Chapter 1: Basic Concepts

Goals

- To provide the tools and techniques necessary to design and implement large-scale computer systems.
 - solid foundation in <u>data abstraction</u>, <u>algorithm specification</u> and <u>performance analysis and measurement</u> provides the necessary methodology.

1.1 SYSTEM LIFE CYCLE

Requirement

- a set of specifications that define the purpose of the project.
- input/output

Analysis

- break the problems down into manageable pieces.
- bottom-up / top-down

Design

- creation of abstract data types
- specification of algorithms and consideration of algorithm design strategies.

(* language independent *)

Refinement and Coding

- choose representations for data objects and write algorithms for each operation on them.
- data object's representation can determine the efficiency of the algorithms related to it.

Verification

- Developing correctness proof for the program
- Testing the program with a variety of input data
- Error removal
- Performance analysis
 - running time
 - amount of memory used

1.3 ALGORITHM SPECIFICATION

1.3.1 Introduction

Definition:

An algorithm is a finite set of instructions that, if followed, accomplishes particular task.

All algorithms must satisfy the following criteria:

- (1) **Input**: There exist quantities that are externally supplied. (zero or more)
- (2) **Output**: At least one quantity is produced.
- (3) **Definiteness**: Each instruction is <u>clear</u> and <u>unambiguous</u>.
- (4) **Finiteness**: The algorithm <u>terminates</u> after a <u>finite</u> <u>number of steps</u>.
- (5) **Effectiveness**: Every instruction must be <u>feasible</u>.
- (6) **Correctness**: The algorithm is correct with respect to a specification.

1.3 ALGORITHM SPECIFICATION

1.3.1 Introduction

- Algorithm vs. Program (procedure)
- → Algorithm always terminates (for all cases).
- → Program does not have to satisfy "finiteness"

How to describe an algorithm

- natural language
- flowchart
- programming language

Example 1.1 [Selection Sort]

Sorting a set of $n \ge 1$ integers

From those integers that are currently unsorted, find the smallest and place it next in the sorted list.

[Program 1.2 Selection sort algorithm]

```
for (i=0; i<n; i++) {
    Examine list[i] to list[n-1]
    and suppose that the smallest integer is at list[min];
    Interchange list[i] and list[min];
}</pre>
```

- first task : finding the smallest integer;
- second task : exchange;

either a function or a macro

[Program 1.3 swap function]

```
void swap(int *x, int *y)
/* both parameters are pointers to ints */
     int temp = *x; /* declares temp as an int and assigns to it
                        the contents of what x points to */
     *x = *y; /* stores what y points to into the location
                       where x points */
     *y = temp; /* place the contents of temp in the location
                       pointed to by y */
Call -- swap(&a, &b);
```

macro version of swap -

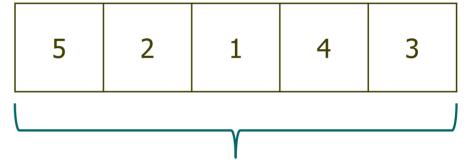
```
#define SWAP(x,y,t) ((t) = (x), (x) = (y), (y) = (t))
```

[Program 1.4 Selection sort]

```
void sort (int list[], int n)
{
    int i, j, min, temp;
    for (i=0; i<n-1; i++) {
        min = i;
        for (j = i+1; j < n; j++)
            if (list[j] < list[min])
            min = j;
        SWAP(list[i], list[min], temp);
    }
}</pre>
```

```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

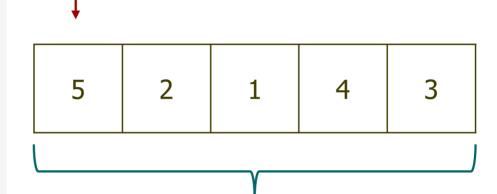
Initial state: unsorted array



```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
i = 0, j = ?, min = 0
```

min



```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

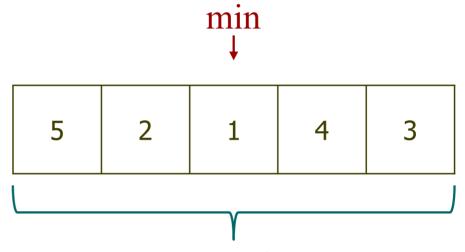
```
i = 0, j = 1, min = 1

min

5 2 1 4 3
```

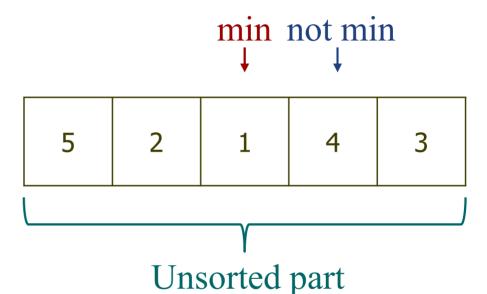
```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
i = 0, j = 2, min = 2
```



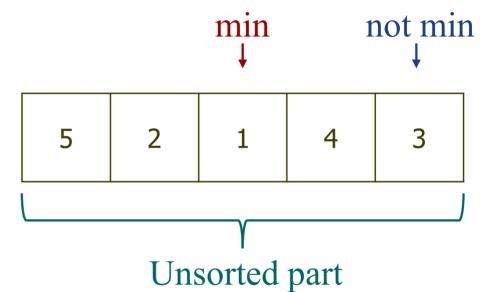
```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
i = 0, j = 3, min = 2
```



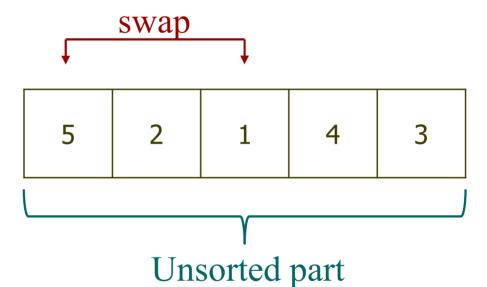
```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
i = 0, j = 4, min = 2
```



```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
i = 0, j = 4, min = 2
```

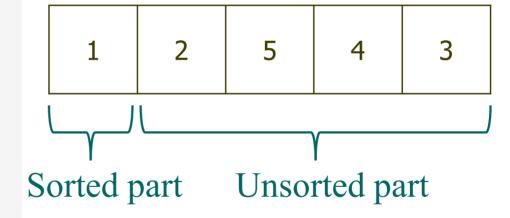


```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
After i=0:

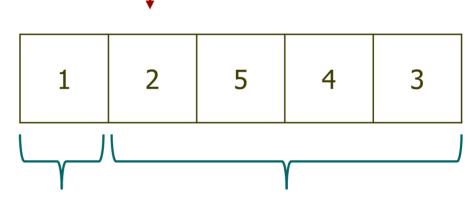
Sorted part: [1]

Unsorted part: [2][5][4][3]
```



```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
i = 1, j = ?, min = 1
min
```



```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
i = 1, j = 2, min = 1

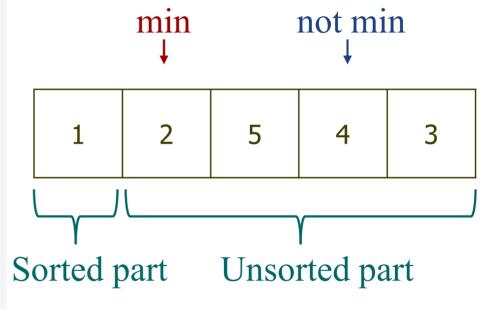
min not min

1 2 5 4 3

Sorted part Unsorted part
```

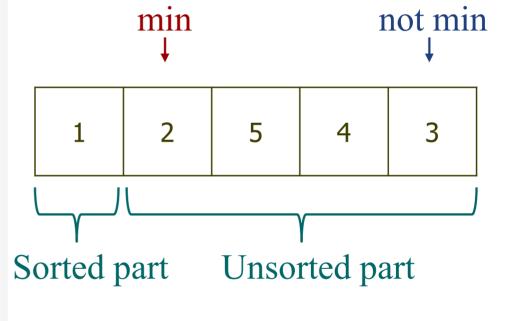
```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
i = 1, j = 3, min = 1
```



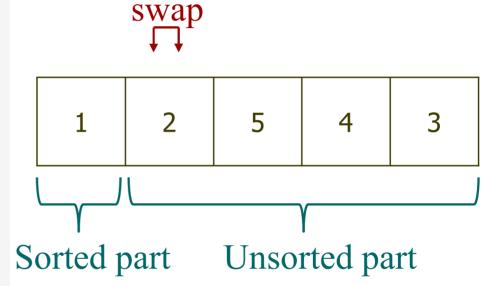
```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
i = 1, j = 4, min = 1
```



```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
i = 1, j = 4, min = 1
```

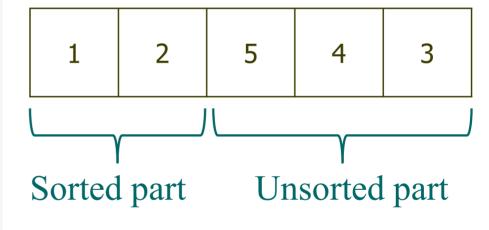


```
void sort (int list[], int n)
   int i, j, min, temp;
  for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
After i=1:

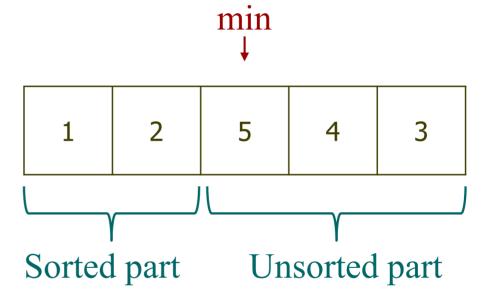
Sorted part: [1][2]

Unsorted part: [5][4][3]
```



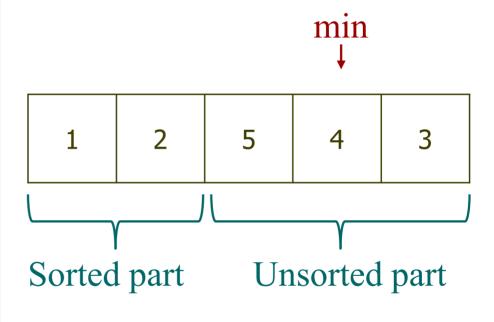
```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
i = 2, j = ?, min = 2
```



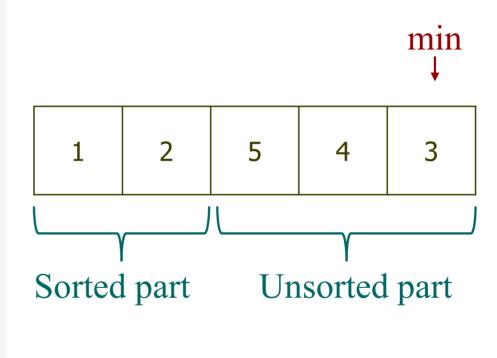
```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
i = 2, j = 3, min = 3
```



