

DAY 4 Theory of Computation

DAY 4 CLASS (23/09/2020, 9AM -10 AM)

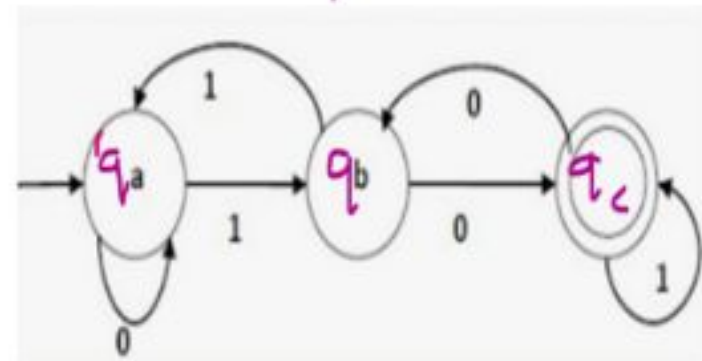
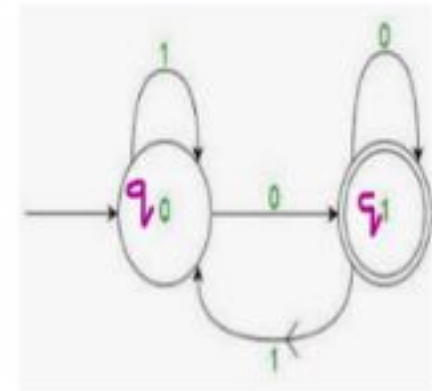
1. Types of Finite Automaton
2. DFA (Deterministic finite Automaton)
3. NFA (Non Deterministic Finite Automaton)
4. Acceptability by NFA.

Types of FA(Finite Automaton)

- 1.DFA (Deterministic FA)
- 2.NFA (Non Deterministic FA)

If every state is applied with all possible inputs in TD, DFA

Example of DFA

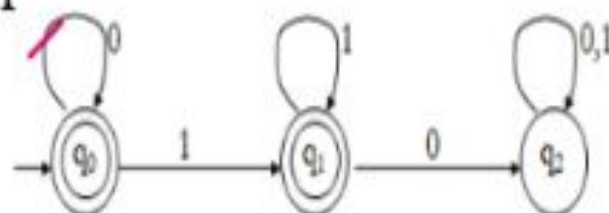


DFA(Deterministic Finite Automaton)

Wednesday, September 23, 2020 9:30 AM

Example

All symbols applied to all states, DFA



DFA
Example

alphabet $\Sigma = \{0, 1\}$
start state $Q = \{q_0, q_1, q_2\}$
initial state q_0
accepting states $F = \{q_0, q_1\}$

transition function δ :

	inputs	
	0	1
states	q_0	q_1
	q_2	q_1
	q_2	q_2

filled
cell

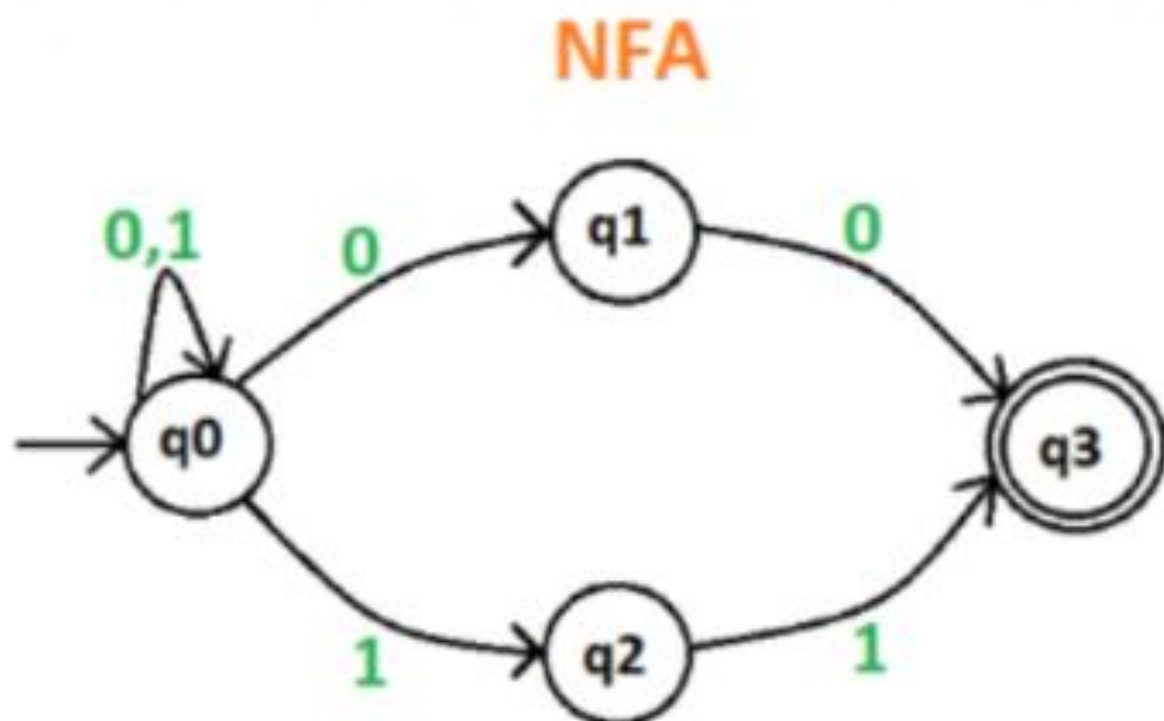
Non Deterministic Finite Automata (NFA):

