

Q=1 Given a string of length  $n$ , calculate if the string is a palindrome or not recursively ??

Ex → NAMAN

Ans → Yes

Iterative

1

$i$   $j$   
NAMAN

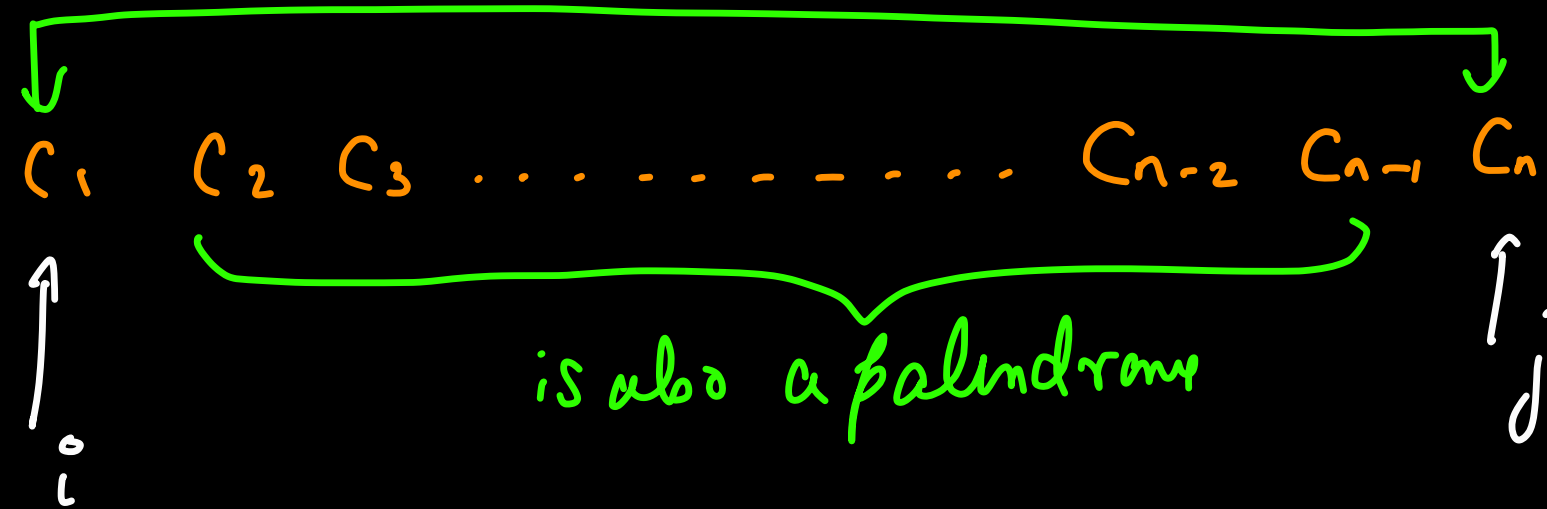
$i$   $j$   
SANKET

```
while (i <= j) {  
    if (s[i] != s[j]) return false;  
    i++;  
    j--;
```

