

# SPECTRUM of MONADS

$(T, \mu, \eta)$  on Set

but

Semantics via coalgebras

defn a comodel is

$$W \in \text{CoProd}(K(T), \text{Set}) =: \text{Comod}(T)$$

$$W \downarrow \quad \uparrow W \times -$$

$$(W \in \text{Set}, \rho: \int_X W \times TX \rightarrow W \times X)$$

write  $\rho(w, t) =: (t)(w)$   
"cointerpretation"

theorem  $\text{Comod}(T) \xrightleftharpoons{\perp} \text{Set}$

(Power, Shkaravskii 2004)

so instructive to look at terminal coalgebra  $B_T$ .

defn (admissible behaviour)

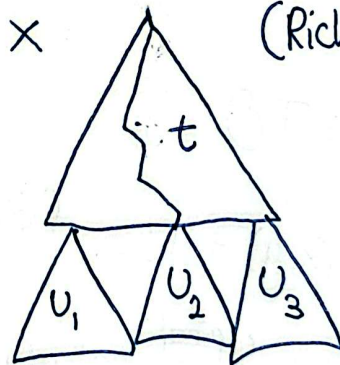
$$\beta: \int_X TX \rightarrow X$$

for  $t \in TX, u: X \rightarrow TY$

$$\beta(t \gg u) = \beta(t \gg u \beta(t))$$

$$t(\lambda x. u_x)$$

$$t(\lambda x. u_{\beta(t)})$$



(Richard, 2022)

Naturality  $\iff \beta(t \gg \underbrace{\text{return } x}_{\eta(x)}) = x$

Comodel

$$B_T \times TX \rightarrow B_T \times X$$

$$(\beta, t) \mapsto (\beta(t \gg -), \beta(t))$$

Terminal

$$W \rightarrow B_T$$

$$w \mapsto \beta_w: t \mapsto (t)(w)$$

# examples

$\beta$  determined by

$$\textcircled{1} T_{\text{inp}} X := \{ \text{binary trees with leaves in } X \} \rightsquigarrow \beta(\underbrace{b \gg b \gg \dots \gg b}_{n\text{-times}})$$

$\iff$  free theory on  $b(x,y)$

$$B_{\text{inp}} \cong 2^{\omega}$$

$$\textcircled{2} T_{\text{out}} X = 2^* \times X$$

every term of form  $t \gg \text{return } x$

$$\iff \text{free theory on } U_0(x), U_1(x) \rightsquigarrow \text{so } \beta = \int_X 2^* \times X \xrightarrow{\pi_x} x$$

unique

$$\textcircled{3} T_{\text{state}} X = (S \times X)^S$$

$$s \in S \mapsto (\beta_s : t \mapsto t(s))$$

$$\beta(\Delta_s) \mapsto \beta \in B_{\text{state}}$$

$\iff$  generated by  $\text{get} \in TS, \forall s \in S. \text{put}_s \in T1.$

$$B_{\text{state}} \cong S$$

$$\textcircled{4} \exists f \in T0$$

$$\rightsquigarrow \beta(f) \in 0 \not\Leftarrow \Rightarrow B \cong \emptyset$$

$$\exists f(x,y) \in T2. f(x,y) = f(y,x)$$

$$\rightsquigarrow \beta(f(x,y)) = x \Rightarrow y = \beta(f(y,x)) = x \not\Leftarrow$$

$$\text{so } B \cong \emptyset$$

(RIP Most <sup>typical</sup> alg. structures)

$$\textcircled{5} \text{ For } B \in BA,$$

$T_B$  generated by  $\forall b \in B. b \in T_B 2.$

(satisfying some equations)

$$\rightsquigarrow \beta \text{ determined by } \beta|_{T_2}$$

ultrafilter

$$B_B \cong |\text{Spec}(B)|$$

$$T_B X \cong \{ d : X \rightarrow B \mid \text{supp}(d) \text{ finite, } d[\text{supp}(d)] \text{ partitions } B \}$$

(Richard 2024)



Example (5) suggests  $B_T = |\text{Spec}(T)|$ . Here is another indicator.

defn (Behaviour Category)

