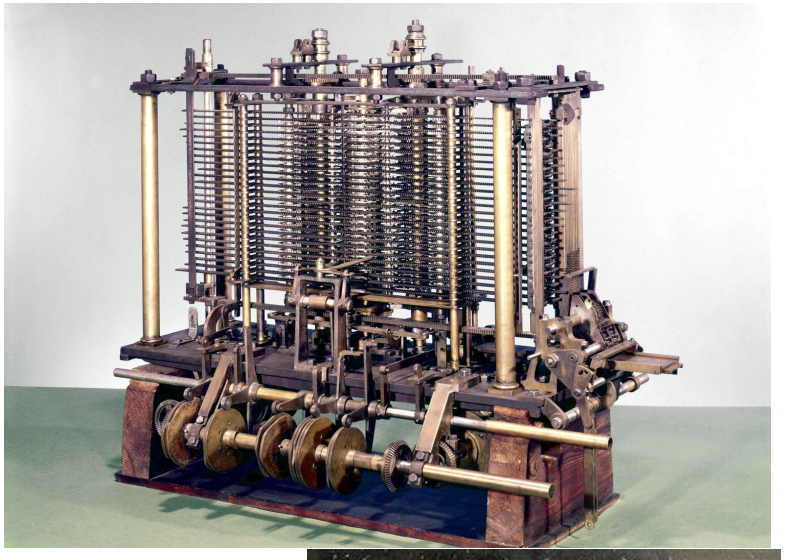
Diagram for the computation by the Engine of the Numbers of Bernoulli. See Note G. (page 722 *et seq.*)

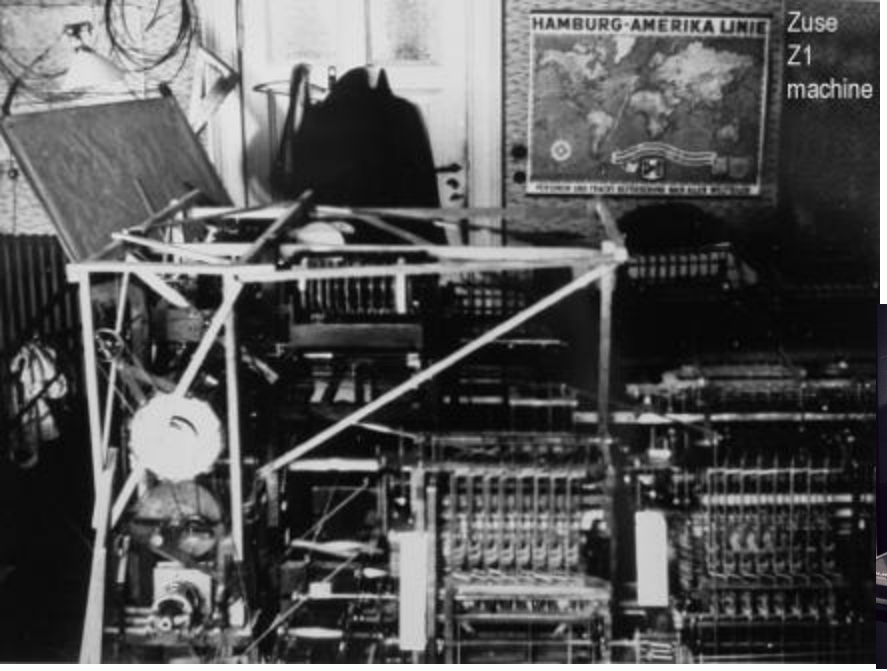
Number of Operation.	Nature of Operation.	Variables acted upon.	Variables receiving results.	Indication of change in the value on any Variable.	Statement of Results.	Data.			Working Variables.										Result Variables.			
						$^1V_1$	$^1V_2$	$^1V_3$	$^0V_4$	$^0V_5$	$^0V_6$	$^0V_7$	$^0V_8$	$^0V_9$	$^0V_{10}$	$^0V_{11}$	$^0V_{12}$	$^0V_{13}$	$^1V_{21}$	$^1V_{22}$	$^1V_{23}$	$^0V_{24}$
						1	2	n											B_1	B_2	B_3	B_7
1	×	$^1V_2 \times ^1V_3$	$^1V_4, ^1V_5, ^1V_6$	$\begin{cases} ^1V_2 = ^1V_2 \\ ^1V_3 = ^1V_3 \end{cases}$	$= 2n$	...	2	n	$2n$	$2n$	$2n$											
2	−	$^1V_4 - ^1V_1$	$^2V_4$	$\begin{cases} ^1V_4 = ^2V_4 \\ ^1V_1 = ^1V_1 \end{cases}$	$= 2n-1$	1	...	...	$2n-1$													
3	+	$^1V_5 + ^1V_1$	$^2V_5$	$\begin{cases} ^1V_5 = ^2V_5 \\ ^1V_1 = ^1V_1 \end{cases}$	$= 2n+1$	1	...	...	...	$2n+1$												
4	+	$^2V_5 \div ^2V_4$	$^1V_{11}$	$\begin{cases} ^2V_5 = ^0V_5 \\ ^2V_4 = ^0V_4 \end{cases}$	$= \frac{2n-1}{2n+1}$	...	...	...	0	0	...	...	...	...	...	$\frac{2n-1}{2n+1}$						
5	+	$^1V_{11} \div ^1V_2$	$^2V_{11}$	$\begin{cases} ^1V_{11} = ^2V_{11} \\ ^1V_2 = ^1V_2 \end{cases}$	$= \frac{1}{2} \cdot \frac{2n-1}{2n+1}$	...	2	...	...	...	...	...	...	...	...	$\frac{1}{2} \cdot \frac{2n-1}{2n+1}$						
6	−	$^2V_{13} - ^2V_{11}$	$^0V_{13}$	$\begin{cases} ^2V_{11} = ^0V_{11} \\ ^0V_{13} = ^1V_{13} \end{cases}$	$= -\frac{1}{2} \cdot \frac{2n-1}{2n+1} = A_0$	...	...	...	...	...	...	...	...	...	...	0	.....	$-\frac{1}{2} \cdot \frac{2n-1}{2n+1} = A_0$				
7	−	$^1V_3 - ^1V_1$	$^1V_{10}$	$\begin{cases} ^1V_3 = ^1V_3 \\ ^1V_1 = ^1V_1 \end{cases}$	$= n-1\ (=3)$	1	...	n	...	...	...	...	...	...	$n-1$							
8	+	$^2V_2 + ^0V_7$	$^1V_7$	$\begin{cases} ^2V_2 = ^1V_2 \\ ^0V_7 = ^1V_7 \end{cases}$	$= 2 + 0 = 2$	...	2	...	...	...	...	2										
9	÷	$^1V_6 \div ^1V_7$	$^2V_{11}$	$\begin{cases} ^1V_6 = ^1V_6 \\ ^0V_{11} = ^2V_{11} \end{cases}$	$= \frac{2n}{2} = A_1$	...	...	...	...	...	$2n$	2	...	...	...	$\frac{2n}{2} = A_1$						
10	×	$^2V_{21} \times ^2V_{11}$	$^1V_{12}$	$\begin{cases} ^2V_{21} = ^1V_{21} \\ ^2V_{11} = ^2V_{11} \end{cases}$	$= B_1 \cdot \frac{2n}{2} = B_1 A_1$	...	...	...	...	...	...	...	...	...	...	$\frac{2n}{2} = A_1$	$B_1 \cdot \frac{2n}{2} = B_1 A_1$	.....	B_1			
11	+	$^1V_{12} + ^1V_{13}$	$^2V_{13}$	$\begin{cases} ^1V_{12} = ^0V_{12} \\ ^1V_{13} = ^2V_{13} \end{cases}$	$= -\frac{1}{2} \cdot \frac{2n-1}{2n+1} + B_1 \cdot \frac{2n}{2}$	...	...	...	...	...	...	...	...	...	...	.....	0	$\left\{-\frac{1}{2} \cdot \frac{2n-1}{2n+1} + B_1 \cdot \frac{2n}{2}\right\}$				
12	−	$^1V_{10} - ^1V_1$	$^2V_{10}$	$\begin{cases} ^1V_{10} = ^2V_{10} \\ ^1V_1 = ^1V_1 \end{cases}$	$= n-2\ (=2)$	1	...	...	...	...	...	...	...	...	$n-2$							
13	−	$^1V_8 - ^1V_1$	$^2V_8$	$\begin{cases} ^1V_8 = ^2V_8 \\ ^1V_1 = ^1V_1 \end{cases}$	$= 2n-1$	1	...	...	...	...	$2n-1$											
14	+	$^1V_1 + ^1V_7$	$^2V_7$	$\begin{cases} ^1V_1 = ^1V_1 \\ ^1V_7 = ^2V_7 \end{cases}$	$= 2 + 1 = 3$	1	...	...	...	...	...	3										
15	÷	$^2V_7 \div ^2V_7$	$^2V_8$	$\begin{cases} ^2V_6 = ^2V_6 \\ ^2V_7 = ^2V_7 \end{cases}$	$= \frac{2n-1}{3}$	...	...	...	...	...	$2n-1$	3	$\frac{2n-1}{3}$									
16	×	$^1V_6 \times ^2V_{11}$	$^4V_{11}$	$\begin{cases} ^1V_6 = ^0V_6 \\ ^2V_{11} = ^4V_{11} \end{cases}$	$= \frac{2n}{2} \cdot \frac{2n-1}{3}$	...	...	...	...	...	...	...	0	...	...	$\frac{2n}{2} \cdot \frac{2n-1}{3}$						
17	−	$^2V_6 - ^1V_1$	$^2V_6$	$\begin{cases} ^2V_6 = ^2V_6 \\ ^1V_1 = ^1V_1 \end{cases}$	$= 2n-2$	1	...	...	...	...	$2n-2$											
18	+	$^1V_1 + ^2V_7$	$^2V_7$	$\begin{cases} ^2V_7 = ^2V_7 \\ ^1V_1 = ^1V_1 \end{cases}$	$= 3 + 1 = 4$	1	...	...	...	...	...	4										
19	÷	$^2V_6 \div ^2V_7$	$^1V_9$	$\begin{cases} ^2V_6 = ^2V_6 \\ ^2V_7 = ^2V_7 \end{cases}$	$= \frac{2n-2}{4}$	...	...	...	...	...	$2n-2$	4	...	$\frac{2n-2}{4}$	...	$\left\{\frac{2n}{2} \cdot \frac{2n-1}{3} \cdot \frac{2n-2}{3} = A_2\right\}$						
20	×	$^1V_9 \times ^4V_{11}$	$^5V_{11}$	$\begin{cases} ^1V_9 = ^0V_9 \\ ^4V_{11} = ^5V_{11} \end{cases}$	$= \frac{2n}{2} \cdot \frac{2n-1}{3} \cdot \frac{2n-2}{4} = A_3$	...	...	...	...	...	...	...	...	0	...	...						
21	×	$^1V_{22} \times ^5V_{11}$	$^6V_{12}$	$\begin{cases} ^1V_{22} = ^1V_{22} \\ ^5V_{12} = ^6V_{12} \end{cases}$	$= B_3 \cdot \frac{2n}{2} \cdot \frac{2n-1}{3} \cdot \frac{2n-2}{3} = B_3 A_3$	...	...	...	...	...	...	...	...	...	...	0	$B_2 A_3$	.....	.....	B_3		
22	+	$^2V_{12} + ^2V_{13}$	$^3V_{13}$	$\begin{cases} ^2V_{12} = ^0V_{12} \\ ^2V_{13} = ^3V_{13} \end{cases}$	$= A_0 + B_1 A_1 + B_2 A_2$	...	...	...	...	...	...	...	...	...	...	0	.....	$\{A_2 + B_1 A_1 + B_2 A_2\}$				
23	−	$^2V_{10} - ^1V_1$	$^2V_{10}$	$\begin{cases} ^2V_{10} = ^2V_{10} \\ ^1V_1 = ^1V_1 \end{cases}$	$= n-3\ (=1)$	1	...	...	...	...	...	...	...	...	$n-3$							