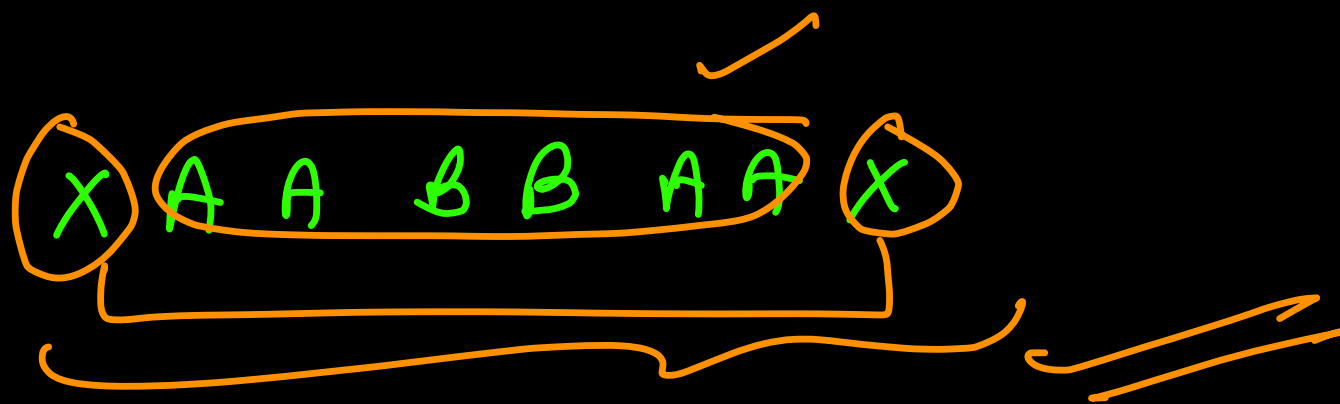
}

JOIN THE DARKSIDE

palindrome



$(C_1 == C_n)$  and remaining string is also a palindrome.  
then we can say the whole string is a palindrome.



$$f(str, i, j) = (str[i] == str[j]) \text{ and } f(str, i+1, j-1)$$

$\downarrow$   
 whether str from  
 index  $i$  to  $j$  is a

