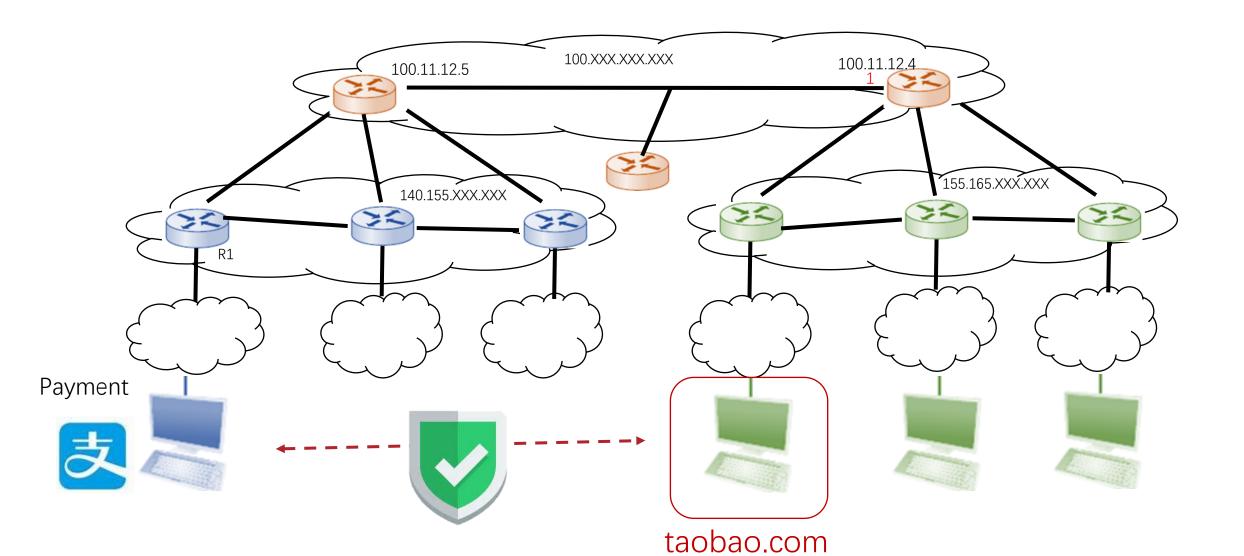


CS120: Computer Networks

Lecture 27. Network Security 1

Zhice Yang



What is Network Security

- Confidentiality
 - To encrypt messages so as to prevent an adversary from understanding the message contents
- Integrity
 - To prevent an adversary from modifying the message contents.
- Availability
 - services must be accessible and available to users
- Authentication
 - To confirm identity of each other
- Timeliness
 - To identify delayed messages

Guarantee	Primitive
Confidentiality	Encryption
Integrity	Hash
Authentication	Digital Signature

Security Risks in Networks

- Eavesdrop
- Injection
- Impersonation
 - can fake (spoof) source address in packet (or any field in packet)
- Hijacking
 - "take over" ongoing connection by removing sender or receiver, inserting himself in place
- Denial of Service (DoS):
 - prevent service from being used by others (e.g., by overloading resources)

• ...

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Cipher

• Cipher: the Cryptographic Algorithm for Encryption or Decryption

HELLO

ABCDEFGHIJKLMNOPQRSTUVWXYZ



RSTUVWXYZABCDEFGHIJKLMNOPQ

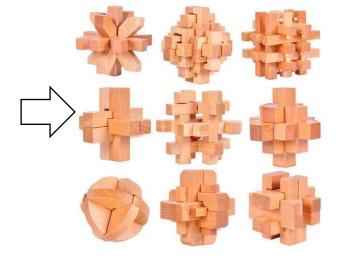
Cipher

- Ciphers are normally parameterized by keys
 - Message: x
 - Key: k1, k2
 - Encryption function: y=En(x, k1)
 - Decryption function: x=De(y, k2)
- Key is the secret
 - The encryption function and decryption function are public known



Cipher as a Secret?

