



CS271 Computer Graphics II

Lecture 6

Mesh Simplification

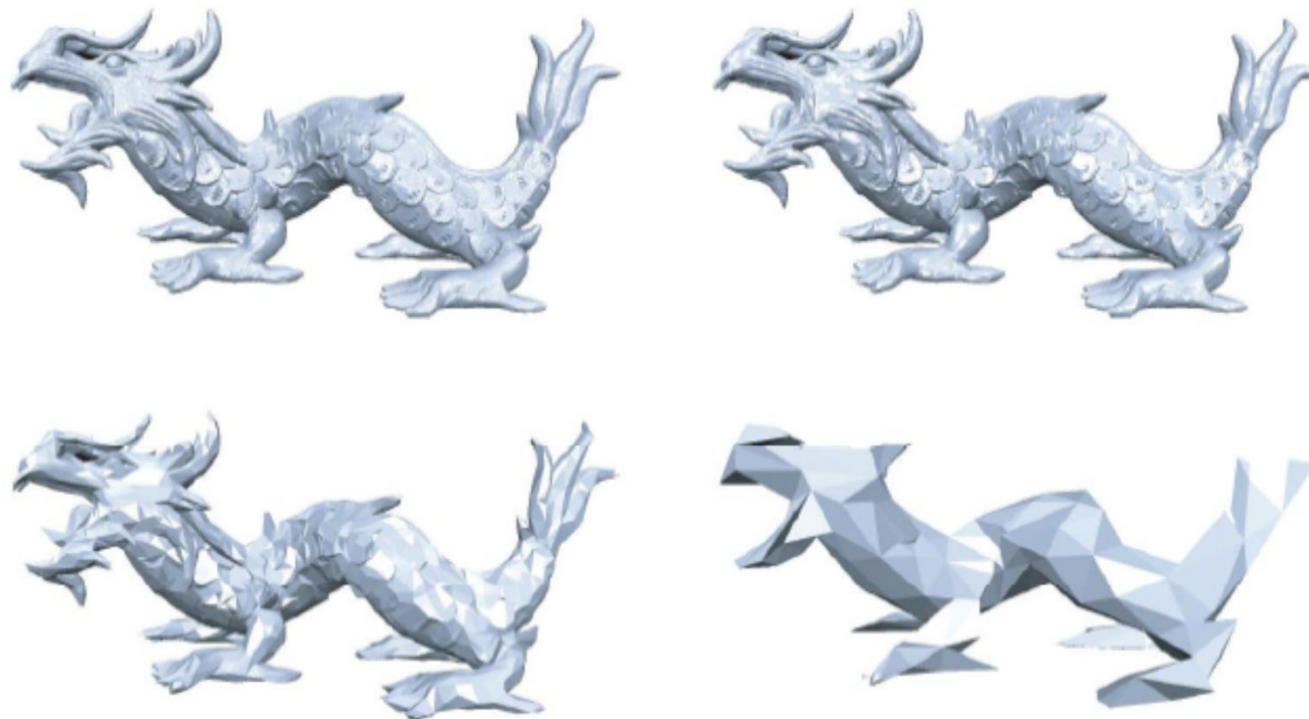


Definition

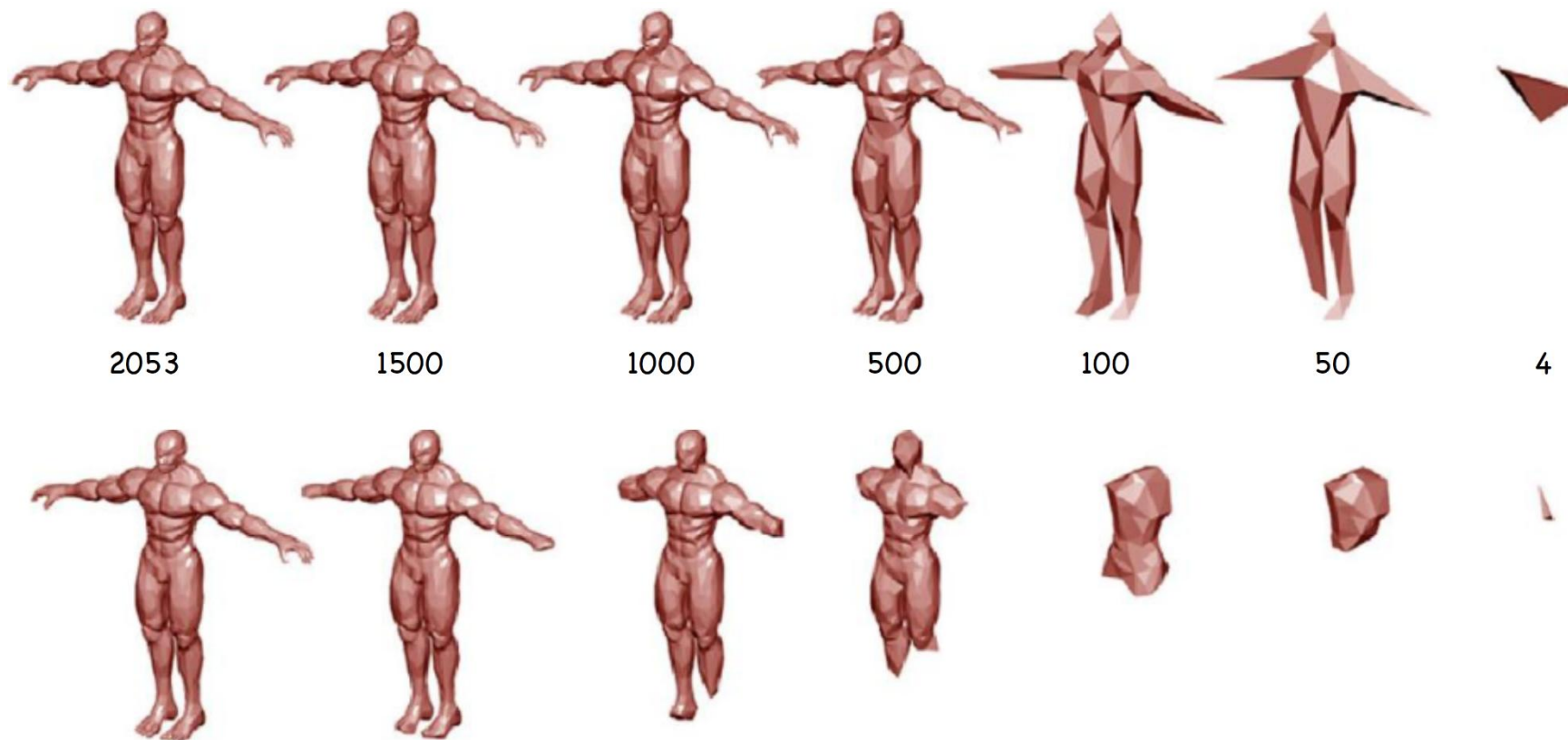


Simplification, decimation, approximation, downsampling

- Transform a given polygonal mesh into another mesh with fewer faces, edges, and vertices
- The simplification or approximation procedure is usually controlled by user-defined quality criteria