palindrome or not ??

$$\hookrightarrow \underline{\underline{if(i \leq j)}}$$

$x \ A \ A \ B \ A \ A \ x$   
 $\quad \quad \quad \uparrow \quad \uparrow$   
 $\quad \quad \quad i \quad j$

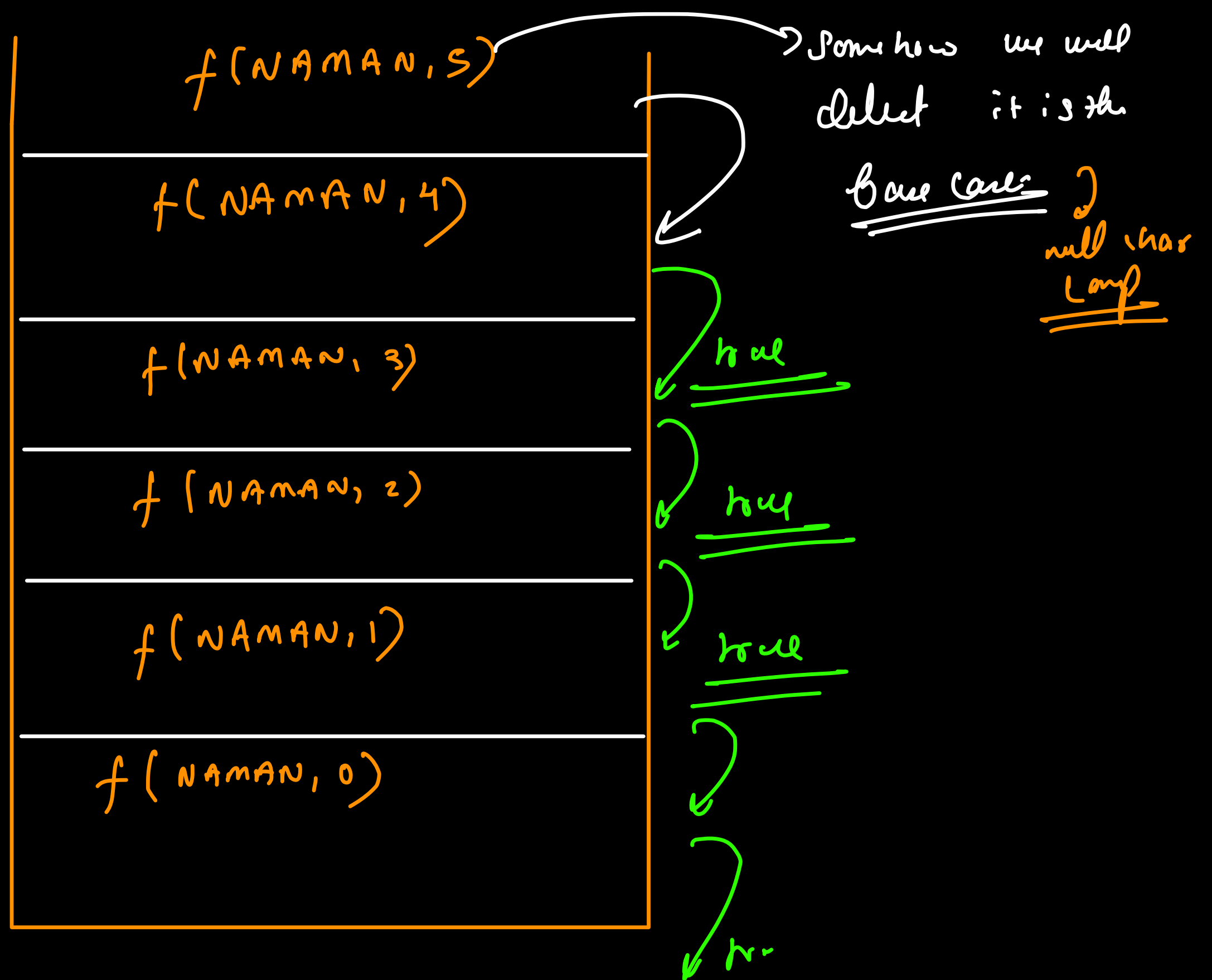
$N \ \boxed{A \ M \ A} \ N$   
 $\uparrow \quad \quad \quad \uparrow$   
 $i \quad \quad \quad j$

