

*From Matrix to Tensor:
The Transition to Numerical Multilinear Algebra*

**Lecture 5. The CP Representation and
Tensor Rank**

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Where We Are

- Lecture 1. Introduction to Tensor Computations
- Lecture 2. Tensor Unfoldings
- Lecture 3. Transpositions, Kronecker Products, Contractions
- Lecture 4. Tensor-Related Singular Value Decompositions
- Lecture 5. The CP Representation and Tensor Rank**
- Lecture 6. The Tucker Representation
- Lecture 7. Other Decompositions and Nearness Problems
- Lecture 8. Multilinear Rayleigh Quotients
- Lecture 9. The Curse of Dimensionality
- Lecture 10. Special Topics

What is this Lecture About?

Sums of Rank-1 Tensors

The SVD of a matrix A expresses A as a very special sum of rank-1 matrices.

Let us do the same thing as much as possible with tensor \mathcal{A} .

This requires an understanding of (a) rank-1 tensors and their unfoldings, (b) the Kruskal tensor format, (c) the alternating least squares framework for multilinear sum-of-squares optimization, and (d) the notion of tensor rank.

We use the order-3 case to motivate the main ideas.

What is this Lecture About?

A Note on Terminology

The central decomposition in this lecture is the **CP** Decomposition.

It also goes by the name of the **CANDECOMP**/**PARAFAC** Decomposition.

CANDECOMP = Canonical Decomposition

PARAFAC = Parallel Factors Decomposition

Rank-1 Tensors (Order-3)

Definition

If $f \in \mathbb{R}^{n_1}$, $g \in \mathbb{R}^{n_2}$, and $h \in \mathbb{R}^{n_3}$, then

$$\mathcal{B} = f \circ g \circ h$$

is defined by

$$\mathcal{B}(i_1, i_2, i_3) = f(i_1)g(i_2)h(i_3).$$

The tensor $\mathcal{B} \in \mathbb{R}^{n_1 \times n_2 \times n_3}$ is a **rank-1 tensor**.

Rank-1 Tensors (Order-3)

The Kronecker Product Connection...

$$\mathcal{B} = \begin{bmatrix} f_1 \\ f_2 \end{bmatrix} \circ \begin{bmatrix} g_1 \\ g_2 \\ g_3 \end{bmatrix} \circ \begin{bmatrix} h_1 \\ h_2 \end{bmatrix} \Leftrightarrow \begin{bmatrix} b_{111} \\ b_{211} \\ b_{121} \\ b_{221} \\ b_{131} \\ b_{231} \\ b_{112} \\ b_{212} \\ b_{122} \\ b_{222} \\ b_{132} \\ b_{232} \end{bmatrix} = \begin{bmatrix} f_1 g_1 h_1 \\ f_2 g_1 h_1 \\ f_1 g_2 h_1 \\ f_2 g_2 h_1 \\ f_1 g_3 h_1 \\ f_2 g_3 h_1 \\ f_1 g_1 h_2 \\ f_2 g_1 h_2 \\ f_1 g_2 h_2 \\ f_2 g_2 h_2 \\ f_1 g_3 h_2 \\ f_2 g_3 h_2 \end{bmatrix} = h \otimes g \otimes f$$

Rank-1 Tensors (Order-3)

The Modal Unfoldings...

If

$$\mathcal{B} = f \circ g \circ h = \begin{bmatrix} f_1 \\ f_2 \end{bmatrix} \circ \begin{bmatrix} g_1 \\ g_2 \\ g_3 \end{bmatrix} \circ \begin{bmatrix} h_1 \\ h_2 \end{bmatrix}$$

then

$$\begin{aligned} \mathcal{B}_{(1)} &= \begin{bmatrix} f_1 g_1 h_1 & f_1 g_2 h_1 & f_1 g_3 h_1 & f_1 g_1 h_2 & f_1 g_2 h_2 & f_1 g_3 h_2 \\ f_2 g_1 h_1 & f_2 g_2 h_1 & f_2 g_3 h_1 & f_2 g_1 h_2 & f_2 g_2 h_2 & f_2 g_3 h_2 \end{bmatrix} \\ &= \begin{bmatrix} f_1 \cdot (h \otimes g)^T \\ f_2 \cdot (h \otimes g)^T \end{bmatrix} \\ &= f \otimes (h \otimes g)^T \end{aligned}$$

Rank-1 Tensors (Order-3)

The Modal Unfoldings...

