



# Welcome to the **Co**Grammar Trees and Heaps

The session will start shortly...

Questions? Drop them in the chat. We'll have dedicated moderators answering questions.



# Coding Interview Workshop Housekeeping

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- The use of disrespectful language is prohibited in the questions, this is a supportive, learning environment for all - please engage accordingly.  
**(Fundamental British Values: Mutual Respect and Tolerance)**
- No question is daft or silly - **ask them!**
- There are **Q&A sessions** midway and at the end of the session, should you wish to ask any follow-up questions. Moderators are going to be answering questions as the session progresses as well.
- If you have any questions outside of this lecture, or that are not answered during this lecture, please do submit these for upcoming Academic Sessions. You can submit these questions here: [Questions](#)

## Coding Interview Workshop Housekeeping cont.

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- For all **non-academic questions**, please submit a query: [www.hyperiondev.com/support](https://www.hyperiondev.com/support)
- Report a **safeguarding** incident: [www.hyperiondev.com/safeguardreporting](https://www.hyperiondev.com/safeguardreporting)
- We would love your **feedback** on lectures: [Feedback on Lectures](#)

# Skills Bootcamp

## 8-Week Progression Overview

### Fulfil 4 Criteria to Graduation

#### ✓ Criterion 1: Initial Requirements

Timeframe: First 2 Weeks

Guided Learning Hours (GLH):

Minimum of 15 hours

Task Completion: First four tasks

**Due Date: 24 March 2024**

#### ✓ Criterion 2: Mid-Course Progress

**60** Guided Learning Hours

Data Science - **13 tasks**

Software Engineering - **13 tasks**

Web Development - **13 tasks**

**Due Date: 28 April 2024**

# Skills Bootcamp Progression Overview

## ✓ Criterion 3: Course Progress

Completion: All mandatory tasks,  
including Build Your Brand and  
resubmissions by study period end  
Interview Invitation: Within 4 weeks  
post-course  
Guided Learning Hours: Minimum of  
112 hours by support end date  
(10.5 hours average, each week)

## ✓ Criterion 4: Demonstrating Employability

Final Job or Apprenticeship  
Outcome: Document within 12  
weeks post-graduation  
Relevance: Progression to  
employment or related  
opportunity

# Lecture Objectives

- Identify components and properties of trees
- Differentiate between various types of trees
- Implement tree operations in Python and JavaScript
- Construct and manipulate heaps
- Analyse and apply tree and heap concepts to solve problems



# Introduction

- ❖ Trees are fundamental data structures in computer science
- ❖ They are used to represent hierarchical relationships between data elements
- ❖ Understanding trees is crucial for solving complex problems and optimizing algorithms

# Trees and Heaps

1. Tree Fundamentals
2. Binary Trees
3. Binary Search Trees (BSTs)
4. AVL Trees
5. Heap Trees
6. Tree Traversal





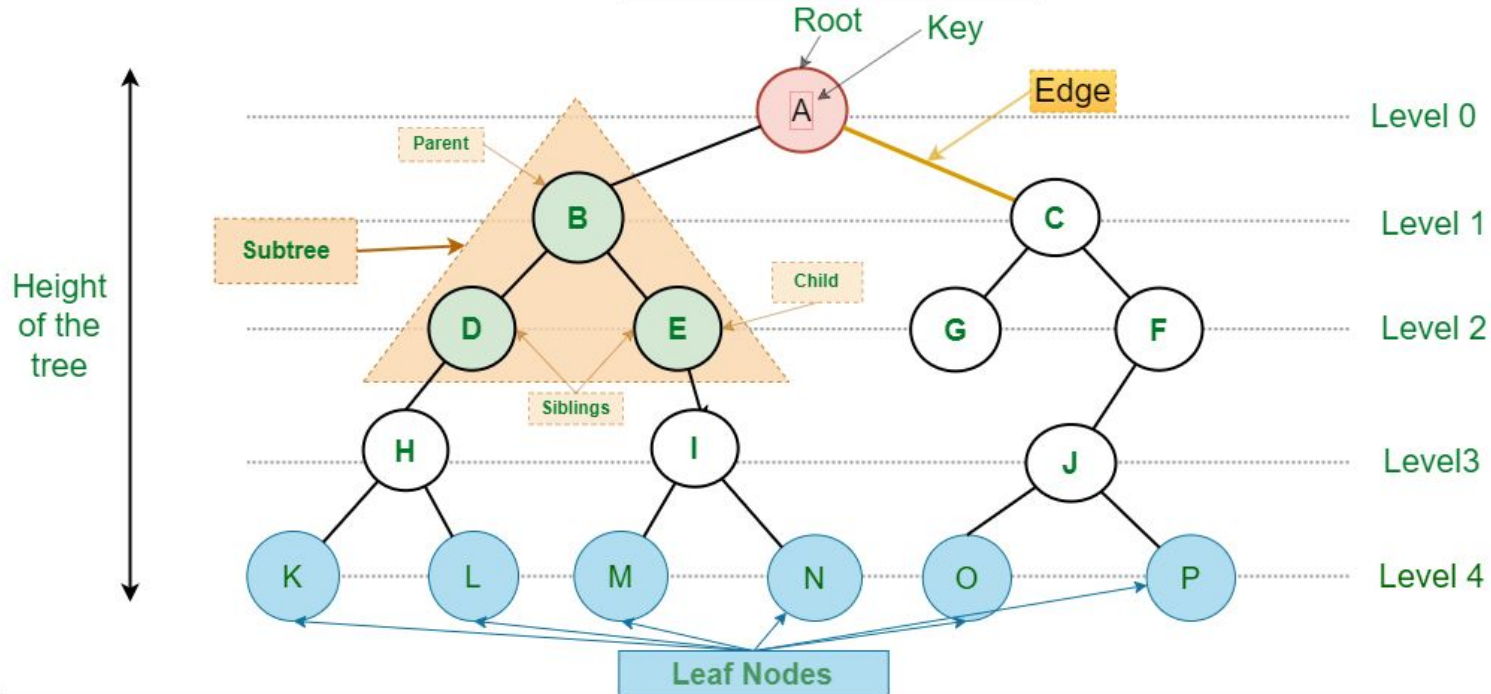


# Tree Fundamentals

**A tree is a non-linear data structure consisting of nodes connected by edges**

- ❖ Components of a tree:
  - **Nodes:** Data elements in the tree
  - **Edges:** Connections between nodes
  - **Root:** The topmost node in the tree
  - **Leaves:** Nodes without child nodes
  - **Depth:** Number of edges from the root to a node
  - **Height:** Maximum depth of the tree

