

# CS208: Applied Privacy for Data Science Membership Attacks

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### Motivation

• Last time: on a dataset with n individuals, releasing m=n counts with error  $E=o(\sqrt{n})$  allows for reconstructing 1-o(1) fraction of sensitive attributes. [Dinur-Nissim `03]

# What is this $\sqrt{n}$ threshold?

• if  $X = X_1 + \cdots + X_n$  for independent random variables  $X_i$  each with standard deviation  $\sigma$ , then the standard deviation of X is  $\Theta(\sqrt{n})$ .

• If the  $X_i$ 's are bounded (or "subgaussian"), then X will have Gaussian-like concentration around its expectation  $n\mu$ :

$$\Pr[|X - n\mu| > t \cdot \sqrt{n}] \le e^{-\Omega(t^2)}$$
 [Chernoff-Hoeffding Bound]

This is why subsampling k out of n rows allows us to approximate m counts each to within  $\pm O(\sqrt{k \log m})$ 

# Normalized Counts (i.e. Averages)

- if  $X = (X_1 + \cdots + X_n)/n$  for independent random variables  $X_i$  each with standard deviation  $\sigma$ , then the standard deviation of X is  $\Theta(1/\sqrt{n})$
- If the  $X_i$ 's are bounded (or "subgaussian"), then X will have Gaussian-like concentration around its mean  $\mu$ :

$$\Pr[|X - \mu| > t/\sqrt{n}] \le e^{-\Omega(t^2)}$$
 [Chernoff-Hoeffding Bound]

This is why subsampling k out of n rows allows us to approximate

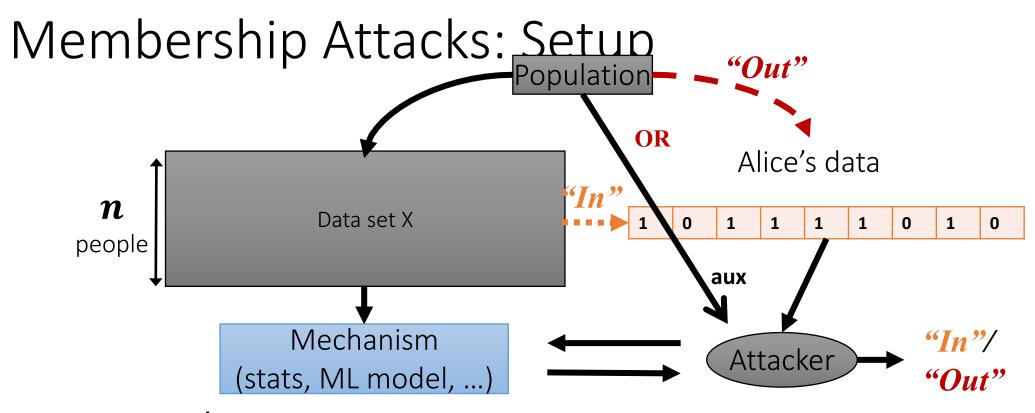
$$m$$
 averages each to within  $\pm O\left(\left(\frac{1}{\sqrt{k}}\right)\cdot\sqrt{\log m}\right)$ 

### Motivation

• Last time: on a dataset with n individuals, releasing m=n averages with error  $E=o(1/\sqrt{n})$  allows for reconstructing 1-o(1) fraction of sensitive attributes.

• Q: what happens if we allow error  $\Omega(1/\sqrt{n}) \le E \le o(1)$ ?

• A (today): if we release  $m=n^2$  counts, can be vulnerable to "membership attacks".



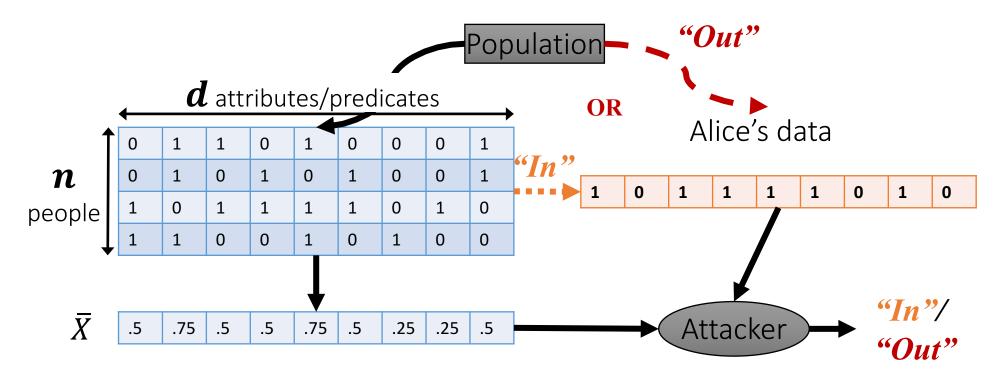
#### Attacker gets:

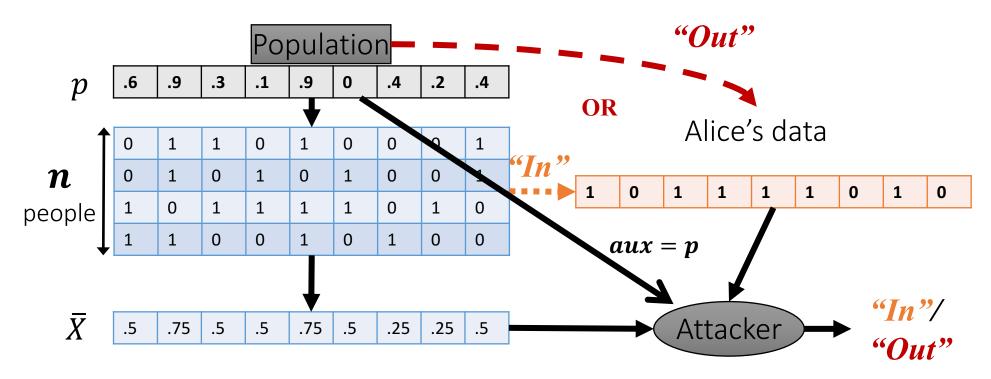
- Access to mechanism outputs
- Alice's data
- (Possibly) auxiliary info about population

#### Then decides: if Alice is in the dataset X

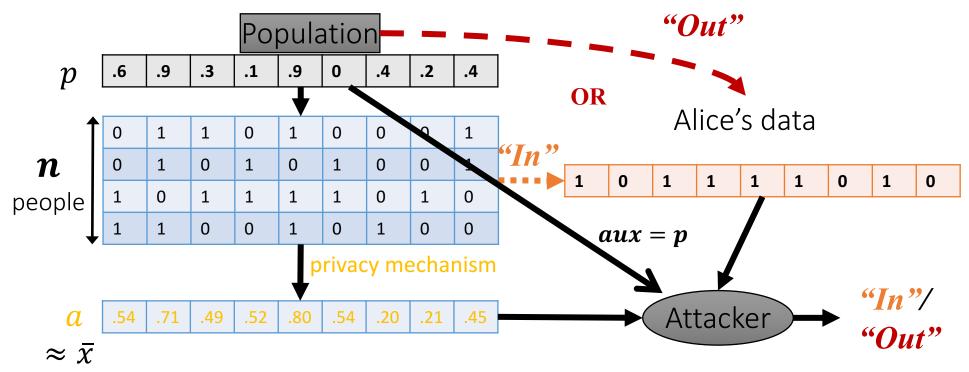
Membership Attacks: Examples Population **OR** Alice's data n Data set X 0 people aux Mechanism Attacker (stats, ML model, ...)

- Genome-wide Association Studies [Homer et al. `08]
  - release frequencies of SNP's (individual positions)
  - determine whether Alice is in "case group" [w/a particular diagnosis]
- ML as a service [Shokri et al. `17]
  - apply models trained on X to Alice's data