If

$$\mathcal{B} = f \circ g \circ h = \begin{bmatrix} f_1 \\ f_2 \end{bmatrix} \circ \begin{bmatrix} g_1 \\ g_2 \\ g_3 \end{bmatrix} \circ \begin{bmatrix} h_1 \\ h_2 \end{bmatrix}$$

then

$$\begin{aligned} \mathcal{B}_{(2)} &= \begin{bmatrix} f_1 g_1 h_1 & f_2 g_1 h_1 & f_1 g_1 h_2 & f_2 g_1 h_2 \\ f_1 g_2 h_1 & f_2 g_2 h_1 & f_1 g_2 h_2 & f_2 g_2 h_2 \\ f_1 g_3 h_1 & f_2 g_3 h_1 & f_1 g_3 h_2 & f_2 g_3 h_2 \end{bmatrix} \\ &= \begin{bmatrix} g_1 \cdot (h \otimes f)^T \\ g_2 \cdot (h \otimes f)^T \\ g_3 \cdot (h \otimes f)^T \end{bmatrix} \\ &= g \otimes (h \otimes f)^T \end{aligned}$$

Rank-1 Tensors (Order-3)

The Modal Unfoldings...

If

$$\mathcal{B} = f \circ g \circ h = \begin{bmatrix} f_1 \\ f_2 \end{bmatrix} \circ \begin{bmatrix} g_1 \\ g_2 \\ g_3 \end{bmatrix} \circ \begin{bmatrix} h_1 \\ h_2 \end{bmatrix}$$

then

$$\begin{aligned} \mathcal{B}_{(3)} &= \begin{bmatrix} f_1 g_1 h_1 & f_2 g_1 h_1 & f_1 g_2 h_1 & f_2 g_2 h_1 & f_1 g_3 h_1 & f_2 g_3 h_1 \\ f_1 g_1 h_2 & f_2 g_1 h_2 & f_1 g_2 h_2 & f_2 g_2 h_2 & f_1 g_3 h_2 & f_2 g_3 h_2 \end{bmatrix} \\ &= \begin{bmatrix} h_1 \cdot (g \otimes f)^T \\ h_2 \cdot (g \otimes f)^T \end{bmatrix} \\ &= h \otimes (g \otimes f)^T \end{aligned}$$

Problem 5.1. Suppose $a \in \mathbb{R}^{n_1 n_2 n_3}$. Show how to compute $f \in \mathbb{R}^{n_1}$ and $g \in \mathbb{R}^{n_2}$ so that $\|a - h \otimes g \otimes f\|_2$ is minimized where $h \in \mathbb{R}^{n_3}$ is given. Hint. It's an SVD problem.

Problem 5.2. Given $\mathcal{A} \in \mathbb{R}^{n_1 \times n_2 \times n_3}$ with positive entries, how would you choose $\mathcal{B} = f \circ g \circ h \in \mathbb{R}^{n_1 \times n_2 \times n_3}$ so that

$$\phi(f, g, h) = \sum_{i=1}^n |\log(\mathcal{A}(i)) - \log(\mathcal{B}(i))|^2$$

is minimized.

The CP Representation (Order-3)

Notation

Given $\lambda \in \mathbb{R}^r$, $F \in \mathbb{R}^{n_1 \times r}$, $G \in \mathbb{R}^{n_2 \times r}$, and $H \in \mathbb{R}^{n_3 \times r}$, we define $[[\lambda; F, G, H]] \in \mathbb{R}^{n_1 \times n_2 \times n_3}$ by

$$[[\lambda, F, G, H]] = \sum_{j=1}^r \lambda_j \cdot F(:,j) \circ G(:,j) \circ H(:,j)$$

A weighted sum of rank-1 tensors where the vectors that specify the rank-1's are columns of the matrices F , G , and H .

