## Cooperative Games in Social Network

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#### Outline

- The Shapley Value
- Invitation Incentive Mechanisms

#### Coalitional/Cooperative Game

- A set of agents N.
- Each subset of agents(coalition)  $S \subseteq N$  cooperate together can generate some value  $v(S) \in R$ . Assume  $v(\emptyset) = 0$ . N is called grand coalition.  $v : 2^N \to R$  is called the characteristic function of the game. v is often assumed to be monotonic:  $S \subseteq T \to v(S) \le v(T)$ .
- The possible outcomes of the game is defined by  $V(S) = \{x \in R^S : \sum_{i \in S} x_i \le v(S)\}.$

### Example

- Three agents {1,2,3}.
- $v(\{1\}) = v(\{2\}) = 10, v(\{3\}) = 1;$   $v(\{1,2\}) = 20, v(\{1,3\}) = v(\{2,3\}) = 12;$  $v(\{1,2,3\}) = 22.$

# Shapley Value: a Fair Distribution of Payoffs

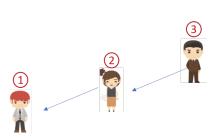
Given a coalitional game (N, v), the Shapley value of each player i is:

$$\phi_i(v) = \sum_{S \subset N \setminus \{i\}} \frac{|S|!(n-|S|-1)!}{n!} (v(S \cup \{i\}) - v(S))$$

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## Shapley Value: a Fair Distribution of Payoffs

 $v(S \cup \{i\}) - v(S)$ :marginal contribution for agent i, which is denoted as  $c_i$ 



- Initially, the coalition is empty and agent 1 comes.  $c_1 = v(\{1\}) - v(\emptyset) = 10$
- Agent 2 comes.  $c_2 = v(\{1,2\}) v(\{1\}) = 10$
- Agent 3 comes.  $c_3 = v(\{1,2,3\}) v(\{1,2\}) = 2$

# Shapley Value: a Fair Distribution of Payoffs

 List all the permutations and get respective marginal contributions.

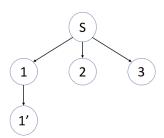
Permutation	<i>C</i> <sub>1</sub>	<i>c</i> <sub>2</sub>	<i>c</i> <sub>3</sub>
1,2,3	10	10	2
1,3,2	10	10	2
2,3,1	10	10	2
2,1,3	10	10	2
3,1,2	11	10	1
3,2,1	10	11	1

Get the final payoff:  $\phi_i = \frac{\sum c_i}{6}$ 

• 
$$\phi_1 = \frac{61}{6}, \phi_2 = \frac{61}{6}, \phi_3 = \frac{10}{6}$$

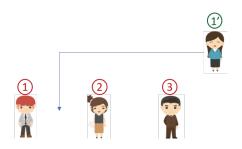
- We want to incentivize participants to invite more participants.
- What if agent 1 invites agent 1' who can complete the same task?

(i.e. 
$$v(\{1\} \cup S) = v(\{1'\} \cup S) = v(\{1, 1'\} \cup S)$$
)

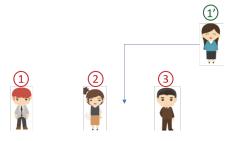




- Without agent 1': $c_1 = 10$ ,  $c_2 = 10$ ,  $c_3 = 2$
- With agent 1': $c'_{1'} = 10, c'_{1} = 0, c'_{2} = 10, c'_{3} = 2$



- Without agent 1': $c_1 = 10$ ,  $c_2 = 10$ ,  $c_3 = 2$
- With agent 1': $c_1' = 10, c_1' = 0, c_2' = 10, c_3' = 2$



- Without agent 1': $c_1 = 10$ ,  $c_2 = 10$ ,  $c_3 = 2$
- With agent 1': $c_1' = 10, c_2' = 10, c_{11}' = 0, c_3' = 2$









- Without agent 1': $c_1 = 10$ ,  $c_2 = 10$ ,  $c_3 = 2$
- With agent 1': $\mathbf{c}'_1 = 10$ ,  $\mathbf{c}'_2 = 10$ ,  $\mathbf{c}'_3 = 2$ ,  $\mathbf{c}'_{1'} = 0$

 Similarly, we can get all the permutations and marginal contributions.

Permutation	<i>C</i> <sub>1</sub>	<b>c</b> <sub>2</sub>	<i>c</i> <sub>3</sub>	C <sub>1′</sub>
1,2,3,1'	10 → 10	10	2	0
1,3,2,1'	10 → 10	10	2	0
3,1′,1,2	11 → 0	10	1	11
1′,3,2,1	10 → 0	10	2	10

• For the permutation where agent 1' is before  $1, c_1 = 0$  since agent 1' performs the same as agent 1.

