

Natural transformations

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1 Basics

Definition 1.1 (Natural transformations). Let $F, G : \mathcal{C} \rightarrow \mathcal{D}$. A natural transformation α between F and G is a collection of morphisms $(\alpha_c \mid c \in \text{obj}(\mathcal{C}))$ satisfying the diagrams

$$\begin{array}{ccc} c & & Fc \xrightarrow{\alpha_c} Gc \\ f \downarrow & & Ff \downarrow \quad \quad \downarrow Gf \\ c' & & Fc' \xrightarrow{\alpha_{c'}} Gc' \end{array}$$

Notation:

$$F \xrightarrow{\alpha} G \qquad \mathcal{C} \begin{array}{c} \xrightarrow{F} \\ \Downarrow \alpha \\ \xrightarrow{G} \end{array} \mathcal{D}.$$

Definition 1.2 (Vertical composition of natural transformations). Let $F, G, H : \mathcal{C} \rightarrow \mathcal{D}$ with natural transformations $\alpha : F \rightarrow G$ and $\beta : G \rightarrow H$. The vertical composition $\beta \circ \alpha$ is defined componentwise:

$$\beta \circ \alpha : F \rightarrow H \qquad (\beta \circ \alpha)_c = \beta_c \circ \alpha_c,$$

$$\begin{array}{ccccc} c & & Fc & \xrightarrow{\alpha_c} & Gc & \xrightarrow{\beta_c} & Hc \\ f \downarrow & & Ff \downarrow & & \downarrow Gf & & \downarrow Hf \\ c' & & Fc' & \xrightarrow{\alpha_{c'}} & Gc' & \xrightarrow{\beta_{c'}} & Hc' \end{array}$$

Notation:

$$\begin{array}{ccc} F & \xrightarrow{\alpha} & G \\ & \searrow \beta \circ \alpha & \downarrow \beta \\ & & H \end{array} \qquad \mathcal{C} \begin{array}{c} \xrightarrow{F} \\ \Downarrow \alpha \\ \xrightarrow{G} \\ \Downarrow \beta \\ \xrightarrow{H} \end{array} \mathcal{D}.$$

Definition 1.3 (Horizontal composition of natural transformations). Let $F_1, F_2 : \mathcal{C} \rightarrow \mathcal{D}$, $G_1, G_2 : \mathcal{D} \rightarrow \mathcal{E}$ with natural transformations $\alpha : F_1 \rightarrow F_2$ and $\beta : G_1 \rightarrow G_2$. The horizontal composition $\beta \bullet \alpha$ is defined as the transformation obtained by applying β to the images of α_c under F_1 and F_2 :

$$\beta \bullet \alpha : G_1 F_1 \rightarrow G_2 F_2 \qquad (\beta \bullet \alpha)_c = G_2 \alpha_c \circ \beta_{F_1 c} = \beta_{F_2 c} \circ G_1 \alpha_c,$$

$$\begin{array}{c}
\begin{array}{ccc}
& & G_1 F_1 c' \xrightarrow{\quad} G_1 \alpha_{c'} \rightarrow G_1 F_2 c' \\
& \nearrow^{G_1 F_1 f} \quad \downarrow \beta_{F_1 c'} \quad \nearrow^{G_1 F_2 f} \\
G_1 F_1 c & \xrightarrow{\quad} & G_1 \alpha_c \rightarrow G_1 F_2 c \\
\downarrow \beta_{F_1 c} & & \downarrow \beta_{F_2 c} \\
G_2 F_1 c & \xrightarrow{\quad} & G_2 \alpha_c \rightarrow G_2 F_2 c \\
& \nearrow^{G_2 F_1 f} \quad \downarrow \beta_{F_1 c'} \quad \nearrow^{G_2 F_2 f} \\
& G_2 F_1 c' \xrightarrow{\quad} G_2 \alpha_{c'} \rightarrow G_2 F_2 c'
\end{array}
\end{array}$$

Note that functors preserve commutative diagrams, so all of the faces of the cube commute. Notation:

$$\mathcal{C} \begin{array}{c} \xrightarrow{F_1} \\ \Downarrow \alpha \\ \xrightarrow{F_2} \end{array} \mathcal{D} \begin{array}{c} \xrightarrow{G_1} \\ \Downarrow \beta \\ \xrightarrow{G_2} \end{array} \mathcal{E} \qquad G_1 F_1 \xrightarrow{\beta \bullet \alpha} G_2 F_2 .$$

Definition 1.4 (Whiskering). Whiskering is defined as a horizontal composition with an identity natural transformation:

$$\begin{array}{ccc}
\mathcal{C} \begin{array}{c} \xrightarrow{F_1} \\ \Downarrow \alpha \\ \xrightarrow{F_2} \end{array} \mathcal{D} & \xrightarrow{G} & \mathcal{E} \qquad \mathcal{C} \xrightarrow{F} \mathcal{D} \begin{array}{c} \xrightarrow{G_1} \\ \Downarrow \beta \\ \xrightarrow{G_2} \end{array} \mathcal{E} ; \\
G\alpha := 1_G \bullet \alpha : G \circ F_1 \rightarrow G \circ F_2 & & (G\alpha)_c = G\alpha_c, \\
\beta F := \beta \bullet 1_F : G_1 \circ F \rightarrow G_2 \circ F & & (\beta F)_c = \beta_{F_c}.
\end{array}$$

Using this notation, we can present any horizontal transformation in the following form:

$$\beta \bullet \alpha = \beta F_2 \circ G_1 \alpha = G_2 \alpha \circ \beta F_1.$$

Proposition 1.5. Let $F_1, F_2, F_3 : \mathcal{C} \rightarrow \mathcal{D}$, $G_1, G_2, G_3 : \mathcal{D} \rightarrow \mathcal{E}$ with natural transformations $\alpha : F_1 \rightarrow F_2$, $\alpha' : F_2 \rightarrow F_3$ and $\beta : G_1 \rightarrow G_2$, $\beta' : G_2 \rightarrow G_3$:

$$\mathcal{C} \begin{array}{c} \xrightarrow{F_1} \\ \Downarrow \alpha \\ \xrightarrow{F_2} \\ \Downarrow \alpha' \\ \xrightarrow{F_3} \end{array} \mathcal{D} \begin{array}{c} \xrightarrow{G_1} \\ \Downarrow \beta \\ \xrightarrow{G_2} \\ \Downarrow \beta' \\ \xrightarrow{G_3} \end{array} \mathcal{E} .$$

Then the following interchange law holds:

$$(\beta' \circ \beta) \bullet (\alpha' \circ \alpha) = (\beta' \bullet \alpha') \circ (\beta \bullet \alpha).$$

Proof.

$$\begin{aligned}
(\beta' \circ \beta) \bullet (\alpha' \circ \alpha) &= \beta' F_3 \circ (\beta F_3 \circ G_1 \alpha') \circ G_1 \alpha = \beta' F_3 \circ G_2 \alpha' \circ \beta F_2 \circ G_1 \alpha, \\
(\beta' \bullet \alpha') \circ (\beta \bullet \alpha) &= \beta' F_3 \circ G_2 \alpha' \circ \beta F_2 \circ G_1 \alpha.
\end{aligned}$$