## Tree Data Structure



Source: <https://www.geeksforgeeks.org/introduction-to-tree-data-structure-and-algorithm-tutorials/>



# Tree Fundamentals


- ❖ Properties of trees:
  - Each node (except the root) has **exactly one parent node**
  - The tree is **acyclic (no cycles)**
  - There is a **unique path from the root to each node**
- ❖ Trees are **recursive data structures**:
  - Each node can be treated as the root of a subtree



# Binary Trees

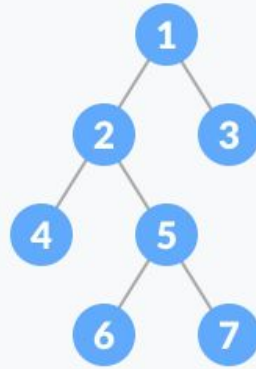
A **binary tree** is a tree in which each node has at most two child nodes (left and right)

❖ Properties:

- There is **no specific ordering or relationship between the values of the nodes.**
  - The **left and right subtrees of a node can have any structure and are not necessarily balanced.**
- 

## 1. Full Binary Tree

A full Binary tree is a special type of binary tree in which every parent node/internal node has either two or no children.

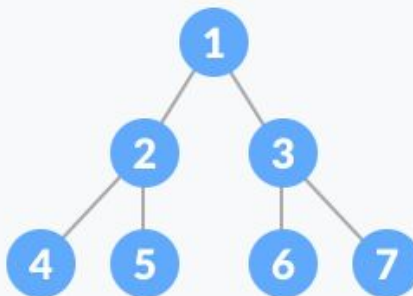


Full Binary Tree

Source: <https://www.programiz.com/dsa/binary-tree>

## 2. Perfect Binary Tree

A perfect binary tree is a type of binary tree in which every internal node has exactly two child nodes and all the leaf nodes are at the same level.



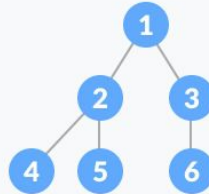
Perfect Binary Tree

Source: <https://www.programiz.com/dsa/binary-tree>

### 3. Complete Binary Tree

A complete binary tree is just like a full binary tree, but with two major differences

1. Every level must be completely filled
2. All the leaf elements must lean towards the left.
3. The last leaf element might not have a right sibling i.e. a complete binary tree doesn't have to be a full binary tree.



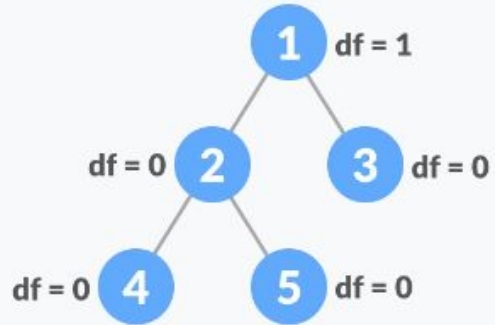
Complete Binary Tree

Source: <https://www.programiz.com/dsa/binary-tree>



## 6. Balanced Binary Tree

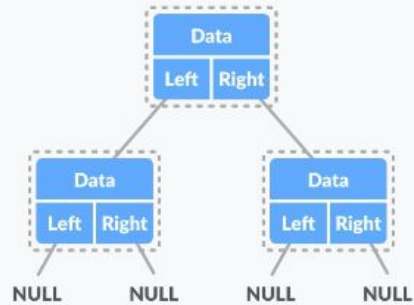
It is a type of binary tree in which the difference between the height of the left and the right subtree for each node is either 0 or 1.



Balanced Binary Tree

Source: <https://www.programiz.com/dsa/binary-tree>





Binary Tree Representation

Source: <https://www.programiz.com/dsa/binary-tree>

*# Binary Tree Node*

*class* *TreeNode*:

```
def __init__(self, val=0, left=None, right=None):  
    self.val = val  
    self.left = left  
    self.right = right
```

*// Binary Tree Node*


*class* *TreeNode* {

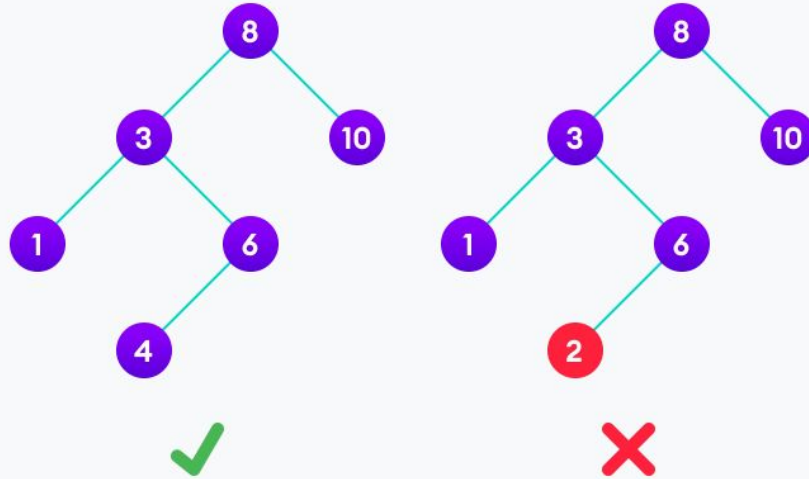
```
    constructor(val = 0, left = null, right = null) {  
        this.val = val;  
        this.left = left;  
        this.right = right;  
    }  
}
```