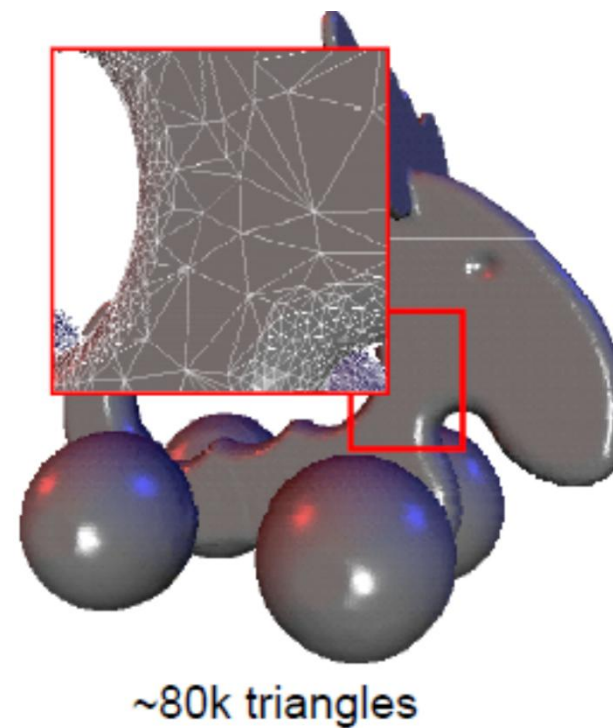
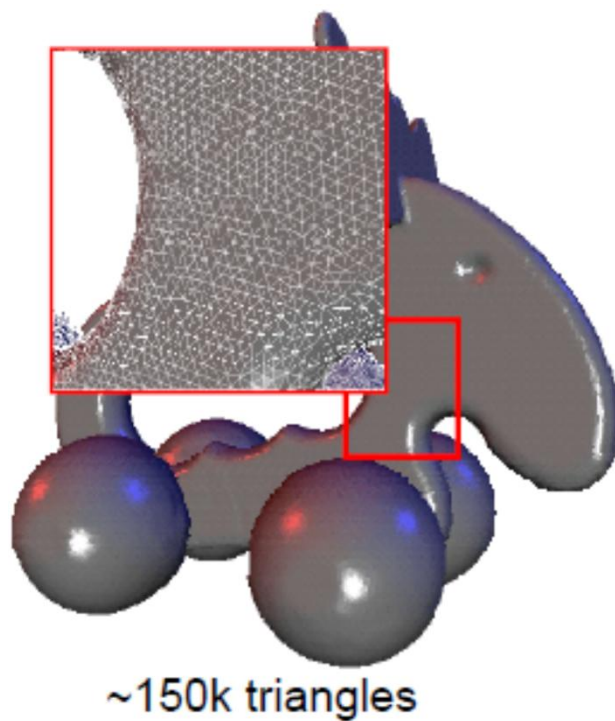
Curvature-preserved vs. Curvature-removed Criteria



