

# Reflection Matrices

A matrix  $\mathbf{H} \in \mathbb{R}^{m \times m}$  is called a **reflection matrix** if

$$\mathbf{H} = \mathbf{I} - 2\mathbf{P},$$

where  $\mathbf{P}$  is an orthogonal projector (symmetric and idempotent)

**Interpretation:** Let  $\mathbf{P}^\perp = \mathbf{I} - \mathbf{P}$  be the orthogonal complement projector

$$\mathbf{x} = \mathbf{P}\mathbf{x} + \mathbf{P}^\perp\mathbf{x}, \quad \mathbf{H}\mathbf{x} = -\mathbf{P}\mathbf{x} + \mathbf{P}^\perp\mathbf{x}$$

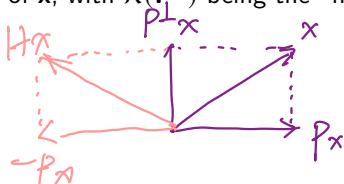
The vector  $\mathbf{H}\mathbf{x}$  is a reflected version of  $\mathbf{x}$ , with  $\mathcal{R}(\mathbf{P}^\perp)$  being the “mirror”

$$\begin{aligned} \mathbf{H}\mathbf{x} &= (\mathbf{I} - 2\mathbf{P})(\mathbf{P}\mathbf{x} + \mathbf{P}^\perp\mathbf{x}) \\ &= \mathbf{P}\mathbf{x} - 2\underbrace{\mathbf{P}\mathbf{P}\mathbf{x}}_{\substack{= \mathbf{P} \text{ idempotent} \\ = \mathbf{P}\mathbf{x}}} + \mathbf{P}^\perp\mathbf{x} - 2\underbrace{\mathbf{P}\mathbf{P}^\perp\mathbf{x}}_{=0} \\ &= -\mathbf{P}\mathbf{x} + \mathbf{P}^\perp\mathbf{x} \end{aligned}$$

A reflection matrix is orthogonal:

$$\mathbf{H}\mathbf{H}^T = \mathbf{H}^T\mathbf{H} = (\mathbf{I} - 2\mathbf{P})(\mathbf{I} - 2\mathbf{P}) = \mathbf{I} - 4\mathbf{P} + 4\mathbf{P}^2 = \mathbf{I} - 4\mathbf{P} + 4\mathbf{P} = \mathbf{I}$$

$\mathbf{H}$  symmetric



# Householder Reflections

**Problem:** Given  $\mathbf{x} \in \mathbb{R}^m$ , find an orthogonal  $\mathbf{H} \in \mathbb{R}^{m \times m}$  such that

$$\mathbf{H}\mathbf{x} = \begin{bmatrix} \beta \\ \mathbf{0} \end{bmatrix} = \beta \mathbf{e}_1, \quad \text{for some } \beta \in \mathbb{R}$$

**Householder reflection:** Let  $\mathbf{v} \in \mathbb{R}^m$ ,  $\mathbf{v} \neq \mathbf{0}$ , and let

$$\mathbf{H} = \mathbf{I} - \frac{2}{\|\mathbf{v}\|_2^2} \mathbf{v}\mathbf{v}^T,$$

which is a reflection matrix with  $\mathbf{P} = \mathbf{v}\mathbf{v}^T / \|\mathbf{v}\|_2^2$

*Handwritten notes:*  $\|\mathbf{v}\|_2^2$  (in red), symmetric and idempotent

$$\mathbf{P} \cdot \mathbf{P} = \frac{\mathbf{v}\mathbf{v}^T}{\|\mathbf{v}\|_2^2} \cdot \frac{\mathbf{v}\mathbf{v}^T}{\|\mathbf{v}\|_2^2} = \frac{\mathbf{v}(\underbrace{\mathbf{v}^T \mathbf{v}}_{\|\mathbf{v}\|_2^2})}{\|\mathbf{v}\|_2^4} = \frac{\mathbf{v}\mathbf{v}^T}{\|\mathbf{v}\|_2^2} = \mathbf{P}$$

