

**Subject** : **Digital Communication**  
**Code** : **UEC 639**  
**Credit** : **4**

**Dr. Amit Mishra**

# Modulation

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- Motivation
- Many modulation formats
  - PSK, FSK, ASK, QAM
- Important criteria
  - Spectral efficiency
  - BER vs. SNR
  - Complexity (Cost)

# Demodulation / Detection

- Two types
  - Coherent detection
    - ❑ phase of the carrier known to the receiver
    - ❑ optimal receptor: correlator
    - ❑ Gaussian noise
  - Noncoherent detection
    - ❑ carrier phase unknown at receiver
    - ❑ optimal receiver : power detection (envelope detection)
    - ❑ non-Gaussian noise

# Carrier

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- Base band
  - no carrier
  - positive frequencies bandwidth
- Pass band
  - with carrier
  - bandwidth twice the base band

# Modulation

- Reasons for carrier modulation
  - Antennas, FDMA, etc.
- Signal space analysis
  - Valid for carrier modulation
  - Basis vectors take into account modulation
  - Often  $\psi_1 = \cos(\omega_0 t)$        $\psi_2 = \sin(\omega_0 t)$ 
    - ❑ Parties **in-phase** et **quadrature**



# Signal space

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- Binary case
  - Signal space - with good normalization
  - Distance between the two bits
  - Probability of error
- M-ary Case
  - Signal space
  - Distance between symbols

**NOTE:** Derive the amplitude value of a transmitted signal in terms of its bit Energy.

$$x(t) = A \cos 2\pi f_c t, \quad 0 \leq t \leq T_b$$

$E_b \leftarrow$  Energy per bit

$$E_b = \int_0^{T_b} x^2(t) dt$$

$$= \int_0^{T_b} A^2 \cos^2 2\pi f_c t dt$$

$$= \frac{A^2}{2} \int_0^{T_b} (1 + \cos 4\pi f_c t) dt$$

$$= \frac{A^2}{2} \int_0^{T_b} (1 + \cos 4\pi f_c t) dt$$

$$= \frac{A^2}{2} T_b + 0$$

$$E_b = \frac{A^2 T_b}{2}$$

$$\boxed{\sqrt{\frac{2E_b}{T_b}} = A}$$

$$\cos^2 \theta = \frac{1 + \cos 2\theta}{2}$$

Therefore,

$$x(t) = \sqrt{\frac{2E_b}{T_b}} \cos(2\pi f_c t)$$

2 dimensional space no matter how many symbols you have

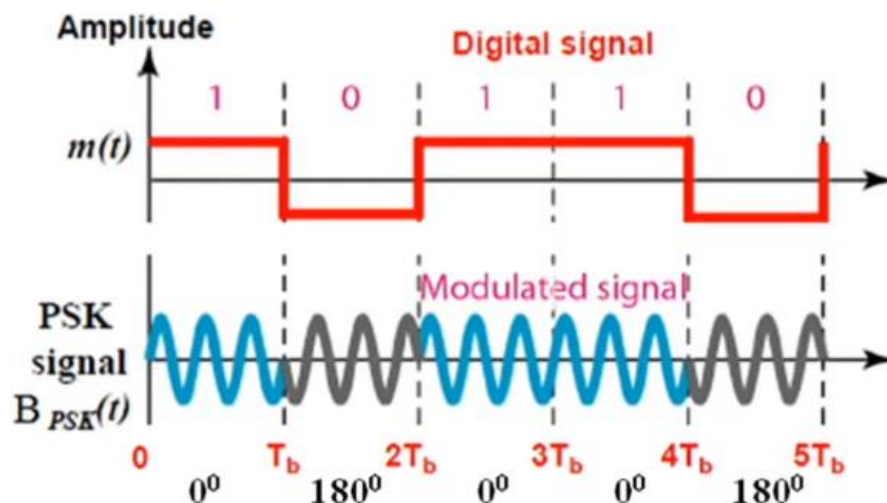
# PSK

- Phase shift keying  $s(t) = \sqrt{\frac{2E}{T}} \cos(\omega_0 t + \phi_i(t))$   $0 \leq t \leq T$  *data*
- Typically  $\phi_i = \frac{2\pi(i-1)}{M}$   $i = 1, 2, \dots, M$
- $E$  energy/symbol,  $T$  symbol time
- Typically  $\psi_1 = \cos(\omega_0 t)$   $\psi_2 = \sin(\omega_0 t)$
- Example: Antipodal (BPSK)

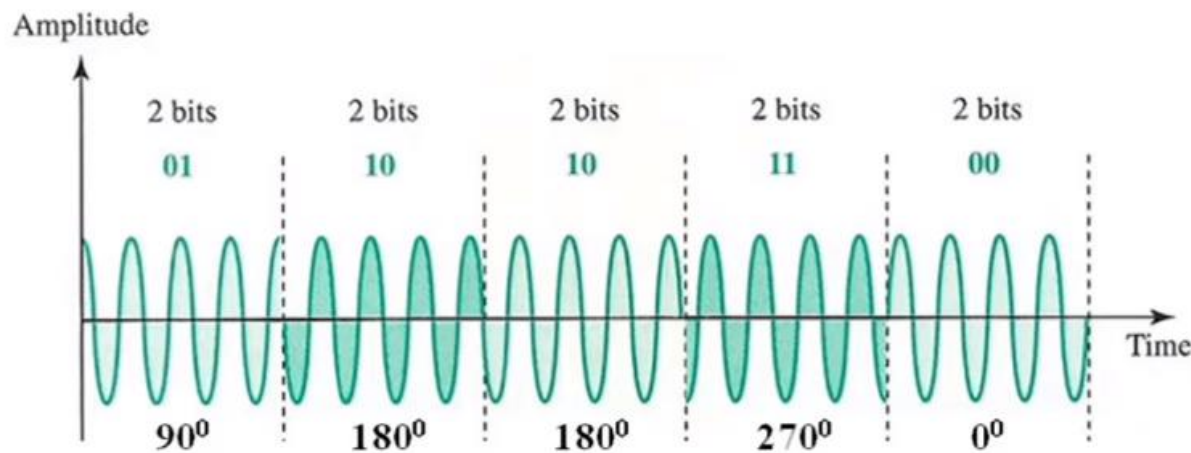


- In M-ary Phase Shift Keying (M-ary PSK),  $M$  different phase angles are used to represent  $M$  symbols.
- The signal is represented by:  

$$s_i(t) = A \cos(2\pi f_c t + \phi_i), \quad 0 \leq t \leq T_s \text{ for } i = 1, 2, \dots, M$$
 where  $\phi_i = \frac{2\pi}{M}(i - 1) + \text{constant}$ , for  $i = 1, 2, \dots, M$
- $M = 4$  and the constant  $= 0$ , then four phases are  $0, \pi/2, \pi, 3\pi/2$ .  
 $M = 4$  and the constant  $= \pi/4$ , then four phases are  $\pi/4, 3\pi/4, 5\pi/4, 7\pi/4$
- Both 4-PSK are called Quadrature Phase Shift Keying or QPSK



BPSK : 2 Level  
 1 bit in 1 symbol duration



M-PSK : 4 Level

2 bit in 1 symbol duration

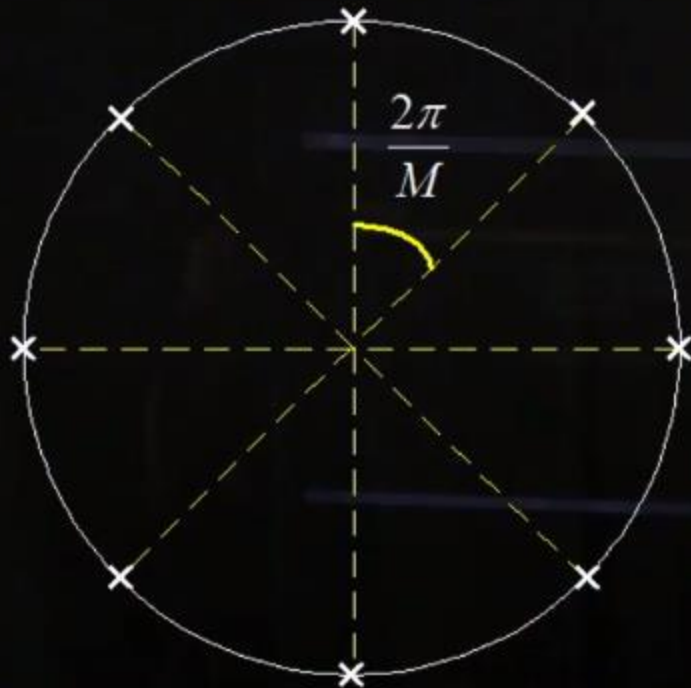
Fig. 6: BPSK and Mary-PSK Signal Representation

- Bandwidth of MPSK ,  $B = \frac{2}{T_s}$  Hz
- Bandwidth Efficiency ,  $\rho = \frac{\log_2 M}{2}$  Bits/ s /Hz

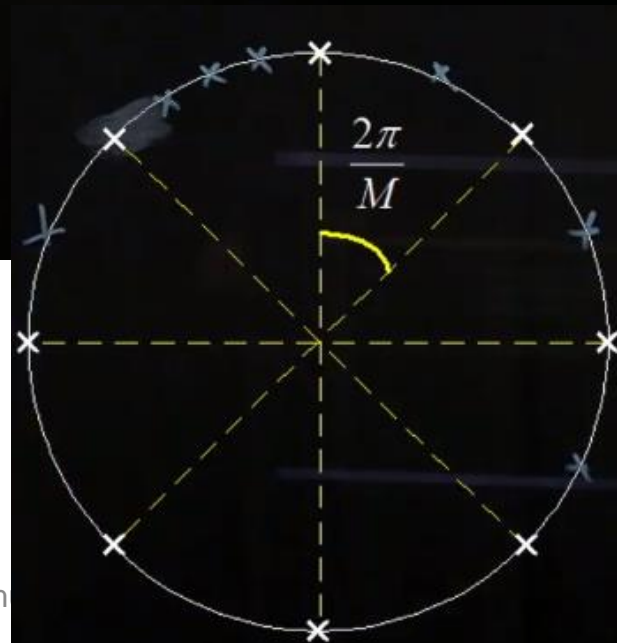
As M increases, Bandwidth Efficiency of Mary PSK increases

- Uses constant envelope so immune to noise.
- There is considerable reduction in bandwidth requirement.
- Better performance than ASK and FSK.
- Increase in probability of errors with increase in number of bits per symbol.
- Design of M-ary PSK modulator and demodulator is complex.

# M-PSK



- Same distance from origin for the entire constellation
- Impact of larger  $M$ ???



# FSK

- Frequency shift keying  $s_i(t) = \sqrt{\frac{2E}{T}} \cos(\omega_i t + \theta) \quad 0 \leq t \leq T$
- $\omega_i$  are  $M$  distinct frequencies  
 $E$  energy/symbol,  $T$  symbol time
- Typically  $\left\{ \psi_i = \cos(\omega_i t) \right\}_{i=1}^M \quad \theta = 0$
- Typically  $M$  is a power of 2



- In M-ary Frequency Shift Keying (M-ary FSK), there are  $M$  different carrier frequencies are used to represent  $M$  symbols.

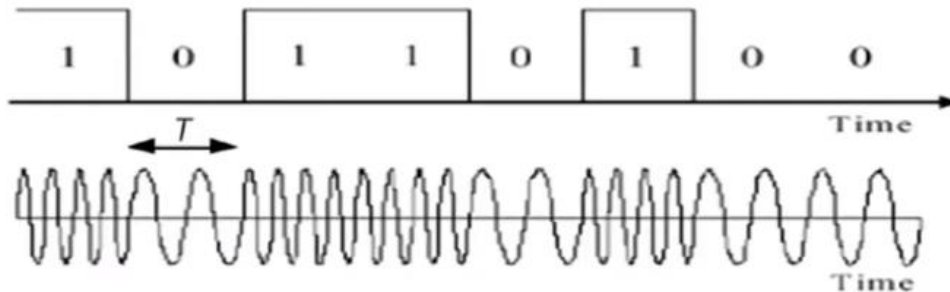
- The signal is represented by:

$$s_i(t) = A \cos(2\pi f_{ci}t), \quad 0 \leq t \leq T_s \text{ for } i = 1, 2, \dots, M$$

where  $M$  different frequencies are  $f_{ci} = f_c + \left(i - \frac{M}{2}\right) \Delta f$ , for  $i = 1, 2, \dots, M$

$\Delta f$  is difference between two adjacent frequencies.

