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Tugas 5

Praktikum Analgo

Kasus 5 Mencari Pasangan Titik Terdekat

```
/*
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Praktikum Analisis Algoritma
Worksheet 5 kasus 5
Closest Pair of Points
*/
```

```
#include <iostream>
#include <float.h>
#include <stdlib.h>
#include <math.h>
using namespace std;
```

```
class Point
{
public:
    int x, y;
};
```

```
int compareX(const void* a, const void* b);
int compareY(const void* a, const void* b);
float dist(Point p1, Point p2);
float bruteForce(Point P[], int n);
float min(float x, float y) ;
float stripClosest(Point strip[], int ukuran, float d);
float closestUtil(Point P[], int n);
float closest(Point P[], int n);
```

```
int main() {
    Point P[] = {{2, 3}, {12, 30}, {40, 50}, {5, 1}, {12, 10}, {3, 4}};
    int n = sizeof(P) / sizeof(P[0]);
```

```
    cout << "Maka jarak terdekatnya adalah " << closest(P,
n)<<endl;
    return 0;
}
```

```
int compareX(const void* a, const void* b) {
    Point *p1 = (Point *)a, *p2 = (Point *)b;
    return (p1->x - p2->x);
}
```

```
int compareY(const void* a, const void* b) {
    Point *p1 = (Point *)a, *p2 = (Point *)b;
    return (p1->y - p2->y);
}
```

```
float dist(Point p1, Point p2) {
    return sqrt( (p1.x - p2.x)*(p1.x - p2.x) +
                (p1.y - p2.y)*(p1.y - p2.y));
}
```

```
float bruteForce(Point P[], int n){
    float min = FLT_MAX;
    for (int i = 0; i < n; ++i)
        for (int j = i+1; j < n; ++j)
            if (dist(P[i], P[j]) < min)
                min = dist(P[i], P[j]);
    return min;
}
```

```
float min(float x, float y) {
    return (x < y)? x : y;
}
```

```
float stripClosest(Point strip[], int ukuran, float d) {
    float min = d;
    qsort(strip, ukuran, sizeof(Point), compareY);

    for (int i = 0; i < ukuran; ++i)
```

```

        for (int j = i+1; j < ukuran && (strip[j].y - strip[i].y)
< min; ++j) {
            if (dist(strip[i],strip[j]) < min)
                min = dist(strip[i], strip[j]);
    }

    return min;
}

float closestUtil(Point P[], int n) {

    if (n <= 3)
        return bruteForce(P, n);

    int mid = n/2;
    Point midPoint = P[mid];

    float dl = closestUtil(P, mid);
    float dr = closestUtil(P + mid, n - mid);

    float d = min(dl, dr);

    Point strip[n];
    int j = 0;
    for (int i = 0; i < n; i++)
        if (abs(P[i].x - midPoint.x) < d)
            strip[j] = P[i], j++;

    return min(d, stripClosest(strip, j, d) );
}

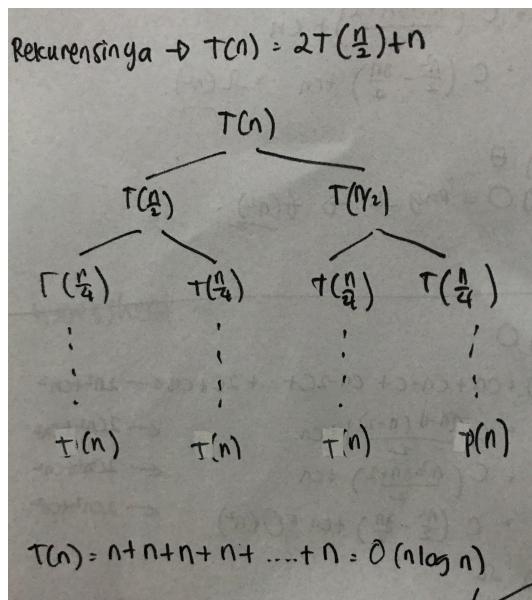
float closest(Point P[], int n) {
    qsort(P, n, sizeof(Point), compareX);

    return closestUtil(P, n);
}

```

Maka jarak terdekatnya adalah 1.41421
alvinas-MacBook-Pro:~ alvinavania\$

Tentukan rekurensi dari algoritma tersebut, dan selesaikan rekurensinya menggunakan metode recursion tree untuk membuktikan bahwa algoritma tersebut memiliki Big-O ($n \lg n$)!