## Householder Reflections (cont'd)

$$\mathbf{H}\mathbf{x} = \left( \mathbf{I} - \frac{2}{\|\mathbf{v}\|_2^2} \mathbf{v}\mathbf{v}^T \right) \mathbf{x} = \mathbf{x} - \frac{2\mathbf{v}^T \mathbf{x}}{\mathbf{v}^T \mathbf{v}} \mathbf{v}$$

Handwritten notes:  $\frac{2}{\|\mathbf{v}\|_2^2} \mathbf{v}(\mathbf{v}^T \mathbf{x})$  is a scalar. The term  $\mathbf{v}^T \mathbf{x}$  is the first element in  $\mathbf{x}$ .

We want  $\mathbf{H}\mathbf{x}$  to be a multiple of  $\mathbf{e}_1$ . Hence, we require  $\mathbf{v} \in \text{span}\{\mathbf{x}, \mathbf{e}_1\}$

Let  $\mathbf{v} = \mathbf{x} + \alpha \mathbf{e}_1$ . Then,

$$\mathbf{v}^T \mathbf{x} = \mathbf{x}^T \mathbf{x} + \alpha x_1, \quad \mathbf{v}^T \mathbf{v} = \mathbf{x}^T \mathbf{x} + 2\alpha x_1 + \alpha^2$$

It follows that

$$\mathbf{H}\mathbf{x} = \frac{\alpha^2 - \mathbf{x}^T \mathbf{x}}{\mathbf{x}^T \mathbf{x} + 2\alpha x_1 + \alpha^2} \mathbf{x} - 2\alpha \frac{\mathbf{v}^T \mathbf{x}}{\mathbf{v}^T \mathbf{v}} \mathbf{e}_1$$

The coefficient of  $\mathbf{x}$  has to be zero, so that  $\alpha^2 = \|\mathbf{x}\|_2^2$ . Therefore,

$$\mathbf{v} = \mathbf{x} \mp \|\mathbf{x}\|_2 \mathbf{e}_1 \implies \mathbf{H}\mathbf{x} = \pm \|\mathbf{x}\|_2 \mathbf{e}_1$$

The sign in the expression of  $\mathbf{v}$  may be determined to be the one that maximizes  $\|\mathbf{v}\|_2$  for the sake of numerical stability.

# Householder QR

1. Let  $\mathbf{H}_1 \in \mathbb{R}^{m \times m}$  be the Householder reflection w.r.t.  $\mathbf{a}_1$ . Transform  $\mathbf{A}$  as

$$\mathbf{A}^{(1)} = \mathbf{H}_1 \mathbf{A} = \begin{bmatrix} \times & \times & \dots & \times \\ 0 & \times & \dots & \times \\ \vdots & \vdots & & \vdots \\ 0 & \times & \dots & \times \end{bmatrix} \quad \mathcal{H}_1 \mathbf{a}_1 = \begin{bmatrix} * \\ 0 \\ \vdots \\ 0 \end{bmatrix}$$

2. Let  $\tilde{\mathbf{H}}_2 \in \mathbb{R}^{(m-1) \times (m-1)}$  be the Householder reflection w.r.t.  $\mathbf{A}^{(1)}(2:m, 2)$  (marked red above). Transform  $\mathbf{A}^{(1)}$  as

$$\mathbf{A}^{(2)} = \underbrace{\begin{bmatrix} 1 & \mathbf{0} \\ \mathbf{0} & \tilde{\mathbf{H}}_2 \end{bmatrix}}_{=\mathbf{H}_2} \mathbf{A}^{(1)} = \begin{bmatrix} \times & \times & \dots & \times \\ 0 & \tilde{\mathbf{H}}_2 \mathbf{A}^{(1)}(2:m, 2:n) \\ \vdots & \vdots & & \vdots \\ 0 & \vdots & & \vdots \end{bmatrix} = \begin{bmatrix} \times & \times & \times & \dots & \times \\ 0 & \times & \times & \dots & \times \\ \vdots & 0 & \times & \dots & \times \\ \vdots & \vdots & \vdots & & \vdots \\ 0 & 0 & \times & \dots & \times \end{bmatrix}$$