Obtain the secret by unlocking the block



Not Scalable Not secure after the cipher is cracked

The mechanism of the locker is public known, but the key unknown



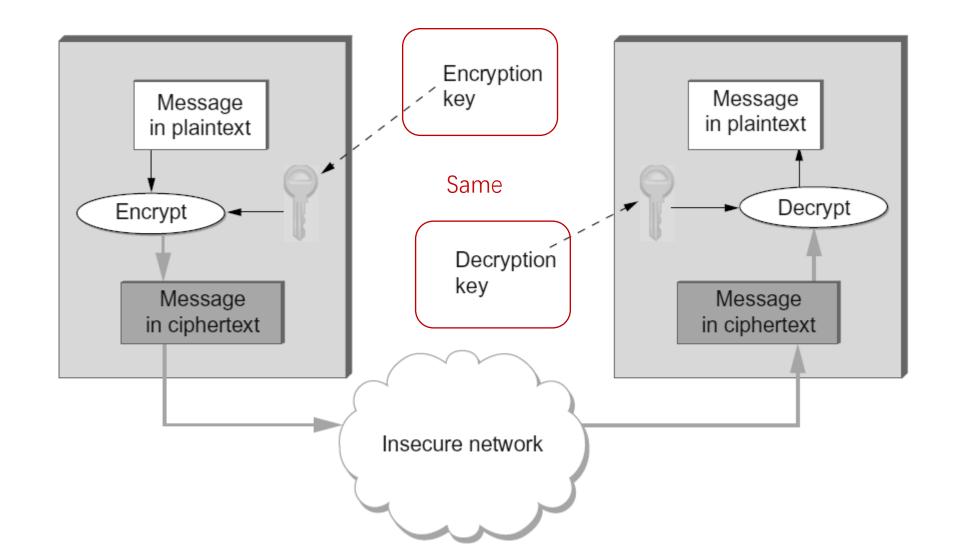






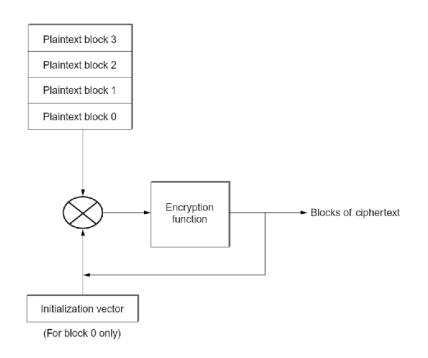


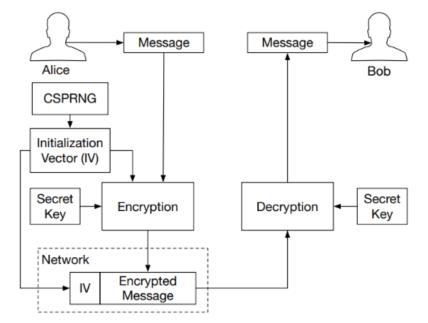




- Examples:
 - Advanced Encryption Standard (AES)
 - Block size: 4*4 = 16 Byte (128 bit)
 - Operation: a permutation of the 128 bits according to the key
 - key size: 128, 192, 256 bit
 - https://aesencryption.net/

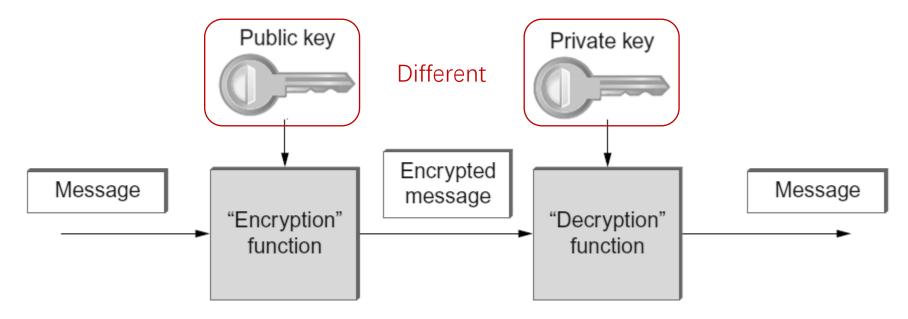
- Ciphers are under various attacks
 - e.g., word frequency, known plaintext, etc.
- Cipher designs
 - Prevent attackers from knowing key even the attacker knows plaintext
 - Initialization Vector (IV)
 - Cipher Block Chaining to prevent same output under same input