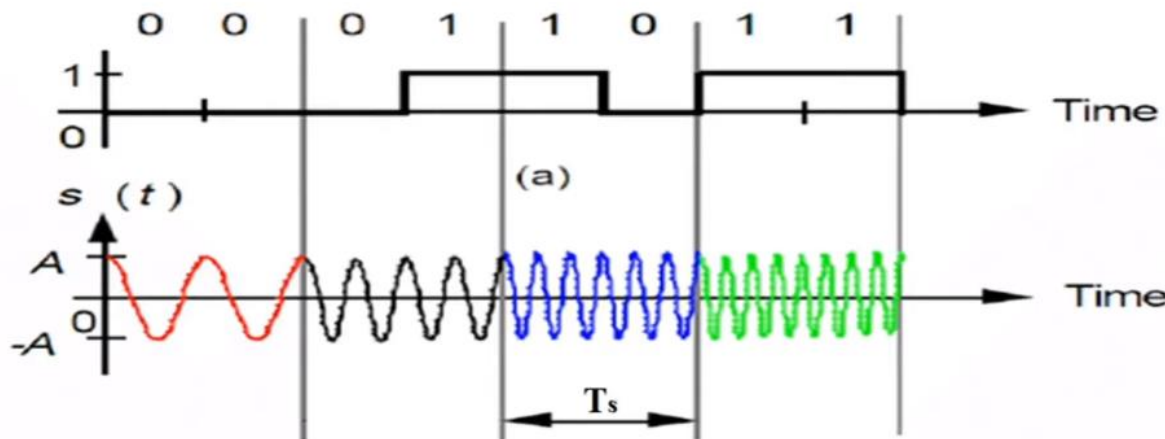
- Frequencies are chosen such that they are orthogonal to each other and there is no interference between adjacent carriers.



BFSK : 2 Level

1 bit in 1 symbol duration





M-FSK : 4 Level  
2 bit in 1 symbol duration

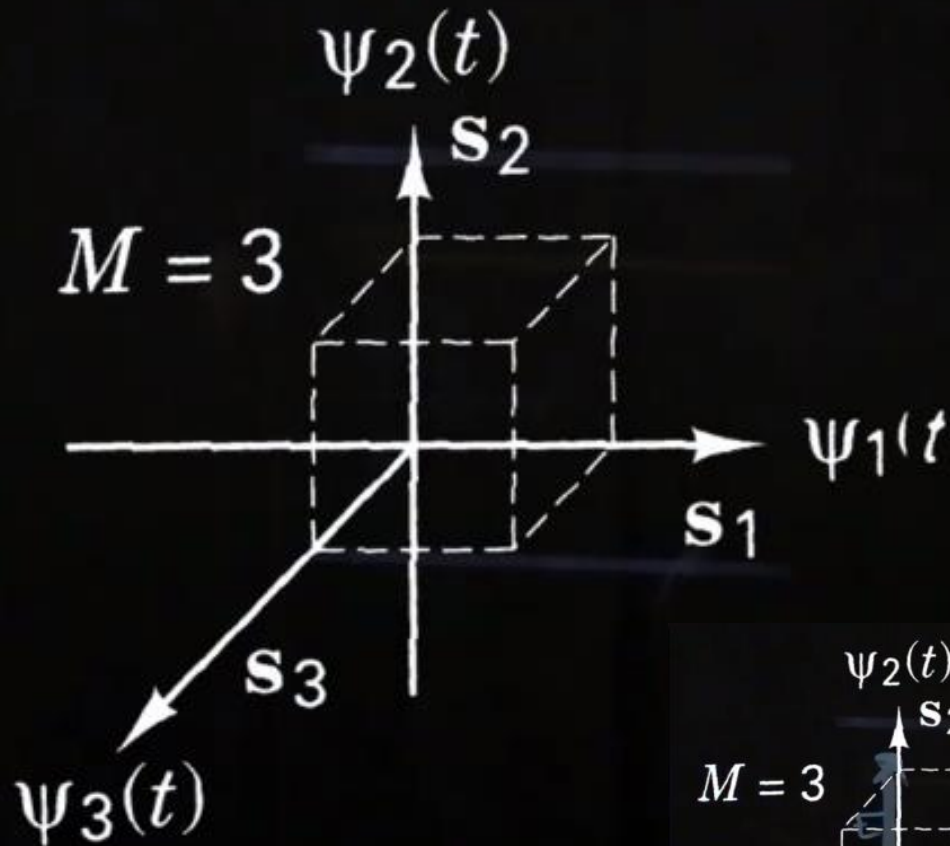
Fig. 4: BFSK and Mary-FSK  
Signal Representation

- Bandwidth of MFSK ,  $B = \frac{M}{2T_s}$  Hz
- Bandwidth Efficiency ,  $\rho = \frac{2\log_2 M}{M}$  Bits/ s /Hz

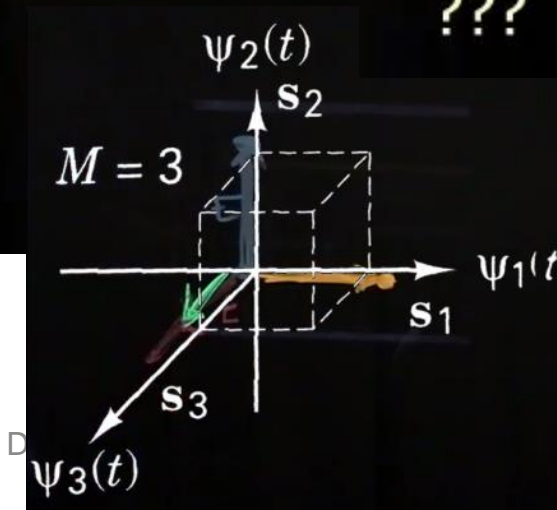
As M increases, Bandwidth Efficiency of MFSK decreases

- Better noise immunity than Mary-ASK.
- The transmitted M number of signals are equal in energy and duration.
- The signals are separated such that they are orthogonal to each other and therefore there is no crowding in the signal space.
- The bandwidth efficiency of M-ary FSK decreases and the power efficiency increases with the increase in M.

# 3-FSK



- Number of frequencies  $M$
- Dimension of space is  $M$
- Impact of larger  $M$  ???



# ASK = PAM

- Amplitude shift keying

$$s_i(t) = \sqrt{\frac{2E_i}{T}} \cos \omega_c t \quad i = 1, 2, \dots, M$$

- Typically evenly spaced discrete levels
- Example: OOK
- Not high performance, so less popular

- The signal is represented by:

$$s_i(t) = A_i \cos(2\pi fct), \quad 0 \leq t \leq T_s \text{ for } i = 1, 2, \dots, M$$

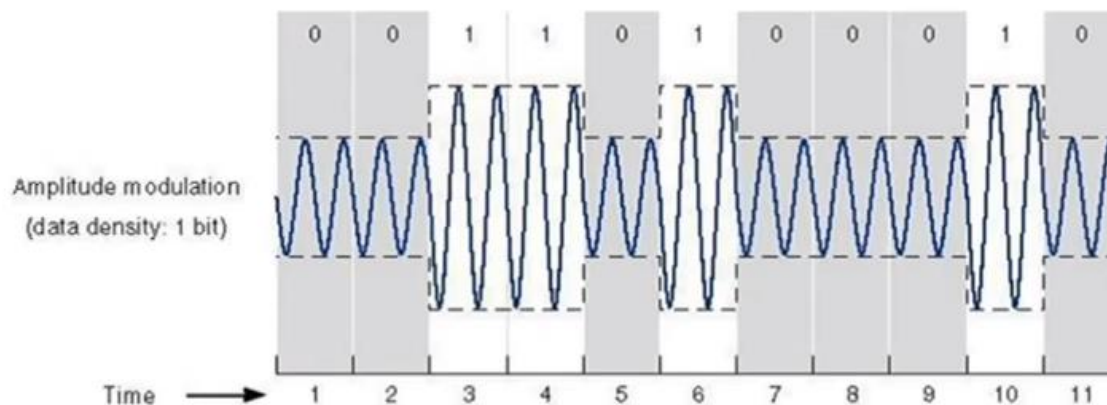
$A_i = (2i - 1 - M)d$ , where  $2d$  is the difference between two consecutive signal amplitudes.

- Let  $M = 4$  and  $d = 1$ , the four signal amplitudes will be  $-3, -1, 1$  and  $3$  V.

The M-ASK signals will be: for  $0 \leq t \leq T_s$

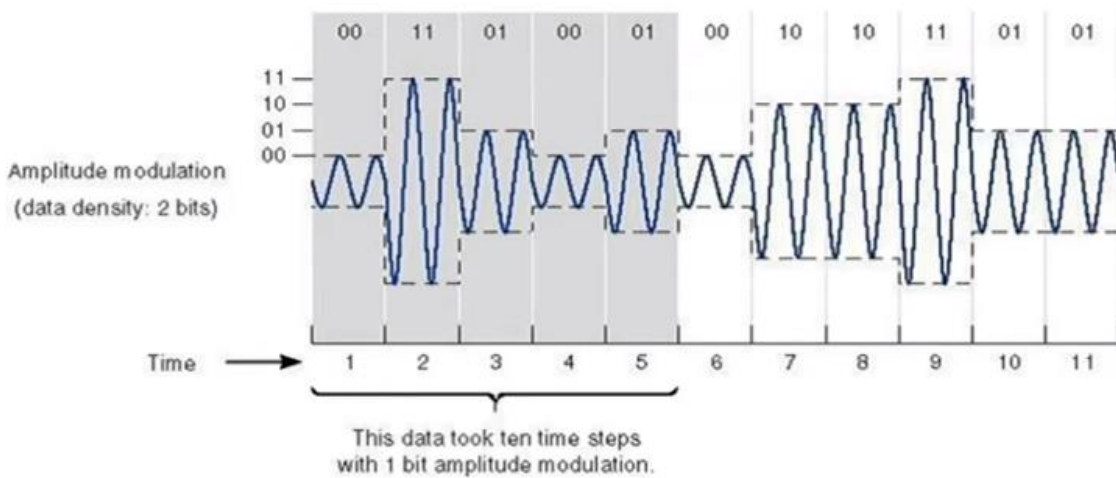
$$s_1(t) = \cos(2\pi fct), \quad s_2(t) = -\cos(2\pi fct),$$

$$s_3(t) = 3 \cos(2\pi fct), \quad s_4(t) = -3 \cos(2\pi fct).$$



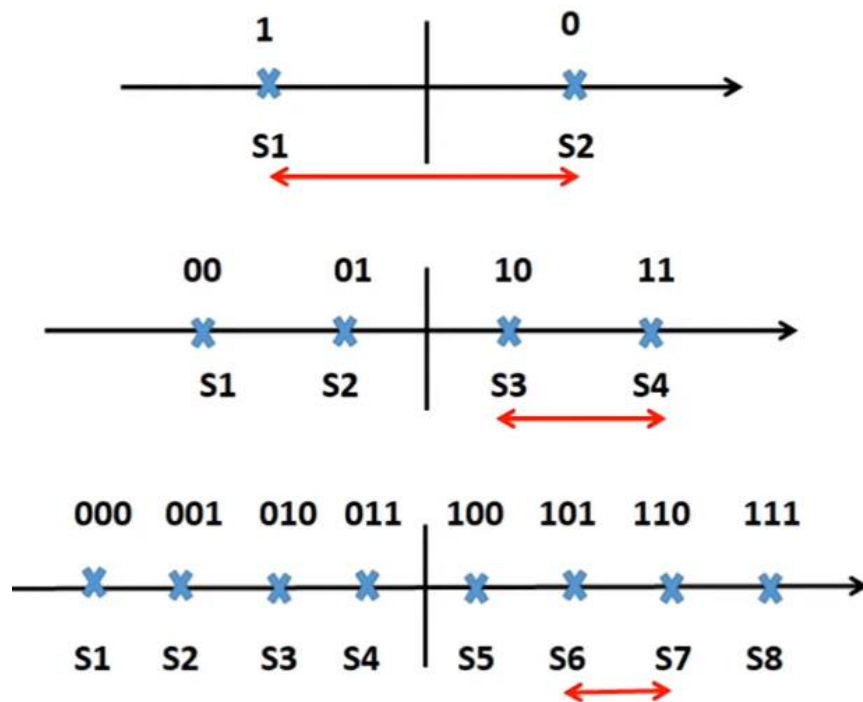
BASK : 2 Level  
1 bit in 1 symbol duration