3. By repeating this,  $\mathbf{A}$  is transformed to  $\mathbf{R}$

## Householder QR (cont'd)

WLOG, assume  $m \geq n$

$$\mathbf{A}^{(0)} = \mathbf{A}$$

for  $k = 1, \dots, n-1$

$$\mathbf{A}^{(k)} = \mathbf{H}_k \mathbf{A}^{(k-1)}, \text{ where}$$

$$\mathbf{H}_k = \begin{bmatrix} \mathbf{I}_{k-1} & \mathbf{0} \\ \mathbf{0} & \tilde{\mathbf{H}}_k \end{bmatrix}$$

$\tilde{\mathbf{H}}_k$  is the Householder reflection of  $\mathbf{A}^{(k-1)}(k:m, k)$

end

$\tilde{\mathbf{H}}_k$  orthogonal matrix  
 $\mathbf{H}_k$  ↗

- The above procedure results in

$$\mathbf{A}^{(n-1)} = \mathbf{H}_{n-1} \cdots \mathbf{H}_2 \mathbf{H}_1 \mathbf{A}, \quad \mathbf{A}^{(n-1)} \text{ is upper triangular}$$

- QR decomposition is obtained by letting

$$\mathbf{R} = \mathbf{A}^{(n-1)} \text{ and } \mathbf{Q} = (\mathbf{H}_{n-1} \cdots \mathbf{H}_2 \mathbf{H}_1)^T$$

- A widely used method for QR decomposition

$$\begin{aligned} \mathbf{A} &= \mathbf{Q}\mathbf{R} \\ &= \mathbf{Q}(\mathbf{H}_{n-1} \cdots \mathbf{H}_1 \mathbf{A}) \\ \mathbf{Q} &= (\mathbf{H}_{n-1} \cdots \mathbf{H}_1)^T \end{aligned}$$

## Householder QR (cont'd)

$$\mathbf{A}^{(0)} = \mathbf{A}$$

for  $k = 1, \dots, n-1$

$$\mathbf{A}^{(k)} = \mathbf{H}_k \mathbf{A}^{(k-1)}, \text{ where}$$

$$\mathbf{H}_k = \begin{bmatrix} \mathbf{I}_{k-1} & \mathbf{0} \\ \mathbf{0} & \tilde{\mathbf{H}}_k \end{bmatrix}$$

$\tilde{\mathbf{H}}_k$  is the Householder reflection of  $\mathbf{A}^{(k-1)}(k:m, k)$   
end

Complexity (for  $m \geq n$ ):

- $O(n^2(m - n/3))$  for  $\mathbf{R}$  only
  - A direct implementation of the above pseudo-code does not lead to this complexity—Need to exploit the structures of  $\mathbf{H}_k$  in the implementations
- $O(m^2n - mn^2 + n^3/3)$  if  $\mathbf{Q}$  is also wanted
- See Section 5.2.2 of textbook

# Givens Rotations

**Example:** Let

$$\mathbf{J} = \begin{bmatrix} c & s \\ -s & c \end{bmatrix}$$

where  $c = \cos(\theta)$ ,  $s = \sin(\theta)$  for some  $\theta$

$$\mathbf{y} = \mathbf{J}\mathbf{x} \iff \begin{bmatrix} y_1 \\ y_2 \end{bmatrix} = \begin{bmatrix} c & s \\ -s & c \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} cx_1 + sx_2 \\ -sx_1 + cx_2 \end{bmatrix}$$

Observe that

*rotate clockwise by  $\theta$*

- $\mathbf{J}$  is orthogonal
- $y_2 = 0$  if  $\theta = \tan^{-1}(x_2/x_1)$ , or if

$$c = \frac{x_1}{\sqrt{x_1^2 + x_2^2}}, \quad s = \frac{x_2}{\sqrt{x_1^2 + x_2^2}}$$