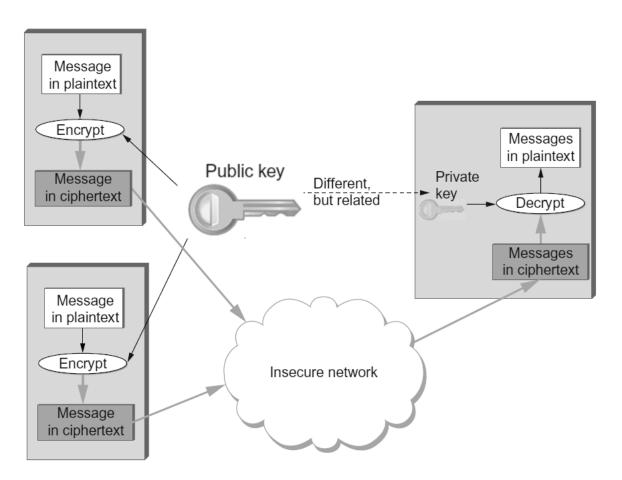


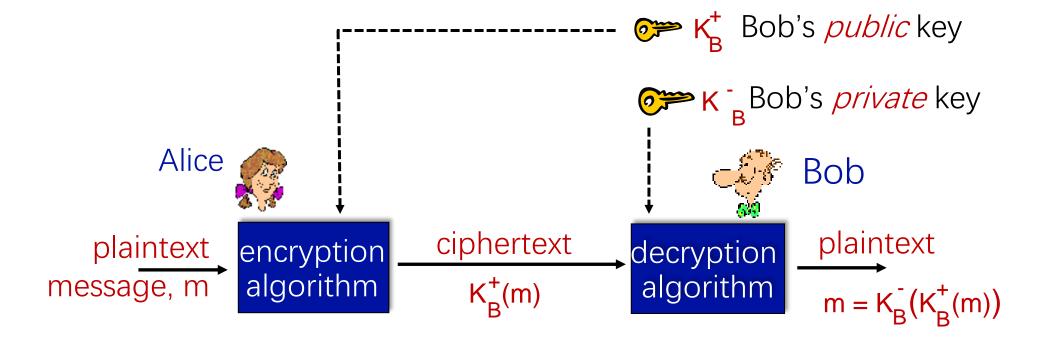
- Problem
 - Sender and receiver have to share the secret key
 - Q: how to agree on the key in first place (particularly if never "met")?
- This problem haven't been solved until very recently (70s)
 - -> Public-Key Cipher

- If the message is encrypted with the public key
 - The message can only be decrypted with the paired private key



• For key sharing: the public key can be released to everyone!





Requirements:

- 1 need $K_B^+(.)$ and $K_B^-(.)$ such that $K_B^-(K_B^+(m)) = m$
- 2 given public key K_B^+ , it should be impossible to compute private key K_B^-

- Example:
 - RSA (Rivest, Shamir, Adelson algorithm)
 - Elliptic Curve Cryptography

What is Network Security

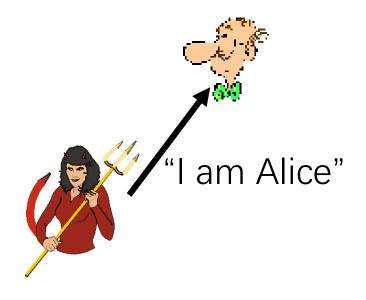
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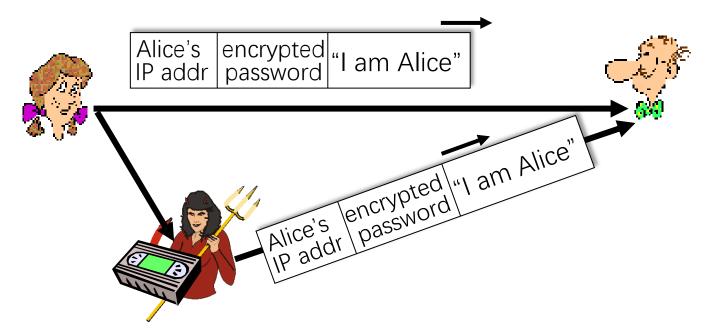
Goal: Bob wants Alice to "prove" her identity to him