M-ASK : 4 Level  
2 bit in 1 symbol duration

Fig. 2: BASK and Mary ASK Signal Representation



BASK :  $M=2$

4-ASK :  $M=4$

8-ASK :  $M=8$

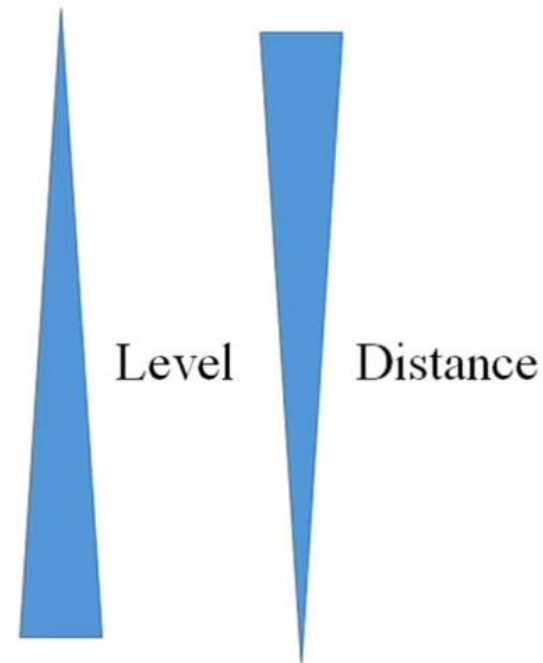


Fig. 3: Constellation Diagram of BASK and Mary-ASK Signal



2 dimensional space no matter how many symbols you have

# APK or QAM

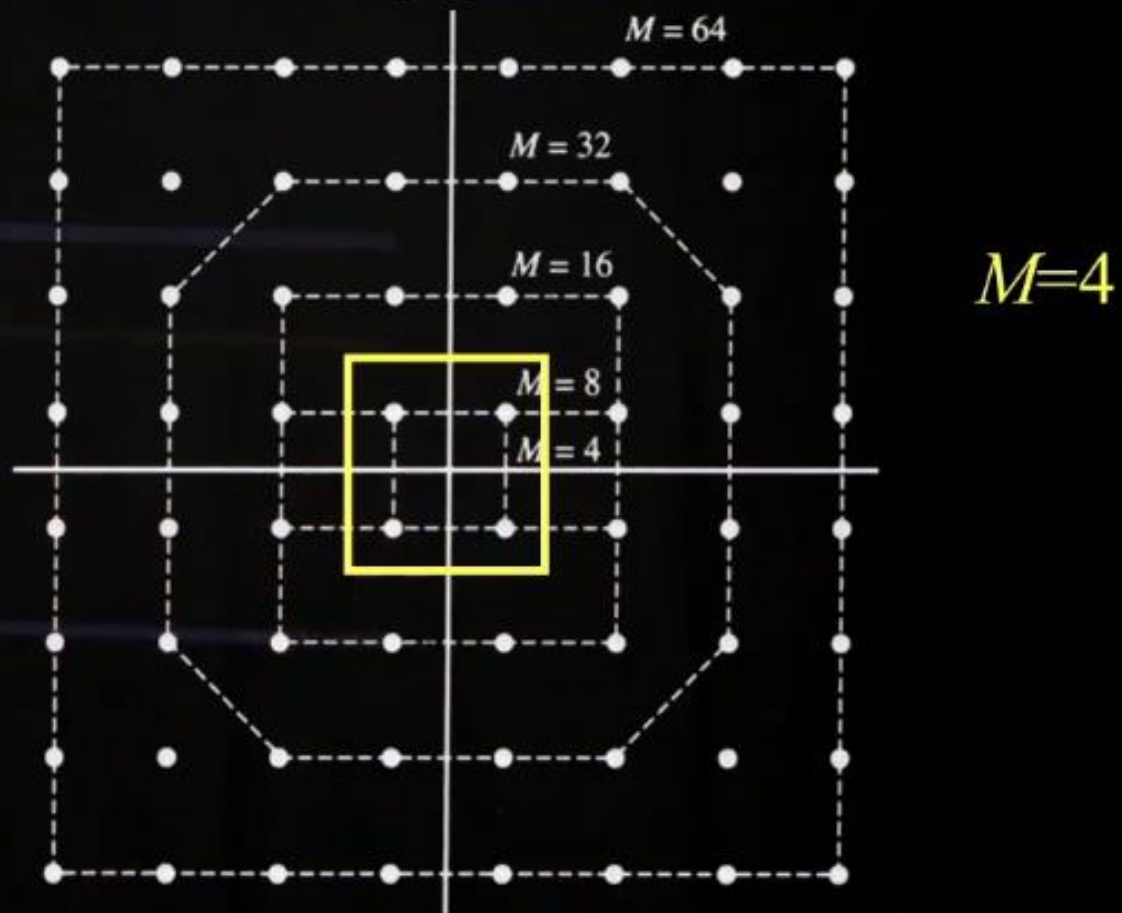
- Amplitude-phase modulation or Quadrature amplitude modulation

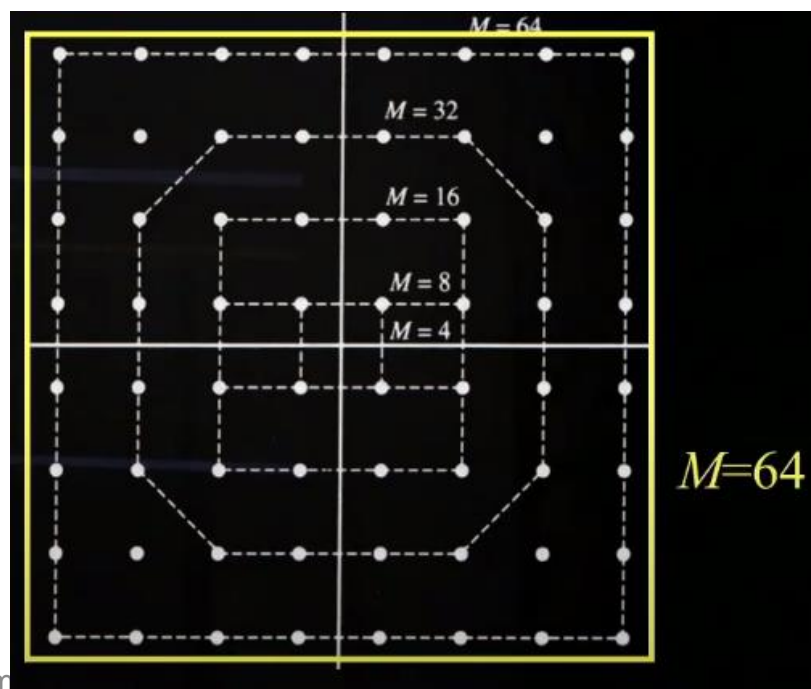
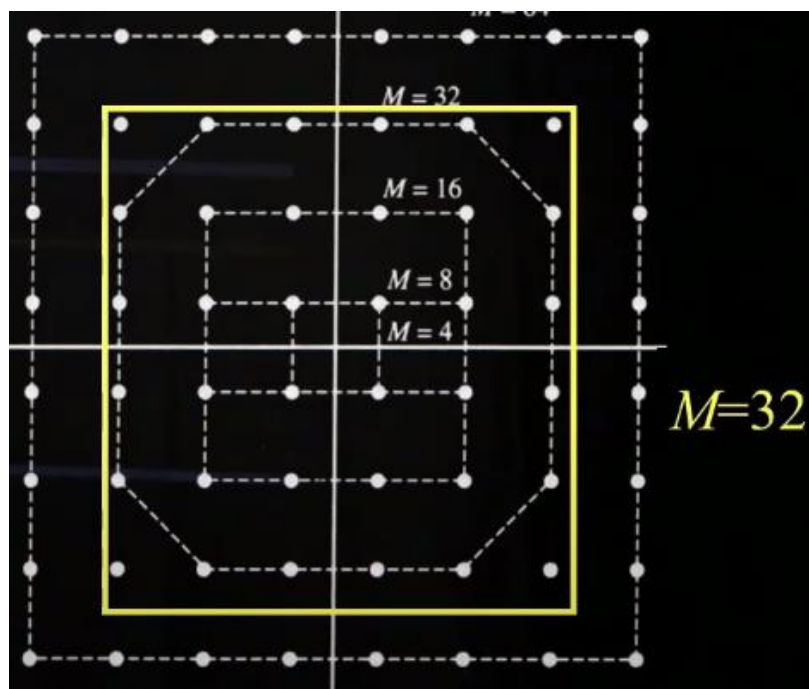
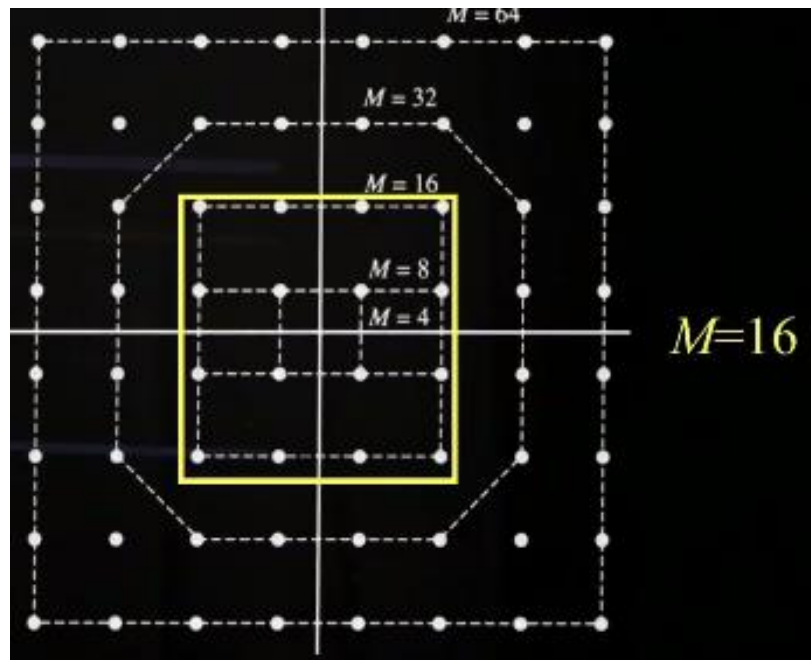
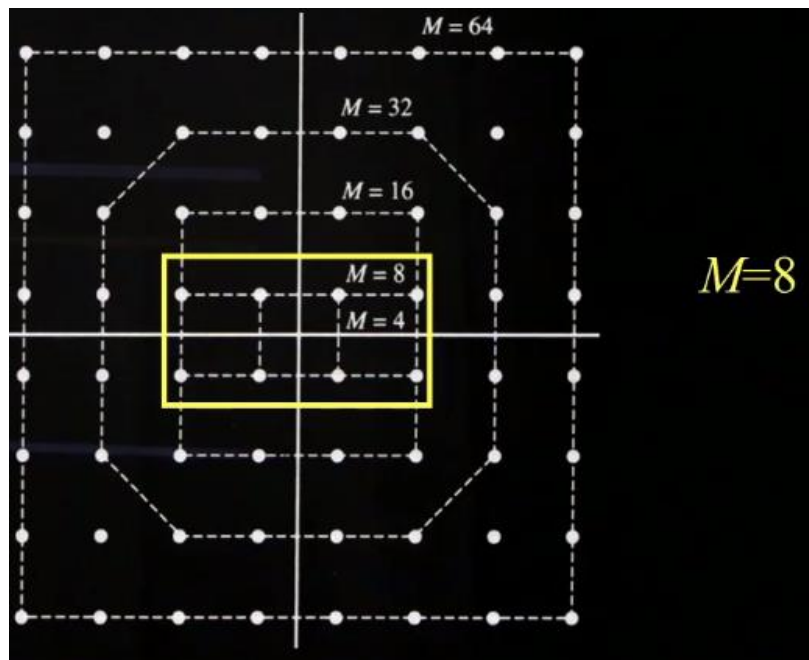
$$s_i(t) = \sqrt{\frac{2E_i}{T}} \cos(\omega_0 t + \phi_i(t)) \quad i = 1, 2, \dots, M$$

