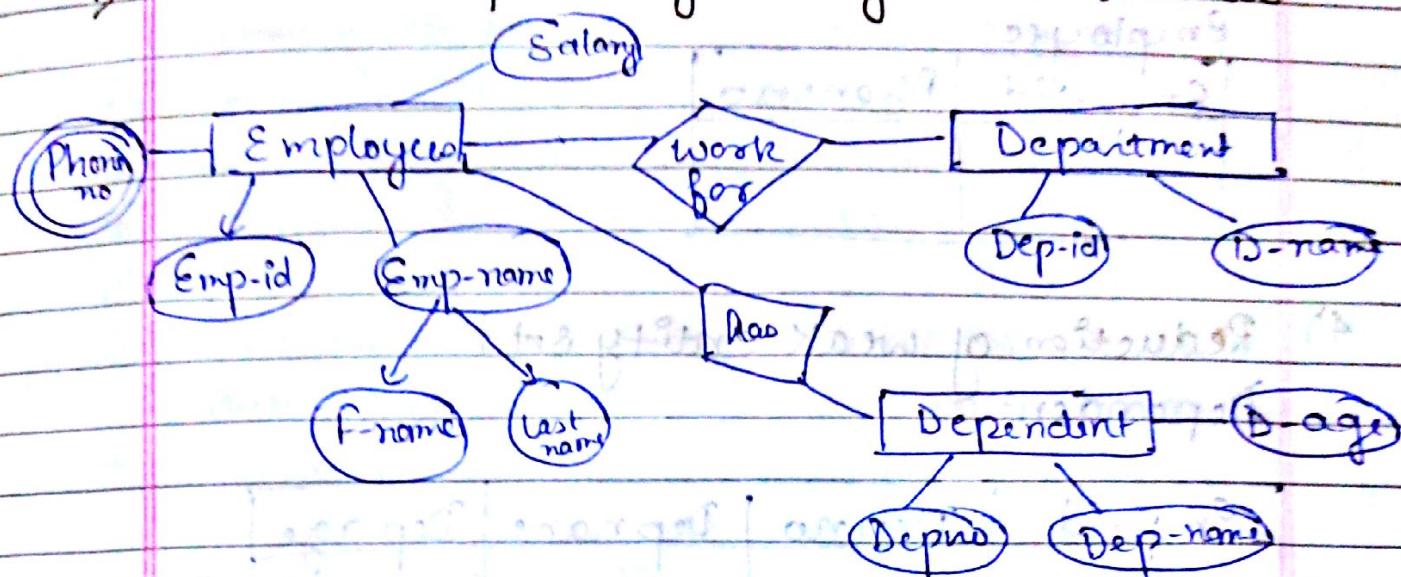


Q5 Construct an E-R diagram of University system.

Reduction of E-R diagram into tables

i) Reduction of Strong entity Set into table



Key attributes of Employee

Empid	E-name	Salary

Department

Dep-id	Dept-name

Reduction of Composite attribute into table

Employee

Emp-id	f-name	l-name	Salary

3) Reduction of multi valued attributes
Phoneno.

Employee

Emp-id	Phoneno

4) Reduction of weak entity set +
Dependent

Emp-id	Dep no	Depname	Dep age

5) Reduction of relationship Sets

work for

Emp-id	Department id

RDBMS :-

- 3 Components :-
- Data Structure
- Data integrity
- Data Manipulation

Candidate key is a subset of superkey.

Key :- It is a set of one or more columns whose combined values are unique among all the

① Candidate Key:

Occurrences in a given table for the all

Types of keys?

- 1) Candidate key.
- 2) Super key
- 3) Primary key
- 4) Foreign key
- 5) Alternate key
- 6) Artificial key
- 7) Composite key

1) Candidate Key: They are those attributes of the relational, which have the properties of uniqueness and irreducibility.

- a) Uniqueness: No legal value of R ever contains two distinct tuples with same values for K.
- b) Irreducibility: No proper subset of K has the uniqueness property.

For example if the combination of (Name, Class) is unique, then it can be identified as the candidate key if and only if Name and Class individually are not unique.

