



Theorems: Undirected Graphs

Theorem 1

The Handshaking theorem:

$$2e = \sum_{v \in V} \deg(v)$$

(why?) Every edge connects 2 vertices

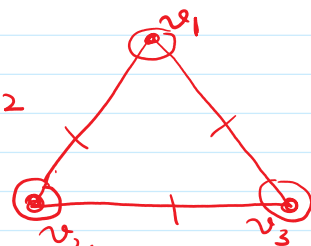


Hand-shaking Theorem.

$$\sum_{i=1}^n \deg(v_i) = 2e$$

$$\deg(v_1) = 2, \deg(v_2) = 2$$

$$\deg(v_3) = 2$$



$$\deg(v_1) + \deg(v_2) + \deg(v_3) = 2 + 2 + 2 = 6 = 2 \times 3$$

$$\sum_{i=1}^3 \deg(v_i) = 2 \times \text{no. of edges} = 2 \times e$$

verify Hand-shaking for the above graph

$$\deg(v_1) = 3, \deg(v_2) = 2, \deg(v_3) = 3, \deg(v_4) = 2$$

$$\deg(v_1) + \deg(v_2) + \deg(v_3) + \deg(v_4)$$

$$= 3 + 2 + 3 + 2$$

$$= 10 = 2 \times 5 = 2e$$



Theorems: Undirected Graphs

Theorem 2:

An undirected graph has even number of vertices with odd degree

Proof: V_1 is the set of even degree vertices and V_2 refers to odd degree vertices

$$2e = \sum_{v \in V} \deg(v) = \sum_{u \in V_1} \deg(u) + \sum_{v \in V_2} \deg(v)$$

$$\Rightarrow \deg(v) \text{ is even for } v \in V_1$$

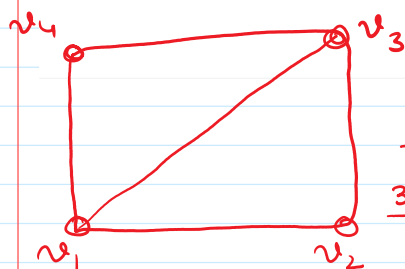
\Rightarrow The first term in the right hand side of the last inequality is even.

\Rightarrow The sum of the last two terms on the right hand side of the last inequality is even since sum is $2e$.

Hence second term is also even

$$\Rightarrow \text{second term } \sum_{v \in V_2} \deg(v) = \text{even}$$

$$3 + 4 + 6 = \text{even}$$



$$3 + 5 + 7 = 15$$

$$3 + 5 + 7 + 9 = 24$$

$$V_1 = \{v_1, v_3\} \checkmark$$

$$V_2 = \{v_2, v_4\} \checkmark$$

@ No. of odd degree vertices in any graph are always even.

Solⁿ Suppose that in any graph G there are n vertices and e edges.

\therefore By Hand-shaking Theorem.

$$\sum_{i=1}^n \deg(v_i) = 2e$$

$$\sum_{v_i \in \text{odd}} \deg(v_i) + \sum_{v_j \in \text{even}} \deg(v_j) = 2e$$

$$\sum_{v_i \in \text{odd}} \deg(v_i) + \text{even} = \text{even}$$

$$\sum_{v_i \in \text{odd}} \deg(v_i) = \text{even} - \text{even} = \text{even}$$

This situation is possible only when no. of odd degree vertices are even.

Is it possible to have a graph with the following degree of vertices:

$$\underbrace{1}_{\text{odd}}, \underbrace{2}_{\text{even}}, \underbrace{3}_{\text{odd}}, \underbrace{4}_{\text{even}}, \underbrace{5}_{\text{odd}}$$

This is not possible.

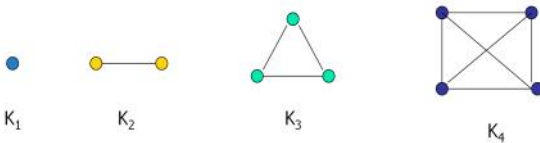
if vertices, $\xrightarrow{\text{odd}}$, $\xrightarrow{\text{odd}}$, $\xrightarrow{\text{odd}}$

This is not possible.

Simple graphs – special cases

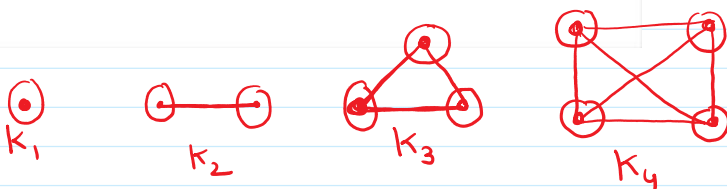
- Complete graph:** K_n is the simple graph that contains exactly one edge between each pair of distinct vertices.