- Discrete levels
- Very BW efficient

unequal energies

# QAM





# Probability of Error for Binary case

## OOK vs. BPSK

- OOK:  $s_1(t)$  NRZ,  $s_0(t) = 0$  **ASK**

$$P_e = Q\left(\sqrt{\frac{E_b}{N_0}}\right)$$



- BPSK:  $s_1(t)$  NRZ,  $s_0(t) = -s_1(t)$  **Antipodal**

$$P_e = Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$$



**HOW ?**



# Noise... key for finding BER

- zero mean Gaussian

$$n_j \sim \left(0, \frac{N_0}{2}\right) \quad \sigma^2 = \frac{N_0}{2}$$

- Independent coefficients

$$En_j n_i = 0 \quad i \neq j$$

- Probability density

$$p_{\underline{n}}(\underline{n}) = \prod_{j=1}^{\infty} p_{n_j}(n_j) = \prod_{j=1}^{\infty} \frac{1}{\sqrt{2\pi\sigma}} e^{-n_j^2/2\sigma^2}$$

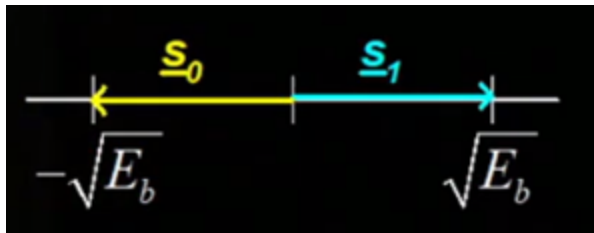


# Binary case

## ➤ Case 1 Antipodal **BPSK**

$$s_1(t) = \begin{cases} A & 0 \leq t \leq T \\ 0 & \text{ailleurs} \end{cases}$$

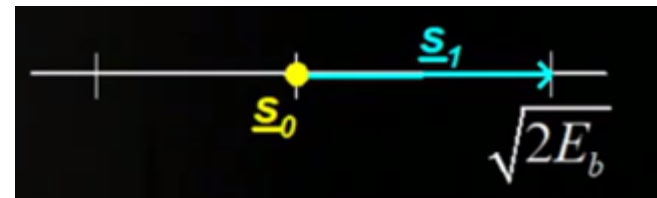
$$s_0(t) = -s_1(t)$$



## ➤ Case 2 On-Off Keying **ASK**

$$s_1(t) = \begin{cases} A & 0 \leq t \leq T \\ 0 & \text{ailleurs} \end{cases}$$

$$s_0(t) = 0$$



# Gram-Schmidt Process

## ➤ Case 1 Antipodal

$$\psi_1(t) = \frac{1}{\sqrt{E_1}} s_1(t)$$

$$E_1 = \int_0^T s_1^2(t) dt = \int_0^T A^2 dt = A^2 T$$

$$\psi_1(t) = \frac{1}{A\sqrt{T}} s_1(t) = \begin{cases} 1/\sqrt{T} & 0 \leq t \leq T \\ 0 & \text{ailleurs} \end{cases}$$

## ➤ Case 2 On-Off Keying

$$\psi_1(t) = \frac{1}{\sqrt{E_1}} s_1(t)$$

$$E_1 = \int_0^T s_1^2(t) dt = \int_0^T A^2 dt = A^2 T$$

$$\psi_1(t) = \begin{cases} 1/\sqrt{T} & 0 \leq t \leq T \\ 0 & \text{ailleurs} \end{cases}$$

# Signal coefficients

## ➤ Case 1 Antipodal

$$s_1(t) = \sqrt{E_1} \psi_1(t) = A\sqrt{T} \psi_1(t)$$

$$s_0(t) = -s_1(t) = -A\sqrt{T} \psi_1(t)$$

$$\underline{s}_1 = \begin{bmatrix} A\sqrt{T} \end{bmatrix} \quad \underline{s}_0 = \begin{bmatrix} -A\sqrt{T} \end{bmatrix}$$

## ➤ Case 2 On-Off Keying

$$s_1(t) = \sqrt{E_1} \psi_1(t) = A\sqrt{T} \psi_1(t)$$

$$s_0(t) = 0$$

$$\underline{s}_1 = \begin{bmatrix} A\sqrt{T} \end{bmatrix} \quad \underline{s}_0 = \begin{bmatrix} 0 \end{bmatrix}$$

*Signal space is of unit dimension*

# One dimensional space

➤ Received signal  $\underline{r} = [r]$

$$\begin{aligned} r &= \langle r(t), \psi_1(t) \rangle = \frac{1}{A\sqrt{T}} \langle r(t), s_1(t) \rangle \\ &= \frac{1}{A\sqrt{T}} \int_0^T r(t) s_1(t) dt \end{aligned}$$

➤ Correlator version

❑ Sampling receiver equivalent to matched filter

# Signal space

## ➤ Case 1 Antipodal

$$\underline{s}_1 = [A\sqrt{T}] \quad \underline{s}_0 = [-A\sqrt{T}]$$



## ➤ Case 2 On-Off Keying

$$\underline{s}_1 = [A\sqrt{T}] \quad \underline{s}_0 = [0]$$



*A fair comparison takes a normalization by the average energy of a bit*



# Normalization of $E_b$

## ➤ Case 1 Antipodal

$$E_1 = \int_0^T s_1^2(t) dt = \int_0^T A^2 dt = A^2 T$$

$$E_0 = \int_0^T s_0^2(t) dt = \int_0^T A^2 dt = A^2 T$$

$$E_b = \frac{1}{2}E_1 + \frac{1}{2}E_0 = A^2 T$$

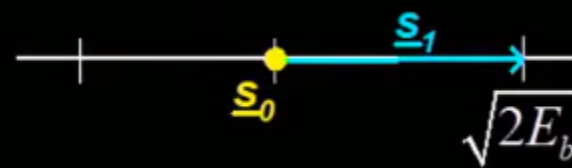


## ➤ Case 2 On-Off Keying

$$E_1 = \int_0^T s_1^2(t) dt = \int_0^T A^2 dt = A^2 T$$

$$E_0 = \int_0^T s_0^2(t) dt = 0$$

$$E_b = \frac{1}{2}E_1 + \frac{1}{2}0 = A^2 T / 2$$



# Distance as a function on $E_b$

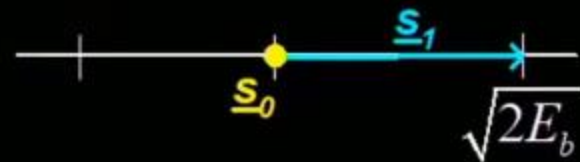
## ➤ Case 1 Antipodal

- ❑ Separation is  $2\sqrt{E_b}$
- ❑ Larger separation, so symbols easier to distinguish



## ➤ Case 2 On-Off Keying

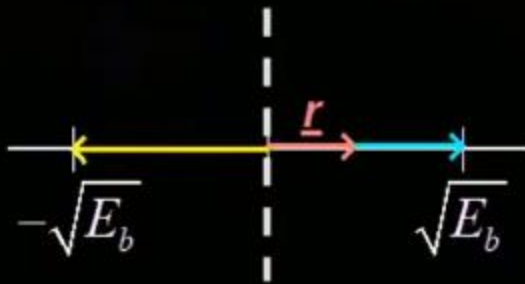
- ❑ Separation is  $\sqrt{2E_b}$
- ❑



# MAP Receiver

## ➤ Case 1 Antipodal

- ❑ Choose the signal closest to received signal



$$r \geq 0 \Rightarrow \text{choose } s_1(t)$$

$$r \leq 0 \Rightarrow \text{choose } s_0(t)$$

## ➤ Case 2 On-Off Keying

- ❑ Choose the signal closest to received signal



$$r \geq \frac{\sqrt{E_b}}{\sqrt{2}} \Rightarrow \text{choose } s_1(t)$$

$$r \leq \frac{\sqrt{E_b}}{\sqrt{2}} \Rightarrow \text{choose } s_0(t)$$

## Calculate $P_e$

$r \geq 0 \Rightarrow$  choisir  $s_1(t)$

$r \leq 0 \Rightarrow$  choisir  $s_0(t)$

### ➤ Case 1 Antipodal

$$P_e = P(0 \text{ sent})P(r \geq 0|0 \text{ sent}) + P(1 \text{ sent})P(r \leq 0|1 \text{ sent})$$

## Calculate $P_e$ equal priors

### ➤ Case 1 Antipodal

$$\begin{aligned} P_e &= P(0 \text{ sent})P(r \geq 0|0 \text{ sent}) + P(1 \text{ sent})P(r \leq 0|1 \text{ sent}) \\ &= \frac{1}{2}P(r \geq 0|r = -\sqrt{E_b} + n) + \frac{1}{2}P(r \leq 0|r = \sqrt{E_b} + n) \end{aligned}$$



# Calculate $P_e$

## ➤ Case 1 Antipodal

$$P_e = P(0 \text{ sent})P(r \geq 0|0 \text{ sent}) + P(1 \text{ sent})P(r \leq 0|1 \text{ sent})$$

$$= \frac{1}{2}P(r \geq 0|r = -\sqrt{E_b} + n) + \frac{1}{2}P(r \leq 0|r = \sqrt{E_b} + n)$$

$$= \frac{1}{2}P(-\sqrt{E_b} + n \geq 0) + \frac{1}{2}P(\sqrt{E_b} + n \leq 0)$$

$$= \frac{1}{2}P(n \geq \sqrt{E_b}) + \frac{1}{2}P(n \leq -\sqrt{E_b})$$

$$= \frac{1}{2}P(n \geq \sqrt{E_b}) + \frac{1}{2}P(n \geq \sqrt{E_b}) \quad \leftarrow \text{Sign changes}$$

$$= P(n \geq \sqrt{E_b}) = \int_{\sqrt{E_b}}^{\infty} \frac{1}{\sqrt{2\pi}\sigma} e^{-n^2/2\sigma^2} dn$$



# Calculate $P_e$

## ➤ Case 1 Antipodal

Quality factor

$$P_e = \int_{\sqrt{E_b}}^{\infty} \frac{1}{\sqrt{2\pi}\sigma} e^{-n^2/2\sigma^2} dn$$

$$Q(x) = \frac{1}{\sqrt{2\pi}} \int_x^{\infty} e^{-x^2/2} dx$$

$$P_e = \int_{\sqrt{E_b}/\sigma}^{\infty} \frac{1}{\sqrt{2\pi}} e^{-u^2/2} du \quad \text{Let, } u = \frac{n}{\sigma}$$

$$= Q\left(\frac{\sqrt{E_b}}{\sigma}\right) = Q\left(\frac{\sqrt{E_b}}{\sqrt{N_0/2}}\right) = Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$$

SNR

$$\frac{E_b}{N_0} = \frac{\text{average energy per bit}}{\text{PSD of Noise}}$$

# Calculate $P_e$

➤ Case 1 Antipodal

➤ Case 2 On–Off Keying

$$r \geq \frac{\sqrt{E_b}}{\sqrt{2}} \Rightarrow \text{choose } s_1(t) \quad r \leq \frac{\sqrt{E_b}}{\sqrt{2}} \Rightarrow \text{choose } s_0(t)$$

$$P_e = \frac{1}{2}P\left(r \geq \sqrt{E_b/2} \mid r = n\right) + \frac{1}{2}P\left(r \leq \sqrt{E_b/2} \mid r = \sqrt{E_b} + n\right)$$

$$= P\left(n \geq \sqrt{E_b/2}\right) = \int_{\sqrt{E_b/2}}^{\infty} \frac{1}{\sqrt{2\pi}\sigma} e^{-n^2/2\sigma^2} du$$

$$= Q\left(\frac{\sqrt{E_b/2}}{\sigma}\right) = Q\left(\frac{\sqrt{E_b/2}}{\sqrt{N_0/2}}\right) = Q\left(\sqrt{\frac{E_b}{N_0}}\right)$$