- 2) Super key: Super key follow the property of uniqueness, but not irreducibility. A Super key has a uniqueness property but not necessarily the irreducibility property. A candidate key is a special case of a Super key.

e.g. If Roll-no is unique in relation STUDENT then, the set of attribute (Roll-no, Name, Class) is a Super key for a relation STUDENT, these set of attributes are also unique, but this combination of key is not having the property of irreducibility.

→ Relation of patient in which Patient-number is unique. The patient-number is a candidate key and (Patient-number, Patient name) is a superkey.

'A superset of a candidate key is a superkey.'

3) Primary Key: The primary key is an attribute or a set of attributes that uniquely identify a specific instance of an entity. Primary key cannot contain any null value because we cannot uniquely identify multiple null values. ~~so~~ primary key is a candidate which is chosen by database designer as primary key.

4) Alternate Key: Exactly one of those candidate keys is chosen as the primary key and the remainder, if any, are then called Alternate Keys. An alternate key is a function of all candidate key minus primary key.

5) Composite keys: A primary key that is made up of more than one attribute is known as Composite key.

Ex :- WORK

Employee ID	Project ID	Hours Worked
01	01	200
01	02	120
02	01	50
02	03	120
03	03	100
03	04	200

It is subset of Superkey.

- 6) Artificial keys :- An artificial key is one that has no meaning to the business or organization and perform function of a primary key in relation.
- Artificial key are permitted when
- no attribute has all the primary key properties
 - the primary key is large and complex

Eg:-

Enrollment		
Student	Class	Row-id
AK	PK	

- 7) Foreign key :- Foreign key are the attributes of the table which refers to primary key of some another table. Foreign key are used to link together two or more different tables which have some form of relationship with each other. These foreign keys is a reference to the tuple of a table from which it was taken, this tuple being called as Reference or target tuple.

Target
Attribute

Cg: Employee

Foreign key

Department → Target table

Emp-id	Name	Salary	Dep-id
1	Abc	30000	10
2	Xyz	40000	20
3	Pqr	50000	30

Dep-id	Dep-Name
10	Sales
20	Market
30	Product

eg: Student

Rno	Name	Class-Code
1	A	2
2	B	1
3	C	-

class

Classcode	Name
1	B.TECH
2	B.TECH
3	BBA

- ② Data integrity :- Basically it consists of two rules
- Entity Integrity Rule
 - Referential Integrity Rule

① Entity Integrity Rule :- This rule states that in a relation, the value of the attribute of a primary key cannot be null.

② Referential Integrity :- It states that if a foreign key exists in a relation, either the foreign key value must match the primary key value of some tuple in its home relation or the foreign key value must be completely null.

Employee				Department	
Emp-id	Name	Salary	Dep-id	Dep-id	Dept name
1	A	3000	10	10	Sales
2	B	4000	20	20	Market
3	C	5000	30	30	Product
4	D	6000	40		

either should match

Enterprise Constraints: These are the additional rules which are given by the user & DBA of a database. e.g. $<$, $>$.

Codd's Rule: Founder of RDBMS is Dr. E. F. Codd
There are 12 rule are as following:

- ① Information rule: All information represented in form of tables
- ② Guaranteed Access rule: To Access any info we use key.
- ③ Comprehensive data Sublanguage rule: There should be a particular language to support RDBMS rule.
- ④ View updating rules: If we Change any data in Database then it should be automatically updated in all records or tables.
- ⑤ High level Insert, update & Delete: The language we used it should contain these rule insert, update & Delete and perform all function.
- ⑥ Physical data independency: Change in lower level not effect higher level and this rule support.
- ⑦ Logical Data independency: Change in conceptual level and does not effect external level.
- ⑧ Integrity Independency: RDBMS support all Integrity rule.
- ⑨ NonSubversion rule: Any language we use to access the database and that language will support our Integrity independency.
- ⑩ Systematic treatment of null Value: There should be special treatment of null value.
- ⑪ Data Description rule: Data we describe it should be in form of table

(12) Distribution & Independence \therefore It should be platform independent.

Data Manipulation \therefore

① Relational Algebra

② Relational calculus.

Q Difference b/w Relational Algebra & Relational Calculus.

Relational Algebra

① It is procedural language

In this we follow the procedure & step

Relational Calculus.

① It is a non procedural language

In this we automatically get output

② We can combine 2 or more table to get an another table

Both are non user friendly.

