

Tema 1: Circuitos Digitales

1.2 Circuitos Combinacionales

Miguel Ángel Otaduy



Universidad
Rey Juan Carlos

Circuitos Combinacionales

- Tipología
 - Salidas = función (entradas)
- Ejemplo: Sumador.
 - Entradas: Número 1, Número 2
 - Salidas: Resultado
- Tecnología de Desarrollo
 - Puertas lógicas
- Herramientas de Diseño
 - Álgebra de Boole, tabla de verdad, mapas de Karnaugh...

Circuitos Secuenciales

- Tipología
 - Salidas = función (entradas, estado)
- Ejemplo: Ascensor
 - Entradas: Botones de pisos, botones de llamada, detectores de pisos
 - Salidas: Motor (subir, bajar, parado)
- Tecnología de Desarrollo
 - Puertas lógicas y biestables
- Herramientas de Diseño
 - Máquinas de estados

Álgebra de Boole: Definición

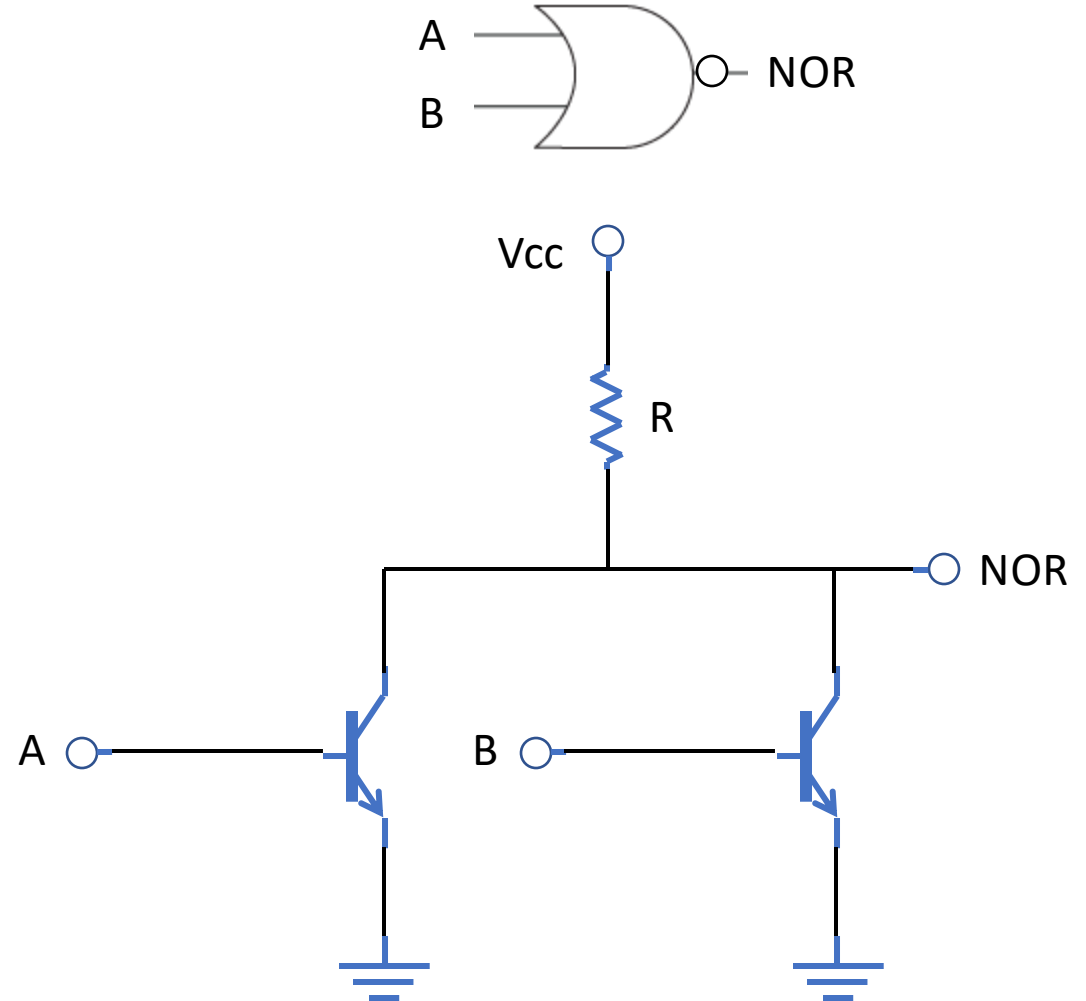
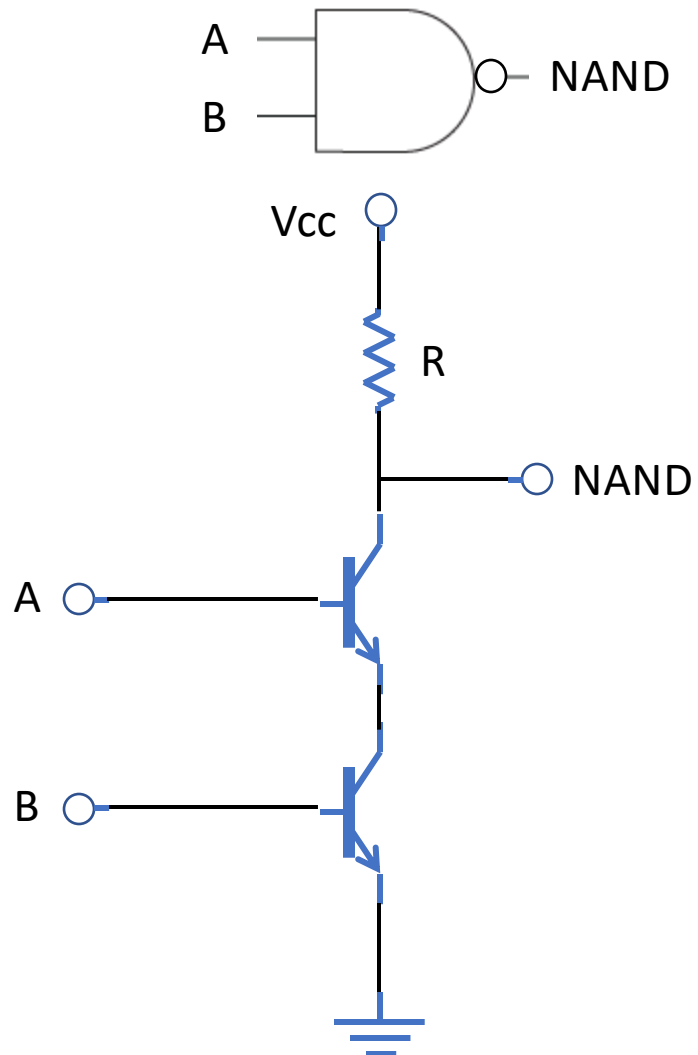
- Cada variable x puede tomar 2 valores $\{0, 1\}$
- Valor complementario
 $\text{NOT}(x)$, x' , \bar{x}
- Función producto Y lógico
 $x \text{ AND } y$, $x * y$
- Función suma O lógica
 $x \text{ OR } y$, $x + y$

