

3a) Assume brute force/backtracking Algorithm
key to the solution is an adjacency matrix

Assume $X[n,n]$ is a connection matrix for the graph G .

backtrack(m)

For $i := 1$ TO N // each row

\rightarrow backtrack($X[i], i, M$) // cur // Check from current node only

\rightarrow backtrack($Kopm$) (return)

\rightarrow backtrack($X[i], K, M$) // each // find m nodes in this row starting at K

 For $j := K$ TO N

 IF $j \neq i$ AND $X[j] == \text{false}$

 // add $x[K, j]$ to solution set

 IF $M == 1$

 output solution

 else // need more

\rightarrow backtrack($X[i], j, M-1$)

 // clear solution

Alternate solution (my preferred) non recursive)

Create Adjacency matrix, assume each node is connected to itself then count the false in each node connections

OR

Create Adjacency List and look for list entry count less than $n-m$. Nodes missing form independent set.

3b) $RT(3a) = \Theta(\sum_{i=1}^n R) = \Theta(n^2)$