

Hermite Interpolation: some points on the poly & some points on its derivatives

Example

$$\left. \begin{array}{l} x=1 \quad f(1)=10 \\ x=3 \quad f(3)=58 \\ x=2 \quad f'(2)=23 \\ x=4 \quad f'(4)=67 \end{array} \right\} \begin{array}{l} f^{(0)}(x) = a_0 + a_1x + a_2x^2 + a_3x^3 \\ \downarrow \\ f^{(1)}(x) = a_1 + 2a_2x + 3a_3x^2 \end{array} \rightarrow \begin{array}{l} 4 \text{ unknowns} \\ 4 \text{ equations} \end{array}$$

$$\left\{ \begin{array}{l} f^{(0)}(1) = a_0 + a_1 + a_2 + a_3 = 10 \\ f^{(0)}(3) = a_0 + 3a_1 + 9a_2 + 27a_3 = 58 \\ f^{(1)}(2) = a_1 + 4a_2 + 12a_3 = 23 \\ f^{(1)}(4) = a_1 + 8a_2 + 48a_3 = 67 \end{array} \right. \rightarrow \begin{array}{l} a_0 = 4 \\ a_1 = 3 \\ a_2 = 2 \\ a_3 = 1 \end{array}$$

$$f(x) = 4 + 3x + 2x^2 + x^3$$

Multi-level access structure

- each player is assigned to a certain level
- each level "L" is associated with a threshold t_L such that $t_0 < t_1 < \dots < t_n$

(a) Disjunctive secret sharing:

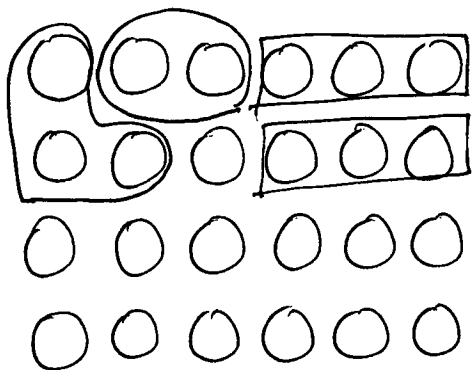
A group of players can recover the secret if the group contains at least t_L players at levels $0 \sim L$ for some level L .
I.e., at least one threshold must be satisfied.

(b) Conjunctive Secret sharing

A group of players can recover the secret if the group contains at least t_L players at levels $0 \sim L$ for every level L .
I.e., all the thresholds must be satisfied.

Example:

$$t_0 < t_1 < \dots < t_n$$



$$\longrightarrow L_0: t_0 = 2$$

$$\longrightarrow L_1: t_1 = 3$$

$$\longrightarrow L_2: t_2 = 4$$

$$\longrightarrow L_3: t_3 = 6 \quad \text{max threshold}$$

Disjunctive 1988

- 2 players at L_0 is enough
- ✓ 3 players at $L_0 | L_1$ "
- ✓ 4 " at $L_0 | L_1 | L_2$ "
- ✓ 6 " at $L_0 | L_1 | L_2 | L_3$ "

Conjunctive 2004

- At least 2 players at L_0
- ∧ " 3 " at $L_0 | L_1$
- ∧ " 4 " at $L_0 | L_1 | L_2$
- ∧ " 6 " at $L_0 | L_1 | L_2 | L_3$

* users/players at lower levels are more important & we don't need the contribution of players at all levels

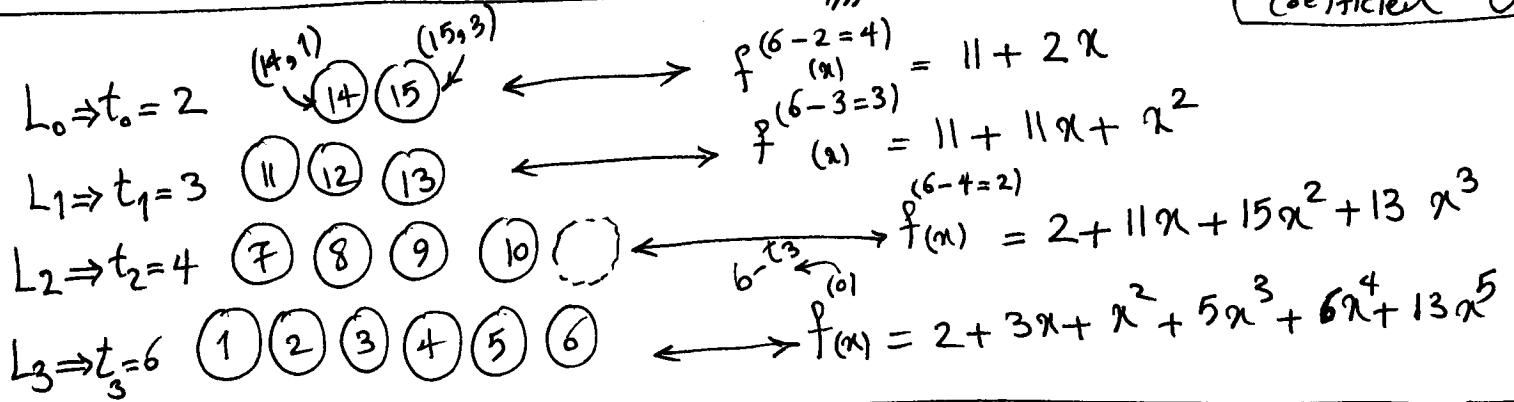
Example of Disjunctive SS

\mathbb{Z}_{19}

3

$t_0 = 2 < t_1 = 3 < t_2 = 4 < t_3 = 6$ max threshold

$f(x) = 2 + 3x + x^2 + 5x^3 + 6x^4 + 13x^5 \rightarrow$ secret is the leading coefficient



dealer does not need this poly $\otimes f^{(1)} = 3 + 2x + 15x^2 + 5x^3 + 8x^4$

$p_{14} \leftarrow (1)$ share

$p_{15} \leftarrow (3)$ share

Ids are in monotonically decreasing order by level (can be random)

