

Rabin public-key Encryption

1

Key Generation Alg.

1. Generate two large random (distinct) primes p & q (almost the same size).
2. Compute $n = p * q$
3. A's public key is " n ", private key is (p, q)

Algorithm of Rabin's Scheme

Encryption by B

1. Obtain A's authentic public-key " n "
2. Represent the message as an integer " m " in $\{0, 1, \dots, n-1\}$
3. Compute ciphertext $c = m^2 \pmod{n}$
4. Send the " c " to another party "A"

Decryption by A

1. use an algorithm to find four square roots m_1, m_2, m_3, m_4 of " c " (\pmod{n}).
2. the actual message " m " is one of m_1, m_2, m_3, m_4

x Dec \rightarrow computationally more expensive
 x \rightarrow "A" has to decide which " m_i " is the message

Algorithm: Finding square roots mod n & its prime factors p & q

Input: n, p, q, a

output: Four square roots of $a \pmod{n}$

1. Use **another algorithm** to find two square roots " r " and " $-r$ " of " a " (\pmod{p})
2. Use **another algorithm** to find two square roots " s " and " $-s$ " of " a " (\pmod{q})
3. Use **EE algorithm** to find

$$cp + dq = 1$$

$$4. \quad x = r \underline{d} q + s \underline{c} p \pmod{n}$$

$$y = r \underline{d} q - s \underline{c} p \pmod{n}$$

5. return $(\pm x, \pm y) \pmod{n}$

Example

cipher $a = 62111$, $p = 277$, $q = 331$, $n = pq = 91687$
 ± 150 (under p), ± 144 (under q)

Using EE Alg. $\rightarrow \underbrace{331}_q \times \underbrace{118}_d + \underbrace{277}_p \times \underbrace{(-141)}_c = 1$

stage 4 $\rightarrow x = 150 \times 118 \times 331 + 144 \times (-141) \times 277 \pmod{91687}$

$y =$

x	$11482908 \pmod{n} = \overbrace{22033}^{m_2}$	$-x \pmod{n} = \overbrace{69654}^{m_1}$
y	$234492 \pmod{n} = \overbrace{51118}^{m_4}$	$-y \pmod{n} = \overbrace{40569}^{m_3}$

$\sqrt{m_3}$

Algorithm X_1 : Find square roots (mod) prime p if $p \equiv 3 \pmod{4}$

Input: odd prime p where $p \equiv 3 \pmod{4}$ & "a" cipher 62111

output: Two square roots of "a" (mod p)

1. Compute $r = a^{\frac{p+1}{4}} \pmod{p}$ S-S-M alg

2. return $(r, -r)$

Example

(a prime #) $\rightarrow q = 331 \rightarrow 331 \equiv 3 \pmod{4}$

$$r = 62111^{\frac{332}{4}} \pmod{331} = 144 \rightarrow \pm 144$$

Algorithm X_2 : Find square roots (mod) prime p if $p \equiv 5 \pmod{8}$

Input: odd prime p where $p \equiv 5 \pmod{8}$ & "a"

output: Two square roots of "a" (mod p)

cipher 62111

1. compute $d = a^{\frac{p-1}{4}} \pmod{p}$

2. if $d=1 \rightarrow r = a^{\frac{p+3}{8}} \pmod{p}$

3. if $d=p-1 \rightarrow r = 2a(4a)^{\frac{p-5}{8}} \pmod{p}$

4. return $(r, -r)$

Example

(a prime #) $\rightarrow p = 277 \rightarrow 277 \equiv 5 \pmod{8}$

$$d = 62111^{\frac{276}{4}} \pmod{277} = 276 \text{ which is } (p-1)$$

$$r = (2 \times 62111) \times (4 \times 62111)^{\frac{272}{8}} \pmod{277} = 150 \rightarrow \pm 150$$

General Algorithm

5

Algo. Finding square roots mod a prime p

Input: odd prime p and a
output: two square roots of $a \pmod{p}$

1. Choose random $b \in \mathbb{Z}_p$ until $b^2 - 4a$ is
a quadratic non-residue \pmod{p}

$$\left(\frac{b^2 - 4a}{p} \right) = -1$$

2. let f be a polynomial $x^2 - bx + a \in \mathbb{Z}_p$

3. Compute $r = x^{\frac{(p+1)}{2}} \pmod{f}$ \rightarrow r is an integer in \mathbb{Z}_p

4. return $(r, -r)$

not necessary for this
course