Exercise 1: Inline Assembly

Refer to file ex1.c

The following inline assembly code will increment the value of x by 1.

```
asm ("incl %0":"+r"(x));
```

- `incl` instruction increments the operand by 1.
- `+r` is used to allocate any free register to the variable x and use that register as both Input and Output
- `%0` corresponds to the register allocated to x.

Exercise 2: GDB

A few starting instructions of `BIOS` are:

```
0xffff0:   ljmp    $0x3630,$0xf000e05b
0xfe05b:   cmpw    $0xffc8,%cs:(%esi)
0xfe060:   jo      0xfe062
0xfe062:   jne     0xd241d416
0xfe068:   mov     %edx,%ss
0xfe06a:   mov     $0x7000,%sp
0xfe06e:   add     %al,(%eax)
0xfe070:   mov     $0x2d4e,%dx
0xfe074:   verw    %cx
0xfe077:   xchg    %ebx,(%esi)
0xfe079:   push    %bp
0xfe07b:   push    %di
0xfe07d:   push    %si
0xfe07f:   push    %bx
0xfe081:   sub     $0x70,%sp
0xfe085:   mov     %ax,%di
0xfe088:   mov     0x4(%bx,%si),%si
0xfe08d:   mov     %cs:0x2c(%bp),%bl
0xfe093:   icebp
0xfe094:   ljmp    *(%esi)
0xfe096:   mov     0x2d(%bp),%al
0xfe09b:   icebp
```

The `BIOS` first initializes all the PCI bus and all other peripheral devices. Then it loads the `bootloader` from the `hardisk` into memory. Finally with a jump statement control goes to the `bootloader`.

Exercise 3: Loading Kernel from Bootloader

Trace: Refer to file Bootloader Trace.pdf

(a) Following instructions change the addressing to 32 bit protected mode.

```
0x7c1d: lgdt    gdtdesc                            # lgdt (%esi)
0x7c22: mov     %cr0,%eax
0x7c25: or      $0x1,%ax
0x7c29: mov     %eax,%cr0
0x7c2c: ljmp    $(SEG_KCODE<<3), $start32    # ljmp $0xb866,$0x87c31
```

After this point processor starts executing 32 bit code. First instructions it executes in 32 bit is:-

```
0x7c31: mov     $0x10,%ax
```

(b) The last instruction that bootloader executed is

```
0x7d87: call    *0x10018
```

This instructions is calling the entry function found in ELF Header. In `bootmain.c` this corresponds to following lines.

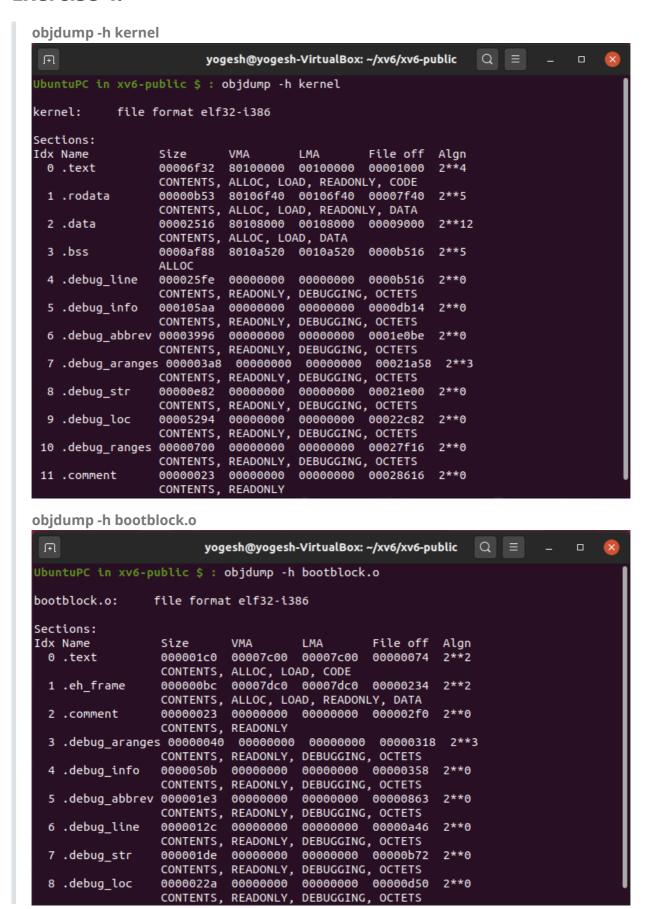```
entry = (void(*)(void))(elf->entry);
entry();
```

First instruction of kernel is

```
0x10000c: mov     %cr4,%eax
```

(c) This information is stored in elf header. First the bootloader loads first 4096 bytes (1st page) into memory. This page contains elf header which has an array of program headers. these program headers contains the size and offset of different segments of kernel which are then loaded into memory.