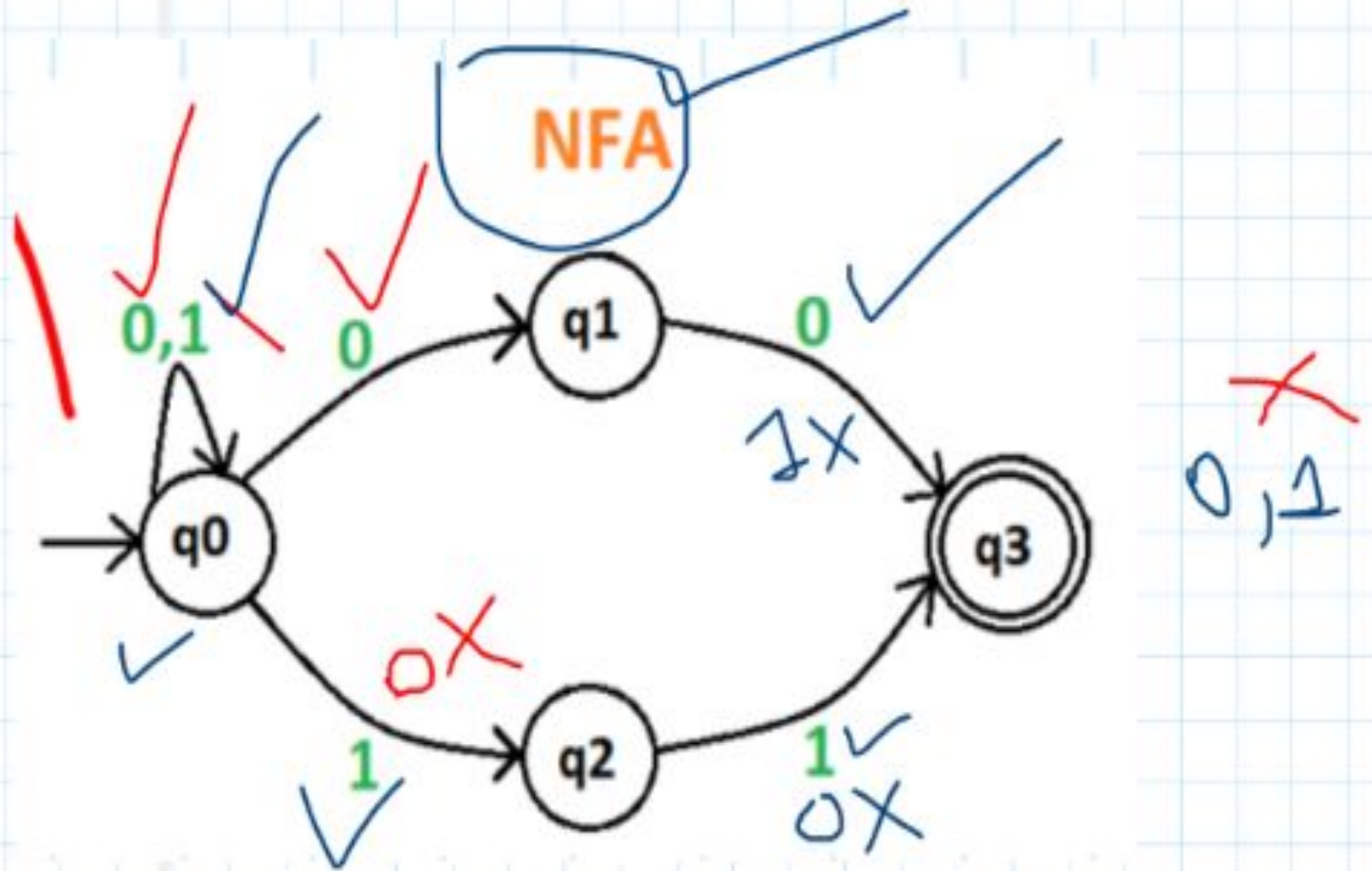
Like the DFA, a NFA has a finite set of states, a finite set of input symbols, one start state, a set of accepting states and transition function (δ).