Mesh Simplification Applications



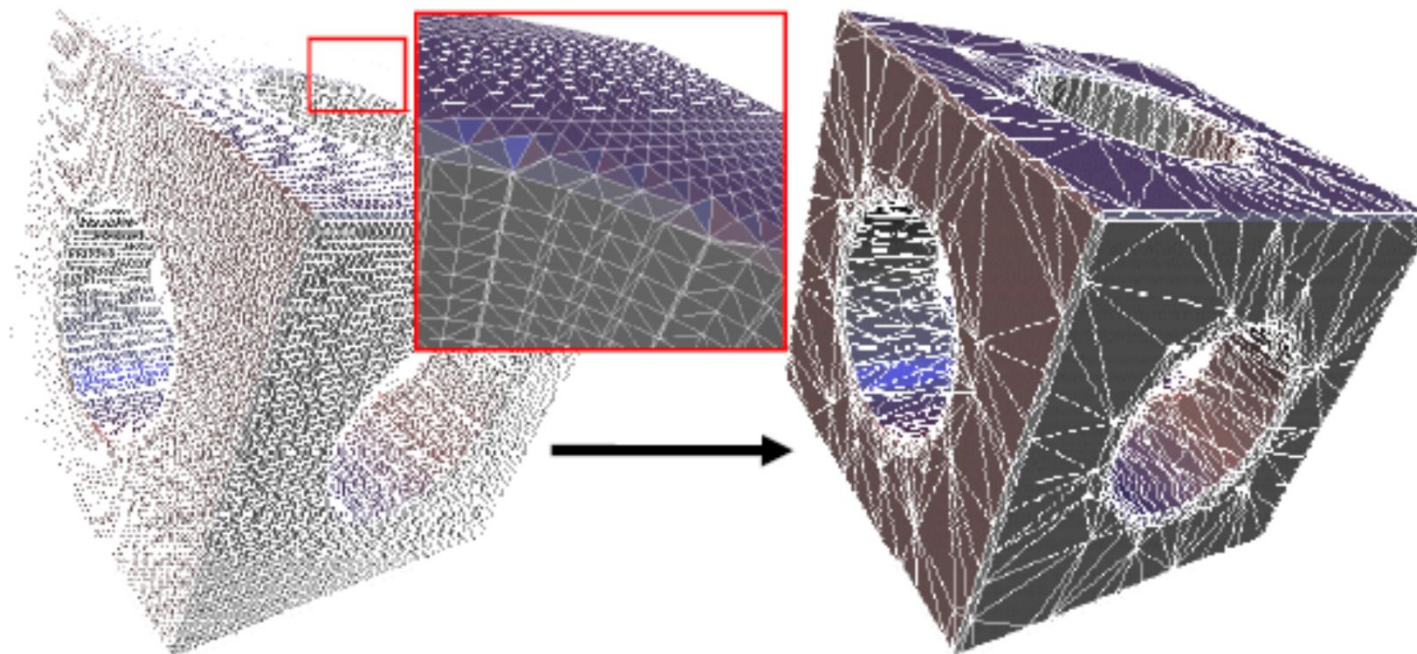
- Oversampled 3D scan data



Mesh Simplification Applications



- Over-tessellation: e.g., iso-surface extraction

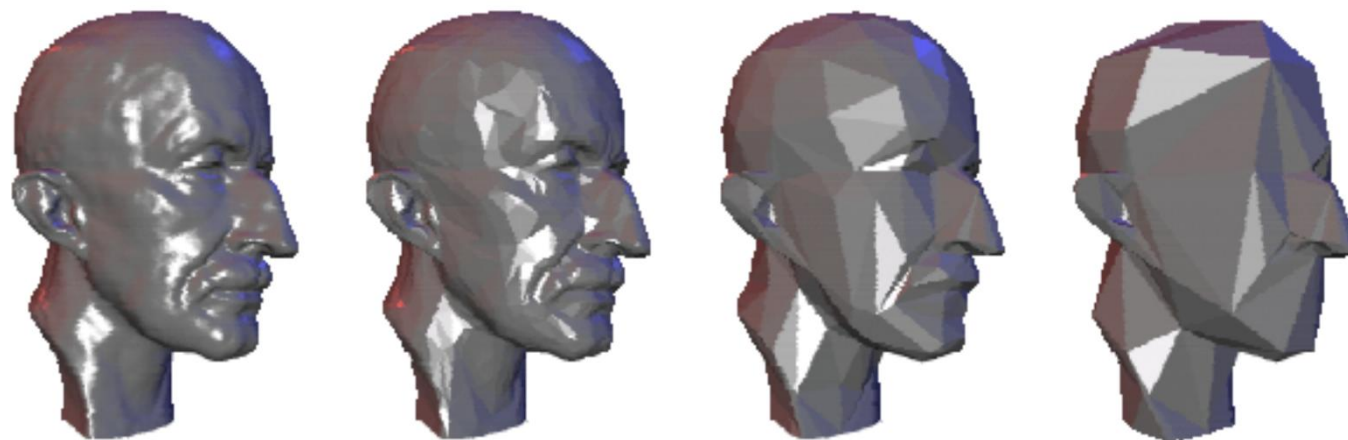


Mesh Simplification Applications



Multi-resolution hierarchies for

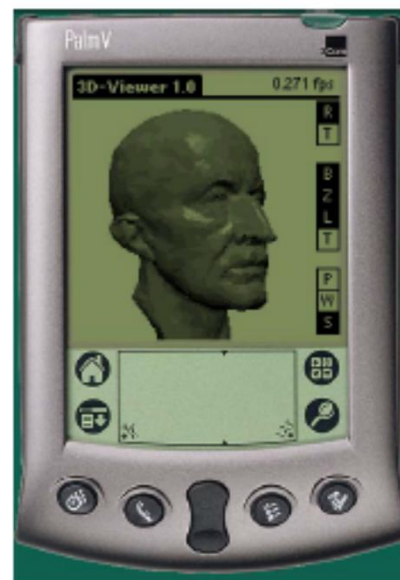
- Efficient geometry processing
- Level-of-detail (LOD) rendering



Mesh Simplification Applications



- Adaption to hardware capabilities





Mesh Simplification



Adjust the complexity of a geometry data set

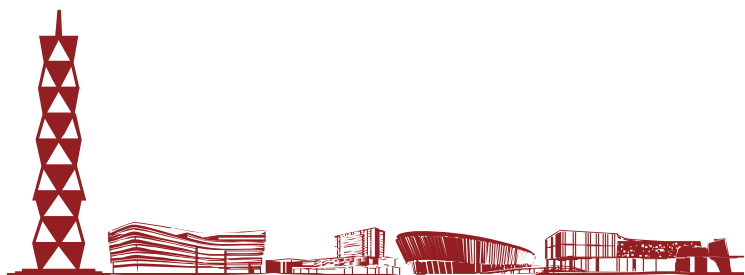
- Since many decimation schemes work iteratively, i.e., they decimate a mesh by removing one vertex at a time, they usually can be inverted
- Hierarchical method

Problem Statement

- Given: $M = (V, F)$
- Find: $M' = (V', F')$ such that
- $|V'| = n < |V|$ and $\|M - M'\|$ is minimal, or
- $\|M - M'\| < \varepsilon$ and $|V'|$ is minimal

Respect additional fairness criteria

- Normal deviation, triangle shape, scalar attributes, etc.



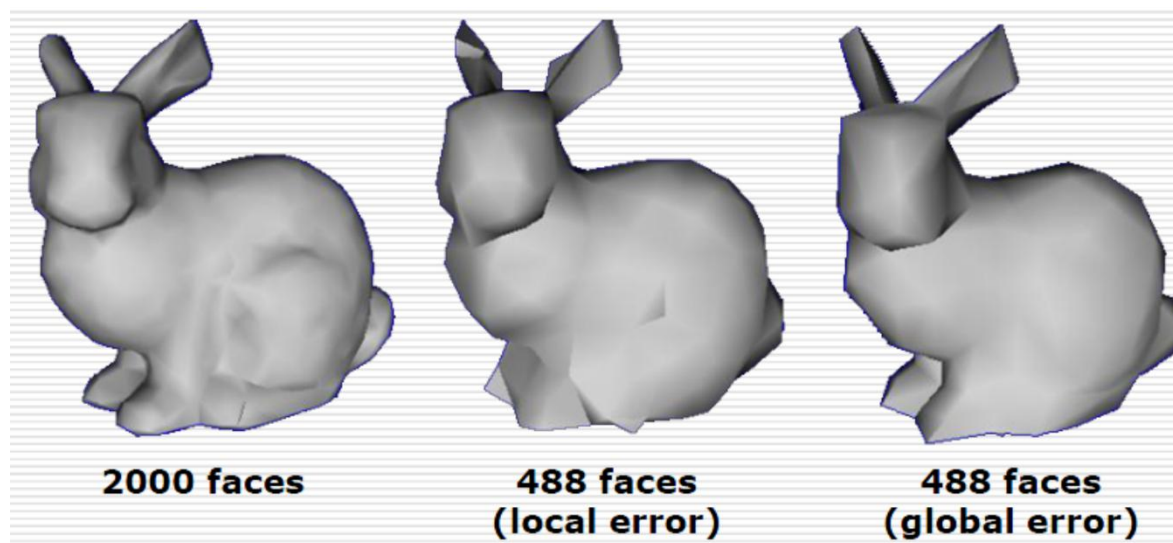


Mesh Simplification



Start with the original fine mesh

- **Simply progressively**
E.g., collapse edges, vertex clustering
- **Aim to keep original appearance**
Normal deviation, triangle shape, scalar attributes
Error control





Local Operations

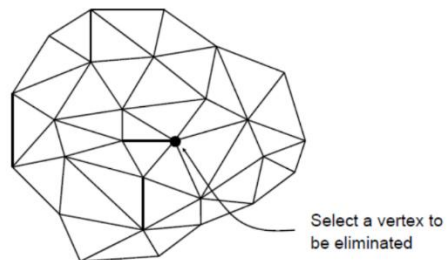
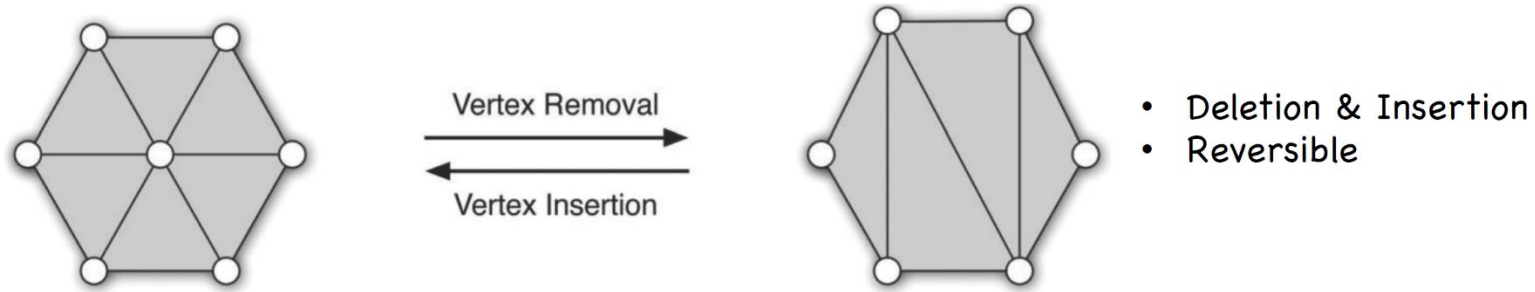