How can we solve the prev prob, without using the str.size() function

↳ let's keep the  $i^{\text{th}}$  index global and push the  $j^{\text{th}}$  index to the last ??

$f(\text{str}, i) \rightarrow f(\text{str}, j+1)$   
 $(\text{if } (\text{str}[i] == \text{str}[j]))$

global  
 $i = 0, 1, 2, 3$



$\therefore \underline{\underline{6 \times 2}}$

$\begin{matrix} i & j \\ 2 & 2 \end{matrix}$

A B B X C B A

(  
    non  
  (  
    true  
  (  
    non  
    (  
      false

Q<sub>2</sub> for an input of  $n$ , print the following pattern  
recursively.

$n = 5$

1	★				
2	★	★			
3	★	★	★		
4	★	★	★	★	
5	★	★	★	★	★

```
for (row = 1; row <= n; row++) {  
    for (col = 1; col <= row; col++) {  
        print (*)  
    }  
    print (\n)  
}
```

for any  $i^{\text{th}}$  row  $\rightarrow \text{col} \in \underline{\underline{[1, i]}}$

$f(n, r) = \text{for } (col = 1; col \leq r; col++) \{$   
 $\text{print}(*)$



this func<sup>n</sup> prints the  
pattern from  $r^{\text{th}}$  row  
to the  $n^{\text{th}}$  row.

$\}$   
 $\text{print}(n)$

$f(n, r+1)$

Base  $\Rightarrow$  if  $(r > n)$   
return

$f(n, 1)$

my self work will be to print stars for the  
 $r^{\text{th}}$  row.



$f(n, r, c)$

prints the pattern  
from  $r^{\text{th}}$  row,  $c^{\text{th}}$  col  
to the  $n^{\text{th}}$  row &  $n^{\text{th}}$  col

	1	2	3	4	5
1	*				
2	*	*			
3	*				
4					
5					

if ( $c > n$ ) { go to a new row  
 print("\n")  
 $f(n, r+1, 1)$   
 return }

print(\*)  
 $f(n, r, c+1)$

Base case

if ( $r > n$ )  
return;

$f(4, 3, 1)$
$f(4, 2, 3)$
$f(4, 2, 2)$
$f(4, 2, 1)$
$f(4, 1, 2)$
$f(4, 1, 1)$
main

$$f(n, r, \text{output}) = \text{print}(\text{output}) \quad \left. \vphantom{f(n, r, \text{output})} \right\} \\ f(n, r+1, \text{output} + "\star")$$

$f(n, 1, "\star")$

$\star$   
 $\star \star$   
 $\star \star \star$

↓

$f$
$f(4, 3, \star \star \star)$
$f(4, 2, \star \star)$
$f(4, 1, \star)$
main

x

(N)  $\Leftarrow$   $N^{\text{th}}$  power of natural no.