Here follows a repetition of Operations thirteen to twenty-three.

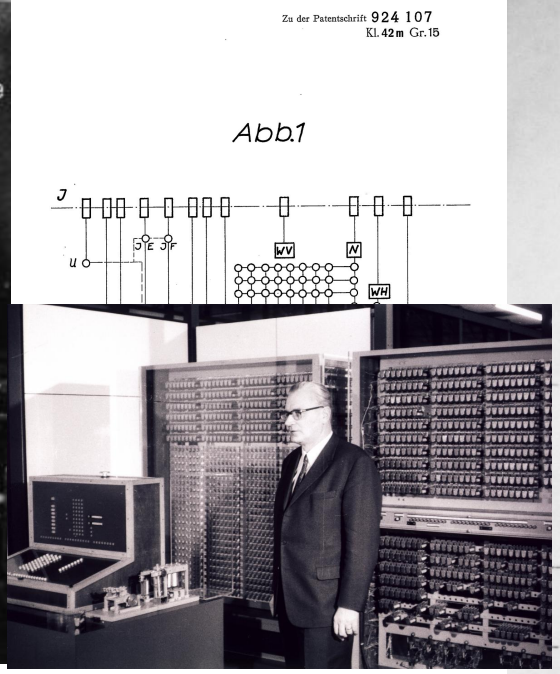
24	+	$^4V_{13} + ^0V_{24}$	$^{24}V_{24}$	$\begin{cases} ^4V_{13} = ^0V_{13} \\ ^0V_{24} = ^{24}V_{24} \end{cases}$	$= B_7$	...	...	...	...	...	...	...	...	...	...	.....	.....	.....	.....	.....	.....	B_7
25	+	$^1V_1 + ^1V_3$	$^1V_3$	$\begin{cases} ^1V_1 = ^1V_1 \\ ^1V_3 = ^1V_3 \\ ^0V_6 = ^0V_6 \\ ^0V_7 = ^0V_7 \end{cases}$	$= n+1 = 4+1 = 5$ ..... by a Variable-card. by a Variable card.	1	...	$n+1$	...	...	0	0										

