C++ Implementation of Cryptographic Algorithms

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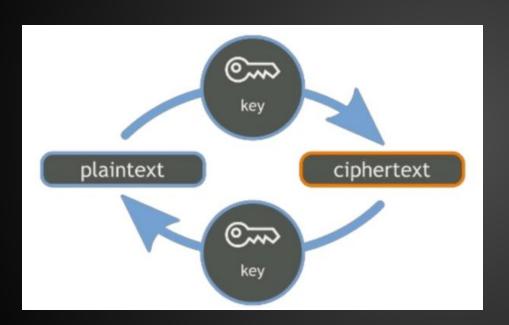
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Overview

- What is cryptography?
- Why is it important?

"Cryptography is the study of secure communications techniques that allow only the sender and intended recipient of a message to view its contents."

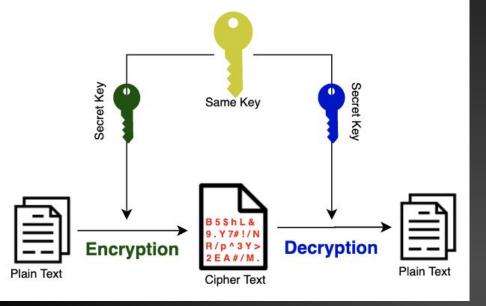


Importance:

- Confidentiality
- Identification and Authentication
- Integrity
- Non-repudiation

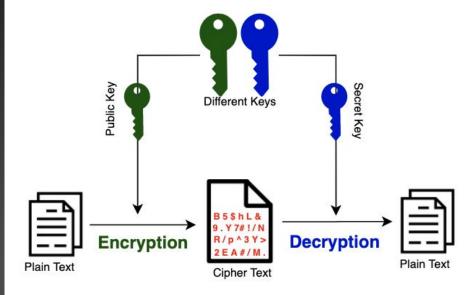
Cryptographic Algorithms

• Asymmetric vs Symmetric



Symmetric Encryption (Private Key)

- Advanced Encryption Standard (AES)
- Data Encryption Standard (DES)

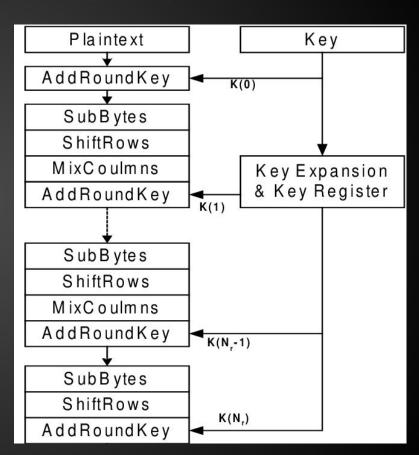


Asymmetric Encryption (Public Key)

- Rivest–Shamir–Adleman (RSA)
- Diffie-Hellman (DH)
- Elliptic Curve Cryptosystem (ECC)
- Digital Signature Algorithm (DSA)

Advanced Encryption Standard (AES)

- Developed by Vincent Rijmen and Joann Daeman
- Recommended by NIST to replace DES in 2001
- Algorithm is referred to as AES-128, AES-192, or AES-256, depending on the key length
- Relatively fast encryption and decryption
- Low power consumption
- Excellently secured



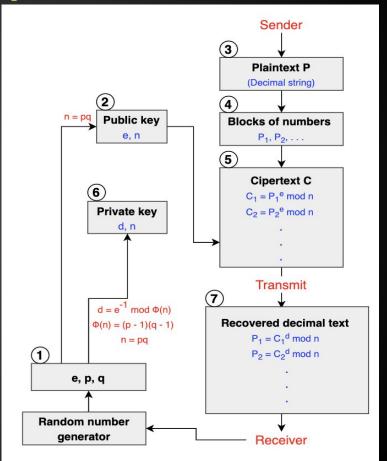


Implementation

- RSA & Diffie-Hellman Key Exchange
- The math behind them

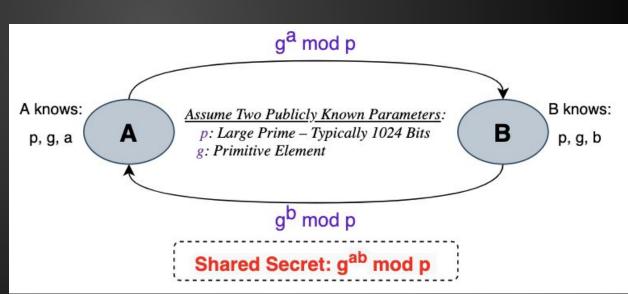
Rivest-Shamir-Adleman (RSA)

- Founded in 1977 by Ronald L. Rivest, Adi Shamir, and Leonard M. Adleman
- High power consumption
- Slow encryption and decryption
- Inherent vulnerability to brute force and oracle attacks
- Algorithm is based on the theory of Prime Numbers



Diffie-Hellman Key Exchange (DH)

