

## Dynamic Programming Lecture 4

Thursday, 29 August 2024 6:01 AM

### 0/1 Knapsack Problem

↳ Given a set of items, each with a weight and a value, determine the maximum value that can be obtained by selecting a subset of items, subject to a total weight constraint. (i.e. the total weight of selected items should not exceed the capacity of the knapsack)

		W	Value/Profit	P/W	
eg.	I1	1	6	$6/1 = 6$	Capacity = 5 units
	I2	2	10	$10/2 = 5$	
	I3	3	12	$12/3 = 4$	

Fractional knapsack:- 
$$\begin{aligned}\text{Max Profit} &= 6 \times 1 + 5 \times 2 + 4 \times 2 \\ &= 6 + 10 + 8 \\ &= \boxed{24}\end{aligned}$$

0/1 Knapsack :-  $I_1 + I_2$ ,  $I_1 + I_3$ ,  $I_2 + I_3$   
                                   16                  18                  22

Max Profit = 22

$$\text{max Profit Knapsack}(n, W) = \begin{cases} 0, & n=0 \text{ or } W=0 \\ \text{MPK}(n-1, W), & w[n] > W \\ \max\left(\text{MPK}(n-1, W), \right. & \\ \left. w[n] + \text{MPK}(n-1, W - w[n])\right), & w[n] \leq W \end{cases}$$

DP approach:-

function maxProfitKnapsack ( $n$ ,  $W$ , weights, values) :

$\downarrow$  # items    $\downarrow$  capacity    $\downarrow$  [i] weight of items    $\downarrow$  [n] value of items

dp = matrix of size  $(n+1) \times (W+1)$  with first row and first column initialized to 0

$\leftarrow$  dp[i][j] :- The max profit if you have first 'i' items & capacity of 'j'

for (i: 1  $\rightarrow$  n) :

for (j: 1  $\rightarrow$  W) :

if weights[i-1]  $\leq$  j :

dp[i][j] = max ( dp[i-1][j],  
                   $\underbrace{\text{values[i-1]}}_{\uparrow \text{ i-th item}} + \underbrace{\text{dp[i-1][j - \text{weights[i-1]}]}}_{\substack{\text{i-1 items} \\ \text{ i-th item}}} )$

else :

dp[i][j] = dp[i-1][j]

return dp[n][W]

eg: 

		w	values
I1	2	w[0]	3 v[0]
I2	3	w[1]	4 v[1]
I3	4	w[2]	5 v[2]
I4	5	w[3]	6 v[3]

 Knapsack capacity (w) = 5

				Cap →						
				dp	0	1	2	3	4	5
Sol.	item	v <sub>i-1</sub>	w <sub>i-1</sub>	0	0	0	0	0	0	0
				1	0	0	3	3	3	3
				2	0	0	3	4	4	7
				3	0	0	3	4	5	7
				4	0	0	3	4	5	7

Max Profit = 7

Items taken: - I1, I2

Time: -  $O(n \cdot W)$   
 ↓      ↑  
 # items    Capacity

Space: -  $O(n \cdot W)$

$dp[3][5]$   $w[2]=4$   $v[2]=5$   
 $dp[3][5] = 5 + dp[2][1]$

$dp[1][4]$   
 ↓  
 one items, capacity = 4  
 first items  
 $w = 5$  |  $w = 3$   
 $v = 6$  |  $v = 5$

①  $w > j$ :  
 $dp[0][4]$

②  $w \leq j$ :

$\max ( dp[0][4] , 5 + dp[0][1] )$   
 ↑    ↑

$dp[3][4]$   
 --- ↑  
 3rd items

eg:

Consider the following set of items, each with a given weight and value:

- Item 1: weight = 1, value = 1
- Item 2: weight = 3, value = 4
- Item 3: weight = 4, value = 5
- Item 4: weight = 5, value = 7

The maximum weight capacity of the knapsack is  $W = 7$ .

$W = 7$   
14, 8, 23  
9

Select all that apply:

- ☒ A. The maximum value that can be achieved with the given capacity is 9.
- ☒ B. If you include Item 4, you cannot achieve the maximum value.
- ☒ C. The optimal solution includes Item 2 and Item 3.
- ☒ D. If you exclude Item 1, the maximum value you can achieve is 9.
- ☒ E. The optimal solution includes Item 1, Item 3, and Item 4.

[GATE CS 2018]

Consider the weights and values of items listed below. Note that there is only one unit of each item.