## Givens Rotations (cont'd)

Givens rotations:

$$\mathbf{J}(i, k, \theta) = \begin{bmatrix} \mathbf{I} & & & \\ & \downarrow & & \\ & c & & s \\ & -s & & c \\ & & \mathbf{I} & \\ & & & \downarrow \\ & & & \mathbf{I} \end{bmatrix} \begin{matrix} \leftarrow i \\ \leftarrow k \end{matrix}$$

where  $c = \cos(\theta)$ ,  $s = \sin(\theta)$

- $\mathbf{J}(i, k, \theta)$  is orthogonal
- Let  $\mathbf{y} = \mathbf{J}(i, k, \theta)\mathbf{x}$ . Then,

$$y_j = \begin{cases} cx_i + sx_k, & j = i \\ -sx_i + cx_k, & j = k \\ x_j, & j \neq i, k \end{cases}$$

- $y_k$  is forced to zero if we choose  $\theta = \tan^{-1}(x_k/x_i)$



# Givens QR

**Example:** Consider a  $4 \times 3$  matrix.

$$\begin{aligned}
 \mathbf{A} = \begin{bmatrix} \times & \times & \times \\ \times & \times & \times \\ \times & \times & \times \\ \times & \times & \times \end{bmatrix} &\xrightarrow{\mathbf{J}_{1,2}} \begin{bmatrix} \times & \times & \times \\ 0 & \times & \times \\ \times & \times & \times \\ \times & \times & \times \end{bmatrix} \xrightarrow{\mathbf{J}_{1,3}} \begin{bmatrix} \times & \times & \times \\ 0 & \times & \times \\ 0 & \times & \times \\ \times & \times & \times \end{bmatrix} \xrightarrow{\mathbf{J}_{1,4}} \begin{bmatrix} \times & \times & \times \\ 0 & \times & \times \\ 0 & \times & \times \\ 0 & \times & \times \end{bmatrix} \\
 &\xrightarrow{\mathbf{J}_{2,3}} \begin{bmatrix} \times & \times & \times \\ 0 & \times & \times \\ 0 & 0 & \times \\ 0 & \times & \times \end{bmatrix} \xrightarrow{\mathbf{J}_{2,4}} \begin{bmatrix} \times & \times & \times \\ 0 & \times & \times \\ 0 & 0 & \times \\ 0 & 0 & \times \end{bmatrix} \xrightarrow{\mathbf{J}_{3,4}} \begin{bmatrix} \times & \times & \times \\ 0 & \times & \times \\ 0 & 0 & \times \\ 0 & 0 & 0 \end{bmatrix} = \mathbf{R}
 \end{aligned}$$

where  $\mathbf{B} \xrightarrow{\mathbf{J}} \mathbf{C}$  means  $\mathbf{C} = \mathbf{J}\mathbf{B}$ ;  $\mathbf{J}_{i,k} = \mathbf{J}(i, k, \theta)$ , with  $\theta$  chosen to zero out the  $(k, i)$ th entry of the matrix transformed by  $\mathbf{J}_{i,k}$

## Givens QR (cont'd)

**Givens QR** ( $m \geq n$ ): Perform a sequence of Givens rotations to annihilate the lower triangular parts of  $\mathbf{A}$

$$\underbrace{(\mathbf{J}_{n,m} \dots \mathbf{J}_{n,n+2} \mathbf{J}_{n,n+1}) \dots (\mathbf{J}_{2m} \dots \mathbf{J}_{24} \mathbf{J}_{23})(\mathbf{J}_{1m} \dots \mathbf{J}_{13} \mathbf{J}_{12})}_{=\mathbf{Q}^T} \mathbf{A} = \mathbf{R}$$

where  $\mathbf{R}$  is upper triangular and  $\mathbf{Q}$  is orthogonal

- Complexity (for  $m \geq n$ ):  $O(n^2(m - n/3))$  for  $\mathbf{R}$  only (see Section 5.2.5 of textbook)
- Not as efficient as Householder QR for general (and dense)  $\mathbf{A}$ 's
  - The flop count for Householder QR is  $2n^2(m - n/3)$
  - The flop count for Givens QR is  $3n^2(m - n/3)$
- Givens QR can be faster than Householder QR if  $\mathbf{A}$  has certain sparse structures and we exploit them

