Data Structures: Stacks, Queues and Lists

Algorithms

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- Sources of information
- 2 Stacks
- Queues
- 4 Lists

- Weiss, M.A. Data structures and algorithms analysis.
 Benjaming/Cummings, 1995. Chapter 3: Lists, stacks and queues (pages 41-86)
- * Brassard, G. and Bratley, P. Fundamentals of algorithmics. Prentice Hall, 1996. Chapter 5: Data structures (pages 147-186)
- Peña Marí, R. Diseño de Programas, Formalismo y Abstracción. Prentice Hall, 1998. Capítulo 7: Implementación de estructuras de datos (pages 257-290)

- Sources of information
- Stacks
 - Definition
 - Pseudocode
 - C Code
- Queues
- 4 Lists

- The element deleted is always the one most recently inserted
- Implements a LIFO Last In First Out policy
- Basic operations are push, pop and top
 - A top or pop on an empty stack is an error in the stack ADT
 - Running out of space performing a push is an implementation error
- All the operations take constant time notwithstanding the number of elements pushed

Array implementation of Stacks

```
type
Stack = record
  TopOfStack : [0.. MaxStackSize];
   StackArray: array [1..MaxStackSize] of ElementType
end record
procedure CreateStack (S)
 S.TopOfStack := 0
end procedure
function EmptyStack (S): boolean
  return S.TopOfStack = 0
end function
```

Array implementation of Stacks (II)

```
procedure Push (x,S)
  if S.TopOfStack = MaxStackSize then
    error "Full stack"
  else
    S.TopOfStack = S.TopOfStack + 1;
    S.StackArray [S.TopOfStack] := x
end procedure
function Top (S): ElementType
  if EmptyStack(S) then
    error "Empty stack"
  else
    return S. StackArray [S. TopOfStack]
end function
procedure Pop (S)
  if EmptyStack(S) then
    error "Empty stack"
  else
    S.TopOfStack = S.TopOfStack - 1
end procedure
```

Stack.h

```
#ifndef MaxStackSize
#define MaxStackSize 10
#endif
typedef int ElementType;
typedef struct {
  int top:
  ElementType array[MaxStackSize];
  } stack;
       void CreateStack(stack *);
        int EmptyStack(stack);
       void Push(ElementType, stack *);
ElementType Top(stack);
       void Pop(stack *);
/* ERRORS: Top or Pop over an empty stack
           Push over a full stack */
```

Stack.c

```
#include <stdlib.h>
#include <stdio.h>
#include "Stack.h"
void CreateStack(stack *S) {
  S -> top = -1:
int EmptyStack(stack S) {
  return (S.top == -1);
void Push(ElementType x, stack *S) {
  if (++S->top = MaxStackSize) {
    printf("error: Full stack\n");
    exit(EXIT FAILURE);
  S \rightarrow array[S \rightarrow top] = x:
```

Stack.c (II)

```
ElementType Top(stack S) {
  if (EmptyStack(S)) {
    printf("error: Empty stack \n");
    exit(EXIT FAILURE);
  return S.array[S.top];
void Pop(stack *S) {
  if (EmptyStack(*S)) {
    printf("error: Empty stack \n");
    exit(EXIT FAILURE);
  S \rightarrow top --;
```

- Sources of information
- Stacks
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 - Definition
 - Circular array implementation
 - Pseudocode
 - C Code
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- The element deleted is always the former inserted
- Implements a FIFO First In First Out policy
- Basic operations are enqueue, dequeue and front
 - A front or dequeue on an empty queue is an error in the queue ADT
 - Running out of space performing an enqueue is an implementation error
- All the operations take constant time notwithstanding the number of elements enqueued

- The positions QFront and QRear represent the front and the end of the queue
- Whenever these positions get the end of the array, they are wrapped around to the beginning
- Status of the queue during some operations:

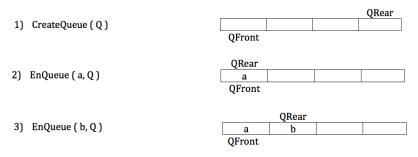
1) CreateQueue (Q)



- The positions QFront and QRear represent the front and the end of the queue
- Whenever these positions get the end of the array, they are wrapped around to the beginning
- Status of the queue during some operations:

		QRear
1) CreateQueue (Q)		
	QFront	
	QRear	
2) EnQueue (a, Q)	a	
	QFront	

- The positions QFront and QRear represent the front and the end of the queue
- Whenever these positions get the end of the array, they are wrapped around to the beginning
- Status of the queue during some operations:



4) EnQueue (c, Q)