Representation Example: K_1, K_2, K_3, K_4



Complete Graph: If \exists an edge between every pair of vertices then that graph is called a complete graph. A complete graph is denoted by K_n .

Where n represents no. of vertices

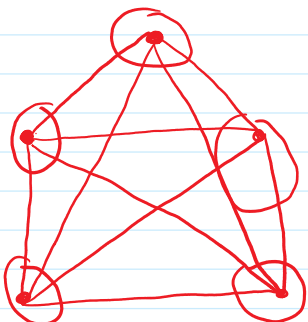


In K_n graph degree of each vertex = $n-1$

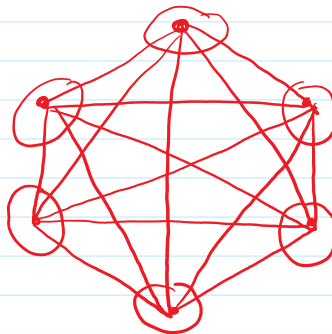
In K_{100} graph what is degree of each vertex?

- (a) 98 (b) 99 (c) 100 (d) 101

Q Draw K_5 and K_6 graph.



K_5



K_6

Simple graphs – special cases

- Cycle:** C_n , $n \geq 3$ consists of n vertices $v_1, v_2, v_3 \dots v_n$ and edges $\{v_1, v_2\}, \{v_2, v_3\}, \{v_3, v_4\} \dots \{v_{n-1}, v_n\}, \{v_n, v_1\}$

Representation Example: C_3, C_4

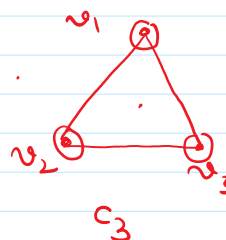


C_3

C_4

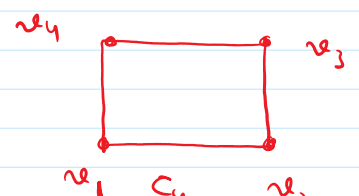
Cycle C_n , $n \geq 3$

$\{v_1, v_2\}, \{v_2, v_3\}, \dots, \{v_{n-1}, v_n\}, \{v_n, v_1\}$



C_3

$\deg(v_i) = 2$



C_4

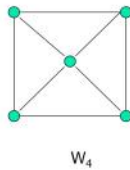
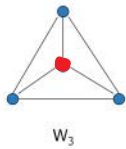
$\deg(v_i) = 2$

Draw C_5, C_6, C_7 & C_8 cycle.

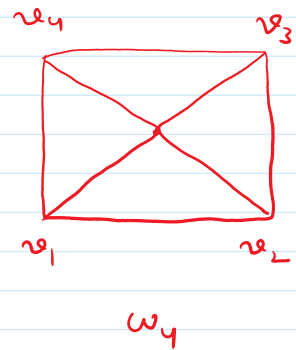
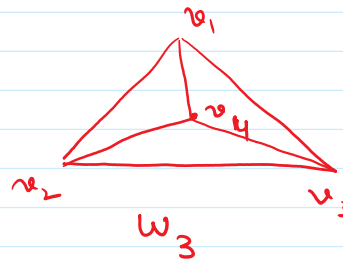
Simple graphs – special cases

- Wheels:** W_n obtained by adding additional vertex to C_n and connecting all vertices to this new vertex by new edges.

Representation Example: W_3, W_4



Wheel: W_n



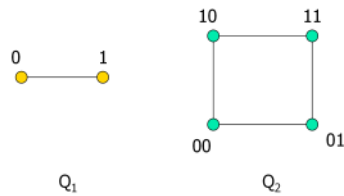
Draw W_5, W_6, W_7, W_8 .



Simple graphs – special cases

- **N-cubes:** Q_n , vertices represented by $2n$ bit strings of length n . Two vertices are adjacent if and only if the bit strings that they represent differ by exactly one bit positions

Representation Example: Q_1 , Q_2

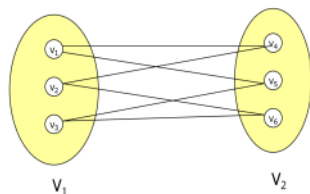


Bipartite graphs

- In a simple graph G , if V can be partitioned into two disjoint sets V_1 and V_2 such that every edge in the graph connects a vertex in V_1 and a vertex V_2 (so that no edge in G connects either two vertices in V_1 or two vertices in V_2)

Application example: Representing Relations

Representation example: $V_1 = \{v_1, v_2, v_3\}$ and $V_2 = \{v_4, v_5, v_6\}$,



Bi-partite graph.