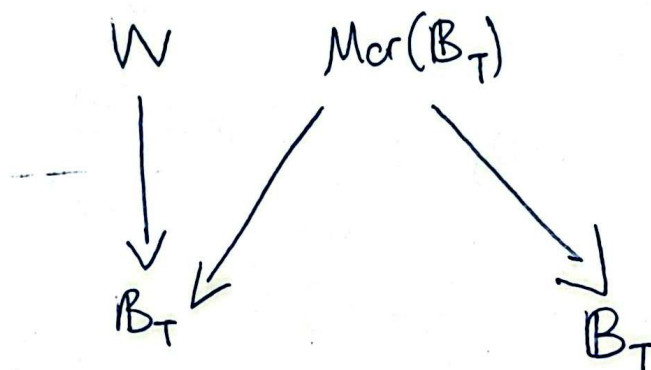
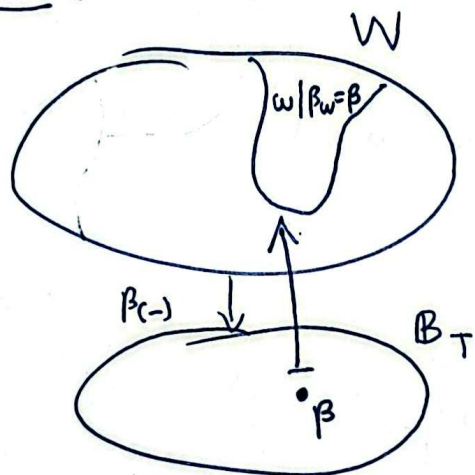
$$\text{Mor}(B_T) = \sum_{\beta \in B_T} T1 / \sim_{\beta}$$

$$t_1 \sim_{\beta} t_2 \iff \triangleleft_{t_1} = \triangleleft_{t_2}$$

i.e. generated by  $t \gg u \sim_{\beta} t \gg u(\beta(t))$

$$\text{dom}([t]_{\beta}) := \beta \quad \text{cod}([t]_{\beta}) := \beta(t \gg -)$$

Theorem (Richard 2022)  $\text{Comod}(T) \cong B_T\text{-Set} \simeq \text{Psh}(B_T)$



But  $\text{Comod}(T) = \text{Coproj}(K2(T), \text{Set}) \xrightarrow{I} [K2(T), \text{Set}] = R\text{Mod}(T)$

so calls to mind / suggests this is very restricted form of:

Serre's theorem Let  $R$  Noetherian CRing,

$$\{\text{Fin-Gen, projective } R\text{-modules}\} \simeq \left\{ \begin{array}{l} \text{locally free sheaves of} \\ \text{structure-sheaf modules} \\ \text{of constant finite rank} \\ \text{on } \text{Spec}(R) \end{array} \right\}$$

I haven't given <sup>the</sup> topology on  $B_T$  yet, but this suggests  $W$  as a presheaf gives the stalks, also  $\text{Mor}(B_T)$ .

In what sense is  $B_T$  a sheaf? First, there is a natural topology given by subbasic open sets

$$[t \mapsto x] = \{\beta \mid \beta(t) = x\} \quad \forall t \in TX, x \in X$$

The structure sheaf  $F_T$  at first approx, should be a sheaf of monads ('just as the ss for a ring is a sheaf of rings) s.t.

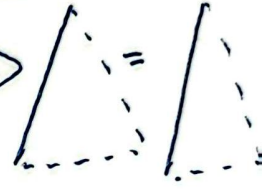
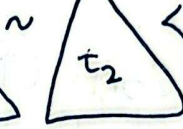
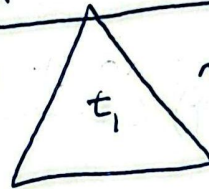
$$\text{im} (B_{F_T} [t \mapsto x] \longrightarrow B_T) = [t \mapsto x]$$

By Duality, we want a quotient  $T \twoheadrightarrow F_T [t \mapsto x]$ , try

$$t \sim t \gg \text{return } x$$

example  $T_{\text{inp}} / b(x,y) \sim b(x,x)$

$$\Rightarrow$$



$$\Leftrightarrow$$

$$\text{So } \text{Spec}(\text{---}) \longrightarrow \approx \{0000\dots\} \quad \parallel$$

problem; we allowed congruence, but the condition  $[b(x,y) \mapsto x]$  only applies to the first input, so cannot be congruent.

Solution; only allow pre congruence. Then  $T/\sim$  right module

$$\mu/\sim: T/\sim T \longrightarrow T/\sim$$

$$\text{So } F_T \cap [t_i \mapsto x_i] = T / \bigvee_i \sim [t_i \mapsto x_i]$$

$\text{Mor}(B_T)$  is the total space of germs of

$$\Rightarrow \text{Stalk over } \beta: T / \bigvee_t \sim [t \mapsto \beta(t)]$$

$$F_T(-)(1)$$

but

$$\bigvee_t \sim [t \mapsto \beta(t)] = \sim_\beta$$



What is <sup>the</sup> right module over  $\beta$ ?

$$TX / \sim_\beta \cong \underbrace{TI / \sim_\beta}_\text{Comod} \times X$$

So Comod $\ell$ s = "local right modules", and story should apply more generally to right modules.