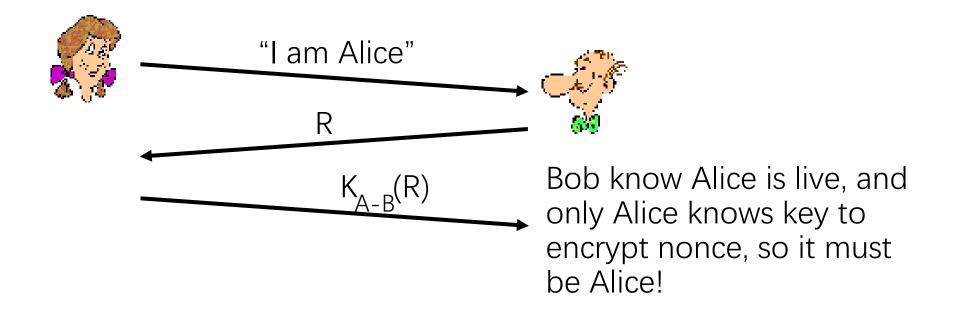




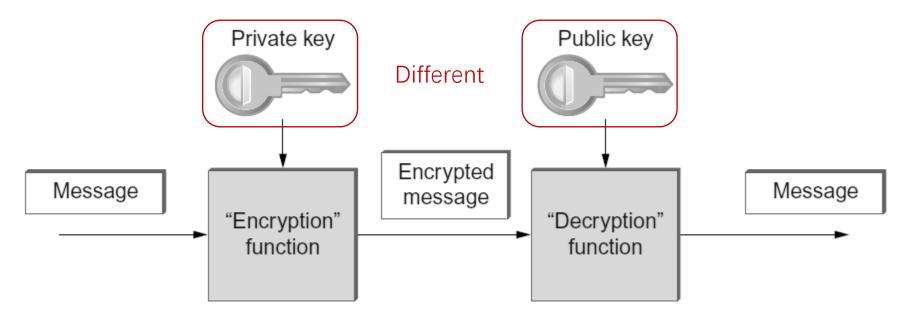
- Solution v1
 - Alice says "I am Alice" and sends her encrypted secret password to "prove" it.
 - Problem: replay



- Solution v2
 - + challenge with a nonce
 - Still need to share a symmetric key first



- If the message is encrypted with the <u>private</u> key
 - The message can only be decrypted with the paired <u>public</u> key



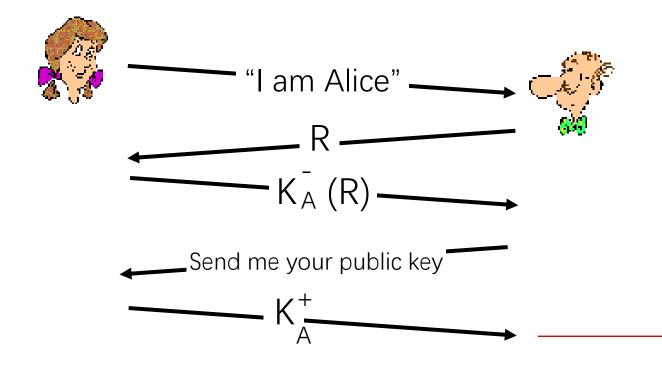
- Solution v3
 - Change to public cypher
 - Fact:

$$K_{\underline{B}}(K_{\underline{B}}(m)) = m = K_{\underline{B}}(K_{\underline{B}}(m))$$

use public key first, followed by private key use private key first, followed by public key

result is the same!

- Solution v3
 - Change to public cypher



Bob computes

$$K_A^+(K_A^-(R)) = R$$

and knows only Alice could have the private key, that encrypted R such that

$$K_A^+(K_A^-(R)) = R$$

- Solution v3
 - Still has a flaw: man in the middle!



Alice's public key

Trudy recovers Bob's m:

$$m = K_A (K_A (m)) \leftarrow K_A (m)$$

and she and Bob meet a week later in person and discuss m, not knowing Trudy knows m

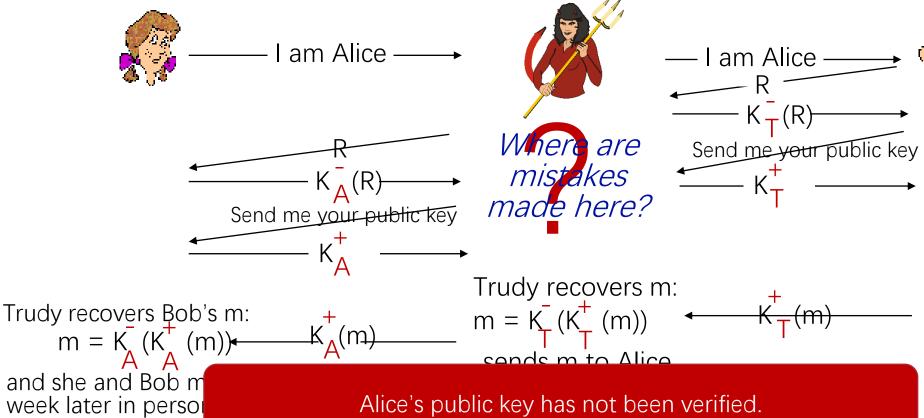
Bob sends a personal message, m to Alice

