

# Natural transformations

Alyson Mei

December 25, 2025

## 1 Basics

**Definition 1.1** (Natural transformations). Let  $F, G : \mathcal{C} \rightarrow \mathcal{D}$ . Natural transformation  $\alpha$  between  $F$  and  $G$  is a collection of morphisms  $(\alpha_c \mid c : \text{obj } \mathcal{C})$  satisfying the diagrams

$$\begin{array}{ccc} c & & Fc \xrightarrow{\alpha_c} Gc \\ f \downarrow & & \downarrow Ff \quad \downarrow Gf \\ c' & & Fc' \xrightarrow{\alpha_{c'}} Gc' \end{array}$$

Notation:

$$F \xrightarrow{\alpha} G \qquad \mathcal{C} \begin{array}{c} \xrightarrow{F} \\ \Downarrow \alpha \\ \xrightarrow{G} \end{array} \mathcal{D}.$$

**Definition 1.2** (Vertical composition of natural transformations). Let  $F, G, H : \mathcal{C} \rightarrow \mathcal{D}$  with natural transformations  $\alpha : F \rightarrow G$  and  $\beta : G \rightarrow H$ . The vertical composition  $\beta \circ \alpha$  of natural transformations  $\beta$  and  $\alpha$  is defined componentwise:

$$\beta \circ \alpha : F \rightarrow H \qquad (\beta \circ \alpha)_c = \beta_c \circ \alpha_c,$$

$$\begin{array}{ccccc} c & & Fc & \xrightarrow{\alpha_c} & Gc & \xrightarrow{\beta_c} & Hc \\ f \downarrow & & \downarrow Ff & & \downarrow Gf & & \downarrow Hf \\ c' & & Fc' & \xrightarrow{\alpha_{c'}} & Gc' & \xrightarrow{\beta_{c'}} & Hc' \end{array}$$

Notation:

$$\begin{array}{ccc} F & \xrightarrow{\alpha} & G \\ & \searrow \beta \circ \alpha & \downarrow \beta \\ & & H \end{array} \qquad \mathcal{C} \begin{array}{c} \xrightarrow{F} \\ \Downarrow \alpha \\ \xrightarrow{G} \\ \Downarrow \beta \\ \xrightarrow{H} \end{array} \mathcal{D}.$$

**Definition 1.3** (Horizontal composition of natural transformations). Let  $F_1, F_2 : \mathcal{C} \rightarrow \mathcal{D}$ ,  $G_1, G_2 : \mathcal{D} \rightarrow \mathcal{E}$  with natural transformations  $\alpha : F_1 \rightarrow F_2$  and  $\beta : G_1 \rightarrow G_2$ . The horizontal composition  $\beta \bullet \alpha$  of natural transformations  $\beta$  and  $\alpha$  is defined as the transformation of images of  $\alpha_c$  under the action of  $F_1$  and  $F_2$  by  $\beta$ :

$$\beta \bullet \alpha : G_1 F_1 \rightarrow G_2 F_2 \qquad (\beta \bullet \alpha)_c = G_2 \alpha_c \circ \beta_{F_1 c} = \beta_{F_2 c} \circ G_1 \alpha_c,$$

$$\begin{array}{c}
\begin{array}{ccc}
& & G_1 F_1 c' \xrightarrow{\quad G_1 \alpha_{c'} \quad} G_1 F_2 c' \\
& \nearrow^{G_1 F_1 f} \quad \downarrow \beta_{F_1 c'} \quad \nearrow^{G_1 F_2 f} \\
G_1 F_1 c & \xrightarrow{\quad G_1 \alpha_c \quad} & G_1 F_2 c \\
\downarrow \beta_{F_1 c} & & \downarrow \beta_{F_2 c} \\
G_2 F_1 c & \xrightarrow{\quad G_2 \alpha_c \quad} & G_2 F_2 c
\end{array} \\
\begin{array}{ccc}
c & \xrightarrow{f} & c' \\
\uparrow F_1 f & & \uparrow F_2 f \\
F_1 c & \xrightarrow{\alpha_c} & F_2 c
\end{array}
\end{array}$$