1, 2, 3, 4, . . . . .

The natural no.'s that can sum upto x, will be  
in the range  $[1, x]$

x  
x = 100

x = 10  $\rightarrow$

2  $\leq N \leq 10$

(12)  $\times$   $[1, \sqrt{x}]$   $\leftarrow$

10  
(12)

$$\underline{\underline{[1, \sqrt{x}]}}$$

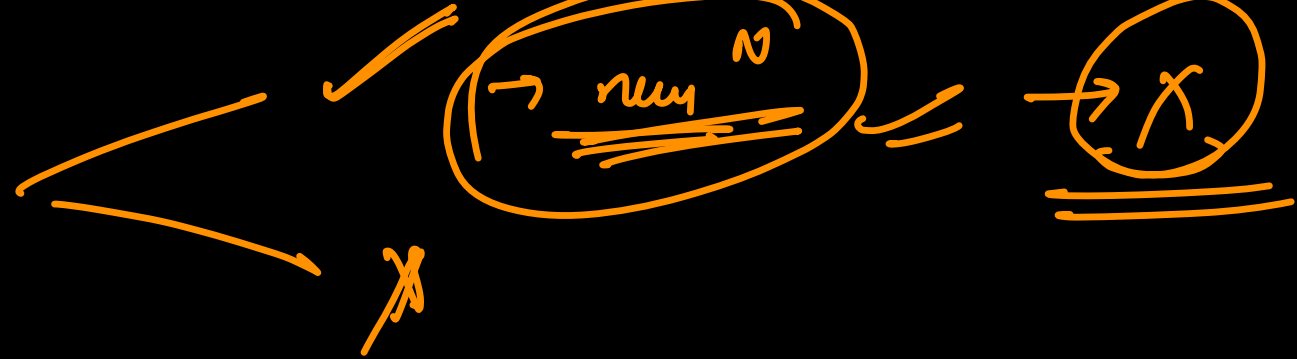
↓

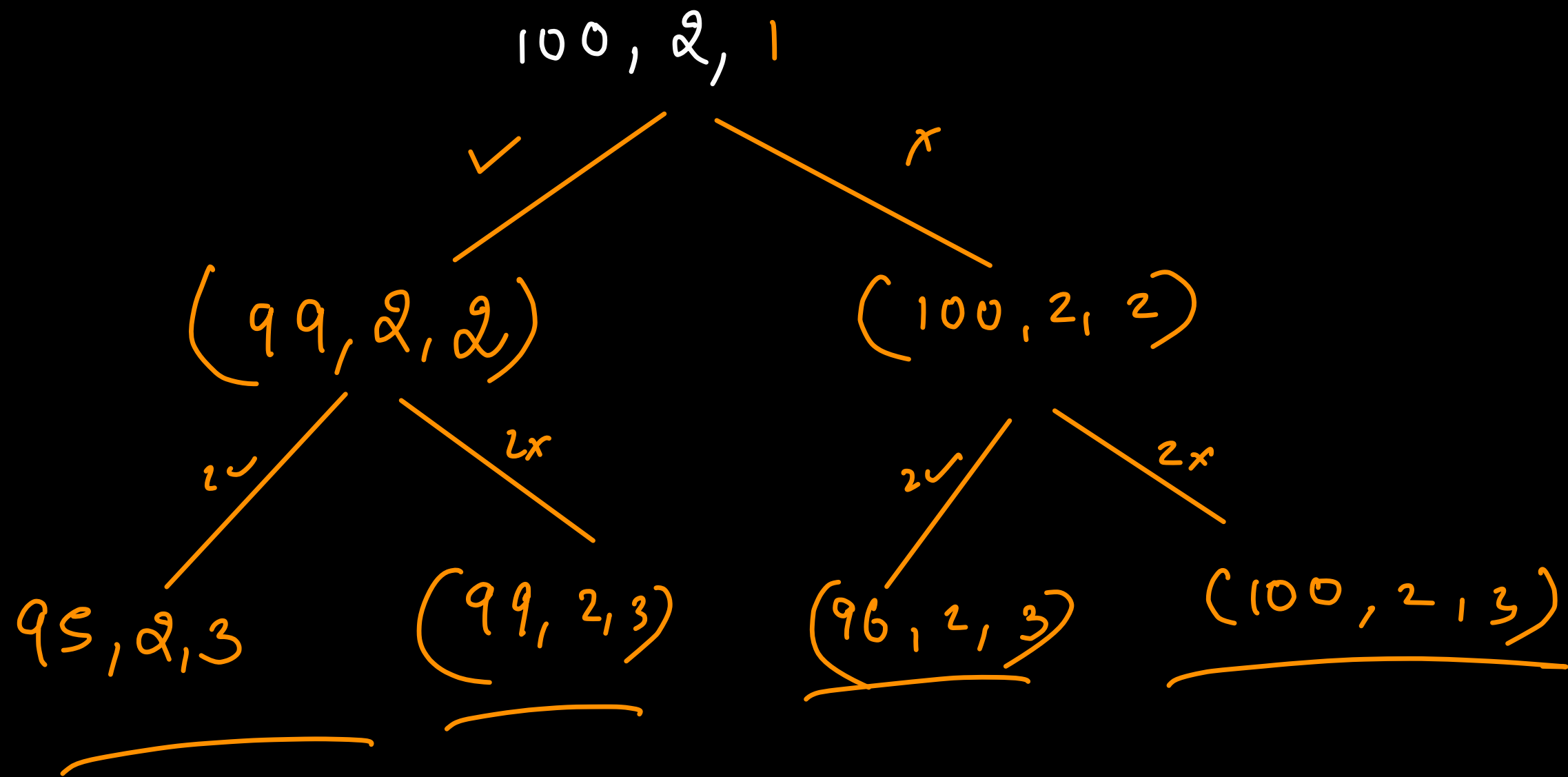
how many combinations of no.'s we can have  
such that sum of  $n^{\text{th}}$  powers of that no.  
is equal to  $x$ .