	QRear		
a	b	С	
OFront			

4) EnQueue (c, Q)

 QRear

 a
 b
 c
 C
 QFront

5) EnQueue (d, Q)

 a
 b
 c
 d

 OFront

- 4) EnQueue (c, Q)
- 5) EnQueue (d, Q)
- 6) DeQueue (Q)

		QRear	
a	b	С	
QFront			

 a
 b
 c
 d

 QFront

		QRear
b	С	d
QFront		

4) EnQueue (c, Q)

	QRear		
a	b	С	
OFront			

5) EnQueue (d, Q)

			QRear
a	b	С	d
QFront			

6) DeQueue (Q)

		QRear
b	С	d
OFront		

7) EnQueue (e, Q)

 QRear
 c
 d

 e
 b
 c
 d

 OFront

```
type
 Queue = record
   QFront, QRear : [0..MaxQSize];
   QSize : [0..MaxQSize];
   QArray: array [1.. MaxQSize] of ElementType
 end record
procedure CreateQueue (Q)
  Q.QSize := 0:
  Q.QFront := 1;
  Q.QRear := MaxQSize
end procedure
function EmptyQueue (Q): boolean
  return Q.QSize = 0
end function
```

```
procedure Increment (x) /* private */
  if x = MaxQSize then
    x := 1
  else
    x := x + 1
end procedure
procedure Enqueue (x,Q)
  if Q.QSize = MaxQSize then
    error "Full Queue"
  else
    Q.QSize := Q.QSize + 1;
    Increment (Q. QRear);
    Q.QArray [Q.QRear] := x
end procedure
```

```
function Dequeue (Q): ElementType
  if EmptyQueue(Q) then
    error "Empty Queue"
  else
   Q.QSize := Q.QSize - 1;
    x := Q.QArray[Q.QFront];
    Increment (Q. QFront);
    return x
end function
function Front (Q) : ElementType
  if EmptyQueue(Q) then
    error "Empty Queue"
  else
    return Q. QArray [Q. QFront]
end function
```

Definition
Circular array implementation
Pseudocode
C. Code

Queue.h

```
#ifndef MaxQSize
#define MaxQSize 5
#endif
typedef int ElementType;
typedef struct {
  int QFront, QRear, QSize;
  ElementType array[MaxQSize];
  } queue:
void CreateQueue(queue *);
int EmptyQueue(queue);
void Enqueue(ElementType, queue *);
ElementType Dequeue(queue *);
ElementType Front(queue);
/* ERRORS: Dequeue or Front over an empty queue
           Enqueue over a full queue */
```

Queue.c

```
#include <stdlib.h>
#include <stdio.h>
#include "Queue.h"
void CreateQueue (Queue *Q) {
  Q \rightarrow QSize = 0:
  Q \rightarrow QFront = 0:
  Q \rightarrow QRear = -1:
int EmptyQueue (Queue Q) {
  return (Q. QSize == 0);
void Increment (int *x) { /* private */
if (++(*x) = MaxQSize)
*x = 0:
```

Queue.c (II)

```
void Enqueue (ElementType x, Queue *Q) {
  if (Q \rightarrow QSize = MaxQSize) {
    printf("error: Full Queue: %d\n", Q->QSize);
    exit(EXIT FAILURE);
  Q->QSize++:
  Increment (\&(Q->QRear));
  Q\rightarrow array[Q\rightarrow QRear] = x;
ElementType Front (Queue Q) {
  if (EmptyQueue(Q)) {
    printf("error: Empty Queue\n");
    exit(EXIT FAILURE);
  return (Q. array [Q. QFront]);
```

Queue.c (III)

```
ElementType Dequeue (Queue *Q) {
    ElementType x;

if (EmptyQueue(*Q)) {
    printf("error: EmptyQueue\n");
    exit(EXIT_FAILURE);
}
Q->QSize--;
x = Q->array[Q->QFront];
Increment(&(Q->QFront));
return x;
}
```

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Lists

- A list is a data structure consisting of a number of elements of the same type
- Basic operations:
 - Print its content
 - Find the position of the first occurrence of an element
 - Insert and Delete some element from some position
 - FindKth element, which returns the element of the indicated position

Simple array

- An array size has to be declared
 - Estimation of the maximum size is required
 - Wastes considerable space
- Computational complexity of the operations
 - FindKth, takes constant time
 - Print and Find, take linear time
 - Insert and Delete are expensive
 - Requires half the list to be moved for either operation, so linear time is required (O(n))
 - Building a list or deleting all its elements would require quadratic time