# Binary Search Trees



- ❖ A binary search tree (BST) is a binary tree with the following properties:
    - The **left subtree of a node contains only nodes with keys less than the node's key**
    - The **right subtree of a node contains only nodes with keys greater than the node's key**
    - Both the left and right subtrees must also be binary search trees
- 



A tree having a right subtree with one value smaller than the root is shown to demonstrate that it is not a valid binary search tree

Source: <https://www.programiz.com/dsa/binary-search-tree>



# Binary Search Trees

- ❖ Operations:
  - **Insertion:** Adding a new node to the tree while maintaining the BST properties
  - **Deletion:** Removing a node from the tree while maintaining the BST properties
  - **Search:** Finding a node with a specific key in the tree



```
# Binary Search Tree Insertion
```

```
def insert(root, val):
```

```
    if not root:
```

```
        return TreeNode(val)
```

```
    if val < root.val:
```

```
        root.left = insert(root.left, val)
```

```
    else:
```

```
        root.right = insert(root.right, val)
```

```
    return root
```

```
// Binary Search Tree Insertion
```

```
function insert(root, val) {
```

```
    if (!root) {
```

```
        return new TreeNode(val);
```

```
    }
```

```
    if (val < root.val) {
```

```
        root.left = insert(root.left, val);
```

```
    } else {
```

```
        root.right = insert(root.right, val);
```

```
    }
```

```
    return root;
```

```
}
```



# Binary Search Trees

- ❖ Time complexity:
  - **Insertion, deletion, and search:  $O(h)$** , where  $h$  is the height of the tree
  - In a balanced BST,  **$h = O(\log n)$** , where  $n$  is the number of nodes



# AVL Trees

**An AVL tree is a self-balancing binary search tree**

- ❖ Balancing property:
  - The heights of the **left and right subtrees of any node differ by at most one**
  - This property ensures that the tree **remains balanced**, preventing degeneration into a linked list



# AVL Trees

- ❖ Rotations:
  - Used to **rebalance the tree when the AVL property is violated after an insertion or deletion**
  - Left rotation and right rotation
- ❖ Applications:
  - Efficient searching and sorting (since the tree is balanced)

# AVL Trees

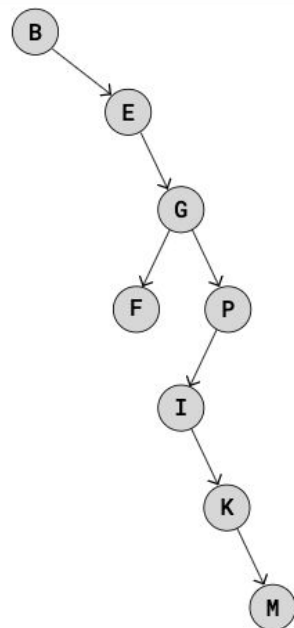
- ❖ The code for AVL trees could become quite long due to the balancing property, so feel free to check it out

- Python:

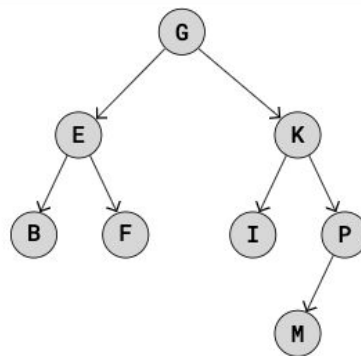
[https://github.com/bfaure/Python3\\_Data\\_Structures/blob/master/AVL\\_Tree/main.py](https://github.com/bfaure/Python3_Data_Structures/blob/master/AVL_Tree/main.py)

- Javascript:

<https://github.com/gwtw/js-avl-tree/tree/master/src>

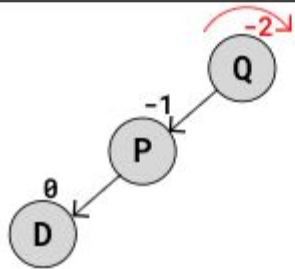


*Binary Search Tree  
(unbalanced)  
Height: 6*

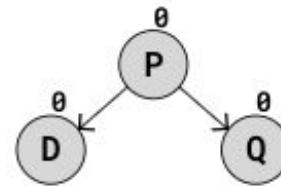


*AVL Tree  
(self-balancing)  
Height: 3*

Source: [https://www.w3schools.com/dsa/dsa\\_data\\_avltrees.php](https://www.w3schools.com/dsa/dsa_data_avltrees.php)

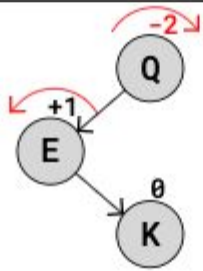


Rotate Right

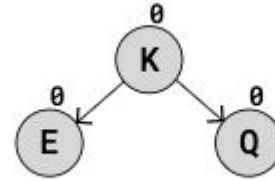


Insert L

Source: [https://www.w3schools.com/dsa/dsa\\_data\\_avltrees.php](https://www.w3schools.com/dsa/dsa_data_avltrees.php)