| x | y | AND |
|---|---|-----|
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |

| x | y | OR |
|---|---|----|
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 1 |

| x | NOT |
|---|-----|
| 0 | 1 |
| 1 | 0 |

Puertas Lógicas



Álgebra de Boole: Teoremas y Propiedades

| | | |
|--------------------------|---|---|
| Propiedad asociativa | $a+(b+c) = (a+b)+c = a+b+c$ | $a \cdot (b \cdot c) = (a \cdot b) \cdot c = a \cdot b \cdot c$ |
| Propiedad conmutativa | $a+b = b+a$ | $a \cdot b = b \cdot a$ |
| Propiedad distributiva | $a+(b \cdot c) = (a+b) \cdot (a+c)$ | $a \cdot (b+c) = a \cdot b + a \cdot c$ |
| Elemento neutro | $0+a=a$ | $1 \cdot a=a$ |
| | $1+a=1$ | $0 \cdot a=0$ |
| Teoremas de identidad | $a+a'=1$ | $a \cdot a'=0$ |
| Teoremas de idempotencia | $a+a=a$ | $a \cdot a=a$ |
| Teorema de involución | $(a')'=a$ | |
| Teoremas de absorción | $a+a \cdot b = a$ | $a \cdot (a+b) = a$ |
| | $a+a' \cdot b=a+b$ | $a \cdot (a'+b)=a \cdot b$ |
| | $a \cdot b+a \cdot b' = a$ | $(a'+b') \cdot (a'+b) = a'$ |
| Teoremas del consenso | $a \cdot b+a' \cdot c+b \cdot c = a \cdot b+a' \cdot c$ | $(a+b) \cdot (a'+c) \cdot (b+c) = (a+b) \cdot (a'+c)$ |
| Teoremas de De Morgan | $(a+b)' = a' \cdot b'$ | $(a \cdot b)' = a'+b'$ |
| Teorema de expansión | $f(a,b) = a' \cdot f(0,b) + a \cdot f(1,b)$ | $f(a,b) = [a+f(0,b)] \cdot [a'+f(1,b)]$ |

Funciones de Conmutación

- Cada salida de un circuito combinatorial se describe mediante una función de conmutación
- Dadas n variables de entrada, una función de conmutación describe el valor (0,1) para cada una de las 2^n combinaciones de valores de entrada
- Representaciones:
 - Tabla de verdad
 - Expresiones de conmutación
 - Mapa de Karnaugh
- Orden de precedencia: NOT > AND > OR

Diseño de un Circuito Combinacional

- Definir la función de conmutación de cada salida
- Obtener su representación mediante una expresión de conmutación concisa
 - Tabla de verdad + teoremas álgebra Boole
 - Mapa de Karnaugh
- Implementar la expresión de conmutación mediante puertas lógicas

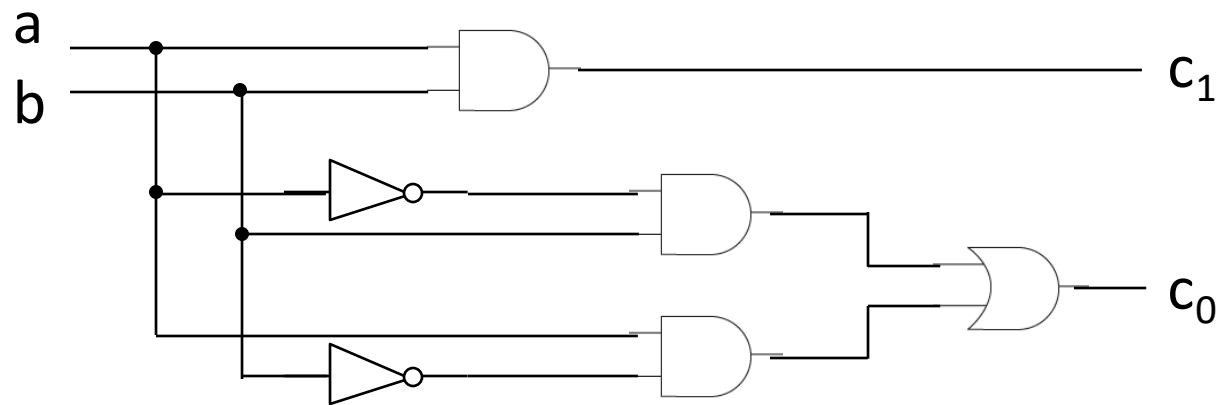
Ejemplo: Sumador de 1 bit

Sumador: $a + b = c_1c_0$

$$c_0 = a' \cdot b + a \cdot b'$$

$$c_1 = a \cdot b$$

| a | b | c_1c_0 |
|---|---|----------|
| 0 | 0 | 00 |
| 0 | 1 | 01 |
| 1 | 0 | 01 |
| 1 | 1 | 10 |