Diers' spectrum for a multi-adjunction

$$I : \text{Comod}(T) \hookrightarrow \text{RMod}(T) := [k^?(T), \text{set}]$$

"  $\text{Coproduct}(k^?(T), \text{set})$

preserves connected limits, so has a left multi-adjoint

$$\sum_{\alpha \in \text{Hom}(M, B_T \times -)} \text{Hom}_{\text{Comod}}(M_\alpha, W \times -) \cong \text{Hom}_{\text{RMod}}(M, W \times -)$$

sanity check  $\text{Spec}(T) = \text{Hom}(T, B_T \times -) \stackrel{\text{spec}(M)}{\cong} B_T$

Okay, (1) Is  $\text{Spec}(M)$  the terminal something?

(2) What is  $M_\alpha$ ?

Another way to view comodels:

$$[T, \text{State}(W)] \cong \int^X TX \rightarrow (W \times X)^W$$

If  $I$  another set, have right module of  $\text{State}(W)$

$$\text{State}(I, W) = (W \times -)^I$$

defn An  $M$ -comodel relative to  $W$  is a set  $I$  with

$$M \rightarrow \text{State}(I, W)$$

s.t.

$$MT \rightarrow \text{State}(I, W) \text{State}(W)$$

$$\begin{array}{ccc} \downarrow & & \downarrow \\ M & \longrightarrow & \text{State}(I, W) \end{array}$$

a map  $(W, I) \rightarrow (W', I')$  is a pair of functions s.t.

$$\begin{array}{ccc} M & \longrightarrow & \text{State}(I, W) \\ \downarrow & & \downarrow \\ \text{State}(I', W') & \longrightarrow & \text{State}(I, W') \end{array}$$

prop  $(\text{Spec}(T), \text{Spec}(M))$  is the terminal  $M$ -comodel.

$$I \longrightarrow \text{Spec}(M)$$

That answers (1).

$$i \longmapsto \alpha_i; m \longmapsto$$

For ②: Category  $B_T$  should be the structure sheaf,  
and  $T_\beta$  should be the stalks of the structure sheaf  
(as a sheaf acted on by itself)

defn (operational topology) On  $\text{Obj}(B_T)$ , given by (Richard 2023)

$$\forall t \in TX, x \in X. [t \mapsto x] := \{ \beta \in B_T \mid \beta(t) = x \}$$

$\mathcal{O}_n \text{ Mor}(B_T)$ , given by

$$\forall m \in T1, t \in TX, x \in X. [m \mid t \mapsto x] := \{ [m]_\beta \mid \beta(t) = x \}$$

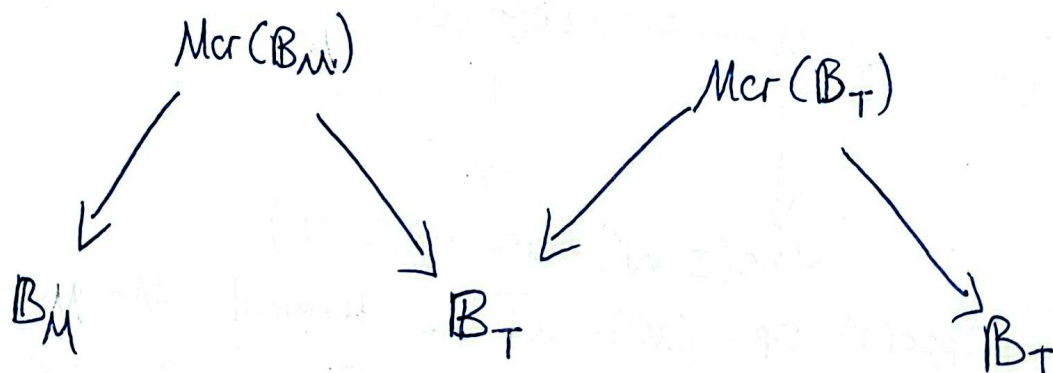
$\hookrightarrow$  dom is a local homeomorphism, hence a sheaf.

So  $T_\beta = T1 / \sim_\beta$ . Generalizing,

$$M_\alpha = M1 / \sim_\alpha \text{ where } m \gg u \sim_\alpha m \gg u(\alpha(m),)$$

We can put operational topology on  $\sum_{\alpha \in \text{Spec}(M)} M_\alpha$  and  $\text{Spec}(M)$ ,  
get "Serre" Theorem:

$$\begin{array}{ccc} \sum_{\alpha \in \text{Spec}(M)} M_\alpha & & \text{Spec}(M) \\ \parallel & & \parallel \\ \text{Mor}(BM) & & BM \end{array}$$



If  $M$  comodel,  $\text{Spec}(M)$  singleton  $\{\alpha\}$  and  $\text{Mor}(B_M) = M1 / \sim_\alpha = M1 =: W$   
to recover the previous picture (is this true for any bimodule?)