Relational Algebra $\therefore J. A \oplus B$

Relational operation \therefore Type of Relational operator

① Traditional set operators

② Special operators

① Traditional Set Operators $\therefore \rightarrow$ Union

\rightarrow Intersection

\rightarrow Difference

\rightarrow Cartesian Product

Union \therefore Combine two table Employee Dep-id.

Emp-id	Name	Dept-id	Depid	Name

→ Intersection: Common element

→ Difference: The Difference between two sets S_1 & S_2 produces a set, which contains all the members of one set, which are not in the other.

R

Cust-name	Cust-status
Ram	Good
Shyam	Excellent

Cust-name	Cust-status
Karan	Average
Ram	Good

① RUS

Cust-name	Cust-status
Ram	Good
Shyam	Excellent
Karan	Average

②

RNS

Cust-name	Cust-status
Shyam	
Ram	Good

③

$R - S \Rightarrow$ Shyam.

④

$R \times S = \{(R, S) : (R, S) \in R \text{ and } (S, R) \in S\}$

Name	Rollno
Ami	10
Mym	8

Name	Roll
Key	11
Shonu	12

$$R \times S = \{(A, B) : A \in R \text{ and } B \in S\}$$

ii)

Special operators

→ Selection

→ Projection

→ Join

→ Division

①

Selection operation: It yields a horizontal subset of a given relation that is the subset of row which should be selected within the given relation for which a particular

Condition is satisfied. Sign of Selection is (σ)

Employee :-

Emp_id	Name	Salary
1001	John	10000
1002	Mike	12000
1003	David	15000
1004	Steve	18000

Pooroint :- $\sigma_{\text{salary} > 10,000}$ (Employee)

② Projection :- The Projection operation on a table, simply form another table by copying Specified columns from the original table. Symbol of projection is (Π)

To Select name of employee :-

$\Pi_{\text{name}} (\text{Employee})$

for salary :-

$\Pi_{\text{salary}} (\text{Employee})$

We want the name of all the employee having salary less than 70000

$\sigma_{\text{salary} < 70000} \Pi_{\text{name}} (\text{Employee})$

Codd's Rules :-

- ① **Information Rule :-** All information in a relational database including tables names, column names is represented in the form of tables. This simple view of data speeds up design and learning process; user productivity is improved since knowledge of only one language is necessary to access all data such as description of the table and attribute definitions, integrity constraints and action are taken when constraint are violated. Access of data is restricted.
- ② **Guaranteed Access Rule :-** Every piece of data in relational database, can be accessed by using a primary key value that identifies the row name and column name. User productivity is improved since there is no need to resort to using physical pointers or address. It provides data independence or makes it possible to retrieve each piece of data in a relational database.
- ③ **Comprehensive Data Sub-language Rule :-** The RDBMS may support several languages. But at least one of them should allow the user to do all following define table^{view}, query & update data, set integrity constraints, set authorizations & define transaction. User productivity is improved since there is just one approach that can be used for all database operations.
- ④ **View Updating Rule :-** Any view that can be updated theoretically can be updated using the RDBMS. Data consistency is ensured since the changes made in the view are

transmitted to base table and vice-versa.

- ⑤ High Level Insert, update & Delete: The RDBMS support insertion, updating and deletion at a table level. The performance is improved since the commands acts on a set of records rather than one record at a time.

- Application program is not affected by changes in the physical access or storage method. Database administrators can make changes to the physical access and storage method which improve performance and do not require changes in the application programs or request.

- ⑥ Logical Data Independence: Logical changes in tables and views such as adding/deleting columns or changing field lengths need not necessitate modifications in the programs.

- ⑦ Integrity Independence: Integrity constraints are stored in the one-line catalog or data dictionary and can therefore be changed without necessitating changes in the application programs.

- ⑧ Non Subversion Rule: If the RDBMS has a language that accesses the information of records at a time, this language should not be used to bypass the integrity constraints. This is necessary for data integrity.

(10) Systematic Treatment of Null values: In RDBMS null values should be supported for the representation of missing information and in applicable information, the database management must have a consistent method for representing null values.

(11) Database Description Rule: The description of database is stored and maintained in the form of tables. It allows the user with appropriate authority to query information in a similar way or language.