# Matrix Computations

## Chapter 3: Least-squares Problems and QR Decomposition

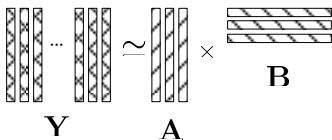
### Section 3.4 Problems Related to Least Squares

Jie Lu  
ShanghaiTech University

# Matrix Factorization

**Matrix Factorization:** Given  $\mathbf{Y} \in \mathbb{R}^{m \times n}$  and a positive integer  $k < \min\{m, n\}$ , solve

$$\min_{\mathbf{A} \in \mathbb{R}^{m \times k}, \mathbf{B} \in \mathbb{R}^{k \times n}} \|\mathbf{Y} - \mathbf{AB}\|_F^2$$



Also called **low-rank matrix approximation**

- $\text{rank}(\mathbf{AB}) \leq k$

# Principal Component Analysis

**Aim:** Given a collection of data points  $\mathbf{y}_1, \dots, \mathbf{y}_n \in \mathbb{R}^m$ , perform a low-dimensional representation

$$\mathbf{y}_i = \mathbf{A}\mathbf{b}_i + \mathbf{c} + \mathbf{v}_i, \quad i = 1, \dots, n,$$

where  $\mathbf{A} \in \mathbb{R}^{m \times k}$  is a basis matrix,  $\mathbf{b}_i \in \mathbb{R}^k$  is the coefficient for  $\mathbf{y}_i$ ,  $\mathbf{c} \in \mathbb{R}^m$  is the base or mean in statistics terms, and  $\mathbf{v}_i$  is noise or modeling error

Principal component analysis (PCA):

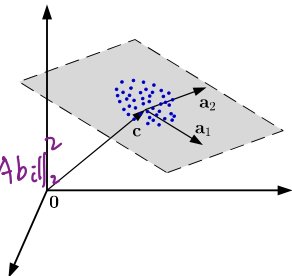
1. Choose  $\mathbf{c} = \frac{1}{n} \sum_{i=1}^n \mathbf{y}_i$
2. Let  $\bar{\mathbf{y}}_i = \mathbf{y}_i - \mathbf{c}$ , and solve

$$\bar{\mathbf{y}}_i = \mathbf{A}\mathbf{b}_i + \mathbf{v}_i$$

$$\downarrow$$
$$\min \|\bar{\mathbf{y}}_i - \mathbf{A}\mathbf{b}_i\|_2^2$$

$$\bar{\mathbf{Y}} = [\bar{\mathbf{y}}_1 \dots \bar{\mathbf{y}}_n] \quad \min_{\mathbf{A}, \mathbf{B}} \|\bar{\mathbf{Y}} - \mathbf{A}\mathbf{B}\|_F^2$$
$$\mathbf{B} = [\mathbf{b}_1 \dots \mathbf{b}_n]$$

3. we may want a semi-orthogonal  $\mathbf{A}$



**Applications:** dimensionality reduction, visualization of high-dimensional data, compression, extraction of meaningful features from data, etc.