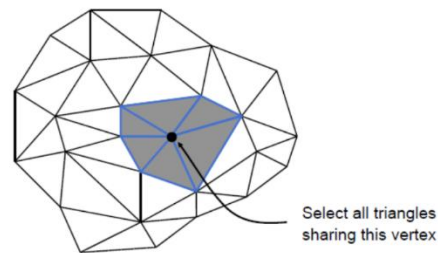
Local Simplification Operator



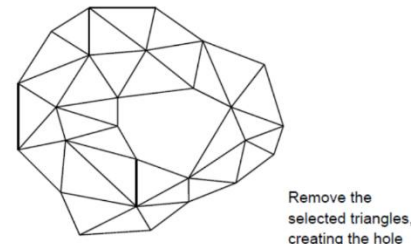
- Vertex Removal



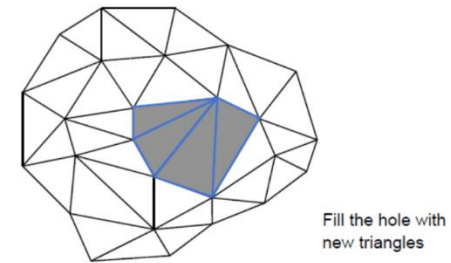
Select a vertex to be eliminated



Select all triangles sharing the vertex



Remove the selected triangles, creating the hole



Fill the hole with new triangles

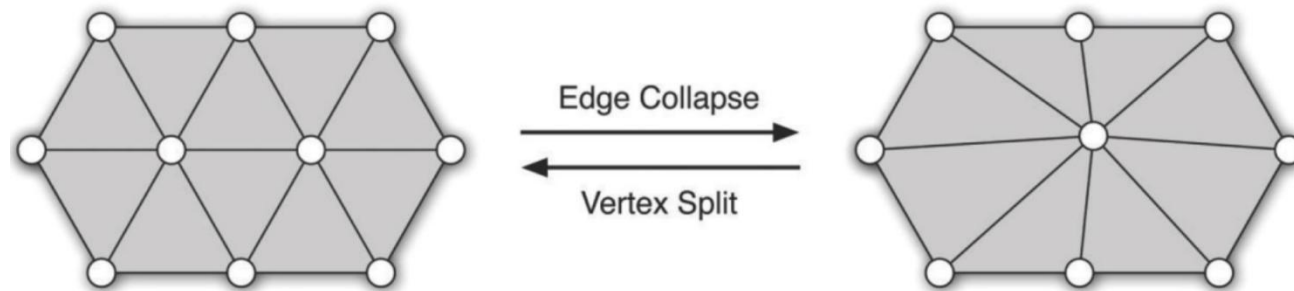




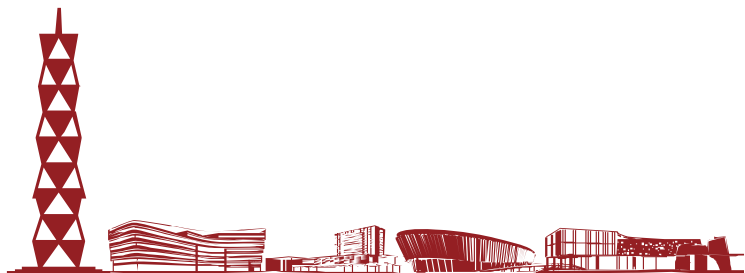
Local Simplification Operator



- Edge Collapse



- Merge two adjacent vertices
- Simple to implement
- Well-suited for implementing geomorphing