(12) Distribution Independence: The RDBMS package must have distribution independence. Thus RDBMS

Packet must make it possible for the database to be distributed across multiple

Relation Algebra	Relational Calculus
① It is a procedural method for solving Queries	① It is non-procedural method of solving the Queries

② The solution to the database access problem using relational algebra is obtained by stating what is required and what are the steps to obtain that information.

It is used as vehicle for implementation of Relational Calculus.

The solution to the database access problem using a relational calculus is obtained simply by stating what is required and letting the system find the answer.

Relation Calculus Queries are converted into equivalent relational algebra format by using Codd's reduction algorithm and then it is implemen-

with the help of relational algebra operators.

- ④ Relational algebra Operators are used as a yardstick for measuring the expressive power of any given language.

A lang. is said to be Complete if it is at least powerful as the Calculus that is, if any relation definable by some expression of the calculus is also definable by some expression of the language in question.

Difference :-

- ① • Relational Algebra is a procedural language that can be used to tell the DBMS how to build a new relation from one or more relations in the database.
- Relational Calculus is non-procedural language that can be used to formulate the definition of relation in terms of one or more database relations.
- ② Relation Algebra : user has to specify what is required and what are the procedure or steps to obtain output.
- Relational Calculus : user just specifies what is required and need not to specify how to obtain it.

Strong Entity Set

- ① It has its own primary key
- ② It is represented by rectangle
- ③ It contains primary key represented by an Underline
- ④ The member of Strong entity set is called dominant entity set
- ⑤ The relationship between two strong entity set is represented by diamond System
- ⑥ The line connecting Strong entity set with the relationship is single
- ⑦ Total participation in the relationship may or may not exist

weak Entity Set

It does not have sufficient attributes to form a primary key on its own.

It is represented by double rectangle.

It contains Partial key represented by dashed underline

The member of weak entity set is called Subordinate entity set

The relationship b/w one strong and ~~one~~ weak entity set is represented by a double diamond sign. i.e Identifying relationship

The line connecting weak entity set with identifying relationship is double

Total participation in the identifying relationship always exists

Relational Algebra:

Joins: General form: $\exists \theta$ such that $R \theta S$

\rightarrow Theta Join.

Theta Joins: It is a Cartesian product operation on the two tables followed by a restriction operation on the resultant tables.

Eg:

Employee Product	Name	Product
	Raja	Pen
	Sparsh	Pen
	Raja	Pencil
	Sparsh	Rubber

Product Customers

C-Product	Customer
Pen	Karan
Pen	Suneet
Pencil	Suneet

$\exists Raja \text{ Pen}$

Name	Product	Customer
Raja	Pen	Karan

Operations

\rightarrow Comparison operator

\rightarrow equal =

\rightarrow Not equal to a

\rightarrow Greater than >

	Name	Product	C-Product	Customer
Equal	Raja	Pen	Pen	Karan
	"	"	Pen	Suneet
Not equal	"	"	Pencil	Suneet
	Sparsh	Pen	Pen	Karan
	"	"	Pen Pencil	Suneet Suneet

Equi Joins :-

Natural Joins :- The projection operation which eliminates one of the duplicated column resulting from the equijoin, the natural join is obtained.

Name	Product	Customer
Raju	Pen	Karan
Sunil	"	Suneet

Inner Joins :-

→ Outer Joins :- If it requires stack the value exist in only one table must appear in the output then the solution is outer joins.

Normal Product

These are of two type

- 1) Left outer joins
- 2) Right outer joins.

Relational Calculus:-

Difference b/w Algebra and Relational calculus

Relational calculus

→ Tuple

→ Domain

Normalization:

- 1) Functional Dependency: In Relation R, X & Y are the two subset are the set of attributes, Y is said to be function dependent on X if a given value of X uniquely determine the value of Y.
- It is denoted by $X \rightarrow Y$ where X is called determinant and Y is called determined.

Employee

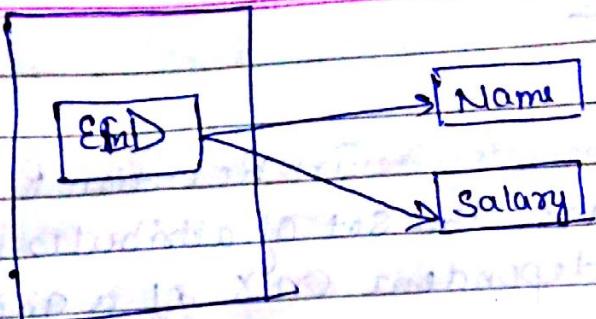
EID	Name	Salary
1	XYZ	30,000
2	Abc	20,000
3	Pqr	10,000

$X(EID) \rightarrow Y(\text{Name}, \text{Salary})$

Functional Dependency chart: It is the graphical representation of functional dependencies among the attributes in any relation.