Solution v3

discuss m, not know

Trudy knows m

Still has a flaw: man in the middle!



Alice's public key has not been verified.

Bob computes authenticating Trudy as Alice

Bob sends a personal message, m to Alice

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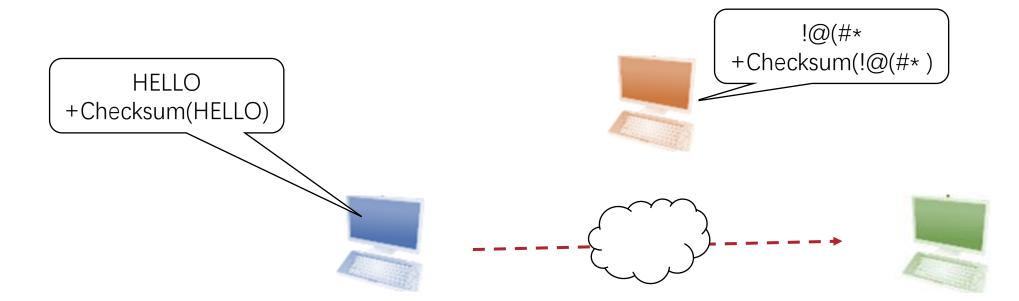
Timeliness

To identify delayed messages

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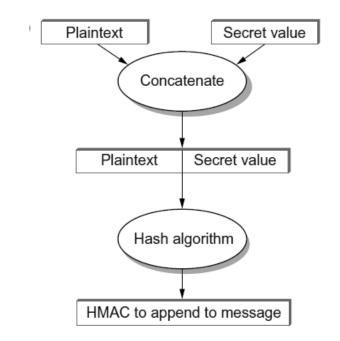
Data Integrity: Checksum

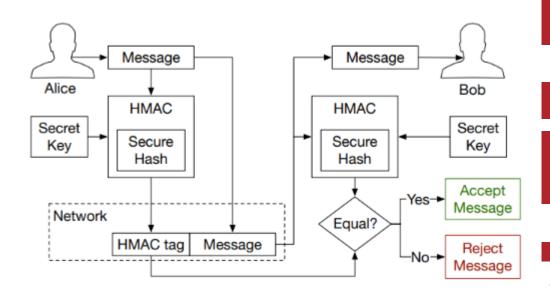
• Checksum can be replicated



Cryptographic Hash

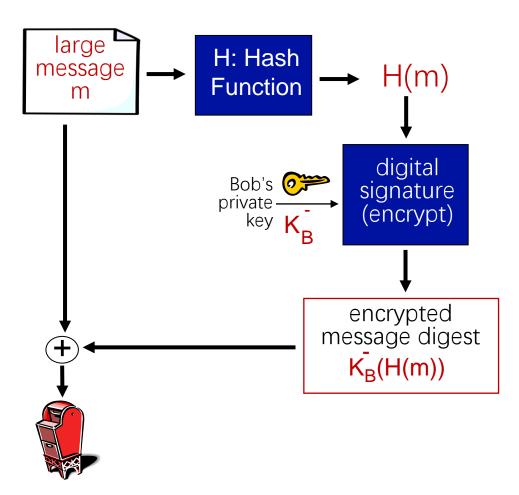
- Cryptographic Hash
 - Example
 - MD5
 - SHA
- HMAC
 - Hash Massage Authentication Code
 - Use Cryptographic Hash Function to generate integrity and authentication check for the message.
- Digital Signature
 - Fixed-length, easy- to-compute digital "fingerprint"
 - Apply hash function H to m, get fixed size message digest, H(m)
 - use private key to encrypt the hash