Suma de Productos Canónica

| a | b | c | out | |
|---|---|---|-----|--------------------------|
| 0 | 0 | 0 | 1 | → $a' \cdot b' \cdot c'$ |
| 0 | 0 | 1 | 0 | |
| 0 | 1 | 0 | 0 | |
| 0 | 1 | 1 | 1 | → $a' \cdot b \cdot c$ |
| 1 | 0 | 0 | 0 | |
| 1 | 0 | 1 | 0 | |
| 1 | 1 | 0 | 0 | |
| 1 | 1 | 1 | 1 | → $a \cdot b \cdot c$ |

Minterms

$$\text{out} = a' \cdot b' \cdot c' + a' \cdot b \cdot c + a \cdot b \cdot c$$

Producto de Sumas Canónico

| a | b | c | out |
|---|---|---|-----|
| 0 | 0 | 0 | 1 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 0 |
| 1 | 1 | 1 | 1 |

→ $a+b+c'$

→ $a+b'+c$

→ $a'+b+c$

→ $a'+b+c'$

→ $a'+b'+c$

Maxterms de
complementos

$$\text{out} = (a+b+c') \cdot (a+b'+c) \cdot (a'+b+c) \cdot (a'+b+c') \cdot (a'+b'+c)$$

Lógica Inversa

| a | b | c | out | out' |
|---|---|---|-----|------|
| 0 | 0 | 0 | 1 | 0 |
| 0 | 0 | 1 | 0 | 1 |
| 0 | 1 | 0 | 0 | 1 |
| 0 | 1 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 | 1 |
| 1 | 0 | 1 | 0 | 1 |
| 1 | 1 | 0 | 0 | 1 |
| 1 | 1 | 1 | 1 | 0 |

→ $a' \cdot b' \cdot c$

→ $a' \cdot b \cdot c'$

→ $a \cdot b' \cdot c'$

→ $a \cdot b' \cdot c$

→ $a \cdot b \cdot c'$

Minterms

$$\text{out}' = a' \cdot b' \cdot c + a' \cdot b \cdot c' + a \cdot b' \cdot c' + a \cdot b' \cdot c + a \cdot b \cdot c'$$

$$\text{out} = (a+b+c') \cdot (a+b'+c) \cdot (a'+b+c) \cdot (a'+b+c') \cdot (a'+b'+c)$$

Mapas de Karnaugh

1 variable

| | |
|-----|---|
| X | |
| 0 | 0 |
| 1 | 1 |

2 variables

| | | |
|----------------------|---|---|
| $X_1 \backslash X_0$ | 0 | 1 |
| 0 | 0 | 1 |
| 1 | 2 | 3 |

4 variables

| | | | | |
|------------------------------|----|----|----|----|
| $X_3 X_2 \backslash X_1 X_0$ | 00 | 01 | 11 | 10 |
| 00 | 0 | 1 | 3 | 2 |
| 01 | 4 | 5 | 7 | 6 |
| 11 | 12 | 13 | 15 | 14 |
| 10 | 8 | 9 | 11 | 10 |

5 variables

3 variables

| | | | | |
|--------------------------|----|----|----|----|
| $X_2 \backslash X_1 X_0$ | 00 | 01 | 11 | 10 |
| 0 | 0 | 1 | 3 | 2 |
| 1 | 4 | 5 | 7 | 6 |

| | | | | |
|------------------------------|----|----|----|----|
| $X_3 X_2 \backslash X_1 X_0$ | 00 | 01 | 11 | 10 |
| 00 | 0 | 1 | 3 | 2 |
| 01 | 4 | 5 | 7 | 6 |
| 11 | 12 | 13 | 15 | 14 |
| 10 | 8 | 9 | 11 | 10 |

$\overline{X_4}$ ($X_4=0$)

| | | | | |
|------------------------------|----|----|----|----|
| $X_3 X_2 \backslash X_1 X_0$ | 00 | 01 | 11 | 10 |
| 00 | 16 | 17 | 19 | 18 |
| 01 | 20 | 21 | 23 | 22 |
| 11 | 28 | 29 | 31 | 30 |
| 10 | 24 | 25 | 27 | 26 |

X_4 ($X_4=1$)

