

## Assignment 4- stacks + queues Amalia Karaman

### 1) Stacking

→ Given an empty stack, what'll be the contents of stack after following operations?

1.  $\text{push}(8)$  →  $[8]$
  2.  $\text{push}(2)$  →  $[8, 2]$
  3.  $\text{pop}()$  →  $[8]$  (remove 2)
  4.  $\text{push}(\text{pop()} \times 2)$  →  $[]$  →  $[16]$  ( $\text{pop } 8, \times 2$ )
  5.  $\text{push}(10)$  →  $[16, 10]$
  6.  $\text{push}(\text{pop}()/2)$  →  $[16, 5]$  ( $\text{pop } 10, \div 2$ )
- Fin:  $[16, 5]$

### 2) Queuing

Given empty queue, what are contents of stack after following operations?

1.  $\text{push}(4)$  →  $[4]$  (pushed 4)
2.  $\text{push}(\text{pop}() + 4)$  →  $[8]$  ( $\text{pop } 4 + \text{add } 4$ )
3.  $\text{push}(8)$  →  $[8, 8]$  (8)
4.  $\text{push}(\text{pop}()/2)$  →  $[8, 4]$  ( $\text{pop } 8 + \text{divide by } 2$ )
5.  $\text{pop}()$  →  $[4]$  (remove 8)
6.  $\text{pop}()$  →  $[]$  (remove 4)

Fin:  $[]$

empty queue, no values remain

#3)

```
int findInDeque(Deque<Integer> q, int x) {
    int n = q.size();
    int leftIndex = 0, rightIndex = n - 1;
    // search from the front
    Iterator<Integer> frontIterator = q.iterator();
    while (frontIterator.hasNext()) {
        if (frontIterator.next() == x) {
            return leftIndex; // Found from the left
        }
        leftIndex++;
    }
    // search from the back
    Iterator<Integer> backIterator = q.descendingIterator();
    while (backIterator.hasNext()) {
        if (backIterator.next() == x) {
            return rightIndex; // Found from the right
        }
        rightIndex--;
    }
    return -1; // If x is not found
}
```

To find the position of  $x$  in the deque, I start by checking from both ends to minimize the number of steps. I set `leftIndex` to 0 (starting from the front) and `rightIndex` to  $n - 1$  (starting from the back). I iterate through the deque from the left, increasing `leftIndex` until I find  $x$ . At the same time, I iterate from the right, decreasing `rightIndex`. If  $x$  is found on either side, I return the corresponding index. If  $x$  is not found, I return -1. Since I only search up to  $n/2$  elements on average, the time complexity is  $O(n/2)$ , which simplifies to  $O(n)$ .



7) For algorithms written in #4-6, explain their time complexity + space complexity in Big-O notation. Explain how.

#### #4) Balanced Brackets

Time:  $O(n)$

Space:  $O(n)$

Time: Input is scanned once by one + each bracket is processed once, + each character is push/popped onto/from the stack. Pushing/popping takes  $O(1)$  time so for  $n$  brackets it's  $O(n)$ .

Space: In worst case all  $n$  brackets are stored in the stack so max stack size is  $O(n)$ .

#### #5) Decode String

Time:  $O(n)$

Space:  $O(n)$

Time: Since stack operations take  $O(1)$  time, this time complexity is  $O(1)$  also b/c each character processed @ MOST twice.

Space: The stack stores nested sequences + repeat counts + in worst case output is  $O(n)$  in size.

#### #6) Infix to Postfix

Time:  $O(n)$

Space:  $O(n)$

Time: It scans each <sup>character</sup> operation once + push/pop operators once as well, and stack operations are  $O(1)$  so we have  $O(n)$  for this algorithm.

Space: It holds the operators TEMPORARILY, + each the stack are processed once. In worst case all  $n$  operators/parentheses could be held at once.