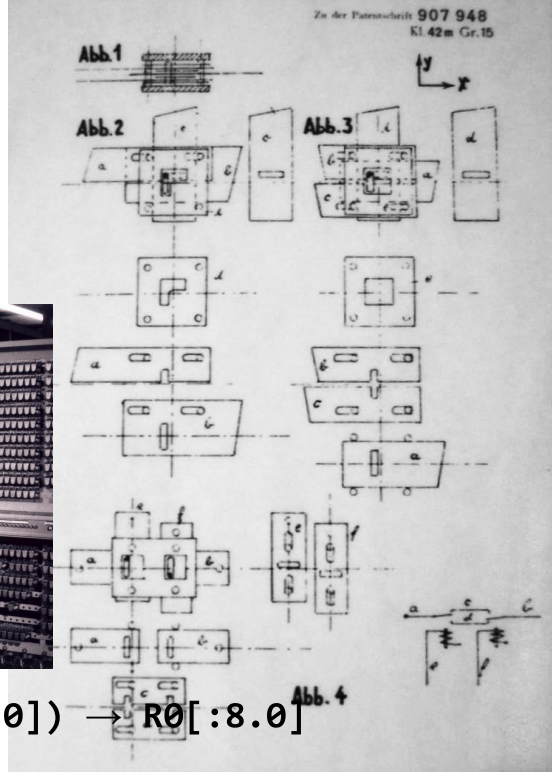


… La nota G estaba dedicada a los números de Bernoulli; en este apartado Ada describe con detalle las operaciones mediante las cuales las tarjetas perforadas "tejerían" una secuencia de números en la máquina analítica. Este código está considerado como el primer algoritmo específicamente diseñado para ser ejecutado por un ordenador, aunque nunca fue probado ya que la máquina nunca llegó a construirse.
(Wikipedia)



Zuse Z1 machine



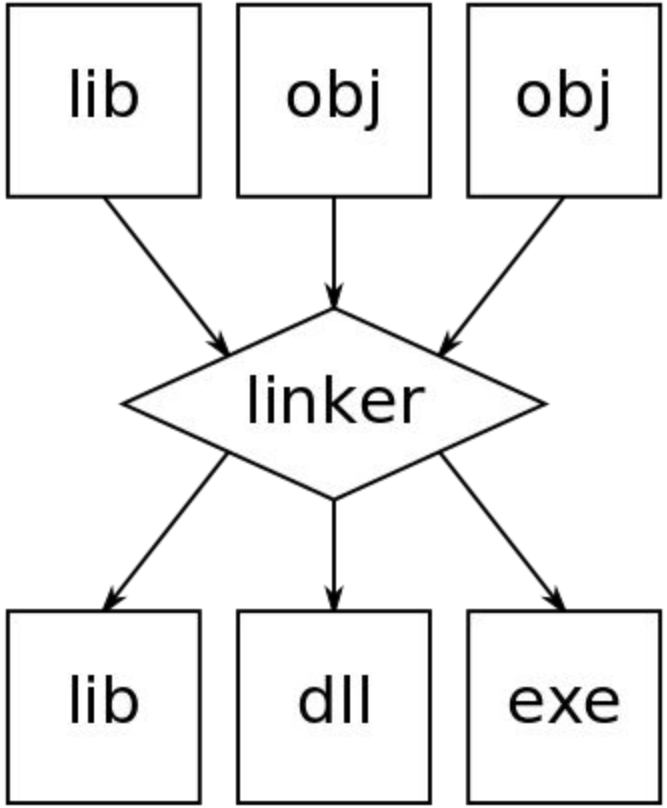


En 1935 Zuse construye la Z1. Leía las instrucciones desde una cinta perforada de 35 mm. No era una máquina Turing completa

```
P2 max (V0[:8.0],V1[:8.0]) → R0[:8.0]
V0[:8.0] → Z1[:8.0]
(Z1[:8.0] < V1[:8.0]) → V1[:8.0] → Z1[:8.0]
Z1[:8.0] → R0[:8.0]
END
```

The Henschel Hs 129B ground attack aircraft

La Z3 (1941) era un computador binario de punto flotante de 22 bits con memoria y unidad de cálculo basada en relés telefónicos. No almacenaba el programa en memoria. A pesar de la ausencia de saltos condicionales, el Z3 era un ordenador Turing-completo



Rear Admiral (then Commodore) Grace M. Hopper, 1984. By 1952, Hopper had finished her program linker. The term compiler was coined by Hopper. It was written for the A-0 System.

The Algol 60 people

GLENN RODNEY ALAN
STUDENT'S NAME

EASTERN WASHINGTON UNIVERSITY

535 68 8500 | SENIOR | WINTER 79 | 12/30/78
STUDENT NO. | CLASS STANDING | QUARTER & YEAR | DATE

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	COURSE SEQUENCE NO.	DEPT. ABBREV.	COURSE TITLE	CRED.	REPEAT (R) WITHDRAWAL (W) FINC (F)	DAY(S) OF THE WEEK	START TIME	END TIME	CLASS ROOM LOCATION
"HAPPINESS IS WINTER QUARTER AT EASTERN WASHINGTON UNIVERSITY"	14 230 01	CS	FORTRAN PROGRAMMIN	3		M,W,F	1200	100	P1103
	14 231 01	CS	COM PROG PROJECTS	2		ARR	ARR	ARR	ARR
	54 212 01	HUM	MUS IN HUMANITIES	5		DAILY	900	1000	MB247
	62 121 01	PHY	DESCRIP ASTRONOMY	5		DAILY	1100	1200	SC151

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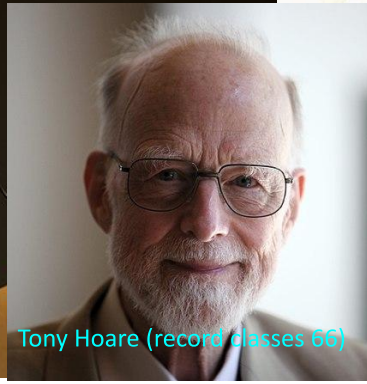
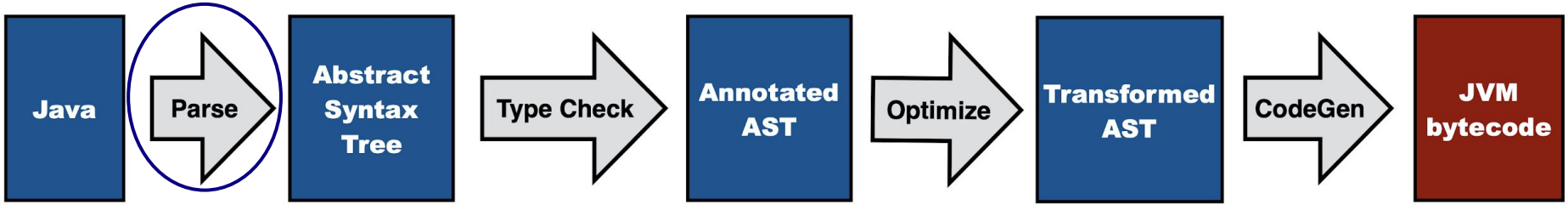
[OO History: Simula and Smallta lk]



The Phases of a Translator



A programming language translator usually consists of a sequence of stages

Lexer:

- Skips the comments and whitespaces and produces the stream of tokens for numbers, identifiers, reserved words, etc

Parser:

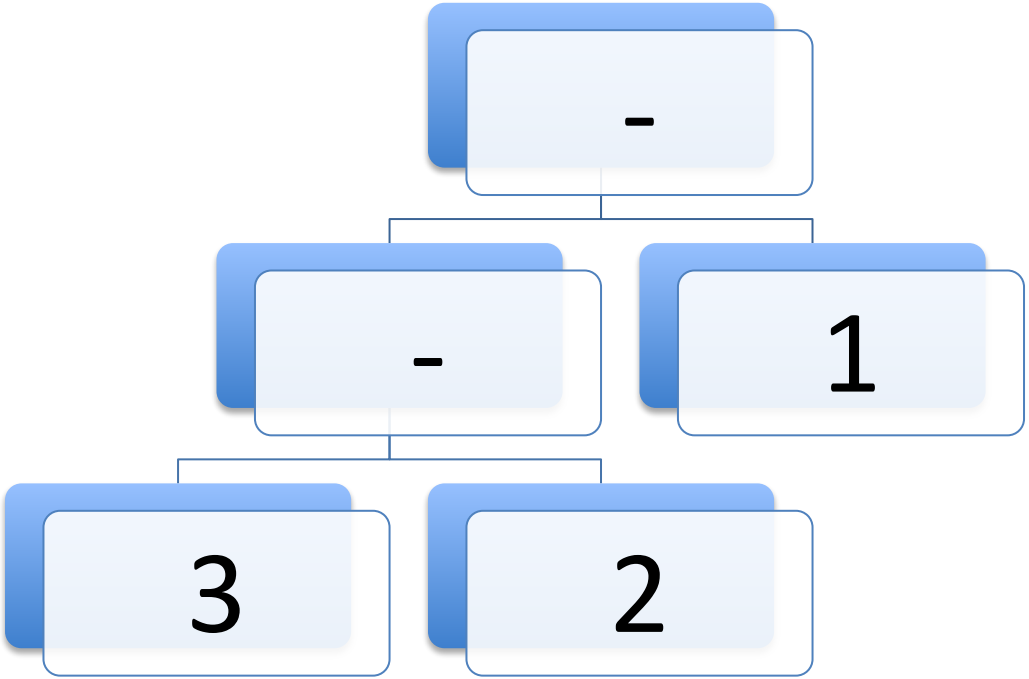
- Reads the stream of tokens, check that it complies with the syntactic rules and produces the *Abstract Syntax Tree*: a data structure representing the underlying syntactic structure of the input program

3 - 2 - 1

Árbol Sintáctico Abstracto

3-2-1

(3-2)-1

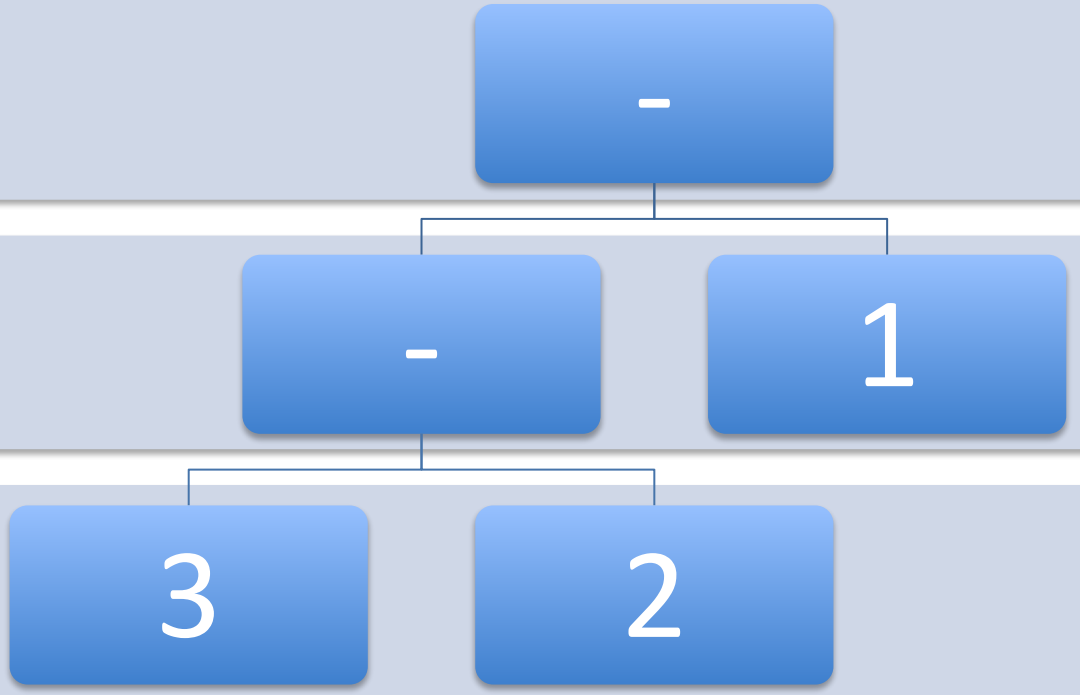


Semántica 3 - 2 - 1

0 = 1 - 1

1 = 3 -2

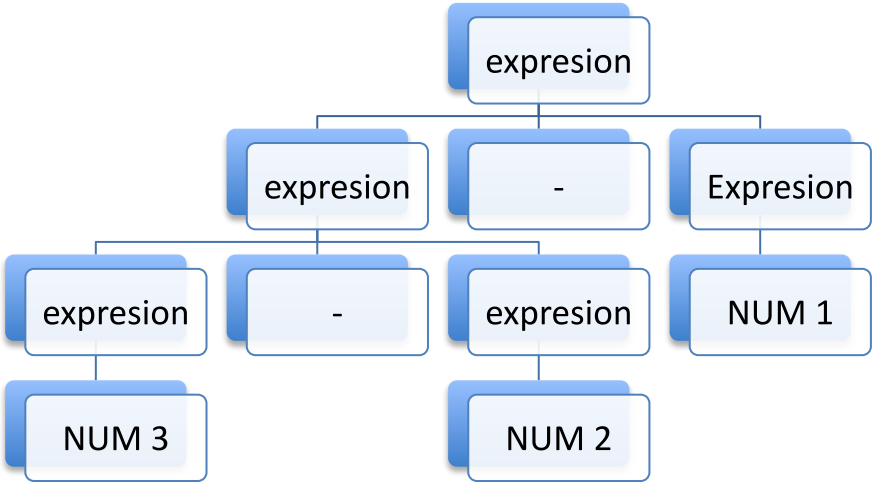
‘3’