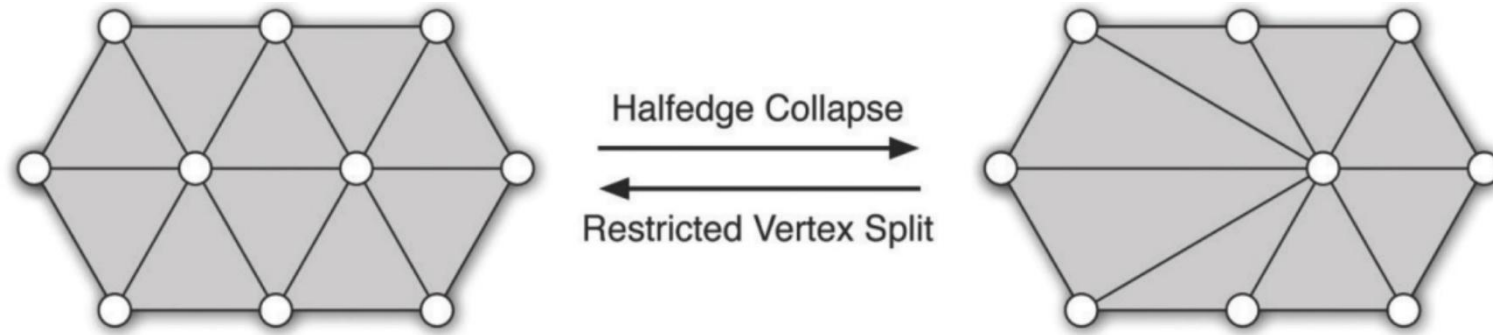




Local Simplification Operator



- Half-edge Collapse



- **Collapse edge into one end point**

Special case of vertex removal

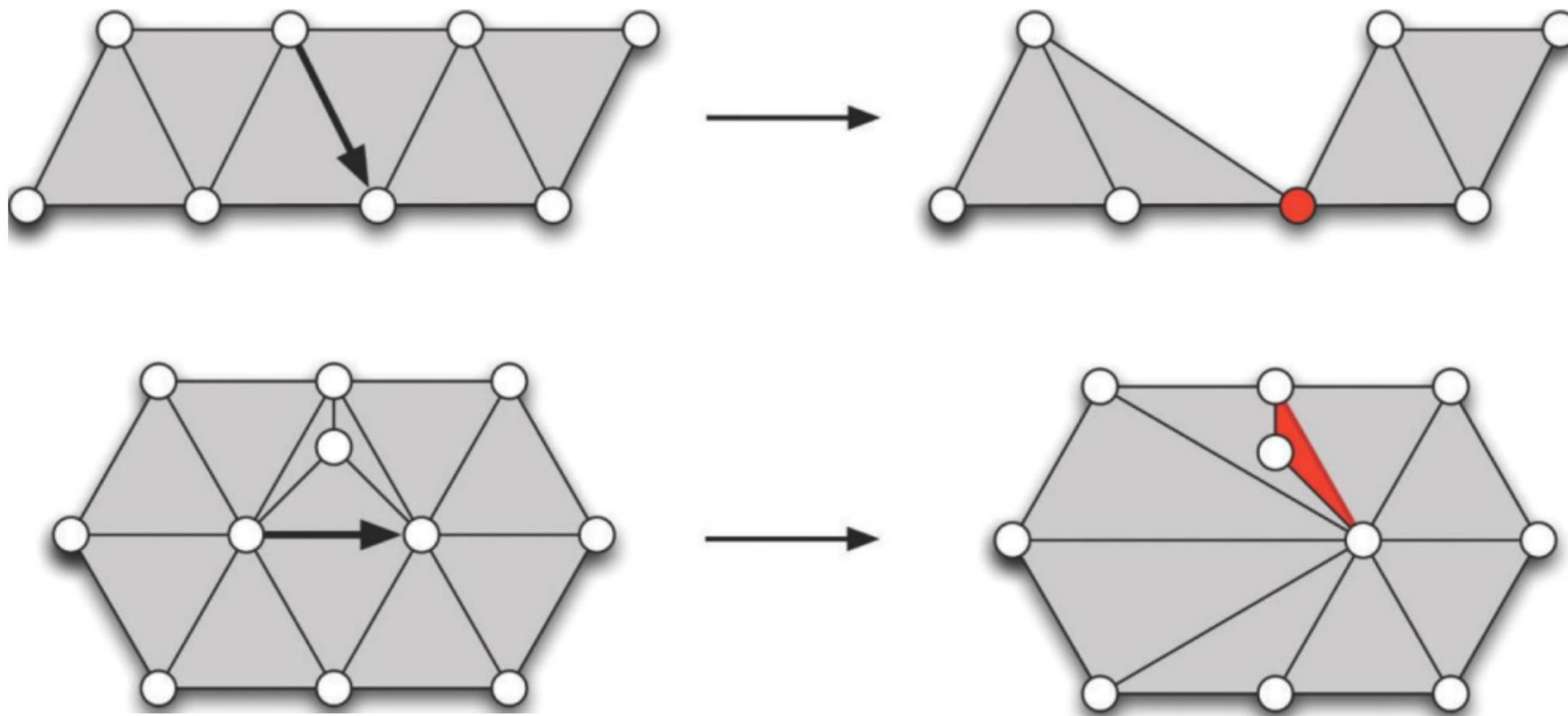
Special case of edge collapse

- After collapse: $n(E) - 3$, $n(V) - 1$, $n(F) - 2$
- According to Euler Formula: unchanged
- Half-edge collapsing would not change the genus of a mesh
- Should determine whether collapse is ok (may introduce non-manifold structure)





Topologically Illegal (half-)edge Collapses





Simplification via Edge Collapse



One popular scheme: iteratively collapse edges

Greedy algorithm from a general overview:

- Assign each edge a cost
- Collapse edge with least cost
- Repeat until target number of elements is reached
- Particularly effective cost function: **Quadric Error Metric**





Quadric Error Metric



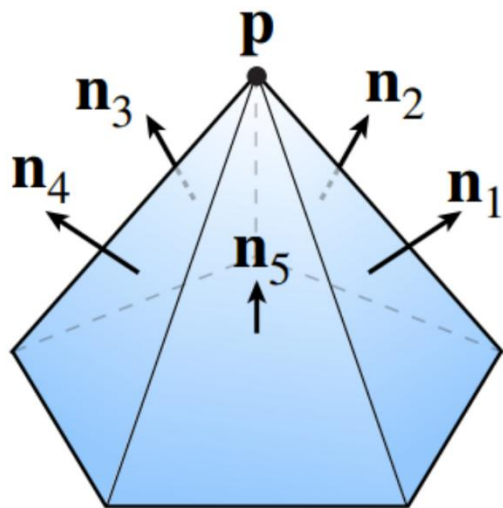
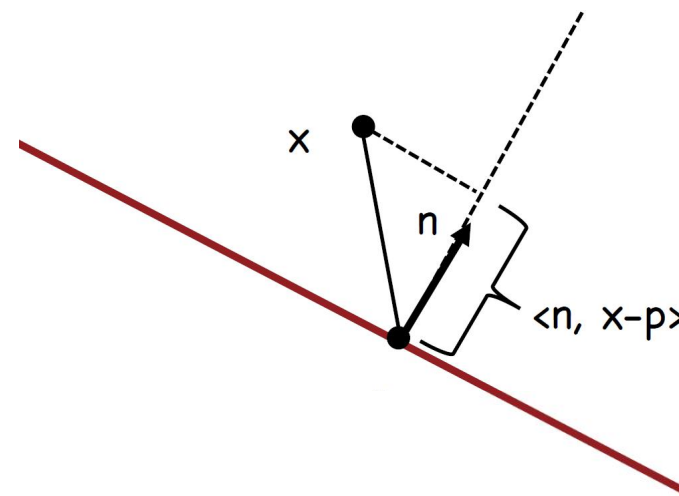


Quadric Error Metric (QEM)



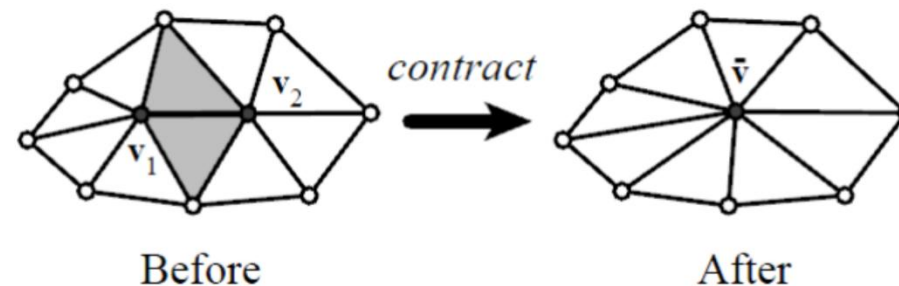
Approximate distance to a collection of triangles