$f^{(0)}(x) = a + bx + cx^2 + dx^3 + ex^4 + gx^5 \rightarrow$ secret: leading coefficient

$f^{(1)}(x) = b + 2cx + 3dx^2 + 4ex^3 + 5gx^4$

$f^{(2)}(x) = 2c + 6dx + 12ex^2 + gx^3$

$f^{(3)}(x) = 6d + 5ex + 3gx^2$

$f^{(4)}(x) = 5e + 6gx \xrightarrow[(15, 3)]{(14, 1)} \begin{cases} 5e + 6g(14) = 1 \\ 5e + 6g(15) = 3 \end{cases} \rightarrow$

2 unknowns & 2 equations $\rightarrow e = 6$
 $\boxed{g = 13} \rightarrow$ secret

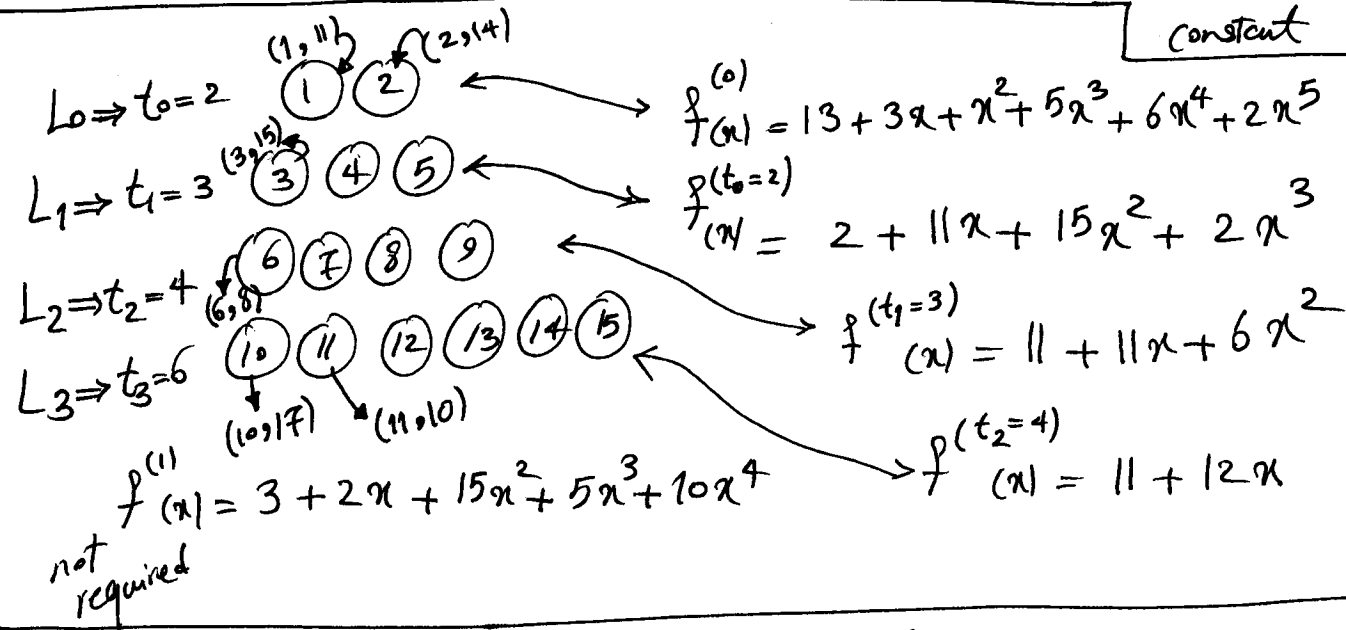
Example of Conjective SS

\mathbb{Z}_{19}

4

$t_0=2 < t_1=3 < t_2=4 < t_3=6$ Ids are in increasing order

$f(x) = \boxed{13} + 3x + x^2 + 5x^3 + 6x^4 + 2x^5 \rightarrow$ secret is the constant term



$$\begin{aligned}
 f^{(0)}(x) &= a + bx + cx^2 + dx^3 + ex^4 + gx^5 \rightarrow \begin{cases} a+b+c+d+e+g=11 \\ a+2b+4c+8d+16e+13g=14 \end{cases} \\
 f^{(2)}(x) &= 2c + 6dx + 12ex^2 + gx^3 \rightarrow \begin{cases} 2c+18d+13e+8g=15 \\ 6d+11e+13g=8 \end{cases} \\
 f^{(3)}(x) &= 6d + 5ex + 3gx^2 \rightarrow \begin{cases} 5e+3g=17 \\ 5e+9g=10 \end{cases} \\
 f^{(4)}(x) &= 5e + 6gx
 \end{aligned}$$

→ $\boxed{a=13}$ ✓

$b=3$ $e=6$
 $c=1$ $g=2$
 $d=5$