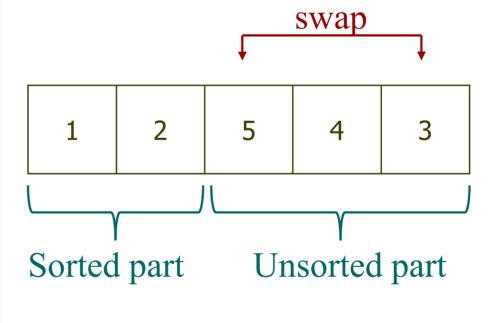
```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
i = 2, j = 4, min = 4
```



```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
i = 2, j = 4, min = 4
```

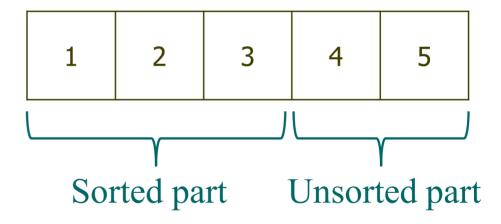


```
void sort (int list[], int n)
   int i, j, min, temp;
  for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
After i=2:

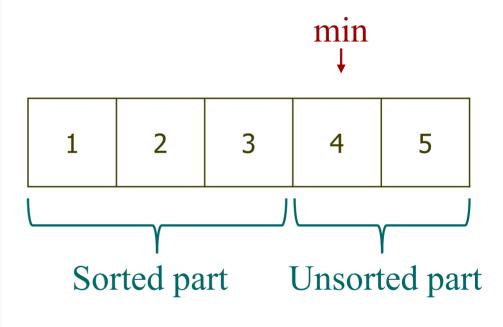
Sorted part: [1][2][3]

Unsorted part: [4][5]
```



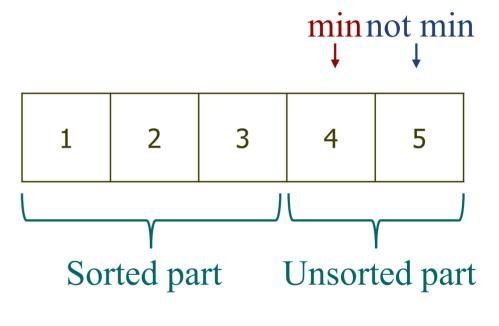
```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
i = 3, j = ?, min = 3
```



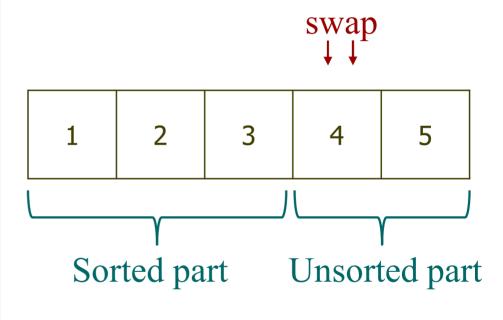
```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
i = 3, j = 4, min = 3
```



```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

```
i = 3, j = 4, min = 3
```

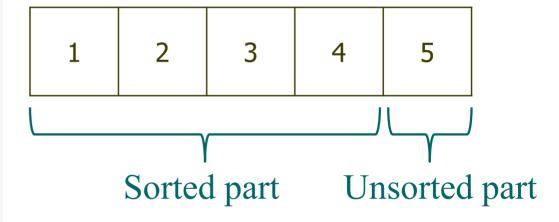


```
void sort (int list[], int n)
   int i, j, min, temp;
   for (i=0; i<n-1; i++) {
      min = i;
      for (j = i+1; j < n; j++)
         if (list[j] < list[min])</pre>
            min = j;
      SWAP(list[i], list[min],
              temp);
```

After i=3:

Sorted part: [1][2][3][4]

Unsorted part: [5]



■ Theorem 1.1:

Function sort(list,n) correctly sorts a set of $n \ge 1$ integers. The result remains in list[0], ..., list[n-1] such that $list[0] \le list[1] \le ... \le list[n-1]$.

proof:

When the outer **for** loop completes its iteration for i = q, we have $list[q] \le list[r], q < r < n$.

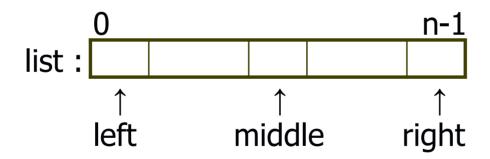
Further, on subsequent iterations, i > q and list[0] through list[q] are unchanged.

Hence following the last iteration of the outer **for** loop (i.e., i = n - 2), we have $list[0] \le list[1] \le \cdots \le list[n - 1]$.

Example 1.2 [Binary Search]

Find out if an integer *searchnum* is in a list. If so, return i such that list[i] = *searchnum*, Otherwise, return -1.

For a sorted list (in nondecreasing order)



$$middle = (left + right) / 2$$

Compare list[middle] with searchnum

searchnum < list[middle]</p>

if *searchnum* is present, it must be in the positions between *left* and *middle-1*.
set *right* to *middle-1*.

- searchnum = list[middle]
 return middle.
- searchnum > list[middle]

if *searchnum* is present, it must be in the positions between *middle+1* and *right*.
set *left* to *middle+1*

Implementing this search strategy : [Program 1.5]

```
while (there are more integers to check) {
    middle = (left + right) / 2;
    if (searchnum < list[middle])
        right = middle - 1;
    else if (searchnum == list[middle])
        return middle;
    else left = middle + 1;
}</pre>
```

```
searchnum = 10
   middle = 1
list
                  23
                           55
         10
              16
                       41
                                64
left = 0 right = 2
middle = (left+right)/2
       = (0+2)/2
```

```
searchnum = 10

middle = 1
list

1 7 10 16 23 41 55 64

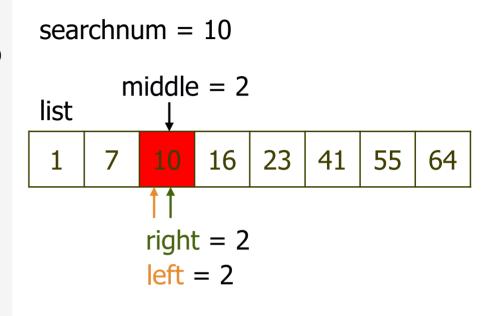
right = 2
left = 2
left = middle + 1
= 1+1
= 2
```

```
searchnum = 10

middle = 2
list

1 7 10 16 23 41 55 64

right = 2
left = 2
middle = (left+right)/2
= (2+2)/2
= 2
```



Handling the comparisons:

function -

```
int compare (int x, int y)
{
    /* compare x and y, return -1 for less than,
    0 for equal, 1 for greater */
    if (x < y) return -1;
    else if (x == y) return 0;
    else return 1;
}</pre>
```

macro -

define COMPARE (x,y) ((x) < (y)) ? -1: ((x) == (y)) ? 0: 1)

[Program 1.7]

```
int binsearch(int list[], int searchnum, int left, int right)
   /* search list[0] \leq list[1] \leq . . . \leq list[n-1] for searchnum.
      Return its position if found. Otherwise return -1 */
   int middle;
                                                             cf. Without COMPARE macro
   while (left <= right) {
                                                             while(left<=right) {</pre>
      middle = (left + right)/2;
                                                               middle = (left + right)/2;
      switch (COMPARE(list[middle], searchnum))
                                                               if(searchnum < list[middle])</pre>
                                                                 right = middle - 1;
         case -1: left = middle + 1;
                                                               else if(searchnum==list[middle])
                    break;
                                                                 return middle;
                                                               else left = middle + 1;
         case 0 : return middle;
         case 1 : right = middle - 1;
```

return -1;

1.3.2 Recursive Algorithms

- **Direct recursion:** Functions that call themselves.
- **Indirect recursion:** Functions that invoke the calling function again.
- Recursion is a general control scheme.
- Often recursive function is easier to understand than its iterative counterpart.
- Many problems can be defined recursively in natural way.

[Binomial Coefficients]

$$\left[\begin{array}{c} n \\ m \end{array}\right] = \frac{n!}{m!(n-m)!}$$

can be recursively computed by the formula:

$$\left[\begin{array}{c} n \\ m \end{array}\right] = \left[\begin{array}{c} n-1 \\ m \end{array}\right] + \left[\begin{array}{c} n-1 \\ m-1 \end{array}\right]$$

Examples:

[factorial]

$$n! = \begin{cases} n * (n-1)! & \text{if } n > 1 \\ 1 & \text{if } n = 1 \end{cases}$$

[Binary search]

[Fibonacci numbers]

$$f_{n} = \begin{cases} 0 & \text{if } n=0 \\ 1 & \text{if } n=1 \\ f_{n-1} + f_{n-2} & \text{if } n>1 \end{cases}$$

[Permutations]

We can construct the set of permutations by printing:

- (1) a followed by all permutations of (b, c, d)
- (2) b followed by all permutations of (a, c, d)
- (3) c followed by all permutations of (a, b, d)
- (4) *d* followed by all permutations of (*a*, *b*, *c*)

```
int fibo(int n)
 int g, h, f, i;
 if (n>1) {
    g = 0;
    h = 1;
    for (i = 2; i \le n; i++) {
        f = g+h;
        g = h;
        h = f;
  else f = n;
 return f;
```

Recursive function

```
int rfibo (int n)
{
   if (n > 1)
     return rfibo(n-1) + rfibo(n-2);
   else
     return n;
}
```

```
int fibo(int n)
 int g, h, f, i;
 if (n>1) {
    g = 0;
    h = 1;
    for (i = 2; i <= n; i++) {
        f = g+h;
        g = h;
        h = f;
  else f = n;
  return f;
```

Initial state

n	5
g	0
h	1
f	?
i	?