Steps:

- ① find out the primary key attribute
- ② Make a rectangle with all the primary key attributes inside it
- ③ write all the non-primary key attribute outside the rectangle
- ④ use the arrows to show the functional dependencies among the attributes



→ Types of functional Dependency:

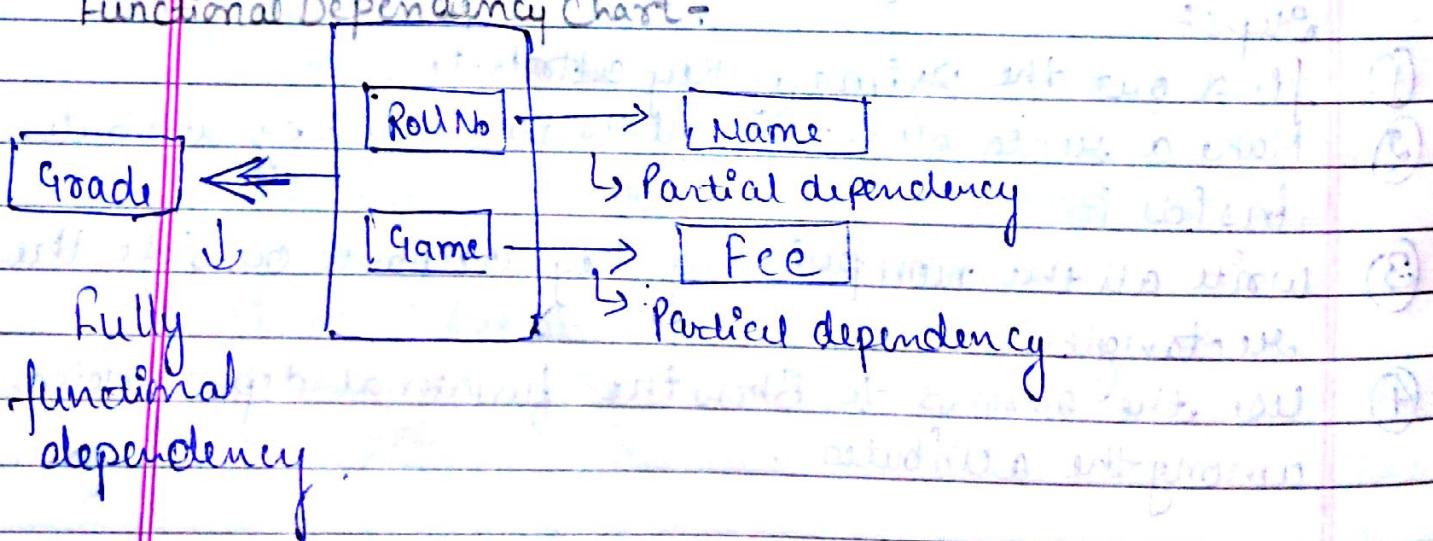
fully

① Partial Dependency:

- Partial Dependency: Suppose we have more than one attribute in primary key. Let A be the non primary key attributes attribute. If A is not dependent upon all the primary key attribute then partial dependency exists.

- Fully functional Dependency: Let A be the non primary key attribute if value of A is dependent upon all the primary key attributes then A is said to be fully functional dependency.