Mapas de Karnaugh: Simplificación

| $X_3 X_2 \backslash X_1 X_0$ | | | | | |
|------------------------------|--|----|----|----|----|
| | | 00 | 01 | 11 | 10 |
| 00 | | 0 | 1 | 3 | 2 |
| 01 | | 4 | 5 | 7 | 6 |
| 11 | | 12 | 13 | 15 | 14 |
| 10 | | 8 | 9 | 11 | 10 |

$$X_0 \cdot X_1' \cdot X_2 \cdot X_3' + X_0 \cdot X_1 \cdot X_2 \cdot X_3' \\ = X_0 \cdot X_2 \cdot X_3'$$

$X_0 \cdot X_2$

| $X_3 X_2 \backslash X_1 X_0$ | | | | | |
|------------------------------|--|----|----|----|----|
| | | 00 | 01 | 11 | 10 |
| 00 | | 0 | 1 | 3 | 2 |
| 01 | | 4 | 5 | 7 | 6 |
| 11 | | 12 | 13 | 15 | 14 |
| 10 | | 8 | 9 | 11 | 10 |

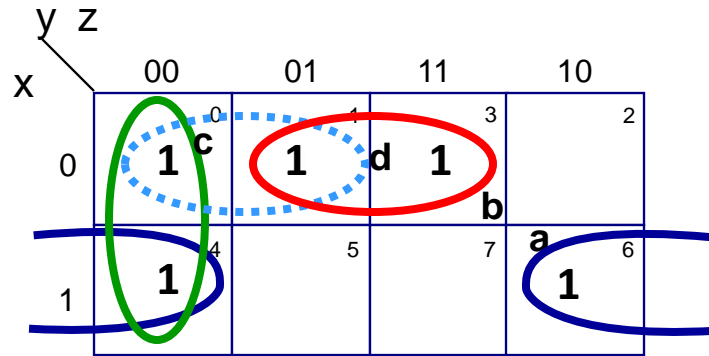
| $X_3 X_2 \backslash X_1 X_0$ | | | | | |
|------------------------------|--|----|----|----|----|
| | | 00 | 01 | 11 | 10 |
| 00 | | 0 | 1 | 3 | 2 |
| 01 | | 4 | 5 | 7 | 6 |
| 11 | | 12 | 13 | 15 | 14 |
| 10 | | 8 | 9 | 11 | 10 |

$$X_0' \cdot X_2'$$

X_2'

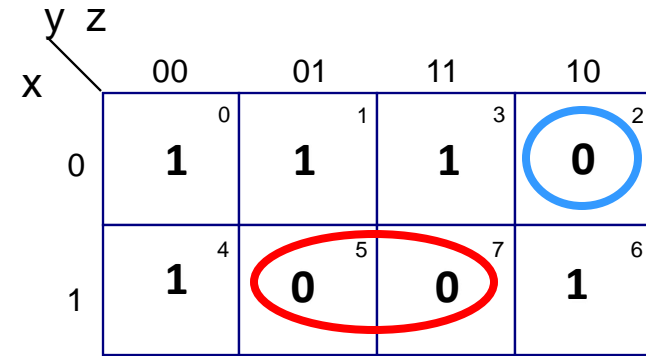
| $X_3 X_2 \backslash X_1 X_0$ | | | | | |
|------------------------------|--|----|----|----|----|
| | | 00 | 01 | 11 | 10 |
| 00 | | 0 | 1 | 3 | 2 |
| 01 | | 4 | 5 | 7 | 6 |
| 11 | | 12 | 13 | 15 | 14 |
| 10 | | 8 | 9 | 11 | 10 |

Mapas de Karnaugh: Solapamientos



Opción a: $x' \cdot z + x \cdot z' + y' \cdot z'$

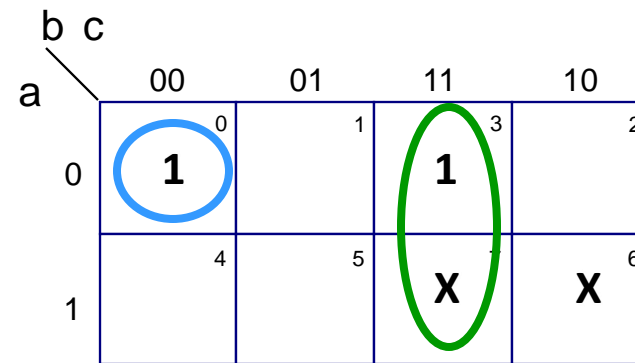
Opción b: $x' \cdot z + x \cdot z' + x' \cdot y'$



