

Longest Common Subsequence

→ given 2 strings → s_1 & s_2

length of ← longest common subsequence

Ex → s_1 → x y z z y ✓
 s_2 → x b y b z y → x, b, x b y, x b z, ...

↓

$x y z y$
 $x y z y$

 → 4

→ Common Subsequence → Common char

↓
ordering also matters

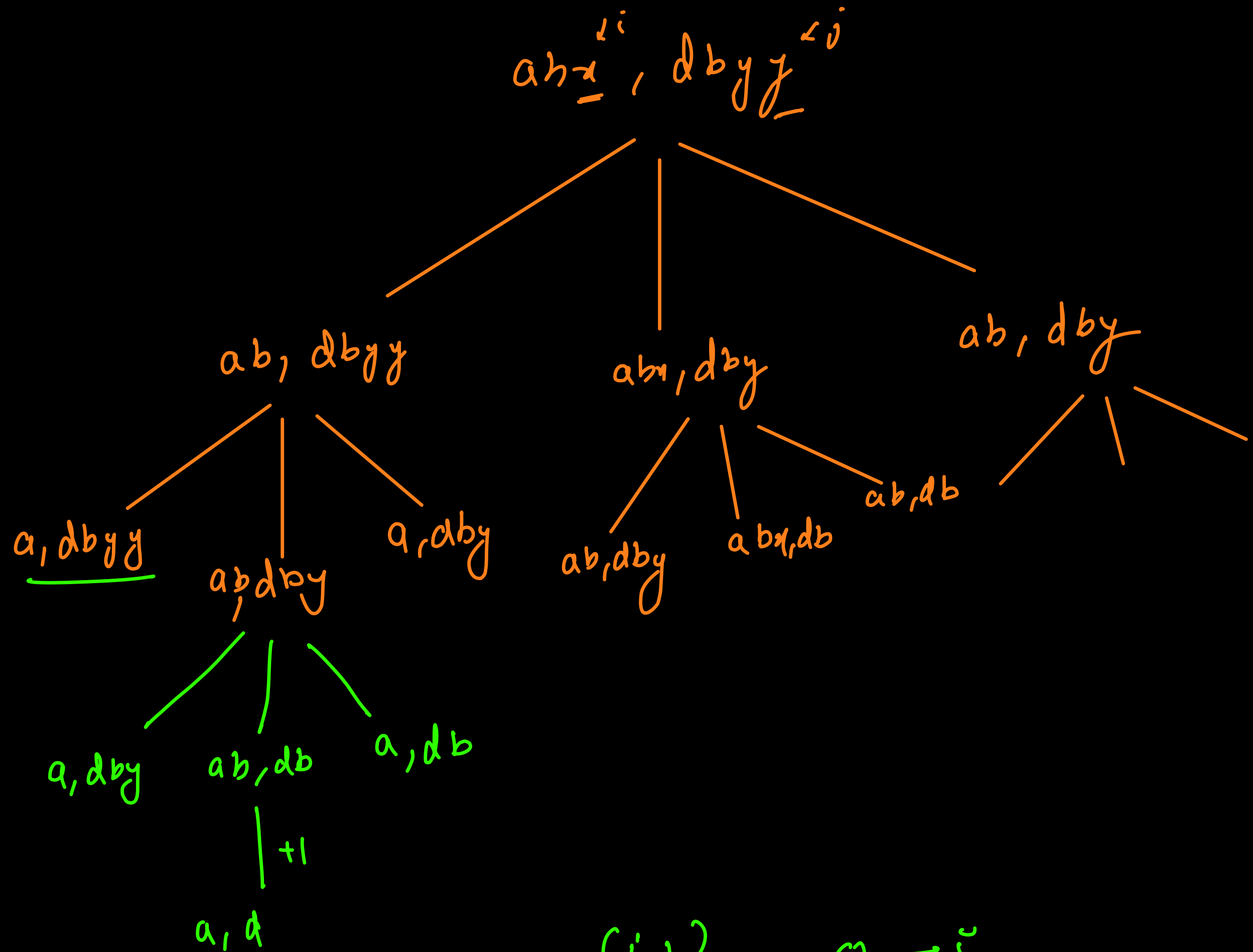
if we go in desc by index, we can maintain order