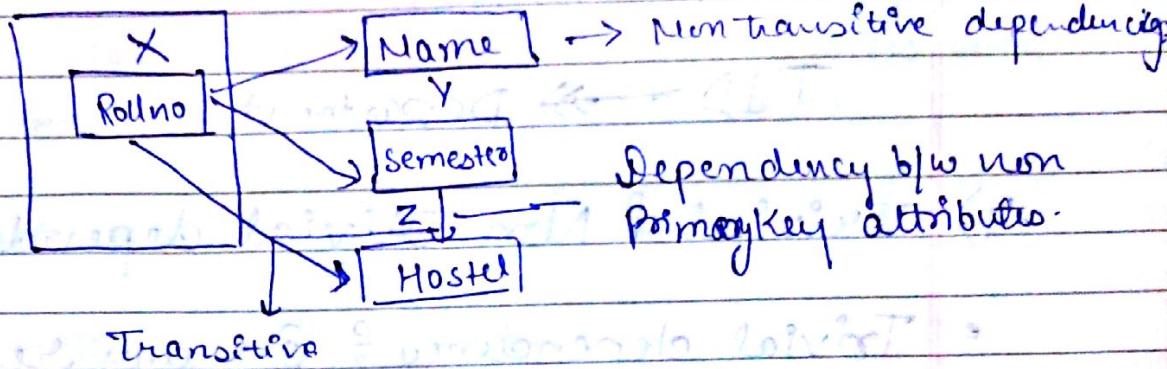
Functional Dependency Chart:



②

Transitive & Non-Transitive dependency :-

- Transitive dependency is due to dependency between the ~~known~~ primary key attributes.
- Suppose in a relation R, $X \rightarrow Y$, $Y \rightarrow Z$ then $X \rightarrow Z$



- Non Transitive :- Any functional dependency which is not transitive is known as Non Transitive dependency.

NOTE :- Non transitive dependency exist if there is no dependency b/w the ~~known~~ primary key attributes

③

Single valued & multiple valued dependency :-

- Single Valued :- In any relation R, if for a particular value of X, Y has single value then it is known as Single valued dependency.

Teacher ID	Name	Department
1	Abc	CSE
1	Abc	ME
2	Xyz	Civil
2	Xyz	ECE
3	Pqr	E.C.E
3	Pqr	CSE

→ for single value

TeacherId → T Name

→ for multiple valued t

TID → Department

4) Trivial & Non Trivial dependency

- Trivial dependency : In any relation R, $X \rightarrow Y$ is trivial if $Y \subseteq X$
 - Non Trivial dependency : In any relation R, $X \rightarrow Y$ if $Y \not\subseteq X$

Anomalies $\hat{=}$ Undesired result

→ || Insertion

→ Deletion

→ update

EID	Name	Salary	DeptNo	Department
4	PQR	39,000	34	Marketing
5.	ABC	20,000	3	Marketing

→ Normalisation : It is a process by which we can decompose or divide any relation into more than one relation to remove the anomalies in the relational database.

Normal forms : It is a step by step process and each step is known as Normal form.

→ Properties of Normalisation :

- ① Remove different anomalies
- ② Decomposition must be lossless
- ③ Preserve the necessary dependency
- ④ Reduce Redundancy

In Normalisation there are 5 Normal form

First Normal Form :

- ↳ Composite attributes.
- ↳ Flattening the table
- ↳ Decomposition of table

Teacher ID	Name	Subject
1	XYZ	ADA DBMS C, C++
2	PQR	Java, Automata ADCE

① Composite attribute

Teacher ID	First Name	Last Name	Subject

(2)

Flattening the table

TeacherID	Name	Subject
1	XYZ	ADA
1	XYZ	DBMS
1	XYZ	C/C++
1	XYZ	C++
2	Pg1	JAVA
2	Pg1	Automation
2	Pg1	ADCP

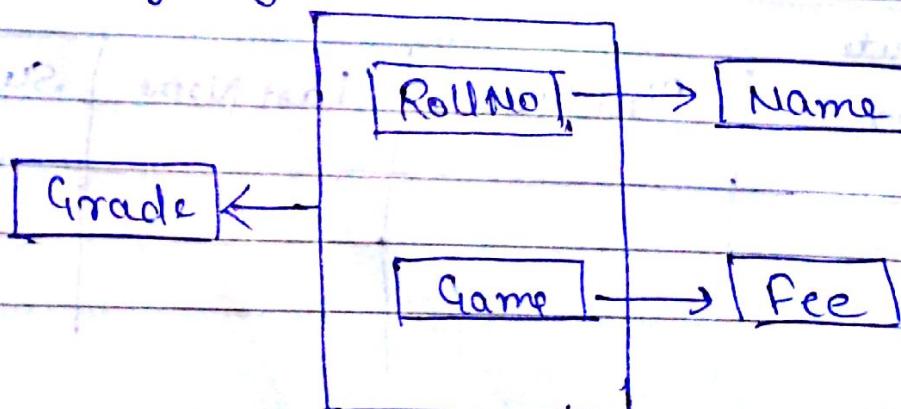
(3) Decomposition of table :

