

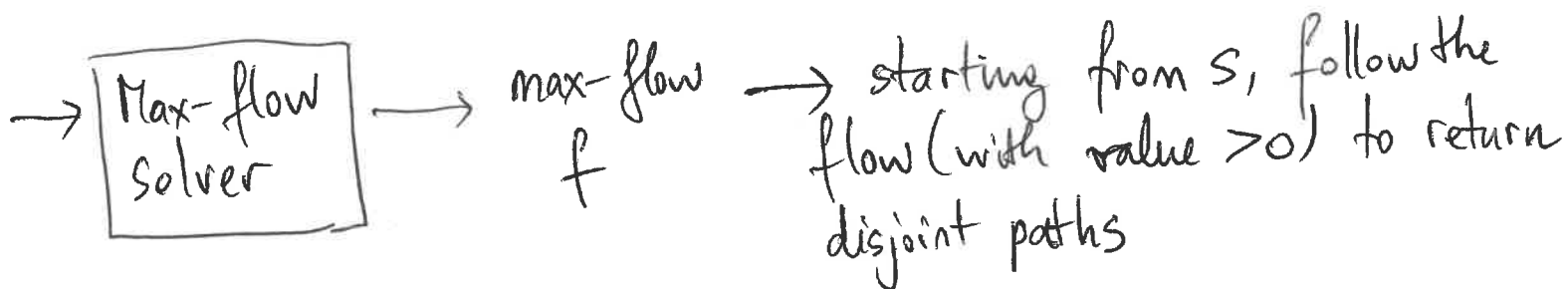
Disjoint paths in a directed graph

Input

graph $G(V, E)$ directed
vertices s and t

Flow Network

- graph $G(V, E)$ directed
- source s and sink t
- capacity 1 on all edges



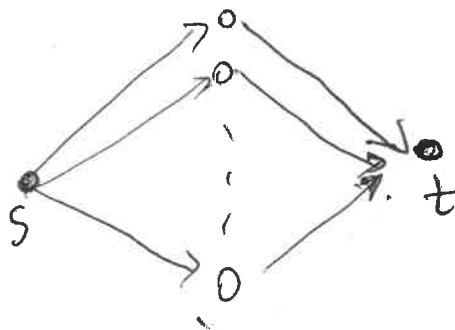
RT analysis

Ford-Fulkerson $\Rightarrow RT = O(|f^*| \cdot E)$

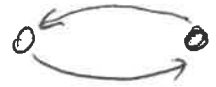
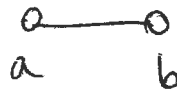
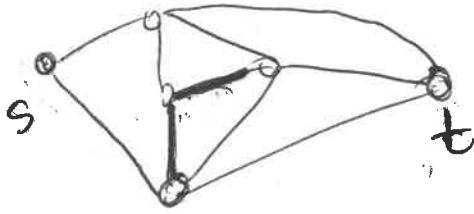
$$|f^*| \leq |V| - 2$$

$$|f^*| = O(V)$$

$RT = O(V \cdot E)$



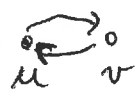
Disjoint paths in undirected graphs



Input

graph $G(V, E)$ undirected
vertices s and t

Flow Network

- build directed graph $G'(V', E')$
- replace each edge (u, v) by two directed edges 
- remove edges into s , and out of t
- remove antiparallel edges
- source s , sink t
- capacity \pm on all edges