$$100 \rightarrow [1, \sqrt{100}]$$

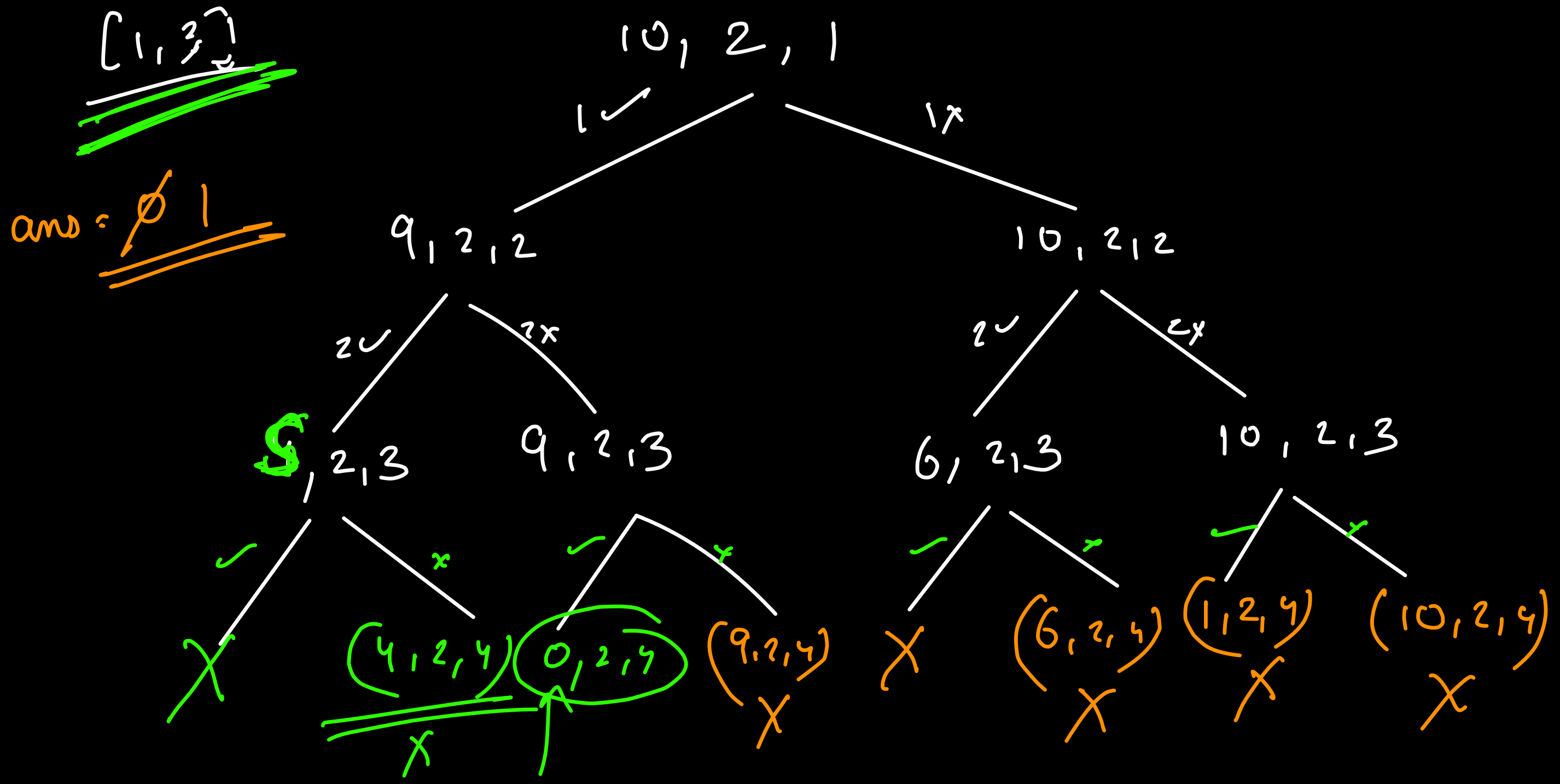
$$\rightarrow \underline{\underline{[1, 10]}}$$