Gramática Independiente del Contexto

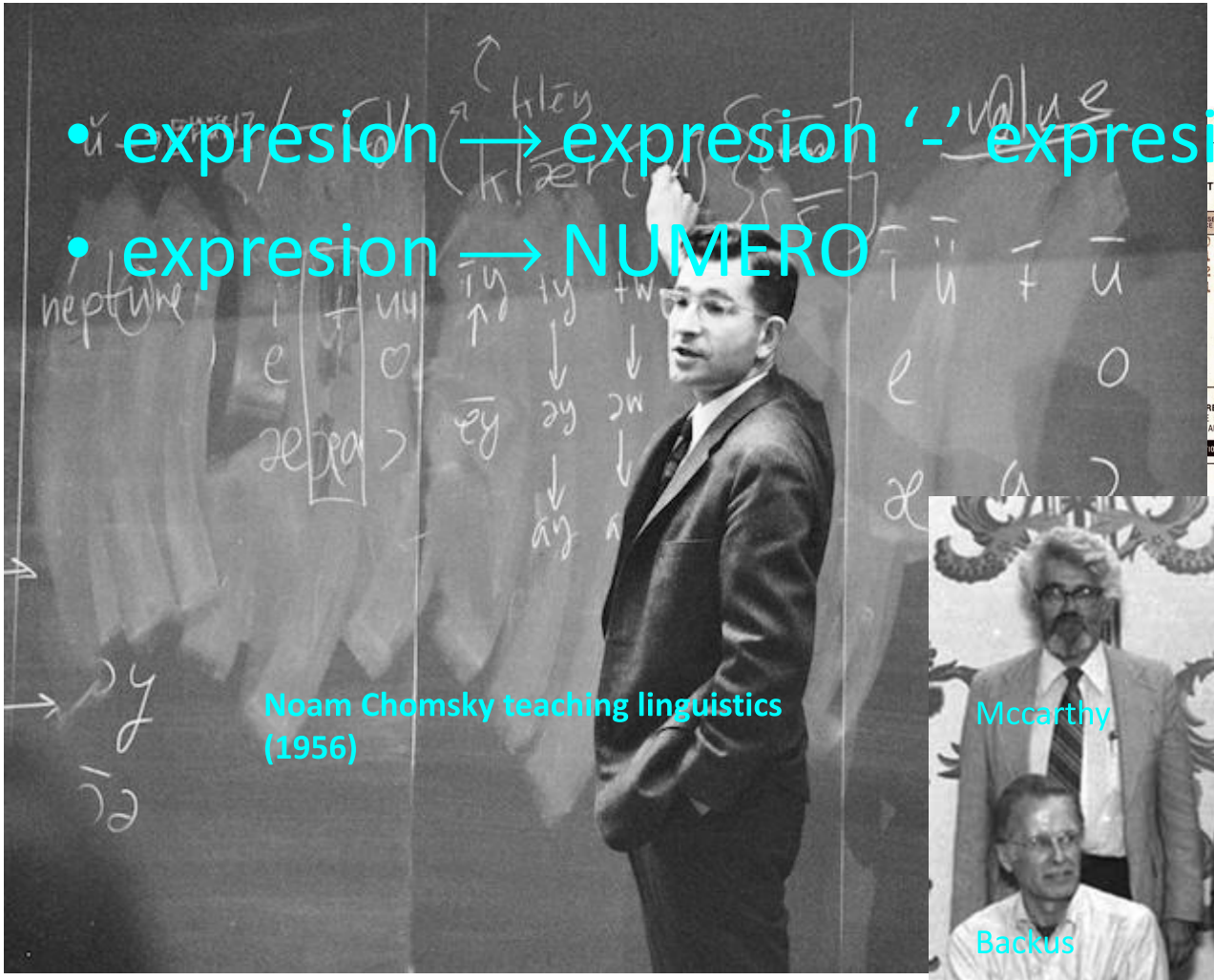
- expresion ⟶ expresion ‘-’ expresion
- expresion ⟶ NUMERO

3-2-1

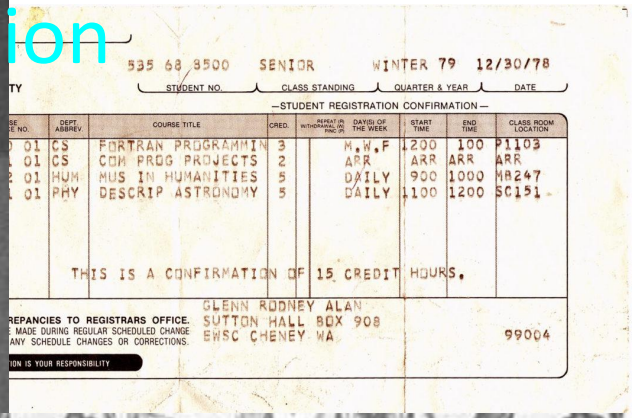


Context Free Grammars

- expresion ⟶ expresion '-' expresion
- expresion ⟶ NUMERO



Noam Chomsky teaching linguistics (1956)



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01	CS	FORTRAN PROGRAMMIN	3		M,W,F	1200	100	P1103
01	CS	COM PROG PROJECTS	2		ARR	ARR	ARR	ARR
01	HUM	MUS IN HUMANITIES	5		DAILY	900	1000	MA247
01	PHY	DESCRIP ASTRONOMY	5		DAILY	1100	1200	SC151

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ION IS YOUR RESPONSIBILITY

Mccarthy

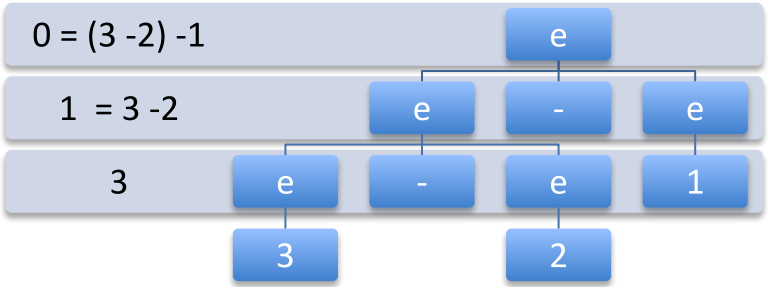
The Algol 60 people

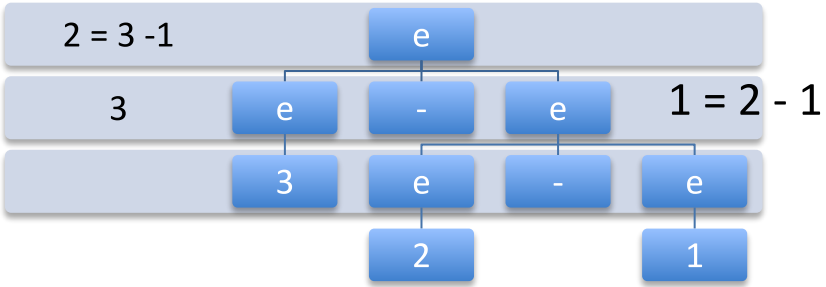
Backus

Naur

Gramática Ambigua

- expresion ⟶ expresion ‘-’ expresion
- expresion ⟶ NUMERO



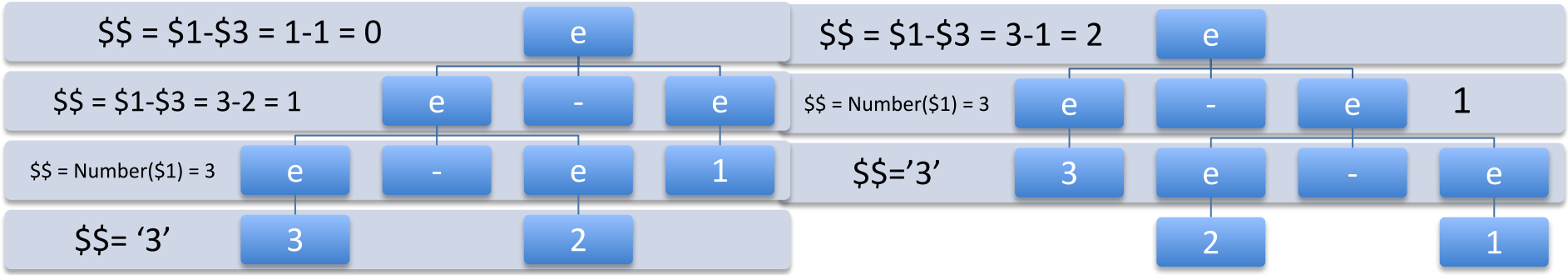


Esquema de Traducción (yacc)