```
int fibo(int n)
 int g, h, f, i;
 if (n>1) {
    g = 0;
    h = 1;
    for (i = 2; i \le n; i++) {
       f = g+h;
        g = h;
        h = f;
  else f = n;
  return f;
```

n	5
g	0
h	1 \ +
f	angle ightarrow 1
i	2

```
int fibo(int n)
 int g, h, f, i;
 if (n>1) {
    g = 0;
    h = 1;
    for (i = 2; i \le n; i++) {
        f = g+h;
        g = h;
        h = f;
  else f = n;
  return f;
```

n	5
g	0→1
h	1→1
f	1
i	2

```
int fibo(int n)
 int g, h, f, i;
 if (n>1) {
    g = 0;
    h = 1;
    for (i = 2; i \le n; i++) {
       f = g+h;
        g = h;
        h = f;
  else f = n;
  return f;
```

n	5
g	1
h	1 \ +
f	1→2
i	3

```
int fibo(int n)
 int g, h, f, i;
 if (n>1) {
    g = 0;
    h = 1;
    for (i = 2; i \le n; i++) {
        f = g+h;
        g = h;
        h = f;
  else f = n;
  return f;
```

n	5
g	1→1
h	1→2
f	2
i	3

```
int fibo(int n)
 int g, h, f, i;
 if (n>1) {
    g = 0;
    h = 1;
    for (i = 2; i \le n; i++) {
       f = g+h;
        g = h;
        h = f;
  else f = n;
  return f;
```

n	5
g	1
h	2 +
f	2→3
i	4

```
int fibo(int n)
 int g, h, f, i;
 if (n>1) {
    g = 0;
    h = 1;
    for (i = 2; i \le n; i++) {
        f = g+h;
        g = h;
        h = f;
  else f = n;
  return f;
```

n	5
g	1→2
h	2→3
f	3
i	4

```
int fibo(int n)
 int g, h, f, i;
 if (n>1) {
    g = 0;
    h = 1;
    for (i = 2; i \le n; i++) {
       f = g+h;
        g = h;
        h = f;
  else f = n;
  return f;
```

n	5
g	2
h	3 \ +
f	3→5
i	5

```
int fibo(int n)
 int g, h, f, i;
 if (n>1) {
    g = 0;
    h = 1;
    for (i = 2; i \le n; i++) {
        f = g+h;
        g = h;
        h = f;
  else f = n;
  return f;
```

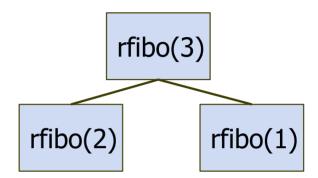
n	5
g	2→3
h	3→5
f	5
i	5

```
int rfibo (int n)
{
  if (n > 1)
    return rfibo(n-1) + rfibo(n-2);
  else
    return n;
}
```

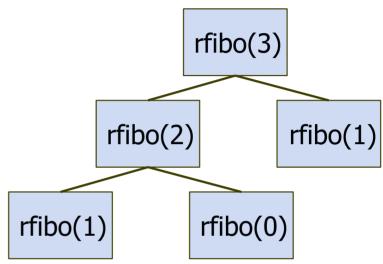
n=3

rfibo(3)

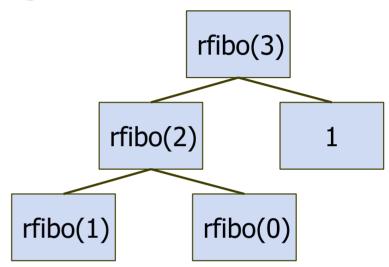
```
int rfibo (int n)
{
   if (n > 1)
      return rfibo(n-1) + rfibo(n-2);
   else
      return n;
}
```



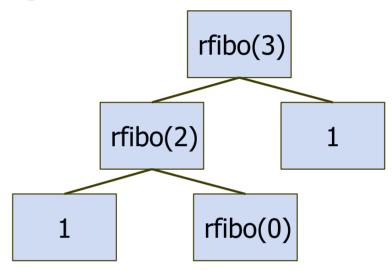
```
int rfibo (int n)
{
   if (n > 1)
      return rfibo(n-1) + rfibo(n-2);
   else
      return n;
}
```



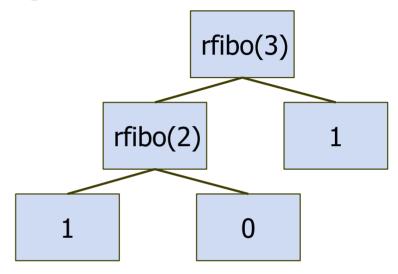
```
int rfibo (int n)
{
   if (n > 1)
     return rfibo(n-1) + rfibo(n-2);
   else
     return n;
}
```



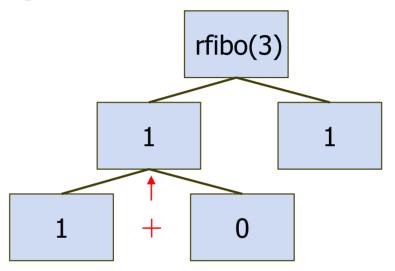
```
int rfibo (int n)
{
   if (n > 1)
     return rfibo(n-1) + rfibo(n-2);
   else
     return n;
}
```



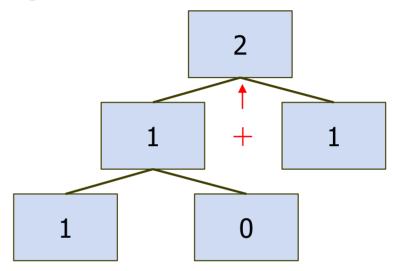
```
int rfibo (int n)
{
  if (n > 1)
    return rfibo(n-1) + rfibo(n-2);
  else
    return n;
}
```



```
int rfibo (int n)
{
   if (n > 1)
     return rfibo(n-1) + rfibo(n-2);
   else
     return n;
}
```



```
int rfibo (int n)
{
   if (n > 1)
      return rfibo(n-1) + rfibo(n-2);
   else
      return n;
}
```



[Program 1.8]

```
int binsearch(int list[], int searchnum, int left, int right)
/* search list[0] <= list[1] <= . . . <= list[n-1] for searchnum.
   Return its position if found. Otherwise return -1 */
int middle;
if (left <= right) {</pre>
  middle = (left + right)/2;
  switch (COMPARE(list[middle], searchnum)) {
      case -1 : return binsearch(list, searchnum, middle + 1, right);
      case 0 : return middle;
      case 1: return binsearch(list, searchnum, left, middle - 1);
 return -1;
```

```
int binsearch(int list[], int searchnum, int left, int right)
        /* search list[0] <= list[1] <= . . . <= list[n-1] for searchnum.
          Return its position if found. Otherwise return -1 */
        int middle;
        if (left <= right) {</pre>
          middle = (left + right)/2;
          switch (COMPARE(list[middle], searchnum)) {
             case -1: return binsearch(list, searchnum, middle + 1, right);
             case 0 : return middle;
             case 1 : return binsearch(list, searchnum, left, middle - 1);
        return -1;
                                             middle = 3
                              list
                                                  16
                                                        23
                                                                     55
searchnum = 10
                                           10
                                                              41
                                                                           64
                              left = 0
                                                                      right = 7
```

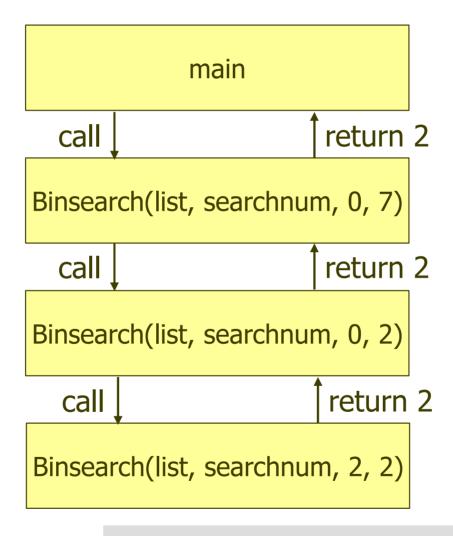
```
int binsearch(int list[], int searchnum, int left, int right)
        /* search list[0] <= list[1] <= . . . <= list[n-1] for searchnum.
           Return its position if found. Otherwise return -1 */
        int middle;
        if (left <= right) {</pre>
          middle = (left + right)/2;
          switch (COMPARE(list[middle], searchnum)) {
             case -1: return binsearch(list, searchnum, middle + 1, right);
             case 0 : return middle;
             case 1 : return binsearch(list, searchnum, left, middle - 1);
        return -1;
                                              middle = 3
                              list
                                                  16
                                                        23
                                                               41
                                                                     55
searchnum = 10
                                           10
                                                                           64
                              left = 0 \quad right = 2
```

```
int binsearch(int list[], int searchnum, int left, int right)
        /* search list[0] <= list[1] <= . . . <= list[n-1] for searchnum.
           Return its position if found. Otherwise return -1 */
        int middle;
        if (left <= right) {</pre>
           middle = (left + right)/2;
          switch (COMPARE(list[middle], searchnum)) {
             case -1: return binsearch(list, searchnum, middle + 1, right);
             case 0 : return middle;
             case 1 : return binsearch(list, searchnum, left, middle - 1);
        return -1;
                                   middle = 1
                              list
                                                        23
                                                               41
                                                                     55
searchnum = 10
                                           10
                                                  16
                                                                            64
                               left = 0 \quad right = 2
```

```
int binsearch(int list[], int searchnum, int left, int right)
        /* search list[0] <= list[1] <= . . . <= list[n-1] for searchnum.
          Return its position if found. Otherwise return -1 */
        int middle;
        if (left <= right) {</pre>
          middle = (left + right)/2;
          switch (COMPARE(list[middle], searchnum)) {
             case -1: return binsearch(list, searchnum, middle + 1, right);
             case 0 : return middle;
             case 1 : return binsearch(list, searchnum, left, middle - 1);
        return -1;
                                   middle = 1
                              list
                                                        23
                                                              41
                                                                     55
searchnum = 10
                                           10
                                                  16
                                                                           64
                                     left = right = 2
```

```
int binsearch(int list[], int searchnum, int left, int right)
        /* search list[0] <= list[1] <= . . . <= list[n-1] for searchnum.
          Return its position if found. Otherwise return -1 */
        int middle;
        if (left <= right) {</pre>
          middle = (left + right)/2;
          switch (COMPARE(list[middle], searchnum)) {
             case -1: return binsearch(list, searchnum, middle + 1, right);
             case 0 : return middle;
             case 1 : return binsearch(list, searchnum, left, middle - 1);
        return -1;
                                     middle = 2
                              list
                                                  16
                                                        23
                                                              41
                                                                     55
searchnum = 10
                                           10
                                                                           64
                                     left = right = 2
```

```
int binsearch(int list[], int searchnum, int left, int right)
        /* search list[0] <= list[1] <= . . . <= list[n-1] for searchnum.
           Return its position if found. Otherwise return -1 */
        int middle;
        if (left <= right) {</pre>
          middle = (left + right)/2;
          switch (COMPARE(list[middle], searchnum)) {
             case -1: return binsearch(list, searchnum, middle + 1, right);
             case 0 : return middle;
             case 1 : return binsearch(list, searchnum, left, middle - 1);
        return -1;
                                      middle = 2
                              list
                                                  16
                                                        23
                                                              41
                                                                     55
searchnum = 10
                                                                           64
                                     left = right = 2
```



[**Program 1.9**]

```
void perm(char *list, int i, int n)
  /* generate all the permutations of list[i] to list[n] */
  int j, temp;
  if (i == n) {
     for (j=0; j<=n; j++) printf("%c", list[j]);
     printf(" ");
  else {
  /* list[i] to list[n] has more than one permutation,
      generate these recursively */
       for (j=i; j<=n; j++) {
          SWAP(list[i], list[j], temp);
          perm(list, i+1, n);
          SWAP(list[i], list[j], temp);
```