$$N=2$$



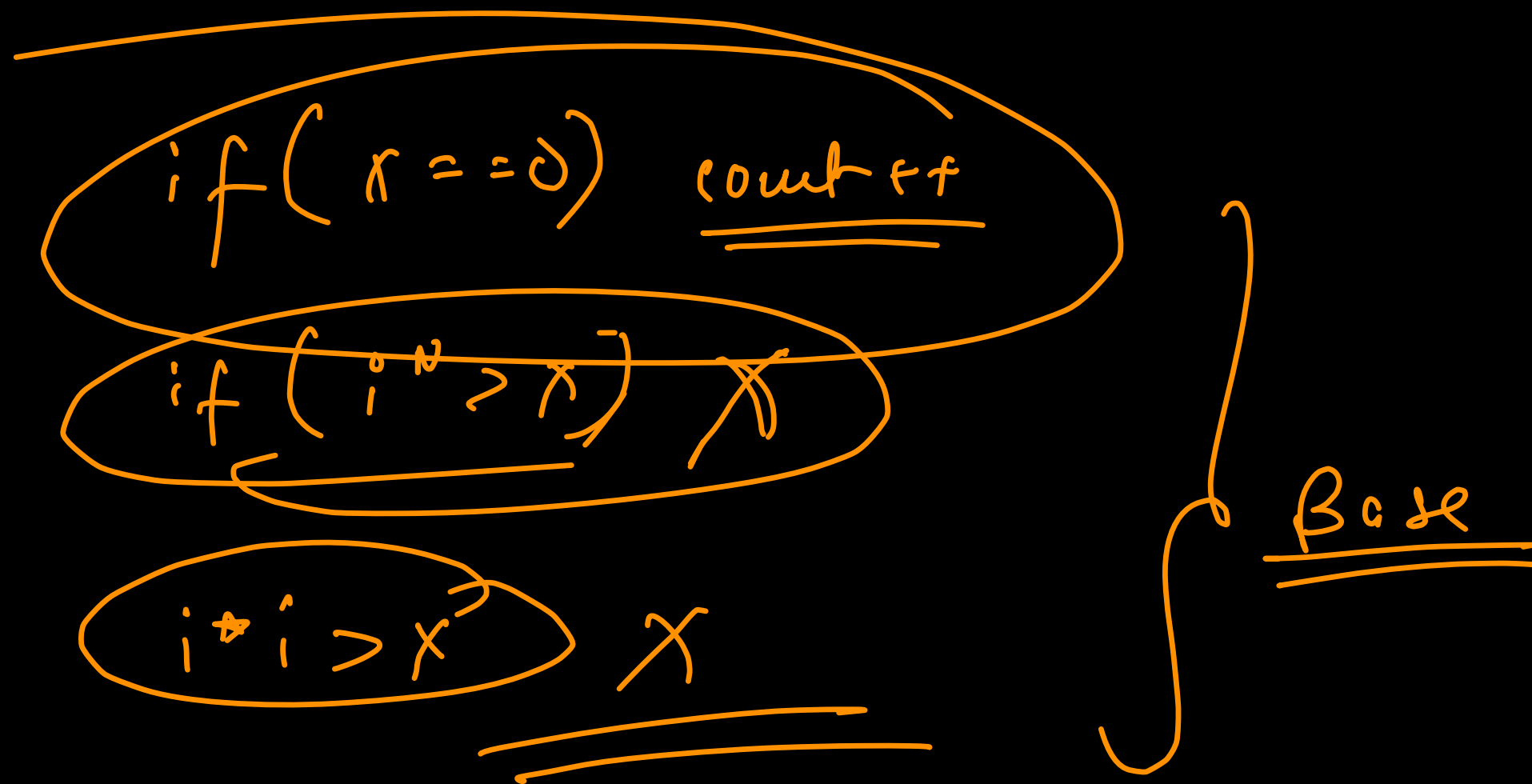


Try to reduce the problem such that it becomes  
Similar to Subsequences



$$f(x, n, i) = f(x - i^n, n, i+1)$$

$$f(x, n, i+1)$$





$$f(x, n, i, out) = f(x - i^n, n, i+1, out + i^n) \\ f(x, n, i+1, out)$$

pow(i, n)

$n=1$   $\rightarrow$   $()$

$n=2$   $() ()$   $(( ))$

$n=3$   $() () ()$

$(( ( ) ) )$

Balanced

$(( ))$

$(( )) )$   $(( ))$

$(( ))$   $(( ))$

Printing

29

$n \rightarrow$  opening  
 $n \rightarrow$  closing

$() (( ))$

for every index we  
can either add an  
 $($  or  $)$

```

f(n, o, c, output) = if (o > c) {
                      f(n, o, c+1, output + ")")
                    }
                    if (o < n) {
                      f(n, o+1, c, output + "(")
                    }