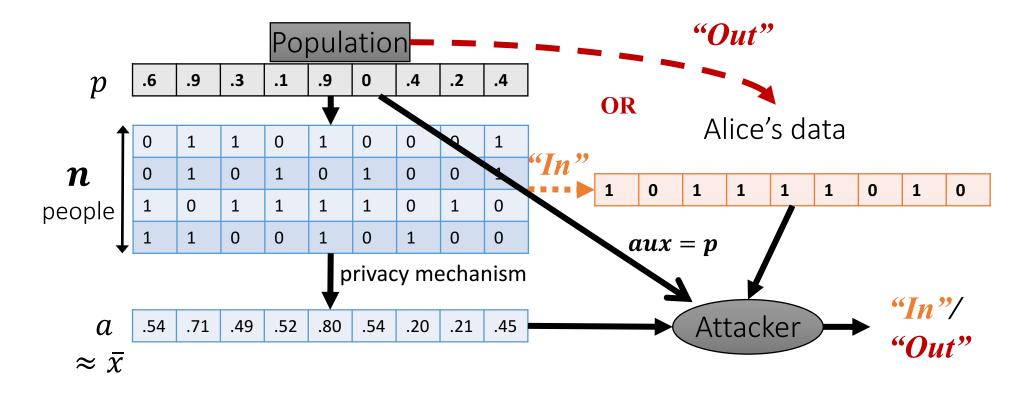




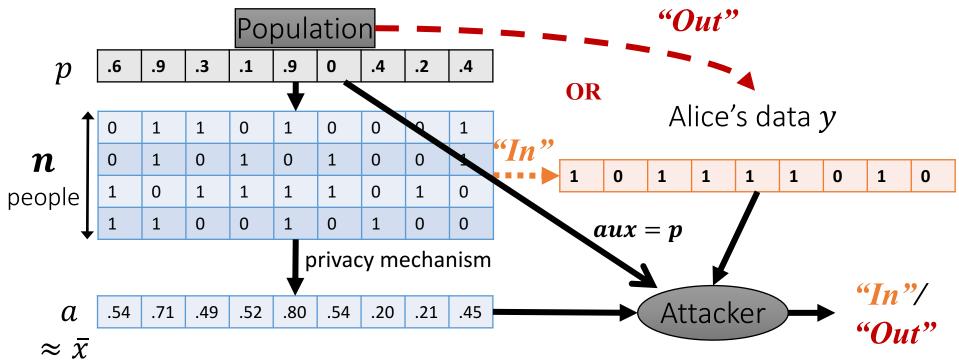
- Population = [vector  $p = (p_1, ..., p_d)$  of probabilities]
  - j'th attribute = iid Bernoulli( $p_i$ ), independent across j
  - Attacker gets p (or a few random draws)



- Population = vector  $p = (p_1, ..., p_d)$  of probabilities
  - j'th attribute = iid Bernoulli( $p_i$ ), independent across j
  - Adversary gets  $a \approx \bar{x}$  and p (or a few random draws)
  - Only assume that a = M(x) has  $|a_j \bar{x}_j| \le \alpha$  whp. ("Noise" need not be independent or unbiased.)



- We are interested in  $\alpha > 1/\sqrt{n}$ .
- In this regime, if p known to mechanism, can prevent attack. (Q: Why?)
- So we will assume random  $p_j$ 's (e.g. iid uniform in [0,1]).



Theorem [Dwork et al. `15]: There is a constant c and an attacker A such that when  $d \ge cn$  and  $\alpha = |a - \bar{x}| < \min \left\{ \sqrt{d/O(n^2 \log(1/\delta))} \right\}$ .

- If Alice is IN, then  $\Pr[A(y, a, p) = IN] \ge \Omega(\frac{1}{\alpha^2 n})$ .
- If Alice is OUT, then  $Pr[A(y, a, p) = IN] \le \delta$ .

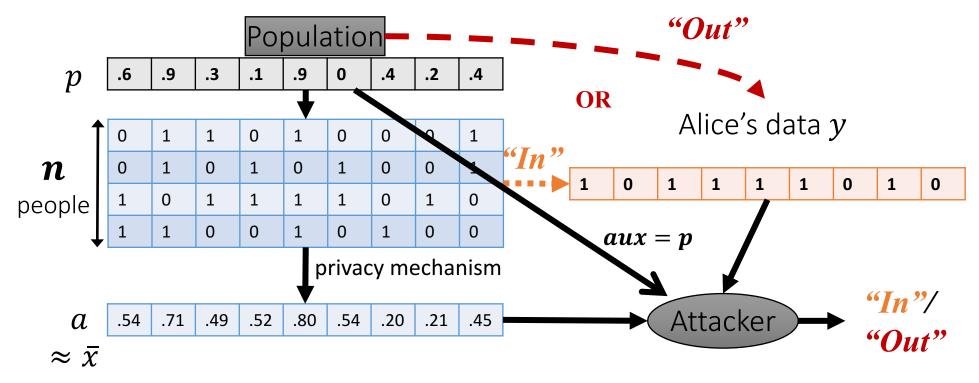
Theorem [Dwork et al. `15]: There is an attacker A such that when  $d \geq O(n)$  and  $\alpha < \min\left\{\sqrt{d/O(n^2\log(1/\delta))}, 1/2\right\}$ :
• If Alice is IN, then  $\Pr[A(y,a,p) = \text{IN}] \geq \Omega\left(\frac{1}{\alpha^2n}\right)$ . (true

- positive)
- If Alice is OUT, then  $\Pr[A(y, a, p) = IN] \leq \delta$ . (false positive)

#### Remarks:

- Only interesting when  $\delta < \Omega\left(\frac{1}{\alpha^2 n}\right)$ .
   On average, successfully trace  $\Omega\left(\frac{1}{\alpha^2}\right)$  members of dataset. This is the best possible. (Why?)
- Gives hope of safely release at most  $\tilde{O}(n^2)$  means!

