

$O(N)$
 $O(1)$ "abacbc" \downarrow i whether every character is occ.
Same no. of times
space

How about we calc the no. of occ. of each char.

Because we have
to keep track
of occ, we
can pop
the freq map

Key value
 \downarrow
char \hookrightarrow frequency
freq = 2

{
'a' : 2
'b' : 2
'c' : 2
}

Power function Recursive



Q → Given two integer values 'a' and 'b', calculate a^b recursively.

Ex → $a = 2$ $b = 3$

ans → 8

Iteratively

```
ans = 1 //  
for (i = 1; i <= b; i++) {  
    ans *= a;  
}
```

$$a^b = \underbrace{a \times a \times a \dots a}_{b \text{ times}}$$

$$4^3 = \underbrace{4 \times 4 \times 4}$$

Recursive

$$a^b = a \times a \times a \times a \dots a \quad (b \text{ times})$$

$$a^{b+1} = \underbrace{a \times a \times a \times a \dots a}_{a^b} \times a \quad (b+1 \text{ times})$$

$$a^{b+1} = a^b \times a$$

Similarly

$$a^b = a^{b-1} \times a$$

The most trivial
(Base Case)

$$\rightarrow b \rightarrow 0$$

$$2^4 = 2 \times 2^3$$

my task

$$2^3 = 2 \times 2^2$$

$$2^3 = 2 \times 2^2$$

$$2^2 = 2 \times 2^1$$

$$2^1 = 2 \times 2^0$$

$$a^b \rightarrow 1$$

$f(a, b)$

returns a
raised to power
 b

$$= a \times f(a, b-1)$$

assumption

Self work

no. of calls $\rightarrow b \times c$
 $\rightarrow O(b)$

70%

if $(b == 0)$
return 1; } Base Case

$$a^b = a + a^{b-1}$$
$$= a + (a + a^{b-2})$$

...

(No. of funcⁿ
calls) \times (no. of calls
in each call)

No. of calls

① Base Case → if ($b == 0$)
return 1;

② Assumption → assume $f(a, b-1)$ works correctly
and a^{b-1} is recursively computed

③ Self work → multiply a with a^{b-1}



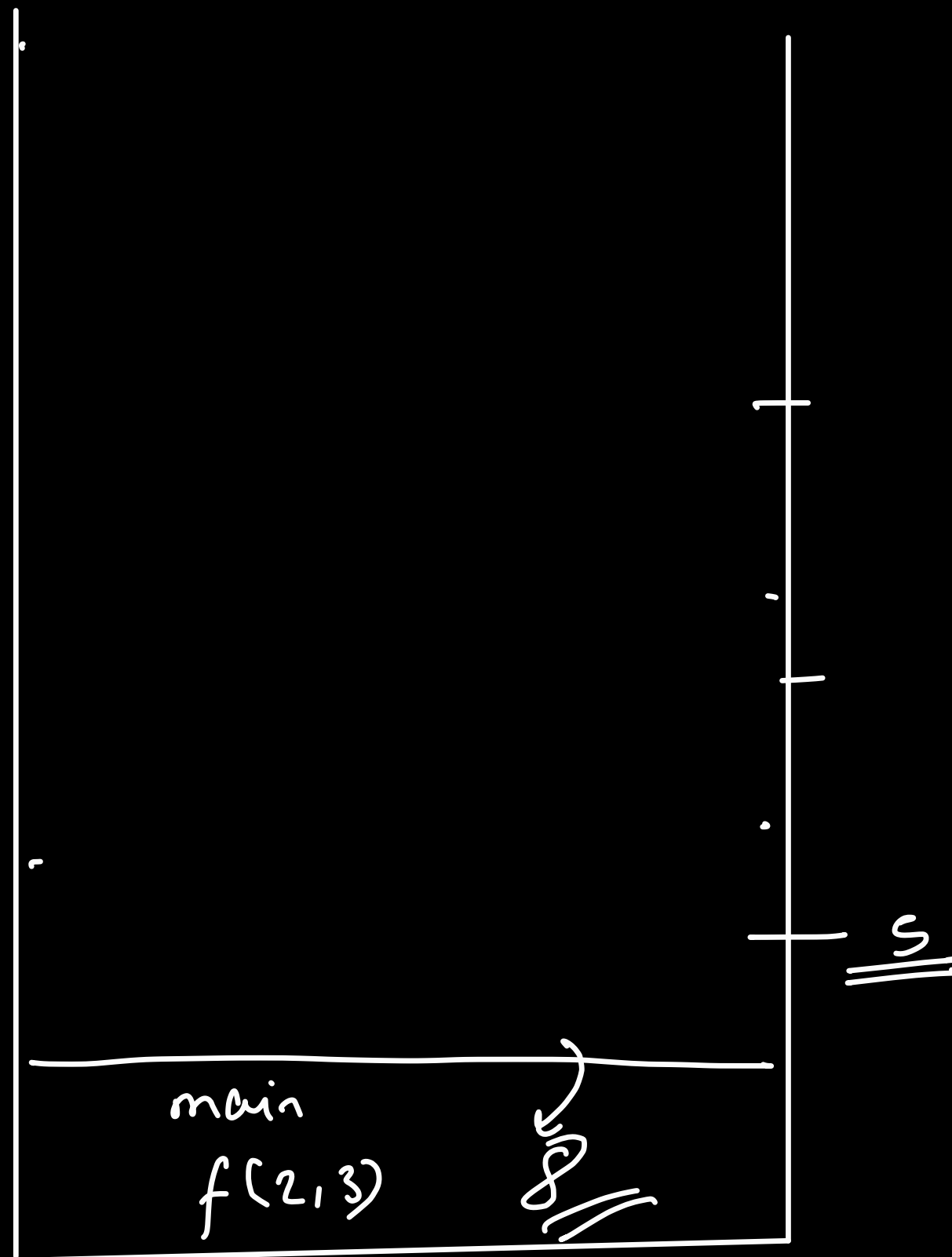
```
1  int f(int a, int b) {  
2      if(b == 0) { // base case  
3          return 1;  
4      }  
5      return a * f(a, b-1);  
6  }  
7
```

main()

f(2,3)