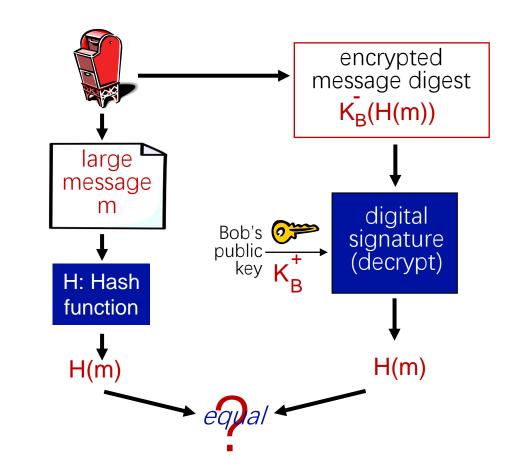


Digital Signature

Bob sends digitally signed message:



Alice verifies signature, integrity of digitally signed message:

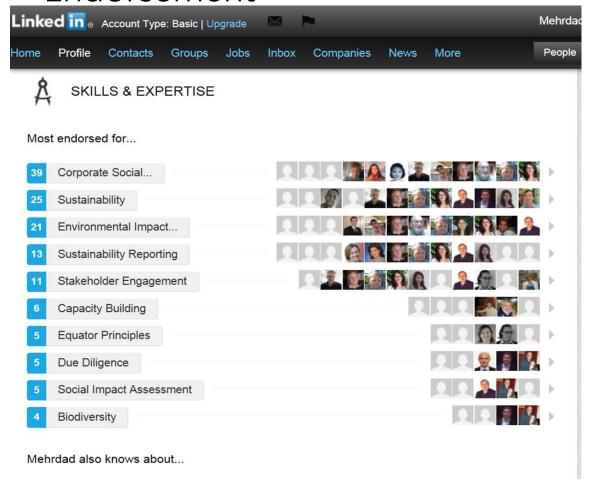


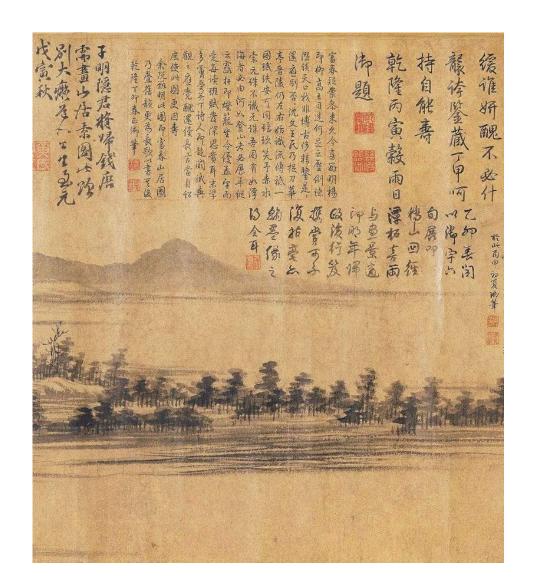
Key Predistribution

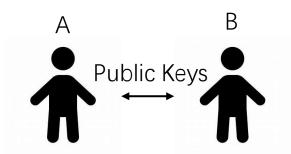
- Distribute through Offline Channel
 - Not scalable



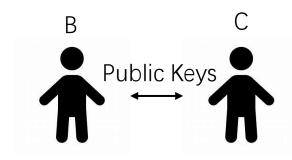
Endorsement



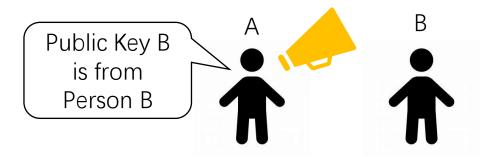




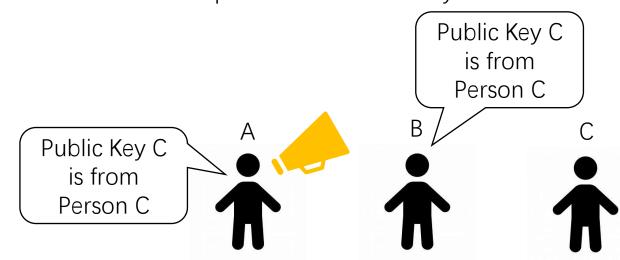
Step 1. Verify Each Other Offline; Exchange Public Keys