$lcs(abxz, dbgyz)$

→ $\begin{matrix} s_1 = abxz \\ s_2 = dbgyz \end{matrix}$

→ $i^{th} \& j^{th}$ char are equal

$1 + lcs(abx, dbgy)$



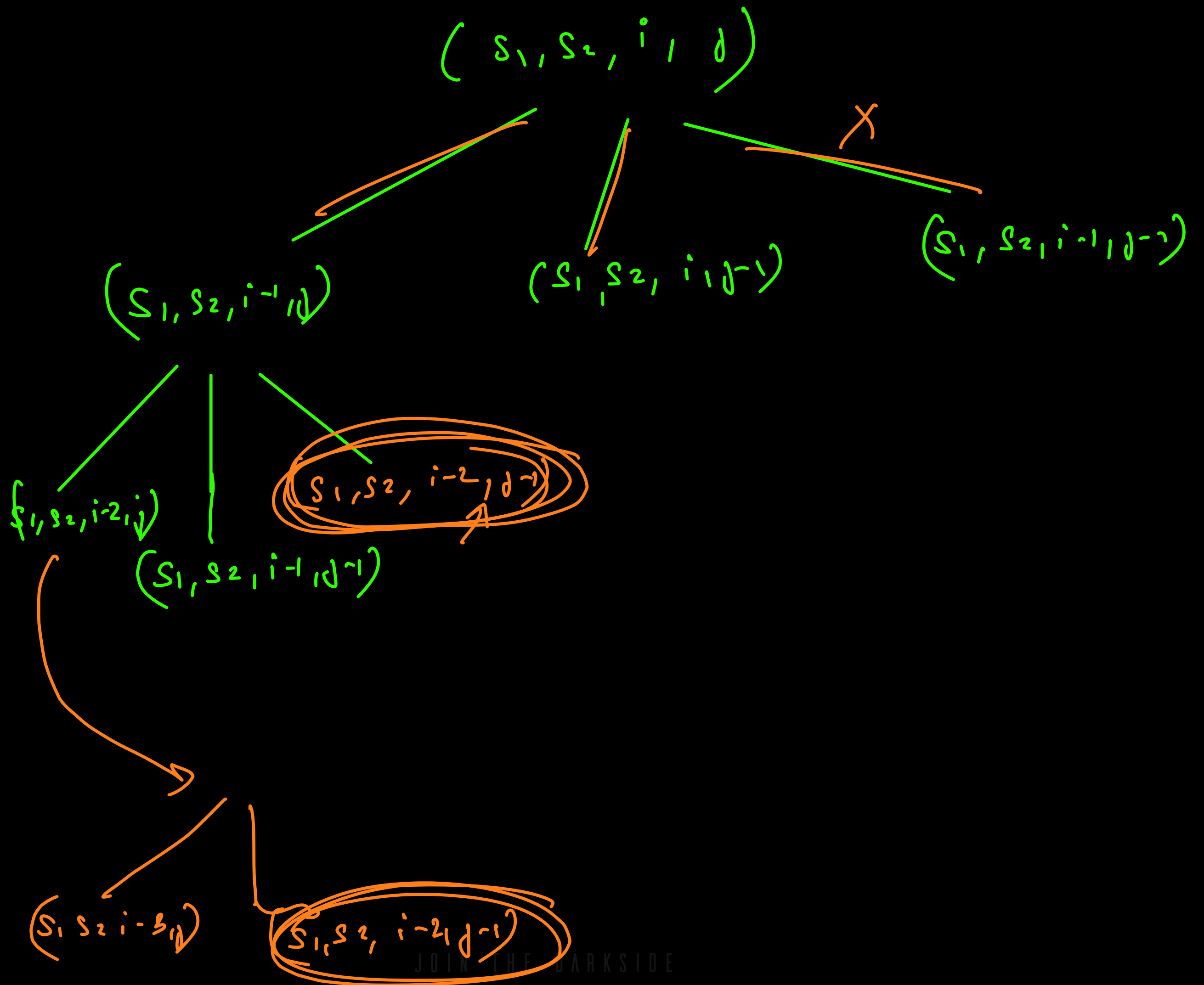
(i, j)

$0 \rightarrow i$
 $0 \rightarrow j$

$f(s_1, s_2, i, j)$
 \downarrow
 length of the LCS
 btw $s_1[0, i]$ and
 $s_2[0, j]$