The CP Representation (Order-3)

Kruskal Form

We say that $\mathcal{X} \in \mathbb{R}^{n_1 \times n_2 \times n_3}$ is in **Kruskal form** if

$$\mathcal{X} = [[\lambda; F, G, H]].$$

where $\lambda \in \mathbb{R}^r$, $F \in \mathbb{R}^{n_1 \times r}$, $G \in \mathbb{R}^{n_2 \times r}$, and $H \in \mathbb{R}^{n_3 \times r}$.

*Can we write a given tensor \mathcal{A} as an illuminating sum of rank-1 tensors?
I.e., Given \mathcal{A} , can we find $\mathcal{X} = [[\lambda; F, G, H]]$ so that $\mathcal{A} \approx \mathcal{X}$ in some
meaningful way?*

The CP Representation (Order-3)

Equivalent Formulations...

If $\mathcal{X} = [[\lambda, F, G, H]] \in \mathbb{R}^{n_1 \times n_2 \times n_3}$, then

$$\mathcal{X}(i_1, i_2, i_3) = \sum_{j=1}^r \lambda_j \cdot F(i_1, j) \cdot G(i_2, j) \cdot H(i_3, j)$$

$$\text{vec}(\mathcal{X}) = \sum_{j=1}^r \lambda_j \cdot H(:, j) \otimes G(:, j) \otimes F(:, j)$$

MATLAB Tensor Toolbox: **ktensor** Set-Up

```
n = [5 8 3]; r = 4;  
% Set up a random, length-r ktensor...  
F = randn(n(1),r); G = randn(n(2),r);  
H = randn(n(3),r); lambda = ones(r,1);  
X = ktensor(lambda,{F,G,H});  
Fsize = size(X.U{1}); Gsize = size(X.U{2});  
Hsize = size(X.U{3});  
L = length(X.lambda); s = size(X);
```

A ktensor is a structure with two fields that is used to represent a tensor in Kruskal form. In the above, `X.lambda` houses the vector of weights while `X.U` is a cell array of the matrices that define the tensor `X`.

Variable	Value
Fsize	[5,4]
Gsize	[8,4]
Hsize	[3,4]
L	4
s	[5 8 3]

MATLAB Tensor Toolbox: Norm of a ktensor

```
function alfa = normKruskal(X)
% X is a ktensor and alfa is the Frobenius norm
% of the tensor it represents.
N = prod(size(X));
% Create a multidimensional array that houses
% the Kruskal tensor...
Xarray = double(X);
% Reshape as a vector and compute its 2-norm...
alfa = norm(reshape(Xarray,N,1));
```

Problem 5.3. Write a MATLAB function $Y = \text{normalize}(X)$ that takes a takes a ktensor X and returns a ktensor Y with the property that (a)

$$Y.U\{j\}(:,k) = X.U\{j\}(:,k)/\text{norm}(X.U\{j\}(:,k))$$

for all appropriate values of k and j and (b) $\text{double}(X) = \text{double}(Y)$.

The CP Representation (Order-3)

The CP Approximation Problem

Given $\mathcal{A} \in \mathbb{R}^{n_1 \times n_2 \times n_3}$ and r , determine $\lambda \in \mathbb{R}^r$, $F \in \mathbb{R}^{n_1 \times r}$, $G \in \mathbb{R}^{n_2 \times r}$, and $H \in \mathbb{R}^{n_3 \times r}$ so that

$$\mathcal{A} \approx [[\lambda; F, G, H]] = \mathcal{X}.$$

Using Least Squares...