TeacherID	Name	Subject

A relation is in the first normal form if the domain of each attribute contains only atomic values it means atomicity must be present in the relation

(4)

Second Normal form : A relation is in the second normal form if it is in the first normal form and all the non-primary key attribute must be fully functionally dependent upon the primary key attribute



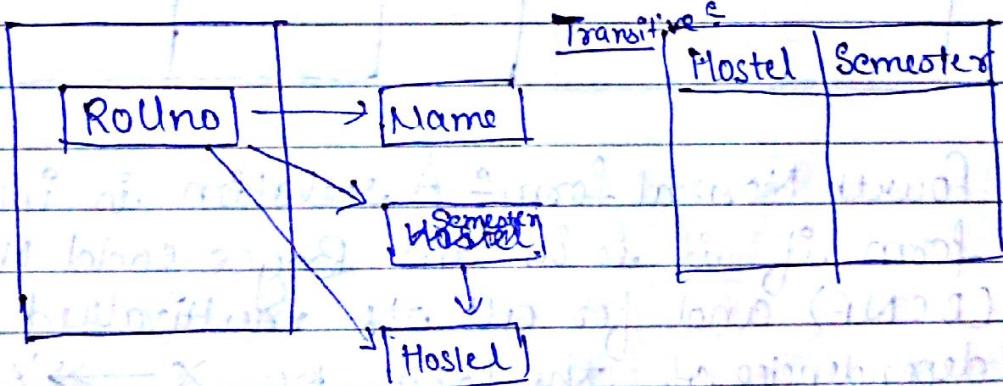
Roll No	Name	Game	Fee	Grade
101	John	Cricket	500	A+
102	Jane	Football	450	B+
103	Mike	Volleyball	400	C+

Roll no	Name	Game	Fee
101	John	Cricket	500
102	Jane	Football	450
103	Mike	Volleyball	400

RollNo	Game	Grade
101	Cricket	A+
102	Football	B+
103	Volleyball	C+

III)

Third Normal form :- A relation is in the third normal form if it is in the Second normal form and non primary key attributes must be non transitively dependent upon primary key attributes.



Nontransitive :-

Rollno	Name	Semester
101	John	1
102	Jane	2
103	Mike	3

BCNF (Boyce Codd Normal form) :-
 A relation is in BCNF if and only if all the determinants are th candidate key

$X \rightarrow Y \rightarrow \text{determined}$
 ↓
 determinant

Roll no.	Teacher	Subject

→ Roll no.
 → (Rollno, Teacher)
 → (Rollno, Subject)

Rollno	TeacherID	TeacherID	Subject

IV: Fourth Normal form: A relation is in fourth form if it is in the Boyce Codd Normal form (BCNF) and for all the multivalued functional dependencies of the form by $X \rightarrow\!\!\! \rightarrow Y$

Programmer	Project	Module

Programmer	Project

Project	Module

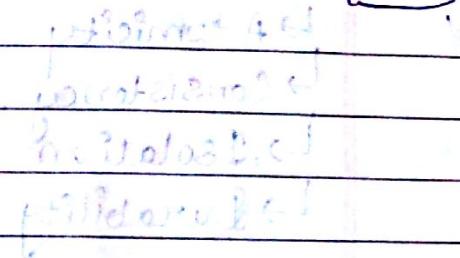
V

Fifth Normal Form (Join dependency):
 Let R be the given relation upto fourth normal form or (4NF). Let R be and it decompose into $(R_1, R_2, R_3, \dots, R_n)$. The Relation R satisfy the join dependencies if and only if joining R_1 to $R_n = R$