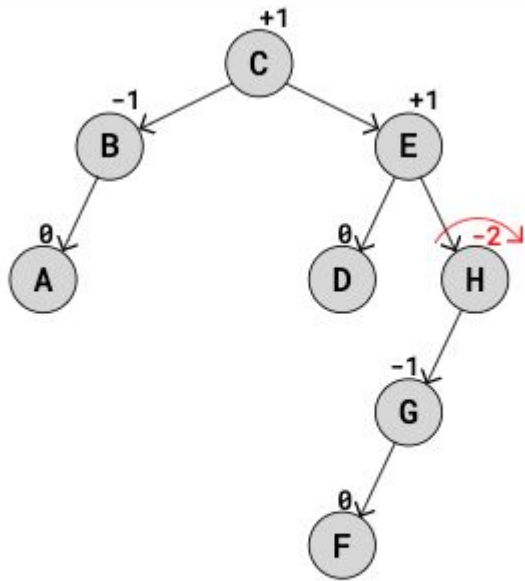


Rotate Left-Right

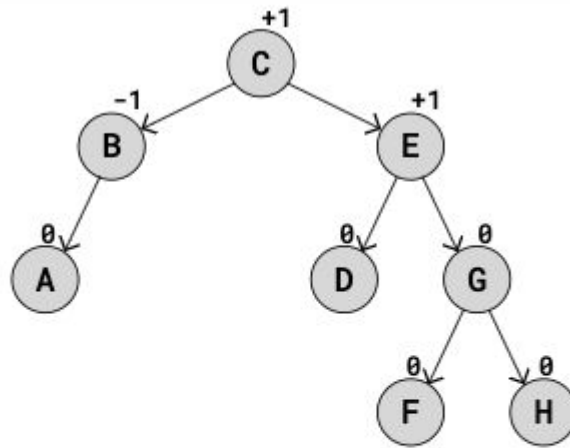


Insert C

Source: [https://www.w3schools.com/dsa/dsa\\_data\\_avltrees.php](https://www.w3schools.com/dsa/dsa_data_avltrees.php)



Rotate Right



Reset

Source: [https://www.w3schools.com/dsa/dsa\\_data\\_avltrees.php](https://www.w3schools.com/dsa/dsa_data_avltrees.php)

# Heap

A heap is a complete binary tree that satisfies the heap property

- ❖ Heap property:
  - For a **min-heap**:  $\text{parent}(i) \leq i$
  - For a **max-heap**:  $\text{parent}(i) \geq i$

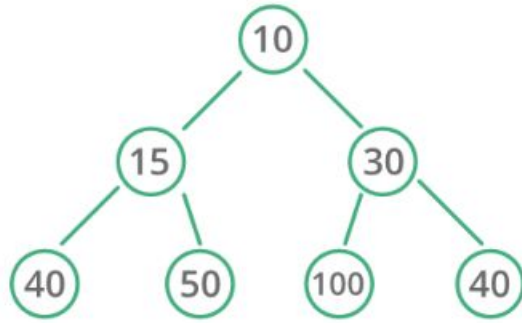


# Heap

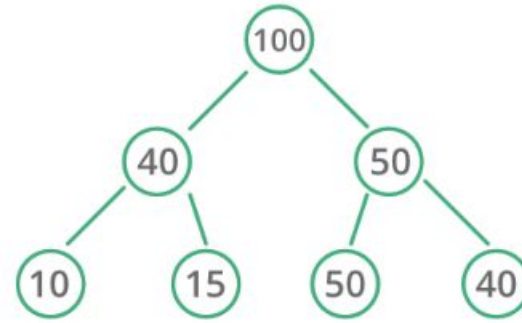
- ❖ Applications:
  - Priority queues
  - Heapsort algorithm



# Heap Data Structure



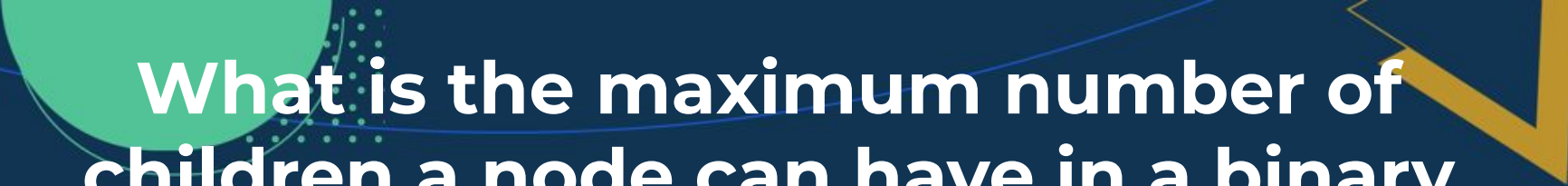
Min Heap



Max Heap

GG

Source: <https://www.geeksforgeeks.org/heap-data-structure/>



# What is the maximum number of children a node can have in a binary tree?

- A. 1
- B. 2
- C. 3
- D. 4



## Answer: B

- ❖ In a binary tree, each node can have at most two children, referred to as the left child and the right child.

# What is the time complexity of searching for a value in a balanced binary search tree?

- A.  $O(1)$
- B.  $O(\log n)$
- C.  $O(n)$
- D.  $O(n^2)$

## Answer: B

- ❖ In a balanced BST, the search operation has a time complexity of  $O(\log n)$  because the tree's height is maintained as  $\log n$ , and the search algorithm eliminates half of the remaining nodes at each step.

# Which of the following is true about a min-heap?

- A. The root node has the minimum value in the heap
- B. The root node has the maximum value in the heap
- C. The heap property is:  $\text{parent}(i) \geq i$
- D. The heap property is:  $\text{parent}(i) > i$

## Answer: A

- ❖ In a min-heap, the heap property ensures that the value of each node is greater than or equal to the value of its parent node. Consequently, the root node always contains the minimum value in the heap.

# Let's Breathe!

Let's take a small break  
before moving on to  
the next topic.