e ⟶ e '-' e { $$ = $1 - $3; }

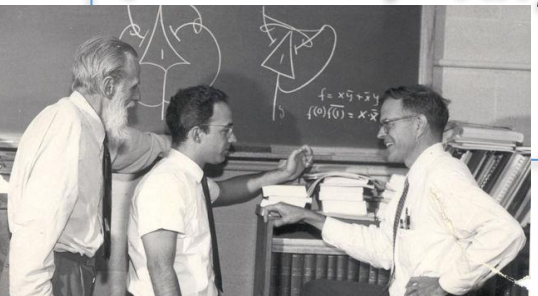
e ⟶ NUM { $$ = Number($1); }

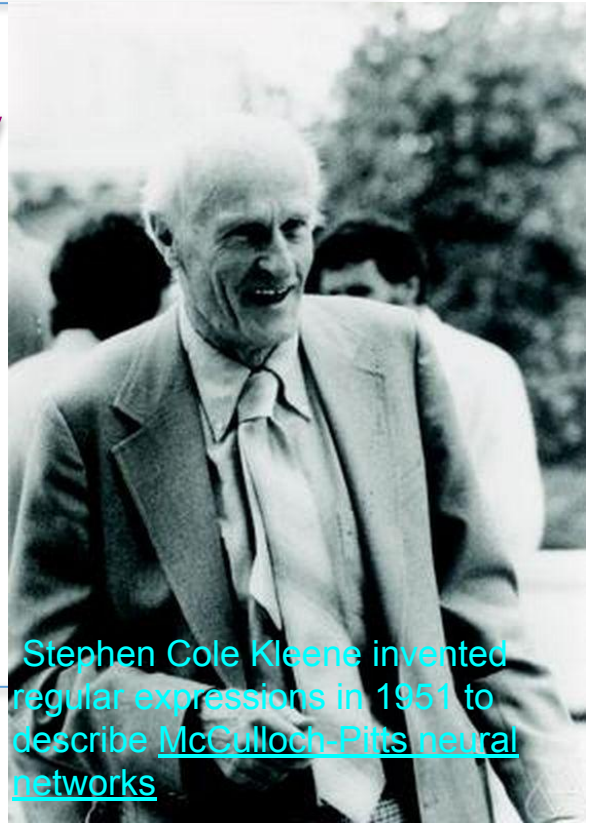
3–2-1



Análisis Léxico y Expresiones Regulares

```
[0-9]+ /* is a Natural Number */

"-"    /* is a '-' */

.      /*Any character but \n*/
```





Stephen Cole Kleene invented regular expressions in 1951 to describe McCulloch-Pitts neural networks

Un Programa que Evalúa Expresiones

https://nolanlawson.github.io/jison-debugger/

```
%lex
%%
[0-9]+          return 'NUMBER'
"-"             return '-'
.               return 'INVALID'
/lex

%%
es: e           {return $1} ;
e : e '-' e     {$$ = $1-$3}
  | NUMBER      {$$ = Number($1)} ;
```

Parser Generators: an example

https://nolanlawson.github.io/jison-debugger/

Jison debugger!

Write your grammar

```
%start er

%%

er:
    er '|' er
  | t
  ;

t: t f
  | f
  ;

f : '(' er ')'
  | f '?'
  | f '*'
  | f '+'
  | '.'
  | '^'
  | '$'
  | CHAR
  ;
```