Por maxterms: $(x+y'+z) \cdot (x'+z')$

Mapas de Karnaugh: “No importa”

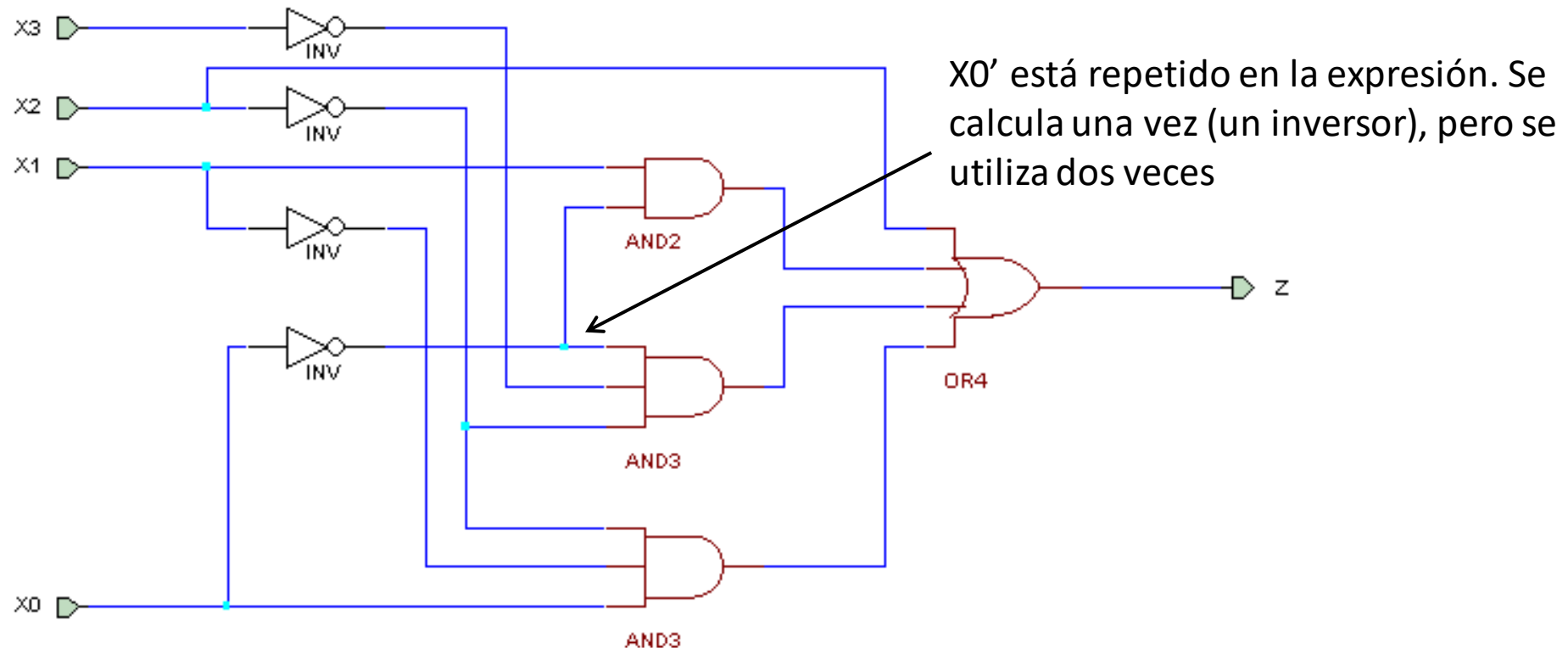
| a | b | c | out |
|---|---|---|-----|
| 0 | 0 | 0 | 1 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | X |
| 1 | 1 | 1 | X |



$$a' \cdot b' \cdot c' + b \cdot c$$

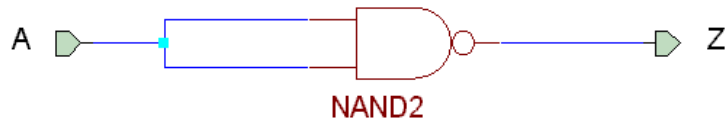
Implementación mediante Puertas

$$Z = X_2 + X_1 \cdot X_0' + X_3' \cdot X_2' \cdot X_0' + X_2' \cdot X_1' \cdot X_0$$