#### Outline

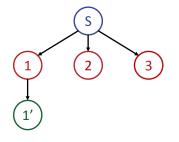
- The Shapley Value
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### A Layer-based Solution

#### Main idea

- Divide the agents into layers in terms of distance
- Lower layers always join the game earlier than higher layers

## Layered Shapley Value



- Recalling the previous example, agents 1,2,3 are in the first layer while agent 1' is in the second layer.
- We only retain the permutation where 1' is after the set {1,2,3} (i.e 1,3,2,1') and remove others

- (1)
- (2)
- 3
- 1'



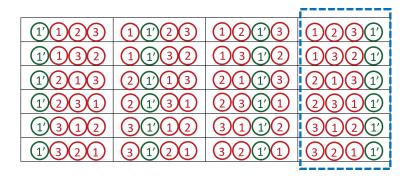
- $\bigcirc$ 1
- 1'

(3)

 $\left(2\right)$ 



### Layered Shapley Value



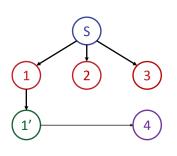
 The permutations in blue dotted box are retained, others are removed.

## Layered Shapley Value

 Similarly, we can get marginal contributions for the remaining permutation.

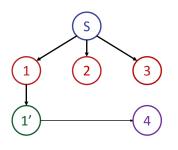
Permutation	<i>C</i> <sub>1</sub>	<b>c</b> <sub>2</sub>	<b>c</b> <sub>3</sub>	C <sub>1′</sub>
1,2,3,1'	10	10	2	0
1,3,2,1'	10	10	2	0
2,1,3,1'	10	10	2	0
2,3,1,1'	10	2	10	0
3,1,2,1'	11	10	1	0
3,2,1,1'	10	11	1	0

## A General Solution: Permission Shapley Value



- In the previous example, suppose agent 1' invites agent 4.
- $v(\{1\}) = v(\{2\}) = 10.$   $v(\{3\}) = 1;$   $v(\{1,2\}) = 20;$   $v(\{1,3\}) = v(\{2,3\}) = 12;$  $v(\{1,2,3\}) = 22.$
- $v(\{1'\} \cup S) = v(\{1\} \cup S) = v(\{1, 1'\} \cup S)$
- $v(\{4\} \cup S) = v(S) + 100$

## A General Solution: Permission Shapley Value

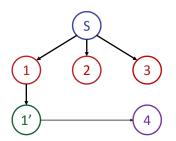


#### Permission structure:

- A permission structure on N is an asymmetric mapping  $p:N \to 2^N$ , i.e.,  $j \in p(i)$  implies that  $i \notin p(j)$
- p(i): the set of players who invited i into the coalition.
- In the graph,  $p(1') = \{1\}, p(4) = \{1'\}.$

#### **Autonomous**

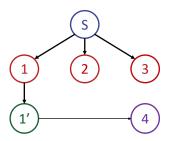
A coalition  $S \subseteq N$  is autonomous in a permission structure p if for all  $i \in S$ ,  $p(i) \subseteq S$ .



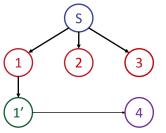
• In this graph,  $S = \{1, 2, 1'\}$  is autonomous, but  $S = \{1, 2, 4\}$  is not autonomous since  $p(4) = \{1'\} \not\subseteq S$ .

#### Largest Autonomous Part

Let p be a permission structure on N. Then the largest automonous part of a coalition  $S \subset N$  is defined by  $\alpha(S) = \bigcup \{T | T \subseteq S \text{ and } T \in A_p\}$  where  $A_p$  denotes the collection of all autonomous coalitions in p.



• In the graph,  $\alpha(\{2,3,1',4\}) = \{2,3\}$ 



Permission Shapley value  $v^p(S)$ .

•  $v^p(S) = v(\alpha(S))$ 

• In the graph, e.g.  $v^p(\{1,2,3,4\}) = v(\{1,2,3\}) = 22$  since  $\alpha(\{1,2,3,4\}) = \{1,2,3\}.$ 

• Use permutation{1,2,4,3,1'} as an example.

Set S	$\alpha(S)$	v(S)	$v^p(S)$
1	1	10	10
1,2	1,2	20	20
1,2,4	1,2	120	20
1,2,4,3	1,2,3	122	22
1,2,4,3,1′	1,2,4,3,1′	122	122

• 
$$\phi_1 = \frac{61}{6} + \frac{100}{3}, \phi_2 = \frac{61}{6}, \phi_3 = \frac{10}{6}$$

• 
$$\phi_{1'} = \frac{100}{3}, \phi_4 = \frac{100}{3}$$