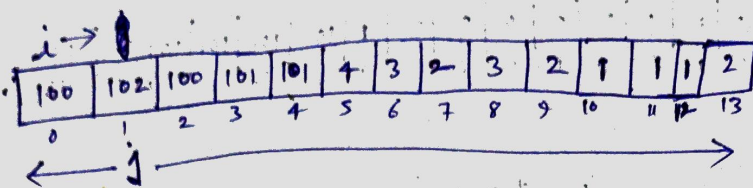


# Longest Consecutive Sequence

↳ Same as "Longest Subsequence of consecutive elements"



## Brute Force :-

1)  $i=0$   
 $current = 100 // arr[i]$   
 $tempLength = 1$      $maxLength = -\infty$

2) Find 101 in the entire array

↳ Found. So,  $current = 101$   
 $tempLength = 2$

3) Find 102 in the entire array.

↳ Found. So,  $current = 102$   
 $tempLength = 3$

4) Find 103 in the entire array.

↳ Not Found. So,  $if (maxLength < tempLength) \{$   
 $maxLength = tempLength;$   
 $\}$

5)  $i=1$      $i++;$   
 $current = 102 // arr[i]$   
 $tempLength = 1$

6) Find 103 in the entire array.

↳ Not Found.  $if (maxLen > tempLen) \{$   
 $maxLen = tempLen;$   
 $\}$

7)  $i=2$      $i++;$   
 $current = 100$   
 $tempLength = 1$

So on....

Here, we are iterating over the array and finding the longest consecutive sequence starting with  $arr[i]$  and keep on noting its length. At last whatever the maximum length we have that is our answer.  $T = O(n) * O(n^2) = O(n^3)$

• Better Solution :-

Sort the array :-

1	1	1	2	2	2	3	3	4	100	100	101	101	102
0	1	2	3	4	5	6	7	8	9	10	11	12	13

current = 1

templen = 1

maxlen = -∞

~~repeat (i=0, i++)~~

for (i=0; i <= n-1; i++) {

if (arr[i] == current + 1) {  
current = current + 1;  
templen ++;

} else if (arr[i] == current) {  
continue;

} else {

if (maxlen < templen) {  
maxlen = templen;

}  
current = arr[i];  
templen = 1;

}

}

if (maxlen < templen) {  
maxlen = templen;

}

$$T = O(n \cdot \log n) + O(n)$$

$$S = O(1)$$

• Best Solution : -

<i>i</i>													
100	102	100	101	101	4	3	2	3	2	1	1	1	2
0	1	2	3	4	5	6	7	8	9	10	11	12	13

1) Iterate and put elements in hashMap

2) Then, let us iterate over the array again with  $i=0$ .

if  $((100-1)$  present in map)

⇒ No ⇒ Then 100 is starting point ⇒ (100, F)

⇒ Then check

3  
marken

100  
↓  
101  
↓  
102  
↓  
103 X

Now, put (100, true) back in map.  
 $i++$

3)  $i=1$

if  $((102-1)$  is present in map)

⇒ Yes ⇒  $i++$

4)  $i=2$

if  $((100-1)$  is present in map)

⇒ No ⇒ Then 100 is starting point ⇒ But (100, true)  
⇒  $i++$

(102, F)	
(101, F)	
(100, F)	(100, T)
(4, F)	
(3, F)	
(2, F)	
(1, F)	

HashMap

(Integer, Boolean)

↓  
Eg.

(100, false)

↓  
It represents length of sequence starting with 100 is not yet found.

5)  $i=3$

if  $((i-1)$  present in map)

$\hookrightarrow$  Yes  $\Rightarrow i++$

6)  $i=4$

if  $((i-1)$  present in map)

$\hookrightarrow$  Yes  $\Rightarrow i++$

7)  $i=5$

if  $((i-1)$  present in map)

$\hookrightarrow$  Yes  $\Rightarrow i++$

8)  $i=6$

⋮  
Similarly

12)  $i=10$

if  $((i-1)$  present in map)

$\hookrightarrow$  No  $\Rightarrow$  Then 1 is starting point

$\downarrow$   
(1, F)

$\downarrow$   
Then check

1  $\downarrow$

2  $\downarrow$

3  $\downarrow$

4  $\downarrow$

5  $\downarrow$   
 $\times$  (5)

4
---

  
maxLen

Then put (1, true) back in the map  
 $i++$



13)  $i = 11$

if  $((i-1))$  is present in map

$\Rightarrow$  No  $\Rightarrow$  Then 1 is starting point

$\Downarrow$   
 $(1, T)$   
 $\Downarrow$   
 $i++$

14)  $i = 12$

Same as (13)

15)  $i = 13$

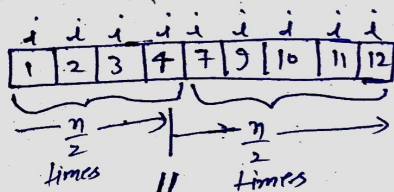
if  $((i-1))$  is present in map

$\Rightarrow$  Yes  $\Rightarrow i++$

So,  $maxLength = 4$

$T = O(n) + O(2n) = O(3n) = O(n)$

$\Downarrow$   
 You have to think logically



$\Downarrow$   
 $O(n) + O(\frac{n}{2}) + O(\frac{n}{2})$   
 $= O(n) + O(n)$   
 $= O(2n)$

So,  $T = O(n) \quad S = O(n)$