- Q: Distance to plane w/ normal n passing through point p ?
- A: $\text{dist}(x) = \langle n, x \rangle - \langle n, p \rangle = \langle n, x - p \rangle$
- Quadric error is then sum of squared point-to-plane distances



$$Q(x) := \sum_{i=1}^k \langle n_i, x - p \rangle^2$$

$$Q^e = Q_1^v + Q_2^v$$





Quadric Error – Homogeneous Coordinates



Suppose in coordinates we have

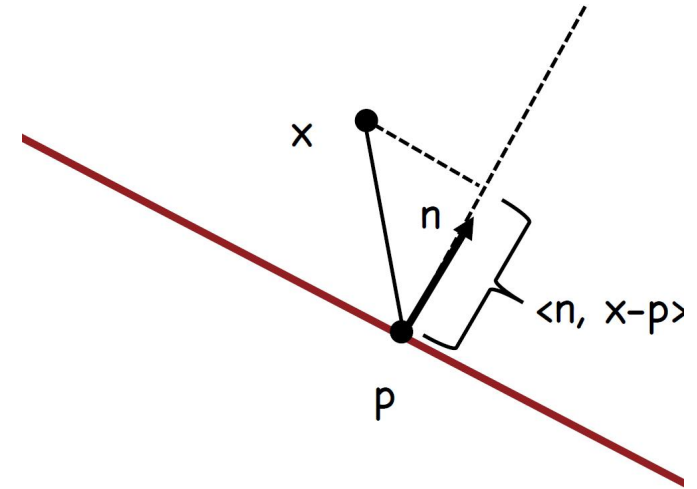
- A query point $\mathbf{x} = (x, y, z)$
- A normal $\mathbf{n} = (a, b, c)$
- An offset $\mathbf{d} := \langle \mathbf{n}, \mathbf{p} \rangle$

In **homogeneous coordinates**, let

- $\mathbf{u} := (x, y, z, 1)$
- $\mathbf{v} := (a, b, c, d)$
- Signed distance to plane is then just $\langle \mathbf{u}, \mathbf{v} \rangle = ax + by + cz + d$
- Squared distance is $\langle \mathbf{u}, \mathbf{v} \rangle^2 = \mathbf{u}^T(\mathbf{v}\mathbf{v}^T)\mathbf{u} =: \mathbf{u}^T\mathbf{K}\mathbf{u}$
- Matrix $\mathbf{K} = \mathbf{v}\mathbf{v}^T$ encodes squared distance to plane

Key idea: sum of matrices \mathbf{K} distance to union of planes

$$\mathbf{u}^T\mathbf{K}_1\mathbf{u} + \mathbf{u}^T\mathbf{K}_2\mathbf{u} = \mathbf{u}^T(\mathbf{K}_1 + \mathbf{K}_2)\mathbf{u}$$



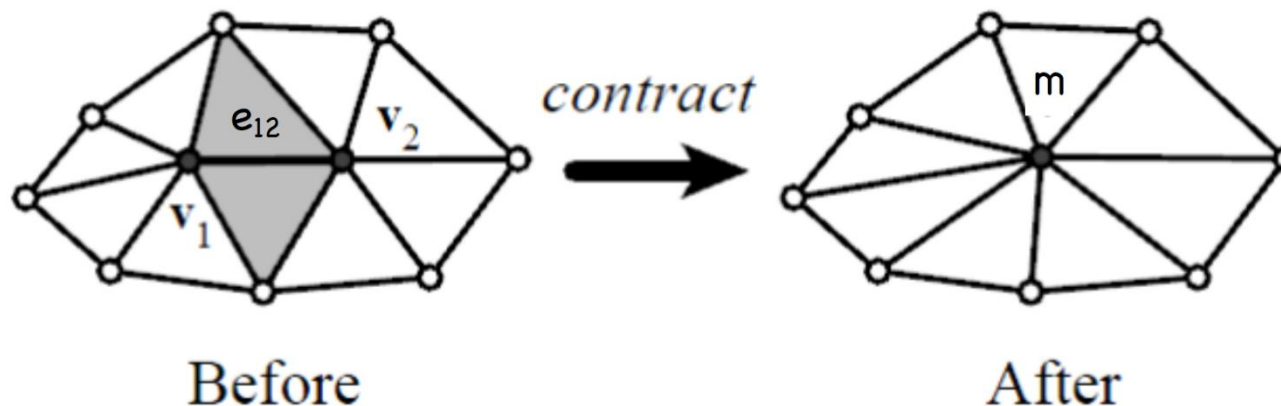
$$\mathbf{K} = \begin{bmatrix} a^2 & ab & ac & ad \\ ab & b^2 & bc & bd \\ ac & bc & c^2 & cd \\ ad & bd & cd & d^2 \end{bmatrix}$$



Quadric Error of Edge Collapse



- How much does it cost to collapse an edge e_{12} ?
Idea: compute midpoint m , measure error $Q(m) = m^T(K_1 + K_2)m$
- Error becomes “score” for e_{12} , determining priority



- Better idea: find point x that minimize error!
- But how to minimize quadric error?



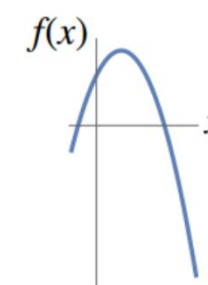
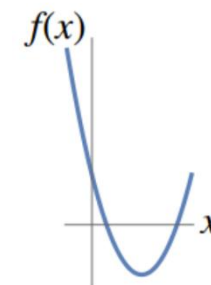
Revisit: Minimizing a Quadratic Function



Suppose you have a function $f(x) = ax^2 + bx + c$

- Q: What does the graph of this function look like?
- Q: How do we find the minimum?
- A: Find where the function looks “flat” if we zoom in really close
- i.e., find point x where 1st derivative vanishes:

$$\begin{aligned}f'(x) &= 0 \\2ax + b &= 0 \\x &= -b/2a\end{aligned}$$





Minimizing Quadratic Polynomial



Not much harder to minimize a quadratic polynomial in **n** variables

- Can always write in terms of a symmetric matrix **A**
- E.g., in 2D: $\mathbf{f}(\mathbf{x}, y) = \mathbf{a}\mathbf{x}^2 + \mathbf{b}\mathbf{x}\mathbf{y} + \mathbf{c}\mathbf{y}^2 + \mathbf{d}\mathbf{x} + \mathbf{e}\mathbf{y} + \mathbf{g}$

$$\mathbf{x} = \begin{bmatrix} x \\ y \end{bmatrix} \quad \mathbf{A} = \begin{bmatrix} a & b/2 \\ b/2 & c \end{bmatrix} \quad \mathbf{u} = \begin{bmatrix} d \\ e \end{bmatrix}$$

$$\mathbf{f}(\mathbf{x}, y) = \mathbf{x}^T \mathbf{A} \mathbf{x} + \mathbf{u}^T \mathbf{x} + \mathbf{g}$$

(will have the same form for any **n**)

- Q: How do we find a critical point (min/max/saddle)?
- A: Set derivative to zero!

$$2\mathbf{A}\mathbf{x} + \mathbf{u} = 0$$
$$\mathbf{x} = -1/2 \mathbf{A}^{-1}\mathbf{u}$$





Positive Definite Quadratic Form

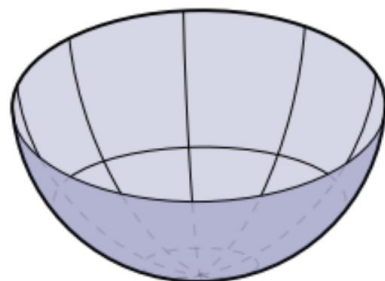


Just like our 1D parabola, critical point is not always a min!