\downarrow
 max \rightarrow 2d
Storage

$$= \begin{cases} 1 + f(s_1, s_2, i-1, j-1) & \text{if } s_1[i] == s_2[j] \\ \max \begin{pmatrix} f(s_1, s_2, i-1, j) \\ f(s_1, s_2, i, j-1) \end{pmatrix} & \text{else} \end{cases}$$



$$(abx, acx, 2, 2) \rightarrow \underline{\underline{2}}$$

$$| +1$$

$$(abx, acx, 1, 1) \rightarrow 1$$

$$\frac{(\dots, ac)}{\dots} \rightarrow \underline{\underline{0}}$$

$$\frac{(\dots, \dots)}{\dots}$$

$$(abx, acx, 0, 1) \rightarrow \underline{\underline{1}}$$

$$(abx, acx, 1, 0) \rightarrow 1$$

$$(abx, acx, -1, 1)$$

$$(abx, acx, 0, 0)$$

$$| +1$$

$$(abx, acx, -1, -1) \leftarrow 0$$

$$(abx, acx, 0, 0)$$

$$(abx, acx, 1, -1)$$

$$\begin{array}{l}
 f(s_1, s_2, i, j) \\
 \downarrow \\
 \text{ LCS } \left(\begin{array}{l} s_1 [i, n-1] \\ s_2 [j, m-1] \end{array} \right)
 \end{array}
 = \begin{cases}
 1 + f(s_1, s_2, i+1, j+1) & \text{if } (s_1[i] == s_2[j]) \\
 \max \left(\begin{array}{l} f(s_1, s_2, i+1, j) \\ f(s_1, s_2, i, j+1) \end{array} \right) & \text{else}
 \end{cases}$$

abbc b, bcbb a

bbb

lcs(a c b d c d a, a d c d b c a)

acbc a

2

a c b d c d a

↑ i ↓ ↓ i

↑ ↑ ↑

(a) c b d c d (9)

↓

dpc

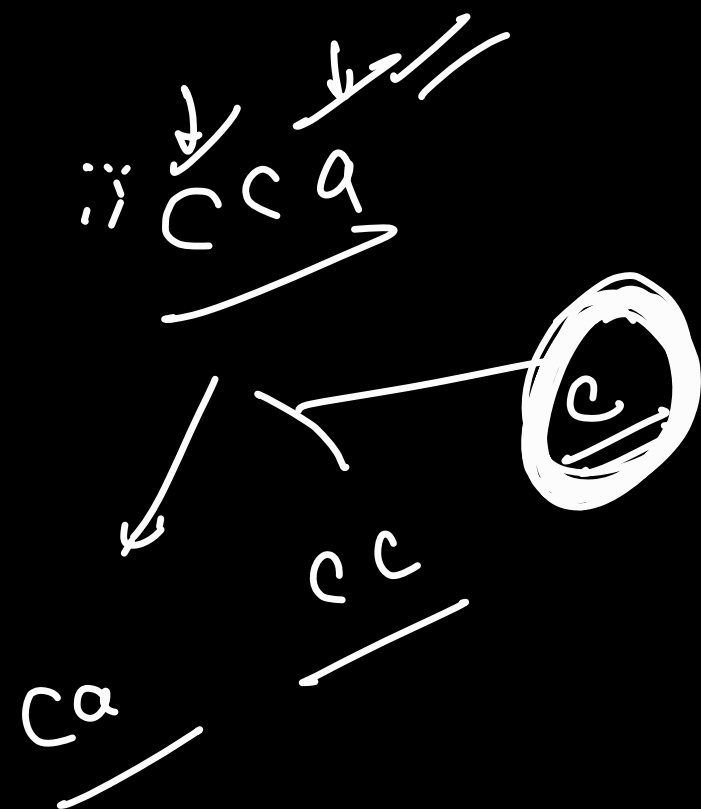
$ch + \underline{\text{palindrome}} + ch \rightarrow \underline{\text{palindrome}}$

c a b a c //

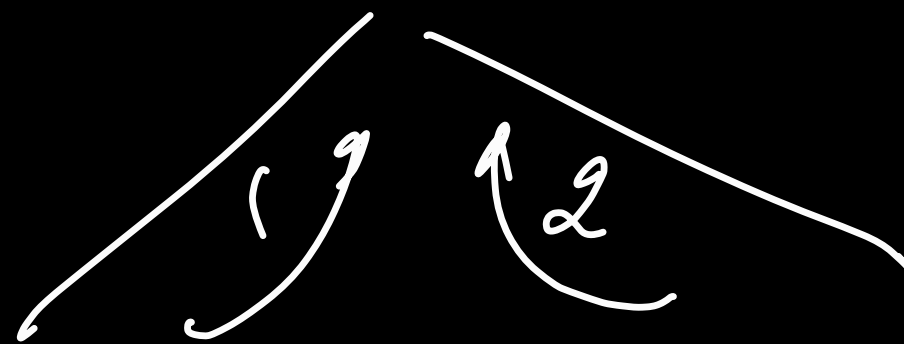
$f(s, i, j)$
 \downarrow
 length of lps
 $s[i, j]$
 \uparrow

$$\begin{aligned}
 &= \begin{cases} 2 + f(s, i+1, j-1) & \text{if } s[i] == s[j] \\ \max \begin{pmatrix} f(s, i+1, j) \\ f(s, i, j-1) \end{pmatrix} & \text{else} \end{cases}
 \end{aligned}$$

