

Natural transformations

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1 Basics

Definition 1.1 (Natural transformations). Let $F, G : \mathcal{C} \rightarrow \mathcal{D}$. Natural transformation α between F and G is a collection of morphisms $(\alpha_c \mid c : \text{obj } \mathcal{C})$ satisfying the diagrams

$$\begin{array}{ccc} c & & Fc \xrightarrow{\alpha_c} Gc \\ f \downarrow & & \downarrow Gf \\ c' & & Fc' \xrightarrow{\alpha_{c'}} Gc' \end{array}$$

Notation:

$$F \xrightarrow{\alpha} G \quad \mathcal{C} \xrightarrow[\mathcal{G}]{}^F \mathcal{D} .$$

Definition 1.2 (Vertical composition of natural transformations). Let $F, G, H : \mathcal{C} \rightarrow \mathcal{D}$ with natural transformations $\alpha : F \rightarrow G$ and $\beta : G \rightarrow H$. The vertical composition $\beta \circ \alpha$ of natural transformations β and α is defined componentwise:

$$\beta \circ \alpha : F \rightarrow H \quad (\beta \circ \alpha)_c = \beta_c \circ \alpha_c,$$

$$\begin{array}{ccccc} c & & Fc \xrightarrow{\alpha_c} Gc \xrightarrow{\beta_c} Hc \\ f \downarrow & & \downarrow Gf & & \downarrow Hf \\ c' & & Fc' \xrightarrow{\alpha_{c'}} Gc' \xrightarrow{\beta_{c'}} Hc' & & \end{array}$$

Notation:

$$F \xrightarrow{\alpha} G \quad \mathcal{C} \xrightarrow[\mathcal{H}]{}^F \mathcal{D} .$$

Definition 1.3 (Horizontal composition of natural transformations). Let $F_1, F_2 : \mathcal{C} \rightarrow \mathcal{D}$, $G_1, G_2 : \mathcal{D} \rightarrow \mathcal{E}$ with natural transformations $\alpha : F_1 \rightarrow F_2$ and $\beta : G_1 \rightarrow G_2$. The horizontal composition $\beta \bullet \alpha$ of natural transformations β and α is defined as the transformation of images of α_c under the action of F_1 and F_2 by β :

$$\beta \bullet \alpha : G_1 F_1 \rightarrow G_2 F_2 \quad (\beta \bullet \alpha)_c = G_2 \alpha_c \circ \beta_{F_1 c} = \beta_{F_2 c} \circ G_1 \alpha_c,$$

$$\begin{array}{ccc}
& & G_1 F_1 c' \xrightarrow{\quad G_1 \alpha_{c'} \quad} G_1 F_2 c' \\
& \nearrow G_1 F_1 f & \downarrow \beta_{F_1 c'} \\
G_1 F_1 c & \xrightarrow{\quad G_1 \alpha_c \quad} & G_1 F_2 c \\
& \downarrow \beta_{F_1 c} & \downarrow \beta_{F_1 c'} \\
& G_2 F_1 c' \xrightarrow{\quad G_2 \alpha_{c'} \quad} G_2 F_2 c' \\
& \searrow G_2 F_1 f & \downarrow \beta_{F_2 c} \\
G_2 F_1 c & \xrightarrow{\quad G_2 \alpha_c \quad} & G_2 F_2 c
\end{array}$$

$\begin{array}{ccc} c & & c' \\ & \nearrow f & \nearrow \\ & F_1 c' & \xrightarrow{\alpha_{c'}} F_2 c' \\ & \swarrow F_1 f & \nearrow \\ F_1 c & \xrightarrow{\alpha_c} & F_2 c \\ & \swarrow & \nearrow F_2 f \end{array}$

Note that functors preserve commutative diagrams, so all of the faces of the cube commute.
Notation:

$$\begin{array}{ccc}
\mathcal{C} \xrightarrow[F_2]{\Downarrow \alpha} \mathcal{D} \xrightarrow[G_2]{\Downarrow \beta} \mathcal{E} & & G_1 F_1 \xrightarrow{\beta \bullet \alpha} G_2 F_2
\end{array}$$

Definition 1.4 (Whiskering). Whiskering is defined as a horizontal composition with identity natural transformation:

$$\begin{array}{ccc}
\mathcal{C} \xrightarrow[F_2]{\Downarrow \alpha} \mathcal{D} \xrightarrow{G} \mathcal{E} & & \mathcal{C} \xrightarrow{F} \mathcal{D} \xrightarrow[G_2]{\Downarrow \beta} \mathcal{E} ; \\
G\alpha := 1_G \bullet \alpha : G \circ F_1 \rightarrow G \circ F_2 & & (G\alpha)_c = G\alpha_c, \\
\beta F := \beta \bullet 1_F : G_1 \circ F \rightarrow G_2 \circ F & & (\beta F)_c = \beta_{Fc}.
\end{array}$$

Using this notation, we can present any horizontal transformation in the following form:

$$\beta \bullet \alpha = \beta F_2 \circ G_1 \alpha = G_2 \alpha \circ \beta F_1.$$

Proposition 1.5. Let $F_1, F_2, F_3 : \mathcal{C} \rightarrow \mathcal{D}$, $G_1, G_2, G_3 : \mathcal{D} \rightarrow \mathcal{E}$ with natural transformations $\alpha : F_1 \rightarrow F_2$, $\alpha' : F_2 \rightarrow F_3$ and $\beta : G_1 \rightarrow G_2$, $\beta' : G_2 \rightarrow G_3$:

$$\mathcal{C} - \xrightarrow[F_2]{\Downarrow \alpha} \mathcal{D} - \xrightarrow[G_2]{\Downarrow \beta} \mathcal{E} .$$

Then there is an interchange law between the natural transformations:

$$(\beta' \circ \beta) \bullet (\alpha' \circ \alpha) = (\beta' \bullet \alpha') \circ (\beta \bullet \alpha).$$

Proof.

$$(\beta' \circ \beta) \bullet (\alpha' \circ \alpha) = \beta' F_3 \circ (\beta F_3 \circ G_1 \alpha') \circ G_1 \alpha = \beta' F_3 \circ G_2 \alpha' \circ \beta F_2 \circ G_1 \alpha,$$

$$(\beta' \bullet \alpha') \circ (\beta \bullet \alpha) = \beta' F_3 \circ G_2 \alpha' \circ \beta F_2 \circ G_1 \alpha.$$

2 Advanced whiskering

The naturality diagram from Def 1.1. basically means, that for given $f : \text{arr}(\mathcal{C})$

$$\alpha_{\text{cod}(f)} \circ \text{dom}(\alpha)f = \text{cod}(\alpha)f \circ \alpha_{\text{dom}(\alpha)}.$$

The question arises: can we write it in the form of

$$\alpha \text{ dom}(\alpha) = \text{cod}(\alpha)\alpha$$

using some polymorphic notation? We already have the notation for whiskering, but so far we only defined it on $\text{obj}(\mathcal{C})$.