If  $T$  the monad, this post hoc explains why  $B_T$  is a category.



# Examples

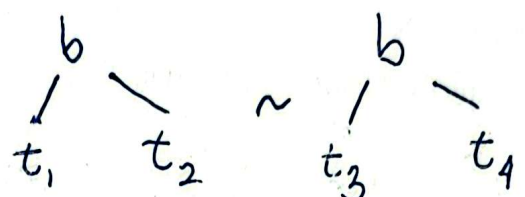
We think of right modules in terms of generators and relations

Defn  $\sim^* \subseteq MX * MX$  a pre-congruence if | Given a set of generators, the free right module consists of

$$\frac{m_1 \sim^* m_2 \quad t: X \rightarrow TY}{m_1 \gg t \sim^* m_2 \gg t}$$

$MX = \{(m, t) \mid t: \alpha(m) \rightarrow TX\}$   
or for an endofunctor  $F, FT$ .

Let  $\sim_{[m \mapsto x]}$  generated by  $m \sim m \gg \text{return } x$

•  $T_{\text{inp}} / \sim_{[b(x,y) \mapsto y]}$  

$$T_1 := T_{\text{inp}} X / \sim = \left\{ \text{leaf}(x) \forall x \in X, \bullet \begin{array}{c} b \\ / \quad \backslash \\ t \end{array} \forall t \in TX \right\} \cong X + TX$$

with right action

$$\text{leaf}(\text{leaf}(x)) \mapsto \text{leaf}(x)$$

$$\text{leaf} \left( \begin{array}{c} b \\ / \quad \backslash \\ t_1 \quad t_2 \end{array} \right) \mapsto \bullet \begin{array}{c} b \\ / \quad \backslash \\ \cdot \quad t_2 \end{array}$$

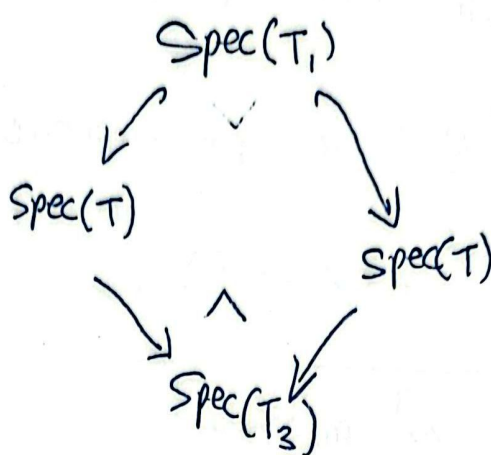
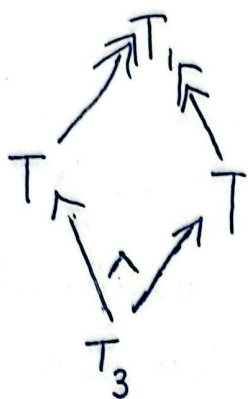
$$\bullet \begin{array}{c} b \\ / \quad \backslash \\ \cdot \quad t \in TX \end{array} \mapsto \bullet \begin{array}{c} b \\ / \quad \backslash \\ \cdot \quad \mu(t) \end{array}$$

$$\text{Spec}(T_{\text{inp}} / \sim) \hookrightarrow \text{Spec}(T)$$

$$\text{Spec}(T_{\text{inp}} / \sim) = \{ \beta \mid \beta(b(x,y)) = y \} = \{ \beta \in 2^\omega \mid \text{head}(\beta) = 1 \} \cong 2^\omega$$



- We can glue two copies of  $T$  along  $T_1$ :



$$T_3 X = \left\{ \underbrace{(\text{return } x, \text{return } x)}_{\text{return } x} \forall x, \underbrace{(b(t_1, t), b(t_2, t))}_{bb(t_1, t_2, t)} \forall t_1, t_2, t \in TX \right\}$$

so  $T_3$  has generators  $\text{return} \in T_3 1$  and  $bb \in T_3 3$ , satisfying

$$(\text{return}, b(x, y)) \sim (bb, (x, x, y))$$

$$\text{Spec}(T_3) = 2^w + 2^w / \sim \text{ where } \text{inl}(\beta) \sim \text{inr}(\beta) \iff \text{head}(\beta) = 1.$$

$$\cong \left\{ \text{inl}(\beta), \text{inr}(\beta) \mid \text{head}(\beta) = 0 \right\} + \left\{ \beta \mid \text{head}(\beta) = 1 \right\}$$

$$\cong 2^w + 2^w + 2^w$$

$$\cong 3 \times 2^w$$

$$(\text{return}, b(x, y))(\alpha) = \begin{cases} x & \text{head } \alpha = 0 \text{ or } 1 \\ y & \text{head } \alpha = 2 \end{cases}$$