Linked

- Each node points to its successor. The last node points to nil
- The list points to the first node (and the last)
- Computational complexity of the operations:
 - Print and Find take linear time
 - Delete makes one pointer change and a dispose operation (O(1))
 - Uses FindPrevious with linear execution time
 - Insert after a position needs a new call and two pointer operations, (O(1))
 - Find the position should take linear time
 - Use a header node makes easier Insert and Delete at the beginning of the list

Doubly linked

- Each nodes points to the previous and the followings ones
- Doubles space memory needed for the pointers
- Doubles the cost of insertions and deletions pointer management
- Simplifies Delete operation
 - FindPrevious is no longer needed

Linked list with a header

```
type
  NodePtr = ^Node:
  List = NodePtr:
  Position = NodePtr;
  Node = record
    Element : ElementType;
    Next : NodePtr
  end record
procedure CreateList ( L )
  new (tmp);
  if tmp = nil then
    error "Out of space"
  else
    tmp^. Element := {Header node };
    tmp^{\cdot}.Next := nil;
    L := tmp
end procedure
```

Linked list with a header (II)

```
function EmptyList ( L ) : boolean
  return L^{\cdot}.Next = nil
end function
function Find (x, L): Position
  p := L^{\cdot}.Next;
  while p \Leftrightarrow nil and p^*. Element \Leftrightarrow x do
    p := p^. Next
  return p
end function
function IsLast ( p ) : boolean /* private */
  return p^.Next = nil
end function
```

Linked list with a header (III)

```
function FindPrevious (x, L): Position /* private */
  p := L:
  while p^. Next ⇔ nil and p^. Next^. Element ⇔ x do
    p := p^. Next
  return p
end function
procedure Delete ( x, L )
  p := FindPrevious (x, L);
  if IsLast (p) then
    error "Not found"
  else
    tmp := p^. Next;
    p^{\cdot}.Next := tmp^{\cdot}.Next;
    dispose (tmp)
end procedure
```

Linked list with a header (IV)

```
procedure Insert ( x, L, p ) /* After p position */
  new ( tmp );
  if tmp = nil then
    error "Out of space"
  else
    tmp^.Element := x;
    tmp^.Next := p^.Next;
    p^.Next := tmp
end procedure
```

Definition
Implementations
Pseudocode
C Code

List.h

```
struct node {
  void *elem; /* "void *" is a generic pointer */
  struct node *next:
};
typedef struct node *position;
typedef struct node *list;
list CreateList ();
int IsEmptyList (list I);
void Insert (void *e, position p);
/*insert e after the node pointed by p*/
position Find (list I, void *e,
            int (*comp)(const void *x, const void *y));
/*this function returns a value greater, equal
  or smaller than 0 in case x should be greater, equal
  or smaller than v */
```

List.h

```
void Delete (list I, void *e,
      int (*comp)(const void *x, const void *y));
position First (list 1);
position Next (position p);
int IsEndOfList (position p);
void *Element (position p);
 /* To go through the list:
  for (p=First(1); !IsLast(p); p=Next(p)) {
   // Do something with element(p)
```

List.c

```
#include <stdlib.h>
#include <stdio.h>
#include "List.h"
static struct node *CreateNode() {
  struct node *tmp = malloc(sizeof(struct node));
  if (tmp = NULL) {
    printf("Out of space \n");
    exit(EXIT FAILURE);
  return tmp;
list CreateList () {
  struct node *I = CreateNode ():
  I \rightarrow Next = NULL:
  return 1:
```

List.c (II)

```
int IsEmptyList (list I) {
  return (I->Next == NULL);
void Insert (void *x, position p) {
  struct node *tmp = CreateNode ();
  tmp \rightarrow Element = x:
  tmp \rightarrow Next = p \rightarrow Next;
  p \rightarrow Next = tmp:
position Find (list I, void *e,
          int (*comp)(const void *x, const void *y)){
  struct node *p = I -> Next:
  while (p != NULL \&\& 0!=(*comp)(p\rightarrow)Element, e))
    p = p -> Next;
  return p;
```

List.c (III)

List.c (IV)

```
void Delete (list I, void *x,
              int (*comp)(const void *, const void *)) {
  struct node *tmp, *p = FindPrevious (I, x, comp);
  if (! IsLast (p)) {
    tmp = p -> Next:
    p \rightarrow Next = tmp \rightarrow Next;
    free(tmp);
position First (list |) { return |->Next; }
position Next(position p) { return p->Next; }
int IsEndOfList (position p) { return (p==NULL); }
void *Element(position p) { return p->Element; }
```