#### The Attacker



Q: How would you do the attack?

$$A(y, a, p) = \begin{cases} IN & \text{if } \langle y, a \rangle - \langle p, a \rangle > T \\ OUT & \text{if } \langle y, a \rangle - \langle p, a \rangle \le T \end{cases}$$

Note: given p,a, can choose  $T=T_{p,a}=O(\sqrt{d\log(1/\delta)})$  to make false positive probability exactly  $\delta$ .

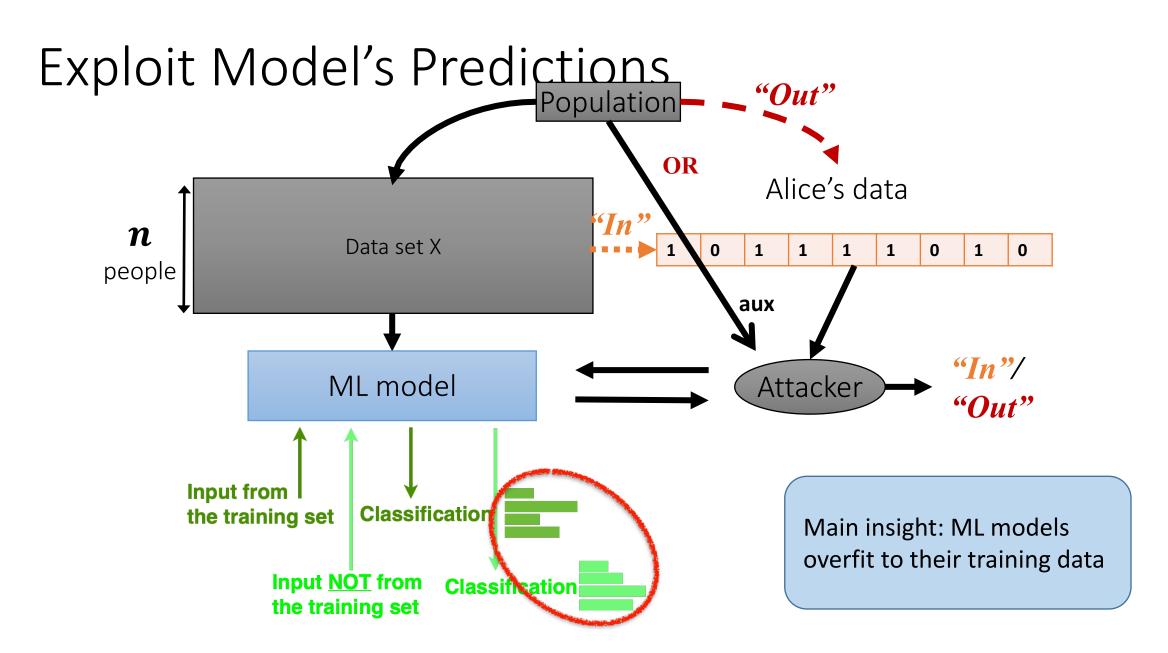
# Attacks on Aggregate Stats (mean)