# OOK vs. BPSK

➤ OOK:  $s_1(t)$  NRZ,  $s_0(t) = 0$

$$P_e = Q\left(\sqrt{\frac{E_b}{N_0}}\right)$$



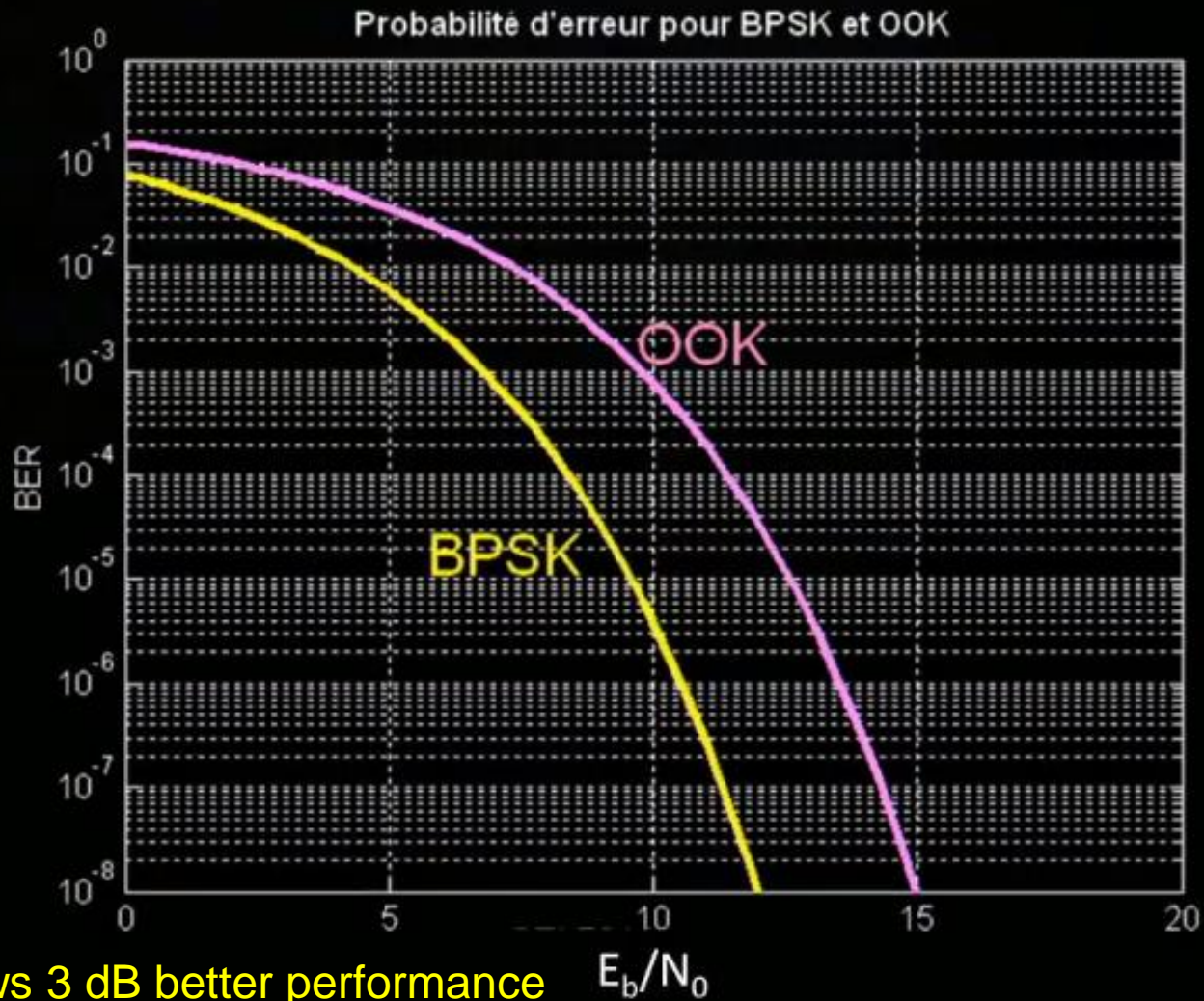
➤ BPSK:  $s_1(t)$  NRZ,  $s_0(t) = -s_1(t)$

$$P_e = Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$$



Large the SNR, lesser the value of Q function, better the signal detection, and it will improve be the performance of the system.

# Probability of error



BPSK shows 3 dB better performance



# Challenges for M-ary cases

- Choosing the symbol
  - MAP most probable (*a posteriori*), and ML most likely
  - modulation determines the form of the receiver
  - MAP/ML determines the receiver decision algorithm

## Analyzing the performance

- exploit the results for the binary case
- exploit the geometric interpretation of the receiver



# Receivers

- Maximizes *a posteriori* probability  
(with *a priori* information)
  - $i$  that minimizes  $\|\mathbf{r} - \mathbf{s}_i\|^2 - N_0 \ln P(\mathbf{s}_i)$
- Maximizes the likelihood ratio
  - $i$  that minimizes  $\|\mathbf{r} - \mathbf{s}_i\|^2$

# Receivers

- MAP

➤  $i$  that minimizes  $\|\mathbf{r} - \mathbf{s}_i\|^2 - N_0 \ln P(\mathbf{s}_i)$

$$\|\mathbf{r}\|^2 - 2\langle \mathbf{r}, \mathbf{s}_i \rangle + \|\mathbf{s}_i\|^2 - N_0 \ln P(\mathbf{s}_i)$$

- ML

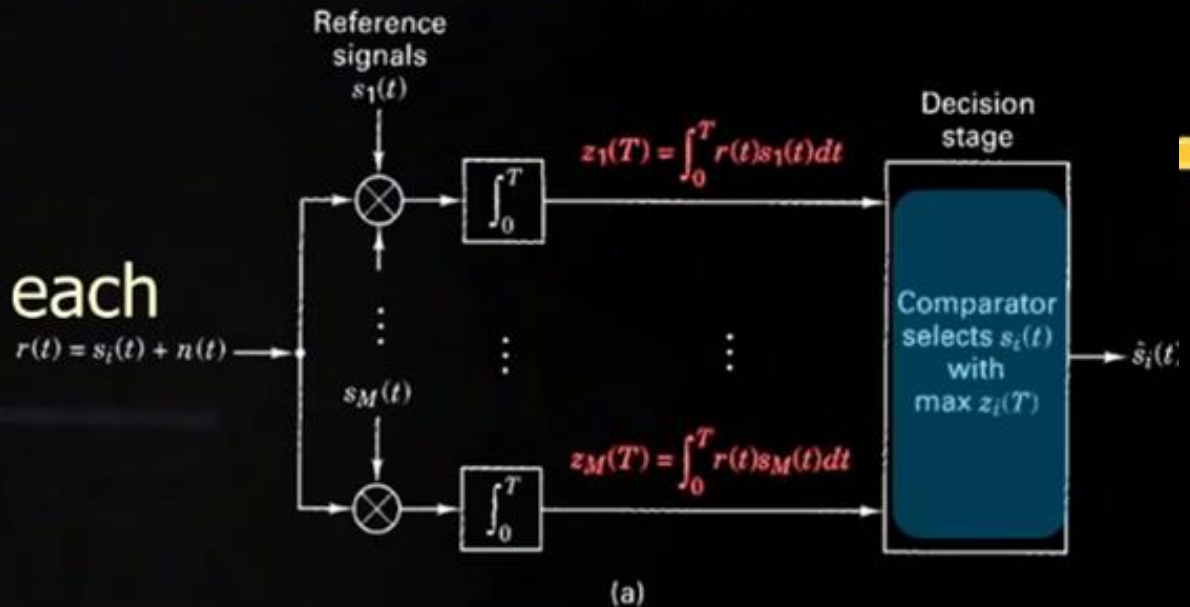
➤  $i$  that minimizes  $\|\mathbf{r} - \mathbf{s}_i\|^2$

$$\|\mathbf{r}\|^2 - 2\langle \mathbf{r}, \mathbf{s}_i \rangle + \|\mathbf{s}_i\|^2$$

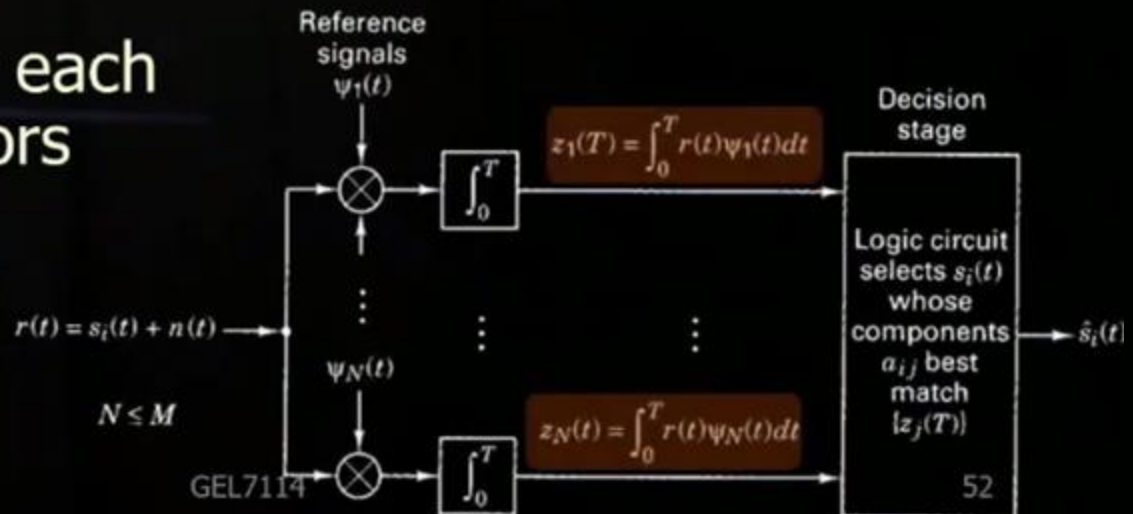
$$\|\mathbf{r}\|^2 - 2\langle \mathbf{r}, \mathbf{s}_i \rangle + \|\mathbf{s}_i\|^2 - N_0 \ln P(\mathbf{s}_i)$$

# Receivers

- One branch for each of  $M$  symbols



- One branch for each of  $N$  basis vectors  
 $N \leq M$



# Receiver vs. modulation format

- MPSK
  - two branches: in-phase et quadrature
- MFSK
  - M branches: one for each symbol
  - basis vector = symbol
  - **orthogonal modulation**
- QAM
  - two branches: in-phase and quadrature