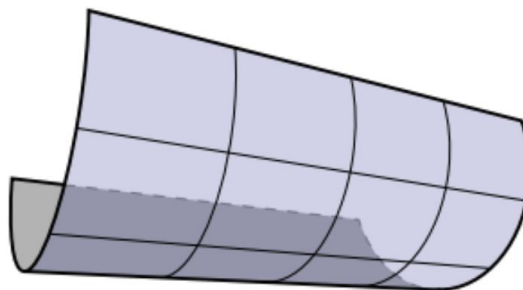
- Q: In 2D, 3D, nD, when do we get a minimum?
- A: When matrix A is positive-definite:

$$\mathbf{X}^T \mathbf{A} \mathbf{X} > 0 \quad \forall \mathbf{X}$$

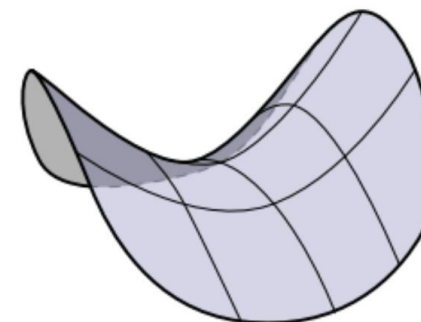
- **1D:** Must have $xax = ax^2 > 0$, i.e., a is positive!
- **2D:** Graph of function looks like a “bowl”:



positive definite



positive semidefinite



indefinite





Minimizing Quadric Error



Find “best” point for edge collapse by minimizing quadratic form

$$\min_{\mathbf{u} \in \mathbb{R}^4} \mathbf{u}^T K \mathbf{u}$$

- Already know 4th (homogeneous) coordinate for a point is 1
- So, break up our quadratic function into two pieces:

$$\begin{bmatrix} \mathbf{x}^T & 1 \end{bmatrix} \begin{bmatrix} B & \mathbf{w} \\ \mathbf{w}^T & d^2 \end{bmatrix} \begin{bmatrix} \mathbf{x} \\ 1 \end{bmatrix} \\ = \mathbf{x}^T B \mathbf{x} + 2\mathbf{w}^T \mathbf{x} + d^2$$

- Now we have a quadratic polynomial in the unknown position $\mathbf{x} \in \mathbb{R}^3$
- Can minimize as before:

$$2B\mathbf{x} + 2\mathbf{w} = 0 \quad \Longleftrightarrow \quad \mathbf{x} = -B^{-1}\mathbf{w}$$

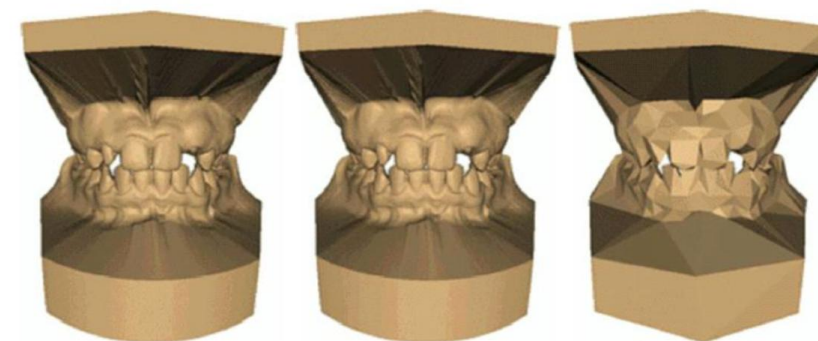




QEM Simplification: Final Algorithm



- Input: a mesh
- Output: a simplified mesh
- Initialization:
 - Compute \mathbf{K} for each triangle (squared distance to plane)
 - Set \mathbf{K}_i at each vertex to sum of \mathbf{K} s from incident triangles
 - For each edge \mathbf{e}_{ij} :
 - Set $\mathbf{K}_{ij} = \mathbf{K}_i + \mathbf{K}_j$
 - Find point x minimizing error, set cost to $\mathbf{K}_{ij}(x)$
- Until we reach target number of triangles:
 - Collapse edge \mathbf{e}_{ij} with smallest cost to optimal point x
 - Set quadric at new vertex to \mathbf{K}_{ij}
 - Update cost of edges touching new vertex



Full
Resolution

60,000
triangles

1000
triangles





Variational Shape Approximation

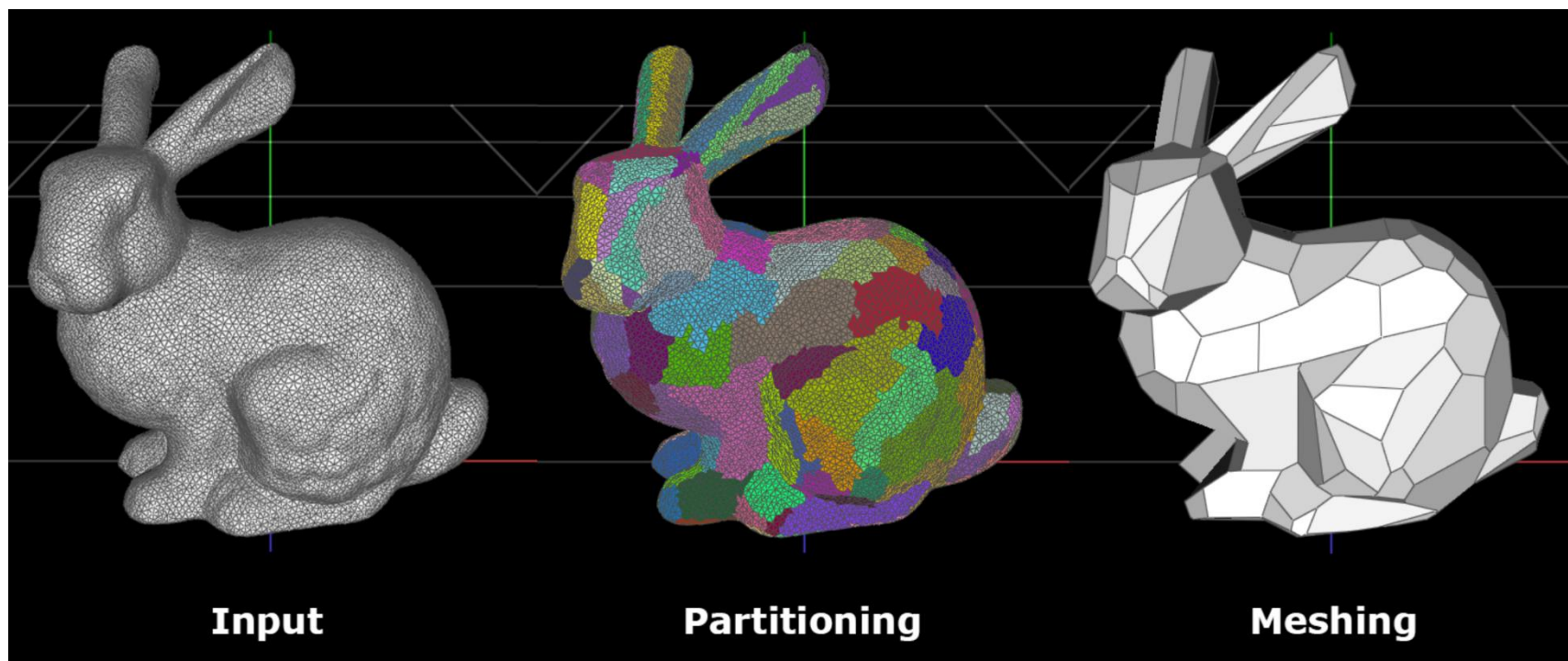




Variational Shape Approximation (VSA)



