

# HF Based on Block Ciphers

1

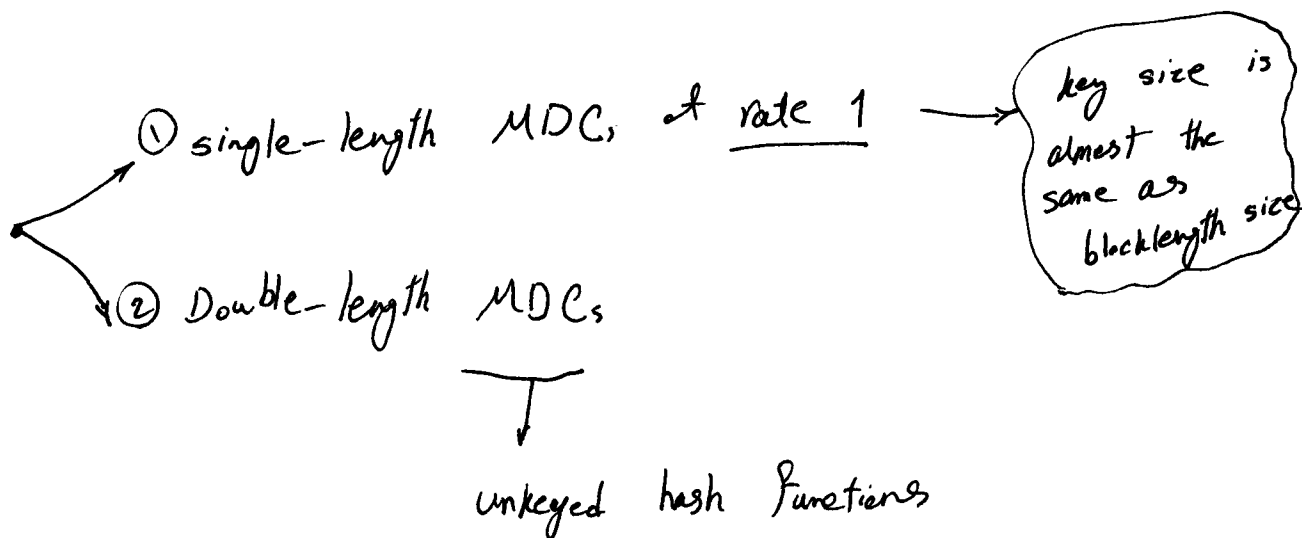
- HF based on Block ciphers
- Customized hash Functions
- HF based on Modular Arithmetic

**Benefit** → There are lots of efficient method to implement them either by hardware / software implementation.

**Def:** An  $(n, r)$  Block cipher is a block cipher defining an invertible function from  $n$ -bit plaintexts to  $n$ -bit ciphertexts using an  $r$ -bit key  $\rightarrow E_k(x)$

$\underbrace{64}_{n\text{-bit}} \leftarrow \text{DES} \rightarrow 56 \underbrace{\text{key}}_{r\text{-bit}}$

**Def:** Let "h" be an iterated hash function constructed from a block cipher, with compression function  $f$  which performs  $s$  block encryptions to process each successive  $n$ -bit message block. Then the rate is  $1/s$ .



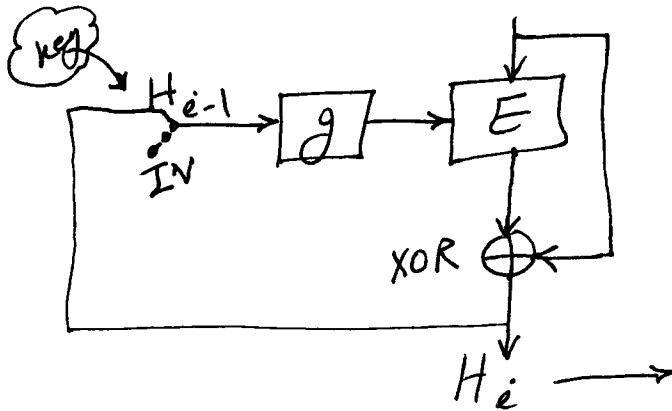
# ① single-length MD Cs at rate 1

2

- Requirements**
- (a) a generic  $n$ -bit block cipher  $E_K$
  - (b) a function " $g$ " which maps  $n$ -bit inputs to keys  $K$
  - (c) a fixed (usually  $n$ -bit) initial value  $IV$