# Tree Traversal

- ❖ Tree traversal is the process of **visiting each node in a tree exactly once**
- ❖ Types of tree traversal:
  - **In-order traversal (Depth-first Search)**
  - **Pre-order traversal (Depth-first Search)**
  - **Post-order traversal (Depth-first Search)**
  - **Level-order traversal (Breadth-first Search)**

# Tree Traversal

- ❖ **In-order traversal** visits nodes in the following order:
  - **Left subtree**
  - **Root**
  - **Right subtree**
- ❖ Useful for:
  - Visiting nodes in sorted order (for BSTs)
  - Creating a copy of the tree




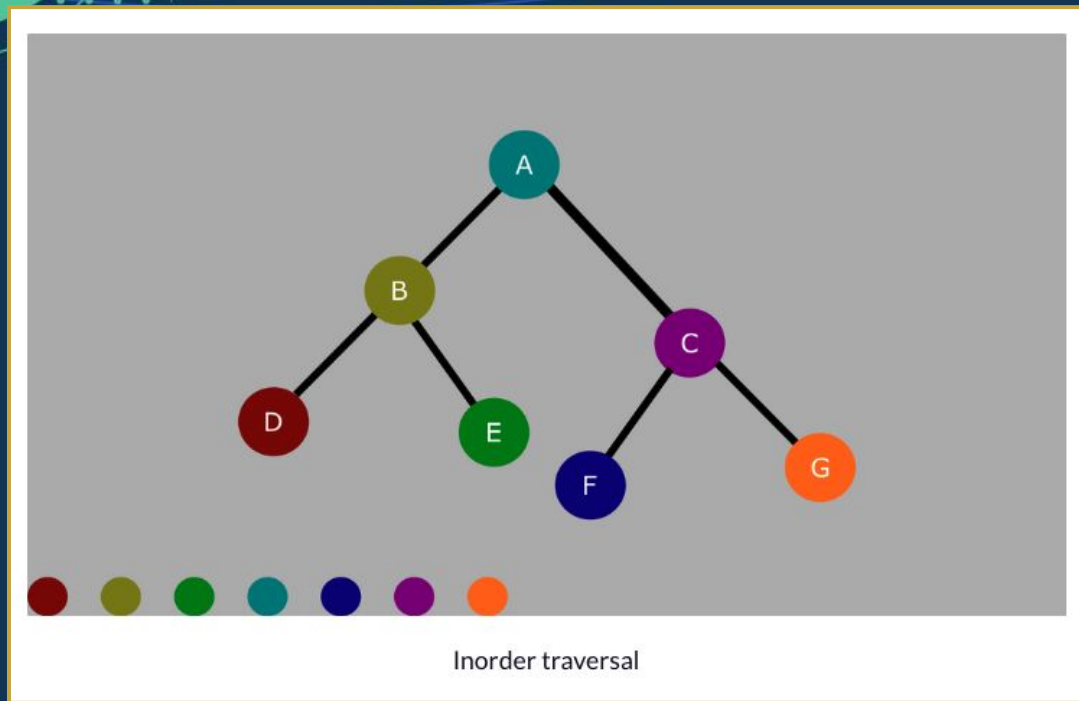
```
# In-order Traversal
```

```
def inorder_traversal(root):  
    if root:  
        inorder_traversal(root.left)  
        print(root.val)  
        inorder_traversal(root.right)
```

```
// In-order Traversal
```

```
function inorderTraversal(root) {  
    if (root) {  
        inorderTraversal(root.left);  
        console.log(root.val);  
        inorderTraversal(root.right);  
    }  
}
```