Space $\rightarrow O(b)$

Time $\rightarrow O(b)$



$$2^8 = 2 \times 2^7$$

$$2^8 = 2^4 \times 2^4$$

$$2^9 = 2 \times (2^4 \times 2^4)$$

$$a^b = \begin{cases} a^{b/2} \times a^{b/2} \\ a \times a^{b/2} \times a^{b/2} \end{cases}$$

$b \rightarrow \underline{\text{even}}$

$b \rightarrow \text{odd}$

$$\begin{array}{l}
 f(a, b) \\
 \downarrow \\
 \text{returns} \\
 \underline{\underline{a^b}}
 \end{array}
 = \begin{cases}
 f(a, b/2) \times f(a, b/2) & \text{if } \underline{b \rightarrow \text{even}} \\
 f(a, b/2) \times f(a, b/2) \times a & \text{if } b \rightarrow \text{odd}
 \end{cases}$$

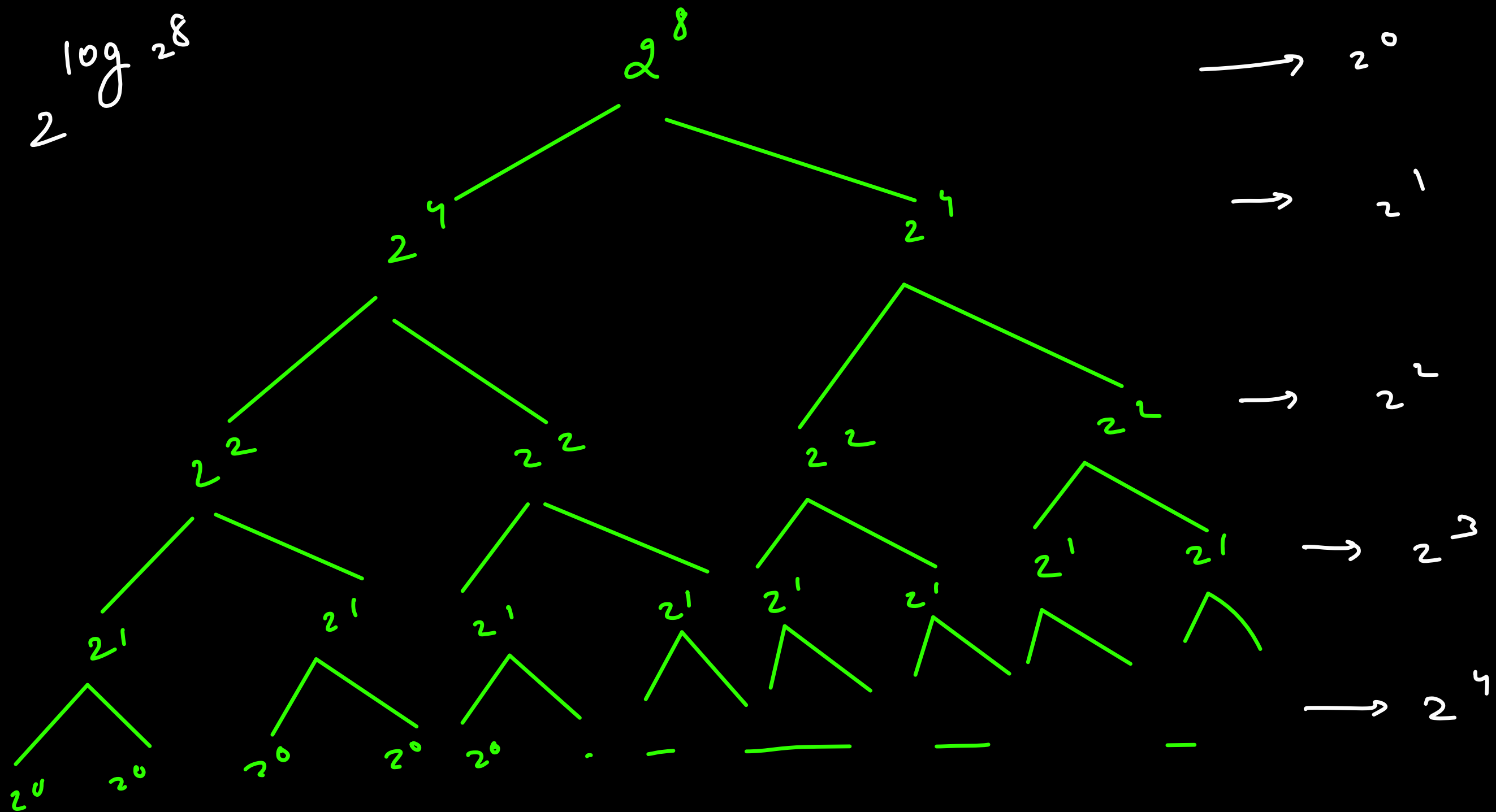
$$(b == 0) \rightarrow \underline{\underline{1}}$$