```
void perm(char *list, int i, int n)
{ /* generate all the permutations of list[i] to list[n] */
   int j, temp;
   if (i == n) {
     for (j=0; j<=n; j++) printf("%c", list[j]);
      printf(" ");
   else {
     /* list[i] to list[n] has more than one permutation,
         generate these recursively */
      for (j=i; j<=n; j++) {
         SWAP(list[i], list[j], temp);
         perm(list, i+1, n);
         SWAP(list[i], list[j], temp);
                             list:
                                              b
```

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perm([a][b][c], 0, 2) // initial call

```
void perm(char *list, int i, int n)
                                                          perm([a][b][c], 0, 2)
{ /* generate all the permutations of list[i] to list[n] */
   int j, temp;
   if (i == n) {
     for (j=0; j<=n; j++) printf("%c", list[j]);
     printf(" ");
   else {
     /* list[i] to list[n] has more than one permutation,
        generate these recursively */
     for (j=i; j<=n; j++) {
        SWAP(list[i], list[j], temp); // j=0, i=0
        perm(list, i+1, n);
        SWAP(list[i], list[j], temp);
                            list:
                                            b
```

```
void perm(char *list, int i, int n)
                                                          perm([a][b][c], 0, 2)
{ /* generate all the permutations of list[i] to list[n] */
                                                          \xrightarrow{i=0, i=0} perm([a][b][c], 1, 2)
   int j, temp;
   if (i == n) {
     for (j=0; j<=n; j++) printf("%c", list[j]);
     printf(" ");
   else {
     /* list[i] to list[n] has more than one permutation,
         generate these recursively */
     for (j=i; j<=n; j++) {
         SWAP(list[i], list[j], temp);
         perm(list, i+1, n);
         SWAP(list[i], list[j], temp);
                            list:
                                             b
```

```
void perm(char *list, int i, int n)
                                                          perm([a][b][c], 0, 2)
{ /* generate all the permutations of list[i] to list[n] */
                                                          \xrightarrow{i=0, i=0} perm([a][b][c], 1, 2)
   int j, temp;
   if (i == n) {
     for (j=0; j<=n; j++) printf("%c", list[j]);
     printf(" ");
   else {
     /* list[i] to list[n] has more than one permutation,
        generate these recursively */
     for (j=i; j<=n; j++) {
        SWAP(list[i], list[j], temp);
        perm(list, i+1, n);
        SWAP(list[i], list[j], temp); // j=0, i=0
                            list:
                                            b
```

```
void perm(char *list, int i, int n)
                                                          perm([a][b][c], 0, 2)
{ /* generate all the permutations of list[i] to list[n] */
                                                          \xrightarrow{i=0, i=0} perm([a][b][c], 1, 2)
   int j, temp;
   if (i == n) {
     for (j=0; j<=n; j++) printf("%c", list[j]);
     printf(" ");
   else {
     /* list[i] to list[n] has more than one permutation,
        generate these recursively */
     for (j=i; j<=n; j++) {
        SWAP(list[i], list[j], temp); // j=1, i=0
        perm(list, i+1, n);
        SWAP(list[i], list[j], temp);
                            list:
                                            a
```

```
void perm(char *list, int i, int n)
                                                         perm([a][b][c], 0, 2)
{ /* generate all the permutations of list[i] to list[n] */
                                                         \xrightarrow{i=0, i=0} perm([a][b][c], 1, 2)
  int j, temp;
   if (i == n) {
     for (j=0; j<=n; j++) printf("%c", list[j]);
     printf(" ");
                                                         \rightarrow perm([b][a][c], 1, 2)
   else {
     /* list[i] to list[n] has more than one permutation,
        generate these recursively */
     for (j=i; j<=n; j++) {
        SWAP(list[i], list[j], temp);
        perm(list, i+1, n);
        SWAP(list[i], list[j], temp);
                            list:
                                            a
                                                   C
```

```
void perm(char *list, int i, int n)
                                                         perm([a][b][c], 0, 2)
{ /* generate all the permutations of list[i] to list[n] */
                                                         \xrightarrow{i=0, i=0} perm([a][b][c], 1, 2)
  int j, temp;
   if (i == n) {
     for (j=0; j<=n; j++) printf("%c", list[j]);
     printf(" ");
                                                         \rightarrow perm([b][a][c], 1, 2)
   else {
     /* list[i] to list[n] has more than one permutation,
        generate these recursively */
     for (j=i; j<=n; j++) {
        SWAP(list[i], list[j], temp);
        perm(list, i+1, n);
        SWAP(list[i], list[j], temp); // j=1, i=0
                           list:
```

```
void perm(char *list, int i, int n)
                                                         perm([a][b][c], 0, 2)
{ /* generate all the permutations of list[i] to list[n] */
                                                         \xrightarrow{i=0, i=0} perm([a][b][c], 1, 2)
  int j, temp;
   if (i == n) {
     for (j=0; j<=n; j++) printf("%c", list[j]);
     printf(" ");
                                                        \rightarrow perm([b][a][c], 1, 2)
   else {
     /* list[i] to list[n] has more than one permutation,
        generate these recursively */
     for (j=i; j<=n; j++) {
        SWAP(list[i], list[j], temp); // j=2, i=0
        perm(list, i+1, n);
        SWAP(list[i], list[j], temp);
                           list:
                                            b
                                                  a
```

```
void perm(char *list, int i, int n)
                                                         perm([a][b][c], 0, 2)
{ /* generate all the permutations of list[i] to list[n] */
                                                         \rightarrow perm([a][b][c], 1, 2)
  int j, temp;
   if (i == n) {
     for (j=0; j<=n; j++) printf("%c", list[j]);
     printf(" ");
                                                        \rightarrow perm([b][a][c], 1, 2)
   else {
     /* list[i] to list[n] has more than one permutation,
        generate these recursively */
     for (j=i; j<=n; j++) {
        SWAP(list[i], list[j], temp);
        perm(list, i+1, n);
                                                         \xrightarrow{j=2, i=0} perm([c][b][a], 1, 2)
        SWAP(list[i], list[j], temp);
                           list:
                                           b
                                                  a
```

```
void perm(char *list, int i, int n)
                                                        perm([a][b][c], 0, 2)
{ /* generate all the permutations of list[i] to list[n] */
                                                        \xrightarrow{i=0, i=0} perm([a][b][c], 1, 2)
   int j, temp;
   if (i == n) {
     for (j=0; j<=n; j++) printf("%c", list[j]);
     printf(" ");
                                                        \rightarrow perm([b][a][c], 1, 2)
   else {
     /* list[i] to list[n] has more than one permutation,
        generate these recursively */
     for (j=i; j<=n; j++) {
        SWAP(list[i], list[j], temp);
        perm(list, i+1, n);
                                                        \rightarrow perm([c][b][a], 1, 2)
        SWAP(list[i], list[j], temp); // j=2, i=0
                           list:
                                           b
```

```
void perm(char *list, int i, int n)
                                                          perm([a][b][c], 0, 2)
{ /* generate all the permutations of list[i] to list[n] */
                                                         \xrightarrow{i=0, i=0} perm([a][b][c], 1, 2)
   int j, temp;
   if (i == n) {
     for (j=0; j<=n; j++) printf("%c", list[j]);
     printf(" ");
                                                         _{j=1, i=0} perm([b][a][c], 1, 2)
   else {
     /* list[i] to list[n] has more than one permutation,
        generate these recursively */
     for (j=i; j<=n; j++) {
        SWAP(list[i], list[j], temp); // j=1, i=1
        perm(list, i+1, n);
                                                          \xrightarrow{j=2, i=0} perm([c][b][a], 1, 2)
        SWAP(list[i], list[j], temp);
                            list:
                                            b
```

```
void perm(char *list, int i, int n)
                                                          perm([a][b][c], 0, 2)
{ /* generate all the permutations of list[i] to list[n] */
                                                          \xrightarrow{i=0, i=0} perm([a][b][c], 1, 2)
   int j, temp;
   if (i == n) {
                                                                      \rightarrow perm([a][b][c], 2, 2)
     for (j=0; j<=n; j++) printf("%c", list[j]);
     printf(" ");
                                                         \xrightarrow{j=1, i=0} perm([b][a][c], 1, 2)
   else {
     /* list[i] to list[n] has more than one permutation,
        generate these recursively */
     for (j=i; j<=n; j++) {
        SWAP(list[i], list[j], temp);
        perm(list, i+1, n);
                                                           \xrightarrow{j=2, i=0} perm([c][b][a], 1, 2)
        SWAP(list[i], list[j], temp);
                            list:
                                            b
```

```
void perm(char *list, int i, int n)
                                                          perm([a][b][c], 0, 2)
{ /* generate all the permutations of list[i] to list[n] */
                                                          \xrightarrow{i=0, i=0} perm([a][b][c], 1, 2)
   int j, temp;
   if (i == n) {
                                                                      \rightarrow perm([a][b][c], 2, 2)
     for (j=0; j<=n; j++) printf("%c", list[j]);
     printf(" ");
                                                         \xrightarrow{j=1, i=0} perm([b][a][c], 1, 2)
   else {
     /* list[i] to list[n] has more than one permutation,
        generate these recursively */
     for (j=i; j<=n; j++) {
        SWAP(list[i], list[j], temp);
        perm(list, i+1, n);
                                                          \xrightarrow{j=2, i=0} perm([c][b][a], 1, 2)
        SWAP(list[i], list[j], temp); // j=1, i=1
                            list:
```

```
void perm(char *list, int i, int n)
                                                          perm([a][b][c], 0, 2)
{ /* generate all the permutations of list[i] to list[n] */
                                                          \xrightarrow{i=0, i=0} perm([a][b][c], 1, 2)
   int j, temp;
   if (i == n) {
                                                                      \rightarrow perm([a][b][c], 2, 2)
     for (j=0; j<=n; j++) printf("%c", list[j]);
     printf(" ");
                                                         \xrightarrow{j=1, i=0} perm([b][a][c], 1, 2)
   else {
     /* list[i] to list[n] has more than one permutation,
        generate these recursively */
     for (j=i; j<=n; j++) {
         SWAP(list[i], list[j], temp); // j=2, i=1
        perm(list, i+1, n);
                                                          \xrightarrow{j=2, i=0} perm([c][b][a], 1, 2)
        SWAP(list[i], list[j], temp);
                                                   b
                            list:
```

```
void perm(char *list, int i, int n)
                                                        perm([a][b][c], 0, 2)
{ /* generate all the permutations of list[i] to list[n] */
                                                        \xrightarrow{i=0, i=0} perm([a][b][c], 1, 2)
   int j, temp;
   if (i == n) {
                                                                    \rightarrow perm([a][b][c], 2, 2)
     for (j=0; j<=n; j++) printf("%c", list[j]);
                                                                    = perm([a][c][b], 2, 2)
     printf(" ");
                                                        \xrightarrow{j=1, i=0} perm([b][a][c], 1, 2)
   else {
     /* list[i] to list[n] has more than one permutation,
        generate these recursively */
     for (j=i; j<=n; j++) {
        SWAP(list[i], list[j], temp);
        perm(list, i+1, n);
                                                         \xrightarrow{i=2, i=0} perm([c][b][a], 1, 2)
        SWAP(list[i], list[j], temp);
                                                  b
                           list:
```