$(b \underset{\uparrow}{c} c a \underset{\uparrow}{b}, 0, 4) \rightarrow 4$



$(b \check{c} c a b, 1, 3) \rightarrow 2$



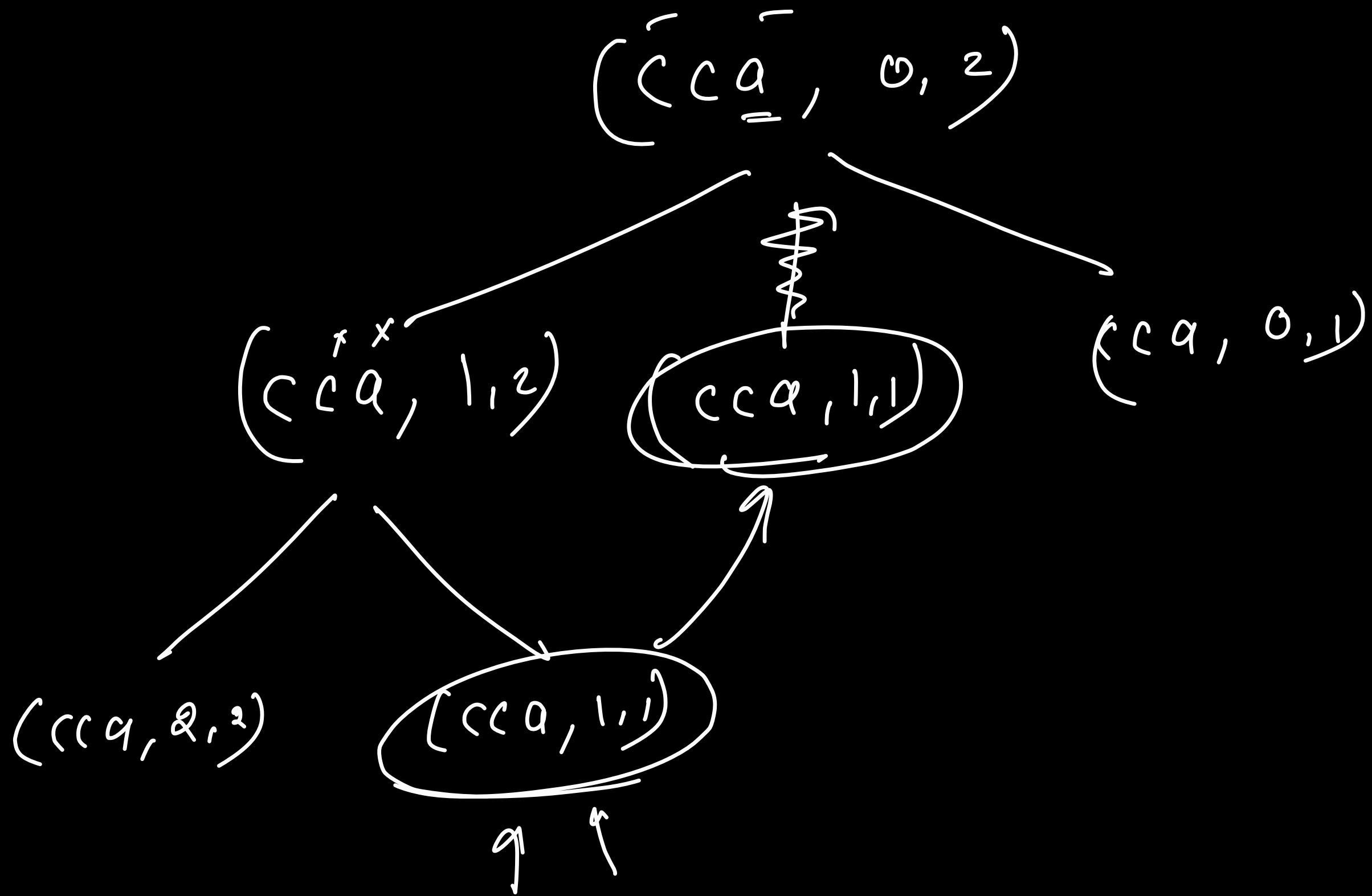
$(b \check{c} c a b, 1, 2) \rightarrow 2$