$$\begin{array}{c}
 2 \\
 \underline{\underline{O(b)}}
 \end{array}$$

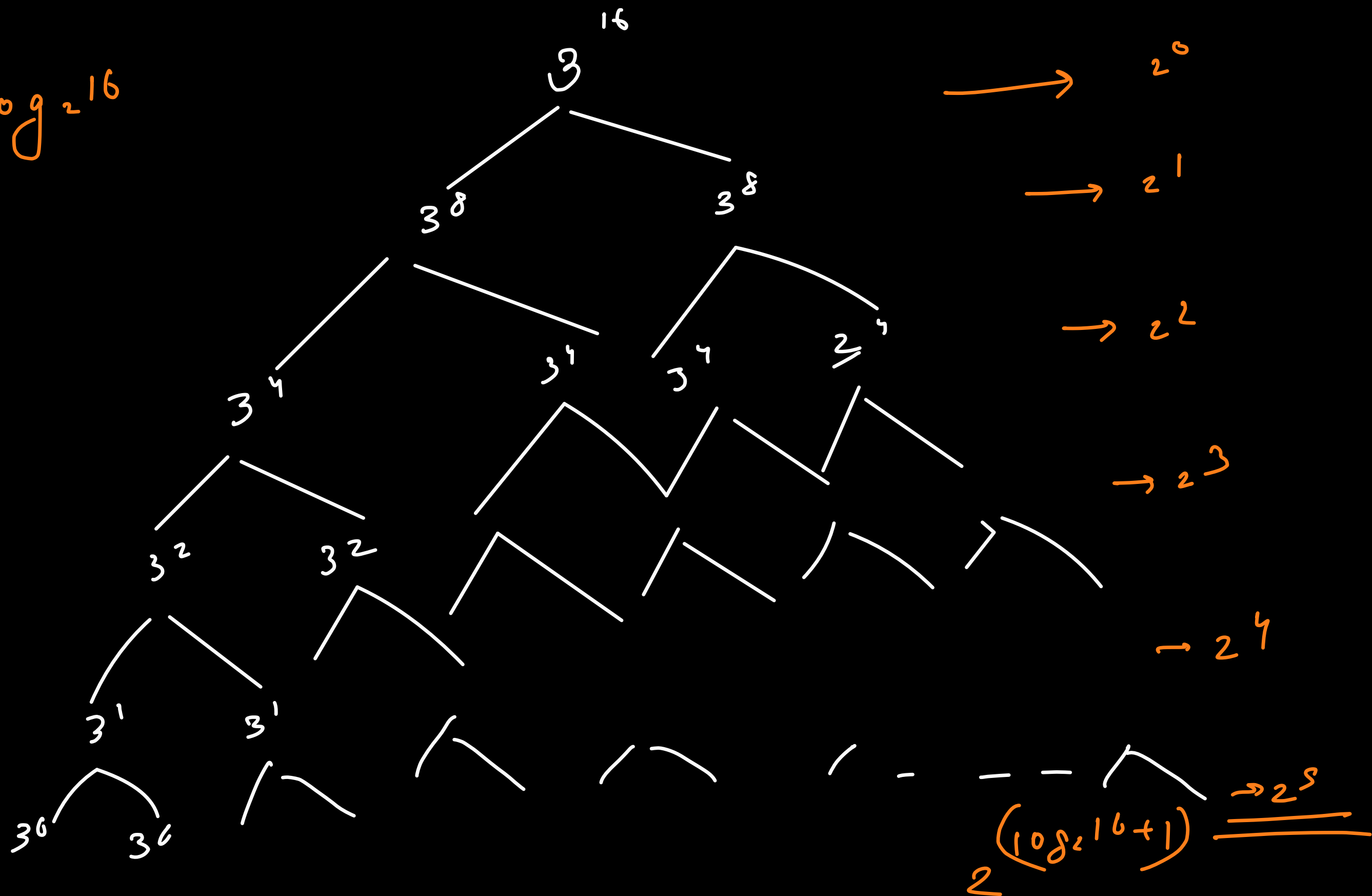
$$\begin{array}{c} 2^8 \\ \downarrow \\ 2^7 \\ \downarrow \\ 2^6 \\ \downarrow \\ 2^5 \\ \downarrow \\ 2^4 \\ \downarrow \\ 2^3 \\ \downarrow \\ 2^2 \\ \downarrow \\ 2^1 \\ \downarrow \\ 2^0 \end{array}$$