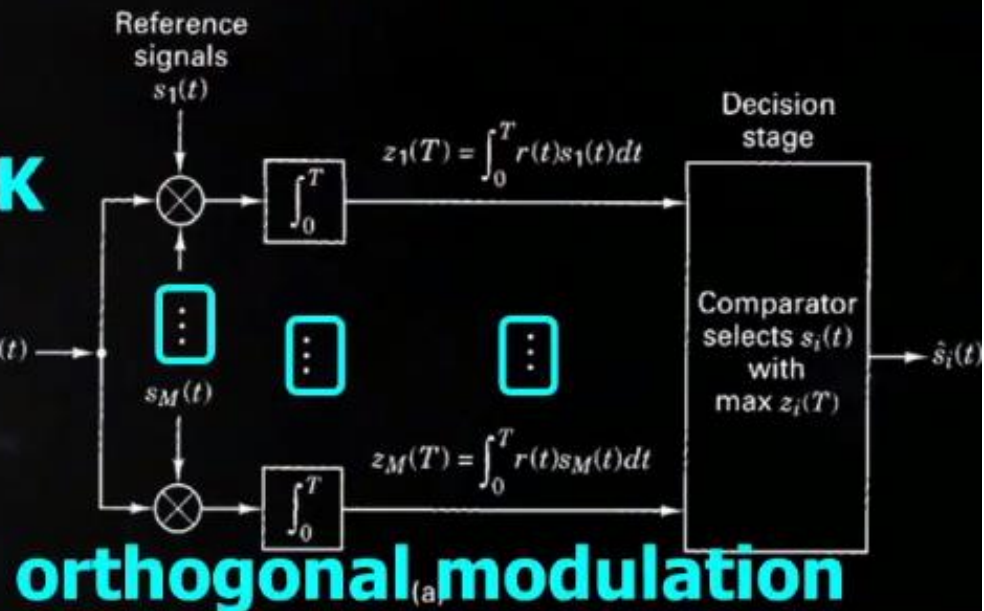


# Receivers

## MFSK

- One branch for each  $M$  symbols

$$r(t) = s_i(t) + n(t)$$

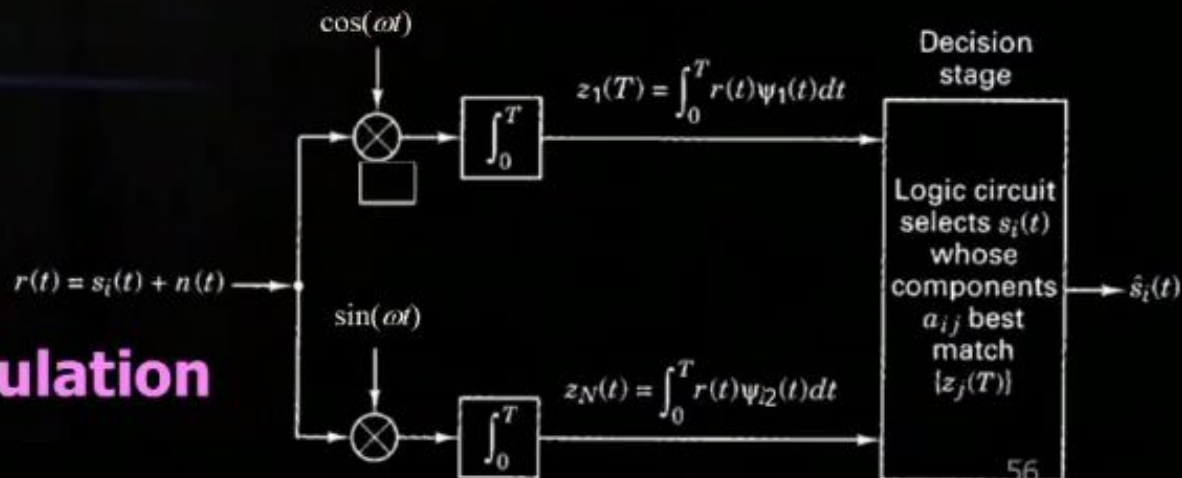


- Two branches

## QAM

## MPSK

## Non orthogonal modulation





# Decision algorithms

- Geometric interpretation

- To reduce the complexity of the decision algorithm
- To provide an intuition of how the receiver works

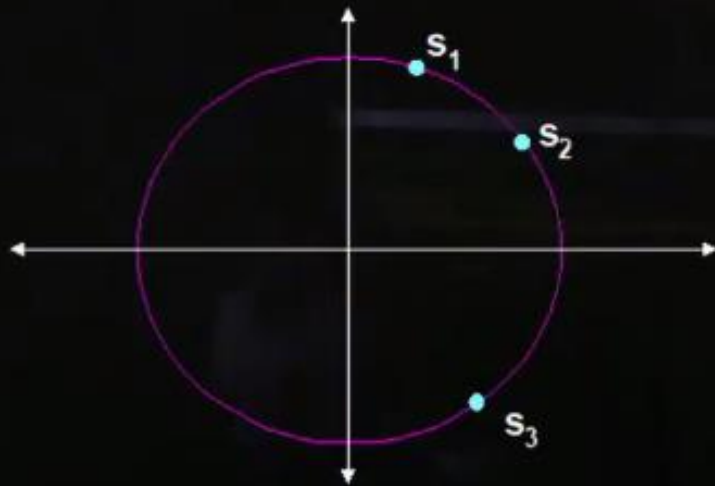
- $i$  that minimizes

$$\|\mathbf{r} - \mathbf{s}_i\|^2$$

- Equivalent to choosing the nearest symbol
- Regions can be determined in advance to simplify the receiver

How to choose the nearest symbol directly from  
vector received?

## Decision regions

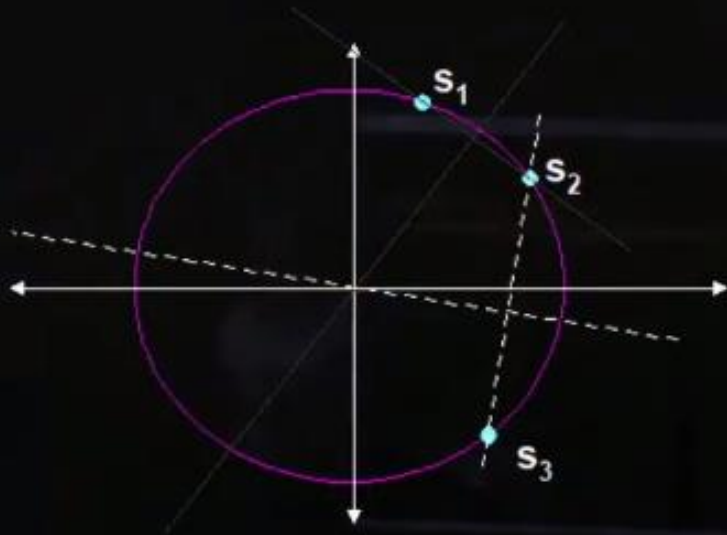


Constellation (equal energies)

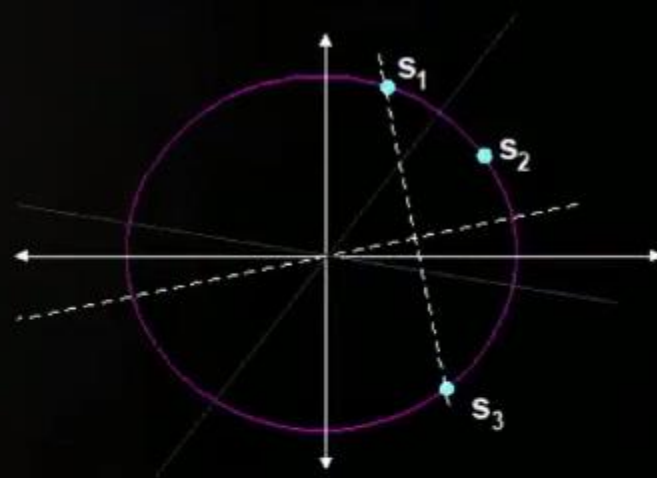


Perpendicular bisector  
between  $s_1$  and  $s_2$

# Decision regions

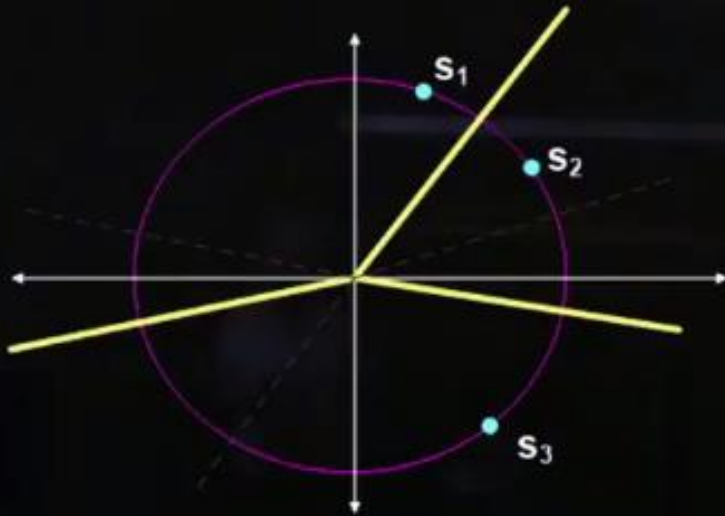


perpendicular bisector  
between  $s_3$  and  $s_2$

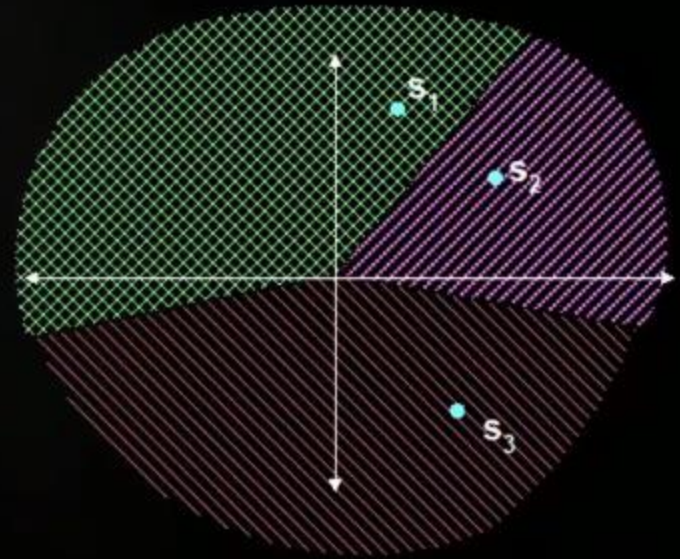


perpendicular bisector  
between  $s_1$  and  $s_3$

# Decision regions



bissecteurs perpendiculaires



régions

How to choose the nearest symbol directly from vector received?  $\Rightarrow$  **regions**



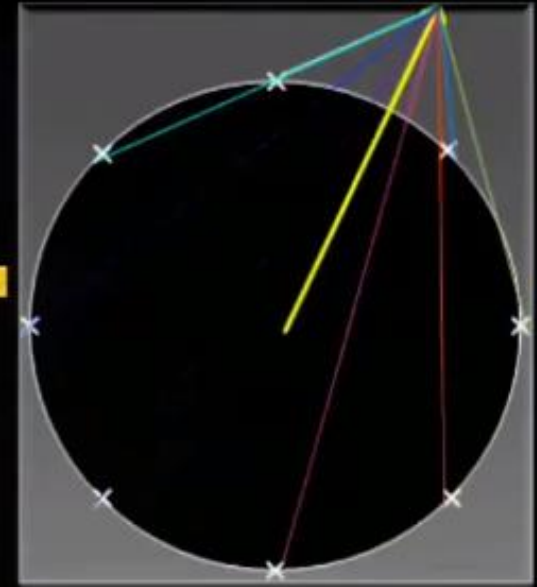
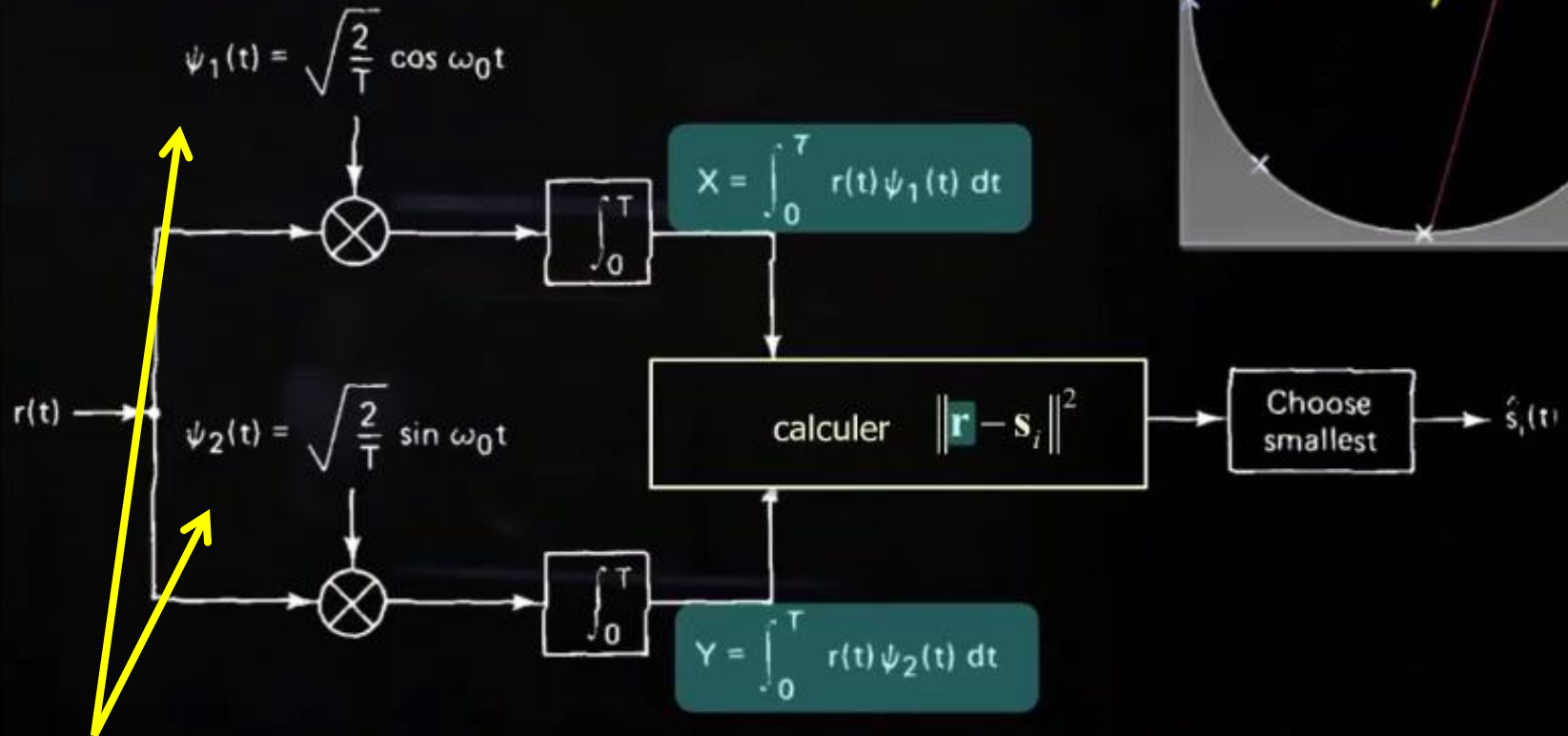
# M-PSK