Load a sample grammar

Choose a grammar…

Parse some text Multiline

a|b*c

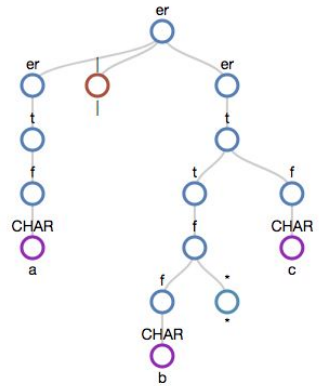
Compiled grammar

Download as JavaScript Download as JSON

Tokens

a	|	b	*	c	EOF
CHAR	|	CHAR	*	CHAR	$end

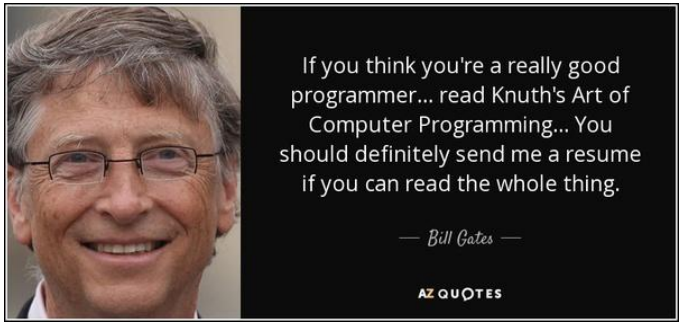
Parse tree Show log

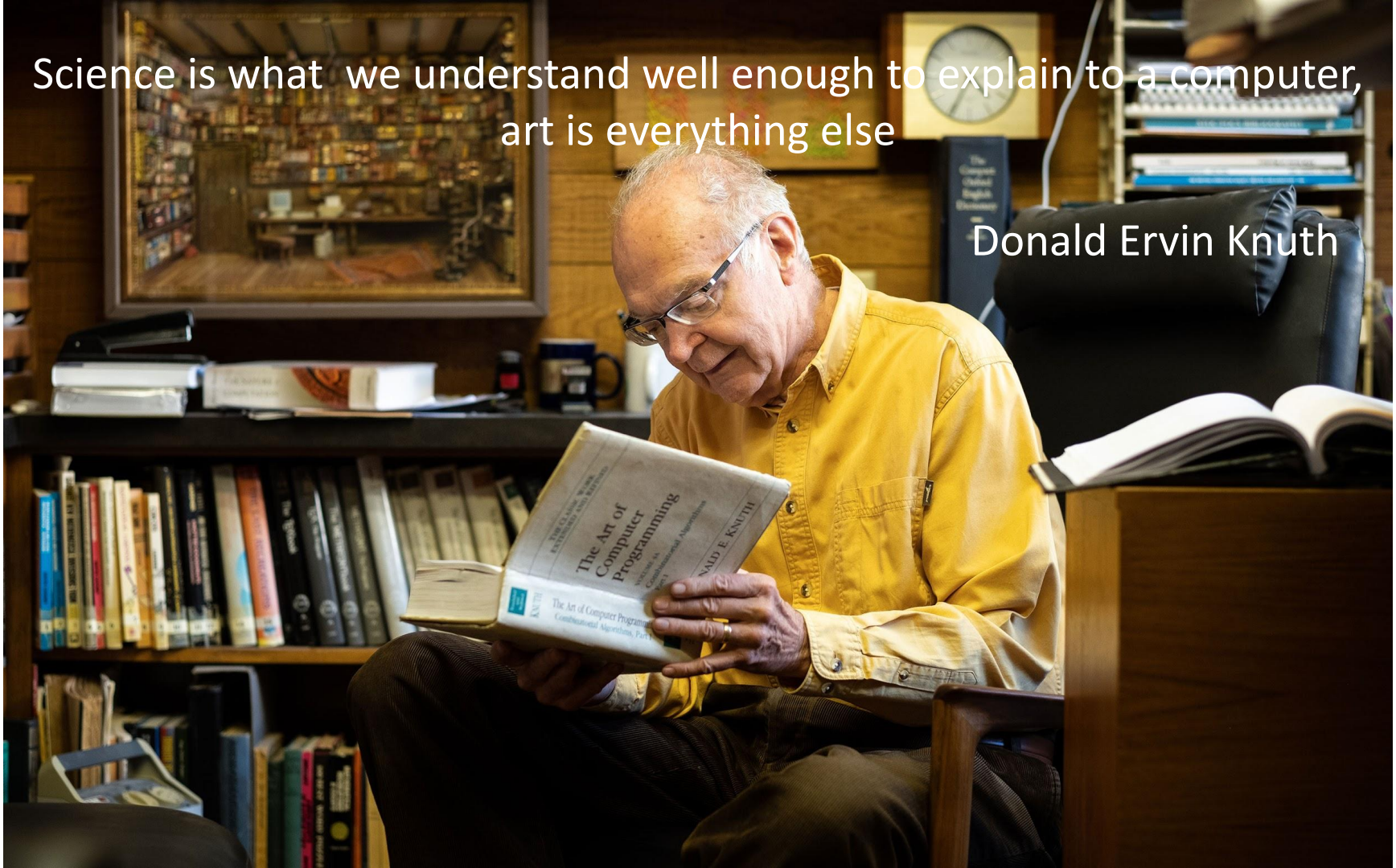


Parser result

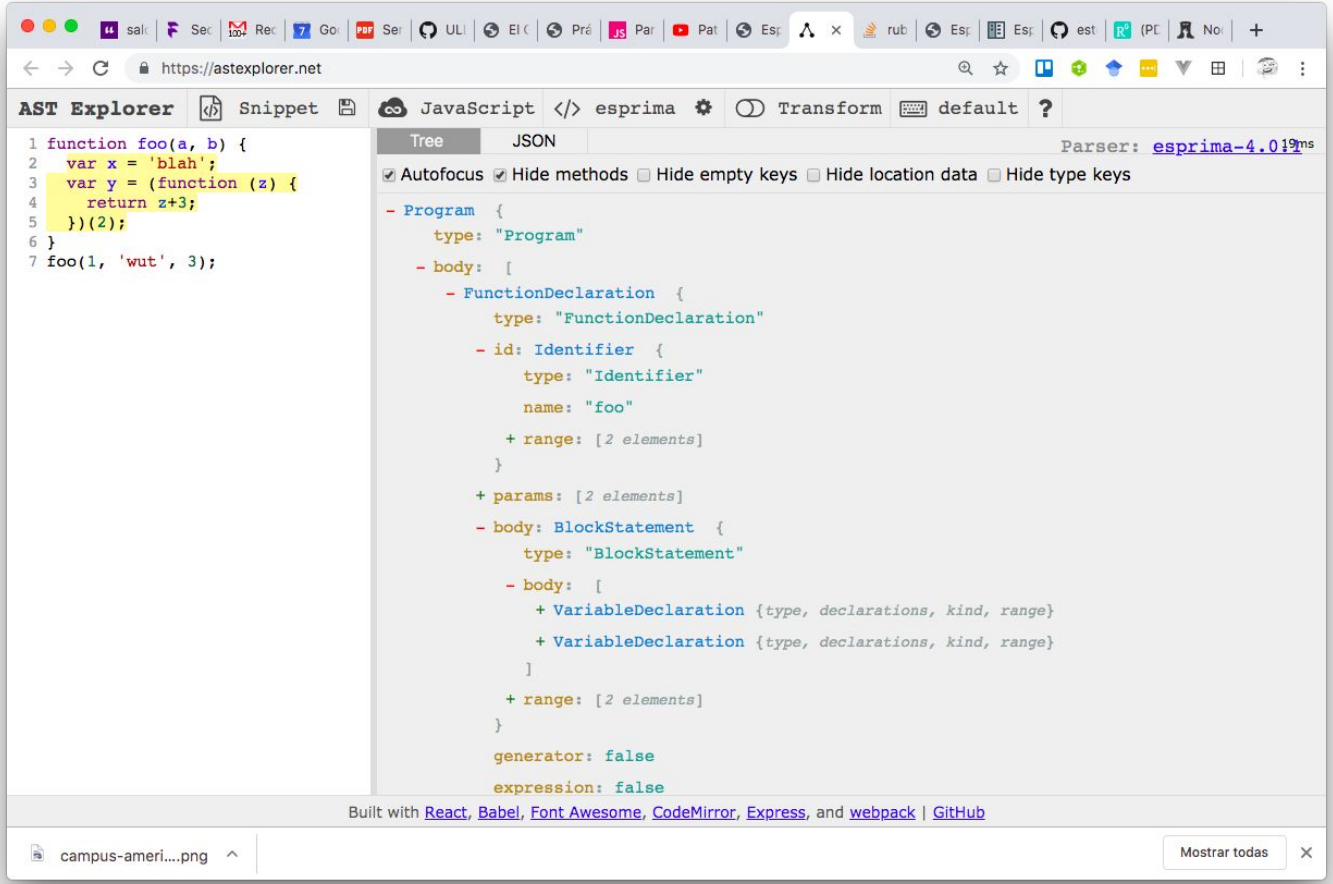
```
true
```

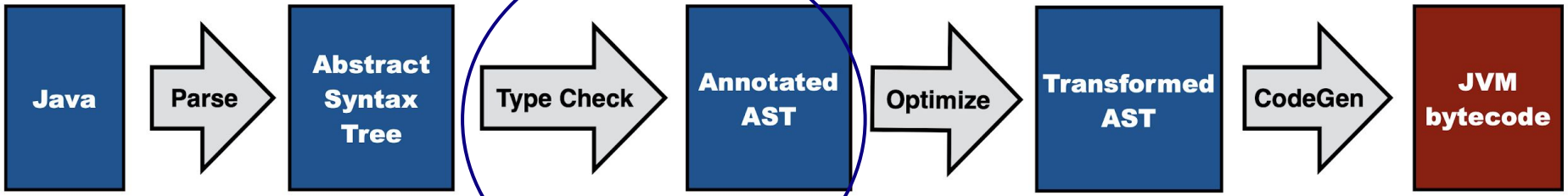
This tool, a parser generator uses a parsing algorithm known as LALR that was invented by Donald Ervin Knuth (1965)




Science is what we understand well enough to explain to a computer,
art is everything else
Donald Ervin Knuth
The Art of Computer Programming

The *Abstract Syntax Tree*: a data structure representing the underlying syntactic structure of the input program: https://astexplorer.net/