Choose λ , F , G , and H so that

$$\|\mathcal{A} - \mathcal{X}\|_F^2 = \left\| \text{vec}(\mathcal{A}) - \sum_{j=1}^r \lambda_j \cdot H(:,j) \otimes G(:,j) \otimes F(:,j) \right\|_2^2$$

is minimized.

A multilinear optimization problem. Reshape using the Khatri-Rao Product...

The Khatri-Rao Product

Definition

If

$$B = [b_1 \mid \cdots \mid b_r] \in \mathbb{R}^{n_1 \times r}$$

$$C = [c_1 \mid \cdots \mid c_r] \in \mathbb{R}^{n_2 \times r}$$

then the **Khatri-Rao product** of B and C is given by

$$B \odot C = [b_1 \otimes c_1 \mid \cdots \mid b_r \otimes c_r].$$

Note that $B \odot C \in \mathbb{R}^{n_1 n_2 \times r}$.

The Khatri-Rao Product

“Fast” Property 1.

If $B \in \mathbb{R}^{n_1 \times r}$ and $C \in \mathbb{R}^{n_2 \times r}$, then

$$(B \odot C)^T (B \odot C) = (B^T B) .* (C^T C)$$

where “ $.*$ ” denotes pointwise multiplication.

Problem 5.4. Prove this property using the Kronecker product facts (i) $(W \otimes X)(Y \otimes Z) = WY \otimes XZ$ and (ii) $(W \otimes X)^T = W^T \otimes X^T$. How many flops are required?

The Khatri-Rao Product

“Fast” Property 2.

If

$$B = [b_1 \mid \cdots \mid b_r] \in \mathbb{R}^{n_1 \times r}$$

$$C = [c_1 \mid \cdots \mid c_r] \in \mathbb{R}^{n_2 \times r}$$

$x \in \mathbb{R}^{n_1 n_2}$, and $y = (B \odot C)^T x$, then

$$y = \begin{bmatrix} c_1^T X b_1 \\ \vdots \\ c_r^T X b_r \end{bmatrix} \quad X = \text{reshape}(x, n_2, n_1)$$

Problem 5.5. Prove this property using $\text{vec}(YXW^T) = (W \otimes Y) \cdot \text{vec}(X)$. How many flops are required?

Problem 5.6. Complete the following function so that it performs as specified.

```
function x = KRLS(B,C,d)
% B is n1-by-1, C is n2-by-1, and d is n1*n2-by-1
% x minimizes norm(Ax - d,2) where A is the Khatri-Rao
% product of B and C.
```

Use the method of normal equations and assume that A has full column rank. Is there an equally efficient way to solve the problem via the QR factorization of A.

The CP Representation (Order-3)

Unfolding Tensors in the Kruskal Form

Given $\mathcal{A} \in \mathbb{R}^{n_1 \times n_2 \times n_3}$ our goal is to minimize

$$\| \mathcal{A} - \mathcal{X} \|_F = \| \mathcal{A}_{(k)} - \mathcal{X}_{(k)} \|_F$$

where

$$\mathcal{X} = [[\lambda; F, G, H]] = \sum_{j=1}^r \lambda_j \cdot f_j \circ g_j \circ h_j$$

with

$$F = [f_1 \mid \cdots \mid f_r] \quad G = [g_1 \mid \cdots \mid g_r] \quad H = [h_1 \mid \cdots \mid h_r]$$

So what do the modal unfoldings of \mathcal{X} look like?

It will be a sum of rank-1 tensor unfoldings...

