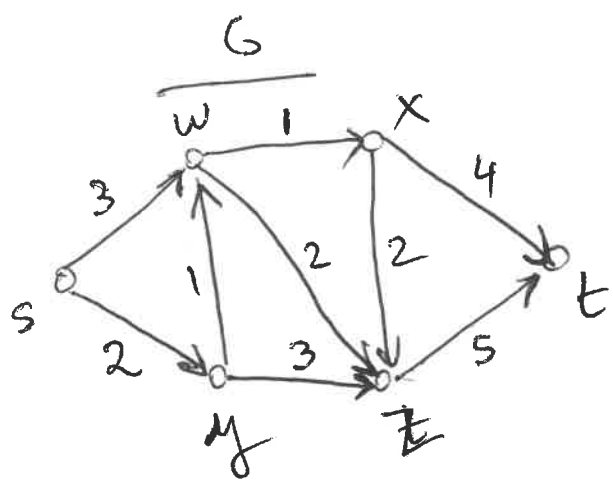
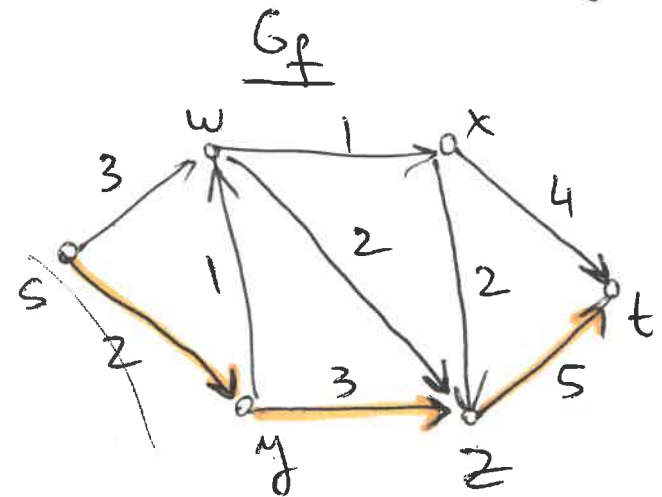


• Use Ford-Fulkerson to compute the maximum flow:



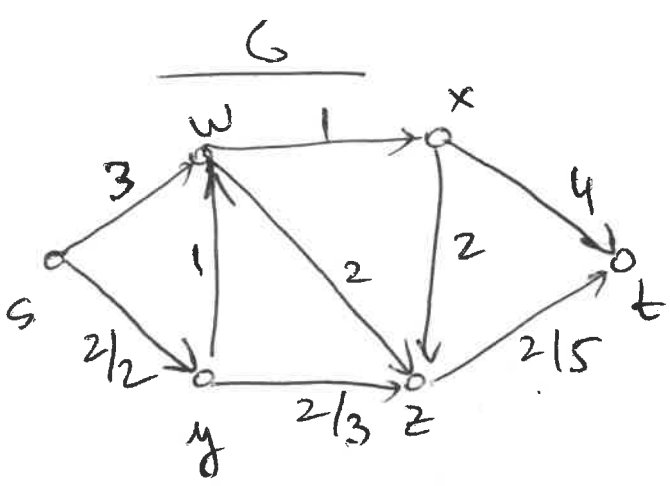
$|f| = 0$



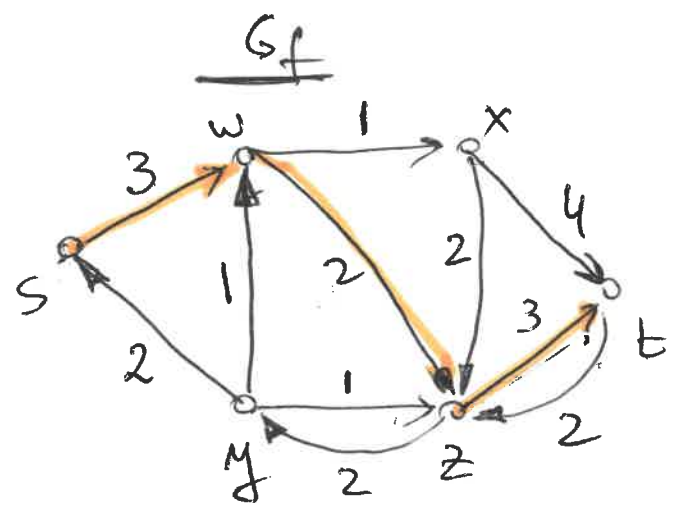
$p = \langle s, y, z, t \rangle$

$c_f(p) = 2$

$|f_p| = 2$



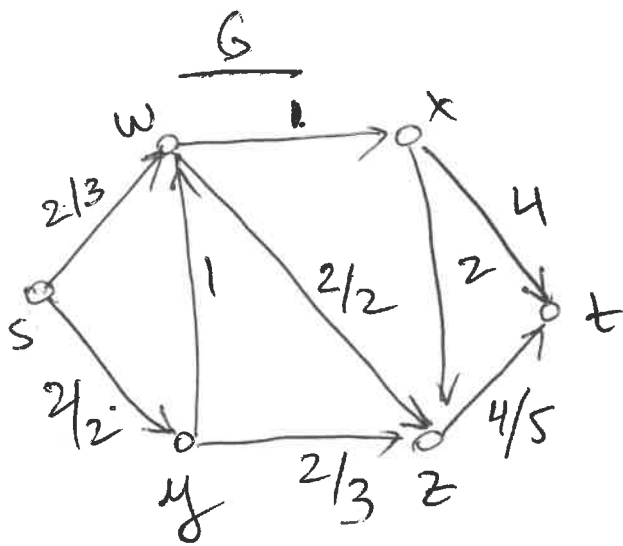
$|f| = 2$



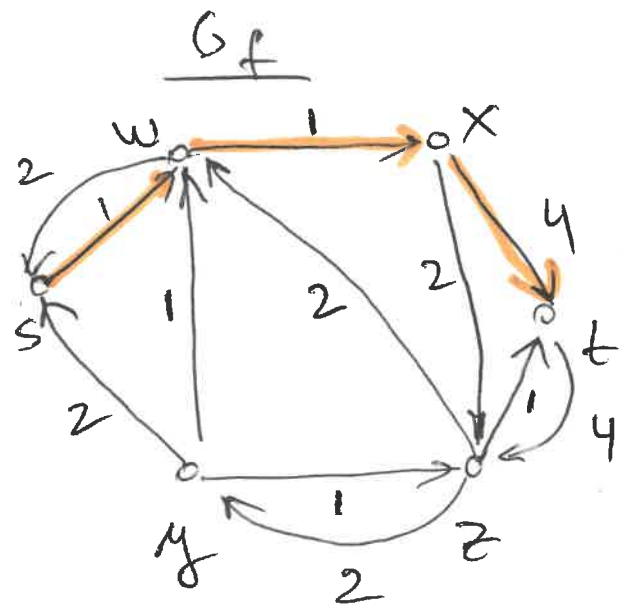
$p = \langle s, w, z, t \rangle$

$c_f(p) = 2$

$|f_p| = 2$



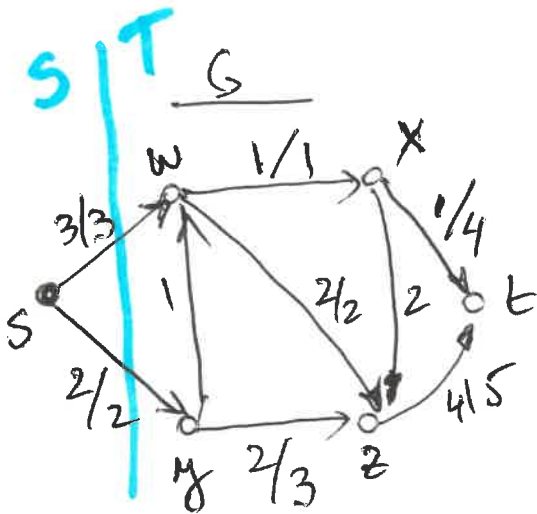
$$|f| = 4$$



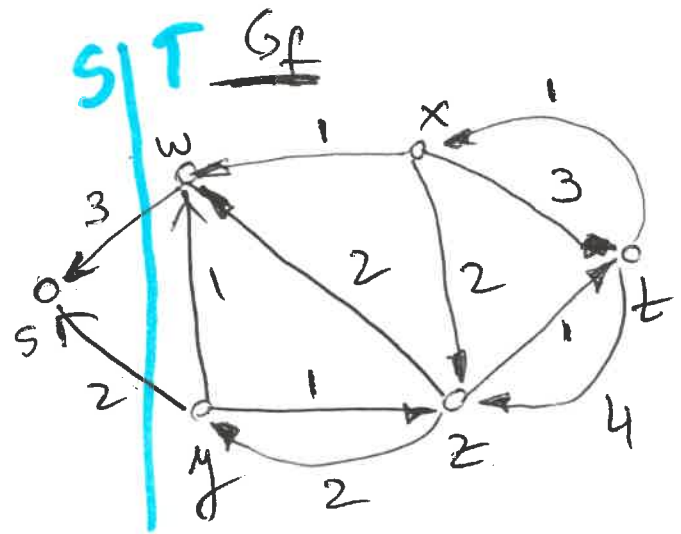
$$p = \langle s, w, x, t \rangle$$

$$c_f(p) = 1$$

$$|f_p| = 1$$



$$|f| = 5$$



- no augmenting path \Rightarrow
max-flow reached

- find a cut (S, T) s.t. $c(S, T) = |f| = 5$ (for max-flow!)

$$S = \{s\}$$

$$T = \{w, x, y, z, t\}$$

$$c(S, T) = 5$$

flow-network

- graph $G(V, E)$ directed
- sources s , sink t
- capacities c

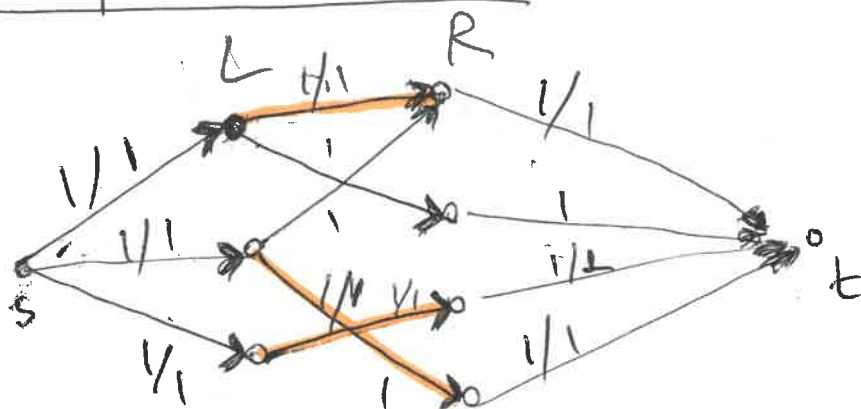


→ maximum-flow f

Ford-Fulkerson : $RT = O(|f^*| \cdot E)$

Edmonds-Karp : $RT = O(V \cdot E^2)$

Maximum bipartite matching



Input

graph $G(V, E)$ bipartite
 $V = L \cup R$



Flow Network

• build directed graph $G'(V', E')$

- add edges in E with direction $L \rightarrow R$

- add sources, sink t

- add edges connecting s to each vertex in L

- add edges from each vertex in R to t

- add capacity 1 on each edge

→ Max-flow
solver

(Ford-Fulkerson
or
Edmonds-Karp)

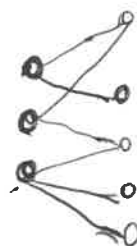
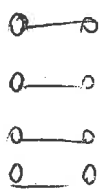
→ Max-flow
f

→ return the edges
in G with flow > 0
as being the edges
in the maximum
matching M

RT analysis

Ford-Fulkerson: $RT = O(|f^*| \cdot |E'|)$

$$|f^*| \leq \frac{|V|}{2}$$



$$|E'| = |E| + |V|$$

$|E| \geq \frac{|V|}{2}$ (each vertex is incident
on at least one edge)

$$\Rightarrow |V| \leq 2 \cdot |E| \quad \} \Rightarrow$$

$$\Rightarrow |E'| \leq |E| + 2|E| = 3|E|$$

$RT = O(V \cdot E)$