## 1-1 Matyas - Meyer - Oseas

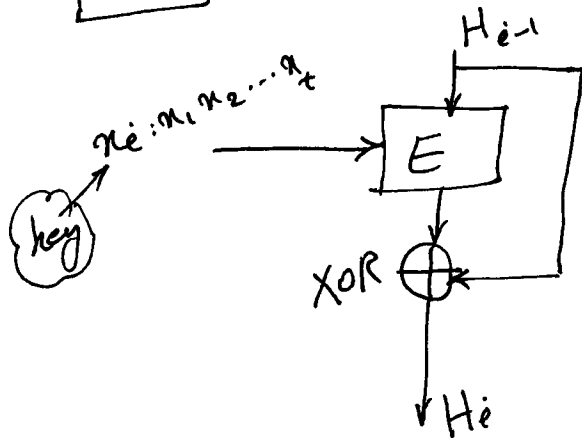
$x_i: x_1, x_2, \dots, x_t$



$t$  blocks of  $n$ -bit message  
Compressed  
hash value  
last ciphertext is a block of  $n$ -bit

## 1-2 Davies - Meyer

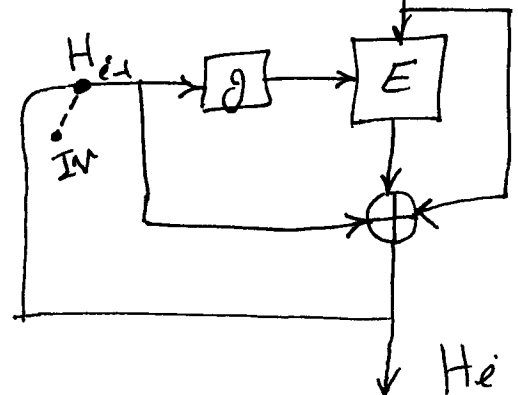
$H_0 = IV$



## 1-3

## Miyaguchi - Preneel

$x_i: x_1, \dots, x_t$



1-1

alg. M-M-O

3

Input: bitstring  $x \xrightarrow{x_i} x_1 x_2 \dots x_t$

output:  $n$ -bit hash code of  $x$

1.  $x$  is divided into  $n$ -bit blocks & padded.

you predefine IV.

2.  $H_t$  :  $H_0 = IV$

$$H_i = E_{g(H_{i-1})}^{(x_i)} \oplus x_i$$

Enc scheme ←  $g(H_{i-1})$  ← key

plaintext ←  $x_i$

$1 \leq i \leq t$

1-2

alg. D-M.

"  
"  
"

2.  $H_t$  :  $H_0 = IV$

$$H_i = E_{x_i} (H_{i-1} \oplus H_{i-1}) \quad 1 \leq i \leq t$$

1-3

alg. M-P

"  
"