Step 3. Verify Each Other Offline; Exchange Public Keys



Step 2. Certifies Public Keys

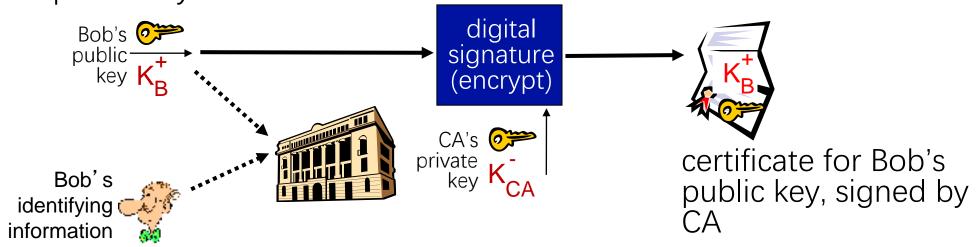


Step 4. Certifies Public Keys from Others

- Certificate Authority (CA)
 - Preinstall trusted public keys
- Web of Trust
 - Collect public keys from known people

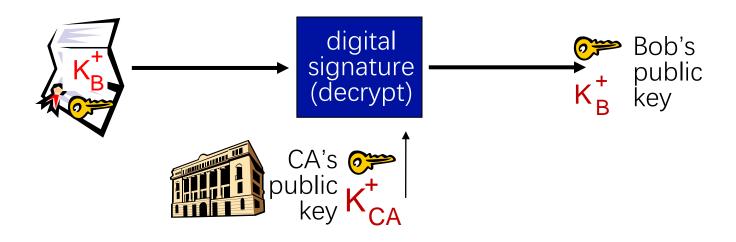
Public-Key Certification Authorities (CA)

- Certification authority (CA): binds public key to particular entity E
- Entity (person, website, router) registers its public key, provides "proof of identity" to CA
 - CA creates certificate binding identity E to E's public key
 - Certificate containing E's public key digitally signed by CA: CA says "this is E's public key"

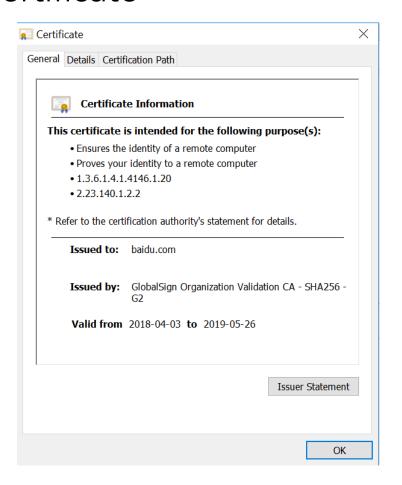


Public-Key Certification Authorities (CA)

- When Alice wants Bob's public key:
 - gets Bob's certificate (from Bob or elsewhere)
 - apply CA's public key to Bob's certificate, get Bob's public key