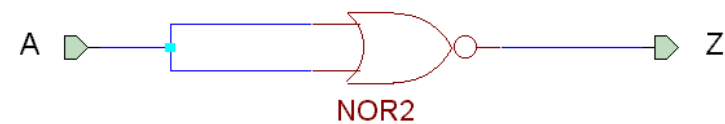


Puertas Universales NAND y NOR

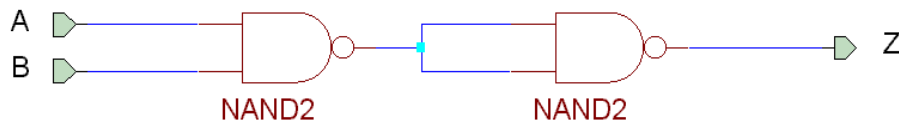
$$\text{NOT}(a) = \text{NOT}(a \text{ AND } a)$$



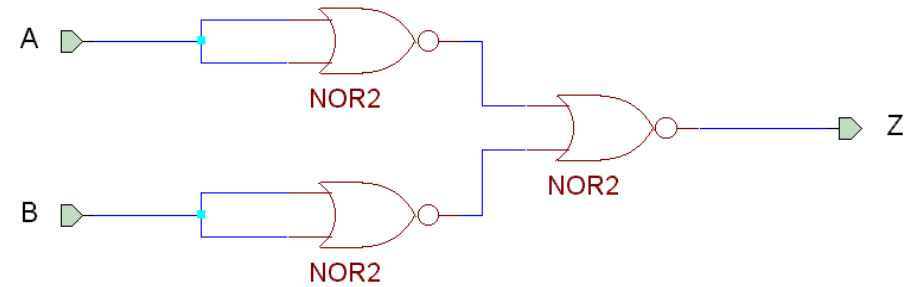
$$\text{NOT}(a) = \text{NOT}(a \text{ OR } a)$$



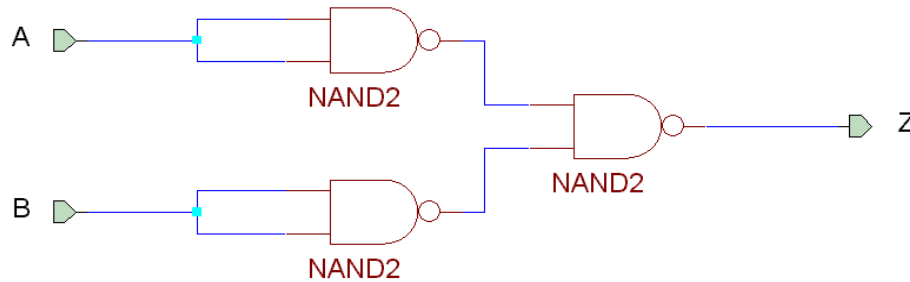
$$a \text{ AND } b = \text{NOT}(\text{NOT}(a \text{ AND } b))$$



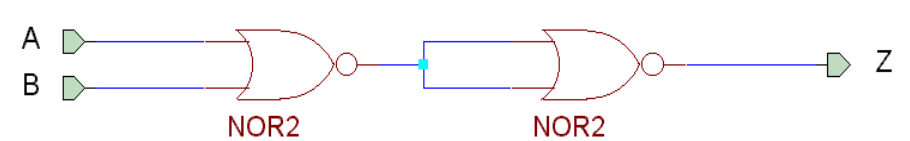
$$a \text{ AND } b = \text{NOT}(\text{NOT}(a) \text{ OR } \text{NOT}(b))$$