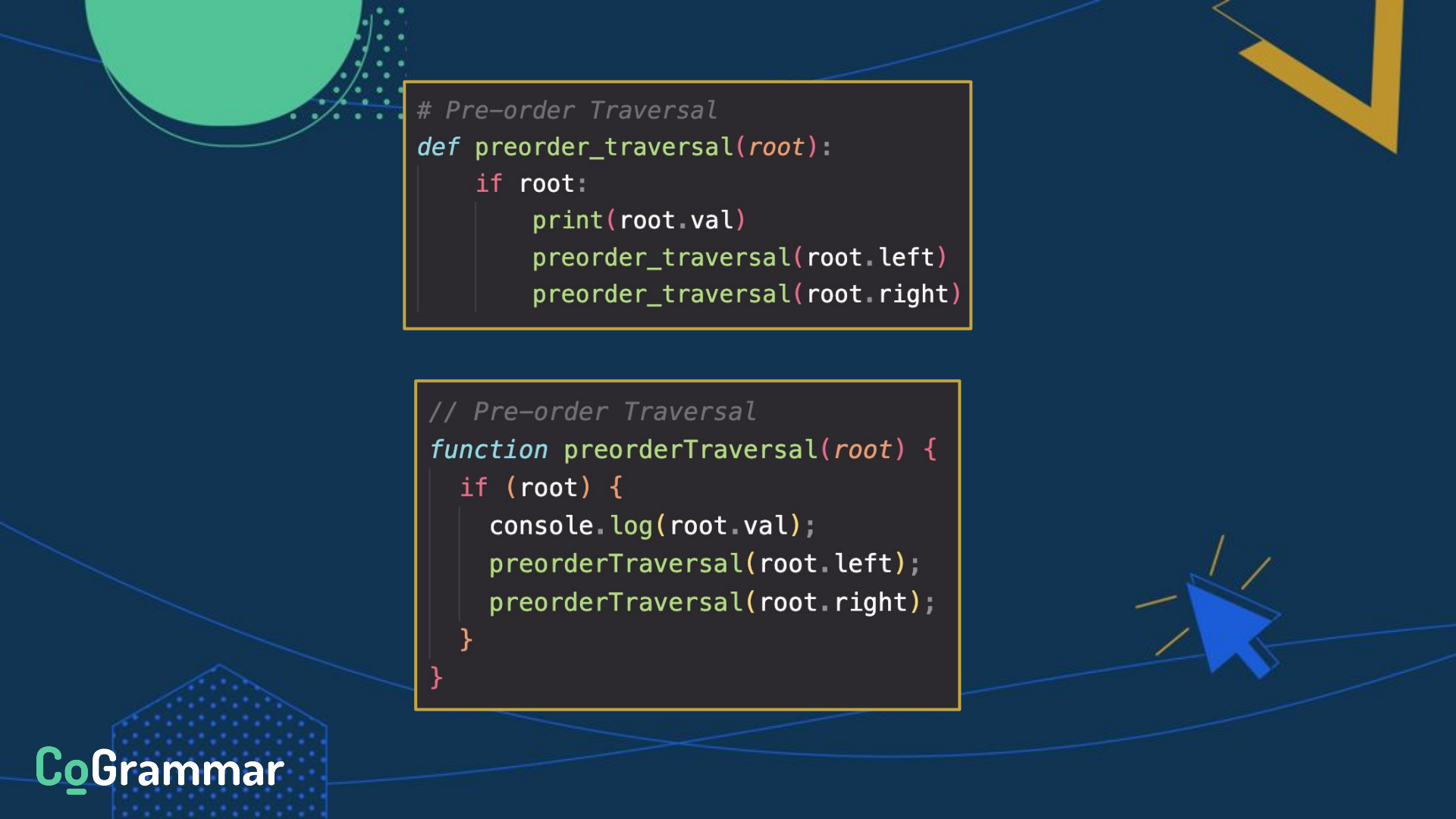




Source: <https://www.freecodecamp.org/news/binary-search-tree-traversal-inorder-preorder-post-order-for-bst/>

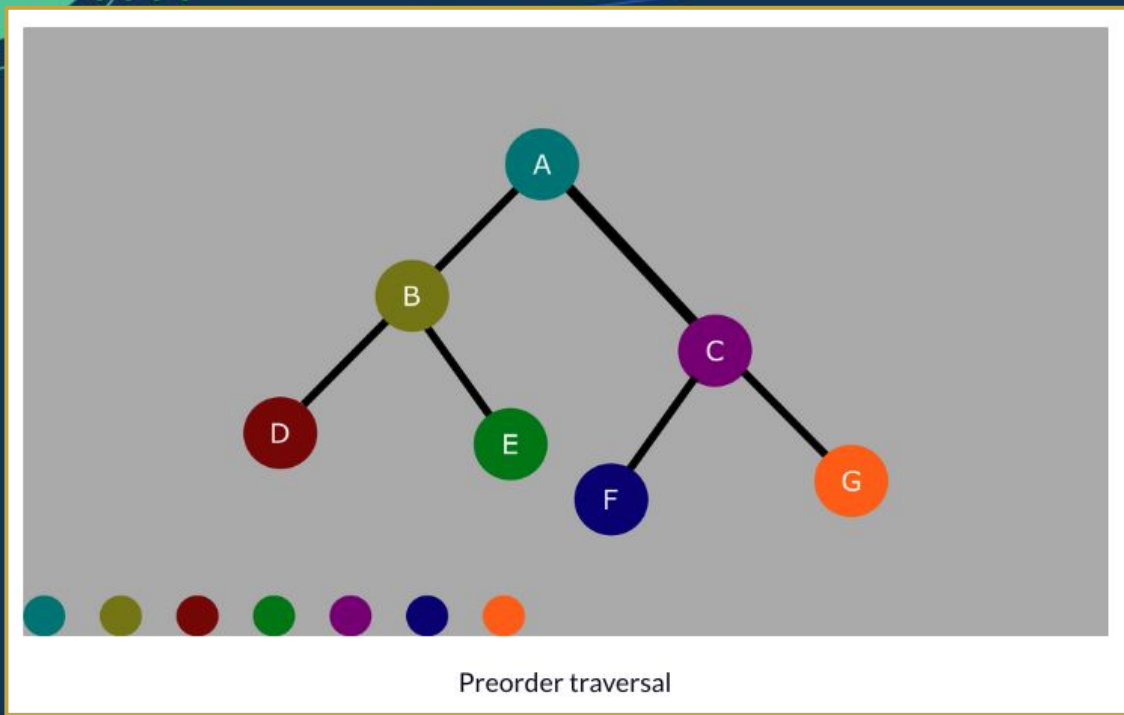
# Tree Traversal

- ❖ **Pre-order traversal** visits nodes in the following order:
  - **Root**
  - **Left subtree**
  - **Right subtree**
- ❖ Useful for:
  - Creating a copy of the tree
  - Prefix expression evaluation



```
# Pre-order Traversal
def preorder_traversal(root):
    if root:
        print(root.val)
        preorder_traversal(root.left)
        preorder_traversal(root.right)
```

```
// Pre-order Traversal
function preorderTraversal(root) {
    if (root) {
        console.log(root.val);
        preorderTraversal(root.left);
        preorderTraversal(root.right);
    }
}
```



Source: <https://www.freecodecamp.org/news/binary-search-tree-traversal-inorder-preorder-post-order-for-bst/>



# Tree Traversal

- ❖ **Post-order traversal** visits nodes in the following order:
  - **Left subtree**
  - **Right subtree**
  - **Root**
- ❖ Useful for:
  - Deleting a tree
  - Postfix expression evaluation