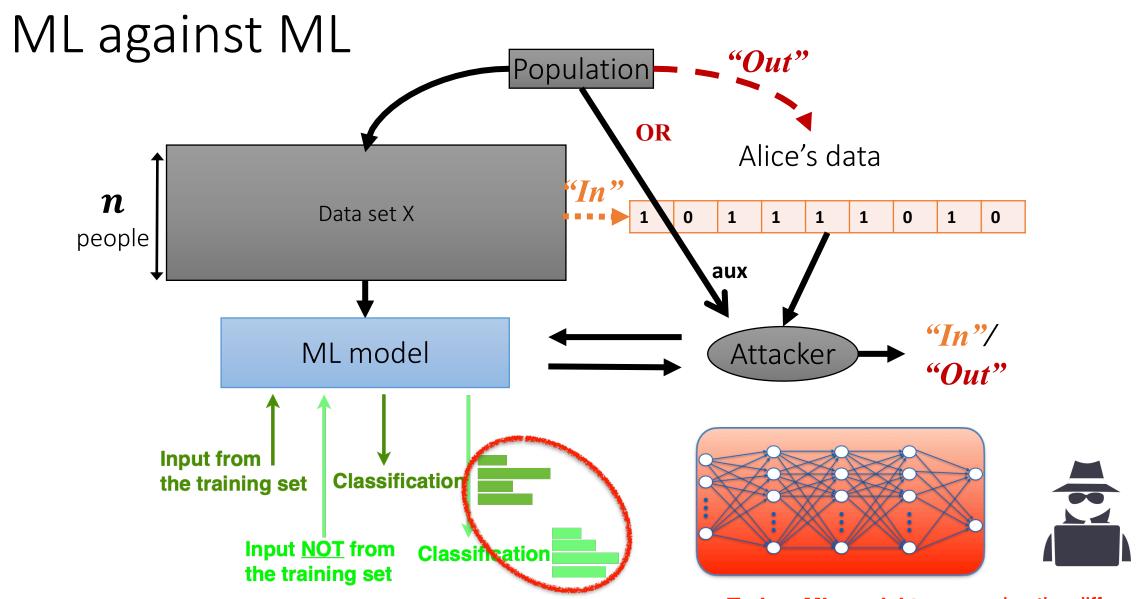
- What error  $\alpha$  makes sense?
  - Estimation error due to sampling  $\approx 1/\sqrt{n}$
  - Reconstruction attacks require  $\alpha \leq 1/\sqrt{n}$ ,  $d \geq n$
  - Robust membership attacks:  $\alpha \lesssim \sqrt{d}/n$
- Lessons
  - "Too many, too accurate" statistics reveal individual data
  - "Aggregate" is hard to pin down



# Membership Attacks on ML

[Shokri et al. 2017]



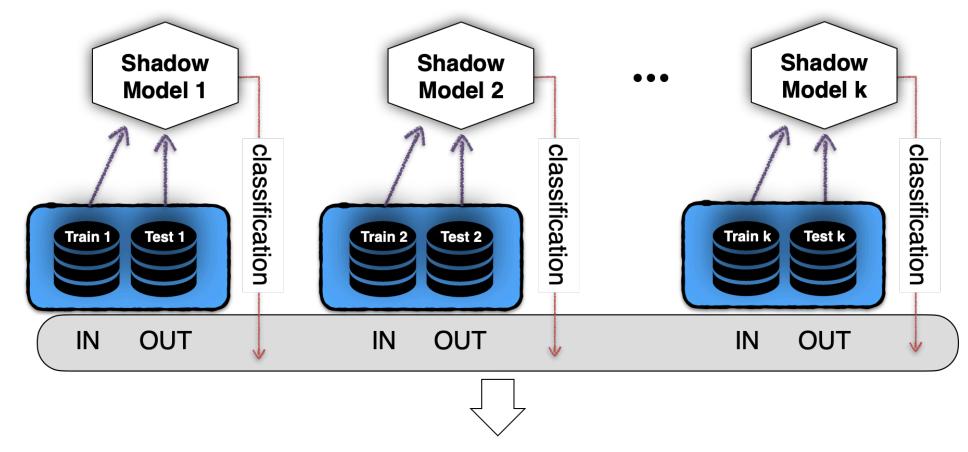


Train a ML model to recognize the difference

Recognize the difference

[slide based on one from Reza Shokri]

## Attack Technique – Shadow Models





#### Train the attack model

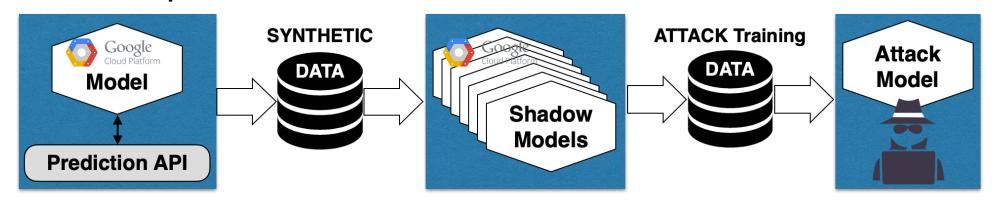
to predict if an input was a member of the training set (in) or a non-member (out)