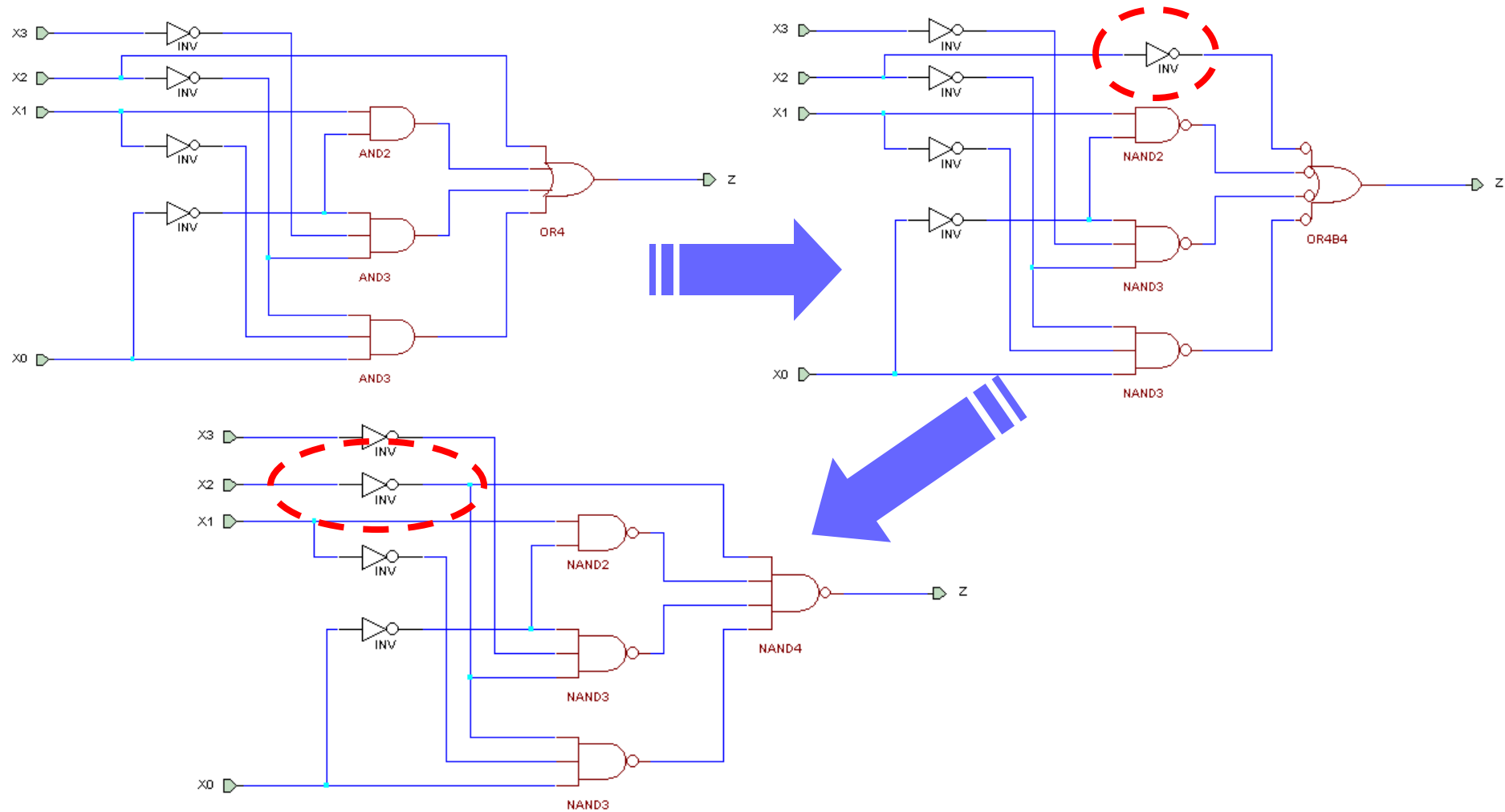
$$a \text{ OR } b = \text{NOT}(\text{NOT}(a) \text{ AND } \text{NOT}(b))$$



$$a \text{ OR } b = \text{NOT}(\text{NOT}(a \text{ OR } b))$$



Puertas Universales NAND



Puertas Universales NOR