Example: M-ary PSK

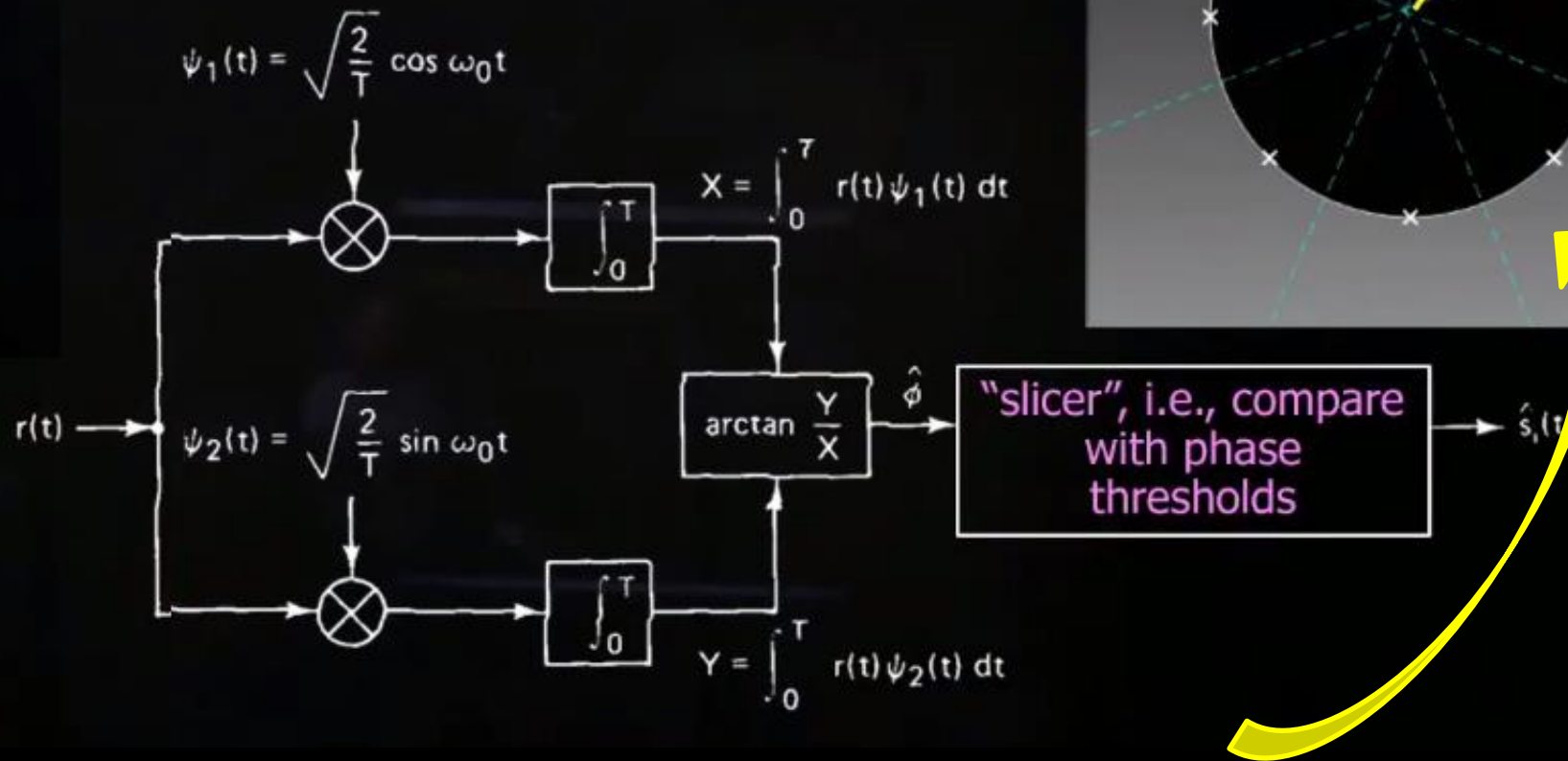


# MPSK receiver



**Complexity reduction: For  $M$  symbols, only two Correlator are sufficient due to basis vector.**

# MPSK receiver



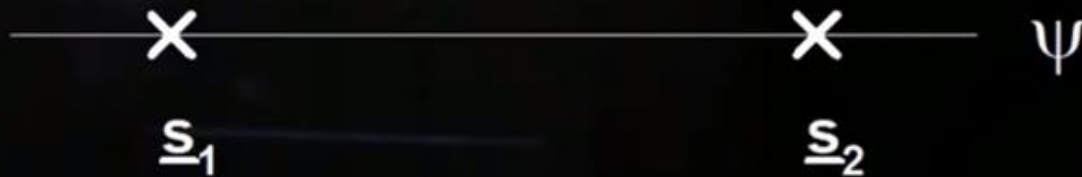
**Complexity reduction:** Now, Instead of calculating  $M$  distance vector corresponding to  $M$  symbols, we need to use only single phase of the received vector for detection of region it falls in.

# Challenges for M-ary cases

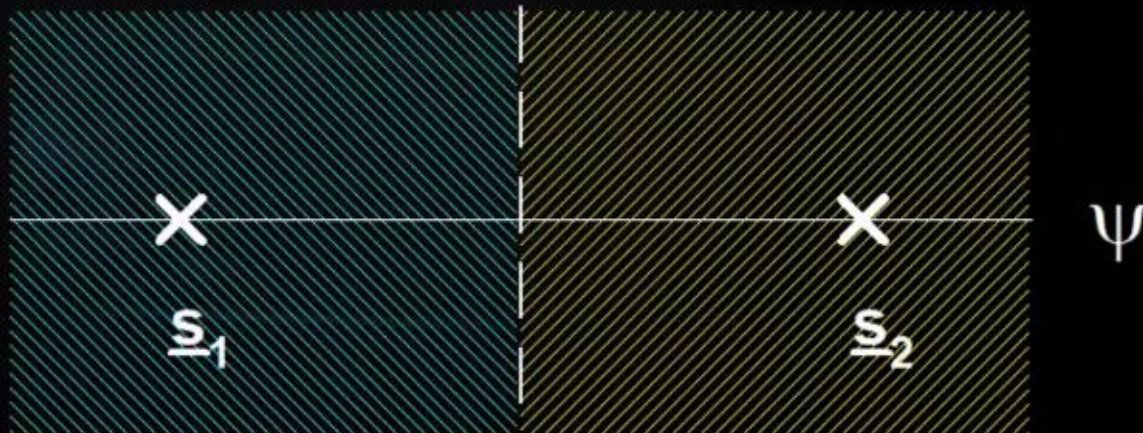
- Choosing the symbol
  - MAP most probable (*a posteriori*), and ML most likely
  - determines the form of the receiver
- Analyzing the performance
  - exploit the results for the binary case
  - exploit the geometric interpretation of the receiver

# $P_e$ from the signal space

- What can we learn from the binary case?
  - Signal space: two points



## Decision regions



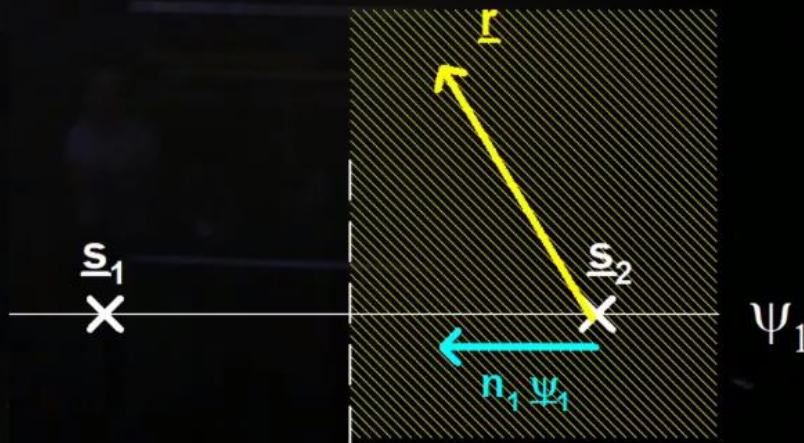


# Received vector

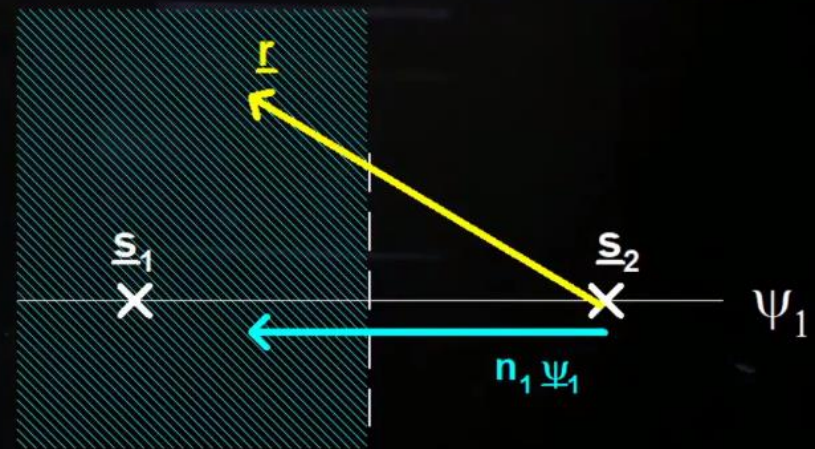


# Noise contribution

- Region two  $\Rightarrow$  choose  $s_2$ , no error



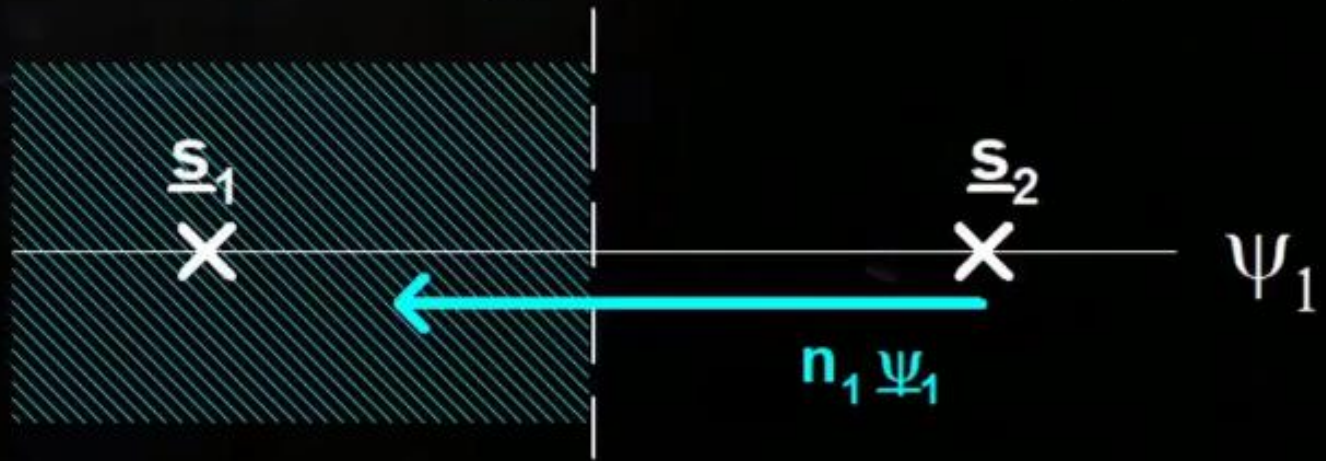
- Region one  $\Rightarrow$  choose  $s_1$ , **error**



# Errors

- When does an error occur?
- When the noise vector exceeds the bisector ...

$$n_1 > \frac{1}{2} \|\mathbf{s}_1 - \mathbf{s}_2\| = \frac{D}{2}$$



**How to calculate  $P_e$  for the M-ary case now that we have the decision regions?**



## **Minimal distance**



- Definition  $D_{\min} = \min_{i \neq k} \|s_i - s_k\|$

➤ closer signals => more errors



# Union bound

A thick, horizontal yellow brushstroke underline that spans the width of the slide, starting below the title and extending to the right edge.

