

Number Theory Lecture 6

Friday, 12 July 2024

8:14 PM

$$\phi(n) = n - \sum_{p|n} \left(\frac{n}{p}\right) \leftarrow O(\sqrt{n})$$

$$\phi(n) = n \prod_{p|n} \left(1 - \frac{1}{p}\right)$$

H/w:- <https://www.spoj.com/problems/ETF/>

Find Euler's totient function of n

$$1 \leq n \leq 10^6$$

$$1 \leq T \leq 2 \times 10^4$$

like
Sieve

$$n \log \log n$$

$$\frac{10^6 \times \log \log 10^6}{\log 6 \cdot \log 10} < \log 6 \cdot 4 < 5$$

$$O(T \cdot \sqrt{n})$$

$$2 \times 10^4 \times 10^3 \sim 2 \times 10^7 \checkmark$$

$$5 \times 10^6 + 2 \times 10^7$$

```

int phi(int n) {
    int ans = n;
    for(int i=2; i*i<=n; i++) {
        if(n%i == 0) {
            while(n%i == 0)
                n /= i;
            ans = ans - ans/i;
        }
    }
    if(n>1) ans = ans - ans/n;
    return ans;
}

```

$$\textcircled{n} = \overline{2^3 \times 3^2 \times 7^4 \times 11^2 \dots}$$

$$\uparrow$$
 ans = \textcircled{n}
 for every prime factor p
 of n, ans = ans - ans/p

$\textcircled{p} \rightarrow \underline{\underline{\text{ans} = \text{ans} - \text{ans}/p}}$

```

N = 1000000
phi = [i for i in range(N+1)]
for i in range(2, N+1):
    if phi[i] == i:
        for j in range(i, N+1, i):
            phi[j] -= phi[j]//i

```

$O(N \log \log N)$

$\phi = [0, 1, \overset{\checkmark}{\underset{\uparrow}{2}}, \overset{\checkmark}{\underset{\uparrow}{3}}, \overset{\checkmark}{\underset{\uparrow}{4}}, \overset{\checkmark}{\underset{\uparrow}{5}}, \overset{\checkmark}{\underset{\uparrow}{6}}, \overset{\checkmark}{\underset{\uparrow}{7}}, \overset{\checkmark}{\underset{\uparrow}{8}}, \overset{\checkmark}{\underset{\uparrow}{9}}, \overset{\checkmark}{\underset{\uparrow}{10}}]$

```

N = 1000000
phi = [i for i in range(N+1)]
for i in range(2, N+1):
    if phi[i] == i:
        for j in range(i, N+1, i):
            phi[j] -= phi[j]//i
t = int(input())
for _ in range(t):
    n = int(input())
    print(phi[n])

```

Gauss Theorem [⊕]

For each integer $n \geq 1$, $\sum_{d|n} \phi(d) = n$

eg. $n = \underline{10}$ 1, 2, 5, 10

$$\phi(1) + \phi(2) + \phi(5) + \phi(10) = 10$$

$$1 + 1 + 4 + 4 = \underline{10}$$

$$\phi(10) = 10 \cdot \frac{\phi(1)}{1} \cdot \frac{\phi(2)}{2} \cdot \frac{\phi(5)}{5}$$

①
②

Proof: $F(n) = \sum_{d|n} \phi(d)$

$$\phi(a \cdot b) = \phi(a) \cdot \phi(b) \text{ if } \gcd(a, b) = 1$$

$$n = p_1^{k_1} \cdot p_2^{k_2} \cdots p_r^{k_r} \quad (\text{prime factorization of } n)$$

$$F(n) = F(p_1^{k_1}) \cdot F(p_2^{k_2}) \cdots F(p_r^{k_r})$$

$$F(p_1^{k_1}) = \sum_{d|p_1^{k_1}} \phi(d)$$

$$= \phi(1) + \phi(p_1) + \phi(p_1^2) + \phi(p_1^3) + \cdots + \phi(p_1^{k_1})$$

$$= 1 + (\cancel{p_1} - 1) + (\cancel{p_1^2} - \cancel{p_1}) + (\cancel{p_1^3} - \cancel{p_1^2}) + \cdots + (p_1^{k_1} - p_1^{k_1-1})$$

$$= p_1^{k_1}$$

$$\therefore F(p_1^{k_1}) = p_1^{k_1}$$

$$\therefore F(n) = p_1^{k_1} \cdot p_2^{k_2} \dots p_r^{k_r} = n \quad \square$$

```
#include<bits/stdc++.h>
using namespace std;
void phi_1_to_n(int n) {
    vector<int> phi(n+1);
    for(int i=0; i<=n; i++) phi[i] = i;
    for(int i=1; i<=n; i++) {
        for(int j=2*i; j<=n; j+=i) {
            phi[j] -= phi[i];
        }
    }
}
```