Item number	Weight (in Kgs)	Value (in rupees)
→ 1	10	60
→ 2	7	28
→ 3	4	20
→ 4	2	24

$$\begin{aligned} v/w \\ 60/10 &= 6 \\ 28/7 &= 4 \\ 20/4 &= 5 \\ 24/2 &= 12 \end{aligned}$$

The task is to pick a subset of these items such that their total weight is no more than 11 Kgs and their total value is maximized. Moreover, no item may be split. The total value of items picked by an optimal algorithm is denoted by  $V_{opt}$ . A greedy algorithm sorts the items by their value-to-weight ratios in descending order and packs them greedily, starting from the first item in the ordered list. The total value of items picked by the greedy algorithm is denoted by  $V_{greedy}$ .

The value of  $V_{opt} - V_{greedy}$  is \_\_\_\_

$$60 - 44 = \boxed{16}$$

$V_{opt}$  :

$$1 \rightarrow \boxed{60}$$

$$\begin{array}{ccc} 2,3 & 2,4 & 3,4 \\ 48 & 52 & 44 \end{array}$$

(A)  $V_{opt} = 60$

(B)  $V_{greedy} = 24 + 20 = 44$

### Subset Sum Problem

↳ Given a set of non-negative integers and a value Sum, determine if there is a subset with sum equal to the given sum.

eg:  $S = \{3, 34, 4, 12, 5, 2\}$  Sum = 9

o/p:- True  $\{4, 5\}$  or  $\{3, 4, 2\}$

Brute Force:- check all possible subsets,  $O(2^n)$



DP approach:-

function isSubsetSum (set, n, sum):

dp = boolean matrix of size  $(n+1) \times (sum+1)$

dp[5][0]

for (i: 0  $\rightarrow$  n): // First Column

dp[i][0] = True

for (i: 1  $\rightarrow$  sum): // First Row

dp[0][i] = False

for (i: 1  $\rightarrow$  n):

for (j: 1  $\rightarrow$  sum):

if set[i-1] > j:  $\rightarrow$  leave

dp[i][j] = dp[i-1][j]

if set[i-1]  $\leq$  j:  $\rightarrow$  take it or leave it

dp[i][j] = dp[i-1][j] or dp[i-1][j - set[i-1]]  
leave take

return dp[n][sum]

$T = O(n \cdot sum)$

$S = O(n \cdot sum)$

$\hookrightarrow$  can be reduced to  $O(\min(n, sum))$



eg.  $S = \{3, 34, 4, 12, 5, 2\}$        $Sum = 9$

		Sum ele dp									
		0	1	2	3	4	5	6	7	8	9
	0	T	F	F	F	F	F	F	F	F	F
3	1	T	F	F	T	F	F	F	F	F	F
34	2	T	F	F	T	F	F	F	F	F	F
4	3	T	F	F	T	T	F	F	T	F	F
12	4	T	F	F	T	T	F	F	T	F	T
5	5	T	F	F	T	T	T	F	T	T	T
2	6	T	F	T	T	T	T	T	T	T	T

$\{5, 4\}$

$\{2, 4, 3\}$

[GATE CS 2008]

The subset-sum problem is defined as follows. Given a set of  $n$  positive integers,  $S = \{a_1, a_2, a_3, \dots, a_n\}$ , and positive integer  $W$ , is there a subset of  $S$  whose elements sum to  $W$ ? A dynamic program for solving this problem uses a 2-dimensional Boolean array,  $X$ , with  $n$  rows and  $W + 1$  columns.  $X[i, j]$ ,  $1 \leq i \leq n$ ,  $0 \leq j \leq W$ , is TRUE, if and only if there is a subset of  $\{a_1, a_2, \dots, a_i\}$  whose elements sum to  $j$ .

Which of the following is valid for  $2 \leq i \leq n$ , and  $a_i \leq j \leq W$ ?

- A.  $X[i, j] = X[i - 1, j] \vee X[i, j - a_i]$
- ☒ B.  $X[i, j] = X[i - 1, j] \vee X[i - 1, j - a_i]$
- C.  $X[i, j] = X[i - 1, j] \wedge X[i, j - a_i]$
- D.  $X[i, j] = X[i - 1, j] \wedge X[i - 1, j - a_i]$

[GATE CS 2008]

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Which entry of the array  $X$ , if TRUE, implies that there is a subset whose elements sum to  $W$ ?

- A.  $X[1, W]$
- B.  $X[n, 0]$
- ☒ C.  $X[n, W]$
- D.  $X[n - 1, n]$