JOIN THE DARKSIDE

$$2^{\log 2^8}$$



$$2^{\log_2 16}$$



$f(a, b)$

$$\rightarrow 2^0 + 2^1 + 2^2 \dots \dots \dots 2^{\log_2 b + 1}$$

~~gp~~ $\rightarrow a \quad ar \quad ar^2 \quad ar^3 \dots \dots \dots ar^{n-1}$

$(r > 1)$

$$\frac{ar(r^n - 1)}{r - 1}$$

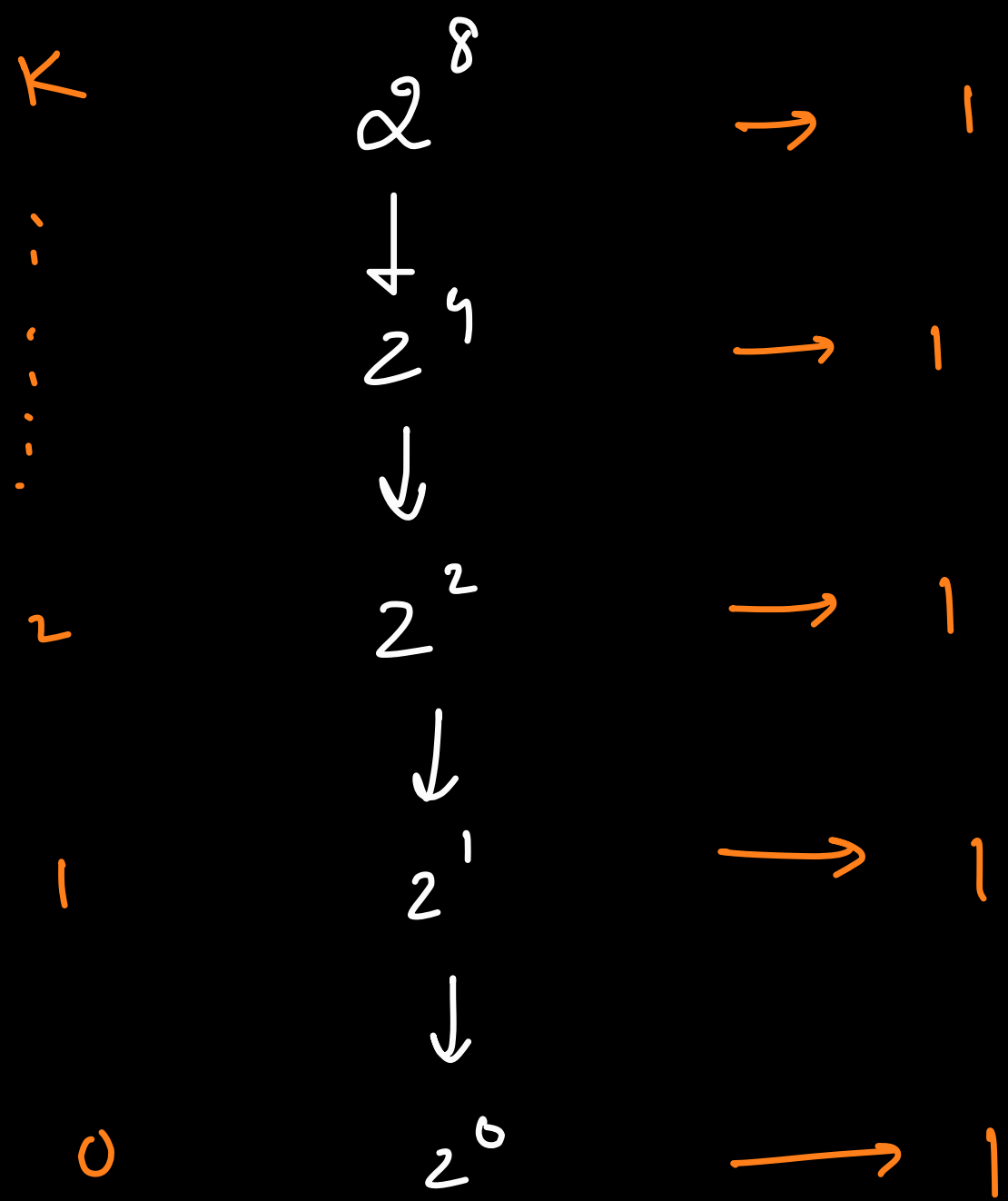
$$\rightarrow \frac{1 \times (2^{\log_2 b + 2} - 1)}{2 - 1} \Rightarrow 2^{\log_2 b + 2} - 1$$
$$\hookrightarrow \approx 2^{\log_2 b} \rightarrow \underline{\underline{b}}$$