X	Y	Z
x_1	y_1	z_1
x_2	y_2	z_2
x_3	y_3	z_3

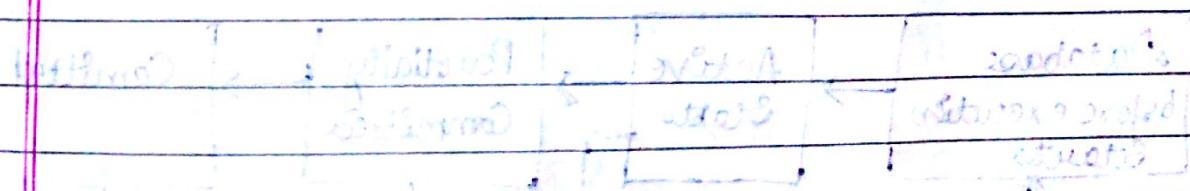
X	Y	Z	X

join dependency



all other forms of dependencies are not sufficient.
 except all functional and multivalued dependencies

footnote position



left hand side
 functional dependency
 multivalued dependency
 partial dependency
 right hand side
 candidate keys
 superkeys
 equivalents

Transaction & concurrency Control

- Transaction: It can be defined as a unit or any part of any program at the time of its execution. During transaction the data items can be read or updated or both.
- Concurrency Control: It is the activity of co-ordinating the action of transactions, that operate simultaneously.

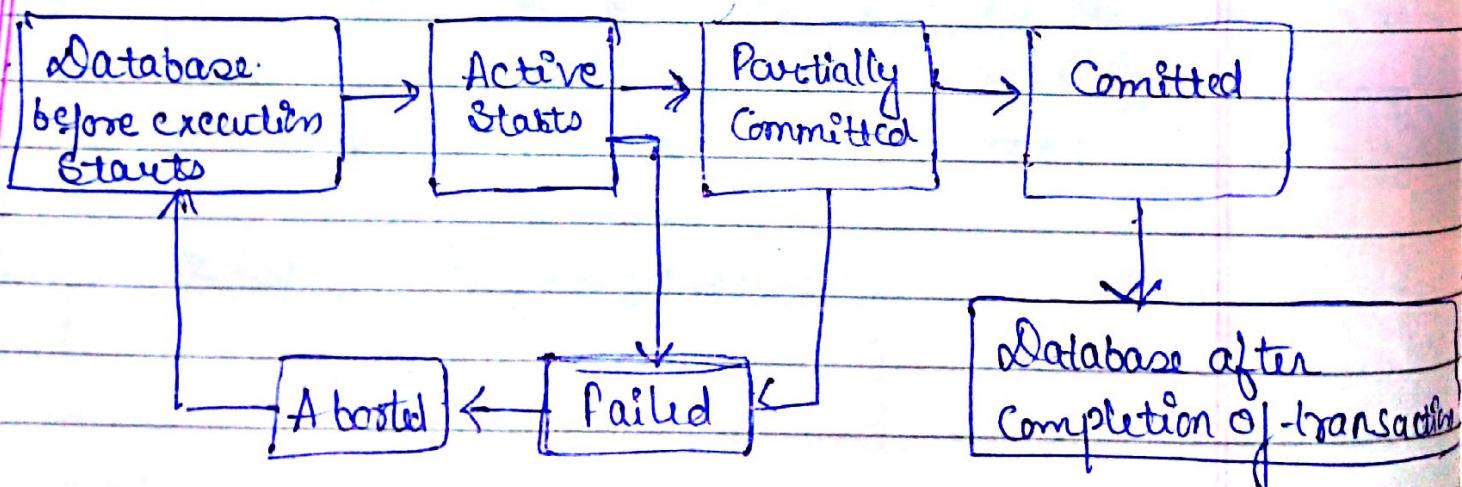
→ Properties of Transactions:

ACID

- ↳ Atomicity
- ↳ Consistency
- ↳ Isolation
- ↳ Durability

- Durability: Changes are made permanent to the database after successful completion of the transaction.

→ Transaction States:



- Serializability: Considered the set of transactions $T_1, T_2, T_3 \dots T_i$. S_1 is the state of database after they are concurrently executed and successfully completed and S_2 is the state of database after they are executed in any serial manner and successfully completed. If S_1 and S_2 are same then the database maintains the serializability.
- Concurrent Execution: If more than one transaction are executed at the same time then they are said to be executed Concurrently.
- Recoverability: To maintain the Atomicity of the database, undo effects of any transaction has to be performed in case of failure of that transaction. If undo effects is successful then that database maintains the recoverability. This process is also known as roll back.

→ Construct an ER diagram for Banking System

Construct an ER diagram for Hospital Management System

Hospital management