2.  $H_t$  :  $H_0 = IV$

$$H_i = E_{g(H_{i-1})}^{(x_i)} \oplus x_i \oplus H_{i-1}$$

$1 \leq i \leq t$

## ② Double-length MDC<sub>s</sub>:

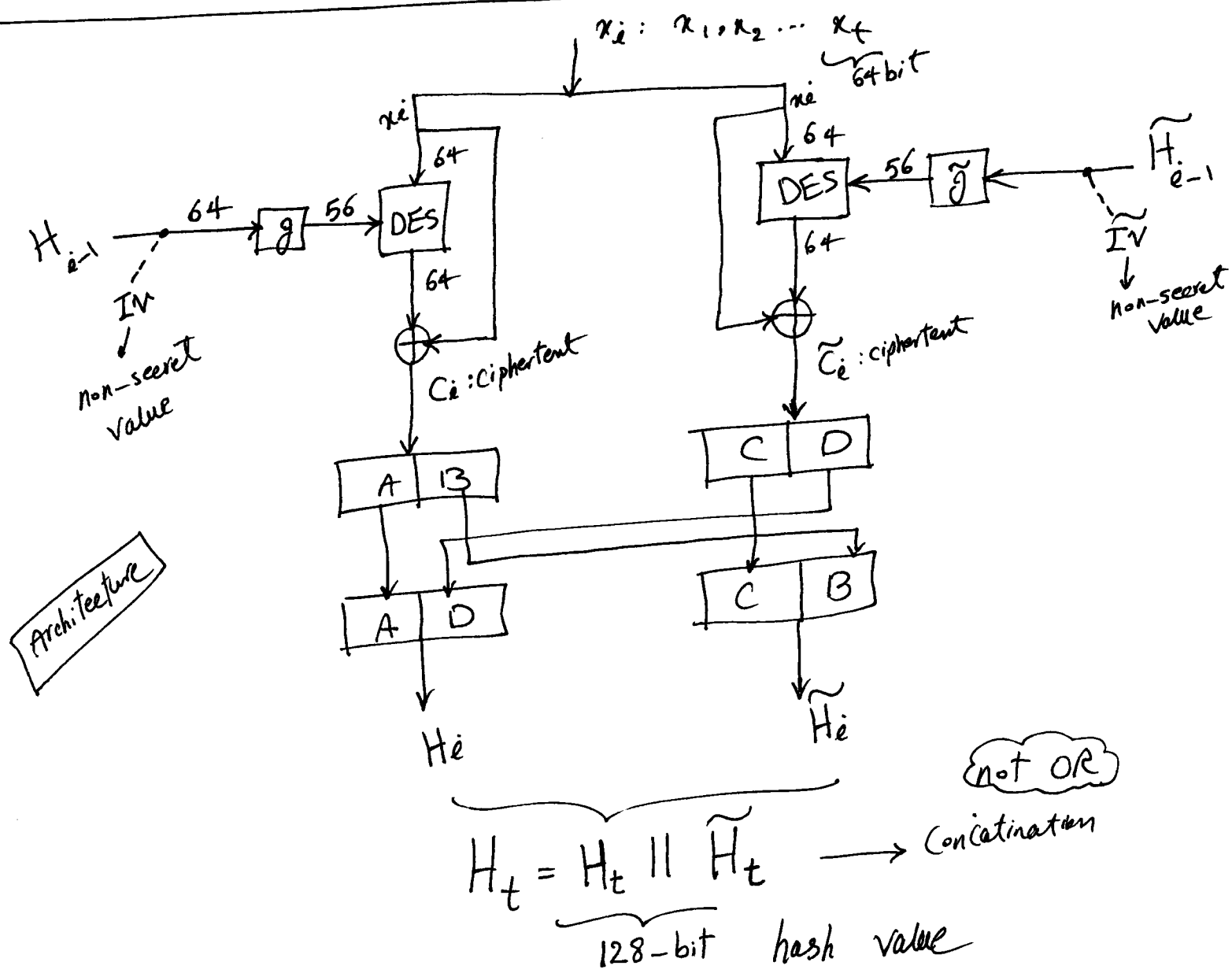
Requirements

- (a) DES bitlength = 64 key size 56
- (b)  $g$  &  $\tilde{g}$  which maps 64-bit to 56-bit

$$g(u) = u_1 \ 1 \ 0 \ u_4 \ \dots \ u_7^x \ u_9 \ \dots \ u_{63}$$

$$\tilde{g}(u) = u_1 \ 0 \ 1 \ u_4 \ \dots \ u_7^x \ u_9 \ \dots \ u_{63}$$

(excluded  $u_8, u_{16}, u_{24}, \dots, u_{64}$ )



## algorithm

$$H_0 = IV \quad k_i = g(H_{i-1}) \quad C_i = E_{k_i}(x_i) \oplus x_i$$

$$\tilde{H}_0 = \tilde{IV} \quad \tilde{k}_i = g(\tilde{H}_{i-1}) \quad \tilde{C}_i = E_{\tilde{k}_i}(x_i) \oplus x_i$$

$$H_i = C_i^L \parallel \tilde{C}_i^R \quad \tilde{H}_i = \tilde{C}_i^L \parallel C_i^R$$