Note that functors preserve commutative diagrams, so all of the faces of the cube commute. Notation:

$$\begin{array}{c}
\mathcal{C} \begin{array}{c} \xrightarrow{F_1} \\ \Downarrow \alpha \\ \xrightarrow{F_2} \end{array} \mathcal{D} \begin{array}{c} \xrightarrow{G_1} \\ \Downarrow \beta \\ \xrightarrow{G_2} \end{array} \mathcal{E}
\end{array}
\quad G_1 F_1 \xrightarrow{\beta \bullet \alpha} G_2 F_2 .$$

**Definition 1.4** (Whiskering). Whiskering is defined as a horizontal composition with identity natural transformation:

$$\begin{array}{c}
\mathcal{C} \begin{array}{c} \xrightarrow{F_1} \\ \Downarrow \alpha \\ \xrightarrow{F_2} \end{array} \mathcal{D} \xrightarrow{G} \mathcal{E}
\end{array}
\quad \mathcal{C} \xrightarrow{F} \mathcal{D} \begin{array}{c} \xrightarrow{G_1} \\ \Downarrow \beta \\ \xrightarrow{G_2} \end{array} \mathcal{E} ;$$

$$\begin{aligned}
G\alpha &:= 1_G \bullet \alpha : G \circ F_1 \rightarrow G \circ F_2 & (G\alpha)_c &= G\alpha_c, \\
\beta F &:= \beta \bullet 1_F : G_1 \circ F \rightarrow G_2 \circ F & (\beta F)_c &= \beta_{F_c}.
\end{aligned}$$

Using this notation, we can present any horizontal transformation in the following form:

$$\beta \bullet \alpha = \beta F_2 \circ G_1 \alpha = G_2 \alpha \circ \beta F_1.$$

**Proposition 1.5.** Let  $F_1, F_2, F_3 : \mathcal{C} \rightarrow \mathcal{D}$ ,  $G_1, G_2, G_3 : \mathcal{D} \rightarrow \mathcal{E}$  with natural transformations  $\alpha : F_1 \rightarrow F_2$ ,  $\alpha' : F_2 \rightarrow F_3$  and  $\beta : G_1 \rightarrow G_2$ ,  $\beta' : G_2 \rightarrow G_3$ :

$$\begin{array}{ccc}
\mathcal{C} & \begin{array}{c} \xrightarrow{F_1} \\ \Downarrow \alpha \\ \xrightarrow{F_2} \\ \Downarrow \alpha' \\ \xrightarrow{F_3} \end{array} & \mathcal{D} \begin{array}{c} \xrightarrow{G_1} \\ \Downarrow \beta \\ \xrightarrow{G_2} \\ \Downarrow \beta' \\ \xrightarrow{G_3} \end{array} & \mathcal{E} .
\end{array}$$

Then there is an interchange law between the natural transformations:

$$(\beta' \circ \beta) \bullet (\alpha' \circ \alpha) = (\beta' \bullet \alpha') \circ (\beta \bullet \alpha).$$

**Proof.**

$$(\beta' \circ \beta) \bullet (\alpha' \circ \alpha) = \beta' F_3 \circ (\beta F_3 \circ G_1 \alpha') \circ G_1 \alpha = \beta' F_3 \circ G_2 \alpha' \circ \beta F_2 \circ G_1 \alpha,$$

$$(\beta' \bullet \alpha') \circ (\beta \bullet \alpha) = \beta' F_3 \circ G_2 \alpha' \circ \beta F_2 \circ G_1 \alpha.$$

## 2 Advanced whiskering

The naturality diagram from Def 1.1. basically means, that for given  $f : \text{arr}(\mathcal{C})$

$$\alpha_{\text{cod}(f)} \circ \text{dom}(\alpha)f = \text{cod}(\alpha)f \circ \alpha_{\text{dom}(\alpha)}.$$

The question arises: can we write it in the form of

$$\alpha \text{ dom}(\alpha) = \text{cod}(\alpha) \alpha$$

using some polymorphic notation? We already have the notation for whiskering, but so far we only defined it on  $\text{obj}(\mathcal{C})$ .