```
void perm(char *list, int i, int n)
                                                        perm([a][b][c], 0, 2)
{ /* generate all the permutations of list[i] to list[n] */
                                                        \xrightarrow{i=0, i=0} perm([a][b][c], 1, 2)
   int j, temp;
   if (i == n) {
                                                                   \rightarrow perm([a][b][c], 2, 2)
     for (j=0; j<=n; j++) printf("%c", list[j]);
                                                                    = perm([a][c][b], 2, 2)
     printf(" ");
                                                        \xrightarrow{j=1, i=0} perm([b][a][c], 1, 2)
   else {
     /* list[i] to list[n] has more than one permutation,
        generate these recursively */
     for (j=i; j<=n; j++) {
        SWAP(list[i], list[j], temp);
        perm(list, i+1, n);
                                                        \xrightarrow{j=2, i=0} perm([c][b][a], 1, 2)
        SWAP(list[i], list[j], temp); // j=2, i=1
                           list:
```

```
void perm(char *list, int i, int n)
                                                           perm([a][b][c], 0, 2)
{ /* generate all the permutations of list[i] to list[n] */
                                                           \xrightarrow{i=0, i=0} perm([a][b][c], 1, 2)
   int j, temp;
   if (i == n) {
                                                                       \xrightarrow{i=1, i=1} perm([a][b][c], 2, 2)
     for (j=0; j<=n; j++) printf("%c", list[i]);
                                                                       \xrightarrow{i=2, i=1} perm([a][c][b], 2, 2)
     printf(" ");
                                                          \xrightarrow{j=1, i=0} perm([b][a][c], 1, 2)
   else {
     /* list[i] to list[n] has more than one permutation,
                                                                       \xrightarrow{i=1, i=1} perm([b][a][c], 2, 2)
        generate these recursively */
     for (j=i; j<=n; j++) {
                                                                       \overline{j=2, j=1} perm([b][c][a], 2, 2)
         SWAP(list[i], list[j], temp);
        perm(list, i+1, n);
                                                           \xrightarrow{j=2, i=0} perm([c][b][a], 1, 2)
        SWAP(list[i], list[i], temp);
                                                                        \rightarrow perm([c][b][a], 2, 2)
                                                                        \xrightarrow{i=2, i=1} perm([c][a][b], 2, 2)
```

```
void perm(char *list, int i, int n)
                                                             perm([a][b][c], 0, 2)
{ /* generate all the permutations of list[i] to list[n] */
                                                             \xrightarrow{i=0, i=0} perm([a][b][c], 1, 2)
   int j, temp;
   if (i == n) {
                                                                         \overline{j=1, i=1} abc
      for (j=0; j<=n; j++) printf("%c", list[j]);
                                                                          \overline{j=2, i=1} acb
      printf(" ");
                                                             \xrightarrow{j=1, i=0} perm([b][a][c], 1, 2)
   else {
      /* list[i] to list[n] has more than one permutation,
                                                                          \overline{j=1, i=1} bac
         generate these recursively */
      for (j=i; j<=n; j++) {
                                                                         \overline{j=2, i=1} bca
         SWAP(list[i], list[j], temp);
         perm(list, i+1, n);
                                                             _{j=2, i=0} perm([c][b][a], 1, 2)
         SWAP(list[i], list[j], temp);
                                                                          \overline{j=1, i=1} cba
                                                                          \overline{j=2, i=1} cab
```

1.4 DATA ABSTRACTION

basic data types of C :

Arrays and Structs

```
char, int, float, double, . . . short, long, unsigned
```

mechanisms for grouping data together :

```
int list[5];
struct student {
    char last_name[10];
    int student_id;
    char grade;
};
```

pointer data type :

```
for every basic data type
there is a corresponding pointer data type, such as pointer-to-an-int,
pointer-to-a-real,
pointer-to-a-char,
and pointer-to-a-float.

int i, *pi;
```

predefined data types / user-defined data types

"What is a data type?"

Definition :

A *data type* is a collection of *objects* and a set of *operations* that act on those objects.

specification of objects

```
e.g., type int, \{0, +1, -1, +2, -2, \ldots, INT_MAX, INT_MIN\} specification of operations e.g., type int, +, -, *, /,  and \%.
```

▶ Algorithms should NOT make use of the representation of objects, but the functions (operations) that are provided.

Definition :

An *abstract data type* (ADT) is a data type that is organized in such a way that the specification of the objects and the specification of the operations on the objects is separated from the representation of the objects and the implementation of the operations.

an abstract data type is implementation independent.

Specification of operations consists of the names of operations, the type of its arguments, and the type of its result. Also a description what the function does without appealing to internal representation details.

e.g., *package* in Ada, *class* in C++

Categories to classify the operations of a data types:

- Constructor/creator: These functions create a new instance of the designated type.
- Transformers: These functions create an instance of the designated type, generally by using one or more other instances.
- Observers/reporters: These functions provide information about an instance of the type.

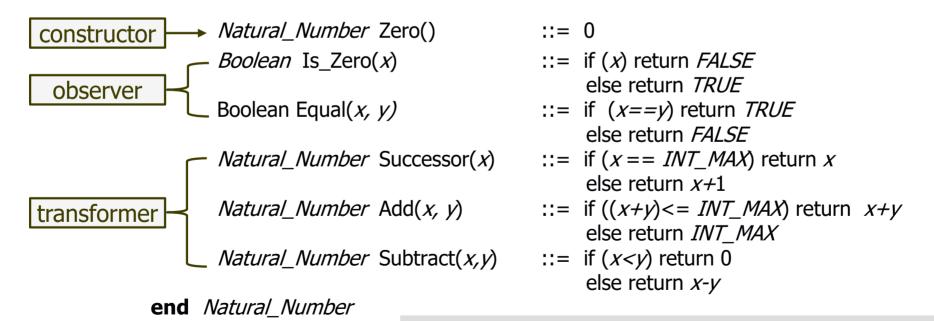
Example 1.5 [Abstract data type Natural_Number]

ADT Natural Number is

object: an ordered subrange of the integers starting at zero and ending at the maximum integer (*INT_MAX*) on the computer

functions:

for all x, y \in *Natural_Number*, *TRUE*, *FALSE* \in *Boolean* and where +, -, <, and == are usual integer operations



1.5 PERFORMANCE ANALYSIS

Criteria of judging a program:

- 1. Does the program <u>meet the original specifications</u> of the task?
- 2. Does it work correctly?
- 3. Is the program <u>well documented</u>?
- 4. Does the program <u>effectively use functions</u> to create logical units?
- 5. Is the program's code <u>readable</u>?

[Performance Evaluation]

- 6. Does the program <u>efficiently use primary and secondary storage?</u>
- 7. Is the program's <u>running time acceptable</u> for the task?

Performance Analysis:

estimates of time and space that are <u>machine independent</u>.

Performance Measurement:

obtaining <u>machine-dependent</u> running times. used to identify inefficient code segments.

Definition :

The *space complexity* of a program is the amount of memory that it needs to run to completion.

The *time complexity* of a program is the amount of computer time that it needs to run to completion.

1.5.1 Space Complexity

Fixed space requirements:

independent from the number and size of the program's inputs and outputs, e.g., the instruction space, space for simple variables, fixed-size structured variables, and constants.

Variable space requirements:

space needed by structured variables whose size depends on the particular instance, I, of the problem being solved.

The variable space requirement of a program P working on an instance I is denoted $S_P(I)$.

We can express the total space requirement S(P) of any program as:

$$S(P) = c + S_P(I)$$

where c is a constant representing the fixed space requirements.

Example 1.6 : [simple arithmetic function]

$$S_{abc}(I) = 0$$
.

[Program 1.10]

```
float abc (float a, float b, float c)
{
   return a+b+b*c + (a+b-c)/(a+b) + 4.00;
}
```

Example 1.7 : [adding a list of numbers iteratively]

[Program 1.11]

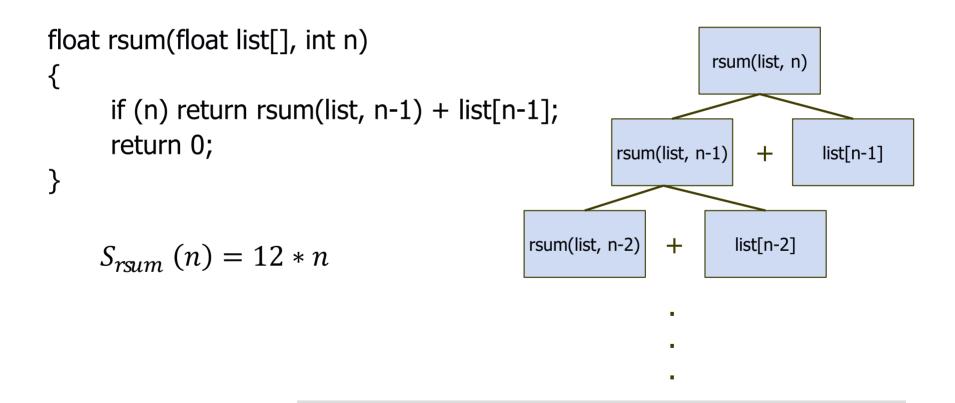
```
float sum(float list[], int n) 

{
	float tempsum = 0;
	int i;
	for (i=0; i<n; i++)
		tempsum += list[i];
	return tempsum;
}

S_{sum}(n) = n if parameters are passed by value.
	S_{sum}(n) = 0 if parameters are passed by reference.
```

Example 1.8: [adding a list of numbers recursively]

[Program 1.12]



Туре	Name	Number of bytes
parameter: array pointer parameter: integer return address: (used internally)	list[] n	4 4 4
TOTAL per recursive call		12

Figure 1.1 : Space needed for one recursive call of program 1.12

1.5.2 Time Complexity

■ The time, T_v , taken by a program, P, is the sum of:

(1) Compile Time

(2) Execution (Running) Time

We are really concerned only with the program's execution time, T_p

Determining the execution time requires a knowledge of:

- the times needed to perform each operation.
- the number of each operation performed for the given instance (dependent on the compiler).

$$T_P(n) = c_a ADD(n) + c_s SUB(n) + c_l LDA(n) + c_{st} STA(n)$$

- *n*: the instance characteristic
- c_a, c_s, c_l, c_{st} : the time needed to perform each operations.
- *ADD*, *SUB*, *LDA*, *STA*: the number of additions, subtractions, loads and stores
- Obtaining such a detailed estimate of running time is rarely worth the effort.
- Counting the number of operations the program performs gives us a machine-independent estimate.

Definition :

A *program step* is a syntactically or semantically meaningful program segment whose execution time is independent of the instance characteristics.

Determining the number of steps that a program or a function needs to solve a particular problem instance by creating a global variable, *count*, and inserting statements that increment count.

[Example 1.9] [Iterative summing of a list of numbers]

[Program 1.13]