```
// Load each program segment (ignores ph flags).
ph = (struct proghdr*)((uchar*)elf + elf->phoff);
eph = ph + elf->phnum;
for(; ph < eph; ph++){
  pa = (uchar*)ph->paddr;
  readseg(pa, ph->filesz, ph->off);
  if(ph->memsz > ph->filesz)
    stosb(pa + ph->filesz, 0, ph->memsz - ph->filesz);
}
```

Exercise 4:

objdump -h kernel

```
UbuntuPC in xv6-public $ : objdump -h kernel

kernel:     file format elf32-i386

Sections:
Idx Name          Size      VMA       LMA       File off  Algn
  0 .text         00006f32  80100000  00100000  00001000  2**4
                  CONTENTS, ALLOC, LOAD, READONLY, CODE
  1 .rodata       00000b53  80106f40  00106f40  00007f40  2**5
                  CONTENTS, ALLOC, LOAD, READONLY, DATA
  2 .data         00002516  80108000  00108000  00009000  2**12
                  CONTENTS, ALLOC, LOAD, DATA
  3 .bss          0000af88  8010a520  0010a520  0000b516  2**5
                  ALLOC
  4 .debug_line   000025fe  00000000  00000000  0000b516  2**0
                  CONTENTS, READONLY, DEBUGGING, OCTETS
  5 .debug_info   000105aa  00000000  00000000  0000db14  2**0
                  CONTENTS, READONLY, DEBUGGING, OCTETS
  6 .debug_abbrev 00003996  00000000  00000000  0001e0be  2**0
                  CONTENTS, READONLY, DEBUGGING, OCTETS
  7 .debug_aranges 000003a8  00000000  00000000  00021a58  2**3
                  CONTENTS, READONLY, DEBUGGING, OCTETS
  8 .debug_str    00000e82  00000000  00000000  00021e00  2**0
                  CONTENTS, READONLY, DEBUGGING, OCTETS
  9 .debug_loc    00005294  00000000  00000000  00022c82  2**0
                  CONTENTS, READONLY, DEBUGGING, OCTETS
 10 .debug_ranges 00000700  00000000  00000000  00027f16  2**0
                  CONTENTS, READONLY, DEBUGGING, OCTETS
 11 .comment      00000023  00000000  00000000  00028616  2**0
                  CONTENTS, READONLY
```

objdump -h bootblock.o

```
UbuntuPC in xv6-public $ : objdump -h bootblock.o

bootblock.o:    file format elf32-i386

Sections:
Idx Name          Size      VMA       LMA       File off  Algn
  0 .text         000001c0  00007c00  00007c00  00000074  2**2
                  CONTENTS, ALLOC, LOAD, CODE
  1 .eh_frame     000000bc  00007dc0  00007dc0  00000234  2**2
                  CONTENTS, ALLOC, LOAD, READONLY, DATA
  2 .comment      00000023  00000000  00000000  000002f0  2**0
                  CONTENTS, READONLY
  3 .debug_aranges 00000040  00000000  00000000  00000318  2**3
                  CONTENTS, READONLY, DEBUGGING, OCTETS
  4 .debug_info   0000050b  00000000  00000000  00000358  2**0
                  CONTENTS, READONLY, DEBUGGING, OCTETS
  5 .debug_abbrev 000001e3  00000000  00000000  00000863  2**0
                  CONTENTS, READONLY, DEBUGGING, OCTETS
  6 .debug_line   0000012c  00000000  00000000  00000a46  2**0
                  CONTENTS, READONLY, DEBUGGING, OCTETS
  7 .debug_str    000001de  00000000  00000000  00000b72  2**0
                  CONTENTS, READONLY, DEBUGGING, OCTETS
  8 .debug_loc    0000022a  00000000  00000000  00000d50  2**0
                  CONTENTS, READONLY, DEBUGGING, OCTETS
```

`objdump -h` displays the header of an executable file. In this case it displays the contents of program section headers of the ELF Binaries.

The important program sections in an ELF Binary -

- `.text` - All the executable instructions of the program
- `.rodata` - The read-only data of the program like the ASCII string constants in C.
- `.data` - The initialized global and static variables in the program.
- `.bss` - The uninitialized global and static variables in the program.

Each section has the following information -

- `LMA(Load memory address)` - The address at which the section is actually loaded in the memory.
- `VMA(Virtual memory address)` - The address at which the binary assumes the section will be loaded.
- `Size` - The size of the section.
- `Offset` - The offset from the beginning of the harddrive where the section is located at.
- `Algn` - The value to which the section is aligned in memory and in the file.
- `CONTENTS, ALLOC, LOAD, READONLY, DATA, CODE` - Flags which gives additional information regarding the section. Eg. Is it READONLY, should it be LOADED etc.

Exercise 5: Bootloader's Link address

If we get wrong `bootloader's` link address, then the 1st instruction that would break is

```
ljmp    $(SEG_KCODE<<3), $start32
```

With correct `bootloader's` link address the output was:

```
[    0:7c2c] => 0x7c2c:  ljmp    $0xb866,$0x87c31
 The target architecture is assumed to be i386
=> 0x7c31:   mov     $0x10,%ax
=> 0x7c35:   mov     %eax,%ds
=> 0x7c37:   mov     %eax,%es
```

The output when `bootloader's` link address is changed to 0x7C04:

```
[    0:7c2c] => 0x7c2c:  ljmp    $0xb866,$0x87c35
[f000:e05b]     0xfe05b: cmpw    $0xffc8,%cs:(%esi)
[f000:e062]     0xfe062: jne     0xd241d416
[f000:d414]     0xfd414: cli
```

The `ljmp` instruction breaks because in the `BIOS` the address 0x7C00 is hard coded, so `BIOS` always loads `bootloader` starting from 0x7C00. But the `linker` converts the code into binary form and assigns addresses in place of labels taking `bootloader's` link address(0x7C04) as the starting address of the `bootloader` in the memory. So the address of the label `$start32` in the `ljmp` instruction doesn't contain the correct instruction and this causes some error. Hence the `BIOS` restarts (execution reaches starting instruction of `BIOS`). This process then repeats and in turn leads to an infinite loop.

The file headers of `kernel` are

```
kernel:     file format elf32-i386
architecture: i386, flags 0x00000112:
EXEC_P, HAS_SYMS, D_PAGED
start address 0x0010000c
```

This shows that entry point of `kernel` is 0x0010000c.

Exercise 6:

After entering the bootloader (at `0x7c00`):

```
(gdb) x/8x 0x00100000
0x100000: 0x00000000   0x00000000   0x00000000   0x00000000
0x100010: 0x00000000   0x00000000   0x00000000   0x00000000
```

After entering the kernel (at `0x10000c`):

```
(gdb) x/8x 0x00100000
0x100000: 0x1badb002   0x00000000   0xe4524ffe   0x83e0200f
0x100010: 0x220f10c8   0x9000b8e0   0x220f0010   0xc0200fd8
```

The code for kernel is stored from memory location `0x00100000`, which is loaded from the disk by the bootloader.

At the point the BIOS enters the bootloader, this loading is not done, hence the main memory does not contain the kernel code. Moreover, it is filled with zeroes because upto this point the system runs in the 20-bit real mode and any memory location from this address onwards is not touched.

At the point the bootloader enters the kernel, the bootloader has already loaded the kernel and there are instructions from that memory location.

The second breakpoint is the entry point of the kernel. The first intructions starting from this location are responsible for turning on paging (which wasn't enabled upto this point).

Exercise 7: Adding System Call

For creating a system call, we need to change 6 files:- <u>user.h</u>, <u>syscall.h</u>, <u>syscall.c</u>, <u>usys.S</u>, <u>defs.h</u>, <u>sysproc.c</u>

```c
// user.h
int wolfie(void* buf, uint size);         // line 26
```

```c
// syscall.h
#define SYS_wolfie 22                      // line 23
```

```c
// syscall.c
extern int sys_wolfie(void);               // line 106
[SYS_wolfie]  sys_wolfie,                  // line 130
```

```asm
// usys.S
SYSCALL(wolfie)                            // line 32
```

```c
// defs.h
int wolfie(void*, uint);                   // line 123
```

```c
// sysproc.c
int sys_wolfie(){                          // line 94
  char* buf;
  uint size;
  if(argptr(0, (void*)&buf, sizeof(buf)) < 0) return -1;
  if(argptr(0, (void*)&size, sizeof(size)) < 0) return -1;

  static char wolf[] = \
"                        \n"
"                   .    \n"
"               / V\\   \n"
"              / `  /   \n"
"             <<   |    \n"
"             /    |    \n"
"           /      |    \n"
"          /       |    \n"
"         /    \\  \\ /    \n"
"        (      ) | |    \n"
" _____|   _/_  | |    \n"
"<_____\\\_____)\\\\__)   \n"
"                        \n";

    static uint wolf_len = sizeof(wolf);
    if(size < wolf_len) return -1;
```

```
    int i = 0;
    while(wolf[i] != '\0'){
      buf[i] = wolf[i];
      ++i;
    }
    buf[i] = '\0';

    return wolf_len;
  }
```

Exercise 8: User Level Application

We created wolfietest.c in which we created a buffer and used system call to fill that buffer with wolf ASCII image. Then we printed this buffer to console using `printf`. 1st parameter in `printf` is file descriptor which is 1 for console out. At the end we used `exit` system call to exit from this program.

```
// wolfietest.c
#include "types.h"
#include "user.h"

int main(int argc, char *argv[]){
  printf(1, "I am a wolf. wooo.....\n\n");
  char wolf[500];
  wolfie(wolf, 500);
  printf(1, wolf);
  exit();
}
```

In Makefile we need to add `_wolfietest\` to `UPROGS` and `wolfietest.c` to `XTRA`.

```
// Makefile
UPROGS=\
  _cat\
  _echo\
  _forktest\
  _grep\
  _init\
  _kill\
  _ln\
  _ls\
  _mkdir\
  _rm\
  _sh\
  _stressfs\
  _usertests\
```

```
    _wc\
    _zombie\
    _wolfietest\                                                    #
 line 184


XTRA=\
    mkfs.c ulib.c user.h cat.c echo.c forktest.c grep.c kill.c\
    ln.c ls.c mkdir.c rm.c stressfs.c usertests.c wc.c zombie.c wolfietest.c\  #
 line 251
```

Output