Kasus 6 Algoritma Karatsuba untuk perkalian cepat

```
/*
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Praktikum Analisis Algoritma
Worksheet 5 kasus 6
Algoritma Karatsuba
*/
```

```
#include <iostream>
using namespace std;

int makeEqualLength(string &a, string &b);
string addBitStrings( string first, string second );
int multiplyiSingleBit(string a, string b);
long int multiply(string X, string Y);

int main()
{
    printf ("%ld\n", multiply("1110", "1010"));
```

```
    printf ("%ld\n", multiply("100", "1010"));
    printf ("%ld\n", multiply("11", "1010"));
    printf ("%ld\n", multiply("0", "1010"));
    printf ("%ld\n", multiply("1", "1010"));
    printf ("%ld\n", multiply("101", "111"));
    printf ("%ld\n", multiply("11", "11"));

}
```

```
int makeEqualLength(string &a, string &b){
    int p1 = a.size();
    int p2 = b.size();
    if (p1 < p2)
    {
        for (int i = 0 ; i < p2 - p1 ; i++)
            a = '0' + a;
        return p2;
    }
    else if (p1 > p2)
    {
        for (int i = 0 ; i < p1 - p2 ; i++)
            b = '0' + b;
    }
    return p1;
}
```

```
string addBitStrings( string first, string second ){
    string result;

    int length = makeEqualLength(first, second);
    int carry = 0;

    for (int i = length-1 ; i >= 0 ; i--)
    {
        int firstBit = first.at(i) - '0';
        int secondBit = second.at(i) - '0';

        int sum = (firstBit ^ secondBit ^ carry)+'0';

        result = (char)sum + result;
```

```
        carry = (firstBit&secondBit) | (secondBit&carry) |  
(firstBit&carry);  
    }  
}
```

```
    if (carry) result = '1' + result;
```

```
    return result;  
}
```

```
int multiplyiSingleBit(string a, string b){  
    return (a[0] - '0')*(b[0] - '0');  
}  
}
```

```
long int multiply(string X, string Y){
```

```
    int n = makeEqualLength(X, Y);
```

```
    if (n == 0) return 0;  
    if (n == 1) return multiplyiSingleBit(X, Y);
```

```
    int fh = n/2;  
    int sh = (n-fh);
```

```
    string Xl = X.substr(0, fh);  
    string Xr = X.substr(fh, sh);
```

```
    string Yl = Y.substr(0, fh);  
    string Yr = Y.substr(fh, sh);
```

```
    long int P1 = multiply(Xl, Yl);  
    long int P2 = multiply(Xr, Yr);  
    long int P3 = multiply(addBitStrings(Xl, Xr),  
addBitStrings(Yl, Yr));  
}
```

```
    return P1*(1<<(2*sh)) + (P3 - P1 - P2)*(1<<sh) + P2;  
}
```

```

Last login: Tue Mar 31 19:29:41 on ttys0
alvinas-MacBook-Pro:~ alvinavania$ /User
iKasus6
140
40
30
0
10
35
9
alvinas-MacBook-Pro:~ alvinavania$
```

Rekurensi dari algoritma tersebut adalah $T(n) = 3T(n/2) + O(n)$, dan selesaikan rekurensinya menggunakan metode substitusi untuk membuktikan bahwa algoritma tersebut memiliki Big-O ($n \lg n$)!

$$\begin{aligned}
& T(n) = 3T\left(\frac{n}{2}\right) + O(n) \quad \text{dgn } O(\text{gebalan}) \neq O(n \lg n) \\
& f(n) = n \lg n \quad \rightarrow n=2 \text{ karena } \frac{2}{2}=1 \\
& T(0) \leq C(f(0)) \\
& T(n) \leq C(n \lg n) \\
& \text{Substitusi:} \\
& T\left(\frac{n}{2}\right) \leq C\left(\left(\frac{n}{2}\right) \lg\left(\frac{n}{2}\right)\right) \\
& T(n) \leq C(n \lg n) + cn \\
& T(n) \leq 3(C\left(\left(\frac{n}{2}\right) \lg\left(\frac{n}{2}\right)\right)) + O(n) \\
& T(n) \leq \frac{3}{2}cn \lg n/2 + O(n) \\
& T(n) = \frac{3}{2}cn \lg n/2 - cn \lg 2 + cn \\
& T(n) = \frac{3}{2}cn \lg n - cn + cn \\
& T(n) = \frac{3}{2}n \lg n \\
& = n \lg n //
\end{aligned}$$

Kasus 7 Permasalahan Tata Letak Keramik Lantai (Tiling Problem)

```
/*
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Praktikum Analisis Algoritma
Worksheet 5 kasus 7
Tiling Problem
*/
```

```
#include <iostream>
using namespace std;

int hitung(int a, int b);

int main()
{
    int a, b;
    cout << "Masukkan angka 1: "; cin>> a;
    cout << "Masukkan angka 2: "; cin>>b;
    cout << "Maka caranya ada "<< hitung(a, b)<< endl;

    return 0;
}

int hitung(int a, int b)
{
    int count[a + 1];
    count[0] = 0;

    for (int i = 1; i <= a; i++) {
        if (i > b)
            count[i] = count[i - 1] + count[i - b];
        else if (i < b)
            count[i] = 1;
        else
            count[i] = 2;
    }
    return count[a];
```

}

```
Last login: Tue Mar 31 19:32:05 on ttys000  
/Users/alvinavania/Desktop/analgo/tugas  
alvinas-MacBook-Pro:~ alvinavania$ /Users  
iKasus7  
Masukkan angka 1:4  
Masukkan angka 2:3  
Maka caranya ada 3  
alvinas-MacBook-Pro:~ alvinavania$
```

Relasi rekurensi untuk algoritma rekursif di atas dapat ditulis seperti di bawah ini. C adalah konstanta. $T(n) = 4T(n/2) + C$. Selesaikan rekurensi tersebut dengan Metode Master !

$$T(n) = \frac{4}{2} T\left(\frac{n}{2}\right) + C \quad f(n) = n$$
$$n^{(\log_2 4)} = n^{(\log_2 4)}$$
$$= \Theta(n^{(\log_2 4 - \epsilon)})$$
$$\epsilon > 0 \quad \Theta(n^{(\log_2 4 - 1)}) \rightarrow \epsilon \approx 1$$
$$= \Theta(n^{(\log_2 3)})$$
$$= \Theta(n^{(5/3)})$$
$$= \Theta(n^2)$$