```

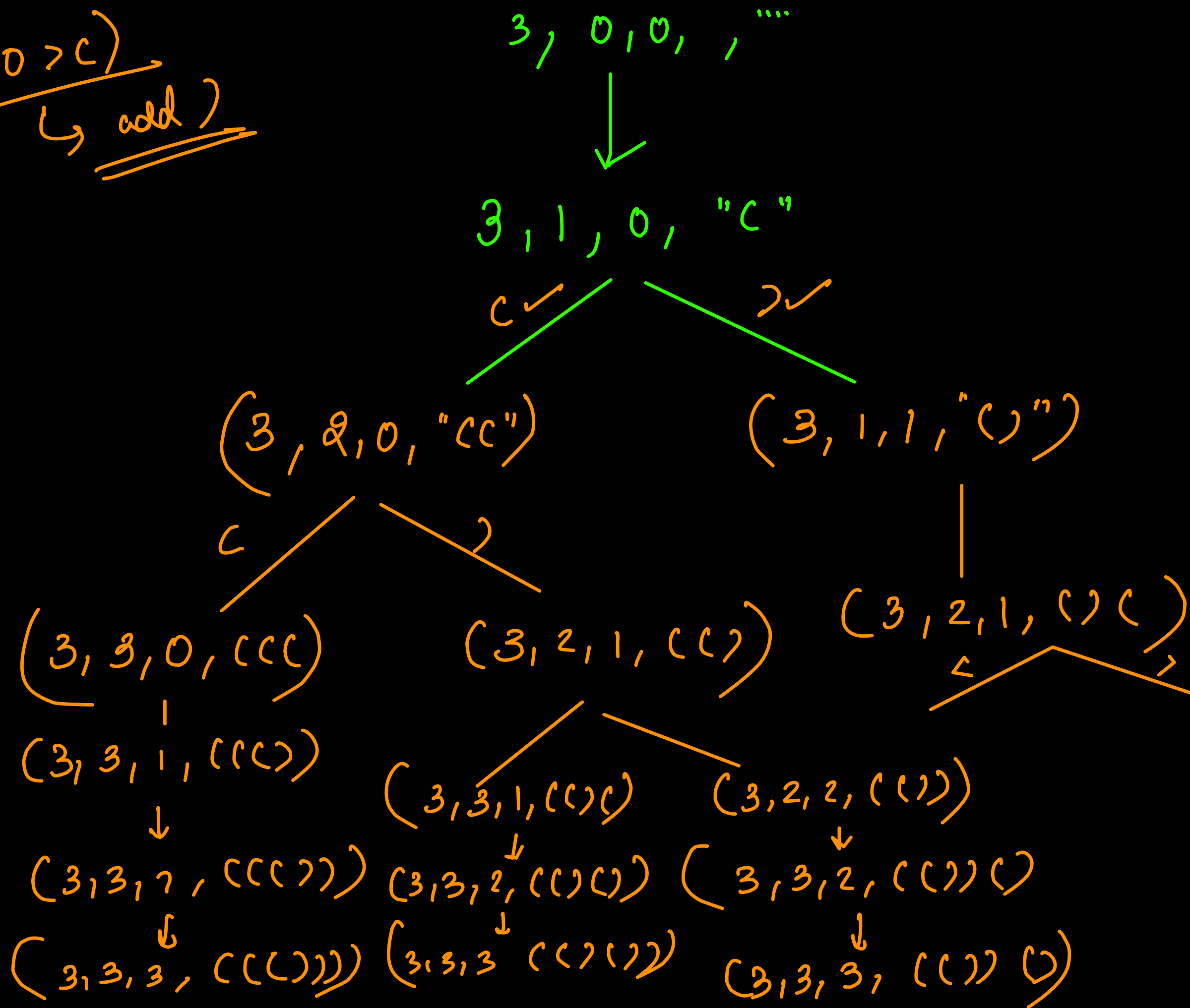
Base Case

```

if (c == n)
  print(output)
  return;

```

if (0 > c) →  
↳ add)



0==c

( ) ( ) X

0>c

( ) ( ) ( ( ) )