The CP Representation (Order-3)

Unfolding Tensors in the Kruskal Form

Since $\mathcal{B} = f \circ g \circ h$ implies

$$\mathcal{B}_{(1)} = f \otimes (h \otimes g)^T$$

$$\mathcal{B}_{(2)} = g \otimes (h \otimes f)^T$$

$$\mathcal{B}_{(3)} = h \otimes (g \otimes f)^T$$

we have

$$\mathcal{X}_{(1)} = \sum_{j=1}^r \lambda_j \cdot f_j \otimes (h_j \otimes g_j)^T = F \cdot \text{diag}(\lambda_j) \cdot (H \odot G)^T$$

$$\mathcal{X}_{(2)} = \sum_{j=1}^r \lambda_j \cdot g_j \otimes (h_j \otimes f_j)^T = G \cdot \text{diag}(\lambda_j) \cdot (H \odot F)^T$$

$$\mathcal{X}_{(3)} = \sum_{j=1}^r \lambda_j \cdot h_j \otimes (g_k \otimes f_j)^T = H \cdot \text{diag}(\lambda_j) \cdot (G \odot F)^T$$

The CP Representation (Order-3)

The Alternating LS Solution Framework...

$$\| \mathcal{A} - \mathcal{X} \|_F$$

=

$$\| \mathcal{A}_{(1)} - F \cdot \text{diag}(\lambda_j) \cdot (H \odot G)^T \|_F$$

\Leftarrow

1. Fix G and H and improve λ and F .

=

$$\| \mathcal{A}_{(2)} - G \cdot \text{diag}(\lambda_j) \cdot (H \odot F)^T \|_F$$

\Leftarrow

2. Fix F and H and improve λ and G .

=

$$\| \mathcal{A}_{(3)} - H \cdot \text{diag}(\lambda_j) \cdot (G \odot F)^T \|_F$$

\Leftarrow

3. Fix F and G and improve λ and H .

The CP Representation (Order-3)

The Alternating LS Solution Framework

Repeat:

1. Let \tilde{F} minimize $\| \mathcal{A}_{(1)} - \tilde{F} \cdot (H \odot G)^T \|_F$ and for $j = 1:r$ set

$$\lambda_j = \| \tilde{F}(:,j) \|_2 \quad \text{and} \quad F(:,j) = \tilde{F}(:,j)/\lambda_j.$$

2. Let \tilde{G} minimize $\| \mathcal{A}_{(2)} - \tilde{G} \cdot (H \odot F)^T \|_F$ and for $j = 1:r$ set

$$\lambda_j = \| \tilde{G}(:,j) \|_2 \quad \text{and} \quad G(:,j) = \tilde{G}(:,j)/\lambda_j.$$

3. Let \tilde{H} minimize $\| \mathcal{A}_{(3)} - \tilde{H} \cdot (G \odot F)^T \|_F$ and for $j = 1:r$ set

$$\lambda_j = \| \tilde{H}(:,j) \|_2 \quad \text{and} \quad H(:,j) = \tilde{H}(:,j)/\lambda_j.$$

These are linear least squares problems. The columns of F , G , and H are normalized.

The CP Representation (Order-3)

Solving the LS Problems

The solution to

$$\min_{\tilde{F}} \| \mathcal{A}_{(1)} - \tilde{F} \cdot (H \odot G)^T \|_F = \min_{\tilde{F}} \| \mathcal{A}_{(1)}^T - (H \odot G) \tilde{F}^T \|_F$$

can be obtained by solving the normal equation system

$$(H \odot G)^T (H \odot G) \tilde{F}^T = (H \odot G)^T \mathcal{A}_{(1)}^T$$

Can be solved efficiently by exploiting the ideas in Problem 5.6.

Problem 5.7. Write a MATLAB function $X = \text{MyKruskal}(A, r, \text{itMax})$ that takes an order-3 tensor A and returns a ktensor X with the property that $\mathcal{A} \approx \mathcal{X}$. $\mathcal{X} = [[\lambda; F, G, H]]$ should be obtained by applying the following improvement steps itMax times:

1. Solve $(H \odot G)^T (H \odot G) \tilde{F}^T = (H \odot G)^T \mathcal{A}_{(1)}^T$ and for $j = 1:r$ set

$$\lambda_j = \|\tilde{F}(:,j)\|_2 \quad \text{and} \quad F(:,j) = \tilde{F}(:,j)/\lambda_j.$$