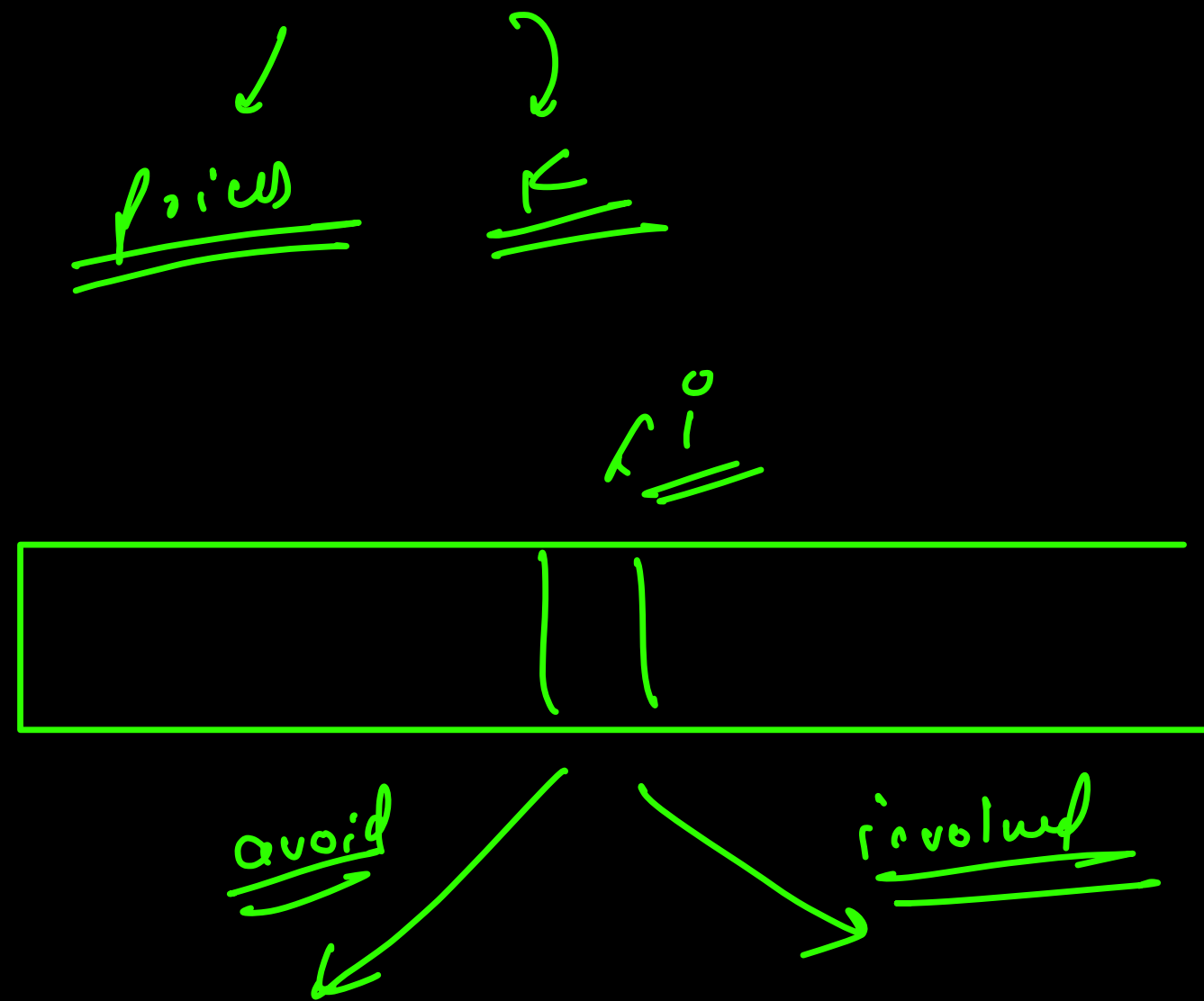
$(b \check{c} c \check{a} b, 2, 3)$

$(b \check{c} c a b, 3, 3)$

$(b \check{c} c a b, 4, 2)$

$(b \check{c} c a b, 2, 1) \rightarrow 0$





if bought → sell
else → buy → if K is left

a txn is going on if you bought already but didn't

$f(i, on, k)$ → 3rd

max profit from
 1st day to $(n-1)^{th}$
 day with k txn

left.

$on \rightarrow$ if a txn is
 going on or not

↓

$f(0, false, k)$
 ↳ ans

$$= \max \begin{cases} f(i+1, on, k) \\ -prices[i] + f(i+1, true, k) \\ +prices[i] + f(i+1, False, k-1) \end{cases}$$

ignore

buy
 ↓
 $k > 0$ and
 $on == \underline{false}$

sell
 ↓
 $on == true$