Max-flow
Solver

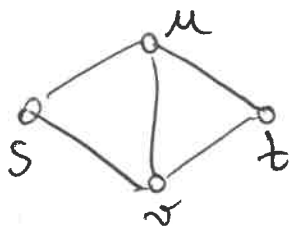
Max-flow
 f

remove flow
on antiparallel
edges

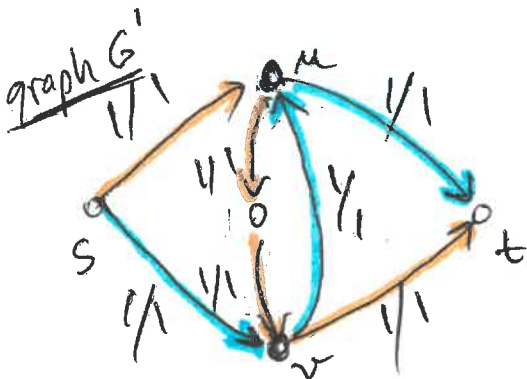
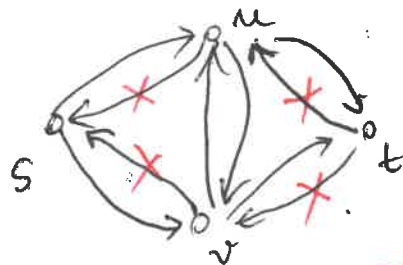
starting from
 s , follow the
flow (with value
 > 0) to return
disjoint paths

example

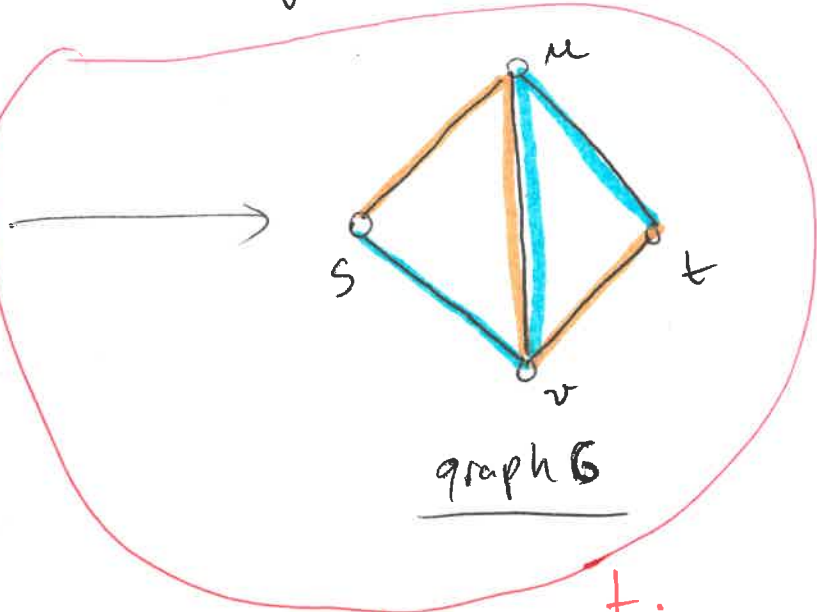
graph $G(V, E)$



graph $G'(V', E')$



- remove antiparallel edges
- add capacity \downarrow

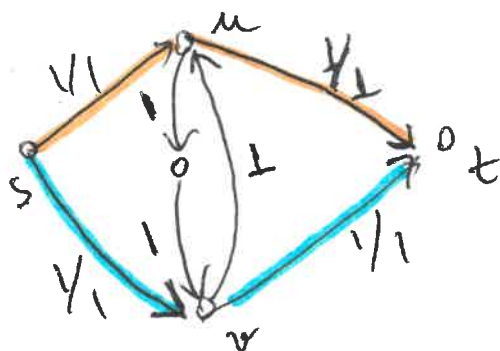


Not correct:
overlapping edges in the resulting paths!

remove flow on antiparallel edges:

$$f = \min\{f(u,v), f(v,u)\}$$

$$f = 1$$



graph G

- edge disjoint paths in G

RT analysis

Ford-Fulkerson $\Rightarrow RT = O(|f^*| \cdot |E'|)$

$$|f^*| \leq |V| - 2 \Rightarrow |f^*| = O(V)$$

$$|E'| \leq 2 \cdot |E| + |E| = 3 \cdot |E| \Rightarrow |E'| = O(E)$$



$$\Rightarrow \boxed{RT = O(V \cdot E)}$$