- Example of senate voting: <http://livebooklabs.com/keepies/c5a5868ce26b8125>

# Topic Modeling

**Aim:** Discover thematic information or topics from a large collection of documents (e.g., books, articles, news, blogs)

**Bag-of-words representation:** Represent each document as a vector of word counts

... In fact, we will soon see that the **implementation** of **SDR** can be very easy, which allows **signal processing** practitioners to quickly test the viability of **SDR** in their applications. Several highly successful **applications** will be showcased as **examples** .....

a document



bag of words



$y =$

count	term
0	efficiency
2	applications
2	SDR
0	communications
1	example
1	signal processing
⋮	⋮
1	implementation

bag-of-words representation

## Topic Modeling (cont'd)

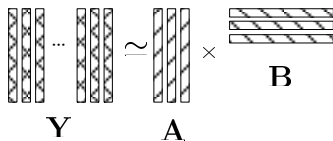
- Let  $n$  be the number of documents
- Let  $\mathbf{y}_i \in \mathbb{R}^m$  be the bag-of-words representation of the  $i$ th document
- $\mathbf{Y} = [\mathbf{y}_1, \dots, \mathbf{y}_n] \in \mathbb{R}^{m \times n}$  is called the term-document matrix
- Hypotheses:<sup>1</sup>
  - If documents have similar columns vectors in  $\mathbf{Y}$  or similar usage of words, they tend to have similar meanings
  - The topic of a document will probabilistically influence the author's choice of words when writing the document

---

<sup>1</sup>P. D. Turney and P. Pantel, "From frequency to meaning: Vector space models of semantics," *Journal of*

## Topic Modeling (cont'd)

**Problem:** Apply matrix factorization to a term-document matrix  $\mathbf{Y}$



The diagram shows the matrix factorization equation  $\mathbf{Y} \approx \mathbf{A} \mathbf{B}$ . Matrix  $\mathbf{Y}$  is represented by two groups of three vertical rectangles, each containing a grid of 'x' marks, with an ellipsis between the groups. Matrix  $\mathbf{A}$  is represented by a single group of three vertical rectangles, each containing diagonal lines. Matrix  $\mathbf{B}$  is represented by a single group of three horizontal rectangles, each containing diagonal lines. The matrices are connected by an approximation symbol ( $\approx$ ) and a multiplication symbol ( $\times$ ).

$\mathbf{A}$  is called a term-topic matrix and  $\mathbf{B}$  is called a topic-document matrix

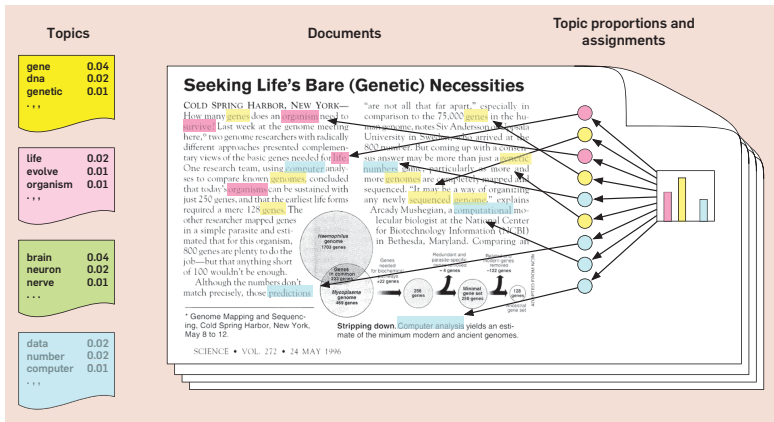
**Interpretation:**

- Each column  $\mathbf{a}_i$  of  $\mathbf{A}$  represents a theme topic (e.g., local affairs, foreign affairs, politics, sports)
- $\mathbf{y}_i \approx \mathbf{A} \mathbf{b}_i$ : each document is postulated as a linear combination of topics
- Matrix factorization aims at discovering topics from the documents

Topic modeling via matrix factorization has been used in or is tightly connected to information retrieval, natural language processing, machine learning; document clustering, classification and retrieval; latent semantic analysis, latent semantic indexing: finding similarities of documents, similarities of terms, etc.