Two partitions of the vertex set

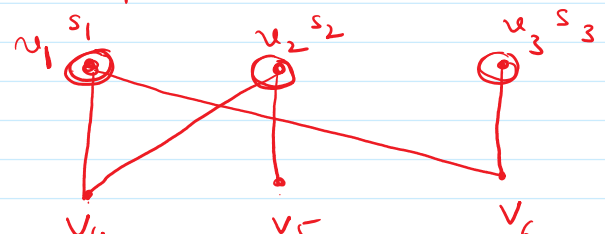
$$V = \{v_1, v_2, v_3, v_4, v_5, v_6\}$$

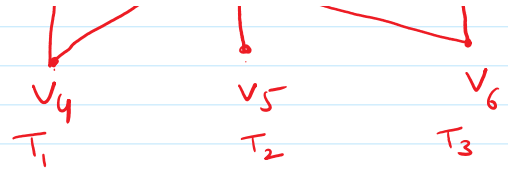
$$V_1 = \{v_1, v_2, v_3\}$$

$$V_1 \cap V_2 = \emptyset$$

$$V_2 = \{v_4, v_5, v_6\}$$

$$V_1 \cup V_2 = V$$





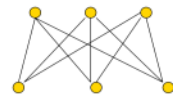
Complete Bipartite graphs

- $K_{m,n}$ is the graph that has its vertex set partitioned into two subsets of m and n vertices, respectively. There is an edge between two vertices if and only if one vertex is in the first subset and the other vertex is in the second subset.

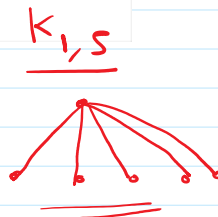
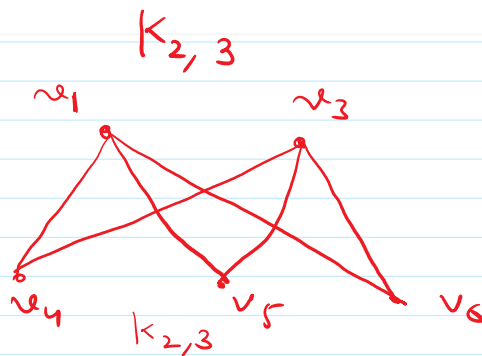
Representation example: $K_{2,3}$, $K_{3,3}$



$K_{2,3}$



$K_{3,3}$

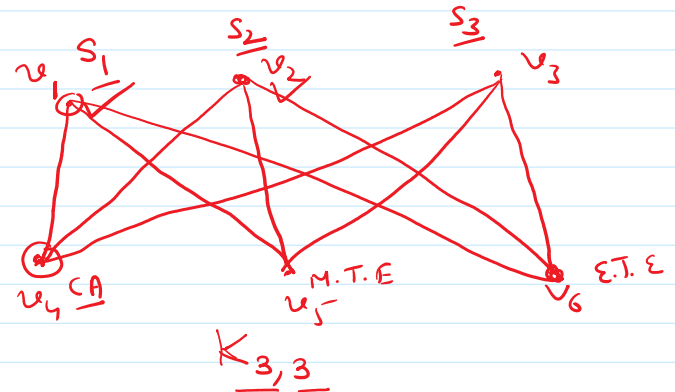


$$V = \{v_1, v_2, v_3, v_4, v_5, v_6\}$$

$$V_1 = \{v_1, v_2, v_3\} \checkmark$$

$$V_2 = \{v_4, v_5, v_6\} \checkmark$$

$$V_1 \cap V_2 = \emptyset, \quad V_1 \cup V_2 = V$$



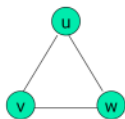


Subgraphs

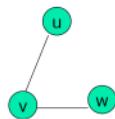
- A subgraph of a graph $G = (V, E)$ is a graph $H = (V', E')$ where V' is a subset of V and E' is a subset of E

Application example: solving sub-problems within a graph

Representation example: $V = \{u, v, w\}$, $E = \{\{u, v\}, \{v, w\}, \{w, u\}\}$, H_1 , H_2



G



H_1

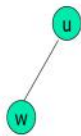


H_2

Subgraphs

- $G = G1 \cup G2$ wherein $E = E1 \cup E2$ and $V = V1 \cup V2$, G , $G1$ and $G2$ are simple graphs of G

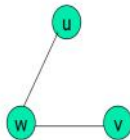
Representation example: $V1 = \{u, w\}$, $E1 = \{\{u, w\}\}$, $V2 = \{w, v\}$, $E2 = \{\{w, v\}\}$, $V = \{u, v, w\}$, $E = \{\{u, w\}, \{w, v\}\}$



G1



G2



G