$(a^{\log_2 b}) \rightarrow b$

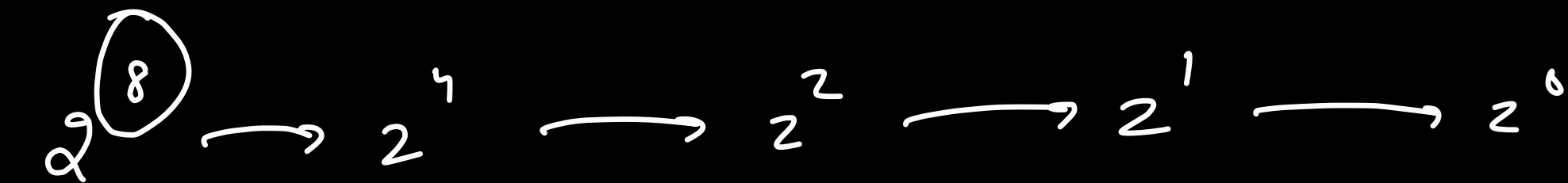
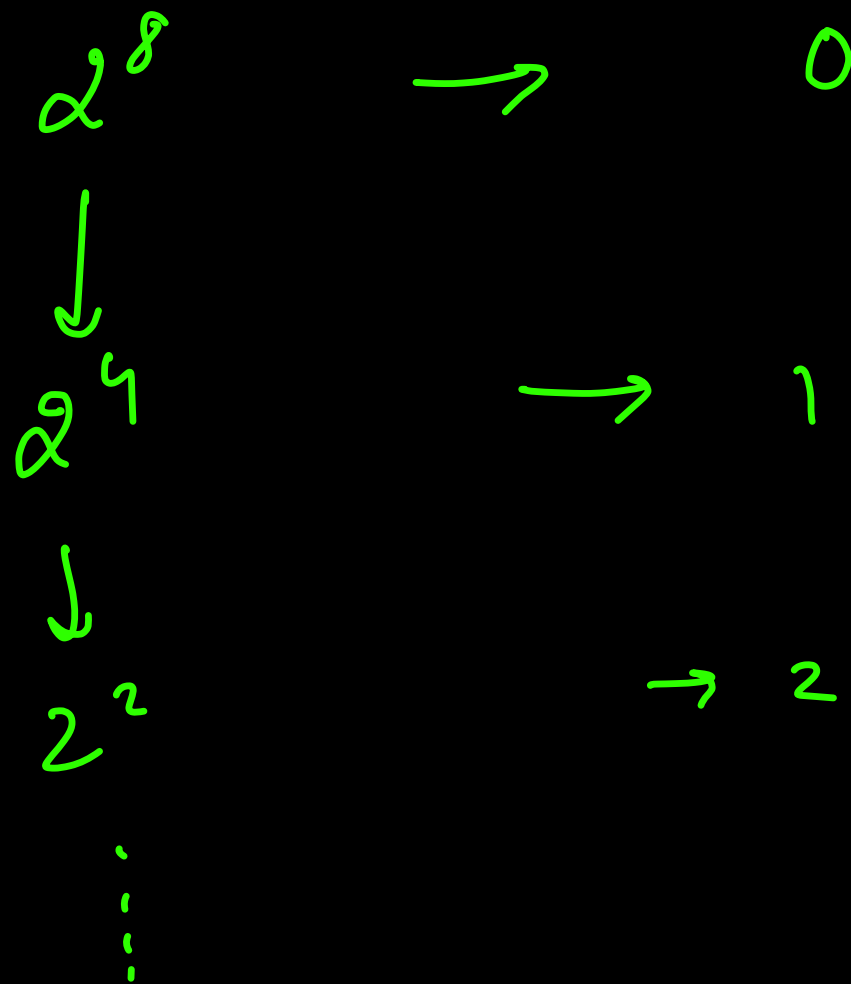
$$2^{\log_2 4b} \rightarrow \frac{\log_2 b + \log_2 4}{\log_2 4b}$$

$$f(a, b) = \begin{cases} f(a, b/2)^2 & \text{if } b \rightarrow \text{even} \\ a \times f(a, b/2)^2 & \text{if } b \rightarrow \text{odd} \end{cases}$$

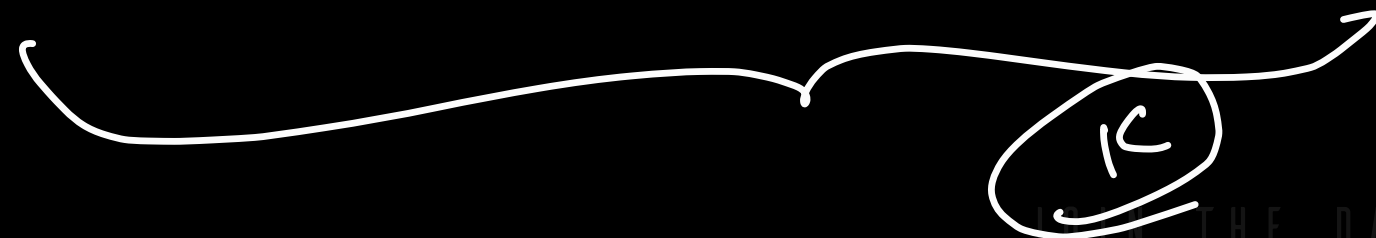
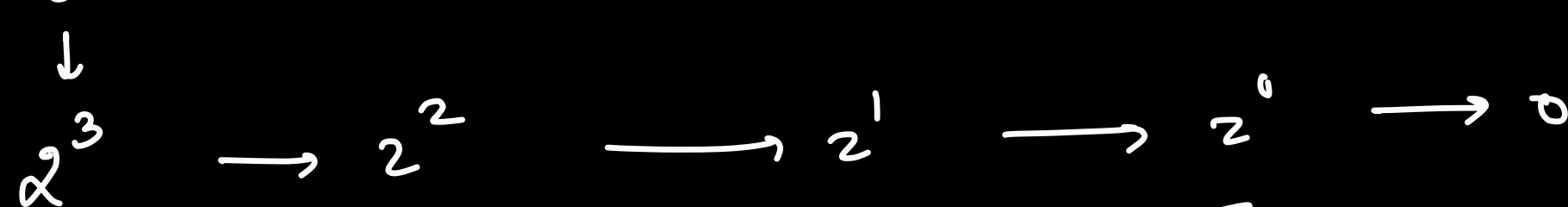
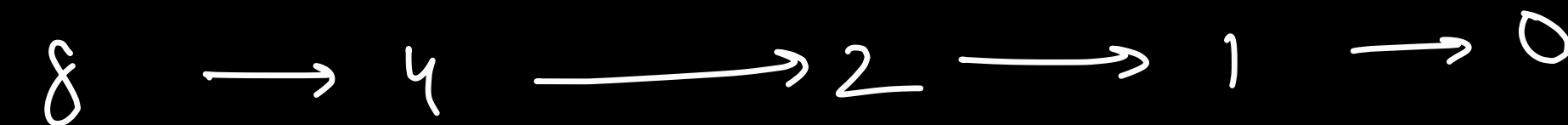
$$\boxed{\uparrow K \approx \log_2 b}$$



$$(1 + 1 + 1 + \dots + 1) \rightarrow \underline{\underline{(\log_2 b)}} \rightarrow \underline{\underline{O(\log b)}}$$