# QPSK

- Quadrature phase shift keying

$$s(t) = \sqrt{\frac{2E}{T}} \cos(\omega_0 t + \phi_i(t)) \quad 0 \leq t \leq T \quad \phi_i \in \left\{ 0, \frac{\pi}{2}, \pi, \frac{3\pi}{2} \right\}$$

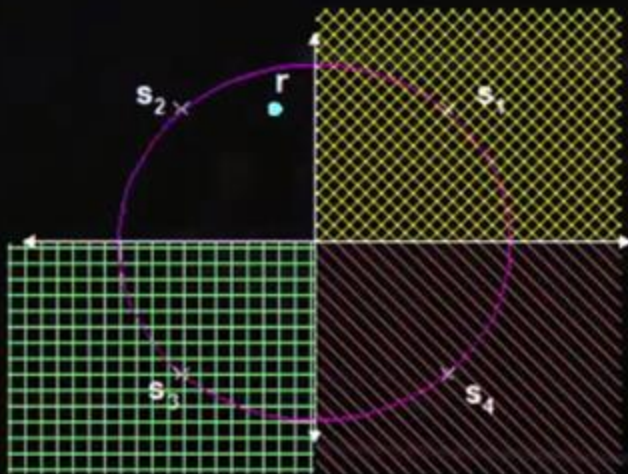
- $E$  energy/symbol,  $T$  symbol time

- Basis vectors

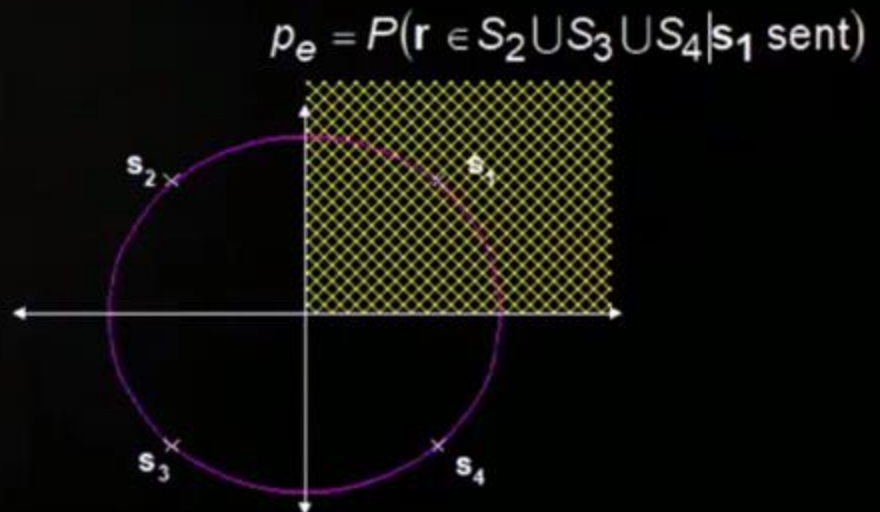
$$\psi_1 = \cos(\omega_0 t) \quad \psi_2 = \sin(\omega_0 t)$$

- Decision regions – four quadrants

# Union bound

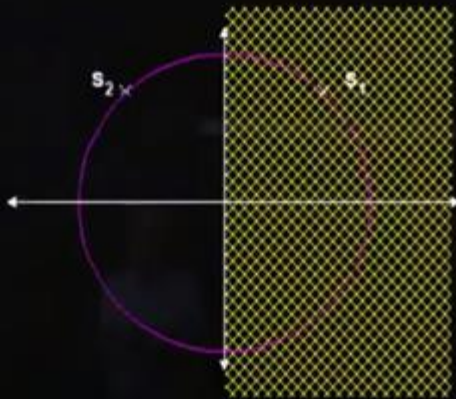


4 decision regions



error region when  
 $s_1$  was sent

# Union bound



$$p_2(s_2 \text{ chosen} | s_1 \text{ sent}) \\ = P(r \notin S_1 | s_1 \text{ sent})$$

only  $s_1$  and  $s_2$  exist



$$p_2(s_4 \text{ chosen} | s_1 \text{ sent}) \\ = P(r \notin S_1 | s_1 \text{ sent})$$

only  $s_1$  and  $s_4$  exist

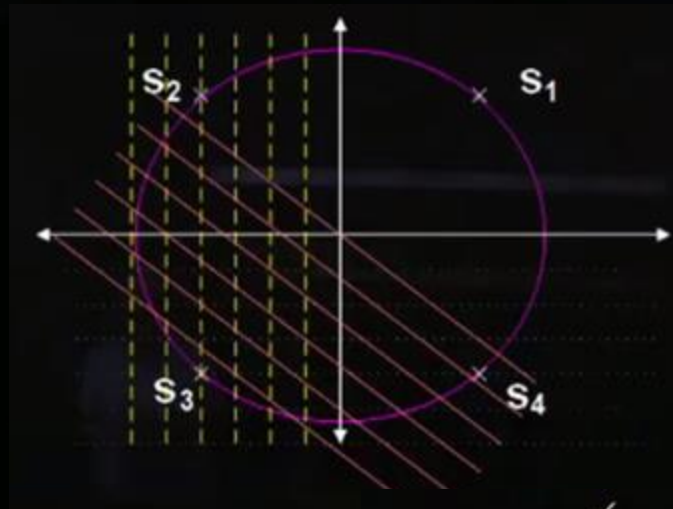


$$P(r \notin S_1 | s_1 \text{ sent}) \\ = p_2(s_3 \text{ choisi} | s_1 \text{ envoyé})$$

only  $s_1$  and  $s_3$  exist



# Union bound

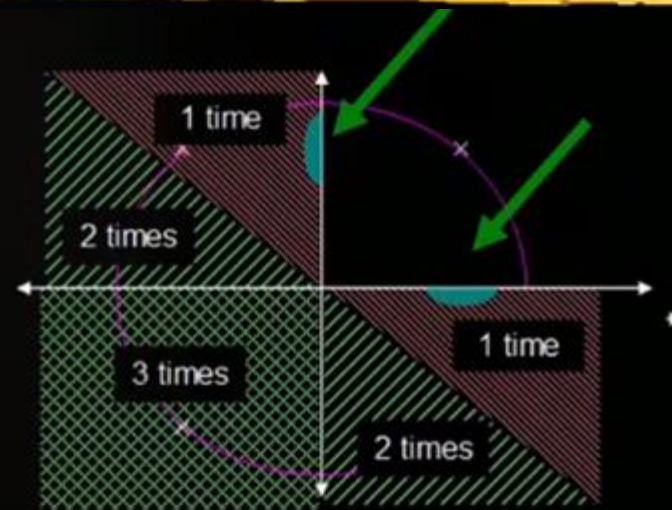


**UNION BOUND**



$$P_e(QPSK) \leq 2Q\left(\sqrt{\frac{2E_b}{N_0}}\right) + Q\left(2\sqrt{\frac{E_b}{N_0}}\right)$$

$$\begin{aligned} p_e &\leq p_2(s_2 \text{ chosen} | s_1 \text{ sent}) \\ &\quad + p_2(s_3 \text{ chosen} | s_1 \text{ sent}) \\ &\quad + p_2(s_4 \text{ chosen} | s_1 \text{ sent}) \end{aligned}$$



**Approximated  
UNION BOUND**



$$P_e(QPSK) \approx 2Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$$

# Union bound

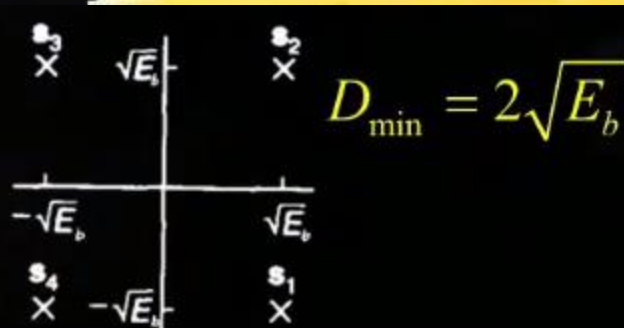
For M-Ary modulation schemes

- Error probability dominated by the nearest points

$$P_e \approx \frac{2K}{M} Q\left(\frac{D_{\min}}{\sqrt{2N_0}}\right) \quad \leftarrow \text{Approximated UNION BOUND}$$

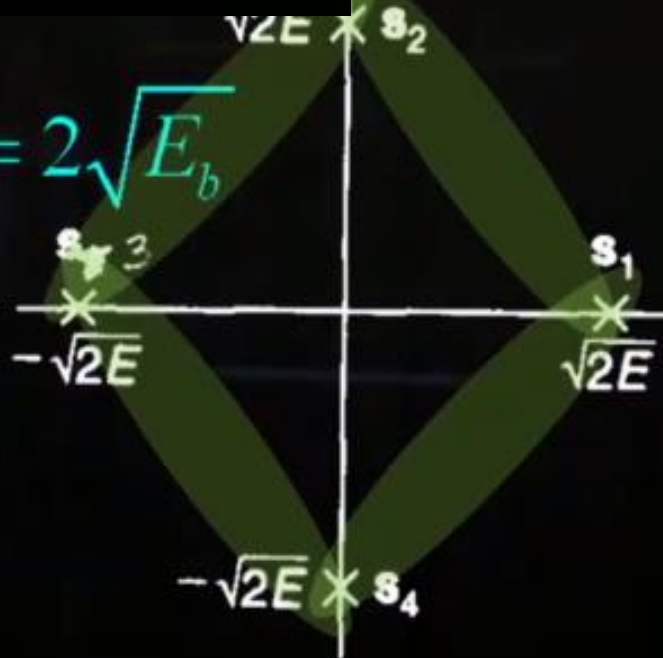
- $K$  is the number of pairs of symbols separated by the minimum distance  $D_{\min}$

# QPSK



$$P_e \approx \frac{2K}{M} Q\left(\frac{D_{\min}}{\sqrt{2N_0}}\right) = \frac{2K}{M} Q\left(\frac{2\sqrt{E_b}}{\sqrt{2N_0}}\right)$$

$$D_{\min} = 2\sqrt{E_b}$$



- $M=4$

- $K=4$

$$P_e \approx 2Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$$

# QPSK

- Error probability coming from the union bound

$$P_e(QPSK) \leq 2Q\left(\sqrt{\frac{2E_b}{N_0}}\right) + Q\left(2\sqrt{\frac{E_b}{N_0}}\right)$$

- Approximation

$$P_e(QPSK) \approx 2Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$$

- Exact error probability

$$P_e(QPSK) = 2Q\left(\sqrt{\frac{2E_b}{N_0}}\right) - Q^2\left(\sqrt{\frac{2E_b}{N_0}}\right)$$



# QPSK

0 dB

0.15963807454081

$$P_e(QPSK) \leq 2Q\left(\sqrt{\frac{2E_b}{N_0}}\right) + Q\left(2\sqrt{\frac{E_b}{N_0}}\right)$$

0.15729920705029

$$P_e(QPSK) \approx 2Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$$

0.15111344691562

$$P_e(QPSK) = 2Q\left(\sqrt{\frac{2E_b}{N_0}}\right) - Q^2\left(\sqrt{\frac{2E_b}{N_0}}\right)$$

# QPSK

	5dB	10dB
$P_e(QPSK) \leq 2Q\left(\sqrt{\frac{2E_b}{N_0}}\right) + Q\left(2\sqrt{\frac{E_b}{N_0}}\right)$	<b>0.01190797987629</b>	<b>0.00000774421643</b>
$P_e(QPSK) \approx 2Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$	<b>0.01190773429556</b>	<b>0.00000774421643</b>
$P_e(QPSK) = 2Q\left(\sqrt{\frac{2E_b}{N_0}}\right) - Q^2\left(\sqrt{\frac{2E_b}{N_0}}\right)$	<b>0.01187228576154</b>	<b>0.00000774420144</b>

Thanks !