

# Natural transformations

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## 1 Basics

**Definition 1.1** (Natural transformations). Let  $F, G : \mathcal{C} \rightarrow \mathcal{D}$ . A natural transformation  $\alpha$  between  $F$  and  $G$  is a collection of morphisms  $(\alpha_c \mid c \in \text{obj}(\mathcal{C}))$  satisfying the diagrams

$$\begin{array}{ccc} c & & Fc \xrightarrow{\alpha_c} Gc \\ f \downarrow & & Ff \downarrow \quad \downarrow Gf \\ c' & & Fc' \xrightarrow{\alpha_{c'}} Gc' \end{array}$$

Notation:

$$F \xrightarrow{\alpha} G \quad \mathcal{C} \xrightarrow[\mathcal{G}]{}^F \mathcal{D} .$$

**Definition 1.2** (Vertical composition of natural transformations). Let  $F, G, H : \mathcal{C} \rightarrow \mathcal{D}$  with natural transformations  $\alpha : F \rightarrow G$  and  $\beta : G \rightarrow H$ . The vertical composition  $\beta \circ \alpha$  is defined componentwise:

$$\beta \circ \alpha : F \rightarrow H \quad (\beta \circ \alpha)_c = \beta_c \circ \alpha_c,$$

$$\begin{array}{ccccc} c & & Fc \xrightarrow{\alpha_c} Gc \xrightarrow{\beta_c} Hc \\ f \downarrow & & Ff \downarrow \quad Gf \downarrow \quad \downarrow Hf \\ c' & & Fc' \xrightarrow{\alpha_{c'}} Gc' \xrightarrow{\beta_{c'}} Hc' \end{array}$$

Notation:

$$F \xrightarrow{\alpha} G \quad \mathcal{C} \xrightarrow[\mathcal{H}]{}^F \mathcal{G} \xrightarrow[\mathcal{G}]{}^{\beta} \mathcal{D} .$$

**Definition 1.3** (Horizontal composition of natural transformations). Let  $F_1, F_2 : \mathcal{C} \rightarrow \mathcal{D}$ ,  $G_1, G_2 : \mathcal{D} \rightarrow \mathcal{E}$  with natural transformations  $\alpha : F_1 \rightarrow F_2$  and  $\beta : G_1 \rightarrow G_2$ . The horizontal composition  $\beta \bullet \alpha$  is defined as the transformation obtained by applying  $\beta$  to the images of  $\alpha_c$  under  $F_1$  and  $F_2$ :

$$\beta \bullet \alpha : G_1 F_1 \rightarrow G_2 F_2 \quad (\beta \bullet \alpha)_c = G_2 \alpha_c \circ \beta_{F_1 c} = \beta_{F_2 c} \circ G_1 \alpha_c,$$

$$\begin{array}{ccc}
& & G_1 F_1 c' \xrightarrow{\quad G_1 \alpha_{c'} \quad} G_1 F_2 c' \\
& \nearrow G_1 F_1 f & \downarrow \beta_{F_1 c'} \\
G_1 F_1 c & \xrightarrow{\quad G_1 \alpha_c \quad} & G_1 F_2 c \\
& \downarrow \beta_{F_1 c} & \downarrow \beta_{F_1 c'} \\
& G_2 F_1 c' \xrightarrow{\quad G_2 \alpha_{c'} \quad} G_2 F_2 c' \\
& \searrow G_2 F_1 f & \downarrow \beta_{F_2 c} \\
G_2 F_1 c & \xrightarrow{\quad G_2 \alpha_c \quad} & G_2 F_2 c
\end{array}$$

$\begin{array}{ccc} c & & c' \\ & \nearrow f & \nearrow \\ & F_1 c' & \xrightarrow{\alpha_{c'}} F_2 c' \\ & \swarrow F_1 f & \nearrow \\ F_1 c & \xrightarrow{\alpha_c} & F_2 c \\ & \swarrow & \nearrow F_2 f \end{array}$

Note that functors preserve commutative diagrams, so all of the faces of the cube commute.  
Notation:

$$\begin{array}{ccc}
\mathcal{C} & \begin{array}{c} \xrightarrow{F_1} \\ \Downarrow \alpha \\ \xrightarrow{F_2} \end{array} & \mathcal{D} \xrightarrow{G_1} \mathcal{E} \\
& \nearrow & \downarrow \beta \\
& G_2 & \mathcal{E}
\end{array}
\qquad
G_1 F_1 \xrightarrow{\beta \bullet \alpha} G_2 F_2 .$$

**Definition 1.4** (Whiskering). Whiskering is defined as a horizontal composition with an identity natural transformation:

$$\begin{array}{ccc}
\mathcal{C} & \begin{array}{c} \xrightarrow{F_1} \\ \Downarrow \alpha \\ \xrightarrow{F_2} \end{array} & \mathcal{D} \xrightarrow{G} \mathcal{E} \\
& \nearrow & \downarrow \beta \\
& G_2 & \mathcal{E}
\end{array}
\qquad
\mathcal{C} \xrightarrow{F} \mathcal{D} \xrightarrow{G_1} \mathcal{E} ;$$

$$\begin{array}{ll}
G\alpha := 1_G \bullet \alpha : G \circ F_1 \rightarrow G \circ F_2 & (G\alpha)_c = G\alpha_c, \\
\beta F := \beta \bullet 1_F : G_1 \circ F \rightarrow G_2 \circ F & (\beta F)_c = \beta_{Fc}.
\end{array}$$

Using this notation, we can present any horizontal transformation in the following form:

$$\beta \bullet \alpha = \beta F_2 \circ G_1 \alpha = G_2 \alpha \circ \beta F_1.$$

**Proposition 1.5.** Let  $F_1, F_2, F_3 : \mathcal{C} \rightarrow \mathcal{D}$ ,  $G_1, G_2, G_3 : \mathcal{D} \rightarrow \mathcal{E}$  with natural transformations  $\alpha : F_1 \rightarrow F_2$ ,  $\alpha' : F_2 \rightarrow F_3$  and  $\beta : G_1 \rightarrow G_2$ ,  $\beta' : G_2 \rightarrow G_3$ :

$$\begin{array}{ccc}
\mathcal{C} & \begin{array}{c} \xrightarrow{F_1} \\ \Downarrow \alpha \\ \xrightarrow{F_2} \\ \Downarrow \alpha' \\ \xrightarrow{F_3} \end{array} & \mathcal{D} \xrightarrow{G_1} \mathcal{E} \\
& \nearrow & \downarrow \beta \\
& G_2 & \mathcal{E} \\
& \Downarrow \beta' & \nearrow
\end{array}$$

Then the following interchange law holds:

$$(\beta' \circ \beta) \bullet (\alpha' \circ \alpha) = (\beta' \bullet \alpha') \circ (\beta \bullet \alpha).$$

**Proof.**

$$(\beta' \circ \beta) \bullet (\alpha' \circ \alpha) = \beta' F_3 \circ (\beta F_3 \circ G_1 \alpha') \circ G_1 \alpha = \beta' F_3 \circ G_2 \alpha' \circ \beta F_2 \circ G_1 \alpha,$$

$$(\beta' \bullet \alpha') \circ (\beta \bullet \alpha) = \beta' F_3 \circ G_2 \alpha' \circ \beta F_2 \circ G_1 \alpha.$$