2. Solve $(H \odot F)^T (H \odot F) \tilde{G}^T = (H \odot F)^T \mathcal{A}_{(2)}^T$ and for $j = 1:r$ set

$$\lambda_j = \|\tilde{G}(:,j)\|_2 \quad \text{and} \quad G(:,j) = \tilde{G}(:,j)/\lambda_j.$$

3. Solve $(G \odot F)^T (G \odot F) \tilde{H}^T = (G \odot F)^T \mathcal{A}_{(3)}^T$ and for $j = 1:r$ set

$$\lambda_j = \|\tilde{H}(:,j)\|_2 \quad \text{and} \quad H(:,j) = \tilde{H}(:,j)/\lambda_j.$$

Choose the initial F , G , and H randomly unless you can think of something more clever.

MATLAB Tensor Toolbox: **The Function** `cp_als`

```
n = [ 5 6 7 ]; rmax = 35;  
% Generate a random tensor...  
A = tenrand(n);  
for r = 1:rmax  
    % Find the closest length-r ktensor...  
    X = cp_als(A,r);  
    % Display the fit...  
    E = double(X)-double(A);  
    fit = norm(reshape(E,prod(n),1));  
    fprintf('r = %1d, fit = %5.3e\n',r,fit);  
end
```

The function `cp_als` returns a ktensor. Default values for the number of iterations and the termination criteria can be modified:

```
X = cp_als(A,r,'maxiters',20,'tol',.001)
```

Problem 5.8. Compare the efficiency of MyKruskal and cp_als.

Rank-1 Tensors: Definition

If $u_k \in \mathbb{R}^{n_k}$ for $k = 1:d$, then

$$\mathcal{B} = u_1 \circ u_2 \circ \cdots \circ u_d,$$

is defined by

$$\mathcal{B}(i_1, \dots, i_d) = u_1(i_1) \cdot u_2(i_2) \cdots u_d(i_d)$$

is a **rank-1 tensor**. Note that $\mathcal{B} \in \mathbb{R}^{n_1 \times \cdots \times n_d}$.

The CP Representation: General Order

Rank-1 Tensors: Modal Unfoldings

If

$$\mathcal{B} = u_1 \circ u_2 \circ \cdots \circ u_d,$$

then

$$\text{vec}(\mathcal{B}) = u_d \otimes \cdots \otimes u_2 \otimes u_1$$

and

$$\mathcal{B}_{(k)} = u_k \otimes (u_d \otimes \cdots \otimes u_{k+1} \otimes u_{k-1} \otimes \cdots \otimes u_1)^T.$$

Problem 5.9. Suppose $\mathcal{B} \in \mathbb{R}^{n_1 \times \cdots \times n_d}$ is the rank-1 tensor defined above. Characterize the unfolding $M = \text{tenmat}(\mathcal{B}, [1:p], [p+1:d])$ where $1 \leq p < d$.

The CP Representation: General Order

Notation

Given $\lambda \in \mathbb{R}^r$ and matrices U_1, \dots, U_d with unit column norms, define

$$[[\lambda; U_1, \dots, U_d]] = \sum_{j=1}^r \lambda_j \cdot U_1(:,j) \circ \dots \circ U_d(:,j)$$

Assume that $U_k \in \mathbb{R}^{n_k \times r}$.

A weighted sum of rank-1 tensors where the vectors that specify the rank-1's are columns of the matrices U_1, \dots, U_d .

The CP Representation: General Order

The Kruskal Form

We say that $\mathcal{X} \in \mathbb{R}^{n_1 \times \dots \times n_d}$ is in **Kruskal form** if

$$\mathcal{X} = [[\lambda; U_1, \dots, U_d]].$$

where $\lambda \in \mathbb{R}^r$ and $U_k \in \mathbb{R}^{n_k \times r}$ for $k = 1:d$.

*Can we write a given tensor \mathcal{A} as an illuminating sum of rank-1 tensors?
I.e., Given \mathcal{A} , can we find a tensor \mathcal{X} in Kruskal form so that $\mathcal{A} \approx \mathcal{X}$ in
some meaningful way?*

The CP Representation: General Order

Equivalent Formulations...