# Obtaining Data for Training Shadow Models

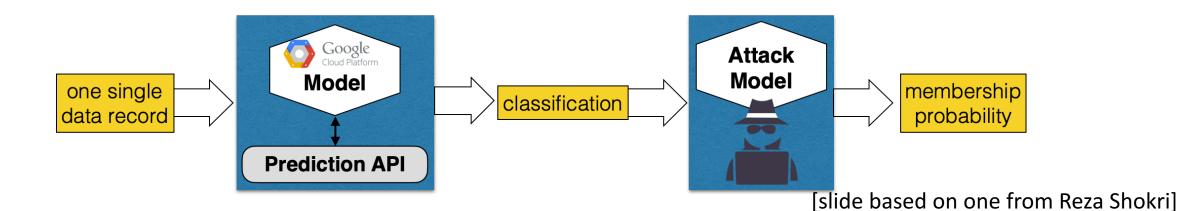
• **Real**: similar to training data of the target model (i.e., drawn from same distribution)

• **Synthetic**: use a sampling algorithm to obtain data classified with high confidence by the target model

### Attack Pipeline

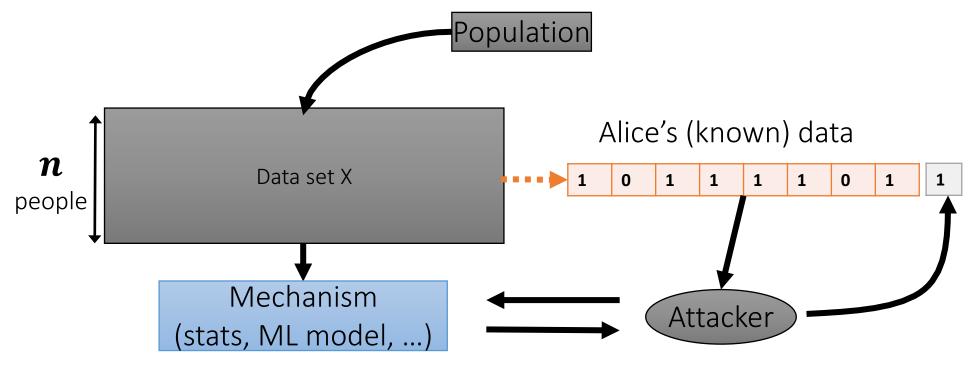


# Using the Attack Model



### Another Attack on ML?

[Frederickson et al. `14, cf. McSherry `16]



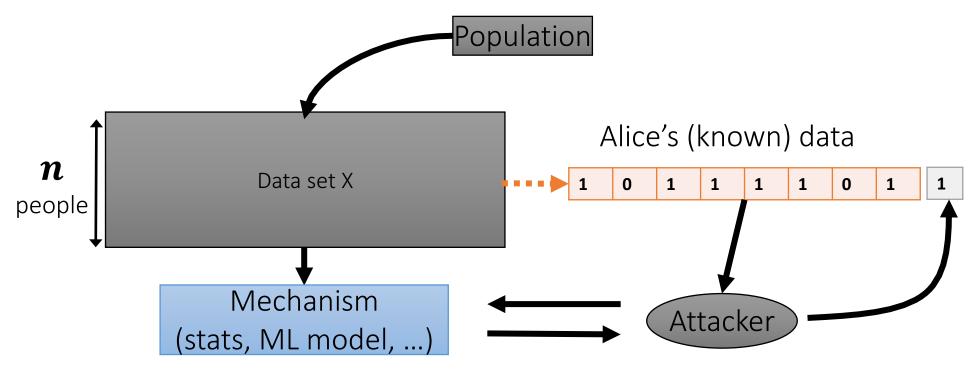
#### Attacker gets:

- Access to mechanism outputs
- Some of Alice's data
- (Possibly) auxiliary info about population

Then computes: a sensitive attribute of Alice

### Another Attack on ML?

[Frederickson et al. `14, cf. McSherry `16]



#### Difference from reconstruction attacks:

- Above attack works even if Alice not in dataset. Based on correlation between known & sensitive attributes.
- Reconstruction attacks work even when sensitive bit uncorrelated.

## Goals of Differential Privacy

- Utility: enable "statistical analysis" of datasets
  - e.g. inference about population, ML training, useful descriptive statistics
- Privacy: protect individual-level data
  - against "all" attack strategies, auxiliary info.