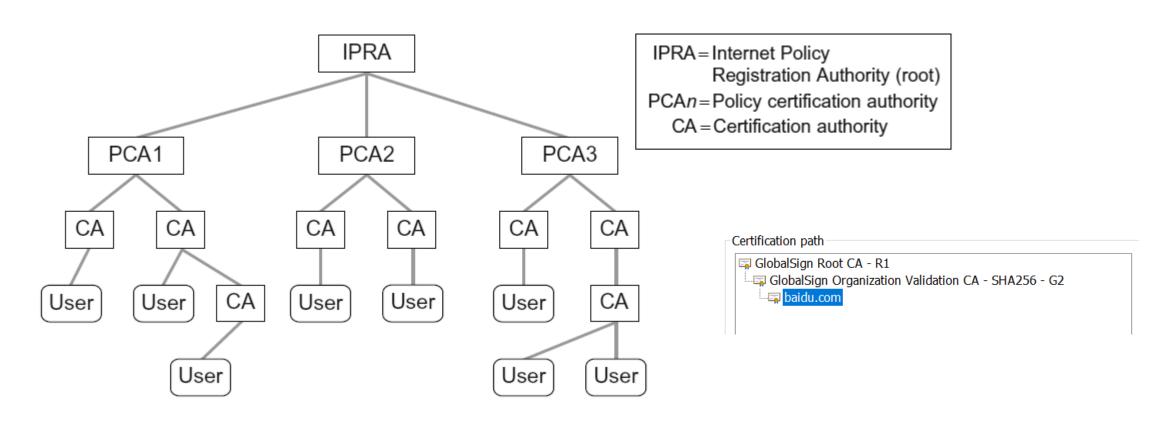
Certificate

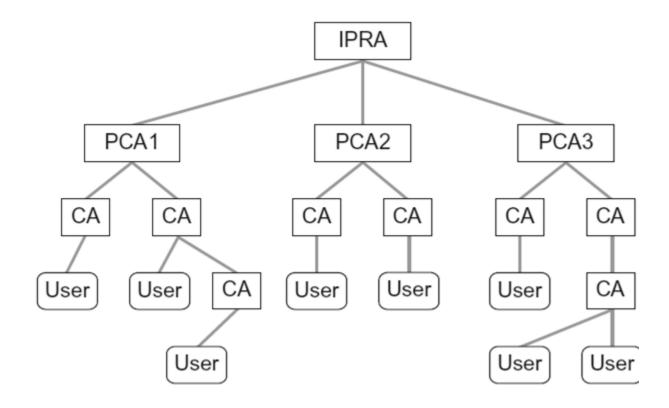


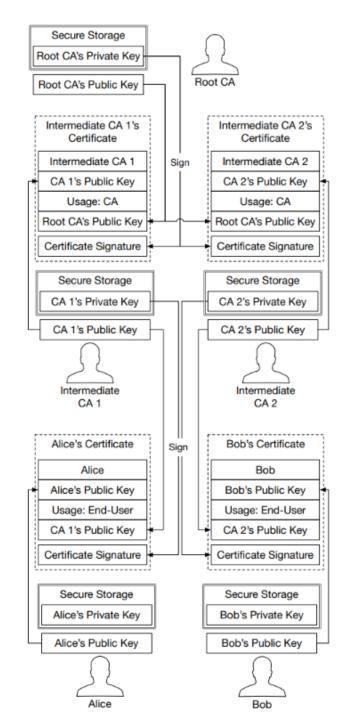
Contains

- The identity of the entity being certified
- The public key of the entity being certified
- The identity of the signer
- The digital signature of the signer
- A digital signature algorithm identifier (which cryptographic hash and which cipher)

Certificate Authority (CA)







Demo

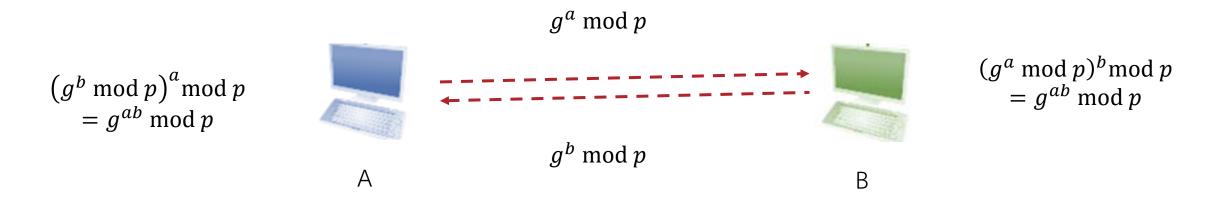
- Certificate Authority (CA)
 - certmgr.msc
 - https://www.sinorailca.com/

Symmetric-Key Predistribution

- Through Trust Server
- Through Public-Key Predistribution

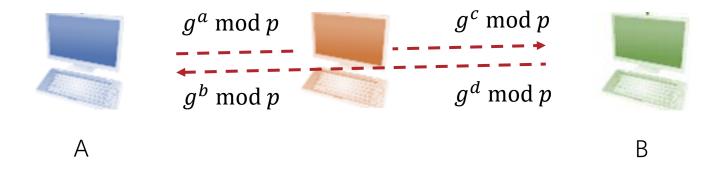
Diffie-Hellman Key Exchange

- Generate shared key without key predistribution
 - a is the secret of A
 - b is the secret of B
 - g and p are publicly known
 - g^ab mod p is the shared key



Diffie-Hellman Key Exchange

- Man in the middle attack
 - A cannot authenticate he is talking with B
- Diffie-Hellman Key Exchange is not secure without authentication



Reference

- Textbook 8.1, 8.2, 8.3
- Some slides are adapted from http://www-net.cs.umass.edu/kurose_ross/ppt.htm by Kurose Ross