If $\mathcal{X} = [[\lambda; U_1, \dots, U_d]] \in \mathbb{R}^{n_1 \times \dots \times n_d}$, then

$$\mathcal{X}(i_1, \dots, i_d) = \sum_{j=1}^r \lambda_j \cdot U_1(i_1, j) \cdots U_d(i_d, j)$$

$$\text{vec}(\mathcal{X}) = \sum_{j=1}^r \lambda_j \cdot U_d(:, j) \otimes \cdots \otimes U_1(:, j)$$

MATLAB Tensor Toolbox: **ktensor** Operations

```
function X = KruskalRandn(n,r)
% Creates a random order-d, length-r
% ktensor having size determined
% by the length-d integer vector n. The
% columns of X.U{1},...,X.U{r} have unit
% 2-norm.
U = cell(r,1);
lambda = ones(r,1);
for k=1:d
    U{k} = randn(n(k),r);
end
X0 = ktensor(lambda,U);
X = arrange(X0);
```

The function `arrange` normalizes the columns of the matrices that define `X0` so that they have unit length. The weight vector is adjusted so that `double(X)` and `double(X0)` have the same value.

Problem 5.10. Implement a function $Y = \text{MyArrange}(X)$ that normalizes a ktensor X in the same way as `arrange`.

The CP Representation: General Order

Unfolding Tensors in Kruskal Form

If

$$\mathcal{X} = [[\lambda; U_1, \dots, U_d]] = \sum_{j=1}^r \lambda_j \cdot U_1(:,j) \circ \dots \circ U_d(:,j)$$

then

$$\mathcal{X}_{(k)} = U_k \cdot \text{diag}(\lambda_i) \cdot (U_d \odot \dots \odot U_{k+1} \odot U_{k-1} \odot \dots \odot U_1)^T$$

Note that Khatri-Rao products can be sequenced:

$$F \odot G \odot H = (F \odot G) \odot H = F \odot (G \odot H).$$

The CP Representation: General Case

The CP Approximation Problem

Given $\mathcal{A} \in \mathbb{R}^{n_1 \times \cdots \times n_d}$ and r , determine

$$\mathcal{X} = [[\lambda; U_1, \dots, U_d]] \in \mathbb{R}^{n_1 \times \cdots \times n_d}$$

so that

$$\|\mathcal{A} - \mathcal{X}\|_F = \|\mathcal{A}_{(k)} - \mathcal{X}_{(k)}\|_F$$

is minimized where

$$\mathcal{X}_{(k)} = U_k \cdot \text{diag}(\lambda_i) \cdot (U_d \odot \cdots \odot U_{k+1} \odot U_{k-1} \odot \cdots \odot U_1)^T$$

The CP Representation: General Case

The Alternating Least Squares Framework

for $k = 1:d$

Fix $U_1, \dots, U_{k-1}, U_{k+1}, \dots, U_d$.

Improve λ and U_k by minimizing

$$\| \mathcal{A}_k - \tilde{U}_k (U_d \odot \dots \odot U_{k+1} \odot U_{k-1} \odot \dots \odot U_1) \|$$

for $j = 1:r$

$$\lambda_j = \| \tilde{U}_k(:, j) \|$$

$$U_k(:, j) = \tilde{U}_k(:, j) / \lambda_j$$

end

end

Problem 5.11. Assume that $B_k \in \mathbb{R}^{n_k \times r}$ for $k = 1:d$. (a) Show how to compute

$$\tilde{d} = (B_1 \odot \cdots \odot B_d)^T d$$

efficiently where $d \in \mathbb{R}^N$ with $N = n_1 \cdots n_d$. (b) Show how to compute efficiently

$$C = (B_1 \odot \cdots \odot B_d)^T (B_1 \odot \cdots \odot B_d).$$

(c) Write a MATLAB function `x = KRLS(B,d)` that solves the least squares problem

$$\min \| (B_1 \odot \cdots \odot B_d)x - d \|$$

Assume that `B` is a cell array that houses the matrices B_1, \dots, B_d . (See Problem 5.6.)