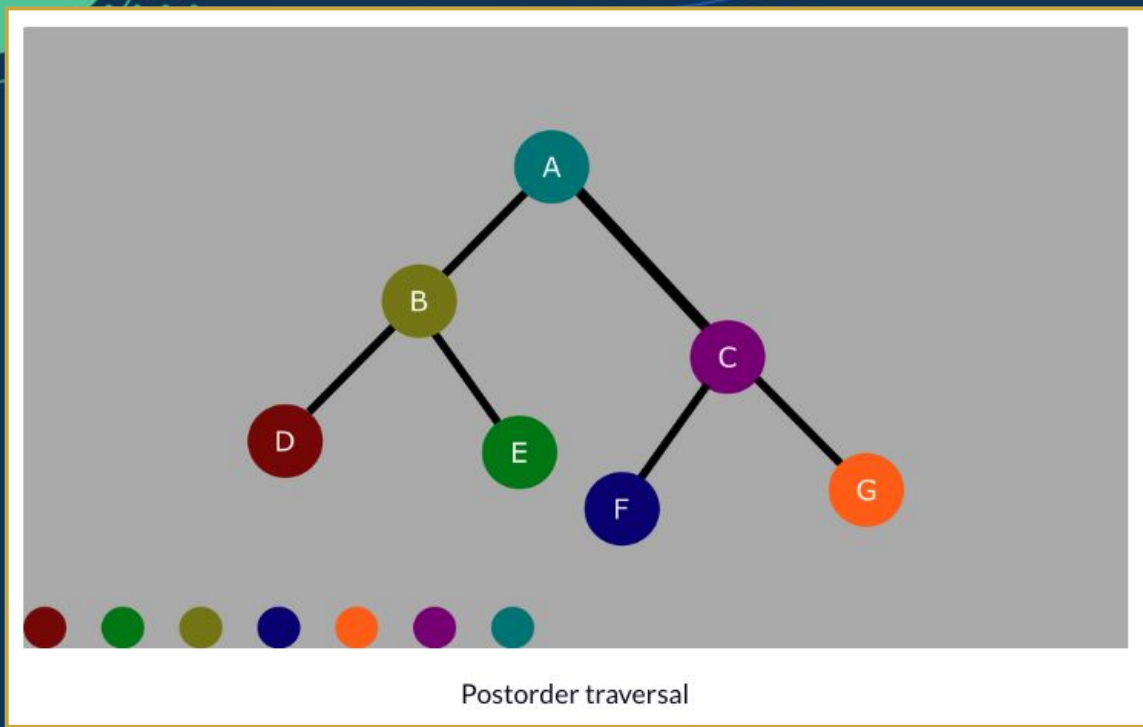


```
# Post-order Traversal
```

```
def postorder_traversal(root):  
    if root:  
        postorder_traversal(root.left)  
        postorder_traversal(root.right)  
        print(root.val)
```

```
// Post-order Traversal
```

```
function postorderTraversal(root) {  
    if (root) {  
        postorderTraversal(root.left);  
        postorderTraversal(root.right);  
        console.log(root.val);  
    }  
}
```



Source: <https://www.freecodecamp.org/news/binary-search-tree-traversal-inorder-preorder-post-order-for-bst/>



# Tree Traversal

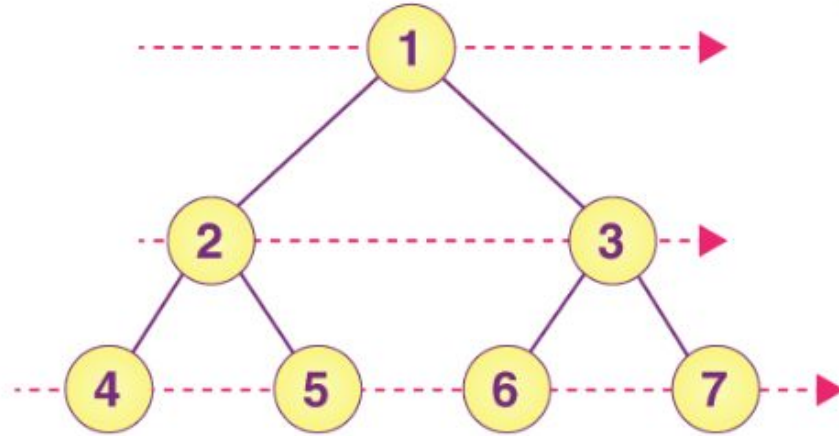
- ❖ **Level-order traversal visits nodes level by level, from left to right**
- ❖ Useful for:
  - Printing the tree level by level
  - Finding the shortest path between two nodes



```
# Level-order Traversal
from collections import deque

def level_order_traversal(root):
    if not root:
        return
    queue = deque([root])
    while queue:
        node = queue.popleft()
        print(node.val)
        if node.left:
            queue.append(node.left)
        if node.right:
            queue.append(node.right)
```

```
// Level-order Traversal
function levelOrderTraversal(root) {
    if (!root) {
        return;
    }
    const queue = [root];
    while (queue.length) {
        const node = queue.shift();
        console.log(node.val);
        if (node.left) {
            queue.push(node.left);
        }
        if (node.right) {
            queue.push(node.right);
        }
    }
}
```



In this case, the level order traversal will give output in the form of the following order: 1, 2, 3, 4, 5, 6, 7.

Source: <https://byjus.com/gate/tree-traversal-notes/>

# Heap Construction

- ❖ To construct a heap from an array of elements:
  - Insert each element into the heap one by one
  - After each insertion, sift up the new element to maintain the heap property
- ❖ This process is called **heapify**
- ❖ **Time complexity:  $O(n \log n)$**  for binary heap since  $n/2$  non-leaf nodes require the heapify operation in worst case and heapify is  $O(\log n)$ , so  $n/2 * O(\log n) = O(n/2 \log n) = O(n \log n)$



```
# Heapify
def heapify(arr):
    heap = MinHeap()
    for val in arr:
        heap.push(val)
    return heap
```



```
// Heapify
function heapify(arr) {
    const heap = new MinHeap();
    for (const val of arr) {
        heap.push(val);
    }
    return heap;
}
```



# Heap Operations

## ❖ Insertion:

- Add a new element to the end of the heap
- Sift up the new element to maintain the heap property
- **Time complexity:  $O(\log n)$**

```
def push(self, val):  
    self.heap.append(val)  
    self._sift_up(len(self.heap) - 1)
```

```
push(val) {  
    this.heap.push(val);  
    this._siftUp(this.heap.length - 1);  
}
```



# Heap Operations

- ❖ **Deletion** (extracting the minimum or maximum element):
  - Replace the root with the last element in the heap
  - Remove the last element
  - Sift down the new root to maintain the heap property
  - Time complexity:  **$O(\log n)$**

```
def pop(self):  
    if not self.heap:  
        return None  
    self.swap(0, len(self.heap) - 1)  
    val = self.heap.pop()  
    self._sift_down(0)  
    return val
```

```
pop() {  
    if (!this.heap.length) {  
        return null;  
    }  
    this.swap(0, this.heap.length - 1);  
    const val = this.heap.pop();  
    this._siftDown(0);  
    return val;  
}
```



## Interview-Style Questions: Social Network

On a social platform like Facebook, you can connect with friends, and each of your friends can also have their own friends. If we consider that both you and each of your friends can have at most 2 friends in total, how many levels of connections beyond your direct friends can you explore?