```
float sum(float list[], int n)
{
    float tempsum = 0; count++; /*for assignment*/
    int i;
    for (i=0; i<n; i++) {
        count++; /*for the for loop */
        tempsum += list[i]; count++; /*for assignment*/
    }
    count++; /* last execution of for */
    count++; /* for return */ return tempsum;
}</pre>
```

```
float sum(float list[], int n)
{
    float tempsum = 0; count++; /*for assignment*/
    int i;
    for (i=0; i<n; i++) {
        count++; /*for the for loop */
        tempsum += list[i]; count++; /*for assignment*/
    }
    count++; /* last execution of for */
    count++; /* for return */ return tempsum;
}</pre>
```

tempsum=0	count++

```
float sum(float list[], int n)
{
    float tempsum = 0; count++; /*for assignment*/
    int i;
    for (i=0; i<n; i++) {
        count++; /*for the for loop */
        tempsum += list[i]; count++; /*for assignment*/
    }
    count++; /* last execution of for */
    count++; /* for return */ return tempsum;
}</pre>
```

tempsum=0	count++
for(i=0;i <n;i++)< td=""><td>count++</td></n;i++)<>	count++

```
float sum(float list[], int n)
{
    float tempsum = 0; count++; /*for assignment*/
    int i;
    for (i=0; i<n; i++) {
        count++; /*for the for loop */
        tempsum += list[i]; count++; /*for assignment*/
    }
    count++; /* last execution of for */
    count++; /* for return */ return tempsum;
}</pre>
```

-	
tempsum=0	count++
for(i=0;i <n;i++)< td=""><td>count++</td></n;i++)<>	count++
tempsum+=list[0]	count++

```
float sum(float list[], int n)
{
    float tempsum = 0; count++; /*for assignment*/
    int i;
    for (i=0; i<n; i++) {
        count++; /*for the for loop */
        tempsum += list[i]; count++; /*for assignment*/
    }
    count++; /* last execution of for */
    count++; /* for return */ return tempsum;
}</pre>
```

tempsum=0	count++		
for(i=0;i <n;i++)< td=""><td>count++</td></n;i++)<>	count++		
tempsum+=list[0]	count++		
for(i=0;i <n;i++)< td=""><td>count++</td></n;i++)<>	count++		

```
float sum(float list[], int n)
{
    float tempsum = 0; count++; /*for assignment*/
    int i;
    for (i=0; i<n; i++) {
        count++; /*for the for loop */
        tempsum += list[i]; count++; /*for assignment*/
    }
    count++; /* last execution of for */
    count++; /* for return */ return tempsum;
}</pre>
```

count++	
count++	
count++	
	 -r
count++	
count++	
	count++ count++

```
float sum(float list[], int n)
{
    float tempsum = 0; count++; /*for assignment*/
    int i;
    for (i=0; i<n; i++) {
        count++; /*for the for loop */
        tempsum += list[i]; count++; /*for assignment*/
    }
    count++; /* last execution of for */
    count++; /* for return */ return tempsum;
}</pre>
```

tempsum=0	count++
for(i=0;i <n;i++)< td=""><td>count++</td></n;i++)<>	count++
tempsum+=list[0]	count++
for(i=0;i <n;i++)< td=""><td>count++</td></n;i++)<>	count++
tempsum+=list[n-1]	count++
for(i=0;i <n;i++)< td=""><td>count++</td></n;i++)<>	count++

- n

i becomes n+1, then breaks the loop

```
float sum(float list[], int n)
{
    float tempsum = 0; count++; /*for assignment*/
    int i;
    for (i=0; i<n; i++) {
        count++; /*for the for loop */
        tempsum += list[i]; count++; /*for assignment*/
    }
    count++; /* last execution of for */
    count++; /* for return */ return tempsum;
}</pre>
```

tempsum=0	count++	
for(i=0;i <n;i++)< td=""><td>count++</td><td></td></n;i++)<>	count++	
tempsum+=list[0]	count++	
		 -
for(i=0;i <n;i++)< td=""><td>count++</td><td></td></n;i++)<>	count++	
tempsum+=list[n-1]	count++	
for(i=0;i <n;i++)< td=""><td>count++</td><td></td></n;i++)<>	count++	
return tempsum	count++	

[Program 1.14 Simplified version of Program 1.13]

```
cf.
float sum(float list[], int n)
                                                    float sum(float list[], int n)
   float tempsum = 0;
                                                      float tempsum = 0; count++;
   int i;
                                                      int i;
   for (i=0; i<n; i++)
                                                      for (i=0; i<n; i++) {
                                                        count++; /*for the for loop */
      /*for the for loop and assignment*
                                                        tempsum += list[i]; count++;
      count += 2;
   count += 3;
                                                      .count++;
   return tempsum;
                                                      count++;
                                                      return tempsum;
```

If the initial value of count is 0, its final value will be 2n+3.

■ [Example 1.10] [*Recursive summing of a list of numbers*]

[Program 1.15]

```
rsum(list, n)
float rsum(float list[], int n)
                                                                          count+=2
   count++; /* for if conditional */
                                                                rsum(list, n-1)
                                                                                     list[n-1]
                                                                              +
   if (n) {
                                                                 count+=2
      count++; /* for return and rsum invocation */
      return rsum(list, n-1) + list[n-1];
                                                       rsum(list, n-2)
                                                                                        n+1
                                                                     +
                                                                            list[n-2]
                                                        count+=2
   count++;
   return list[0];
                                               rsum(list, 0)
                                                             +
                                                                     list[0]
the step count is 2n+2.
                                                count+=2
```

[Example 1.11] : [Matrix addition]

[Program 1.16]

[Program 1.17]

```
void add(int a[][MAX_SIZE], int b[][MAX_SIZE],
         int c[][MAX_SIZE], int rows, int cols)
  int i, j;
  for (i=0; i<rows; i++) {
     count++; /* for i for loop */
     for (j=0; j<cols; j++) {
        count++; /* for j for loop */
        c[i][j] = a[i][j] + b[i][j];
        count++; /* for assignment statement */
     count++; /* last time of j for loop */
  count++; /* last time of i for loop */
```

[Program 1.18 Simplification of Program 1.17]