$T = ?$

$i=1$	n
$i=2$	$n/2$
$i=3$	$n/3$
\vdots	
$i=n$	n/n

② ④ ④
 ✓ 2 6 8

[illegible]

$$T = n + \frac{n}{2} + \frac{n}{3} + \frac{n}{4} + \dots + \frac{n}{n}$$

$$= n \left[1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n} \right]$$

$$= \frac{1}{2} \sqrt{\frac{5}{2}}$$

$$\approx n \int_1^n \frac{1}{i} di$$

$$= n \cdot [\log i]^n$$

$$= n \cdot \log n$$

$$n + \frac{n}{2} + \frac{n}{3} + \frac{n}{5} + \frac{n}{7} - \boxed{O(n \log \log n)}$$

$$O(n \log n)$$

① Find $\sum_{k=1}^n \gcd(n, k)$

$n=5: \gcd(1,5) + \gcd(2,5) + \gcd(3,5) + \gcd(4,5) + \gcd(5,5)$

assuming $T(\gcd(a,b)) = O(\log b)$

Brute force $T = O(n \cdot \log n)$

$$\sum_{k=1}^n \gcd(n, k) = \sum_{d|n} \underbrace{d}_{\uparrow} \cdot \underbrace{\phi(n/d)}_{\uparrow}$$

Proof: $N=12$, divisors of 12: $\boxed{1, 2, 3, 4, 6, 12}$

$\gcd(\underline{1}, 12) + \gcd(\underline{2}, 12) + \gcd(\underline{3}, 12) + \dots + \gcd(\underline{12}, 12)$

$\gcd(\underline{k}, \underline{12}) = \underline{g} \quad g|k, g|12$

①- $\gcd(k, 12) = 1 \quad k \in U(12) \quad \phi(12) = 4 \quad \left(\phi\left(\frac{12}{1}\right)\right)$
 $\{1, 5, 7, 11\}$

$$\textcircled{2} \quad \gcd(k, 12) = 2 \Rightarrow \gcd\left(\frac{k}{2}, 6\right) = 1$$

$$\phi(6) = 2 \quad \underline{\underline{\phi\left(\frac{12}{2}\right)}}$$

$$\frac{k}{2} \in U(6) \quad \underline{\underline{\phi(6)}}$$

$$\frac{k}{2} \in \{1, 5\}$$

$$k \in \{\underline{\underline{2}}, \underline{\underline{10}}\}$$

$$\textcircled{3} \quad \gcd(k, 12) = 3 \Rightarrow \gcd\left(\frac{k}{3}, 4\right) = 1$$

$$\phi(4) = 2 \quad \underline{\underline{\phi\left(\frac{12}{3}\right)}}$$

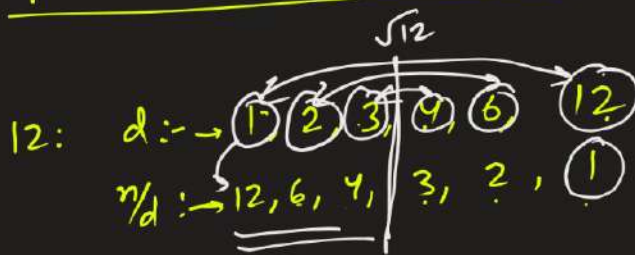
$$\frac{k}{3} \in U_4 = \{1, 3\}$$

$$\Rightarrow k \in \{\underline{\underline{3}}, \underline{\underline{9}}\}$$

$$\textcircled{d} \quad \gcd(k, n) = d \Rightarrow \gcd\left(\frac{k}{d}, \frac{n}{d}\right) = 1$$

$$\hookrightarrow \underline{\underline{\phi\left(\frac{n}{d}\right)}}$$

$$\therefore \boxed{\sum_{k=1}^n \gcd(n, k) = \sum_{d|n} \underbrace{d}_{\frac{n}{\frac{n}{d}}} \cdot \phi\left(\frac{n}{d}\right)} \quad \square$$



$$\underline{\underline{2\sqrt{n}}}$$

$$\frac{2 \cdot \sqrt{n} \cdot \sqrt{n}}{O(n)}$$

for ($i \rightarrow 1$ to \sqrt{n}) {
 if ($n \% i = 0$) {
 3

$$\textcircled{2} \sum_{k=1}^n \text{lcm}(n, k) = \frac{n}{2} \left[1 + \sum_{d|n} d \cdot \phi(d) \right]$$

Proof :- Hints :-

- ① $\text{lcm}(a, b) = \frac{a \times b}{\text{gcd}(a, b)}$
- ② $\text{gcd}(a, n) = \text{gcd}(n-a, n) \leftarrow \text{Euclidean}$
- ③ $\text{lcm}(a, n) + \text{lcm}(n-a, n) = \frac{n^2}{\text{gcd}(a, n)}$

Lemma 1: $\text{lcm}(a, n) + \text{lcm}(n - a, n) = \frac{an}{\text{gcd}(a, n)} + \frac{(n-a)n}{\text{gcd}(n-a, n)} = \frac{n \times n}{\text{gcd}(a, n)}.$

Lemma 2: $\sum \frac{n}{\text{gcd}(a, n)} = \sum_{f \mid n} \frac{n}{f} \times \phi(\frac{n}{f}) = \sum_{d \mid n} d \phi(d),$

Proof: consider what happens if $\text{gcd}(a, n) = f \mid n$. It appears $\phi(\frac{n}{f})$ times on the LHS, and each time it has value of $\frac{n}{f}$. Now substitute $d = \frac{n}{f}$, which is also a divisor of n .

Now, to your problem, pull out $\text{lcm}(n, n) = n$.

We have $2 \sum_{a=1}^{n-1} \text{lcm}(a, n) = \sum [\text{lcm}(a, n) + \text{lcm}(n - a, n)] = n \sum \frac{n}{\text{gcd}(a, n)} = n \times \sum_{d \mid n} d \phi(d).$

Add back $\text{lcm}(n, n) = n$, and you get the formula in OEIS.