The difference between the DFA and the NFA is in the type of δ . For the NFA, δ is a function that takes a state and input symbol as arguments as like the DFA's transition function, but returns a set of **zero, one or more state (DFA returns exactly one state)**.

So, we can conclude in a NFA, there **may be zero no. of arcs** out of each state for each input symbol. While in DFA, it **has exactly one arc** out of each state for each input symbol.

Example of NFA





Nondeterministic Finite Automaton

$$M = (Q, \Sigma, \delta, q_0, F)$$

where

$$Q = \{q_0, q_1, q_2, q_3, q_4\}$$

$$\Sigma = \{0, 1\}$$

$$F = \{q_2, q_4\}$$

δ	0	1
q_0	$\{q_0, q_3\}$	$\{q_0, q_1\}$
q_1	\emptyset	$\{q_2\}$
q_2	$\{q_2\}$	$\{q_2\}$
q_3	$\{q_4\}$	\emptyset
q_4	$\{q_4\}$	$\{q_4\}$

More Than One entry

one entry

No entry



Non-deterministic Finite Automata (NFA)

- A Non-deterministic Finite Automaton (NFA) consists of:
 - $Q \Rightarrow$ a finite set of states
 - $\Sigma \Rightarrow$ a finite set of input symbols (alphabet)
 - $q_0 \Rightarrow$ a start state
 - $F \Rightarrow$ set of accepting states
 - $\delta \Rightarrow$ a transition function, which is a mapping between $Q \times \Sigma \Rightarrow$ subset of Q
- An NFA is also defined by the 5-tuple:
 - $\{Q, \Sigma, q_0, F, \delta\}$

Why Subset of Q

as incase of NFA there may be zero, one or more number next states for same input symbol and

Let $Q = \{q_0, q_1, q_2\}$

Subset of $Q = \{\text{null}, \{q_0\}, \{q_1\}, \{q_2\}, \{q_0, q_1\}, \{q_1, q_2\}, \{q_0, q_2\}, \{q_0, q_1, q_2\}\}$

1. No next state: null

2. One next state: q_0 / q_1 / q_2

3. More than one next state for same input symbol like q_0 and q_1 , q_1 and q_2 , q_0 and q_1 and q_2

Non Deterministic Finite Automata (NFA):

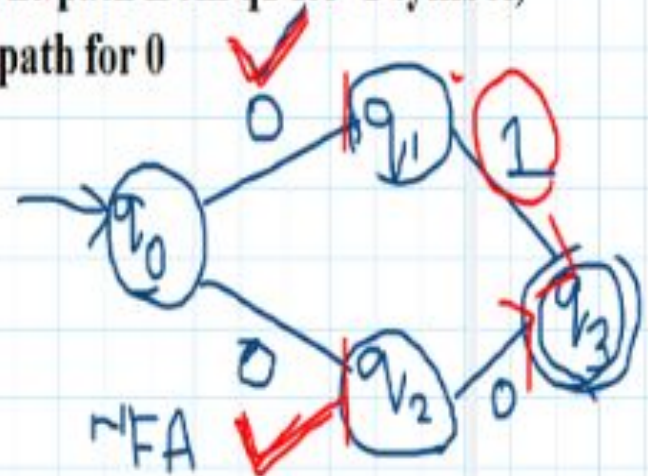
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NFA, Multiple path for same input

1. One path from q_1 for 1 symbol, no path for 0



1. For state q_0 , same input 0 is applied twice, there are more than one next state (q_1, q_2)

3. no next state from previous state q_3

The string $w = 0110$ is accepted by the NFA through the path $q_0 \rightarrow q_0 \rightarrow q_0 \rightarrow q_3 \rightarrow q_4$