```
void add(int a[][MAX_SIZE], int b[][MAX_SIZE],
                                                           cf.
                                                           void add(...)
         int c[][MAX_SIZE], int rows, int cols)
                                                             int i, j;
   int i, j;
                                                             for (i=0; i<rows; i++) {
                                                               count++;
   for (i=0; i<rows; i++) {
                                                               for (j=0; j<cols; j++) {
      for (j=0; j<cols; j++)
                                                                 - count++;
         count += 2;←
                                                                  c[i][j] = a[i][j] + b[i][j];
      count += 2;
                                                                 count++;
                                                               count++;
   count++; ←
                                                             count++;
```

The step count will be 2rows*cols + 2rows + 1

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■ Tabular method: *steps/execution*

[Figure 1.2] Step count table for Program 1.11

Statement	s/e	Frequen	cy Total steps
float sum(float list[], int n)	0	0	0
{	0	0	0
float tempsum=0;	1	1	1
int i;	0	0	0
for (i=0; i <n; i++)<="" td=""><td>1</td><td>n+1</td><td>n+1</td></n;>	1	n+1	n+1
tempsum += list[i];	1	n	n
return tempsum;	1	1	1
}	0	0	0
Total			2n+3

[Example 1.13]

[Figure 1.3] Step count table for Program 1.12

s/e	Freque	ncy Total steps
0	0	0
0	0	0
1	n+1	n+1
1	n	n
1	1	1
0	0	0
		2n+2
	0 0 1 1 1	0 0 0 0 1 n+1 1 n 1 1

[Example 1.14]

[Figure 1.4] Step count table for Program 1.16

Statement	s/e	Frequency	Total Steps
void add(int a[][MAX_SIZE],)	0	0	0
{	0	0	0
int i, j;	0	0	0
for (i=0; i <rows; i++)<="" td=""><td>1</td><td>rows + 1</td><td>rows + 1</td></rows;>	1	rows + 1	rows + 1
for (j=0; j <cols; j++)<="" td=""><td>1</td><td>rows·(cols+1)</td><td>rows·cols + rows</td></cols;>	1	rows·(cols+1)	rows·cols + rows
c[i][j]=a[i][j]+b[i][j];	1	rows·cols	rows·cols
}	0	0	0
Total			2rows·cols + 2rows + 1

Summary

- Time complexity of a program is given by the number of steps taken by the program to compute the function it was written for.
- The number of steps is itself a function of the instance characteristics.
 - e.g., the number of inputs, the number of outputs, the magnitudes of the inputs and outputs, etc.
- Before the step count of a program can be determined, we need to know exactly which characteristics of the problem are to be used.

- For many programs, the time complexity is not dependent solely on the characteristics specified.
- The step count varies for different inputs of the same size.
 - Best case: the minimum number of steps that can be executed for the given parameters.
 - **Worst case:** the maximum number of steps that can be executed for the given parameters.
 - **Average:** the average number of steps executed on instances with the given parameters.

1.5.3 Asymptotic Notation (O, Ω , Θ)

Our motivation to determine step counts:

to compare the time complexities of two programs for the same function, and

to predict the growth in run time as the instance characteristics change.

- Determining the exact step count (either worst case or average)
 of a program can prove to be an exceedingly difficult task.
- Expending immense effort to determine the step count exactly isn't a worthwhile endeavor as the notion of a step is itself inexact.

(e.g.,
$$x = y$$
 and $x = y+z+(x/y)+(x*y*z-x/t)$ count as one step)

 Because of the inexactness of what a step stands for, the exact step count isn't very useful for comparative purposes. For most situations, step counts can be represented as a function of instance characteristics, such as

$$C_1 n^2 \le T_P(n) \le C_2 n^2 \text{ or } T_Q(n, m) = C_1 n + C_2 m.$$

What if the difference of two step counts are large? e.g., 3n + 3 versus 100n + 100.

What if two step counts are of different orders? e.g., $c_1n^2 + c_2n$ versus c_3n .

 \rightarrow If n is very large, the program with complexity c_3n will be faster than the one with complexity $c_1n^2+c_2n$, no matter what the values of c_1,c_2 and c_3 .

break even point:

- The value of n beyond which one program will be faster than the other program.
- The exact break even point cannot be determined analytically.
- The programs have to be run on a computer in order to determine the break even point.

Some terminology:

Definition : [Big "oh"]

$$f(n) = O(g(n))$$

iff there exist positive constants c and n_0 such that $f(n) \le cg(n)$ for all $n, n \ge n_0$.

$$3n + 2 = 0(n)$$

$$3n + 3 = 0(n)$$

$$100n + 6 \approx 0(n)$$

$$10n^2 + 4n + 2 \approx 0(n^2)$$

$$1000n^2 + 100n - 6 = 0(n^2)$$

$$6*2^n + n^2 = 0(2^n)$$

$$3n + 3 = 0(n^2)$$

$$10n^2 + 4n + 2 \approx 0(n^4)$$

$$3n + 2 \neq 0(1)$$

$$10n^2 + 4n + 2 \neq 0(n)$$

0(1)	a constant	$O(n^2)$	quadratic
$O(\log n)$	logarithm	$O(n^3)$	cubic
0(n)	linear	$0(2^n)$	exponential

In order for the statement f(n) = O(g(n)) to be informative, g(n) should be as small a function of n as one can come up with for which f(n) = O(g(n)).

■ Theorem 1.2:

If
$$f(n) = a_m n^m + \dots + a_1 n + a_0$$
, then $f(n) = O(n^m)$.

Proof :

$$f(n) \leq \sum_{i=0}^{m} |a_i| n^i$$

$$\leq n^m \sum_{i=0}^{m} |a_i| n^{i-m}$$

$$\leq n^m \sum_{i=0}^{m} |a_i|, \text{ for } n \geq 1.$$

So,
$$f(n) = O(n^m)$$
.

Definition : [Omega]

$$f(n) = \Omega(g(n))$$

iff there exist positive constants c and n_0 such that $f(n) \ge cq(n)$ for all $n, n \ge n_0$.

Example 1.16:

$$3n + 2 = \Omega(n)$$

$$100n + 6 = \Omega(n)$$

$$6*2^n + n^2 = \Omega(2^n)$$

$$3n + 3 = \Omega(n)$$

$$10n^2 + 4n + 2 = \Omega(n^2)$$

$$10n^2 + 4n + 2 = \Omega(n)$$

$$6*2^n + n^2 = \Omega(n^{100})$$

$$6*2^n + n^2 = \Omega(n)$$

$$10n^2 + 4n + 2 = \Omega(1)$$

$$6*2^n + n^2 = \Omega(n^2)$$

$$6*2^n + n^2 = \Omega(1)$$

In order for the statement $f(n) = \Omega(g(n))$ to be informative, g(n) should be as large a function of n as possible for which $f(n) = \Omega(g(n))$ is true.

■ Theorem 1.3 :

If
$$f(n) = a_m n^m + \dots + a_1 n + a_0$$
 and $a_m > 0$,
then $f(n) = \Omega(n^m)$.

Proof :

From **Theorem 1.2** and **Definition of big-O**,

$$\exists c, n_0 \text{ with } c>0 \text{ such that } \sum_{i=0}^{m-1} |a_i| n^i \leq c n^{m-1} \text{ for } \forall n \geq n_0.$$

Note that

$$\begin{split} \sum_{i=0}^{m-1} -a_i n^i &\leq \sum_{i=0}^{m-1} |a_i| n^i \leq c n^{m-1} \text{ and} \\ \sum_{i=0}^{m-1} a_i n^i &\geq -\sum_{i=0}^{m-1} |a_i| n^i \geq -c n^{m-1} \text{ for } \forall n \geq n_0. \end{split}$$

Proof (cont.):

Thus,

$$f(n) = a_m n^m + \sum_{i=0}^{m-1} a_i n^i \ge a_m n^m - c n^{m-1}$$

$$\ge \left(a_m - \frac{c}{n}\right) n^m \text{ for } \forall n \ge n_0.$$

Let n_1 be $\max\left\{n_0, ceiling\left(\frac{c}{a_m}\right)\right\} + 1$ and let c_1 be $a_m - \frac{c}{n_1}$. Then $c_1 > 0$ and $f(n) \ge c_1 n^m$ for $\forall n \ge n_1$. \square

Definition : [Theta]

$$f(n) = \Theta(g(n))$$

iff there exist positive constants c_1, c_2 and n_0 such that $c_1g(n) \le f(n) \le c_2g(n)$ for all $n, n \ge n_0$.

Example 1.17:

$$3n + 2 \approx \Theta(n)$$

$$10n^2 + 4n + 2 = 0(n^2)$$

$$10 * \log n + 4 = \Theta(\log n)$$

$$3n + 3 = \Theta(n)$$

$$6*2^n + n^2 \approx \Theta(2^n)$$

$$3n + 2 \neq \Theta(1)$$

$$10n^2 + 4n + 2 \neq 0(n)$$

$$6*2^n + n^2 \neq \Theta(n^{100})$$

$$6*2^n + n^2 \neq \Theta(1)$$

$$3n + 2 \neq \Theta(n^2)$$

$$10n^2 + 4n + 2 \neq 0(1)$$

$$6*2^n + n^2 \neq \Theta(n^2)$$

Theorem 1.4 :

If
$$f(n) = a_m n^m + \dots + a_1 n + a_0$$
 and $a_m > 0$,
then $f(n) = \Theta(n^m)$.

Proof:

By Theorem 1.2,

 $\exists c_2, n_2 \text{ such that } c_2 > 0 \text{ and } f(n) \leq c_2 n^m \text{ for } \forall n \geq n_2.$

Likewise, by **Theorem 1.3**,

 $\exists c_1, n_1 \text{ such that } c_1 > 0 \text{ and } f(n) \ge c_1 n^m \text{ for } \forall n \ge n_1.$

Let n_0 be $\max(n_1, n_2)$, then

$$c_1 n^m \le f(n) \le c_2 n^m \text{ for } \forall n \ge n_0. \square$$

Example 1.18: [Complexity of matrix addition]

Statement	Asymptotic complexity
void add(int a[][MAX_SIZE])	0
{	0
int i, j;	0
for (i=0; i <rows; i++)<="" td=""><td>Ø (rows)</td></rows;>	Ø (rows)
for (j=0; j <cols; j++)<="" td=""><td>Ø (rows · cols)</td></cols;>	Ø (rows · cols)
c[i][j] = a[i][j] + b[i][j];	Ø(rows · cols)
)	0
Total	Ø (rows · cols)

Example 1.19 : [Binary Search]

[Program 1.7] (Refer to next page)

The instance characteristic -- number of elements in the list.

Each iteration of *while* loop takes $\Theta(1)$ time.

The *while* loop is iterated at most $\lceil \log_2(n+1) \rceil$ times.

Worst case - the loop is iterated $\Theta(\log n)$ times Best case - $\Theta(1)$.

[Program 1.7]

```
int binsearch(int list[], int searchnum, int left, int right)
   /* search list[0] \le \text{list}[1] \le \dots \le \text{list}[n-1] for searchnum.
      Return its position if found. Otherwise return -1 */
   int middle;
   while (left <= right) {
      middle = (left + right)/2;
      switch (COMPARE(list[middle], searchnum))
         case -1: left = middle + 1;
                   break;
         case 0 : return middle;
         case 1 : right = middle - 1;
   return -1;
```

Example 1.21 : [Magic square]

The magic square is an $n \times n$ matrix of integers from 1 to n^2 such that the sum of each row and column and two major diagonals is the same.

When n=5: the common sum is 65.

15	8	1	24	17
16	14	7	5	23
22	20	13	6	4
3	21	19	12	10
9	2	25	18	11

Coxeter's rule :

Put a one in the middle of the top row. Go up and left assigning numbers in increasing order to empty boxes. If your move cause you to jump off the square (that is, you go beyond the square's boundaries), figure out where you would be if you landed on a box on the opposite side of the square. Continue with this box. If a box is occupied, go down instead of up and continue.

[Program 1.23]

```
printf ("Enter the size of the square: ");
scanf("%d", &size);
/* check for input errors */
 if (size<1 || size>MAX_SIZE+1) {
         fprintf(stderr, "Error! Size is out of range₩n");
         exit(1);
 if (!(size % 2)) {
         fprintf(stderr, "Error! Size is even₩n");
         exit(1);
for (i=0; i<size; i++)
   for (j=0; j < size; j++)
         square[i][j] = 0;
```

0	0	0	0	0
0	0	0	0	0
0	0	0	0	0
0	0	0	0	0
0	0	0	0	0

```
square[0][(size-1)/2] = 1; /* middle of first row */
/* i and j are current position */
 i = 0;
 i = (size-1) / 2;
  for (count = 2; count <= size * size; count++) {
       row = (i-1 < 0)? (size-1): (i-1); /* up */
        column = (j-1 < 0)? (size-1): (j-1); /* left */
        if (square[row][column]) /* down */
         i = (++i) \% size;
                /* square is unoccupied */
        else {
         i = row;
         j = column;
        square[i][j] = count;
```

```
i = 0

j = (5-1)/2 = 2

row = ?

column = ?
```

0	0		0	0
0	0	0	0	0
0	0	0	0	0
0	0	0	0	0
0	0	0	0	0

```
square[0][(size-1)/2] = 1; /* middle of first row */
/* i and j are current position */
 i = 0;
 j = (size-1) / 2;
  for (count = 2; count <= size * size; count++) {
       row = (i-1 < 0)? (size-1): (i-1); /* up */
        column = (j-1 < 0)? (size-1): (j-1); /* left */
        if (square[row][column]) /* down */
         i = (++i) \% size;
                /* square is unoccupied */
        else {
         i = row;
         j = column;
        square[i][j] = count;
```

```
i = 0

j = 2

row = 5-1 = 4

column = j-1 = 1
```

0	0	1	0	0
0	0	0	0	0
0	0	0	0	0
0	0	0	0	0
0	0	0	0	0

```
square[0][(size-1)/2] = 1; /* middle of first row */
/* i and j are current position */
 i = 0;
 j = (size-1) / 2;
 for (count = 2; count <= size * size; count++) {
       row = (i-1 < 0)? (size-1): (i-1); /* up */
       column = (j-1 < 0)? (size-1): (j-1); /* left */
        if (square[row][column]) /* down */
         i = (++i) \% size;
               /* square is unoccupied */
       else {
         i = row;
         j = column;
       square[i][j] = count;
```

```
    i = row = 4
    j = column = 1
    row = 4
    column = 1
```

0	0	1	0	0
0	0	0	0	0
0	0	0	0	0
0	0	0	0	0
0		0	0	0

```
square[0][(size-1)/2] = 1; /* middle of first row */
/* i and j are current position */
 i = 0;
 j = (size-1) / 2;
  for (count = 2; count <= size * size; count++) {
       row = (i-1 < 0)? (size-1): (i-1); /* up */
        column = (j-1 < 0)? (size-1): (j-1); /* left */
        if (square[row][column]) /* down */
         i = (++i) \% size;
                /* square is unoccupied */
        else {
         i = row;
         j = column;
        square[i][j] = count;
```