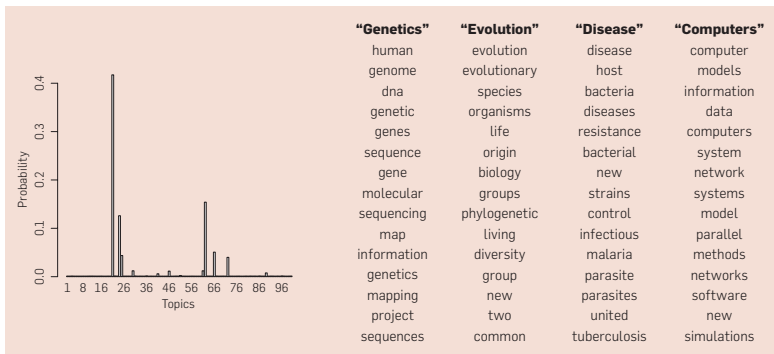


# Topic Modeling (cont'd)



Source: D. Blei, “Probabilistic topic models,” *Communications of the ACM*, vol. 55, no. 4, pp. 77–84, 2012.

# Topic Modeling (cont'd)



Topics found in a real set of documents. The document set consists of 17,000 articles from the journal *Science*. The topics are discovered using a technique called *latent Dirichlet allocation*, which is not the same as, but has strong connections to, matrix factorization [Blei'12]

# Matrix Factorization

**Problem:**

$$\min_{\mathbf{A} \in \mathbb{R}^{m \times k}, \mathbf{B} \in \mathbb{R}^{k \times n}} \|\mathbf{Y} - \mathbf{AB}\|_F^2$$

The problem has non-unique solutions

- If  $(\mathbf{A}^*, \mathbf{B}^*)$  is an optimal solution to the problem, then  $(\mathbf{A}^* \mathbf{Q}^{-1}, \mathbf{Q} \mathbf{B}^*)$  is also an optimal solution for any nonsingular  $\mathbf{Q} \in \mathbb{R}^{k \times k}$
- The non-uniqueness of solution makes it a bad formulation for problems such as topic modeling

The problem is non-convex, but can be solved by singular value decomposition (beautifully)

It can also be solved by LS approach

# Alternating LS for Matrix Factorization

Alternating LS (ALS): Given a starting point  $(\mathbf{A}^{(0)}, \mathbf{B}^{(0)})$ , do

$$\mathbf{A}^{(i+1)} = \arg \min_{\mathbf{A} \in \mathbb{R}^{m \times k}} \|\mathbf{Y} - \mathbf{A}\mathbf{B}^{(i)}\|_F^2 \quad (1)$$

$$\mathbf{B}^{(i+1)} = \arg \min_{\mathbf{B} \in \mathbb{R}^{k \times n}} \|\mathbf{Y} - \mathbf{A}^{(i+1)}\mathbf{B}\|_F^2 \quad (2)$$

for  $i = 0, 1, 2, \dots$ , and stop when a termination criterion is satisfied

Make a mild assumption that  $\mathbf{A}^{(i)}, \mathbf{B}^{(i)}$  have full rank at every  $i$

Look at (2) first.

$$\min_{\mathbf{B} \in \mathbb{R}^{k \times n}} \|\mathbf{Y} - \mathbf{A}^{(i+1)}\mathbf{B}\|_F^2 = \min_{b_1, \dots, b_n \in \mathbb{R}^k} \sum_{j=1}^n \|\mathbf{y}_j - \mathbf{A}^{(i+1)}\mathbf{b}_j\|_2^2$$

$$\Leftrightarrow \min_{\mathbf{b}_j \in \mathbb{R}^k} \|\mathbf{y}_j - \mathbf{A}^{(i+1)}\mathbf{b}_j\|_2^2, \quad \forall j = 1, \dots, n$$

$$\mathbf{b}_j^{(i+1)} = \left[ (\mathbf{A}^{(i+1)})^T \mathbf{A}^{(i+1)} \right]^{-1} (\mathbf{A}^{(i+1)})^T \mathbf{y}_j$$

$$\mathbf{B}^{(i+1)} = \left[ (\mathbf{A}^{(i+1)})^T \mathbf{A}^{(i+1)} \right]^{-1} (\mathbf{A}^{(i+1)})^T \mathbf{Y}$$