Special Topics in Security ECE 5698

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Unix Security



- Multi-user operating system
- Operating system functionality
 - process management
 - (virtual) memory management
 - file system management
 - I/O management



- Structure
 - operating system kernel
 - user-space programs (daemons, applications, shell)

Kernel

- provides a hardware abstraction layer for user-space programs
- complete access to all (physical) resources
- trusted computing base
- provides services via system calls



System call

- performs a transition from user mode to privileged (kernel) mode
- this crosses the border between two security domains
- usually implemented with hardware (processor) support
 - processor interrupt
 - x86 call gates

- Kernel vulnerability
 - usually leads to complete system compromise
 - attacks performed via system calls
 - e.g., a famous one appeared in February 2009, vmsplice
- Solaris / NetBSD call gate creation input validation problem
 - malicious input when creating a LDT (x86 local descriptor table)
 - used in 2001 by Last Stage of Delirium to win Argus Pitbull
 Competition
- Kernel Integer Overflows
 - FreeBSD procfs code (September 2003)
 - Linux brk() used to compromise debian.org (December 2003)
 - Linux setsockopt() (May 2004)

More Linux vulnerabilities

- Linux message interface (August 2005, CAN-2005-2490)
- race condition proc and prct1 (July 2006, CVE-2006-3626)
- local privilege escalation (September 2007, CVE 2007-4573)

Device driver code is particularly vulnerable

- (most) drivers run in kernel mode, either kernel modules or compiled-in
- often not well audited
- very large code based compared to core services

Examples

- aironet, asus_acpi, decnet, mpu401, msnd, and pss (2004)
 found by sparse (tool developed by Linus Torvalds)
- remote root (MadWifi 2006, Broadcom 2006)

- Code running in user mode is always linked to a certain identity
 - security checks and access control decisions are based on user identity
- Unix is user-centric
 - no roles
- User
 - identified by user name (UID), group name (GID)
 - authenticated by password (stored encrypted)
- User root
 - superuser, system administrator
 - special privileges (access resources, modify OS)
 - cannot decrypt user passwords

Process Management

Process

- implements user-activity
- entity that executes a given piece of code
- has its own execution stack, memory pages, and file descriptors table
- separated from other processes using the virtual memory abstraction

Thread

- separate stack and program counter
- share memory pages and file descriptor table

Process Management

- Process Attributes
 - process ID (PID)
 - uniquely identified process
 - user ID (UID)
 - ID of owner of process
 - effective user ID (EUID)
 - ID used for permission checks (e.g., to access resources)
 - saved user ID (SUID)
 - to temporarily drop and restore privileges
 - lots of management information
 - scheduling
 - memory management, resource management



User Authentication

- How does a process get a user ID?
- Authentication (login)
- Passwords
 - Traditional: user passwords are used as keys for crypt() function
 - runs DES algorithm 25 times on a block of zeros
 - 12-bit "salt"
 - 4096 variations
 - chosen from date, not secret
 - prevent same passwords to map onto same string
 - make dictionary attacks more difficult



- Password cracking
 - dictionary attacks
 - Crack, JohnTheRipper

User Authentication

- Shadow passwords
 - password file is needed by many applications to map user ID to user names
 - encrypted passwords are not
- /etc/shadow
 - holds encrypted passwords
 - account information
 - last change date
 - expiration (warning, disabled)
 - minimum change frequency
 - readable only by superuser and privileged programs
 - MD5 hashed passwords (default) to slow down guessing



DEMO, Shadow Passwds...

Group Model

- Users belong to one or more groups
 - primary group (stored in /etc/password)
 - additional groups (stored in /etc/group)
 - possibility to set group password
 - and become group member with newgrp



/etc/group

```
groupname : password : group id : additional users
root:x:0:root
bin:x:1:root,bin,daemon
users:x:100:ek
```

DEMO, groups...

File System

File tree

- primary repository of information
- hierarchical set of directories
- directories contain file system objects (FSO)
- root is denoted "/"



File system object

- files, directories, symbolic links, sockets, device files
- referenced by *inode* (index node)

File System

- Access Control
 - permission bits
 - chmod, chown, chgrp, umask
 - file listing:



Type	r	W	X	S	t
File	read access	write access	execute	suid / sgid inherit id	sticky bit
Directory	list files	insert and remove files	stat / execute files, chdir	new files have dir-gid	files only delete- able by owner

Shell

Shell

- one of the core Unix application
- both a command language and programming language
- provides an interface to the Unix operating system
- rich features such as control-flow primitives, parameter passing, variables, and string substitution
- communication between shell and spawned programs via redirection and pipes
- different flavors
 - bash and sh, tcsh and csh, ksh

Shell Attacks

Environment Variables

- + SHOME and SPATH can modify behavior of programs that operate with relative path names
- \$IFS internal field separator
 - used to parse tokens
 - usually set to [\t\n] but can be changed to "/"
 - "/bin/ls" is parsed as "bin ls" calling bin locally



- preserve attack (/usr/lib/preserve is SUID)
 - called "/bin/mail" when vi crashed to preserve file
 - change IFS, create bin as link to /bin/sh, kill vi

Shell Attacks

- Control and escape characters
 - can be injected into command string
 - modify or extend shell behavior
 - user input used for shell commands has to be rigorously sanitized
 - easy to make mistakes
 - classic examples are `;' and `&'
- Applications that are invoked via shell can be targets as well
 - increased vulnerability surface
- Restricted shell
 - invoked with -r
 - more controlled environment

DEMO, Restricted Shell, creating a chroot environment

Shell Attacks

- system(char *cmd)
 - function called by programs to execute other commands
 - invokes shell
 - executes string argument by calling /bin/sh -c string
 - makes binary program vulnerable to shell attacks
 - especially when user input is utilized
- popen (char *cmd, char *type)
 - forks a process, opens a pipe and invokes shell for cmd

DEMO, system() issue...

File Descriptor Attacks

- SUID program opens file
- forks external process
 - sometimes under user control
- on-execute flag
 - if close-on-exec flag is not set, then new process inherits file descriptor
 - malicious attacker might exploit such weakness
- Linux Perl 5.6.0
 - getpwuid() leaves /etc/shadow opened (June 2002)
 - problem for Apache with mod_perl