- Founded between 1969-1973; published in 1976 by Whitfield Diffie and Martin Hellman
- Some common DH protocols attacks are
 - denial-of-service attacks
 - man-in-the-middle attacks
 - outsider attacks
 - insider attacks



Designs

- Language: C++
- Using BigInteger Library GMP

Design: Diffie-Hellman using GMP

```
// Below:
// Alice choose a = kpr, A \in \{2, \ldots, p-2\}
// Alice compute A = kpub,A ≡αa mod p
// a is Alice's secret key
// Alice computes A , and sent it to Bob
    mpz_t a;
    mpz_t A;
    mpz_inits(a, A, NULL);
    gmp_randstate_t state_a;
    gmp_randinit_mt (state_a);
    gmp_randseed_ui(state_a, time(NULL) );
    mpz_urandomm(a, state_a, n2);
                                          //choose a = kpr, A \in \{2, ..., p-2\}
    mpz_powm(A, alpha, a, p);
                                       //compute A = kpub, A \equiv \alpha^a \mod p
    cout << "Alice secret key a: " << endl;</pre>
    gmp_printf(" %Zd \n", a);
    cout << endl;</pre>
    cout << "Alice public key A: " << endl;</pre>
    gmp_printf(" %Zd \n", A);
    cout << endl;</pre>
```

```
// Bob choose b = kpr, B \in \{2, \ldots, p-2\}
// Bob compute B = kpub, B \equiv \alpha b \mod p
// b is Bob's secret key
// Bob computes B , and sent it to Alice
    mpz t b;
    mpz_t B;
    mpz_inits(b, B, NULL);
    gmp_randstate_t state_b;
    gmp_randinit_mt (state_b);
    gmp_randseed_ui(state_b, (time(NULL) + 50) );
    mpz_urandomm(b, state_b, n2);
                                           //choose b = kpr, B \in \{2, ..., p-2\}
    mpz_powm(B, alpha, b, p);
                                       //compute B = kpub_B \equiv \alpha b \mod p
    cout << "Bob secret key b: " << endl;</pre>
    gmp_printf(" %Zd \n", b);
    cout << endl;</pre>
    cout << "Bob public key B: " << endl;</pre>
    qmp printf(" %Zd \n", B);
    cout << endl:
```

Design: RSA using GMP

* These are only code snippets. Full code on GitHub*
We also wrote RSA using Boost Library but GMP is more efficient.

```
void generate_keys(mpz_t p, mpz_t q) {
    if (mpz_probab_prime_p(p, 50) == 2 \&\& mpz_probab_prime_p(q, 50) == 2){}
        if (p == a){
           cout << " Cannot have identical prime numbers"<<endl;</pre>
           return;
       mpz_t phi, check;
       mpz_inits(phi, e, check, d, NULL);
                                                       void encryptMsg(mpz t pbk, string msg, mpz t n){
       mpz_sub_ui(p, p, 1);
                                                             For encryption, c = m^e \mod n, where m = original message.
       mpz_sub_ui(q, q, 1);
                                                           mpz_t ciphertxt, x;
       mpz_mul(phi, p, q);
                                                           unsigned long int tst;
                                                           mpz_init(ciphertxt);
       gmp_randstate_t rd;
       gmp_randinit_mt(rd);
                                                           for (int i = 0; i < msg.length(); i++)
       unsigned long int seed = 20201204;
       qmp randseed ui (rd, seed);
                                                               mpz_init_set_ui(x, int(msg.at(i)));
                                                               mpz_powm(ciphertxt, x, e, n);
       mpz urandomm(e, rd, phi);
                                                               tst = mpz get ui(ciphertxt);
       while (mpz cmp(e. phi) < 0){
                                                               arr.push back(int(tst));
           mpz_gcd(check, e, phi);
                                                               cout << tst;
           if (mpz cmp ui(check, 1) == 0){
                break:
           else {
                                                           cout << endl:
                seed = seed + 3:
                                                           mpz_clears(x, ciphertxt, NULL);
                gmp_randseed_ui (rd, seed);
                                                           cout << endl:
                mpz_urandomm(e, rd, phi);
       mpz_invert(d, e, phi);
       mpz_clears(check, phi, NULL);
```

```
void decryptMsg(mpz_t prvk, mpz_t n){
     For decryption, m = c^d \mod n.
   mpz_t OG, c, temp;
   mpz_t N1, N2;
   mpz inits(N1, N2, NULL);
   mpz init set(N1, prvk);
   mpz_init_set(N2, n);
   mpz init(c);
    long int temp2 = 0;
    for (int i = 0; i < arr.size(); i++){
       temp2 = arr[i];
       mpz_init_set_ui(temp, temp2);
       mpz set ui(c, temp2);
       mpz_powm(OG, temp, N1, N2);
   unsigned long int tst2;
   tst2 = mpz get ui(0G);
   cout << static cast<char>(tst2);
   cout <<endl;</pre>
```