(K steps)



$$2^0 \rightarrow 2^1 \rightarrow 2^2 \dots \dots \dots \underline{\underline{2^k}}$$

$$2^k \approx b$$

$$\log_2 b \approx k$$

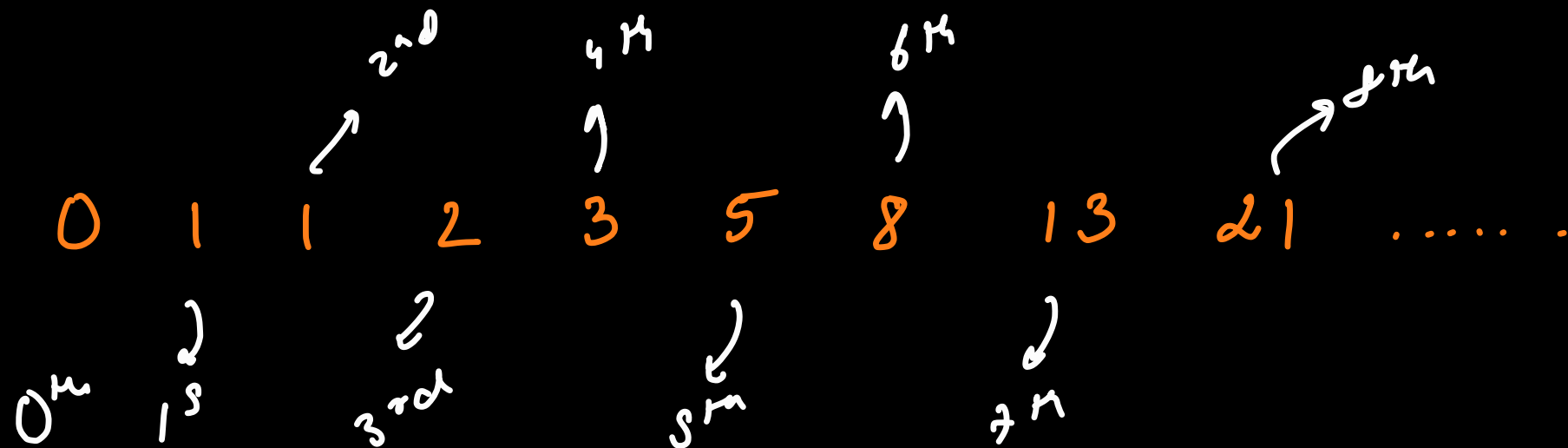
Q-1 Given a non-negative integer value n , Calculate the n^{th} fibonacci. (recursively)