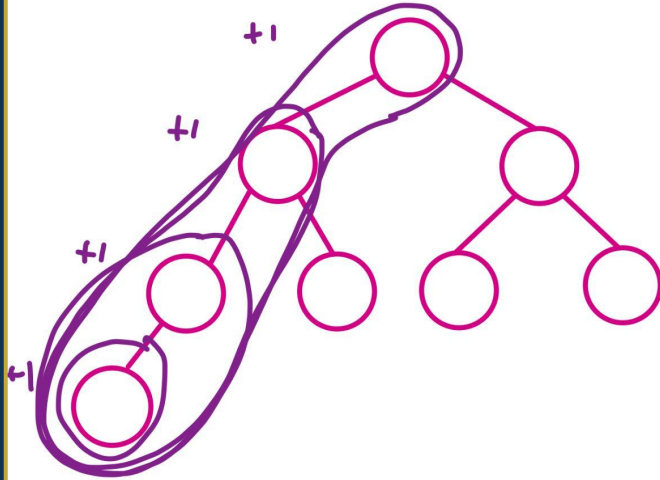


# Interview-Style Questions

- ❖ Find the height of a binary tree
- ❖ Check if a binary tree is a BST
- ❖ Find the kth smallest element in a BST

# Interview-Style Question 1

- ❖ Given the root of a binary tree, find its height (maximum depth)



check height recursively  
for left & right subtrees.

4 depth we get from  $1 + \max(\text{left height, right height})$

*# Find the height of a binary tree*

```
def height(root):  
    if not root:  
        return 0  
    return 1 + max(height(root.left), height(root.right))
```

*// Find the height of a binary tree*

```
function height(root) {  
    if (!root) {  
        return 0;  
    }  
    return 1 + Math.max(height(root.left), height(root.right));  
}
```



## Interview-Style Question 2

- ❖ Given the root of a binary tree, determine if it is a valid binary search tree (BST)
  - Basically, are all left nodes less in value than the right nodes
  - We can just check this for the left and right subtrees recursively



*# Check if a binary tree is a BST*

```
def is_bst(root, min_val=float('-inf'), max_val=float('inf')):  
    if not root:  
        return True  
    if root.val <= min_val or root.val >= max_val:  
        return False  
    return is_bst(root.left, min_val, root.val) and is_bst(root.right, root.val, max_val)
```

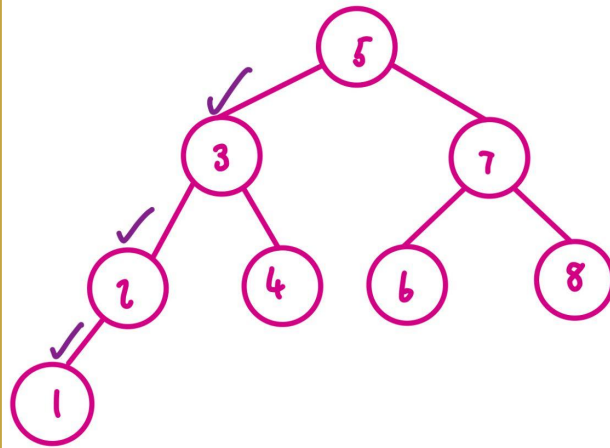
*// Check if a binary tree is a BST*

```
function isBST(root, minVal = -Infinity, maxVal = Infinity) {  
    if (!root) {  
        return true;  
    }  
    if (root.val <= minVal || root.val >= maxVal) {  
        return false;  
    }  
    return (  
        isBST(root.left, minVal, root.val) && isBST(root.right, root.val, maxVal)  
    );  
}
```



## Interview-Style Question 3

- ❖ Given the root of a binary search tree and an integer  $k$ , find the  $k$ th smallest element in the BST



Use DFS

eg. 3rd smallest ( $k=3$ )

stack = [5, 3, 2, 1]

case: 1

$k=2$

case: 2

$k=1$

case: 3

$k=0$

return 3 correct



*# Find the kth smallest element in a BST*

```
def kth_smallest(root, k):  
    stack = []  
    while True:  
        while root:  
            stack.append(root)  
            root = root.left  
        root = stack.pop()  
        k -= 1  
        if not k:  
            return root.val  
        root = root.right
```

*// Find the kth smallest element in a BST*

```
function kthSmallest(root, k) {  
    const stack = [];  
    while (true) {  
        while (root) {  
            stack.push(root);  
            root = root.left;  
        }  
        root = stack.pop();  
        k--;  
        if (!k) {  
            return root.val;  
        }  
        root = root.right;  
    }  
}
```

# Which of the following is the correct order of nodes visited in a post-order traversal of a binary tree?

- A. Root, Left, Right
- B. Left, Right, Root
- C. Right, Left, Root
- D. Left, Root, Right

## Answer: B

- ❖ In a post-order traversal, the left subtree is visited first, then the right subtree, and finally the root node.

## Bonus Interview question

Let's assume a that you would like to insert data quickly and retrieve it fast. What data structure would you use and why?

## Bonus Interview question

What is the complexity of such data structure?

## Bonus Interview question

Implement that in your language of choice.

## Bonus Interview question

Let's add in this order: 2,1,3,4,5,6

## Bonus Interview question

Is it still efficient?



## Bonus Interview question

What would you do to optimise that?

# CoGrammar

## Q & A SECTION

**Please use this time to ask  
any questions relating to the  
topic, should you have any.**

# Thank you for attending



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