Problem 5.12. Refer to Problem 5.7 and develop a general order version of `X = MyKruskal(A,r,itMax)` based on the preceding alternating least squares framework. Take advantage of the ideas in Problem 5.11. Compare your implementation with `cp_als`.

What About r ?

In the CP approximation problem we have assumed that r , the length of the approximating ktensor, is given:

$$\mathcal{A} \approx \mathcal{X} = \sum_{j=1}^r \lambda_j U_1(:,j) \circ \cdots \circ U_d(:,j)$$

We can think of \mathcal{X} as a rank- r approximation to \mathcal{A} .

Departure from Matrix Case...

Suppose

$$\mathcal{X}_r = \sum_{j=1}^r \lambda_j U_1(:,j) \circ \cdots \circ U_d(:,j)$$

is the best rank- r approximation of \mathcal{A} and

$$\mathcal{X}_{r+1} = \sum_{j=1}^{r+1} \lambda_j \hat{U}_1(:,j) \circ \cdots \circ \hat{U}_d(:,j)$$

is the best rank- $(r+1)$ approximation of \mathcal{A} .

IT DOES NOT FOLLOW THAT \mathcal{X}_{r+1} is \mathcal{X}_r plus a rank-1.

In this regard, the best Kruskal approximation is not SVD-like.

Definition

The rank of a tensor \mathcal{A} is the smallest number of rank-1 tensors that sum to \mathcal{A} .

This agrees with the definition for matrices. But there are some differences that make tensor rank a more complicated issue...

Anomaly 1

The largest rank attainable for an n_1 -by-...- n_d tensor is called the **maximum rank**. It is *not* a simple formula that depends on the dimensions n_1, \dots, n_d . Indeed, its precise value is only known for small examples.

Maximum rank does not equal $\min\{n_1, \dots, n_d\}$ unless $d \leq 2$.

Anomaly 2

If the set of rank- k tensors in $\mathbb{R}^{n_1 \times \dots \times n_d}$ has positive Lebesgue measure, then k is a **typical rank**.

Size	Typical Ranks
$2 \times 2 \times 2$	2,3
$3 \times 3 \times 3$	4
$3 \times 3 \times 4$	4,5
$3 \times 3 \times 5$	5,6

For n_1 -by- n_2 matrices, typical rank and maximal rank are both equal to the small of n_1 and n_2 .

Anomaly 3

The rank of a particular tensor over the real field may be different than its rank over the complex field.

Anomaly 4

A tensor with a given rank may be approximated with arbitrary precision by a tensor of lower rank. Such a tensor is said to be **degenerate**.

Problem 5.13. For various small choices for $\mathbf{n} = [n_1, \dots, n_d]$, see if you can discover the typical rank possibilities by using `cp_als`. In particular, for a randomly generated \mathcal{A} , compute the smallest r such that

$$\| \mathcal{A} - \text{cp_als}(\mathcal{A}, r) \|_F \leq 10^{-6}$$

By running a sufficient number of examples, see what you can deduce about the typical rank of tensors in $\mathbb{R}^{n_1 \times \dots \times n_d}$.

Summary of Lecture 5.

Key Words

- An order- d **rank-1 tensor** is the outer product of d vectors.
- A tensor in **Kruskal form** has length r , if it is the sum of r rank-1 tensors. For an order- d tensor, the vectors that make up the rank-1's are specified as columns from d matrices.
- The **CP approximation** problem for a given tensor \mathcal{A} and a given integer r involves finding the nearest length- r k-tensor to \mathcal{A} in the Frobenius norm.
- The **alternating least squares** framework is used by `cp_als` to solve the CP approximation problem. It proceeds by solving a sequence of structured linear least squares problems.