Output: RSA using GMP

```
qhaith@ubuntu:~/Documents/CPE 593/Final Project$ g++ main.cpp -lgmpxx -lgmp
ghaith@ubuntu:~/Documents/CPE 593/Final Project$ ./a.out
211 229 48319
*** Generating public/private key ***
Public Key: 19307
Private Key: 7043
Original Message:
THIS IS A TEST
... Encrypting message ...
Encrypted Message: 4653941936271043814276927104381427691158727694653986964381446539
... Decrypting message ...
Decrypted Message:
THIS IS A TEST
Private Key: 7043
ghaith@ubuntu:~/Documents/CPE_593/Final_Project$ ./a.out
211 59 12449
*** Generating public/private key ***
Public Key: 2923
Private Key: 4867
Original Message:
THIS IS A TEST
... Encrypting message ...
Encrypted Message: 511060491732244262617322442626109262651101210822445110
... Decrypting message ...
Decrypted Message:
THIS IS A TEST
Private Key: 4867
ghaith@ubuntu:~/Documents/CPF 593/Final Project$ /a.out
```

Discussion

- Weaknesses in the algorithms
- Weaknesses in the GMP Library

Mathematical Attacks against RSA

- By determining the prime factors p and q of the modulus n, an attacker can find out $\phi(n) = (p-1)(q-1)$, which is in turn enables to determine $d = e \land -1 \pmod{\phi(n)}$.
- By figuring out the totient $\phi(n)$ directly without first determining p and q from which d = e \land -1 (mod $\phi(n)$) can be determined.
- By calculating d directly without first determining $\phi(n)$.

Padding

- RSA encryption is deterministic, i.e., for a specific key, a particular plaintext is always mapped to a particular ciphertext. An attacker can derive statistical properties of the plaintext from the ciphertext
- A possible solution to all these problems is the use of padding, which embeds a random structure into the plaintext before encryption and avoids the above mentioned problems. There is many techniques for padding.

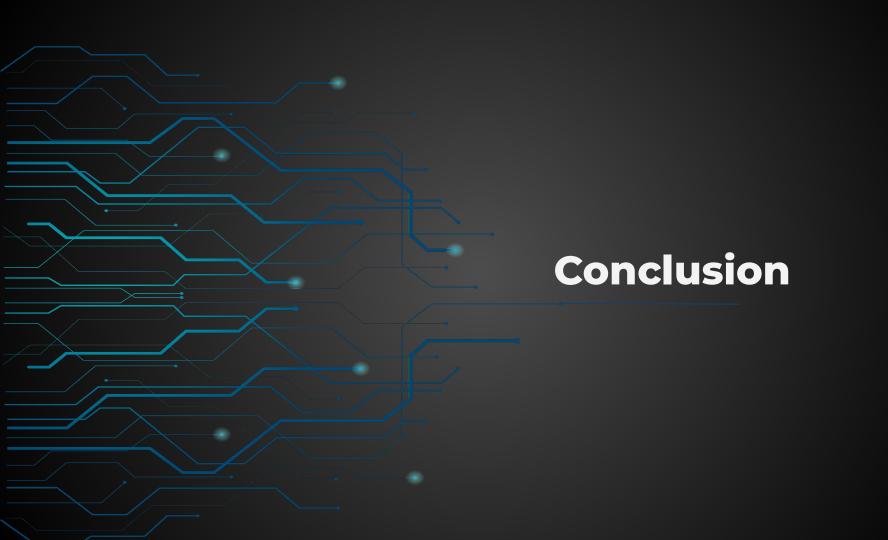
Side-Channel Attacks (Timing Attack)

- They exploit information about the private key which is leaked through physical channels such as the power consumption or the timing behavior.
- There are several countermeasures available to prevent the attack. A simple one is
 to execute a multiplication with dummy variables after a squaring that
 corresponds to an exponent bit 0. This results in a power profile (and a run time)
 which is independent of the exponent.
- Countermeasures against more advanced side-channel attacks are not as straightforward.
- GMP Library provides some low-level functions for cryptography where resilience to side-channel attacks is desired.

- In general, the security of a cryptosystem is a function of two things: the strength of the algorithm and the length of the key.
- A perfect cryptography algorithm is the one that no way to break it than with a brute-force attack.
- RSA with short secret key is proven **insecure** against brute force attack.
- This attack can be easily circumvented by choosing large key.
- However, the larger key will make the encryption and decryption process little **slow** as it will require **greater computations** in key generation as well as in encryption/decryption algorithm.
- As of 2020, the largest publicly known factored RSA number was 829 bits (250 decimal digits, RSA-250).[29] Its factorization, by a state-of-the-art distributed implementation, took approximately 2700 CPU years.
- Quantum Computing is the current biggest threat to RSA and many other cryptography algorithms.

Recommendations:

- In practice, RSA keys are typically 1024 to 4096 bits long.
- As of 2020, it is not known whether such 1024-bit keys can be cracked, but minimum recommendations have moved to at least 2048 bits.



Future Works

- RSA is not robust as it should be due to the deterministic property, it would be interesting to read and work on different methods to make it more secure.
- Implement Encryption methods on different file types.
- Hybrid encryption method utilizing multiple algorithms to improve the security.

THANKS!

* For more information, review the notes section of each slide or our paper*