Ex $n = 5$

ans \rightarrow 5

$n = 8$

ans \rightarrow 21

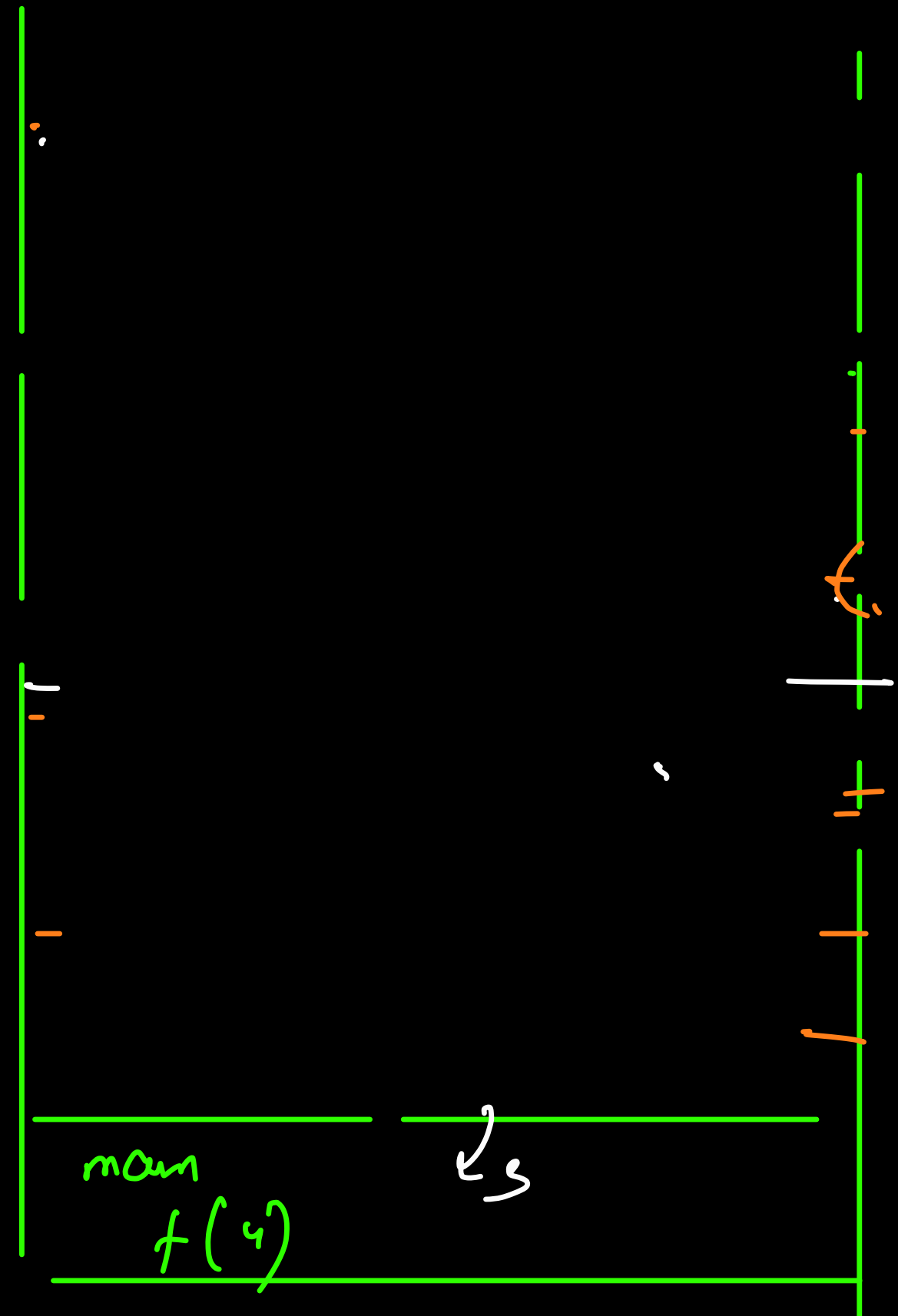
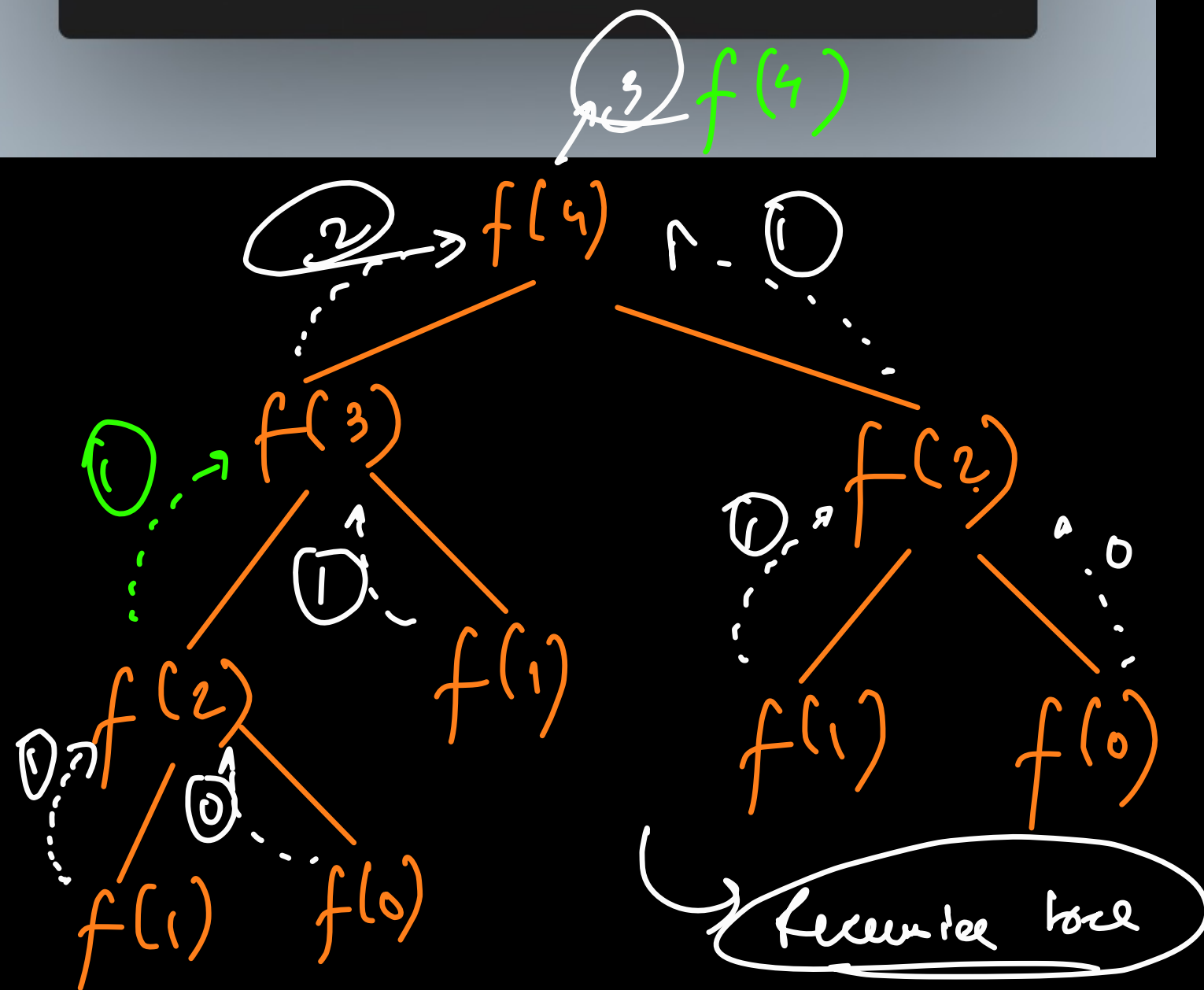
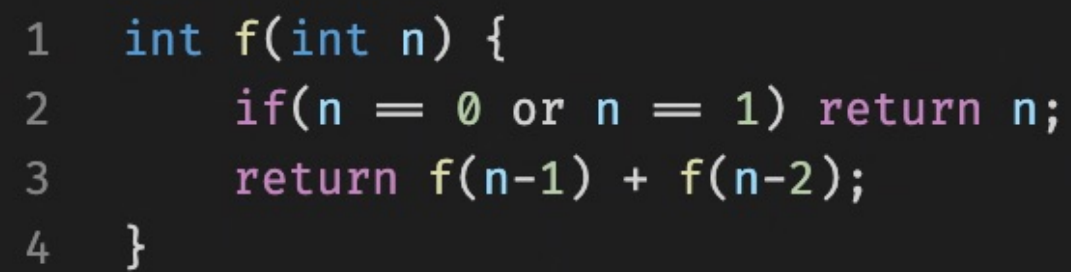


any i^{th} fib can be computed using the sum of $(i-1)^{\text{th}}$ & $(i-2)^{\text{th}}$ fib.