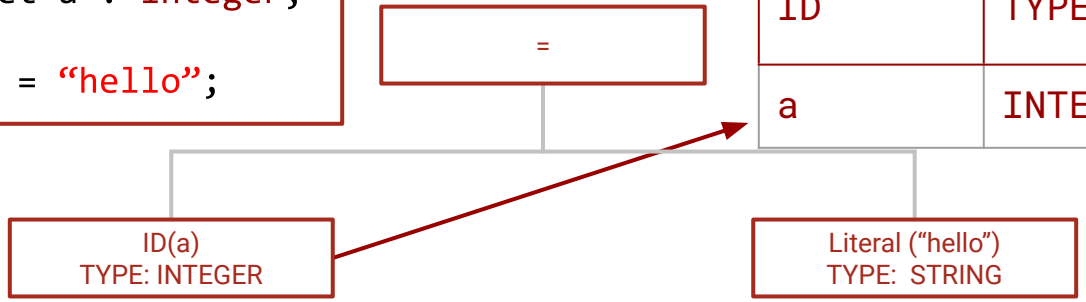


- Receives as input the abstract syntax tree
- Checks that the program complies with the static semantic rules of the language
- Performs name analysis, relating uses of names to declarations of names
- Checks that the types of arguments of operations are consistent with their specification

Input Program

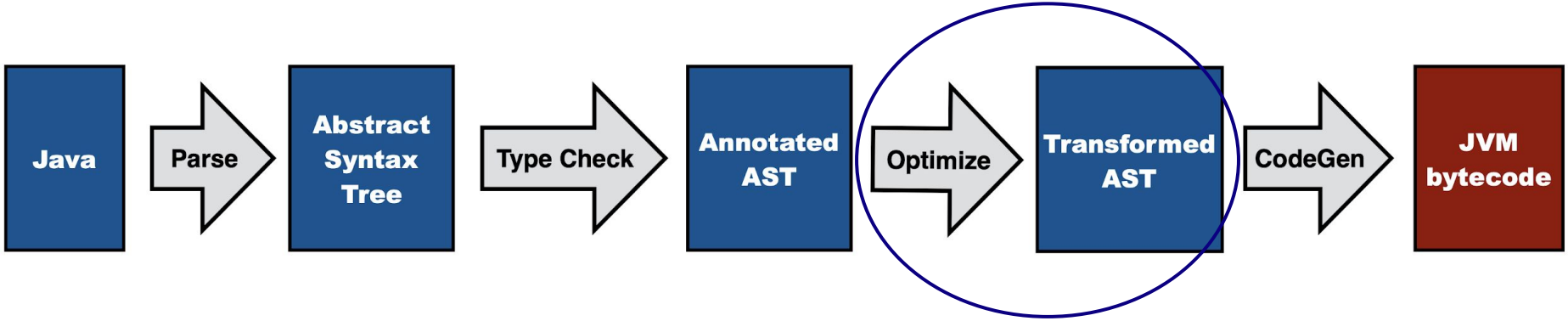
```
let a : integer;

a = “hello”;
```

AST



Symbol Table

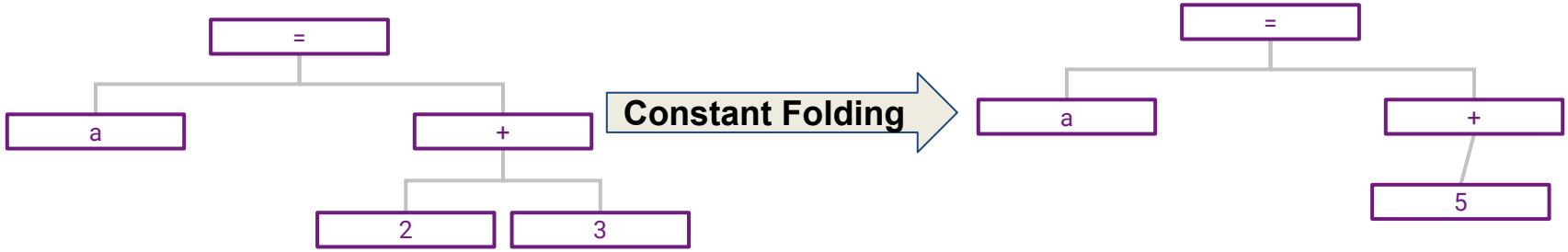
ID	TYPE
a	INTEGER

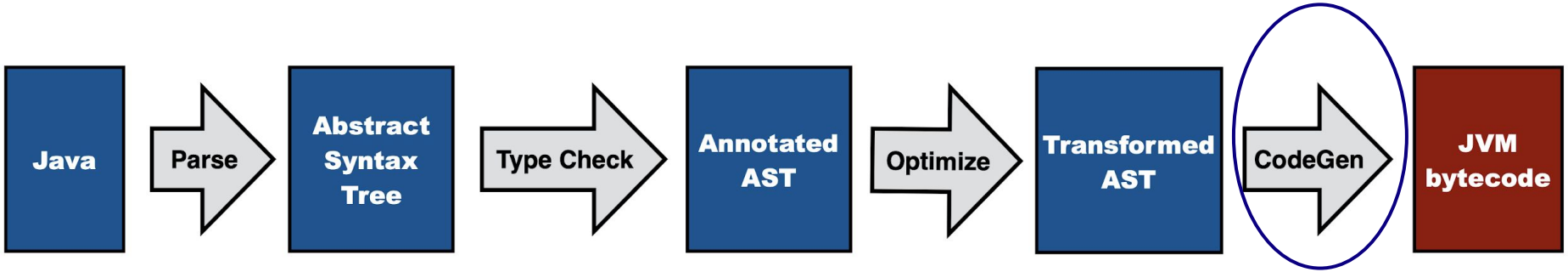


- Applies transformations that improve the program in various goals
- Goals: execution time, memory consumption, energy consumption, etc.
- Examples of transformations: Constant folding, Constant propagation, Loop invariants

Input Program

```
a = 2+3;
```

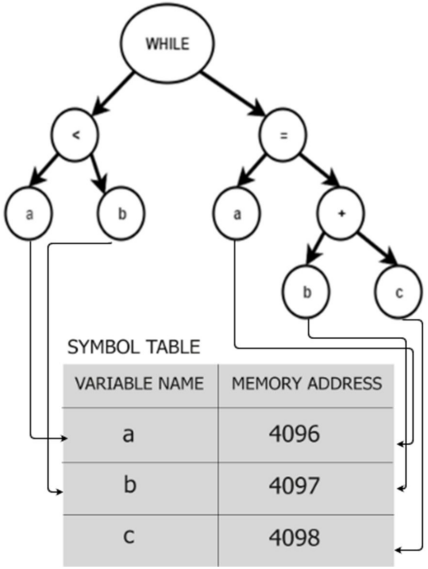




- Transforms abstract syntax tree to instructions for a particular computer architecture

Input Program

```
while (a < b) do
  a = b + c
end while
```



SYMBOL TABLE

VARIABLE NAME	MEMORY ADDRESS
a	4096
b	4097
c	4098

CODE GENERATION

```
//Translating guard
L1:
MOV R0,[4096]
MOV R1,[4097]
LT R0,R1
JZ R0,L2
```

```
//Translating body
MOV R0,[4097]
MOV R1,[4098]
ADD R0,R1
MOV [4096],R0
JMP L1
L2:
```