VSA is highly sensitive to features and symmetries and produces anisotropic meshes of high approximation quality

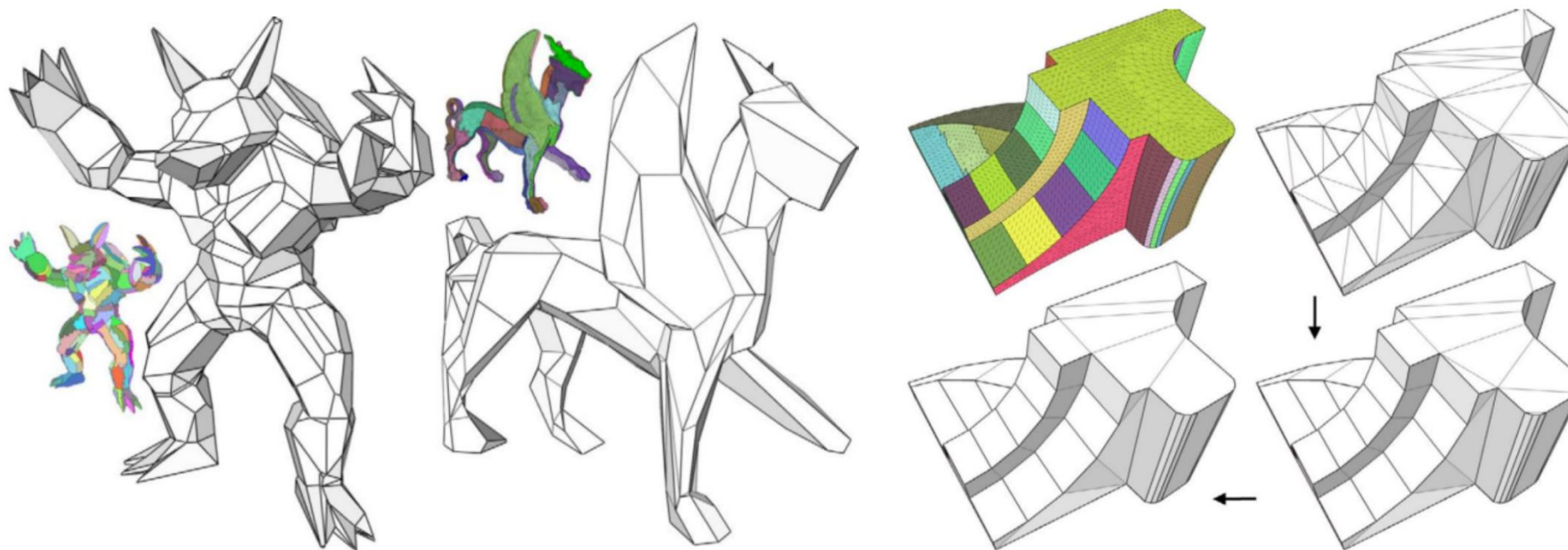




Variational Shape Approximation (VSA)



- The input shape is approximated by a set of proxies
- A plane in space through the point x_i with normal direction n_i





Region Representation



- **M**: a triangle mesh
- **R** = {**R**₁, ..., **R**_k}: a partition of **M** into **k** regions $R_1 \cup \dots \cup R_k = M$
- Proxies: **P** = {**P**₁, ..., **P**_k}, **P**_i = (**x**_i, **n**_i)

Distance metrics between **R**_i and **P**_i

- The squared orthogonal distance of **x** from the plane **P**_i

$$L^2(R_i, P_i) = \int_{x \in R_i} (n_i^T x - n_i x_i)^2 dA$$

- A measure of the normal field:

$$L^{2,1}(R_i, P_i) = \int_{x \in R_i} \|n(x) - n_i\|^2 dA$$





Goal of VSA



Given a number k and an error metric $E(\mathbf{L}_2 \text{ or } \mathbf{L}_{2,1})$, find a set $\mathbf{R} = \{\mathbf{R}_1, \dots, \mathbf{R}_k\}$ of regions and a set $\mathbf{P} = \{\mathbf{P}_1, \dots, \mathbf{P}_k\}$ of proxies such that the global distortion

$$E(R, P) = \sum_{i=1}^k E(R_i, P_i)$$

is minimized





Lloyd's Clustering Algorithm



- The algorithm iteratively alternates between a **geometry partitioning** phase and a **proxy fitting** phase
- **Geometry partitioning phase**
 - A set of regions that best fit a **given** set of proxies
 - Modifies the set R of regions to achieve a lower approximation error while keeping the proxies P fixed
- **Proxy fitting phase**
 - The partitioning is kept **fixed**, and the proxies are adjusted to minimize approximation error
 - L^2 metric: the best proxy is the least-squares fitting plane
 - $L^{2,1}$ metric: the proxy normal n_i is just the area-weighted average of the triangle normals
- **Initialization**
 - Randomly picking k triangles as R
 - The planes of k triangles are used to initialize P





Related Papers



- Garland M, Heckbert P S. **Surface simplification using quadric error metrics**[C]//Proceedings of the 24th annual conference on Computer graphics and interactive techniques. 1997: 209-216.
- Cohen-Steiner D, Alliez P, Desbrun M. **Variational shape approximation**[M]//ACM SIGGRAPH 2004 Papers. 2004: 905-914.
- Skrodzki M, Zimmermann E, Polthier K. **Variational shape approximation of point set surfaces**[J]. Computer Aided Geometric Design, 2020, 80: 101875.
- Liu Y J, Xu C X, Fan D, et al. **Efficient construction and simplification of Delaunay meshes**[J]. ACM Transactions on Graphics (TOG), 2015, 34(6): 1-13.
- Yi R, Liu Y J, He Y. **Delaunay mesh simplification with differential evolution**[J]. ACM Transactions on Graphics (TOG), 2018, 37(6): 1-12.
- Liang Y, He F, Zeng X. **3D mesh simplification with feature preservation based on whale optimization algorithm and differential evolution**[J]. Integrated Computer-Aided Engineering, 2020, 27(4): 417-435.
- Liu H T D, Gillespie M, Chislett B, et al. **Surface simplification using intrinsic error metrics**[J]. ACM Transactions on Graphics, 2023, 42(4).