$$\begin{array}{c}
 f(n) \\
 \downarrow \\
 \text{nth fibonacci}
 \end{array}
 =
 \begin{array}{c}
 \overbrace{f(n-1) + f(n-2)}^{\text{self work}} \\
 \downarrow \qquad \qquad \downarrow \\
 \text{assume } f \text{ returns} \quad \text{assume } f \text{ returns} \\
 (n-1)^{\text{th}} \text{ fib correctly} \quad (n-2)^{\text{th}} \text{ fib correctly}
 \end{array}$$

Base \rightarrow

$$\left. \begin{array}{l}
 f(0) \rightarrow 0 \\
 f(1) \rightarrow 1
 \end{array} \right\}$$



Q Given a number n , print the first n natural no,
in increasing order recursively.

Ex $\rightarrow n = 4$

Ans \rightarrow 1
2
3
4

$f(n) \rightarrow$
 \downarrow

print first
 n natural no.

recursively

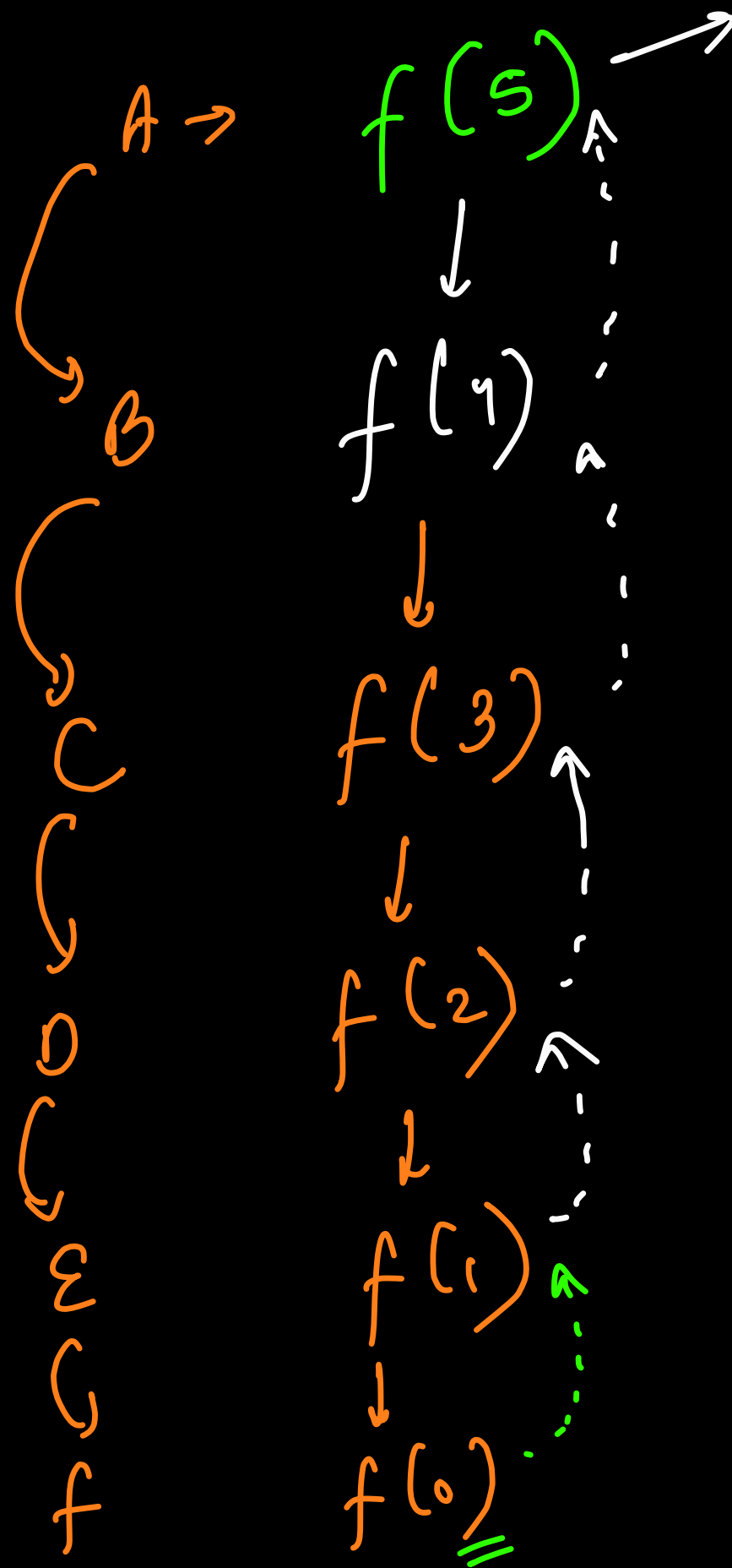
if ($n \rightarrow 0$)
 \downarrow
don't do

any res

assume f works correctly
for $n-1$ and print first $n-1$ natural
no.

$f(n-1);$

print(n);



console

1
2
3
4
5