```
i = 4

j = 1

row = 4-1 = 3

column = 1-1 = 0
```

0	0	1	0	0
0	0	0	0	0
0	0	0	0	0
0	0	0	0	0
0	2	0	0	0

```
square[0][(size-1)/2] = 1; /* middle of first row */
/* i and j are current position */
 i = 0;
 j = (size-1) / 2;
 for (count = 2; count <= size * size; count++) {
       row = (i-1 < 0)? (size-1): (i-1); /* up */
       column = (j-1 < 0)? (size-1): (j-1); /* left */
        if (square[row][column]) /* down */
         i = (++i) \% size;
               /* square is unoccupied */
       else {
         i = row;
         j = column;
       square[i][j] = count;
```

```
    i = row = 3
    j = column = 0
    row = 3
    column = 0
```

0	0	1	0	0
0	0	0	0	0
0	0	0	0	0
	0	0	0	0
0	2	0	0	0

```
square[0][(size-1)/2] = 1; /* middle of first row */
/* i and j are current position */
 i = 0;
 j = (size-1) / 2;
  for (count = 2; count <= size * size; count++) {
       row = (i-1 < 0)? (size-1) : (i-1); /* up */
        column = (j-1 < 0)? (size-1): (j-1); /* left */
        if (square[row][column]) /* down */
         i = (++i) \% size;
                /* square is unoccupied */
        else {
         i = row;
         j = column;
        square[i][j] = count;
```

```
i = 3

j = 0

row = 3-1 = 2

column = 5-1 = 4
```

0	0	1	0	0
0	0	0	0	0
0	0	0	0	0
3	0	0	0	0
0	2	0	0	0

```
square[0][(size-1)/2] = 1; /* middle of first row */
/* i and j are current position */
 i = 0;
 j = (size-1) / 2;
 for (count = 2; count <= size * size; count++) {
       row = (i-1 < 0)? (size-1): (i-1); /* up */
       column = (j-1 < 0)? (size-1): (j-1); /* left */
        if (square[row][column]) /* down */
         i = (++i) \% size;
               /* square is unoccupied */
       else {
         i = row;
         j = column;
       square[i][j] = count;
```

```
    i = row = 2
    j = column = 4
    row = 2
    column = 4
```

0	0	1	0	0
0	0	0	0	0
0	0	0	0	
3	0	0	0	0
0	2	0	0	0

```
/* output the magic square */
  printf("Magic Square of the size %d : \text{\psi}n\text{\psi}n", size);
  for (i = 0; i < size; i++) {
    for (j = 0; j < size; j++)
        printf ("%5d", square[i][j];
    printf("\text{\psi}n");
    }
  printf("\text{\psi}n \text{\psi}n");
}</pre>
```

Result magic square:

15	8	1	24	17
16	14	7	5	23
22	20	13	6	4
3	21	19	12	10
9	2	25	18	11

instance characteristic -- *n* denoting the size of the magic square.

```
the nested for loops -- \Theta(n^2) for (i=0; i<size; i++) for (j=0; j<size; j++) square[i][j] = 0;

Others --- \Theta(1)

Total asymptotic complexity is \Theta(n^2).
```

1.5.4 Practical Complexities

- The time complexity of a program is generally some function of the instance characteristics.
- This complexity function:
 - is very useful in determining how the time requirements vary as the instance characteristics changes, and
 - may also be used to compare two programs P and Q that perform the same task.

Assume that program P has complexity $\Theta(n)$ and program Q has complexity $\Theta(n^2)$.

We can assert that program P is faster than program Q for *sufficiently large* n.

How the various functions grow with n?

Instance characteristic n							
Time	Name	1	2	4	8	16	32
1	Constant	1	1	1	1	1	1
log n	Logarithmic	0	1	2	3	4	5
n	Linear	1	2	4	8	16	32
nlog n	Log linear	0	2	8	24	64	160
n²	Quadratic	1	4	16	64	256	1024
n³	Cubic	1	8	64	512	4096	32768
2 ⁿ	Exponential	2	4	16	256	65536	4294967296
n!	Factorial	1	2	24	40326	20922789888000	26313×10 ³³

Figure 1.7 Function values

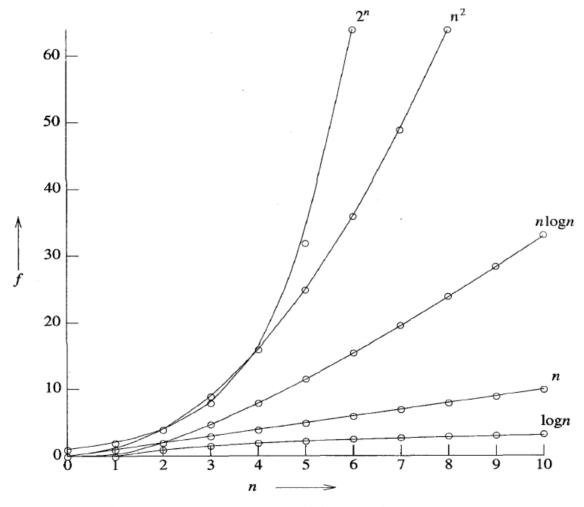


Figure 1.8 Plot of function values

	f(n)							
n	n	$n\log_2 n$	n^2	n^3	n^4	n^{10}	2^n	
10	.01 μs	.03 µs	.1 μs	1 μs	10 μs	10 s	1 μs	
20	.02 μs	.09 μs	.4 μs	8 μs	160 μs	2.84 h	1 ms	
30	.03 μ	.15 μ	.9 μ	27 μ	810 μ	6.83 d	1 s	
40	.04 μs	.21 μs	1.6 µs	64 µs	2.56 ms	121 d	18 m	
50	.05 μs	.28 μs	2.5 μs	125 μs	6.25 ms	3.1 y	13 d	
100	.10 µs	.66 μs	10 μs	1 ms	100 ms	3171 y	4*10 ¹³ y	
10^3	1 μs	9.96 µs	1 ms	1 s	16.67 m	3.17*10 ¹³ y	32*10 ²⁸³ y	
104	10 μs	130 µs	100 ms	16.67 m	115.7 d	3.17*10 ²³ y		
10 ⁵	100 μs	1.66 ms	10 s	11.57 d	3171 y	3.17*10 ³³ y		
10 ⁶	1 ms	19.92 ms	16.67 m	31.71 y	3.17*10 ⁷ y	3.17*10 ⁴³ y		

 μs = microsecond = 10⁻⁶ seconds; ms = milliseconds = 10⁻³ seconds s = seconds; m = minutes; h = hours; d = days; y = years

Figure 1.9 Times on a 1 billion instruction per second computer

1.6 PERFORMANCE MEASUREMENT

- How to measure real execution time.
 - Use of C's standard library.
 Functions are accessed through the statement:
 #include <time.h>.
 - Inaccurate results can be produced for small data
 (e.g. if the value of CLOCKS_PER_SEC is 18 on our computer,
 and the number of clock ticks is less than 10)

	Method 1	Method 2	
Start timing	start=clock();	start=time(NULL);	
Stop timing	stop=clock();	stop=time(NULL);	
Type returned	clock_t	time_t	
Result in seconds	duration= ((double)(stop-start))/ CLOCKS_PER_SEC;	duration= (double) difftime(stop, start);	

Figure 1.10: Event timing in C

[Program 1.24 First timing program for selection sort]

```
#include <stdio.h>
#include <time.h>
#include "selectionSort.h"
#define MAX_SIZE 1001
void main(void)
    int i, n, step=10;
    int a[MAX_SIZE];
    double duration;
    clock_t start;
```

```
/* times for n=0, 10, ..., 100, 200, ..., 1000 */
printf("
                         time₩n")
         n
for(n=0; n<=1000; n+=step)
{ /* get time for size n */
  /* initialize with worst-case data */
  for(i=0; i<n; i++)
     a[i] = n-i;
  start = clock();
  sort(a, n);
  duration = ((double)(clock()-start)) / CLOCKS_PER_SEC;
   printf("%6d %f₩n", n, duration);
  if(n==100) step=100;
```

- Although Program 1.24 is logically correct, it can fail to measure run times accurately because the events we are trying to time are too short.
- In **Program 1.25**, for each n, we do the sort as many times as needed to bring the total time up to 1 second. (refer to the following page)

[Program 1.25 More accurate timing program for selection sort]

```
#include <stdio.h>
#include <time.h>
#include "selectionSort.h"
#define MAX_SIZE 1001

void main(void)
{
   int i, n, step=10;
   int a[MAX_SIZE];
   double duration;
```

```
/* times for n=0, 10, ..., 100, 200, ..., 1000 */
printf("
                          repetitions
                                               time₩n")
         n
for(n=0; n<=1000; n+=step)
{ /* get time for size n */
   long repetitions = 0;
  clock_t start = clock();
   do
     repetitions++;
     /* initialize with worst-case data */
     for(i=0; i<n; i++)
        a[i] = n - i;
     sort(a, n);
   } while(clock()-start < 1000); /* repeat until enough time has elapsed */
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```

```
duration = ((double)(clock()-start)) / CLOCKS_PER_SEC;
duration /= repetitions;
printf("%6d %9d %f₩n", n, repetitions, duration);
if(n==100) step = 100;
}
```

n	repetitions	time	
0	8690714	0.000000	
10	2370915	0.000000	
20	604948	0.000002	
30	329505	0.000003	
40	205605	0.000005	
50	145353	0.000007	
60	110206	0.000009	
70	85037	0.000012	
80	65751	0.000015	
90	54012	0.000019	
100	44058	0.000023	
200	12582	0.000079	
300	5780	0.000173	
400	3344	0.000299	
500	2096	0.000477	
600	1516	0.000660	
700	1106	0.000904	
800	852	0.001174	
900	681	0.001468	
1000	550	0.001818	

Figure 1.11: Worst case performance of selection sort (in seconds)

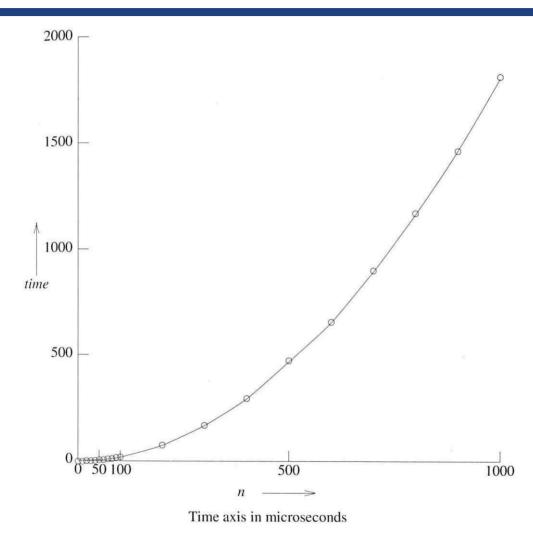


Figure 1.12: Graph of worst case performance of selection sort

Generating Test Data

- Generating a data set that results in the worst case performance of a program isn't always easy.
- We may generate a suitably large number of random test data.
- Obtaining average case data is usually much harder.
- It is desirable to analyze the algorithm being tested to determine classes of data that should be generated for the experiment - algorithm specific task.